

# CSE-321 Assignment 2 - *More Fun with Objective CAML*

(100 points)

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Due at 11:59pm, March 5

Welcome to the second assignment of CSE-321 Programming Languages! In this assignment, you will further familiarize yourself with functional programming in Objective CAML (OCAML) by implementing various tail-recursive functions, sorting algorithms, and structures.

In order to assist in grading your assignments, you should *strictly* follow the submission instruction.

## 1 Submission instruction

Download the zip file `hw2.zip` from the course webpage or from `/home/class/cs321/` on `programming2.postech.ac.kr`, and unzip it:

```
lngsg@programming2:~ $ unzip hw2.zip
Archive:  hw2.zip
   creating: hw2/
  inflating: hw2/hw2.mli
  inflating: hw2/Makefile
  inflating: hw2/.depend
  inflating: hw2/hw2.ml
```

You will write code in `hw2.ml` and never touch other files. The stub file `hw2.ml` looks like:

```
exception NotImplemented

type 'a tree = Leaf of 'a | Node of 'a tree * 'a * 'a tree

(** Recursive functions **)

let rec lrevrev _ = raise NotImplemented

let rec lfoldl _ _ _ = raise NotImplemented

(** Tail recursive functions **)

let fact _ = raise NotImplemented

...
```

1. Fill the function body with your own code *only if you have a correct implementation of the function*. This is absolutely crucial; if you leave code that does not compile, you will receive no credit. If you cannot implement a function, just leave it intact! Make sure that your program compiles by running `make`:

```
lngsg@programming2:~/hw2 $ ls
hw2.ml  hw2.mli  Makefile
lngsg@programming2:~/hw2 $ make
ocamlc -c hw2.mli -o hw2.cmi
ocamlc -c hw2.ml -o hw2.cmo
ocamlc -o hw2 hw2.cmo
```

2. To run your program on the OCAML interpreter, use the command `#use`. Here is a sample session:

```
lngsg@programming2:~/hw2 $ ocaml
OCaml version 4.05.0

# #use "hw2.ml";;
exception NotImplemented
type 'a tree = Leaf of 'a | Node of 'a tree * 'a * 'a tree
val lrevrev : 'a list list -> 'a list list = <fun>
val lfoldl : 'a -> 'b -> 'c -> 'd = <fun>
val fact : 'a -> 'b = <fun>
val fib : 'a -> 'b = <fun>
val alterSum : 'a -> 'b = <fun>
val ltabulate : 'a -> 'b -> 'c = <fun>
val lfilter : 'a -> 'b -> 'c = <fun>
val union : 'a -> 'b -> 'c = <fun>
...
# lrevrev [[1]; [2; 3]; [4; 5; 6]];;
- : int list list = [[6; 5; 4]; [3; 2]; [1]]
```

3. When you have the file `hw2.ml` ready for submission, copy it to your hand-in directory on `programming2.postech.ac.kr`. For example, if your Hemos ID is `foo`, copy it to:

```
/home/class/cs321/handin/foo/
```

## 2 Recursive functions

For this part, do not use any library functions provided by OCAML.

### 2.1 `lrevrev` for returning a reversed list of reversed lists [5 points]

(Type) `lrevrev : 'a list list -> 'a list list`

(Description) `lrevrev l` returns a reversed list of reversed lists.

(Example) `lrevrev [[1, 2, 3], [4, 5, 6], [7]]` returns `[[7], [6, 5, 4], [3, 2, 1]]`.

### 2.2 `lfoldl` for left folding a list [5 points]

(Type) `lfoldl : ('a * 'b -> 'b) -> 'b -> 'a list -> 'b`

(Description) `lfoldl f e l` takes  $e$  and the first item of  $l$  and applies  $f$  to them, then feeds the function with this result and the second argument and so on.

`lfoldl f e [x1; x2; ...; xn]` returns  $f(x_n, \dots, f(x_2, f(x_1, e)) \dots)$  or  $e$  if the list is empty.

(Note) You may not use `List.fold_left`.

### 3 Tail recursive functions

For each description below, give a tail recursive implementation. Usually you want to introduce a tail recursive helper function; the main function is not recursive but just invokes the helper function with appropriate arguments. For example, a tail recursive implementation of `fact` may look like

```
let fact n =  
  let rec fact_aux n acc =  
    ...  
  in fact_aux n 1
```

where `fact_aux` is tail-recursive.

You should fill in the code **only if it is a correct tail recursive implementation**. If you are unsure that it is a correct tail recursive implementation, do not fill in the code.

For this part, do not use any library functions provided by OCAML.

#### 3.1 `fact` for factorials [3 points]

(Type) `fact: int -> int`

(Description) `fact n` returns  $\prod_{i=1}^n i$ .

(Invariant)  $n \geq 0$ .

#### 3.2 `fib` for Fibonacci numbers [3 points]

(Type) `fib: int -> int`

(Description)

`fib n` returns `fib (n - 1) + fib (n - 2)` when  $n \geq 2$ .

`fib n` returns 1 if  $n = 0$  or  $n = 1$ .

(Invariant)  $n \geq 0$ .

(Hint) Perhaps you want to use the idea of dynamic programming?

#### 3.3 `alterSum` for applying addition or subtraction to the elements of a given `int list` alternately [3 points]

(Type) `alterSum: int list -> int`

(Description) `alterSum t` returns an integer that is the result of applying addition and subtraction to the elements of the list  $l$  alternately. The first operation, if applicable, is subtraction. For an empty list, `alterSum` returns 0. For a singleton list, `alterSum` returns its element.

(Example) `alterSum [3, 2, 7, 3]` returns  $3 - 2 + 7 - 3 = 5$ .

`alterSum [1]` returns 1.

`alterSum []` returns 0.

(Hint) You can use a flag variable to indicate which operation should be applied.

### 3.4 ltabulate [3 points]

(Type) `ltabulate : int -> (int -> 'a) -> 'a list`

(Description) `ltabulate n f` applies `f` to each element of a list `[0; 1; ...; n-1]`.

(Example) `ltabulate 4 (fun x -> x * x)` returns `[0; 1; 4; 9]`.

(Invariant)  $n \geq 0$

### 3.5 lfilter for filtering a list [3 points]

(Type) `lfilter : ('a -> bool) -> 'a list -> 'a list`

(Description) `lfilter p l` returns all elements of `l` that satisfies the predicate `p`.

(Example) `lfilter (fun x -> x > 2) [0; 1; 2; 3; 4; 5]` returns `[3; 4; 5]`.

### 3.6 union for union of two sets [5 points]

(Type) `union : 'a list -> 'a list -> 'a list`

(Description) `union S T` returns a set that includes all elements of `S` and `T` without duplication of any element. Note that all list elements have an equality type as indicated by equality type variable `'a`. The order of elements in the return value does not matter.

(Invariant) Each input set consists of distinct elements.

(Example) `union [1; 2; 3] [2; 4; 6]` returns `[3; 1; 2; 4; 6]`.

(Hint) You can implement `union` without introducing an auxiliary tail recursive function. That is, `union` itself can be implemented as a tail recursive function.

### 3.7 inorder for an inorder traversal of binary trees [8 points]

(Type) `inorder : 'a tree -> 'a list`

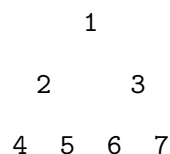
(Description) `inorder t` returns a list of elements produced by an inorder traversal of the tree `t`.

(Example) `inorder (Node (Node (Leaf 1, 3, Leaf 2), 7, Leaf 4))` returns `[1; 3; 2; 7; 4]`.

(Hint) `inorder` can be implemented as follows:

```
let inorder t =  
  let rec inorder' (t' : 'a tree) (post : 'a list) : 'a list = ...  
  in  
    inorder' t [ ]
```

`post` will be a list of elements to be appended to the result of an inorder traversal of `t'`. For example, when `inorder'` visits the node marked 2 in the tree below, `post` will be bound to `[1; 6; 3; 7]`.



### 3.8 postorder for a postorder traversal of binary trees [8 points]

(Type) `postorder: 'a tree -> 'a list`

(Description) `postorder t` returns a list of elements produced by a postorder traversal of the tree *t*.

(Example) `postorder (Node (Node (Leaf 1, 3, Leaf 2), 7, Leaf 4))` returns `[1; 2; 3; 4; 7]`.

### 3.9 preorder for a preorder traversal of binary trees [8 points]

(Type) `preorder: 'a tree -> 'a list`

(Description) `preorder t` returns a list of elements produced by a preorder traversal of the tree *t*.

(Example) `preorder (Node (Node (Leaf 1, 3, Leaf 2), 7, Leaf 4))` returns `[7; 3; 1; 2; 4]`.

## 4 Sorting in the ascending order

For this part, do not use any library functions provided by OCAML.

### 4.1 quicksort for quick sorting [8 points]

(Type) `quicksort: 'a list -> 'a list`

(Description) `quicksort l` implements quick sorting by selecting the first element of  $l$  as a pivot. To compare elements of type `'a`, use the primitive operators `<`, `>`, and `=`.

(Example) `quicksort [3; 7; 5; 1; 2]` selects 3 as a pivot to obtains two sublists `[1; 2]` and `[5; 7]` to be sorted independently.

### 4.2 mergesort for merge sorting [8 points]

(Type) `mergesort: 'a list -> 'a list`

(Description) `mergesort l` divides  $l$  into two sublists, sorts each sublist, and then merges the two sorted sublists. If the length of  $l$  is even, then the two sublists are of equal length. If not, one sublist has one more element than the other. To compare elements of type `'a`, use the primitive operators `<`, `>`, and `=`.

## 5 Structures

The goal of this part is to learn *modular programming* in OCAML — structures and signatures. We will first implement a structure for heaps. You should keep in mind that this data structure is not an ordinary heap data structure. You had better think of it as a mechanism for dynamic memory allocation. See the explanation below carefully.

**Remark.** In the signatures `HEAP` and `DICT`, `empty` is given types `unit -> 'a heap` and `unit -> 'a dict`, respectively. A better design would be to use types `'a heap` and `'a dict`, but in order to facilitate grading, we decided to change it into a function.

For this part, you may use the `List` library of OCAML.

### 5.1 Heap for heaps [10 points]

The structure `Heap` conforms to the signature `HEAP`. A heap is a mechanism for dynamic memory allocation.

```
module type HEAP =
sig
  exception InvalidLocation
  type loc
  type 'a heap
  val empty : unit -> 'a heap
  val allocate : 'a heap -> 'a -> 'a heap * loc
  val dereference : 'a heap -> loc -> 'a
  val update : 'a heap -> loc -> 'a -> 'a heap
end
```

- `loc` is the internal representation of location, which is similar to the *pointer* of C language. `type loc` is not visible to the outside of the structure.
- `'a heap` is a heap for the type `'a`.
- `empty ()` returns an empty heap.
- `allocate h v` allocates the given value `v` to a fresh heap cell and returns the pair  $(h', l)$  of the updated heap  $h'$  and the location  $l$  of this cell.
- `dereference h l` fetches the value `v` stored in the heap cell at location  $l$ . `InvalidLocation` is raised if the  $l$  is an invalid `loc`.
- `update h l v` updates the heap cell at location  $l$  with the given value `v` and returns the updated heap  $h'$ . `InvalidLocation` is raised if the  $l$  is an invalid `loc`.



## 5.2 Signature DICT

DICT is a signature for dictionaries.

```
module type DICT =
sig
  type key
  type 'a dict
  val empty : unit -> 'a dict
  val lookup : 'a dict -> key -> 'a option
  val delete : 'a dict -> key -> 'a dict
  val insert : 'a dict -> key * 'a -> 'a dict
end
```

- `empty ()` returns an empty dictionary.
- `lookup d k` searches the key  $k$  in the dictionary  $d$ . If the key is found, it returns the associated item. Otherwise, it returns `None`.
- `delete d k` deletes the key  $k$  and its associated item in the dictionary  $d$  and returns the resultant dictionary  $d'$ . If the key does not exist in the dictionary  $d$ , it returns the given dictionary  $d$  without any modification.
- `insert d (k, v)` inserts the new key  $k$  and its associated item  $v$  in the dictionary  $d$ . If the key  $k$  already exists in the dictionary  $d$ , it just updates its associated item with the given item  $v$ .

### 5.2.1 Structure DictList [10 points]

Implement the structure `DictList` of signature `DICT` with the definition `'a dict = (key * 'a) list`.

The structure `DictList` uses a list of pairs as the representation of a dictionary. The implementation should be straightforward because a list of pairs itself may be thought of as a dictionary.

### 5.2.2 Structure DictFun [10 points]

Implement the structure `DictFun` of signature `DICT` with the definition `'a dict = key -> 'a option`.

The structure `DictFun` uses a “functional representation” of dictionaries. The idea is that we represent a dictionary as a function that, given a key, returns an associated item. The implementation of `DictFun` may be either very difficult or just a piece of cake depending on how familiar you are with “functional thinking.” Our advice is: forget about everything that you have learned so far about imperative programming; just “think functionally!” You will be amazed at the conciseness of your code once you figure it out.