Bernie Pope, bjpope@unimelb.edu.au

Implementing Python in Haskell, twice!

Melbourne Python Users Group, Monday 7 July 2014

Is this about Haskell or Python?



It's a bit of both

Overview

- * Motivation
- Parsing Python
- Translating Python to Haskell
- Bytecode compilation and interpretation
- * Future directions

Motivation

- * A fun way to kill time.
- * Haskell is particularly good for writing compilers.
- * Python is a (mostly) simple language to implement.

Lexical analysis

- * I use Alex (like Lex/Flex) for lexical analysis.
- * Lexical analysis takes the input source program text, the filename (for error messages), and breaks the source up into a sequence of tokens:

lex :: String -> Filename -> Either ParseError [Token]

* Python's lexical structure is well defined:

https://docs.python.org/3/reference/lexical_analysis.html

Parsing

- * I use Happy (like Yacc/Bison) for parsing.
- * Parsing takes the input source program text, the filename (for error messages), and builds an abstract syntax tree:

* Python's grammar is well defined:

```
https://docs.python.org/3/reference/grammar.html
```

Parsing

* Internally the parser calls the lexer and processes the generated sequence of tokens.

language-python

- * Eventually I'd written a lexer, parser and pretty printer.
- * Support for both Python 2.x and 3.x.
- * Could parse the CPython test suite.
- * The result is a library called language-python: https://github.com/bjpop/language-python
- * Now what?

- * Thought experiment: what would it take to translate Python into Haskell?
- Okay, what is the semantics of Python?

```
while True:
    try:
        1/0
    except:
        break
    finally:
    continue
```

- * Thought experiment: what would it take to translate Python into Haskell?
- * Okay, what is the semantics of Python?

```
while True:

try:

1/0

except:

break

finally:

continue
```

SyntaxError:
'continue' not
supported inside
'finally' clause

- * The real trick is in encoding Python's imperative effects into a pure language (state, control flow, mutable values, input/output).
- * Haskell's monad (transformers) provide an elegant way to combine different effects together.

* An example, recursive factorial using an accumulator:

```
def fac(n, acc):
    if n == 0:
        return acc
    else:
        return fac(n - 1, n * acc)
```

```
def s fac 2 none
   (\ [ s n, s acc] ->
       ifThenElse
           (do t 6 <- read s n
              t 6 == 0)
           (do t 7 <- read s acc
              ret t 7)
           (do t 0 <- read s fac
              t 1 <- read s n
              t 2 <- t 1 - 1
              _t_3 <- read s n
              t 4 <- read s acc
              t 5 <- t 3 * t 4
               tailCall t 0 [ t 2, t 5]))
```

def fac(n, acc):

* translated into Haskell by berp:

def s fac 2 none

```
def s fac 2 none
   (\ [ s n, s acc] ->
       ifThenElse
           (do t 6 <- read s n
              t 6 == 0)
           (do t 7 <- read s acc
              ret t 7)
           (do t 0 <- read s fac
              t 1 <- read s n
              t 2 <- t 1 - 1
              _t_3 <- read s n
              t 4 <- read s acc
              t 5 <- t 3 * t 4
               tailCall t 0 [ t 2, t 5]))
```

* translated into Haskell by berp:

```
def s fac 2 none
    (\ [ s n, s acc] \rightarrow
       ifThenElse
           (do t 6 <- read s n
               _t_6 == 0)
           (do t 7 <- read s acc
               ret t 7)
           (do t 0 <- read s fac
               t 1 <- read s n
               t 2 <- t 1 - 1
               _t_3 <- read s n
               t 4 <- read s acc
               t 5 <- t 3 * t 4
               tailCall t 0 [ t 2, t 5]))
```

n == 0:

```
def s fac 2 none
   (\ [ s n, s acc] ->
       ifThenElse
           (do t 6 <- read s n
              t 6 == 0)
           (do t 7 <- read s acc
              ret t 7)
           (do t 0 <- read s fac
              t 1 <- read s n
              t 2 <- t 1 - 1
              _t_3 <- read s n
              t 4 <- read s acc
              t 5 <- t 3 * t 4
               tailCall t 0 [ t 2, t 5]))
```

* translated into Haskell by berp:

```
def s fac 2 none
   (\ [ s n, s acc] ->
       ifThenElse
           (do t 6 <- read s n
              t 6 == 0)
           (do t 7 <- read s acc
              ret t 7)
          (do t 0 <- read s fac
              t 1 <- read s n
              t 2 <- t 1 - 1
              _t_3 <- read s n
              t 4 <- read s acc
              t 5 <- t 3 * t 4
              tailCall t 0 [ t 2, t 5]))
```

return acc

```
def s fac 2 none
   (\ [ s n, s acc] ->
       ifThenElse
           (do t 6 <- read s n
              t 6 == 0)
           (do t 7 <- read s acc
              ret t 7)
           (do t 0 <- read s fac
              t 1 <- read s n
              t 2 <- t 1 - 1
              _t_3 <- read s n
              t 4 <- read s acc
              t 5 <- t 3 * t 4
               tailCall t 0 [ t 2, t 5]))
```

```
def s fac 2 none
   (\ [ s n, s acc] \rightarrow
       ifThenElse
           (do t 6 <- read s n
              t 6 == 0)
           (do t 7 <- read s acc
              ret t 7)
           (do t 0 <- read s fac
              t 1 <- read s n
              t 2 <- t 1 - 1
                                      return fac(n - 1, n * acc)
              t 3 <- read s n
              t 4 <- read s acc
              t 5 <- t 3 * t 4
              tailCall t 0 [ t 2, t 5]))
```

```
def s fac 2 none
   (\ [ s n, s acc] ->
       ifThenElse
           (do t 6 <- read s n
              t 6 == 0)
           (do t 7 <- read s acc
              ret t 7)
           (do t 0 <- read s fac
              t 1 <- read s n
              t 2 <- t 1 - 1
              _t_3 <- read s n
              t 4 <- read s acc
              t 5 <- t 3 * t 4
               tailCall t 0 [ t 2, t 5]))
```

- * Berp has some cute party tricks:
 - Tail call optimisation, fac runs in constant stack space.
 - callcc (call with current continuation, borrowed from Scheme).

* callCC example in berp:

```
>>> def f():
        count = 0
        k = callCC(lambda x: x)
   print(count)
        if count < 3:
            count = count + 1
            k(k)
>>> f()
0
2
```

Is Haskell a good target for Python compilation?

* Pros:

• We get to use the Haskell runtime features for free: garbage collection, threads, I/O.

* Cons:

- The runtime representation of Python state (in berp) is pretty heavy weight (slow).
- Python uses a lot of mutation. Haskell compilers are not optimised for this.

- * Having pursued the berp thought experiment far enough I decided to try making a bytecode compiler and interpreter.
- * I started by writing a program to read .pyc files generated by CPython. Then used it to reverse engineer the meaning of the bytecode.

```
0 LOAD FAST 0
3 LOAD CONST 1
6 COMPARE OP 2
9 POP_JUMP_IF_FALSE 19
12 LOAD_FAST 1
15 RETURN_VALUE
16 JUMP_FORWARD 21
19 LOAD GLOBAL 0
22 LOAD_FAST 0
25 LOAD_CONST 2
28 BINARY SUBTRACT
29 LOAD_FAST 0
32 LOAD FAST 1
35 BINARY MULTIPLY
36 CALL_FUNCTION 2
39 RETURN_VALUE
40 LOAD CONST 0
43 RETURN_VALUE
```

```
    Bytecode for the factorial function:

       0 LOAD_FAST 0
       3 LOAD_CONST 1
       6 COMPARE_OP 2
       9 POP_JUMP_IF_FALSE 19
       12 LOAD_FAST 1
       15 RETURN_VALUE
       16 JUMP_FORWARD 21
       19 LOAD GLOBAL 0
       22 LOAD_FAST 0
       25 LOAD_CONST 2
       28 BINARY SUBTRACT
       29 LOAD_FAST 0
       32 LOAD FAST 1
       35 BINARY MULTIPLY
       36 CALL_FUNCTION 2
       39 RETURN_VALUE
       40 LOAD CONST 0
```

43 RETURN_VALUE

```
0 LOAD FAST 0
3 LOAD CONST 1
6 COMPARE OP 2
9 POP_JUMP_IF_FALSE 19
12 LOAD_FAST 1
15 RETURN_VALUE
16 JUMP_FORWARD 21
19 LOAD GLOBAL 0
22 LOAD_FAST 0
25 LOAD_CONST 2
28 BINARY SUBTRACT
29 LOAD_FAST 0
32 LOAD FAST 1
35 BINARY MULTIPLY
36 CALL_FUNCTION 2
39 RETURN_VALUE
40 LOAD CONST 0
43 RETURN_VALUE
```

```
0 LOAD FAST 0
3 LOAD CONST 1
6 COMPARE_OP 2
9 POP JUMP IF FALSE 19
12 LOAD_FAST 1
                           return acc
15 RETURN VALUE
16 JUMP_FORWARD 21
19 LOAD_GLOBAL 0
22 LOAD_FAST 0
25 LOAD_CONST 2
28 BINARY SUBTRACT
29 LOAD_FAST 0
32 LOAD FAST 1
35 BINARY MULTIPLY
36 CALL_FUNCTION 2
39 RETURN_VALUE
40 LOAD CONST 0
43 RETURN_VALUE
```

```
0 LOAD FAST 0
3 LOAD CONST 1
6 COMPARE OP 2
9 POP_JUMP_IF_FALSE 19
12 LOAD_FAST 1
15 RETURN_VALUE
16 JUMP_FORWARD 21
19 LOAD GLOBAL 0
22 LOAD_FAST 0
25 LOAD_CONST 2
28 BINARY SUBTRACT
29 LOAD_FAST 0
32 LOAD FAST 1
35 BINARY MULTIPLY
36 CALL_FUNCTION 2
39 RETURN_VALUE
40 LOAD CONST 0
43 RETURN_VALUE
```

```
0 LOAD FAST 0
3 LOAD CONST 1
6 COMPARE_OP 2
9 POP JUMP IF FALSE 19
12 LOAD_FAST 1
15 RETURN_VALUE
16 JUMP_FORWARD 21
19 LOAD_GLOBAL 0
22 LOAD FAST 0
25 LOAD CONST 2
28 BINARY SUBTRACT
29 LOAD_FAST 0
32 LOAD_FAST 1
35 BINARY_MULTIPLY
36 CALL_FUNCTION 2
39 RETURN VALUE
40 LOAD CONST 0
43 RETURN_VALUE
```

```
return fac(n - 1, n * acc)
```

```
0 LOAD FAST 0
3 LOAD CONST 1
6 COMPARE OP 2
9 POP_JUMP_IF_FALSE 19
12 LOAD_FAST 1
15 RETURN_VALUE
16 JUMP_FORWARD 21
19 LOAD GLOBAL 0
22 LOAD_FAST 0
25 LOAD_CONST 2
28 BINARY SUBTRACT
29 LOAD_FAST 0
32 LOAD FAST 1
35 BINARY MULTIPLY
36 CALL_FUNCTION 2
39 RETURN_VALUE
40 LOAD CONST 0
43 RETURN_VALUE
```

Bytecode for the factorial function:

```
0 LOAD FAST 0
3 LOAD CONST 1
6 COMPARE_OP 2
9 POP_JUMP_IF_FALSE 19
12 LOAD_FAST 1
15 RETURN_VALUE
16 JUMP_FORWARD 21
19 LOAD_GLOBAL 0
22 LOAD FAST 0
25 LOAD CONST 2
28 BINARY SUBTRACT
29 LOAD_FAST 0
32 LOAD FAST 1
35 BINARY MULTIPLY
36 CALL_FUNCTION 2
39 RETURN_VALUE
40 LOAD_CONST 0
```

43 RETURN_VALUE

return None

Bytecode for the factorial function:

```
0 LOAD FAST 0
3 LOAD_CONST 1
6 COMPARE_OP 2
9 POP JUMP IF FALSE 19
12 LOAD_FAST 1
15 RETURN_VALUE
16 JUMP_FORWARD 21
19 LOAD_GLOBAL 0
22 LOAD FAST 0
25 LOAD CONST 2
28 BINARY SUBTRACT
29 LOAD FAST 0
32 LOAD FAST 1
35 BINARY MULTIPLY
36 CALL FUNCTION 2
39 RETURN_VALUE
40 LOAD_CONST 0
43 RETURN_VALUE
```

return None

dead code

* It turns out that CPython uses a very straightforward compilation scheme. Easy to emulate in Haskell:

```
compileExpr :: ExprSpan -> Compile ()
compileExpr (CondExpr {..}) = do
    compile ce_condition
    falseLabel <- newLabel
    emitCodeArg POP_JUMP_IF_FALSE falseLabel
    compile ce_true_branch
    restLabel <- newLabel
    emitCodeArg JUMP_FORWARD restLabel
    labelNextInstruction falseLabel
    compile ce_false_branch
    labelNextInstruction restLabel</pre>
```

- * Blip generates bytecode which is compatible with CPython.
- * I originally used CPython to test the generated byte code.
- * Then I decided to write a bytecode interpreter in Haskell too.

* Evaluating the bytecode is reasonably straightforward:

```
evalOneOpCode :: HeapObject -> Opcode -> Word16 -> Eval ()
evalOneOpCode (CodeObject {..}) opcode arg =
   case opcode of
      CALL FUNCTION -> do
         functionArgs <- replicateM (fromIntegral arg)</pre>
                              popValueStack
         functionObjectID <- popValueStack</pre>
         functionObject <- lookupHeap functionObjectID</pre>
         callFunction functionObject $ reverse functionArgs
      JUMP ABSOLUTE -> setProgramCounter $ fromIntegral arg
      ... etcetera ...
```

- * I've implemented about half of the bytecode instructions so far. Many of them are simple manipulations of the stack.
- * However, implementing attribute lookup completely and correctly is quite difficult.

* Time for a little demo:

```
$ blip
Berp version 0.2.1, type control-d to exit.
>>> def fac(n):
        if n <= 1:
            return 1
        else:
            return fac(n - 1) + fac(n - 2)
>>> x = 0
>>> while x < 10:
        print(fac(x))
        x = x + 1
1
2
5
8
21
34
55
```

Future directions

- Complete the byte code interpreter
- * Add language extensions:
 - * Tail call optimisation?
 - Algabraic types, pattern matching?
- * Possibly optimise execution by compiling to machine code.

Source code

- https://github.com/bjpop/language-python
- https://github.com/bjpop/berp
- https://github.com/bjpop/blip