## Dive deeper into Hashing! With each hash, we're carving a path through the jungle of problems!



# Hashing 2



### Agenda:

- 1. Target Sum Pair
- 12. Count no of pair with sum= K
  - 2. Suborray with sum = K
  - 4. No. of subarray having sum = k
  - a. Distint no. in every window of size k.

```
Ouick Revision:
I. Hash Map ?
         HashMap < Key, Value> map = new HashMap< >C);
* create
                      Doubtypes
* Insert map. put ( key, Value);
   Accels map.get(key); -> 9+ will return value or null.
* Check Key map contains Key (Key)1 - True / False
  Remove key-volue map. remove (key);
   Size of hmap map size(); - no of keys in map.
   All keys from Set -> map. keySet(); -> 91 returns all keys in Set-
                             gterate with help of for-each loop.
2. Hash Set:
 * create Hash Set < Key dottatype > set = new Hash Set <> C);
 * Insert Set add ( key);
  * Check key set contains (key); - For False
 * Remove Key Set ramove ( Key);
 * Size of heet set size():
```

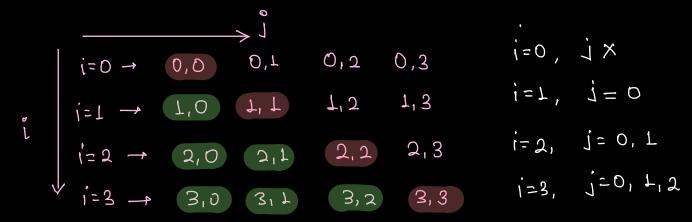
#### Problem 1:

Given an array A[] and a value K. Check if there exist a Pair (i, j) such that A[i] + A[j] = K and i != j.





all possible pairs;



```
boolean pairsum ( int | am, int k) {

int n= archength;

forcint i=0: i<n; i+t) {

int tanget = k-anci];

forcint j=0; j<i; j+t) {

if anci] == tanget ) {

Te: o(n²)

3

3

return false;
```

pair Sm ( int[] om, int K) { booleam int no arriength; K=6 for (int i = 0: i<n; i++) { 8 -2 7 11 int tanget = K-amii]; cm: 0 2 3 Ч for (int j=0; j<1; j++){ if( am[j] = = target )} K-amrij retorn true; 0 tenget 3 -2 Q O'TX L ر ع return false; 0, 1,2X ЬØ 0- retur B പ്പ me:

# optimis ation using Hashset:

cm: 8 | 3 | 1 | -2 | 11 | K= 6

Q I	tenget-≥K-ami)	Search that in left
0	6-8=-2	×
L	6-9 = -3	×
2	6-1=5	×
$\mathcal{M}$	6-(-2)=8	- redum true.

Hashset 8 9

bookern Pairsum ("int() am, Prot K) {

int n= am.length;

Hashset < Integer > set = new Hashset > ();

for (int i=0; i<n; i+1) {

int tenget = K-amii;

if ( set · contains (tenget)) {

I return toue;

}

set · add (amii);

Topo: Dry Rum.

return folse;

#### Problem 2:

Given an array A[] and a value K. Count total number of pairs (i, j). Such that A[i] + A[j] = K and i != j.

 $(i,j) \rightarrow (1,3)$ , (0,4), (2,4) (1,5), (3,6), (4,6), (0,7), (2,7), (6,7)count = 9 paix

NOTE: same logic of lost problem, only difference is insteadle of presence of element we have to marge court of occurrence of element.

am;	ಭ	S	7	2	T	2
	0	Τ	2	3	Ч	S

K=3

0	ATiJ	tenget = K-amrij	
٥	s W	3-3=0	Count = &
1	rs,	3-5=-2	2
28	$\mathcal{T}$	3-1=2	4
$\mathcal{A}$	2	3-2-1	
4	L	3-1 = 2	

3-2 = 1

B

2

hashmays. — Koy + Elent

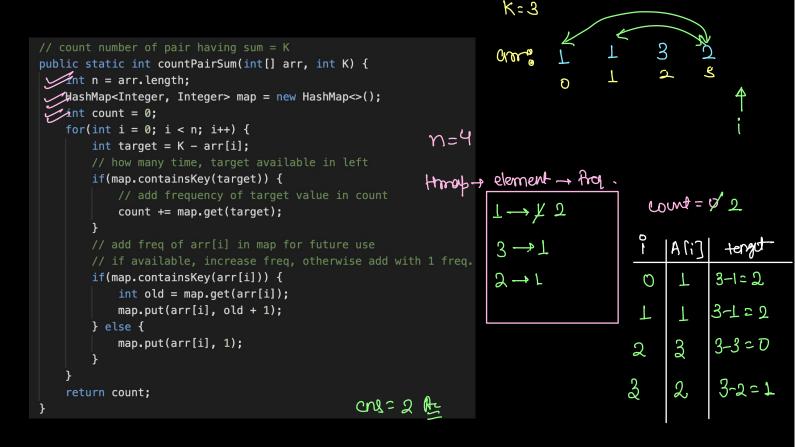
3>1

5-1

1-1 × 2

2-1 × 2

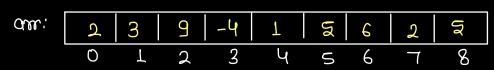
```
// count number of pair having sum = K
public static int countPairSum(int[] arr, int K) {
    int n = arr.length;
    HashMap<Integer, Integer> map = new HashMap<>();
    int count = 0;
    for(int i = 0; i < n; i++) {
        int target = K - arr[i];
                                                          T.C: O(n)
        // how many time, target available in left
                                                           S.C:0(n)
        if(map.containsKey(target)) {
            // add frequency of target value in count
            count += map.get(target);
        // add freq of arr[i] in map for future use
        // if available, increase freq, otherwise add with 1 freq.
        if(map.containsKey(arr[i])) {
            int old = map.get(arr[i]);
            map.put(arr[i], old + 1);
        } else {
            map.put(arr[i], 1);
    return count;
```



#### Problem 3:

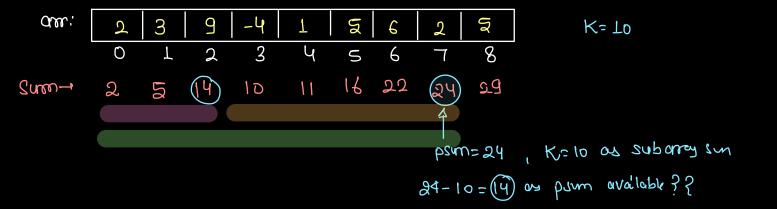
Given an array A[] and a value K. Check if there exist a Subarray with Sum = K.

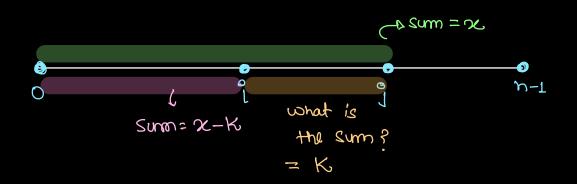
Continue part of

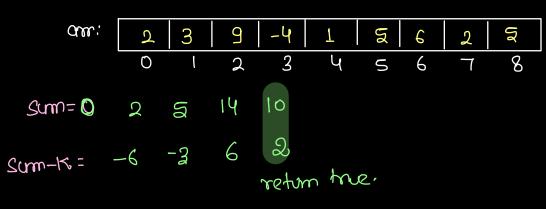


om: 
$$5$$
 10 20 100 105  $K=110$ 

NOTE: I. We com do it in T.C: O(n2) {Refer Suborry + prefix sum}





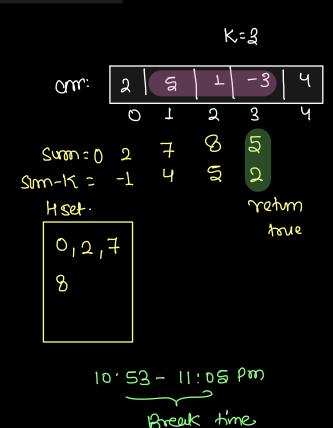


```
K=8
Hachset
O - for edge ~
2,5,14
```

```
// subarray having sum = K
public static boolean subArraySumK(int[] arr, int k) {
    int n = arr.length;
    HashSet<Integer> set = new HashSet<>();
    int sum = 0;
    set.add(0);
                                          T.C: 0(n)
    for(int i = 0; i < n; i++) {
        sum += arr[i];
                                          S.C: O(D)
        int target = sum - k;
        if(set.contains(target)) {
            return true;
        set.add(sum);
    return false;
}
```

```
// subarray having sum = K
public static boolean subArraySumK(int[] arr, int k) {
    int n = arr.length;
    HashSet<Integer> set = new HashSet<>();

    int sum = 0;
    set.add(0);
    for(int i = 0; i < n; i++) {
        sum += arr[i];
        int target = sum - k;
        if(set.contains(target)) {
            return true;
        }
        set.add(sum);
    }
    return false;
}</pre>
```



#### Problem 4:

Given an array A[] and a value K. Count total number of Subarray with Sum = K.

am: \_4 2 5 0 2 3 T6 7

> pairs - [1-4] count=2 Ps [6-6]

am:

(2-3) (4-4) (0-0) (0-3)(L-4)

Cout = S N

K= 8

-8 B B am: 6 2 L 4 2 3

16 6 જ 8 0 Sum=0 8  $\bigcirc$ -2 -8  $\bigcirc$ sum-K=

K=8

courd = &

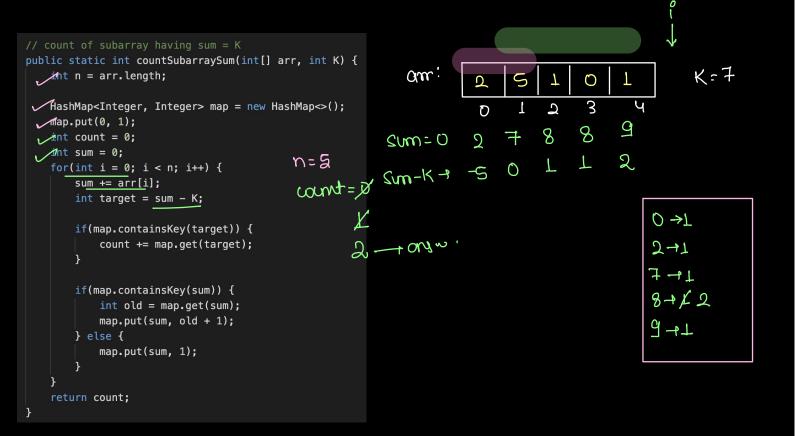
psum vs. foeg.

0 - X 2 8 <del>-></del> ¥2

K=6

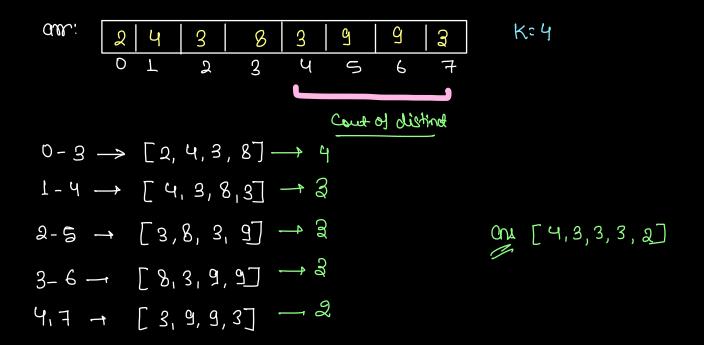
6 -> T 16-1 L

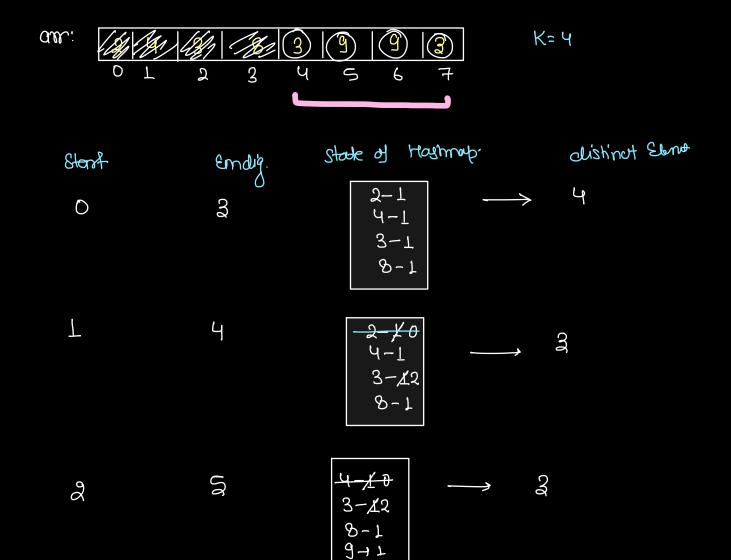
```
// count of subarray having sum = K
public static int countSubarraySum(int[] arr, int K) {
    int n = arr.length;
   HashMap<Integer, Integer> map = new HashMap<>();
   map.put(0, 1);
    int count = 0;
    int sum = 0;
    for(int i = 0; i < n; i++) {
        sum += arr[i];
        int target = sum - K;
        if(map.containsKey(target)) {
            count += map.get(target);
        if(map.containsKey(sum)) {
            int old = map.get(sum);
            map.put(sum, old + 1);
        } else {
            map.put(sum, 1);
    return count;
}
```

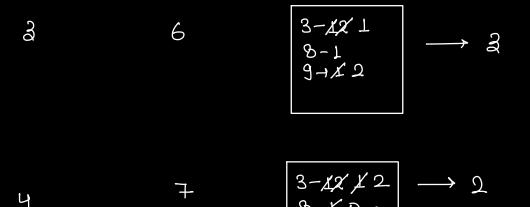


### Problem 5:

Given an array A[] and a value K. Calculate no. of distinct elements in every window of size K.







5 (3) -1 100/2 stol.

```
8-4+1=5
public static int[] distinctElement(int[] arr, int K) {
                                                                                                indx
  int n = arr.length;
 int[] ans = new int[n-K+1];
                                                         ONY:
 int indx = 0;
                                                                                     3
                                                                  \mathcal{O}
                                                                              Q
HashMap<Integer, Integer> map = new HashMap<>();
                                     am:
                                                                                                            K= 4
                                                                                      (9)
                                                                        (3)
                                                                                            (3)
     for(int i = 0; i < K; i++) {
        if(map.containsKey(arr[i])) {
                                                                   2
                                                                          ч
                                                                                       6
            int old = map.get(arr[i]);
            map.put(arr[i], old + 1);
        } else {
            map.put(arr[i], 1);
                                    n=8
  //mns[indx] = map.size();
                                                                                       map "
    indx++;
                                                                                       2-/0
    int s = 1;
                                                                                       4-10
    int e = K;
                                                                                        3-12 X2
        if(map.containsKey(arr[e])) {
            int old = map.get(arr[e]);
            map.put(arr[e], old + 1);
        } else {
            map.put(arr[e], 1);
                                                    T.C:0(n)
        int old = map.get(arr[s-1]);
                                                     5.C:0(n)
        map.put(arr[s-1], old-1);
        if(map.get(arr[s-1]) == 0) {
            map.remove(arr[s-1]);
        ans[indx] = map.size();
        idnx++;
       Stt, ett;
    return ans;
```

