

## Lecture 3a: Message-Passing Computing

Parallell Programming (INF-3201)

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## Message-Passing Programming using User-level Message-Passing Libraries

Two primary mechanisms needed:

1. A method of creating separate processes  
- typically executing on different computers
2. A method of sending and receiving messages

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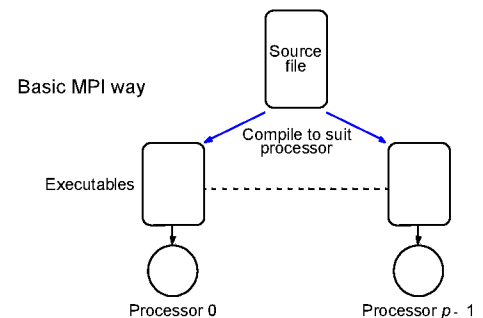
## Process creation

- Static
  - All processes are specified **before execution**
  - The number of processes is fixed during execution.
- Dynamic
  - Processes can be created and their execution initiated **during the execution** of other processes.
  - The number of processes may vary during execution
- The code for processes is normally written and compiled before the execution of any process

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## Single Program Multiple Data (SPMD) model

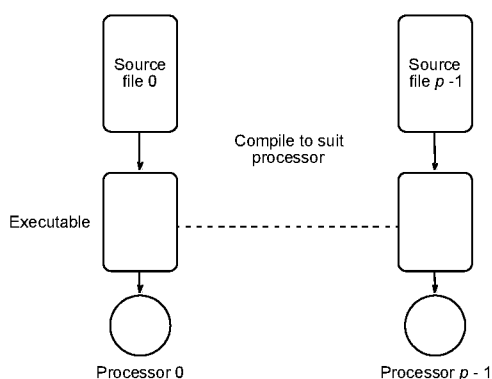
Different processes merged into one program. Control statements select different parts for each processor to execute. All executables started together - static process creation



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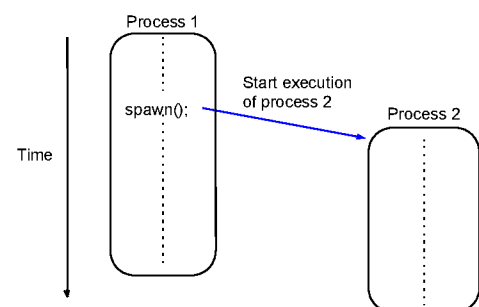
## Multiple program, multiple data (MPMD) model



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## Multiple Program Multiple Data (MPMD) Model

Separate programs for each processor. One processor executes master process. Other processes started from within master process - dynamic process creation.

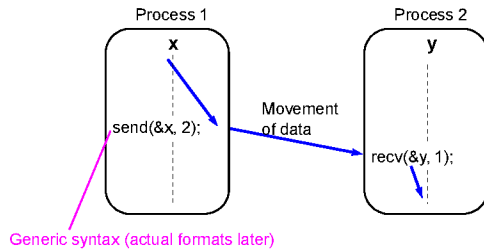


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## Basic “point-to-point” Send and Receive Routines

Passing a message between processes using send() and recv() library calls:



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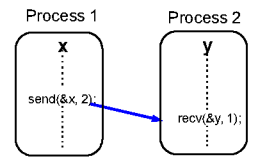
## Synchronous Message Passing

Performs two actions:

- Transfer data
- Synchronize processes.

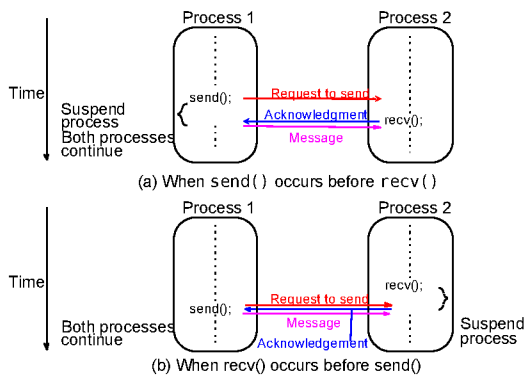
**Send** - returns when message can be accepted by receiver.

**Receive** - returns when message received.



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## Synchronous send() and recv() using 3-way protocol



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## Asynchronous Message Passing

- Routines that do **not wait** for actions to complete before returning. Usually require local storage for messages.
- More than one version depending upon the actual semantics for returning.
- In general, they do **not synchronize processes** but allow processes to move forward sooner. Must be used with care.
  - Process 1:
 

```
x = a
Isend(&x, 2)
x = b
```
  - Which value will be sent to process 2, a or b?

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## MPI Definitions of Blocking and Non-Blocking

- **Blocking** - return after their **local actions** complete, though the message transfer may not have been completed.
- **Non-blocking** - return immediately.
  - Assumes that data storage used for transfer not modified by subsequent statements prior to being used for transfer, and it is left to the programmer to ensure this.

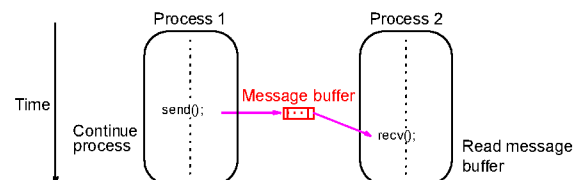
*These terms may have different interpretations in other systems.*

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## How message-passing routines return before message transfer completed

Message buffer needed between source and destination to hold message:



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## Asynchronous (blocking) routines changing to synchronous routines

- Once local actions completed and message is safely on its way, sending process can continue with subsequent work.
- Buffers only of finite length and a point could be reached when send routine held up because all **available buffer space exhausted**.
- Then, send routine will **wait** until storage becomes re-available - i.e then routine behaves as a synchronous routine.

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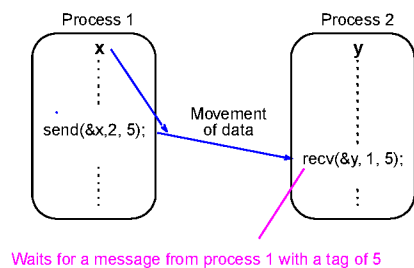
## Message selection - Message Tag

- Used to differentiate between different types of messages being sent.
  - $P_1$ : `send(&x, 2, 5)`
  - $P_2$ : `recv(&y, 1, 5)`
- Message tag is carried within message.
- If special type matching is not required, a **wild card** message tag is used, so that the `recv()` will match with any `send()`.

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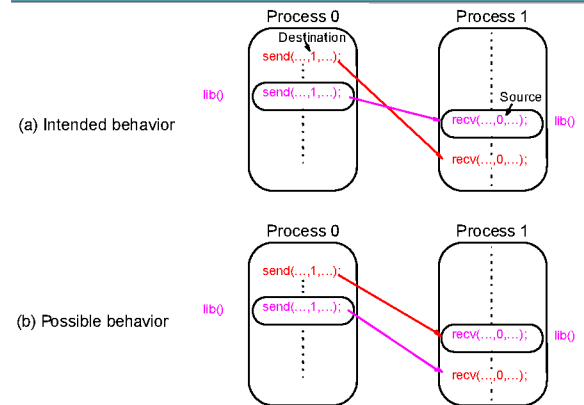
## Message Tag Example

To send a message, x, with message tag 5 from a source process, 1, to a destination process, 2, and assign to y:



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## Unsafe message passing - Example



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## "Group" message passing routines

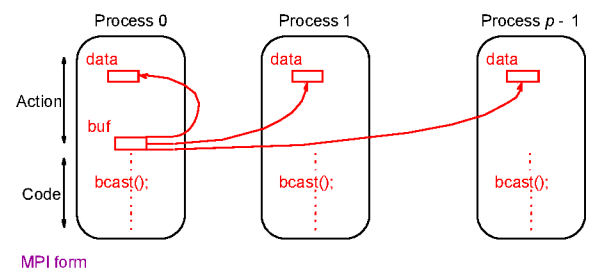
- Have routines that send message(s) to a group of processes or receive message(s) from a group of processes
- Higher efficiency than separate point-to-point routines although not absolutely necessary.
  - Example: IP-multicast

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## Broadcast

Sending same message to all processes concerned with problem.

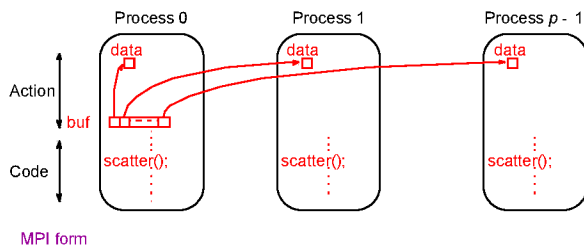
**Multicast** - sending same message to **defined group** of processes.



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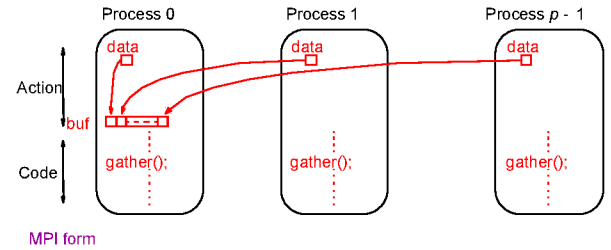
## Scatter

Sending each element of an array in root process to a separate process. Contents of  $i$ th location of array sent to  $i$ th process.



## Gather

Having one process collect individual values from set of processes.



## Reduce

Gather operation combined with specified arithmetic/logical operation.

Example: Values could be gathered and then added together by root:

