

CS 3311 Formal Models of Computation — Sample Exam 2

Question 1. (20 points) Give a regular expression for the following sets:

Part a. $L = \{w \mid w \in \{a, b\}^* \text{ and } w \text{ begins and ends with } aa \text{ or } bb\}$

Part b. $L = \{w \mid w \in \{a, b\}^* \text{ and the number of } a\text{'s in } w \text{ is divisible by three } \}$

Question 2. (10 points) Consider the following grammar G :

$$S \rightarrow aSb \mid A$$

$$A \rightarrow cAd \mid B$$

$$B \rightarrow eBf \mid \lambda$$

Part a. Give a derivation for a terminal string such that the $S \rightarrow aSb$ rule is used exactly twice, and the $A \rightarrow cAd$ rule is used exactly once during the derivation.

Part b. Give a derivation for a terminal string such that the $S \rightarrow aSb$ rule is used exactly once, the $A \rightarrow cAd$ rule is not used, and the $B \rightarrow eBf$ rule is used exactly twice during the derivation.

Part c. Use set notation to define the language generated by the grammar.

Question 3. (10 points) Construct a **context-free grammar** over $\{a, b, c\}$ whose language is $\{a^n b^m c^{2n+m} \mid n, m > 0\}$. Explain how you construct the grammar.

Question 4. (10 points)

Part a. Construct a **context-free grammar** over $\{a, b, c\}$ whose language is $\{a^i b^j c^k \mid i = j \text{ or } j = k \text{ where } i, j, k \geq 0\}$. Explain how you construct the grammar.

Part b. Show that your grammar is ambiguous. Present a comprehensive proof with complete sentences.

Question 5. (10 points) Give a **context-free grammar** for the set of strings over $\{a, b\}$ where each string has an odd length, and contains exactly one 'b'. Explain how the grammar generates the strings.

Question 6. (10 points) Consider the CFG G defined by the following productions. Prove by induction that $L(G) = \{a^n b^m \mid 0 \leq n < m\}$. In other words, prove by induction that every string in $L(G)$ has the form $a^n b^m$ where $0 \leq n < m$.

$$\begin{aligned} S &\rightarrow aSb \mid B \\ B &\rightarrow bB \mid b \end{aligned}$$

Question 7. (20 points) Consider the following grammar G with a non-recursive start symbol (S):

$$\begin{aligned} S &\rightarrow T & A &\rightarrow aA \mid BC & B &\rightarrow bB \mid \lambda \\ T &\rightarrow ABC \mid aBC & C &\rightarrow cC \mid \lambda \end{aligned}$$

Part a. Show the set of nullable variables in G .

Part b. Construct an essentially noncontracting grammar G_L (with a non-recursive start symbol) equivalent to G .