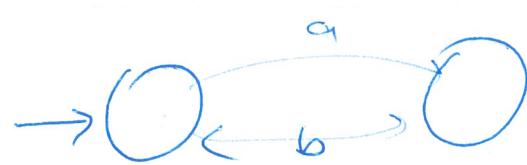


Last time: DFAs



$a^i b^i$



even a's odd b's

DFA

machines
to process
regular expressions

state
based
abstraction

shopp'

checklist

a ... - -

$0 \rightarrow 0$

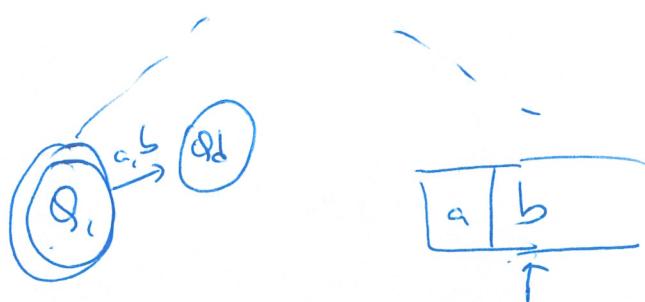
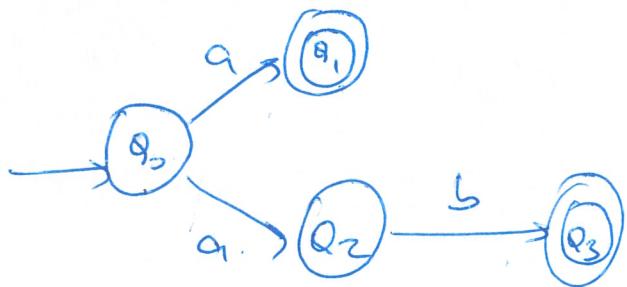
$0 \rightarrow 0$

0lab 1lab

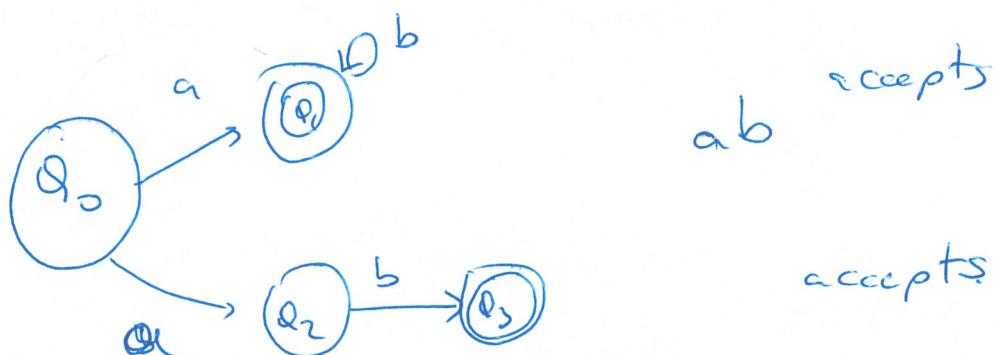


(2)

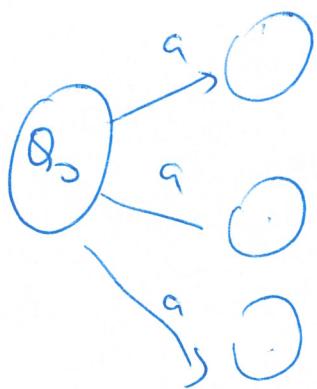
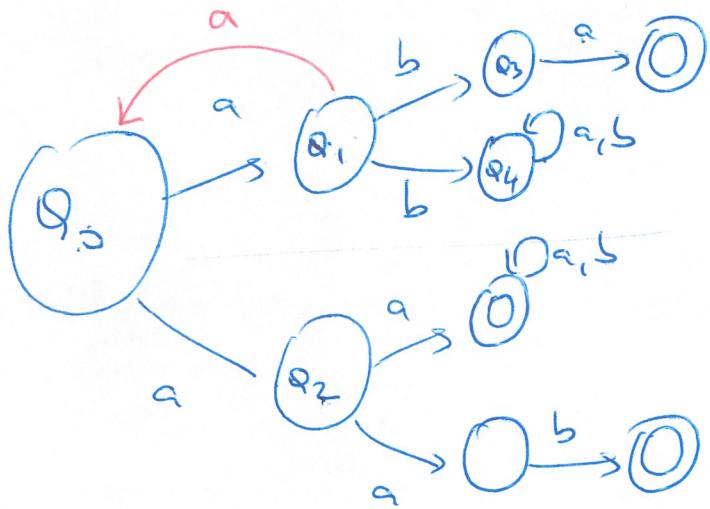
Nondeterministic finite automata NFA



If a "path" leads to an accepting state
 then the NFA accepts.
 still the entire input has to be read.



(3)



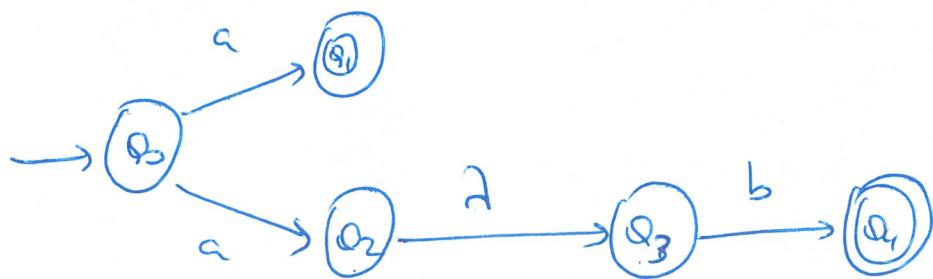
NFA \leftrightarrow DFA conversions.

A DFA is an NFA. T F

An NFA is a DFA : is not true in general.

Can we convert an NFA into a DFA yes
wait for next week please

(4)



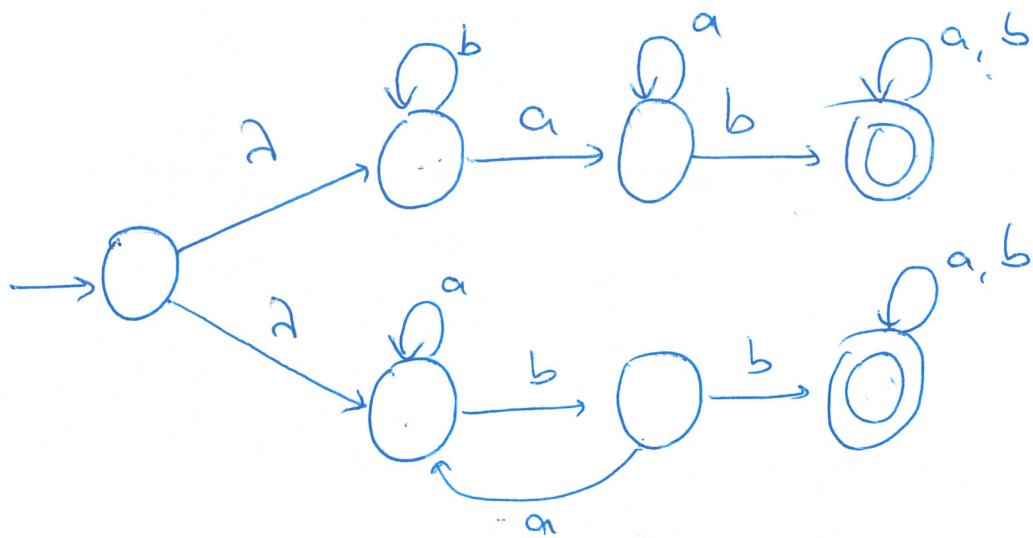
NFA - 2

is an $\begin{cases} \text{DFA} \\ \text{NFA} \end{cases}$ $\xrightarrow{\text{can convert}}$ deterministic
 is an $\begin{cases} \text{NFA - 2} \\ \text{NFA - 1} \end{math}$ $\xrightarrow{\text{non deterministic and}}$ non-deterministic but λ -transitions are not allowed.
 $\xrightarrow{\text{non deterministic and}}$ λ -transitions are allowed

is a $\begin{cases} \text{DFA} \\ \text{NFA} \end{cases}$ $\xrightarrow{\text{can convert}}$
 is a $\begin{cases} \text{NFA - 2} \\ \text{regular expression} \end{cases}$ $\xrightarrow{\text{can convert}}$
 equivalent

$L_1 = \{ w \mid w \text{ contains an "ab", or "bb"} \}$

(5)

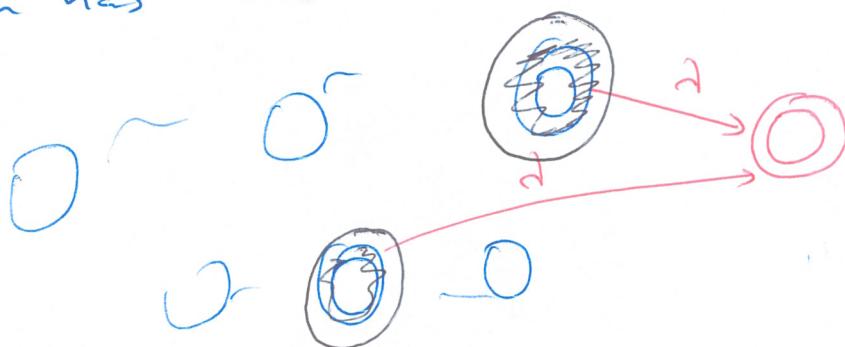


$$[(a \cup b)^* ab (a \cup b)^*] \cup [(a \cup b)^* bb (a \cup b)^*]$$

Consider an NFA-2 which has 2

accepting states.

Can we convert this machine into an NFA-2
which has one accepting state?



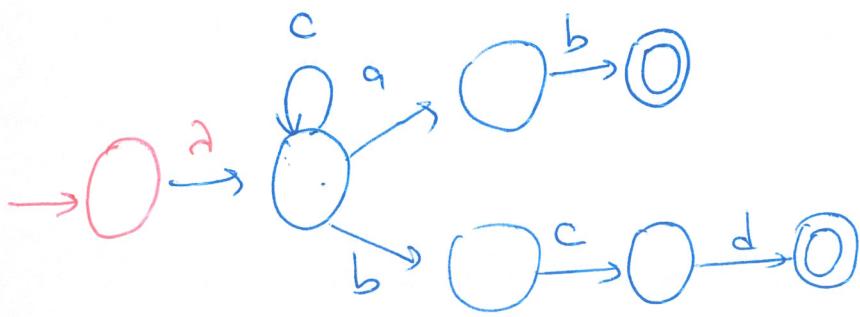
maybe

yes

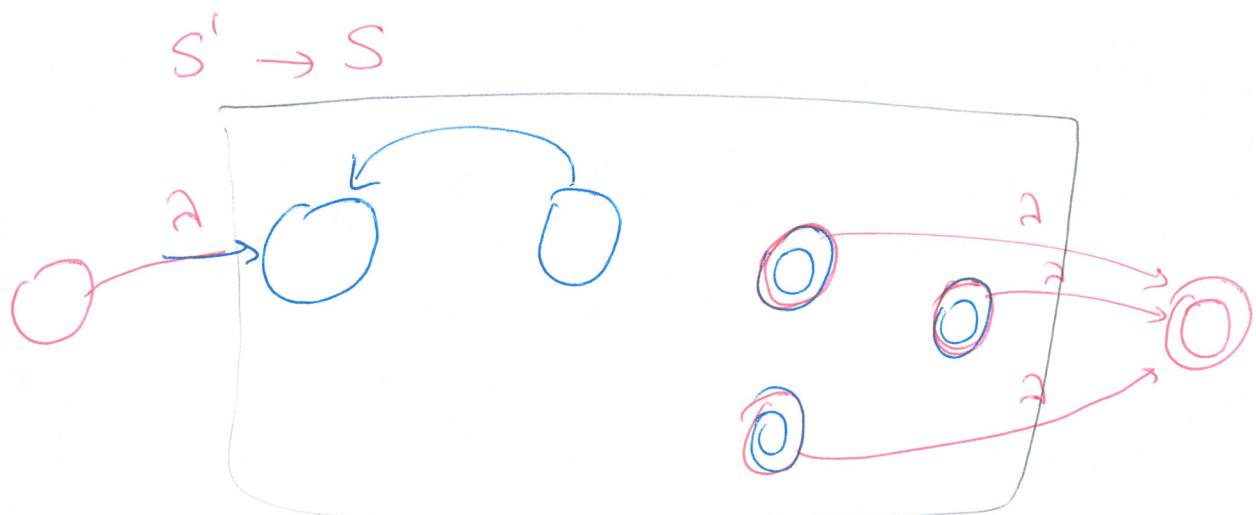
no

TGIF, idk, idc

⑥



We don't want incoming transition to the start state.



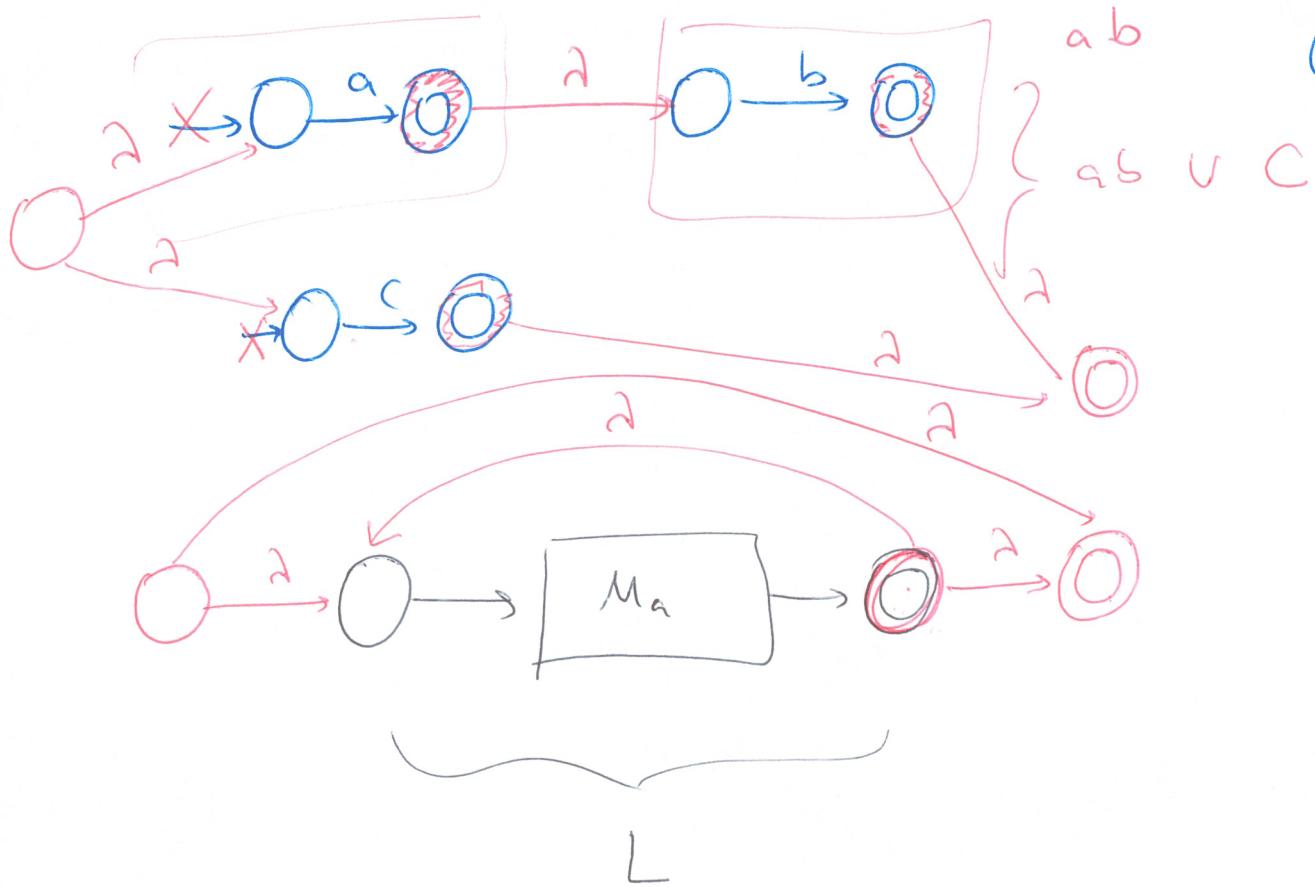
Regular expression

$$\begin{array}{ll}
 \alpha & r_1 \cup r_2 \\
 \alpha^* & r_1 \cdot r_2 \\
 \emptyset & r_1^*
 \end{array}$$

$a b \cup c$

$$(a \cup b)^* (b a^* b a^* b)^*$$

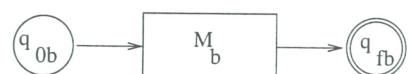
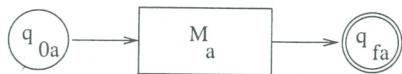
⑦



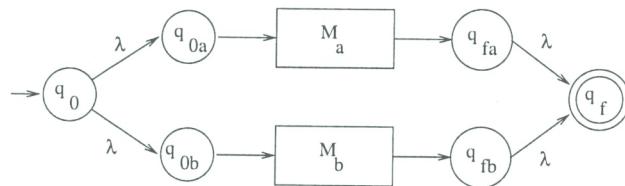
CS3311: How to convert a regular expression into an NFA- λ

Theorem 5.5.3 Let M_1 and M_2 be two NFA- λ s. There are NFA- λ s that accept $L(M_1) \cup L(M_2)$, $L(M_1)L(M_2)$, and $L(M_1)^*$.

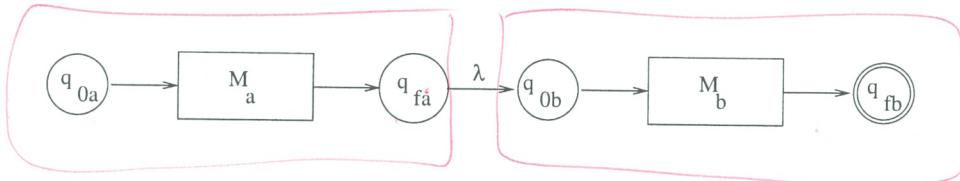
Consider two machines corresponding to two regular expressions:



The following machine represents the union of the two regular expressions:



The following machine represents the concatenation of the two regular expressions:



The following machine represents the Kleene star of a regular expression:

