# Writing a self-hosting compiler for a purely functional language

(work in progress)

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### What is this?

This is an **experiment** to write a minimalistic self-hosting compiler in a purely functional language. It's intended both as a learning experience (for myself) and an experiment in bootstrapping.

#### Goals:

- self-hosting compiler
- very small code-base
- bootstrappable using a Haskell implementation
- generated code should have "acceptable" quality

#### Non-goals:

- a fully-featured language
- efficient compilation (the compiler can be slow)
- bootstrapping "real" compilers from this code base



# Why self-host?

Why do bootstrapping / self-hosting?

- I like simple, self-contained stuff
- self-hosting is a philosophically desirable property
- future-proofness: If you can bootstrap from something simple, it won't bit-rot that easily
- it's a test of the maturity of the compiler
- the compiler itself is a very good test program!

Exhibit A: Maru, a tiny self-hosting LISP dialect was successfully revived and extended by Attila.

Exhibit B: The York Haskell Compiler (forked from nhc98) is now dead, apparently because nobody knows how to build it anymore.

# The language

The syntax and semantics is based on (strict) Haskell - the idea is that the compiler should be compilable by both GHC and itself. This makes development much easier!

The language is basically untyped call-by-value lambda calculus with data constructors and recursive lets. During parsing, Haskell features like type annotations, data type declarations, imports etc are simply ignored - but the compiler itself is a well-typed program!

In fact, a type checker should be probably added, after all; but I wanted to keep it simple.

# Language (non-)features

- no static type system (untyped lambda calculus)
- na data type declarations (constructors are arbitrary capitalized names)
- no module system instead, we have C-style includes
- strict language (exceptions: if-then-else, logical and/or)
- ML-style side effects, but wrapped in a monad
- only simple pattern matching + default branch (TODO: nested patterns)
- no infix operators
- list construction syntax [a,b,c] is supported
- no indentation syntax (only curly braces), except for top-level blocks
- only line comments, starting at the beginning of the line
- no escaping in character / string constants yet
- (universal polymorphic equality comparison primop)



# Compilation pipeline

- 1. lexing and parsing
- 2. collecting string literals
- 3. partition recursive lets into strongly connected components
- 4. TODO: eliminate pattern matching into simple branching on constructors
- 5. recognize primops
- 6. collect data constructors
- 7. scope checking & conversion to core language
- 8. inline small functions + some basic optimizations
- closure conversion
- 10. TODO: compile to an SSA intermediate language
- 11. final code generation

## The runtime system

#### The runtime system is relatively straightforward:

- there is a heap, which can only contain closures and data constructors
- there is stack (separate from the C stack) which contains pointers to the heap
- heap pointers are tagged using the lowest 3 bits
- heap objects fitting into 61 bits (ints, chars, nullary constructors, static functions) are not allocated, just stored in registers / on stack
- statically unknown application is implemented as a runtime primitive
- there is a simple copying GC the GC roots are simply the content of the stack (plus statically known data)
- ▶ there are some debug features like printing heap objects
- primops: integer arithmetic; integer comparison; bitwise operations; lazy ifte/and/or; basic IO.



# Current state of the experiment

Repo: https://github.com/bkomuves/nanohs

#### Current state:

- it can compile itself successfully;
- the source code is not as nice as it could be;
- the generated code is pretty horrible;
- there are some desirable features missing.

#### Code size:

- ▶ the compiler:  $\sim 1900$  lines (required code)
- ightharpoonup type annotations: +525 lines (ignored)
- the C runtime:  $\sim 600$  lines (including some debug features)

#### Mistakes I made

I made a huge amount of mistakes during the development, which made it much harder and longer than I expected.

- Trying to self-host before the compiler actually works. Writing a compiler is hard work, you don't want to do it in a limited language!
- Not figuring out the precise semantics upfront:
  - exact details of lazyness vs. strictness
  - controlling exactly where heap / stack allocations can happen
  - recursive lets
- ▶ Not having debugging support from early on (for example: good error messages, threading names & source locations, RTS features, pretty printing AST-s)
- Trying to target assembly first, instead of a simpler target like C; it's just extra cognitive load
- Generatic code directly, without having a low-level IL
- Using de Bruijn indices (as opposed to levels) in the generated code



# Example GC bug - Marshalling from C strings

```
heap_obj marshal_from_cstring(char *cstring) { ...
v1:
      obj = heap_allocate(Cons,2);
      obj[1] = cstring[0];
      obj[2] = marshal_from_cstring(cstring+1);
      return obj;
v2:
      rest = marshal_from_cstring(cstring+1);
      obj = heap_allocate(Cons,2);
      obj[1] = cstring[0];
      obj[2] = rest;
      return obj;
v3:
      rest = marshal_from_cstring(cstring+1);
               push(rest);
      obj = heap_allocate(Cons,2);
      obj[1] = cstring[0];
      obi[2] = pop();
      return obj;
```

#### Future work

I'm not convinced that it is worth hacking on this toy compiler (as opposed to work on a full-featured "real" compiler), but in any case, here is a TODO list:

- add nested patterns, and patterns in lambda arguments
- do some optimizations, so that the generated code is better
- tail call elimination
- escaping in string constants
- compile to a low-level intermediate language first, and do codegen from that
- add a type system
- more backends: x86-64 assembly and/or machine code; some simple VM