

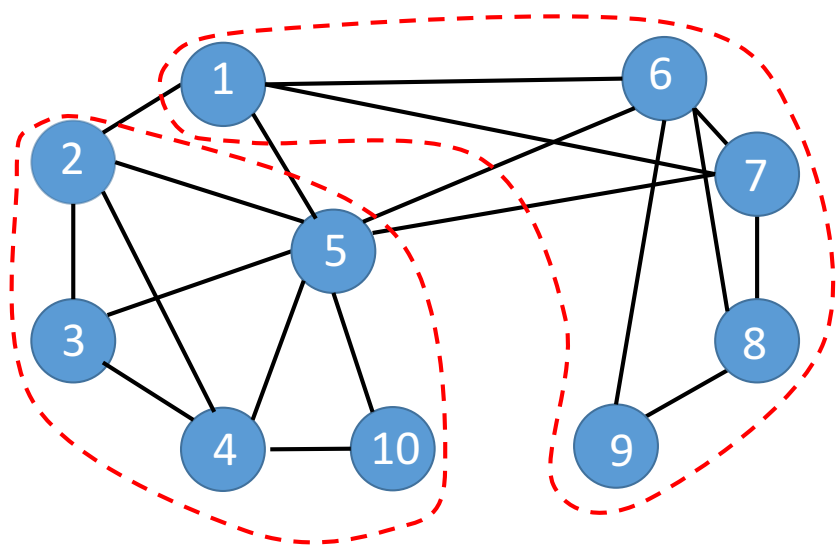
Inhomogeneous Hypergraph Clustering with Applications

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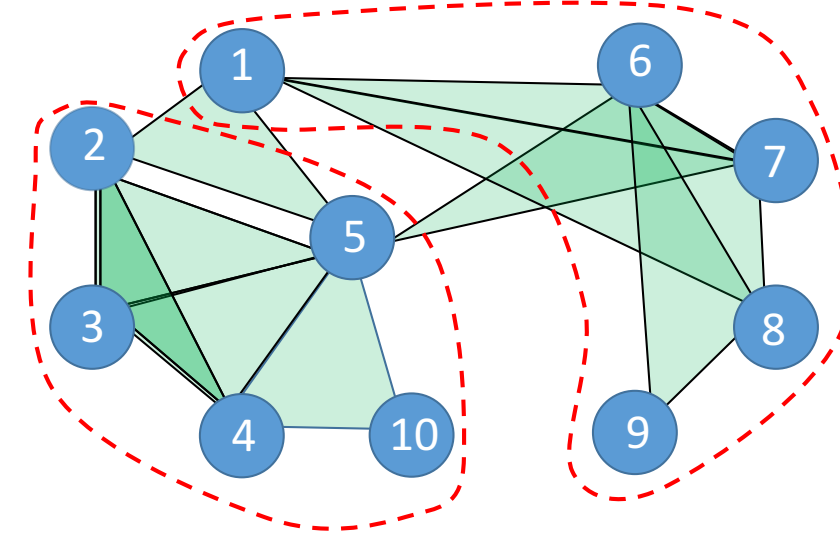
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Introduction

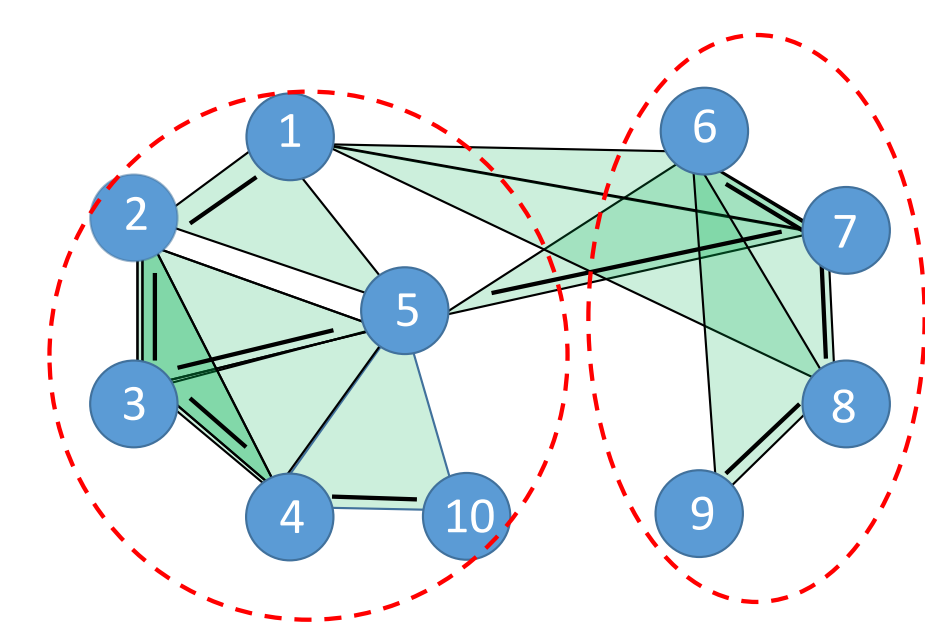
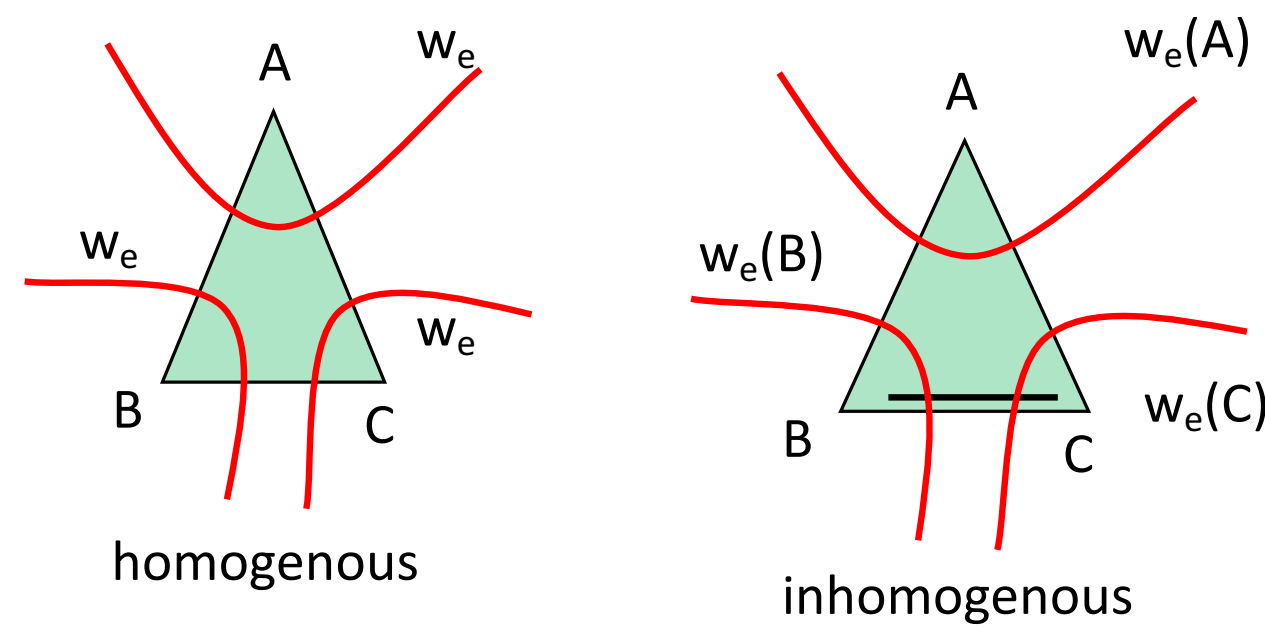
► **Graph clustering:** Group the vertices into clusters such that there should be many edges within each cluster and relatively few between the clusters.



► **Hypergraph clustering:** Relationships among the vertices are more complex than pairwise, modeled as hyperedges.



► **Inhomogeneous Hypergraph (V, E, w) Clustering:** Different subsets of vertices within the same hyperedge may have different structural importance. The weight of hyperedges can be a set function: for $e \in E$, $w_e(\cdot) : 2^e \rightarrow \mathbb{R}_{\geq 0}$, $w_e(S) = w_e(\bar{S})$, $w_e(\emptyset) = 0$.



Balanced partition and normalized cut:

► Partition of V : $V = S \cup \bar{S}$.

► Cost of cut (inhomogeneous case): for $S \subseteq V$,

$$\text{Cut}(S) = \sum_{e \in \partial S} w_e(S \cap e).$$

► Volume of set S :

$$\text{Vol}(S) = \sum_{v \in S} d_v = \sum_{v \in S} \sum_{e: v \in e} w_e(\{v\}).$$

► Inhomogeneous Hypergraph Partitioning: **Minimize normalized hypergraph cut**

$$\arg \min_S \text{NCut}(S) = \text{Cut}(S) \times \left(\frac{1}{\text{Vol}(S)} + \frac{1}{\text{Vol}(\bar{S})} \right).$$

► Generalization of graph/hypergraph normalized cut.

Algorithm

► **Projection:** Use a complete graph to “approximate” each hyperedge. $\{w_{v\bar{v}}^{(e)}\}_{v, \bar{v} \in e}$ stands for projected edge weights from hyperedge e .

$$w_e(\cdot) \rightarrow \{w_{uv}^{(e)}\}_{u, v \in e},$$

e.g. $e = \{u, v, s\}$, $w_e(u) = w_e(v) = 3$, $w_e(s) = 2$.

► **Merging:** Merge subgraphs obtained from each hyperedge.

$$w_{vu} = \sum_{e \in E} w_{vu}^{(e)}$$

► **Spectral partitioning:** Run standard spectral clustering algorithm.

Projection and Infeasibility

► The projection step is critical.

$$\mathbf{P1:} \quad \min_{w^{(e)}} \beta^{(e)} \quad \text{s.t.} \quad w_e(S) \leq \sum_{v \in S, u \in e/S} w_{vu}^{(e)} \leq \beta^{(e)} w_e(S),$$

for all $S \in 2^e$ for which $w_e(S)$ is defined,

► Quadratic optimality: If there exist feasible constants $\beta^{(e)}$ for all hyperedges e and $w_{vu} = \sum_{e \in E} w_{vu}^{(e)} \geq 0$ for all $\{v, u\}$, then $\alpha = \text{NCut}$ of spectral clustering satisfies

$$(\beta^*)^3 \alpha^* \geq \frac{\alpha^2}{8} \geq \frac{(\alpha^*)^2}{8},$$

where α^* is the optimal normalized cut and $\beta^* = \max_{e \in E} \beta^{(e)}$.

► $w_{v\bar{v}}^{(e)}$ may be negative or even the optimization problem is infeasible, e.g. $e = \{1, 2, 3, 4\}$, with $w_e(\{1, 4\}) = w_e(\{2, 3\}) = 1$ and $w_e(S) = 0$ for all other choices of sets S .

Performance guarantee with submodular weights

► Submodular weights:

$$w_e(S_1) + w_e(S_2) \geq w_e(S_1 \cap S_2) + w_e(S_1 \cup S_2) \quad \text{for all } S_1, S_2 \in 2^e.$$

► If $w_e(\cdot)$ is submodular, then there are nonnegative $\{w_{vu}^{(e)}\}_{v, u \in e}$ and $\beta^{(e)}$ feasible to **P1**. For general $|e|$, $\beta^{(e)} \leq |e| - 1$ [1].

► Optimal linear projection (OLP):

$$w_{vu}^{*(e)} = \sum_{S \in 2^e / \{\emptyset, e\}, v \in S} \left[\frac{w_e(S)}{|S|(\delta(e) - |S|)} 1_{u \notin S} - \frac{w_e(S)}{(|S| - 1)(\delta(e) - |S| + 1)} 1_{u \in S} \right].$$

► OLP is proper for practical use.

a) Best known approximation ratio $\beta^{(e)}$ when $|e| \leq 7$.

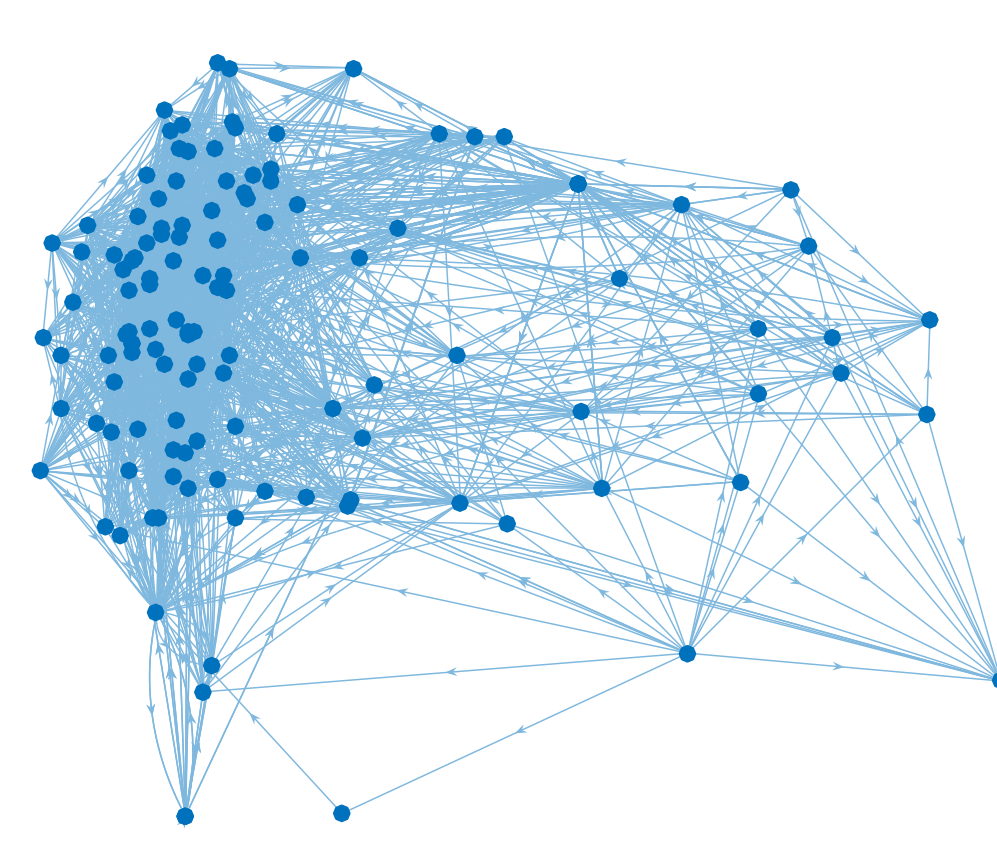
$ e $	2	3	4	5	6	7
β	1	1	1.5	2	4	6

b) Consistency: if $w_e(\cdot)$ is from graph cuts, then OLP recovers the underlying graph.

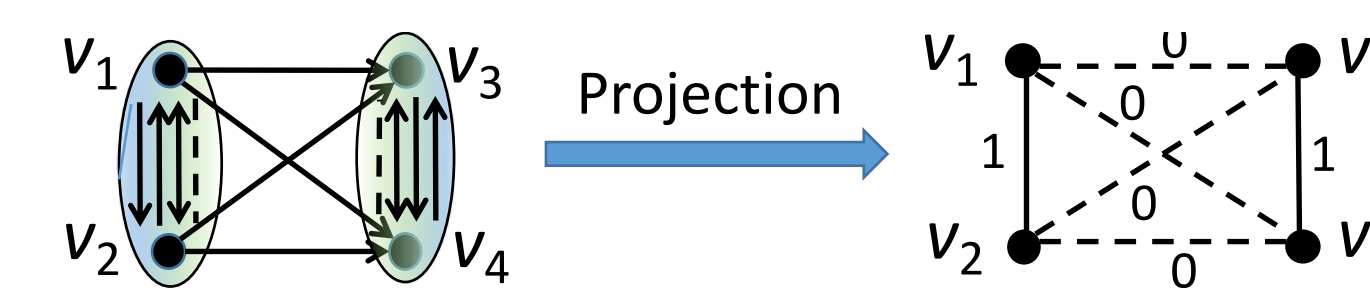
c) Efficiently computable.

Application I: Foodweb hierarchical clustering

► Florida Bay food web



► The network motif to construct the inhomogeneous hypergraph:

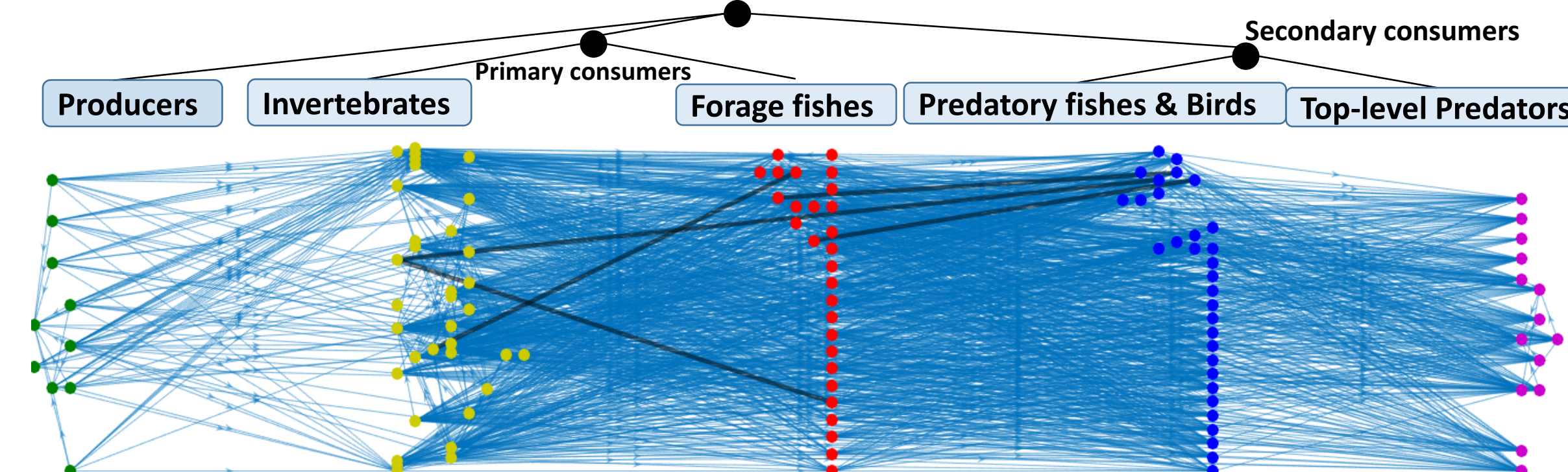


$$w_e(\{v_i\}) = 1 \text{ for } i = 1, 2, 3, 4$$

$$w_e(\{v_1, v_3\}) = 2, \text{ and } w_e(\{v_1, v_4\}) = 2$$

$$w_e(\{v_1, v_2\}) = 0$$

► 5 clusters; Only 6 carbon flows are in inverse directions.



Application II: Category Learning in Rankings

► Rriffled independence: Rankings of candidates within different categories are independent from each other [2].

► Goal: Detecting the riffled independent category S .
 Q : Set of candidates to be ranked.

$S \subseteq Q$, a specific category of candidates.

$\sigma : Q \rightarrow [n]$, a full ranking over Q .

► Define

$$I_{i,j,k} \triangleq I(\sigma(i); 1_{\sigma(j) < \sigma(k)}) = \sum_{\sigma(i)} \sum_{1_{\sigma(j) < \sigma(k)}} \mathbb{P}(\sigma(i), 1_{\sigma(j) < \sigma(k)}) \log \frac{\mathbb{P}(\sigma(i), 1_{\sigma(j) < \sigma(k)})}{\mathbb{P}(\sigma(i)) \mathbb{P}(1_{\sigma(j) < \sigma(k)})},$$

$I_{i,j,k} = 0$ indicates the position of i is independent from the rankings of j and k .

► The underlying optimization problem:

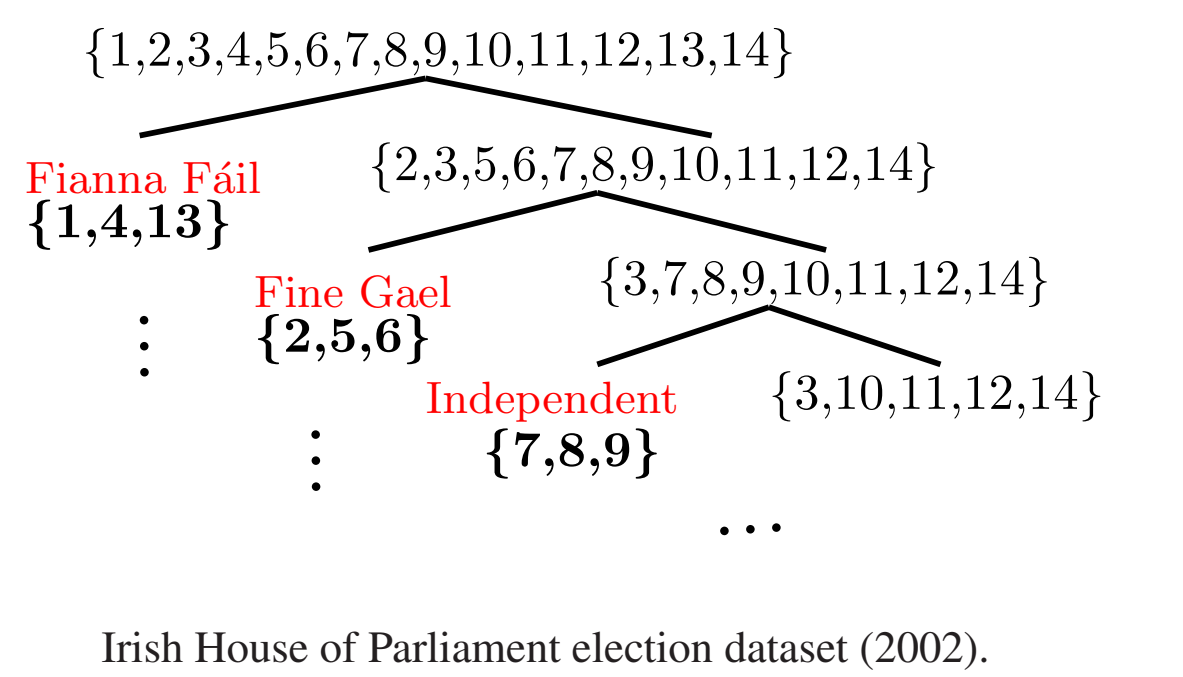
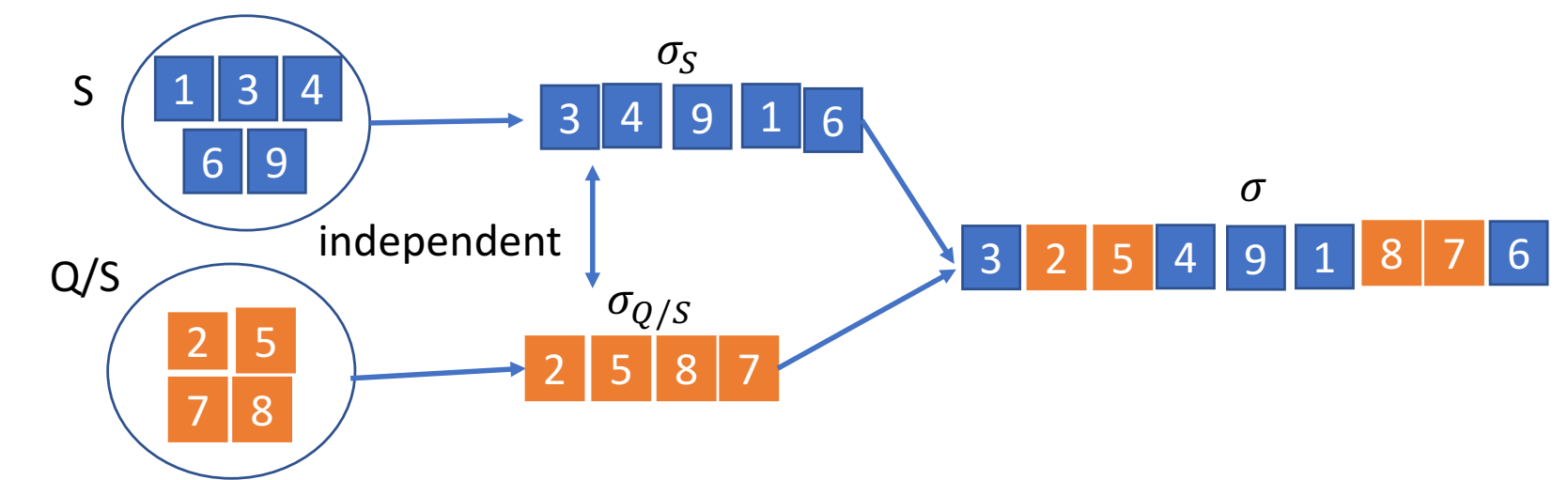
$$\arg \min_{S \subseteq Q} \mathcal{F}(S) \triangleq \sum_{(i,j,k) \in \Omega_{S,S}^{cross}} I_{i,j,k} + \sum_{(i,j,k) \in \Omega_{S,S}^{cross}} I_{i,j,k},$$

$\Omega_{A,B}^{cross} \triangleq \{(i,j,k) | i \in A, j, k \in B\}$, A, B different categories.

► A normalized form of above problem equivalent to inhomogeneous hypergraph partitioning.

► The correct categories of candidates can be detected with only 70% of samples as opposed to the benchmark.

► Only 30% of triples of mutual information $I_{i,j,k}$ requires to be estimated.



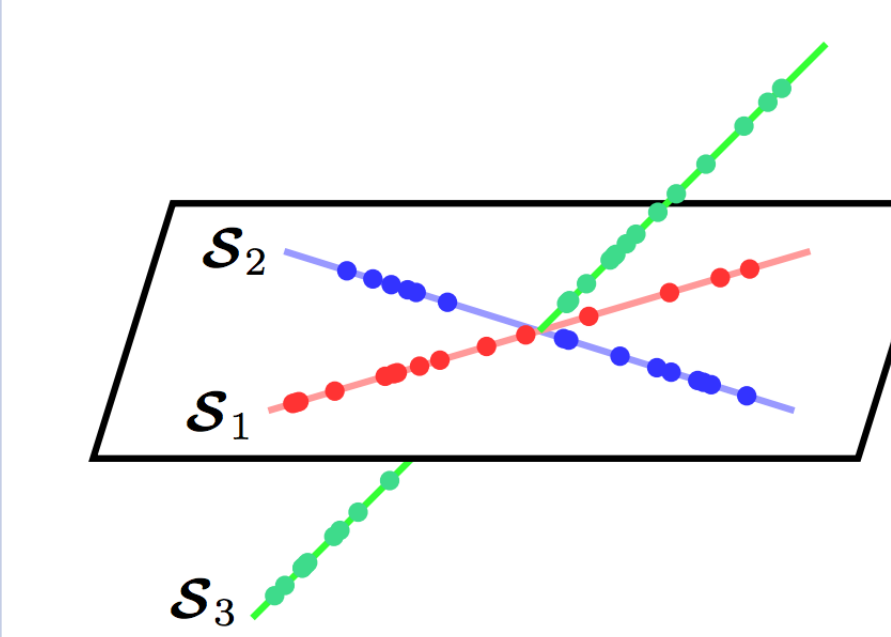
Application III: Subspace clustering

► Subspace segmentation has the goal to partition data according to their intrinsically embedded subspaces.

► The intrinsic affine space: p -dimensional.

► Construct ψ -uniform ($\psi > p + 1$, typically set to $p + 2$) hypergraphs $\mathcal{H} = (V, E, w)$.

► w_e , the weight of hyperedge e , is to characterize how likely the corresponding points in e fit into one affine space.



Union of three subspaces ($p = 1$)

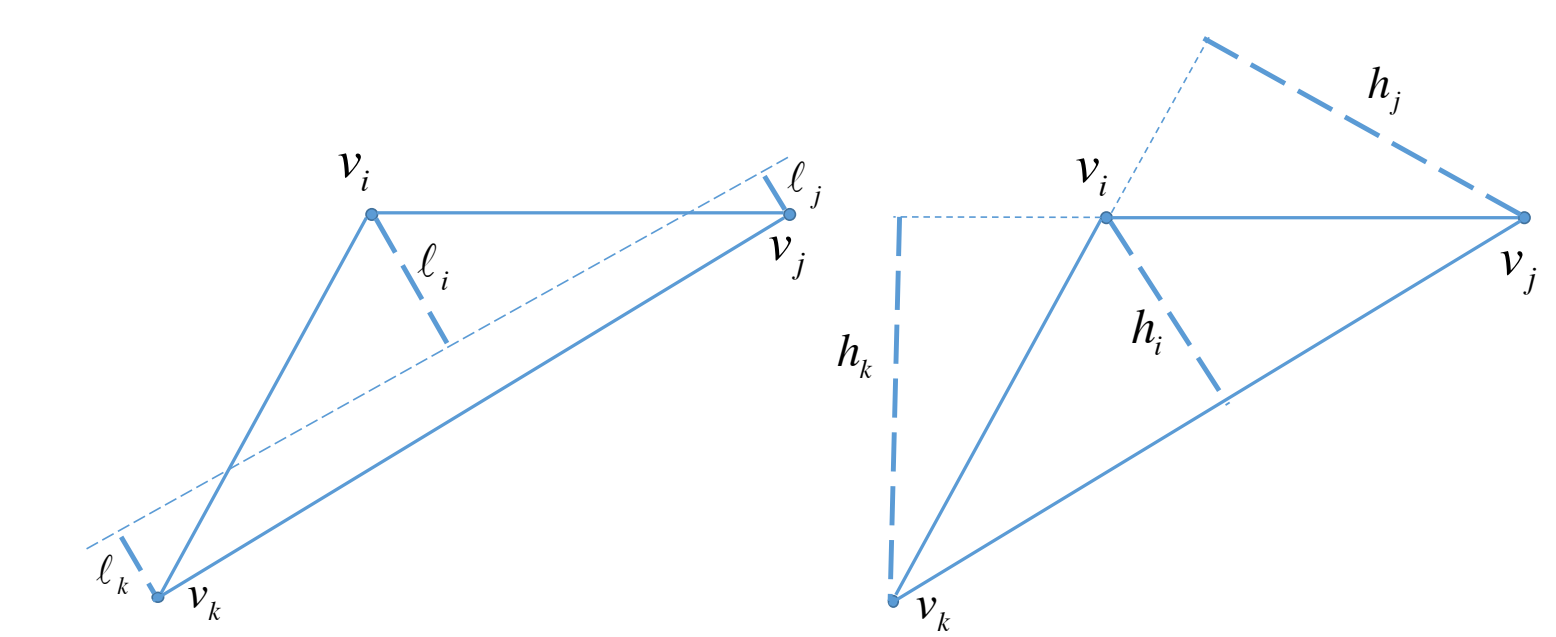


Illustration of the weight ($p = 1$) used for subspace segmentation.

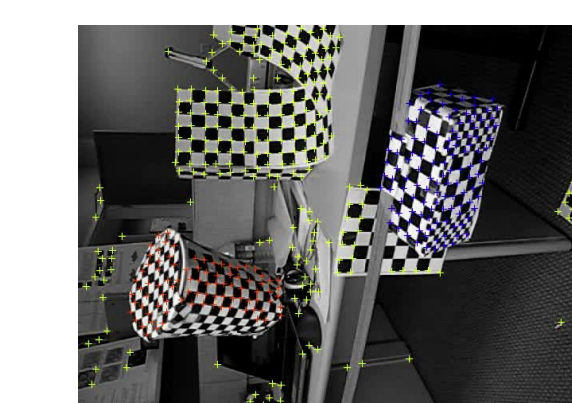
$$d_e^{H-1} = (\ell_i + \ell_j + \ell_k)/3$$

$$d_e^{H-2} = (h_i + h_j + h_k)/3$$

$$d_e^{\text{InH}}(\{v_i\}) = h_i$$

$$w_e = \exp(-d_e^2/\theta^2)$$

► Evaluation over the Hopkins 155 dataset:



Method	Misclassification rates e% (MN: mean; MD: median)									
	Two Motions					Three Motions				
	Chck.(78)	Trfc.(31)	Arnc.(11)	All(120)		Chck.(26)	Trfc.(7)	Arnc.(2)	All(115)	
	MN MD	MN MD	MN MD	MN MD		MN MD	MN MD	MN MD	MN MD	
H+ d_e^{H-1}	12.27 5.06	14.91 9.94	12.85 3.66	12.92 6.01		22.13 23.98	21.99 18.12	19.79 19.79	21.97 20.45	
H+ d_e^{H-2}	4.20 0.43	0.33 0.00	1.53 0.10	2.93 0.06		7.05 2.22	7.02 3.98	6.47 6.47	7.01 2.12	
InH	1.69 0.00	0.61 0.22	1.22 0.62	1.40 0.04		4.82 0.69	2.46 0.60	4.23 4.23	4.06 0.65	

References

- [1] N. R. Devanur, S. Dughmi, R. Schwartz, A. Sharma, and M. Singh, “On the approximation of submodular functions,” *arXiv preprint arXiv:1304.4948*, 2013.
- [2] J. Huang, C. Guestrin *et al.*, “Uncovering the riffled independence structure of ranked data,” *Electronic Journal of Statistics*, vol. 6, pp. 199–230, 2012.