# Medusa

## Mutant Equivalence Detection Using SAT Analysis



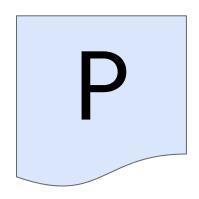
Benjamin Kushigian<sup>1,2</sup>, Amit Rawat<sup>1</sup>, René Just<sup>2</sup>

<sup>1</sup>UMass Amherst

<sup>2</sup> University of Washington



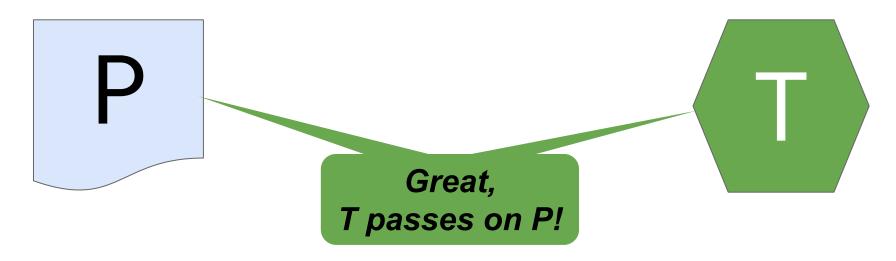
## **Question:** Is our program correct?



## Let's test its correctness

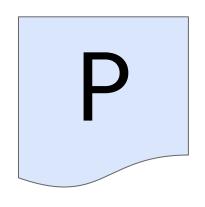


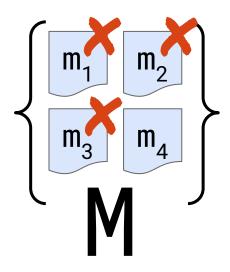
## Question: how good is our test suite?

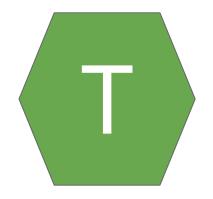


Let's analyze T's effectiveness

## Let's analyze T's effectiveness







#### **Mutation Analysis**

- Generate mutant set M
- Run *T* on mutants in *M*
- Each mutant that fails a test in T is killed

## **Mutation Score**

For test suite *T* and mutant set *M*, the *mutation score* of *T* is

$$\mu(T) = \frac{\left| \left\{ \text{ Killed mutants in } M \right\} \right|}{\left| \left\{ \text{ Nonequivalent mutants in } M \right\} \right|}$$

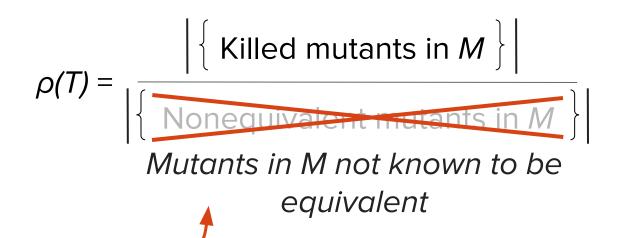
## **Mutation Score**

For test suite *T* and mutant set *M*, the *mutation score* of *T* is

$$\mu(T) = \frac{\left| \left\{ \text{ Killed mutants in } M \right\} \right|}{\left| \left\{ \text{ Nonequivalent mutants in } M \right\} \right|}$$

## A common proxy: Mutation Kill Ratio

For test suite *T* and mutant set *M*, the *mutant kill ratio* of *T* is



# Mutation Kill Ratio vs. Mutation Score

- ρ(T) approximates μ(T)
   Undetected Equivalents are detected malysis
   Mutants Skew Amalysis
   Mutants Skew Amalysis
   proximates μ(T) Resulting tants, the worse ρ(T) approximates μ(T) Resulting tants

## Say we have an equivalent mutant



- A human developer tries to kill m<sub>A</sub>
  - This wastes the developer's time

# Detecting Equivalent Mutants

## **Detecting Equivalent Mutants**

There are a number of ways to detect equivalent mutants

- 1. Hand proof
- 2. Compiler techniques
- 3. Automated reasoning tools

## **Detecting Equivalent Mutants**

Two Main Challenges

- Applicability: What programs can a system reason about?
- Efficiency: How many resources does a system need?

## Applicability: What can be reasoned about?

#### **Applicability challenges:**

- Loops and recursion
  - Loop unrolling can witness non-equivalence but cannot prove equivalence in general

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## Applicability: What can be reasoned about?

#### **Applicability challenges:**

- Loops and recursion
  - Loop unrolling can witness non-equivalence but cannot prove equivalence in general
- Heap space

#### Approaches to increase applicability (future work described in the paper):

- Middle Out Constraint Generation
- Exception Abstraction
- Foldability: is (fold f xs acc) equivalent to (fold f' xs acc)?

## Efficiency: How many resources are needed?

#### Efficiency challenge:

Solving NP-hard problems such as SAT

### Medusa improves efficiency with constraint forking

- Decreases problem size
- Reuses work done by SMT solver

# Proving Program Equivalence

**Question: are these equivalent?** 



```
int squarePos(int x) {
   if (x > 0) x = x * x;
   return x;
```

```
int squarePos(int x) {
   if (x \le 0) x = x * x;
  return x;
```





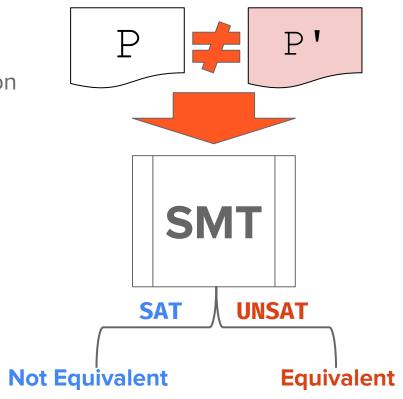


ORIGINAL

NOT EQUIVALENT MUTANT

## Question: How to (dis)prove equivalence?

- Transform query to SMT constraints
  - Constrain both programs' execution
  - Assert inputs are equal
  - Assert outputs are different
- Ask SMT if constraints are satisfiable
  - o SAT => Not Equivalent
  - OUNSAT => Equivalent
  - May also be UNKNOWN



## **Question: How to constrain Java?**

#### **Observation**

- Java is complicated
- No formal semantics
- Semantics defined by Javac/JVM

#### We Should:

- Compile program
- Constrain bytecode

## Question: How to constrain bytecode?

Two types of bytecode instructions

### **Straight Line**

- Operand Stack Instructions iadd, imul, push, pop...
- Variables Table Instructions istore, iload, ...

### **Branching**

• Control Flow Instructions ifgt, iflt, if cmpgt, ...

Distinction captured by Control Flow Graphs

```
int squarePos(int x) {
   if (x > 0) x = x * x;
   return x;
}
```

```
int squarePos(int x) {
   if (x <= 0) x = x * x;
   return x;
}</pre>
```

Given a program and a mutant:

1. Compile to bytecode

```
0: iload_1
1: ifle 8
4: iload_1
5: iload_1
6: imul
7: istore_1
8: iload_1
9: ireturn
```

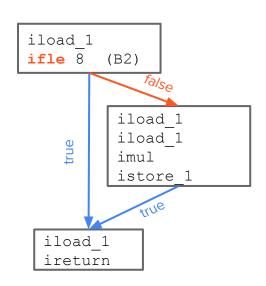
```
0: iload_1
1: ifgt 8
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```

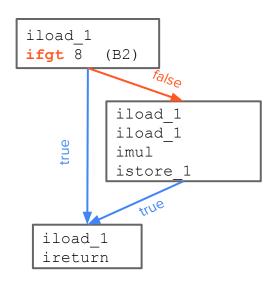
ORIGINAL

MUTANT

#### Given a program and a mutant:

- 1. Compile to bytecode
- 2. Construct CFG

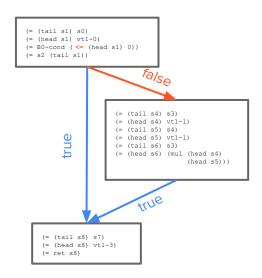


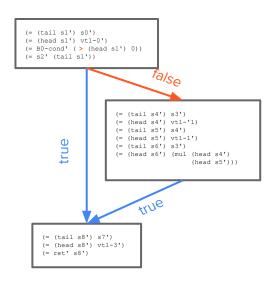


ORIGINAL



- 1. Compile to bytecode
- 2. Construct CFG
- 3. Generate constraints

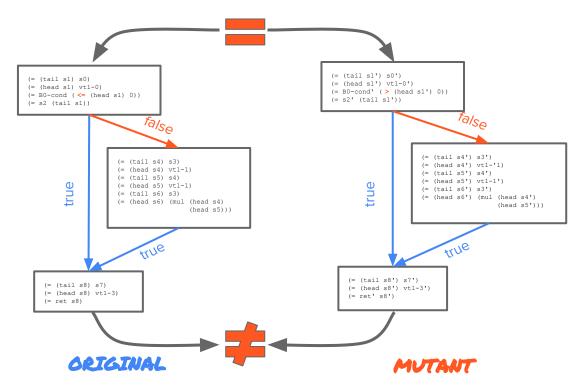




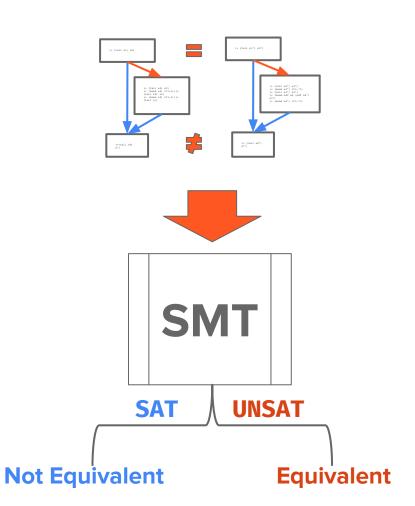




- 1. Compile to bytecode
- 2. Construct CFG
- 3. Generate constraints
- 4. Equivalence Query
  - Assert inputs are equal
  - Assert outputs aren't equal

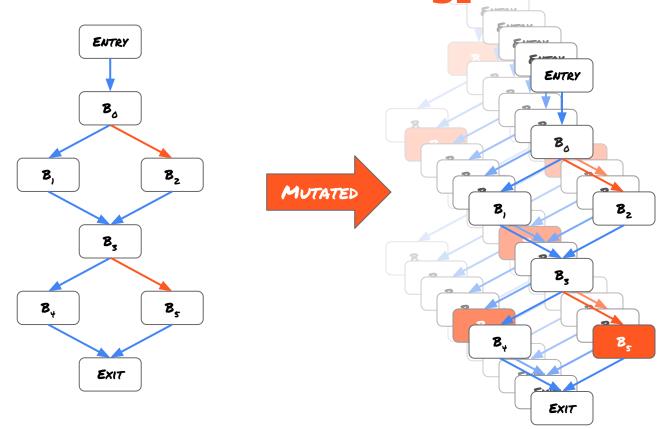


- 1. Compile to bytecode
- Construct CFG
- 3. Generate constraints
- 4. Equivalence Query
  - Assert inputs are equal
  - Assert outputs aren't equal
- Ask SMT if equivalence query is satisfiable

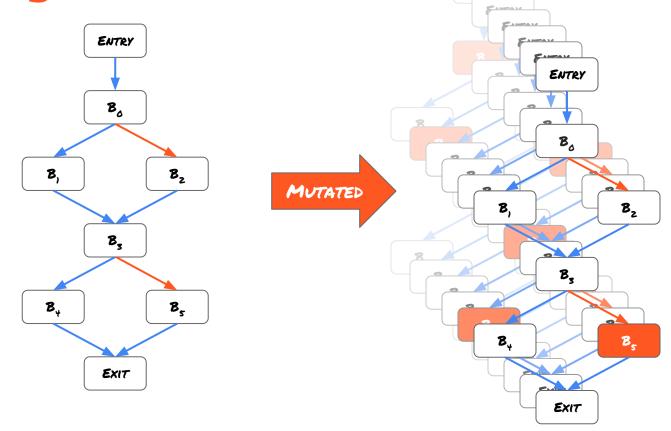


# Scaling Up

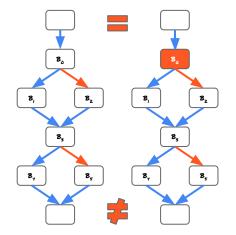
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Strategies: Reason about batches

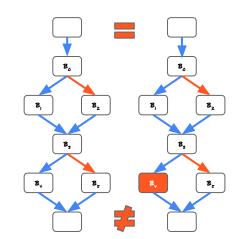


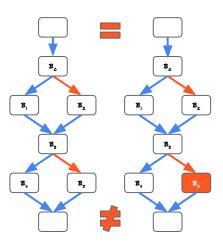
## **Naive Strategy**



#### **Naive Strategy**

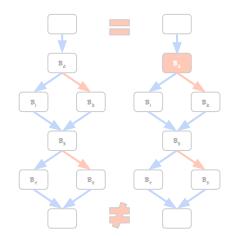
for each mutant m of program P:
 constrain P
 constrain m
 assert inputs are equal
 assert outputs are not equal
 check sat

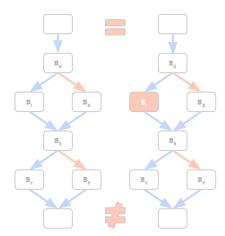




## **Naive Strategy**

**Observation:** naive recomputes *P*'s constraints every time



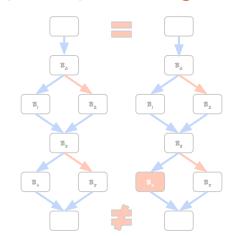


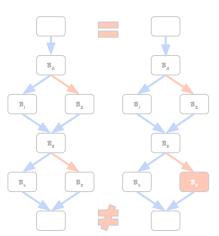
## Naive SSCOPES TO THE RESCUE!

for each mutant m of program P:

#### constrain P

constrain m
assert inputs are equal
assert outputs are not equal
check sat





## **SMT Scopes**

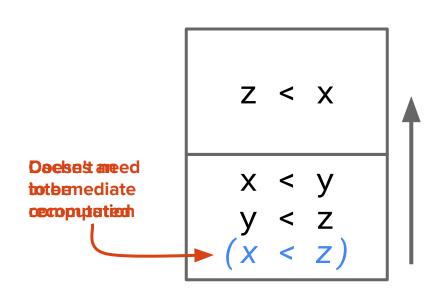
Scopes allow SMT to cache computations

Pushed and popped to a stack

```
(assert (< x y))
(assert (< y z))
(check-sat) ; SAT

(push) ; New scope
(assert (< z x))
(check-sat) ; UNSAT

(pop)
(check-sat) ; SAT</pre>
```

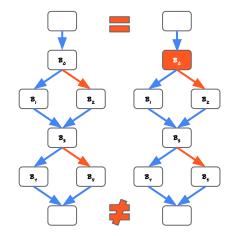


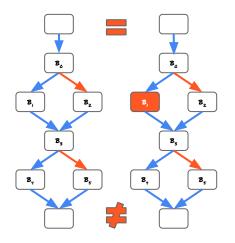
scope stack

## **Naive Strategy**

**Observation:** naive recomputes *P*'s constraints every iteration

**Solution:** use scopes!



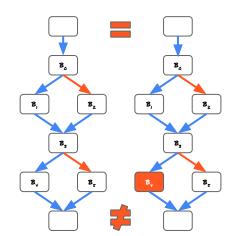


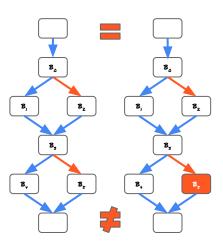
#### **Naive Strategy**

for each mutant m of program P:

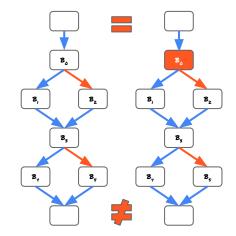
### constrain P

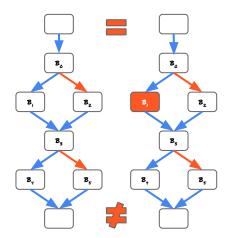
constrain m
assert inputs are equal
assert outputs are not equal
check sat





## **Cache Strategy**

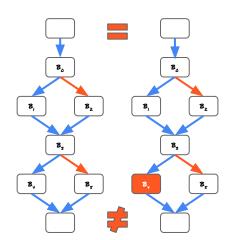


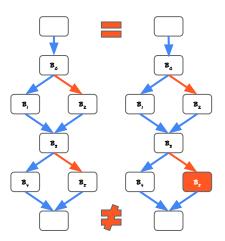


#### **Cache Strategy**

#### constrain P

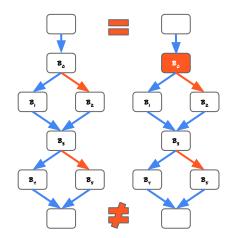
for each mutant m of program P:
 push scope
 constrain m
 assert inputs are equal
 assert outputs are not equal
 check sat
 pop scope

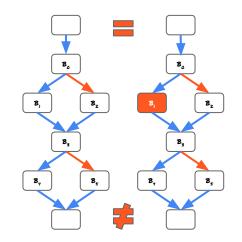




# **Cache Strategy**

**Observation:** mutant's basic blocks are reconstrained each iteration

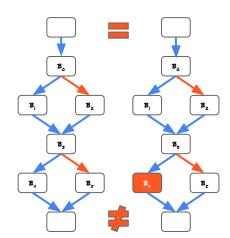


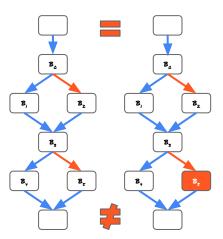


### **Cache Strategy**

### constrain P

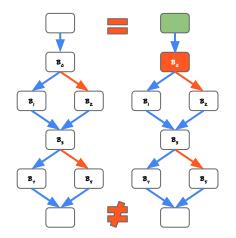
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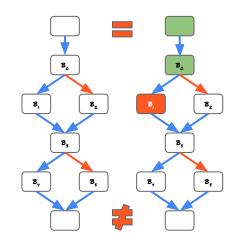




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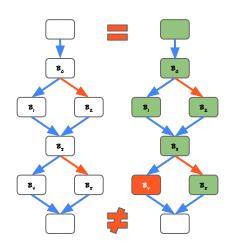


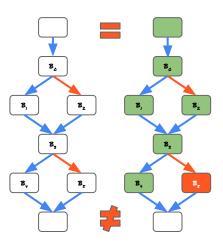


### **Cache Strategy**

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 assert inputs are equal
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 check sat
 pop scope





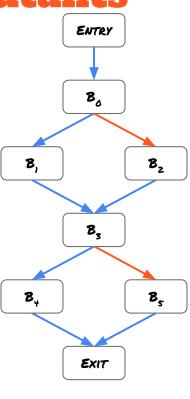
Atomic CFG Mutants
A Special Case

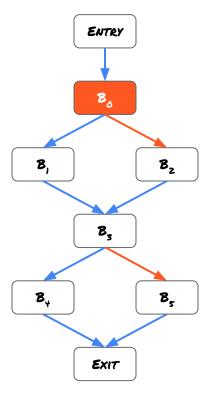
Atomic CFG Mutants are mutants that

- have the same CFG structure as the original program
- 2. mutate at most one basic block

Medusa uses this structural information to perform **constraint forking** 

252 of 488 (52%) of studied mutants were atomic CFG mutants

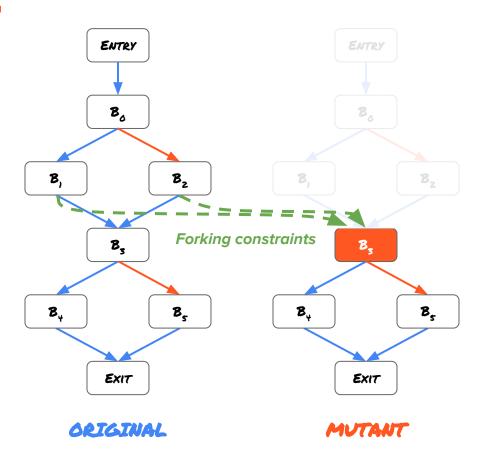






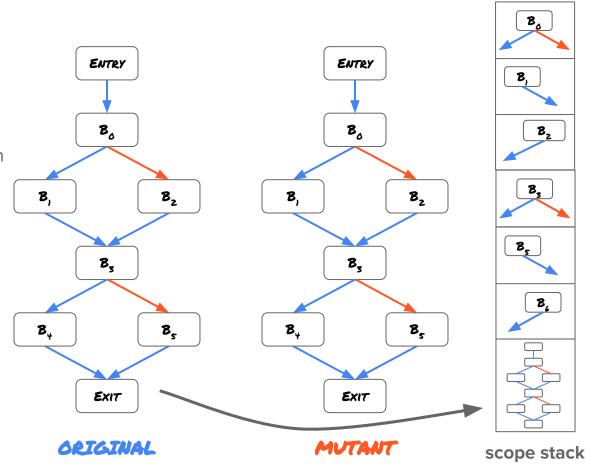


- Execution starts the same
- Use original program's constraints until mutated block is reached
- This
  - Reuses P's constraints for the beginning of execution
  - Makes formulas smaller

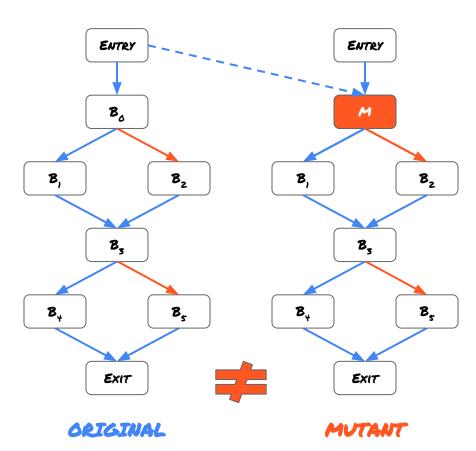


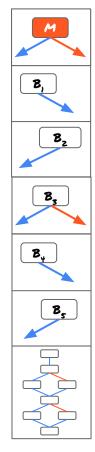
# Fork Strategy Setup

- 1. Order basic blocks  $B_0$ ,  $B_1$ , ...,  $B_n$ :
  each block comes before its children
- Make two copies of the original constraints
- 3. Push a scope, and assert all of the first copy of the constraints
- Bottom up, assert each block constraint and its outgoing edges into new scopes

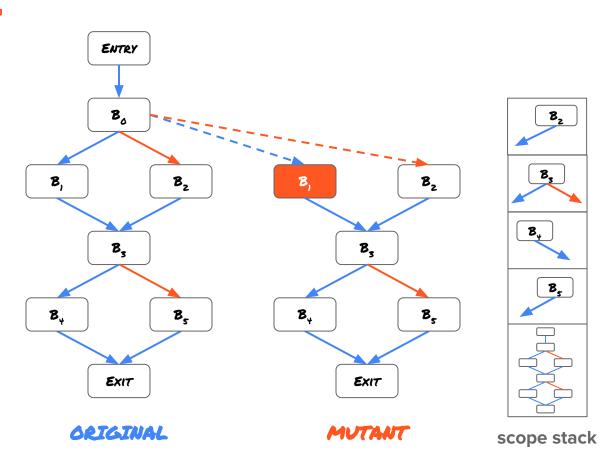


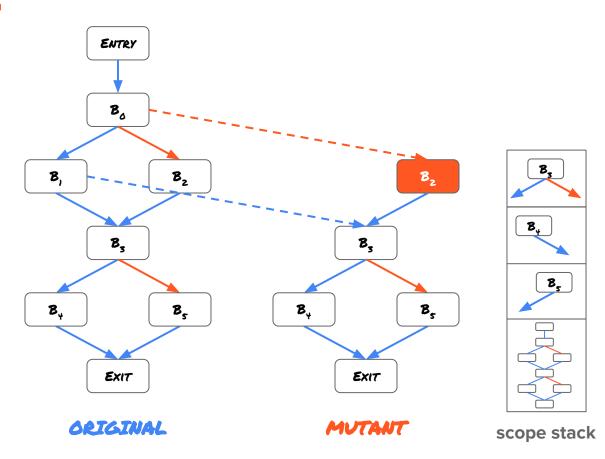
- Pop top scope,, clearing B<sub>0</sub>'s constraints.
- For each mutated version M of B<sub>0</sub>:
  - Push a new scope; assert
     M's block/edge constraints
  - Fork constraints
  - Check satisfiability
  - Pop the scope

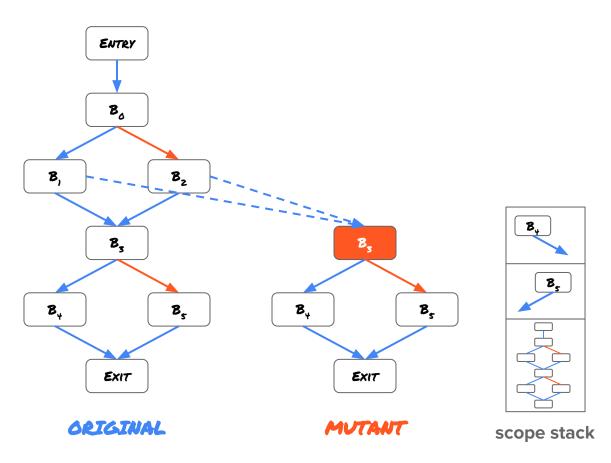


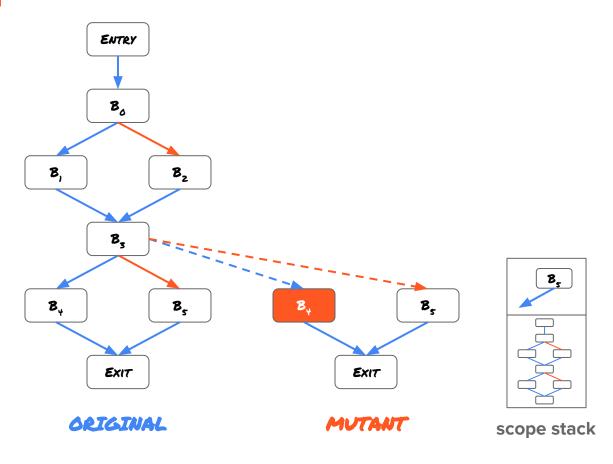


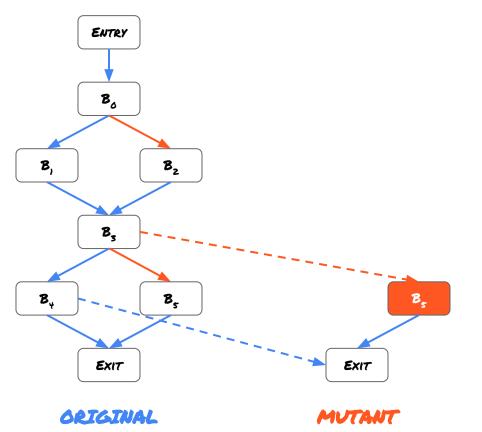
scope stack

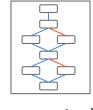






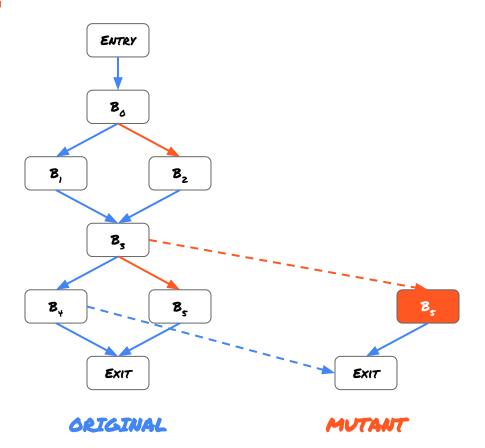


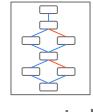




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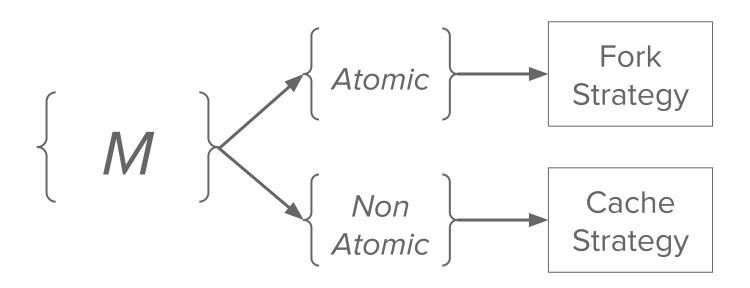
- Fork Strategy is only applicable to atomic CFG mutants
- The strategy can be generalized (future work)
- Medusa currently uses a hybrid strategy





scope stack

# **Hybrid Strategy**



# Preliminary Results

# **Three Test Subjects**

- 1. **Tax** computes single-payer tax amount for a given income.
- 2. **TicTacToe** checks the win condition for the game of Tic Tac Toe, including bounds checking on inputs.
- 3. **Triangle** classifies a triangle as equilateral, isosceles, scalene, or invalid

### **Mutants Generated by MAJOR**

### MAJOR generated 488 mutants

- Ground truth calculated by hand
- 37 equivalent (7.6%)
- Medusa correctly classified all mutants
- 252 atomic (52%)
  - Atomic mutants are common

Subjects		All	Atomic
Tax	total	99	84
	equiv	10	10
ТісТасТое	total	267	98
	equiv	23	19
Triangle	total	122	70
	equiv	4	3
All	total	488	252
	equiv	37	32

### **Mutants Generated by MAJOR**

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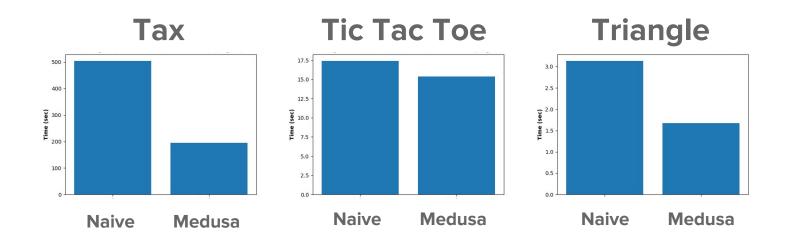
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**TicTacToe** has lower ratio of atomic mutants (large number of short circuiting boolean operators from bounds checking)

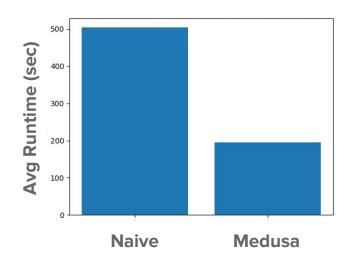
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### Summary of Medusa's performance

- **Medusa** was run 5 times on each subject and mutant (run times are averaged)
- Tax was long-running due to floating point computations.



### **Runtime results: Tax**



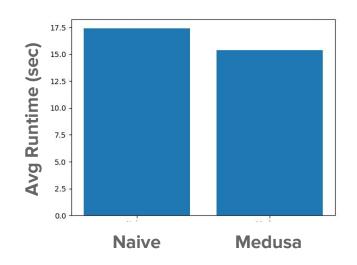
	#	Naive	Cache (Improve)	Fork (Improve)
Atomic	84	521	258 <b>(50.5%)</b>	197 <b>(62.2%)</b>
Non Atomic	15	410	185 <b>(54.9%)</b>	N/A
All	99	504	247 <b>(51.0%)</b>	195 <b>(61.3%)</b>

HYBRID-

### Results

- Atomic mutants tend to be longer running than non atomic mutants
- Cache has a 51% improvement over naive
- **Hybrid** has a **61% improvement** over naive

### **Runtime results: Tic Tac Toe**



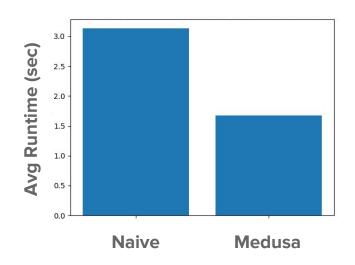
	#	Naive	Cache (Improve)	Fork (Improve)
Atomic	98	17.9	16.5 (7.8%)	13.9 <b>(22.5%)</b>
Non Atomic	169	17.0	16.3 ( <b>4.1</b> %)	N/A
All	267	17.4	16.4 ( <b>6.3%)</b>	15.4 (11.5%)

HYBRID-

### Results

- Atomic mutants tend to be longer running than non atomic mutants
- Cache has a 6% improvement over naive
- Hybrid has a 12% improvement over naive

# **Runtime results: Triangle**



	#	Naive	Cache (Improve)	Fork (Improve)
Atomic	70	3.72	3.26 (12.4%)	1.3 ( <b>65.1%)</b>
Non Atomic	52	2.34	2.16 <b>(7.7%)</b>	N/A
All	122	3.13	2.79 <b>(10.9%)</b>	1.67 (46.6%)

HYBRID-

### Results

- Atomic mutants tend to be longer running than non atomic mutants
- Cache has a 11% improvement over naive
- Hybrid has a 47% improvement over naive

# Medusa: Mutant Equivalence Detection Using SAT Analysis

