Todays Content: a. Double linked list

b. LRV Cache

Double linked list

```
class Node of int data

Node nent / // obj reference can hold

Node prev / address of node objects

Noac Cint n) {

data = n

nent = null

prev = null

3
```



obs: We can travel from left - Right a Right - left.

Note: Noch reference is given

Note: Noch reference is given

Note: Given noch is not head (tax) noch. Inf: No need to worm there edge cases

Notes: Linked list is not null

Ent: Delek #ads //direct address is given.

#adi #ads #ads #ady #ads #ads

Null

Null

Pads Delek Noch (Name town) & To: 0000 cases

Null

Void Delete Noce (Noce temp) & Tc: O(1) Sc: O(1)

Noce ti= temp. prev;

Noce ti= temp. nent;

ti. nent = ta;

ta. prev = ti;

temp. nent = temp. prev = NULL;

}

Obs: In Double lighted list, given a nocle address, delete it: 0(1) unlike in single linked list

Delete it valu: Note: Assume data are distinct.

#adi #adi #adi #adi #adi #adi #adi

NULL NULL

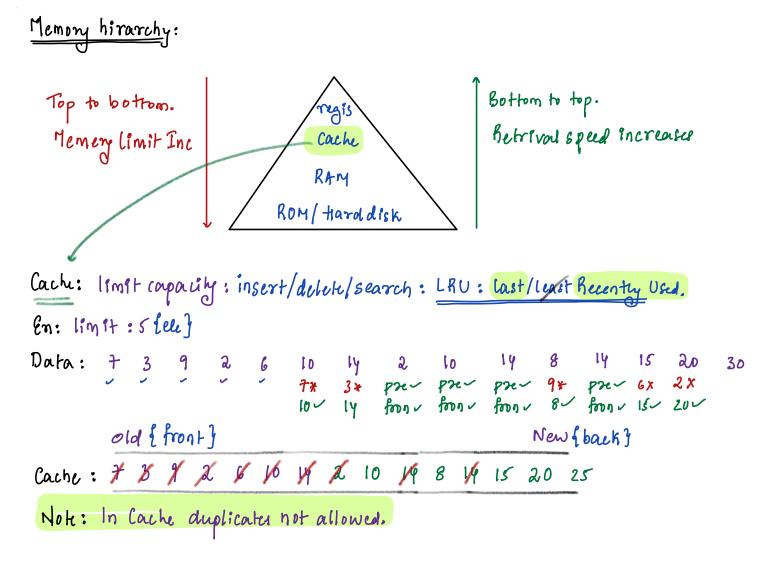
NULL

HM: { LIU, adı > L20, adz > L30, ad3 > L40, ad4 > L50, ads > L60, ad6 > }

Delete: 40, get address of 40 using hashmap = ady.

Delete ad4 in double linkedlist.

20. Insert a new node Just before toud of a Double Linked list Moke: Tail ref is given in Input Notel: No: of nody > = 2 Notes: New nocle is already created a address given. #ad1 #ada NULL NULL nn:#ads 80 p new node. Insert befre tail (Noue nn, Noce tail) { TC: O(1) SC: O(1) 4/1 tail nocce Node ti= tail.prev tail.prev=nn; nn.prev= tij ti. nent = nnj nn.nent = tail ;



Search(n)

insert-back(n)

Size cache == limit

delete-front()

insert-back(n)

Insent n: -

delete(n)

insert-back(n)

Flow chart:

Design and implement a data structure for Least Recently Used (LRU) cache. It should support the following operations: get and set.

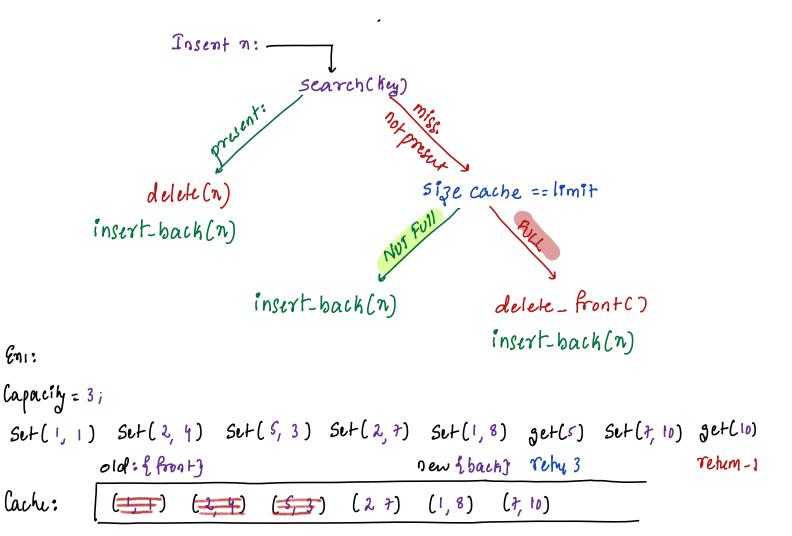
get (key) - Get the value (will always be positive) of the key if the key exists in the cache, otherwise return -1. set (key, value) - Set or insert the value if the key is not already present. When the cache reaches its capacity, it should invalidate the least recently used item before inserting the new item.

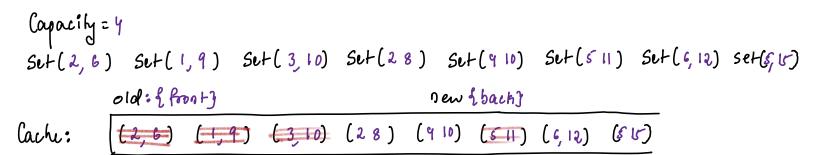
The LRUCache will be initialized with an integer corresponding to its capacity. Capacity indicates the maximum number of unique keys it can

Definition of "least recently used": An access to an item is defined as a get or a set operation of the item. "Least recently used" item is the one with the oldest access time.

NOTE: If you are using any global variables, make sure to clear them in the constructor.

En:





```
Obs:
```

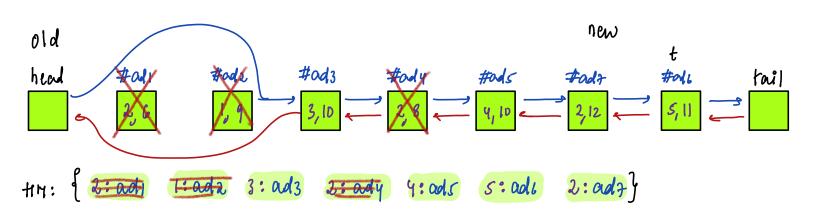
1: Order of Insertin is need:

2: Deleting from middle - of Double linked list?
Note: Because we are only given key to delete
We also need to store & key, address in DIL > in hashmap.

LAU Couche using DLL+ Hash Marp

Note: To avoid edge cases, Create head a tail.

Capacity=4
Set(2,6) Set(1,9) Set(3,10) Set(28) Set(410) Set(511) Set(2,12) get(5)



Noa t = hmoget(s) // ad6

delelelt);

before-tail(t);

return toval;

```
1 - public class Solution {
2 -
        class Node{ // Node object
3
            int key;
4
            int value;
5
            Node prev, next;
6 -
            Node(int k, int v){
7
                key = k;
8
                value = v;
9
                prev = null;
10
                next = null;
11
12
        Node head = new Node(-1,-1);
13
        Node tail = new Node(-1,-1);
14
15
        HashMap<Integer, Node> hm = new HashMap<>();
16
        int cap = 0; // Global variable.
17 -
        public Solution(int capacity) {
18
            cap = capacity; // Initialize Global with local.
19
            head.next = tail;
20
            tail.prev = head;
21
22 -
        public int get(int k) {
23 -
            if(hm.containsKey(k) == true){
                // K is present, we acesses its adrees from hashmap.
24
25
                // Where ever k is present delete it and add it at back{before tail}.
26
                Node t = hm.get(k); // current adress
                delete(t); // delete current t
27
                before_tail(t); // adding t node before tail
28
29
                return t.value;
30
31 -
            else{
32
                return -1;
33
34
35 -
        public void set(int k, int v) {
36 -
            if(hm.containsKey(k) == true){
37
                // K is present, we acesses its adrees from hashmap.
38
                // Where ever k is present delete it and add it at back{before tail}.
39
                Node t = hm.get(k); // current node adress
40
                delete(t); // delete current t
41
                before_tail(t);
42
                t.value = v; // we are updating value;
43
44 -
            else{ // K is not present
45
46 -
                if(hm.size() == cap){ // We need to delete from old.
47
                    Node t = head.next;
48
                    delete(t); // Delete t from linkedlist
49
                    hm.remove(t.key); // Delete k and value from hashmap/
50
51
52
                Node nn = new Node(k,v); // New node we are creating
                before_tail(nn);
53
54
                hm.put(k,nn); // insert key and address in the hashmap
55
56
57
        public void delete(Node temp){ // Delete the given temp node.
58 -
59
            Node t1 = temp.prev;
60
            Node t2 = temp.next;
61
            t1.next = t2;
62
            t2.prev = t1;
63
            temp.next = null;
64
            temp.prev = null;
65
66
        public void before_tail(Node nn){ // Adding node before tail.
67 -
            Node t1 = tail.prev;
68
69
            tail.prev = nn;
70
            nn.prev = t1;
71
            t1.next = nn;
72
            nn.next = tail;
73
74
```