

Class:	CPE-300L	Semester:	Spring 2020
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		Document topic:	Postlab 8
Instructor's comments:			

1. Introduction / Theory of Operation

- For this lab experiment we familiarized ourselves with the TimeQuest Analyzer and used the GDP and altered various components to implement the algorithm that calculates the greatest common divisor.

a. *What is slack in time quest analysis?*

- Slack is a range value which indicates if the timing requirements are met or not met. Positive slack is when the timing requirements are met (Data received before Required). Negative slack is when the timing requirements are not met (Data received after Required).

b. *Why is it important to consider timing analysis in a digital circuit?*

- It is important to consider timing analysis in a digital circuit because it gives you the needed amount of time in order to get data to various components. If the timing is too fast, then the data may not be able to be received before required in the given time period. Therefore, the timing is crucial for all the components to work in sync.

2. Prelab report

- Attached below

3. Experiment Results

Experiment #1: General Datapath

Verilog Code

```

module GDP(Result, start, restart, clk, n, done);
// top module combines all other modules, including DP and CU
// Note: AlB checks if A < B
//      AeqB checks if A = B

input start, clk, restart;
input [7:0] n;
output [7:0] Result; // 1 or 0
output done;

wire WE, RAE, RBE, OE, AlB, IE;
wire [1:0] WA, RAA, RBA, SH;
wire [2:0] ALU;

CU control (IE, WE, WA, RAE, RAA, RBE, RBA, ALU, SH, OE, start, clk, restart, AlB, AeqB);

DP datapath (AlB, AeqB, Result, n, clk, IE, WE, WA, RAE, RAA, RBE, RBA, ALU, SH, OE);

assign done = OE;

endmodule

module CU(IE, WE, WA, RAE, RAA, RBE, RBA, ALU, SH, OE, start, clk, reset, AlB, AeqB);

```

```

input start, clk, reset;
output IE, WE, RAE, RBE, OE;
output [1:0] WA, RAA, RBA, SH;
output [2:0] ALU;

input wire AlB, AeqB;
reg [2:0] state;
reg [2:0] nextstate;

parameter S0 = 0;
parameter S1 = 1;
parameter S2 = 2;
parameter S3 = 3;
parameter S4 = 4;
parameter S5 = 5;
parameter S6 = 6;

initial
    state = S0;

// State register
always @ (posedge clk)
begin
    state <= nextstate;
end

// NS logic
always @ (*)
case(state)
S0:    if(start) nextstate = S1;
        else    nextstate = S0;
S1:    nextstate = S2;
S2:    nextstate = S3;
S3:    if (AeqB) nextstate = S6;
        else if (~AlB) nextstate = S4;
        else nextstate = S5;
S4:    nextstate = S3;
S5:    nextstate = S3;
S6:    if (reset) nextstate = S0;
        else nextstate = S6;

default: nextstate = S0;
        endcase

// output logic
assign IE = (state == S1) || (state == S2);
assign WE = (state == S0) || (state == S1) || (state == S2) || (state == S4) || (state == S5);
assign WA[1] = (state == S2) || (state == S5);
assign WA[0] = (state == S1) || (state == S4);
assign RAE = (state == S0) || (state == S3) || (state == S4) || (state == S5) || (state == S6);
assign RAA[1] = (state == S5);
assign RAA[0] = (state == S3) || (state == S4) || (state == S6);

assign RBE = (state == S0) || (state == S3) || (state == S4) || (state == S5);
assign RBA[1] = (state == S3) || (state == S4);
assign RBA[0] = (state == S5);

assign ALU[2] = (state == S0) || (state == S3) || (state == S4) || (state == S5);
assign ALU[1] = 0;
assign ALU[0] = (state == S0) || (state == S3) || (state == S4) || (state == S5);
assign SH[1] = 0;
assign SH[0] = 0;
assign OE = (state == S6);
endmodule

```

```

module DP(A1B, AeqB, Result, nIn, clk, IE, WE, WA, RAE, RAA, RBE, RBA, ALU, SH, OE);

    input clk, IE, WE, RAE, RBE, OE;
    input [1:0] WA, RAA, RBA, SH;
    input [2:0] ALU;
    input [7:0] nIn;

    output A1B, AeqB;
    output wire [7:0] Result;

    reg [7:0] rfIn;
    wire [7:0] RFa, RFb, aluOut, shOut, n;

    initial
        rfIn = 0;

    always @ (*)
        rfIn = n;

    mux8 mux (n, shOut, nIn, IE);
    Regfile RF (clk, RAA, RFa, RBA, RFb, WE, WA, rfIn, RAE, RBE);
    alu theALU (aluOut, RFa, RFb, ALU);
    shifter SHIFT (shOut, aluOut, SH);
    buff buffer1 (Result, shOut, OE);

    assign A1B = (aluOut[7]) & (!aluOut[6:0]);
    assign AeqB = ~(!aluOut[7:0]);

endmodule

// final buffer
module buff(output reg [7:0] result, input[7:0] a, input buf1);
    always @(*)
        if(buf1 == 1)
            result = a;
        else
            result = 8'bzzzz_zzzz;
endmodule

// Shifter
module shifter (out,a,sh);
    input [7:0] a;
    input [1:0] sh;
    output reg [7:0] out;

    always @ (*)
    begin
        case(sh)
            3'b00: out=a;
            3'b01: out=a << 1;
            3'b10: out=a >> 1;
            3'b11: out={a[6],a[5],a[4],a[3],a[2],a[1],a[0], a[7]} ;
        endcase
    end
endmodule

```

```

// 2-to-1 mux (sel = 0 -> choose a)
module mux8(result, a, b, sel);

    output reg[7:0] result;
    input[7:0] a;
    input[7:0] b;
    input sel;

    always @(*)
        if(sel == 0)
            result = a;
        else
            result = b;

endmodule

// Regfile for GDP
module Regfile(
    input clk,
    input [1:0] RAA, // Port A Read address
    output [7:0] ReadA, // Port A
    input [1:0] RBA, // Port B Read address
    output [7:0] ReadB, // Port B
    input WE, // Write enable
    input [1:0] WA, // Write port register //address
    input [7:0] INPUT_D, // Write data port
    input RAE, // Port A decoder enable
    input RBE // Port B decoder enable
);
    // width      depth
    reg [7:0] REG_F [0:3];

// Write only when WE is asserted
initial
begin
    REG_F[0] = 0;
    REG_F[1] = 0;
    REG_F[2] = 0;
    REG_F[3] = 0;
end

always @(posedge clk)
if (WE == 1) REG_F[WA] <= INPUT_D;

//reading to Port A and B, combinational
assign ReadA = (RAE)? REG_F [RAA]:0;
assign ReadB = (RBE)? REG_F [RBA]:0;
endmodule
// ALU
module alu (out,a,b,sel);
    input [7:0] a,b;
    input [2:0] sel;
    output [7:0] out;
    reg [7:0] out;

    always @ (*)
    begin
        case(sel)
            3'b000: out=a;
            3'b001: out=a&b;
            3'b010: out=a|b;
            3'b011: out=!a;
            3'b100: out=a+b;
            3'b101: out=a-b;
            3'b110: out=a+1;
            3'b111: out=a-1;
        endcase
    end
endmodule

```

TestBench Code

```
module GDP_tb;
    wire displayRes;
    reg start, restart, clk;
    wire [7:0] A;
    reg [7:0] In;

    initial
    begin
        start = 0;
        restart = 0;
        clk = 0;

        forever
            #2 clk = ~clk;
    end

    always
    begin
        In = 9;
        #5;
        start = 1;
        $write("Input A: %2d " , In);
        #5;
        In = 18;
        #5;
        $write("Input B: %2d " , In);
        while(displayRes != 1) // Wait for completion signal
            #5 begin end

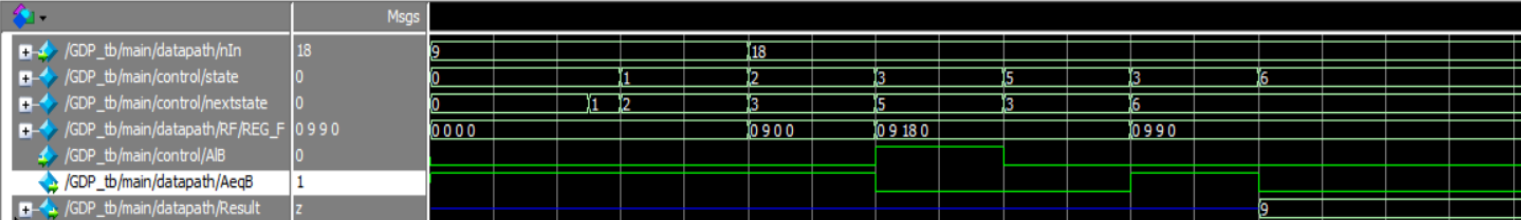
        $write("Output: %2d, ", A);

        if(A == 9) // calculates if A = 9
            $display("Correct!");
        else
            $display("Incorrect!");
        #5;
        restart = 1;
        start = 0;
        #5;

        $stop;
    end
end

//Initialize GDP
GDP main (A, start, restart, clk, In, displayRes);
endmodule
```

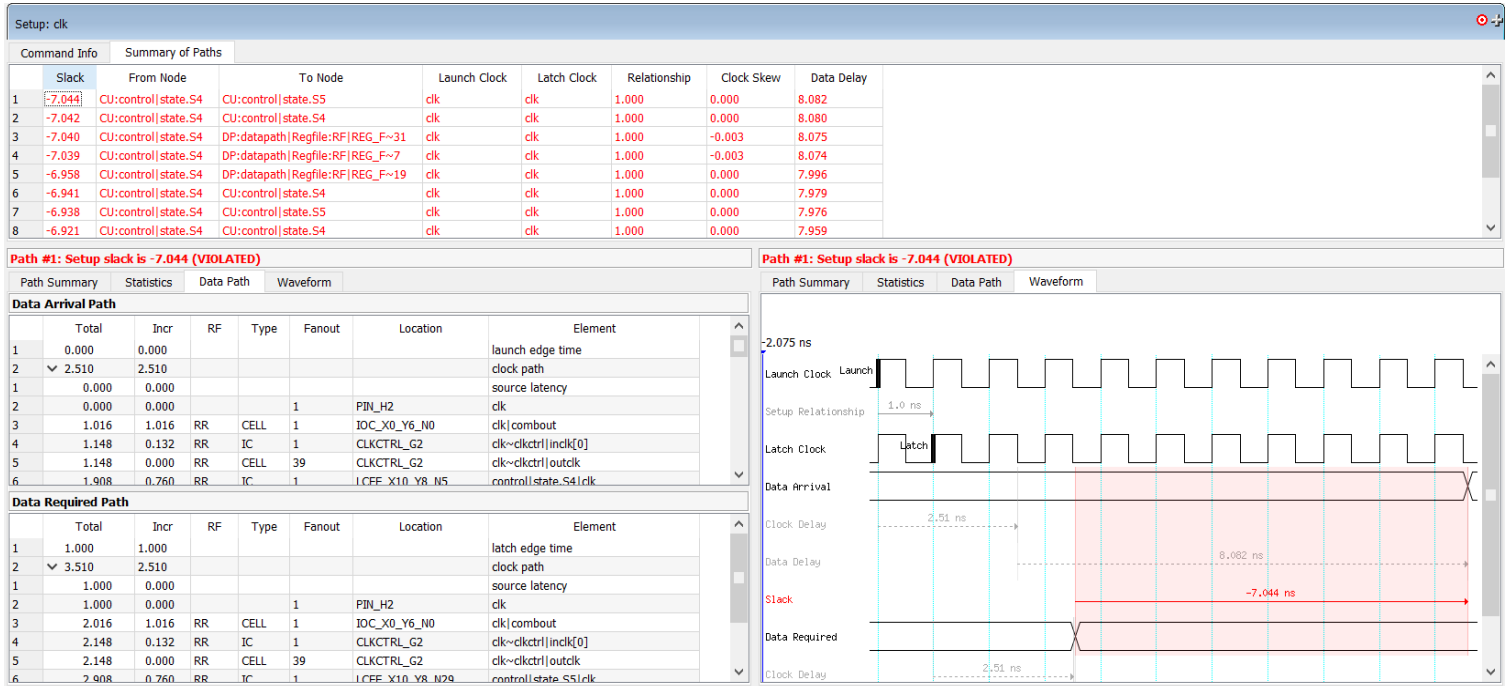
Waveform



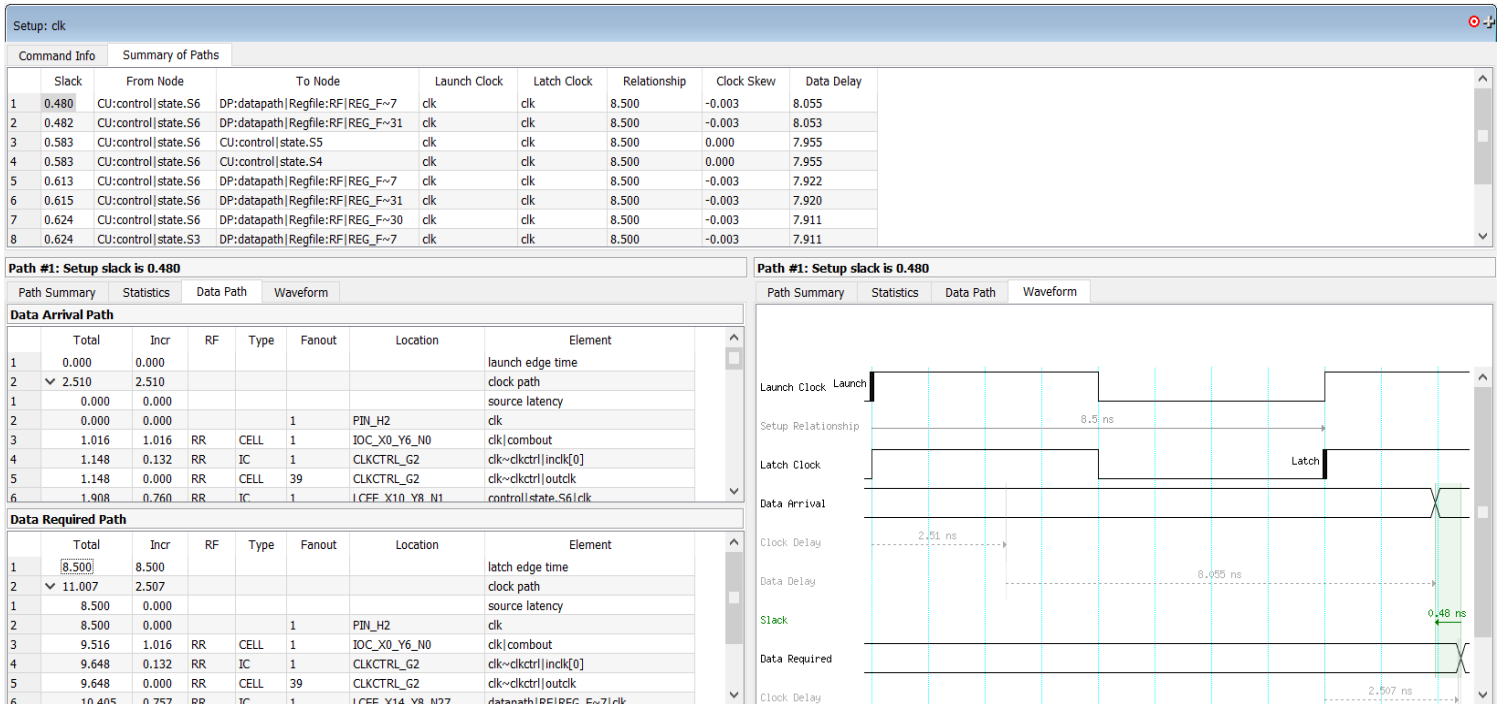
Unable to grab DE2 board due to lab being closed during COVID-19 shutdown

Experiment #2: TimeQuest analysis

- Results of TimeQuest with 1ns clock period. Received a skew of **-7.044**, therefore the timing requirements are not met. Resulting in data being received **after** it is required.



- Results of TimeQuest with 8.5ns clock period. Received a skew of **0.480**, therefore the timing requirements are met. Resulting in all data being received **before** it is required.



4. Questions

1. *Find out the information about the TimeQuest analyzer.*

- TimeQuest analyzer has many powerful tools to specify timing constraints for simple and complex digital designs. The TimeQuest timing analyzer compilation report shows you various tools and setting which you can select and view. For this lab experiment we selected the clock and viewed the TimeQuest GUI. The TimeQuest GUI allowed us to see our slack, from node to node, setup relationship, latch clock, arrival of data, clock delay, clock skew, data delay and when data is required. One also has the ability to see any desired path from the design and calculate the proper clock period and frequency in order for the data to be successfully received before being required for that path.

The TimeQuest analyzer GUI allowed us to visually see when our data is being received and when it is required. We used the slack timing from the TimeQuest analyzer report to see exactly what the clock period should be in order to fulfil the timing requirements for our design. Our original clock period was 1ns but after looking at the TimeQuest analyzer, we concluded that the proper clock period was 8.5ns for our design to fulfil the timing requirements.

2. *List the purposes of initializing the memory in Quartus II*

- It is important to initialize memory in Quartus II because it allows the storage of the data in the memory to be neater. Another important purpose to initialize memory in Quartus II is so that when you want to see how where in memory data is being stored, it is easier to reference when you initialize the memory at a specific location.

5. Conclusions

- For this lab experiment we used the TimeQuest analyzer and altered various components of the GDP in order to implement the algorithm that calculates the greatest common divisor. A problem that we encountered was that we did not finish within the allocated time, so we had to meet up outside of the lab to finish. Another encountered problem was trying to find the proper amount of delay within the test bench for verifying if our algorithm worked. But after messing around with the different delay amounts, we were able to properly verify the testbench with the waveform. Overall this lab taught us how to use the TimeQuest analyzer and how to manipulate the GDP and control unit in order to implement the algorithm that calculates the greatest common divisor.

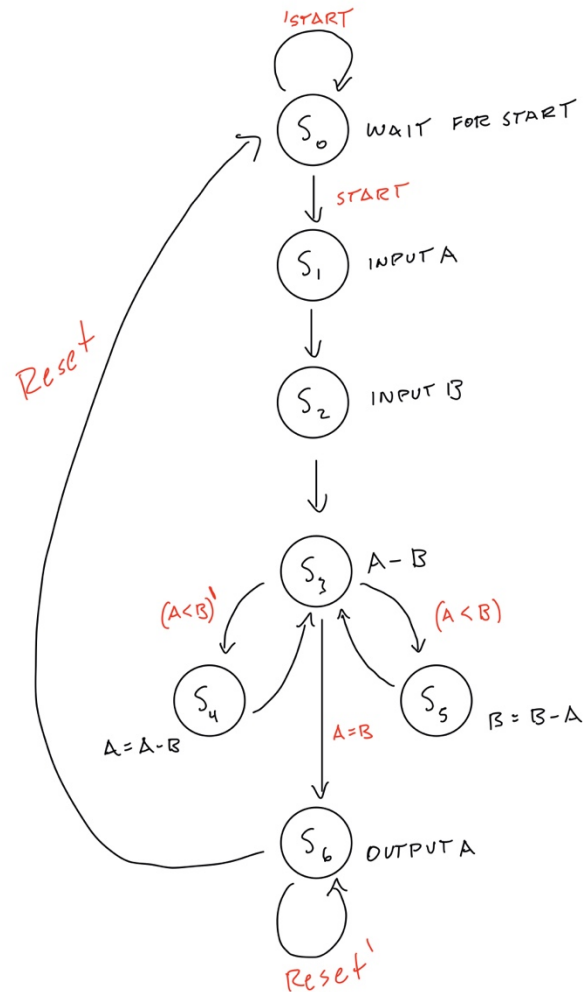
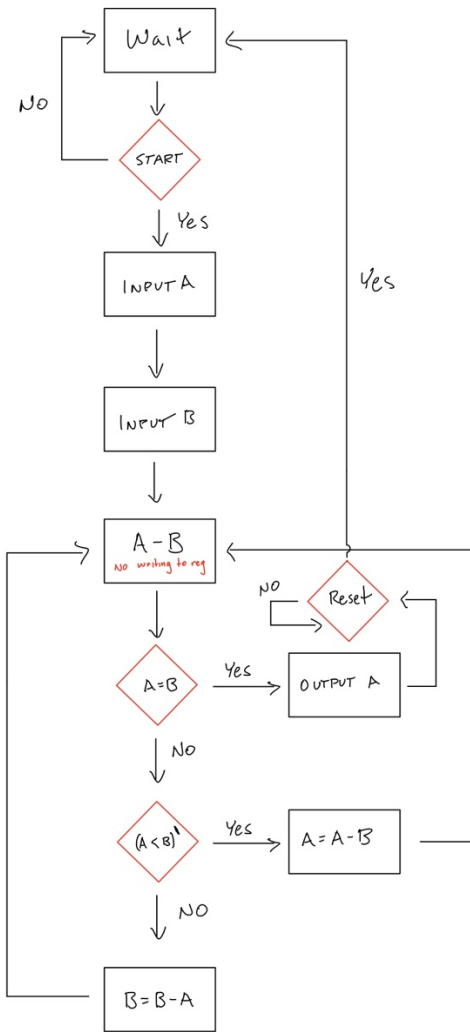
- For this lab experiment we will be using the GDP that was created in lab experiment 6, to implement the algorithm that calculates the greatest common divisor of 2 natural numbers. We will also be learning how to set up timing constraints and how to obtain timing information with the TimeQuest analyzer.

1. *Screenshots of the timing report.*

Setup: dclk								
Command Info		Summary of Paths						
	Slack	From Node	To Node	Launch Clock	Latch Clock	Relationship	Clock Skew	Data Delay
1	0.691	reg_B[1]	reg_sum[9]~reg0	dclk	dclk	4.000	0.324	3.630
2	0.697	reg_B[1]	reg_sum[9]~reg0	dclk	dclk	4.000	0.324	3.624
3	0.698	reg_B[2]	reg_sum[8]~reg0	dclk	dclk	4.000	0.324	3.623
4	0.698	reg_B[1]	reg_sum[9]~reg0	dclk	dclk	4.000	0.324	3.623
5	0.698	reg_B[2]	reg_sum[8]~reg0	dclk	dclk	4.000	0.324	3.623
6	0.701	reg_A[1]	reg_sum[9]~reg0	dclk	dclk	4.000	0.324	3.620
7	0.707	reg_A[1]	reg_sum[9]~reg0	dclk	dclk	4.000	0.324	3.614
8	0.708	reg_A[1]	reg_sum[9]~reg0	dclk	dclk	4.000	0.324	3.613

Path #1: Setup slack is 0.691

2. Flowchart, state graph and control words for the algorithm



state	IE	WE	WA(10)	RAE	RAA(10)	RBE	RBA(10)	ALU(210)	SH(10)	OE	
s0	0	1	00	1	00	1	00	101	xx	0	//Wait for start
s1	1	1	01	0	xx	0	xx	000	xx	0	//Input A
s2	1	1	10	0	xx	0	xx	xxx	xx	0	//Input B
s3	0	0	xx	1	01	1	10	101	00	0	//A-B no write to Reg
s4	0	1	01	1	01	1	10	101	00	0	//A = A-B
s5	0	1	10	1	10	1	01	101	00	0	//B = B-A
s6	x	0	xx	1	01	0	xx	000	00	1	//Output A