

Introduction to Programming Languages

- Today
 - Syllabus
 - Overview

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Office hours

- TR 2-3
 - Occasionally a few minutes late due to meetings between 1 and 2
 - Email me to schedule another time if you cannot make office hours
- See me in person especially for programming questions
 - I will not debug your code through emails
- Start working your homework and projects early

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Grading

- Homework, Surveys, Quizzes, Group assignments (25%)
 - Around 5 written assignments (15%)
 - Surveys (2%)
 - Pop Quizzes (3%)
 - Group presentation (5%)
- Projects (40%)
 - Compile SomeLife
 - Four phases
 - Scheme programming
- Exams (35%)

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Introduction to Programming Languages (Objectives)

- Given a language the student will be able to evaluate the language using common design criteria
- The student will be able to name and describe the different phases of a compiler.

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Reasons for Studying Programming Languages

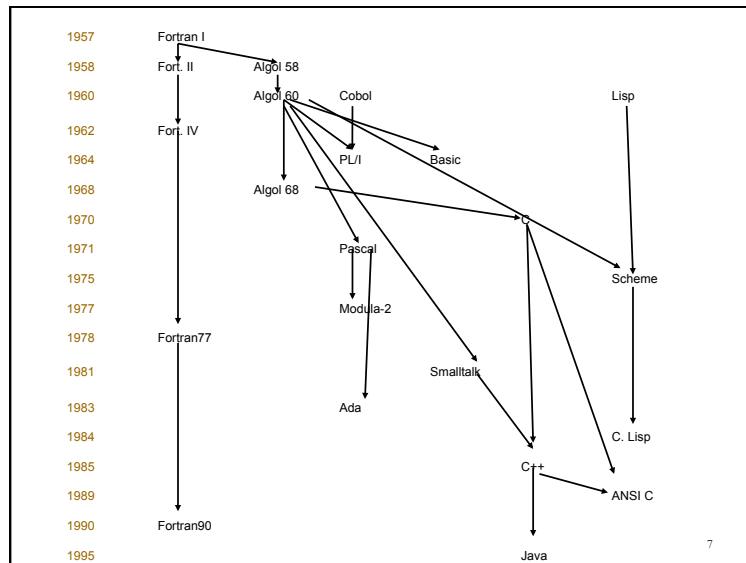
- Increased capacity to express ideas
 - new paradigms → new problem solving skills
- Improved background for choosing an appropriate language
- Increased ability to learn new languages
 - some languages are similar
 - understand obscure language features
- Better understanding of implementation of languages
 - understand implementation costs
 - figure out how to do things in languages that don't support them explicitly
 - Simulate language features
- Increased ability to design new languages
 - Or make better use of language technology wherever it appears

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Programming Domains

- Scientific Applications
 - Computationally intensive
- Business Applications
 - I/O intensive
- Artificial Intelligence
 - Use of symbolics
- Systems software
 - More low-level interactions
- Scripting
 - Lists of commands often doing simple processing
- Special purpose
 - Verilog

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Language Evaluation Criteria

- Readability
- Writability
- Reliability

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Language Evaluation Criteria

- Readability
 - Simplicity
 - Small number of basic components
 - Examples against simplicity
 - feature multiplicity
 - operator overloading
 - Orthogonality
 - A relatively small set of primitive constructs can be combined in a relatively small number of ways to build control and data structures.
 - Consistent simple rules
 - Functions that can't return any data type
 - Related to simplicity

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Language Design Criteria

- Readability (cont.)
 - Control statements
 - Structured control flow (no gotos)
 - Data types and structures
 - Adequate facilities for user-defined types
 - No records in FORTRAN 77
 - No Boolean type in C
 - Syntax Considerations
 - Do not restrict identifier forms (FORTRAN77)
 - Six characters at most
 - Use keywords (none in PL/I)
 - if then = else then else = then else then = else;
 - Appearance indicates their purpose
 - For example, **static** in C

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Language Design Criteria

- Writability
 - Simplicity and orthogonality
 - Not a large number of constructs
 - Consistent set of rules
 - Support for abstraction
 - Data abstraction and process abstraction
 - Ability to define and use complicated structures
 - Not Fortran 77
 - Expressivity
 - Ability to conveniently express common functionality
 - Power of a language
 - Dynamic types, first-class functions in Scheme.

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Language Design Criteria

- Reliability
 - Type checking
 - All type errors should be caught at compile- or run-time.
 - Anything in C
 - Exception handling
 - Ability to intercept run-time errors and take corrective action
 - Aliasing
 - Do not have two or more distinct references to the same memory cell
 - Security
 - Do not allow access to non-user data
 - Stack overflow in C

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Influences on Language Design

- Computer Architecture
 - von Neumann
 - Parallel machines
- Programming methodologies
 - Data abstraction
 - Object oriented vs. procedure oriented

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Paradigms

- A programming paradigm is a way of conceptualization of what it means to perform computation, of structuring and organizing how tasks are carried out in a computer.
- Example
 - A **block-structured paradigm** is a set of programming languages that support nested block structures including procedures.

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Example Paradigms

- | | |
|--|---|
| <ul style="list-style-type: none">■ Imperative<ul style="list-style-type: none">□ Program = steps of computation via state changes□ traditional model of program modifying memory□ features include variables, assignments, arrays, ...□ C, Pascal, Fortran | <ul style="list-style-type: none">■ Object oriented<ul style="list-style-type: none">□ everything is an object□ an object has its own memory□ computation performed by communicated objects□ objects are an instance of class having both data and methods□ Java, Eiffel, Part of C++ |
|--|---|

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Example Paradigms

- | | |
|---|---|
| <ul style="list-style-type: none">■ Functional (Applicative)<ul style="list-style-type: none">□ Values are single entities and are not "stored" in memory□ Computation performed by applying functions□ Functions are 1st-class values which means they can be used like data (created, returned from function calls, ...)□ Parts of Scheme, Lisp | <ul style="list-style-type: none">■ Logic (Declarative)<ul style="list-style-type: none">□ A program is a declaration of facts, rules of inference and queries□ Computation is done by a (backtracking) inference engine that tries to do a "proof". |
|---|---|

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This Class

- Syntax analysis
 - scanners
 - parsers
- Semantic analysis
 - compilers
- Functional programming
 - study the features of Scheme by writing programs
- Language Security

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Why Scanners and Parsers?

- To understand programming languages we need something expressing syntactic structure.
 1. English
 - not precise
 2. Grammar
 - precise but hard to use
 3. Scanner and Parser
 - precise and automatic

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Why Compilers?

- To understand the meaning of a programming language, we need a tool to express that meaning.
- We interact with compilers all the time
 - Java
 - C, C++
- Develops skills for you to reason about program behavior.
- Specify language semantics
 - English
 - not precise
 - awkward

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Why Compilers?

- Specify language semantics
 - Denotational Semantics
 - precise
 - mathematically elegant
 - hard to read
 - Interpreter
 - precise
 - elegant
 - useful beyond theory
 - high-level specification
 - Compiler
 - precise
 - useful beyond theory
 - understanding of implementation increases understanding of concept
 - low-level specification

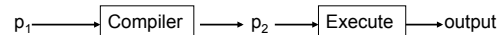
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Compilation vs. Interpretation

- An **interpreter** is a program that takes another program, p_1 , and evaluates p_1 to determine its meaning.



- A **compiler** is a program that takes second program, p_1 , and produces a third program, p_2 , which when evaluated gives the meaning of p_1 . (note that p_2 does not need to be machine language)



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Compilation vs. Interpretation

- **Interpretation:**

- Greater flexibility
- Portability (Java)

- **Compilation**

- Better performance

- **Most languages are a combination of both**

- Java (compilation to Java byte code, which is interpreted and possibly compiled into machine code)
- C (mostly compiled, but I/O formats interpreted)

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What is a compiler?

- A compiler is just a program that takes other programs and converts them into another language. That language is often assembler so that it can be assembled, linked and run on a computer



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Principles of Compiler Design

- The compiler must **preserve** the meaning of the program being compiled
 - The compiler must faithfully implement the defined semantics of a programming language.
- The compiler must **improve** the source code in a discernable way
 - A direct translation of a source program results in highly inefficient code.

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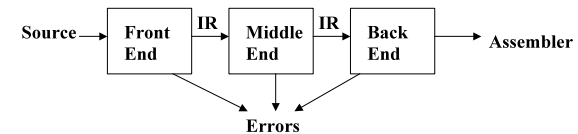
Some Possible Constraints

- Code speed
 - Very fast code might be the highest need of an application
 - e.g., weather
- Code size
 - How much space the object code requires
 - e.g., embedded systems
- Feedback
 - How much feedback is given to the user when an error is encountered
- Compile time
 - programs need to be compiled as fast as possible
- Debugging support
 - code improvements may make debugging difficult

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Overview of Compiler Phases

■ Basic phases

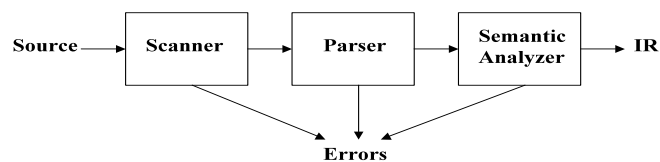


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Front End

■ Syntax Analysis

- determine if programs made up of valid sentences
 - scanner - valid words (reserved words, variables, etc.)
 - parser - valid sentence structure (*if* statements, etc.)
 - semantic analyzer - determine if sentences have meaning (type checking)



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Front End

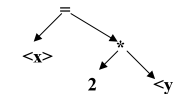
■ Lexical Analysis

- convert words in a program into tokens
 - <var>
 - +, -, =
 - IF, THEN, FOR

■ Parsing

- convert sentences into their structure

$x = 2 * y$



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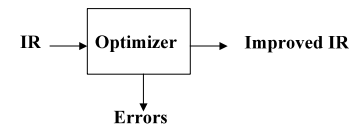
Front End

- Semantic Analysis (context-sensitive analysis)
 - after the program is parsed it is checked to make sure its meaning can be determined
 - type checking
 - number of parameters
 - variables declared
 - functions have prototypes
 - recursion supported or not
- If the program passes semantic analysis, it is converted into an intermediate representation (IR) that is used by the middle end and back end

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Middle End

- Basic structure
 - often the IR is like assembler



- Optimizer is machine-independent

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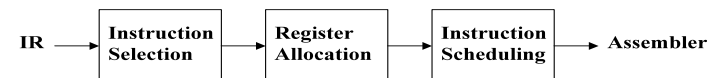
Middle End

- Optimization
 - remove redundancy
 - remove useless code
 - move code out of loops
 - use constants where possible
 - use less expensive operations
 - more ... (cs4130/cs5130 in Spring 2013)

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Back End

- Basic phases



- Order of allocation and scheduling may be different

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Example

- Consider the following

$w = w * 2 * x * y * w * 2$

- Compiler first must recognize
 - variable names
 - =, *
- Next it must determine that the statement is in the source language
- Then, it must make sure the types are correct

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Example

- Next it must allocate space for the variables
 - stack?
 - data segment?
 - registers?
- Then, the front end might generate

```
loadAl    rsp, @w → rw
loadl     2 → r2
mult      rw, r2 → rt1
loadAl    rsp, @x → rx
mult      rt1, rx → rt2
loadAl    rsp, @y → ry
mult      rt2, ry → rt3
loadAl    rsp, @w → rw
mult      rt3, rw → rt4
loadl     2 → r2
mult      rt4, r2 → rt5
storeAl   rt5 → rsp, @w
```

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Example

- Optimization is next
 - eliminate extra loads of w and 2 and extra multiplication of w*2

```
loadAl    rsp, @w → rw
loadl     2 → r2
mult      rw, r2 → rt1
loadAl    rsp, @x → rx
mult      rt1, rx → rt2
loadAl    rsp, @y → ry
mult      rt2, ry → rt3
mult      rt3, rt1 → rt4
storeAl   rt4 → rsp, @w
```

- Code generation converts the intermediate into the target machines assembler.

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Why Scheme?

- In this class we will introduce Scheme. Why?
- Developing skills in a functional language improves overall programming skills.
 - Gives more tools to solve problems
 - Recursion is important and powerful
- Scheme is small and simple
- Functional languages are used in AI, natural language recognition, vision systems, expert systems, rapid prototyping, studies of languages, editors (emacs), ...
- Exposes students to a different way to think about programming.

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Scheme's Distinguishing Features

- Mostly functional
- Expression oriented
- Recursion
- Automatic storage allocation and collection
- PROGRAMS = DATA
- dynamic type checking