

Name: _____

Date: _____

CSC 17C
Midterm

Problem 1) Analyze and compare the linear and binary search

Show $O()$ by mathematical analysis

Show $O()$ by timing comparison and graph results

Show $O()$ by operational analysis

Problem 2) Compare and contrast bubble sort with selection sort

Show $O()$ by mathematical analysis

Show $O()$ by timing comparison and graph results

Show $O()$ by operational analysis

Problem 3) Compare insertions i.e. push method with Simple Vector using arrays, with Optimized Simple Vector using arrays, and Simple Vector with Linked list.

Show $O()$ by mathematical analysis

Show $O()$ by timing comparison and graph results

Show $O()$ by operational analysis

Problem 4) Given $X = B^Y = H^Z$

$$\text{Prove } \log_B(x) = \frac{\log_H(x)}{\log_H(B)}$$

Problem 5) Prove the summation formula

$$\sum_{i=1}^N i = 1 + 2 + 3 + \dots + (N-1) + N = \frac{N \cdot (N+1)}{2}$$

Problem 6) Derive the order of the error with respect to the exponential approximation.

$$e^{-1/N} \approx (1 - 1/N) \text{ where}$$

$$e^x = \sum_{i=0}^{\infty} \frac{x^i}{i!}$$

Problem 7) Derive the $O()$ for the Recursive vs. non-Recursive Fibonacci function. Simple logic is enough using the fact that the Recursive Fibonacci takes Fibonacci time.

Name: _____ Date: _____

Problem 8) Given 4 dice calculate the probability of 1 pair, 2 pair, 3 of a kind and 4 of a kind. Simulate the results and compare to calculation.

Problem 9) Given a biased coin analogy, if a bit vector is 30% full, what are the odds that 5 bits randomly chosen all fall within the filled section. Simulate the results and compare to calculation.

Problem 10) Recursive function. Provide code and test sufficiently.

Derive a power function that is $O(\log_2 N)$
where $Z = X^N$, $X \in \mathbb{R}$, $N \in \mathbb{N} \rightarrow Z^+$

Hint: Split N for odd or even conditions

Problem 11) Recursive function. Provide code and test sufficiently.

$$\text{let } g(2x) = \frac{2g(x)}{1+g^2(x)}, \quad -1 \leq x \leq 1$$

Base Condition

$$\text{for } |x| < \epsilon, \quad \epsilon \approx 10^{-6} \text{ where } g(x) = x - x^3/6$$

Problem 12) Mutual Recursion. Provide code and test sufficiently. $\approx x = \pi/4$

$$\text{let } C(2x) = \frac{1}{2} C(x) S(x) \text{ and}$$

$$S(2x) = \frac{C^2(x) S^2(x)}{C^2(x) - S^2(x)}$$

Base Condition

$$\text{for } |x| < \epsilon, \quad \epsilon \approx 10^{-6} \quad \begin{aligned} S(x) &= 1 + x^2/2 \\ C(x) &= 1/x + x/6 \end{aligned}$$

Problem 13) Code the mode problem to utilize the Set and Map containers in the STL to solve for the number of modes in an array. Test and show correct solutions for no modes, 1 mode and multiple modes. Code with the Set to reduce the line count and number of for-loops necessary. Then utilize the Map to reduce the line count further.

Problem 1)

1. Binary Search

① W_n : the number of iterations of while loop in a worst-case execution

② If $2^{\bar{i}} \leq n < 2^{\bar{i}+1}$, then $W_n = \bar{i} + 1$

$$\rightarrow \bar{i} = \log_2 n, \quad W_n = \log_2 n + 1$$

③ $\boxed{O(\log_2 n)}$

2. Linear Search

$$\sum_{i=0}^{n-1} \underbrace{PO_i}_{\text{If statement}} = n \cdot PO_0 \Rightarrow \boxed{O(n)}$$

Time Comparison

Type 1 Linear Search
Type 2 Binary Search

(Unit: micro second)

n	Type 1	Type 2
100	4	3
200	3	2
400	4	2
800	3	2
1600	21	3
3200	5	5
6400	3	2
12800	3	3
25600	3	2
51200	4	2

Operation Comparison

Type 1

Linear Search

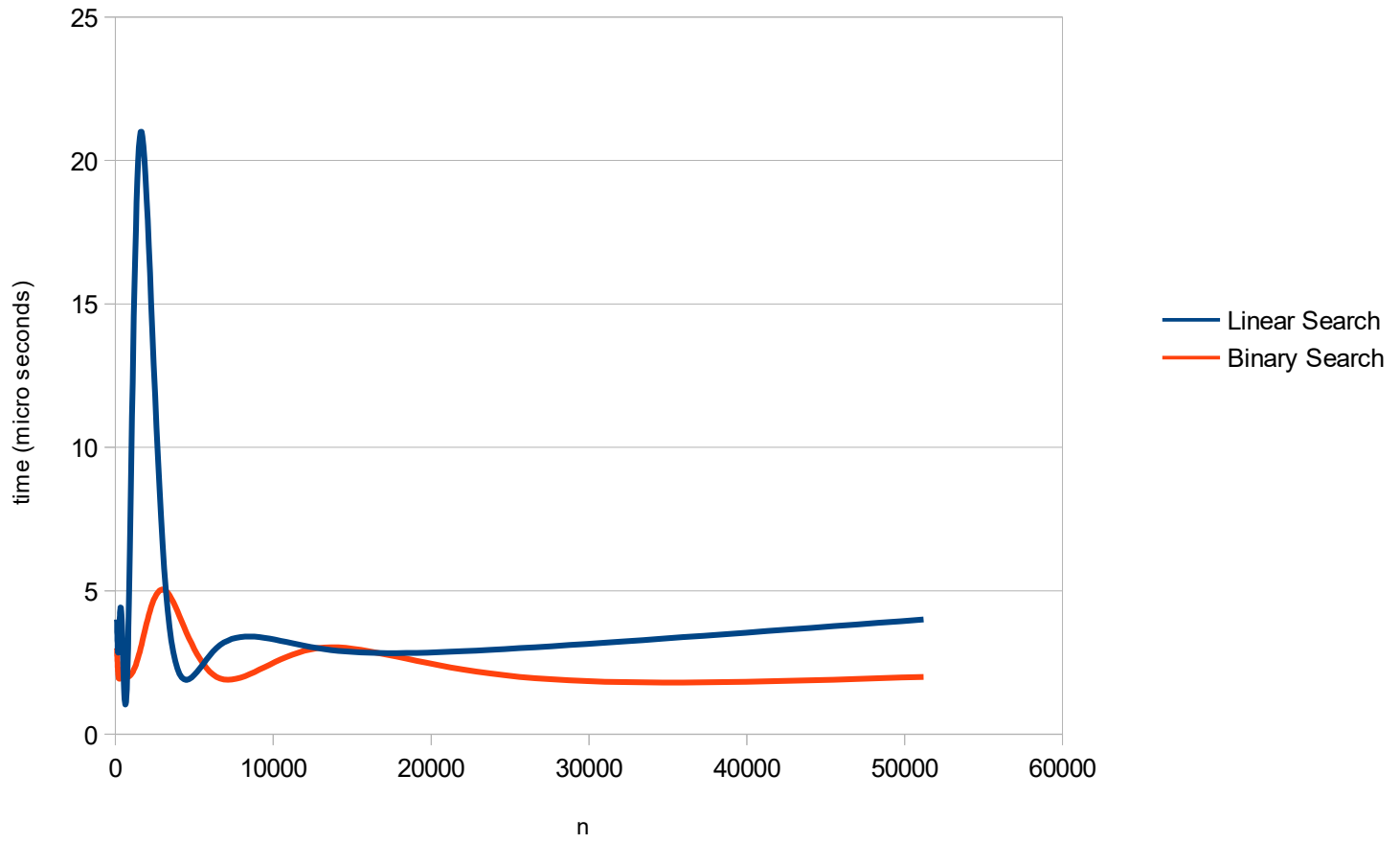
Type 2

Binary Search

n	Type 1	Type 2
100	77	38
200	116	8
400	26	68
800	113	58
1600	1,004	68
3200	176	58
6400	263	58
12800	263	68
25600	203	48
51200	539	68

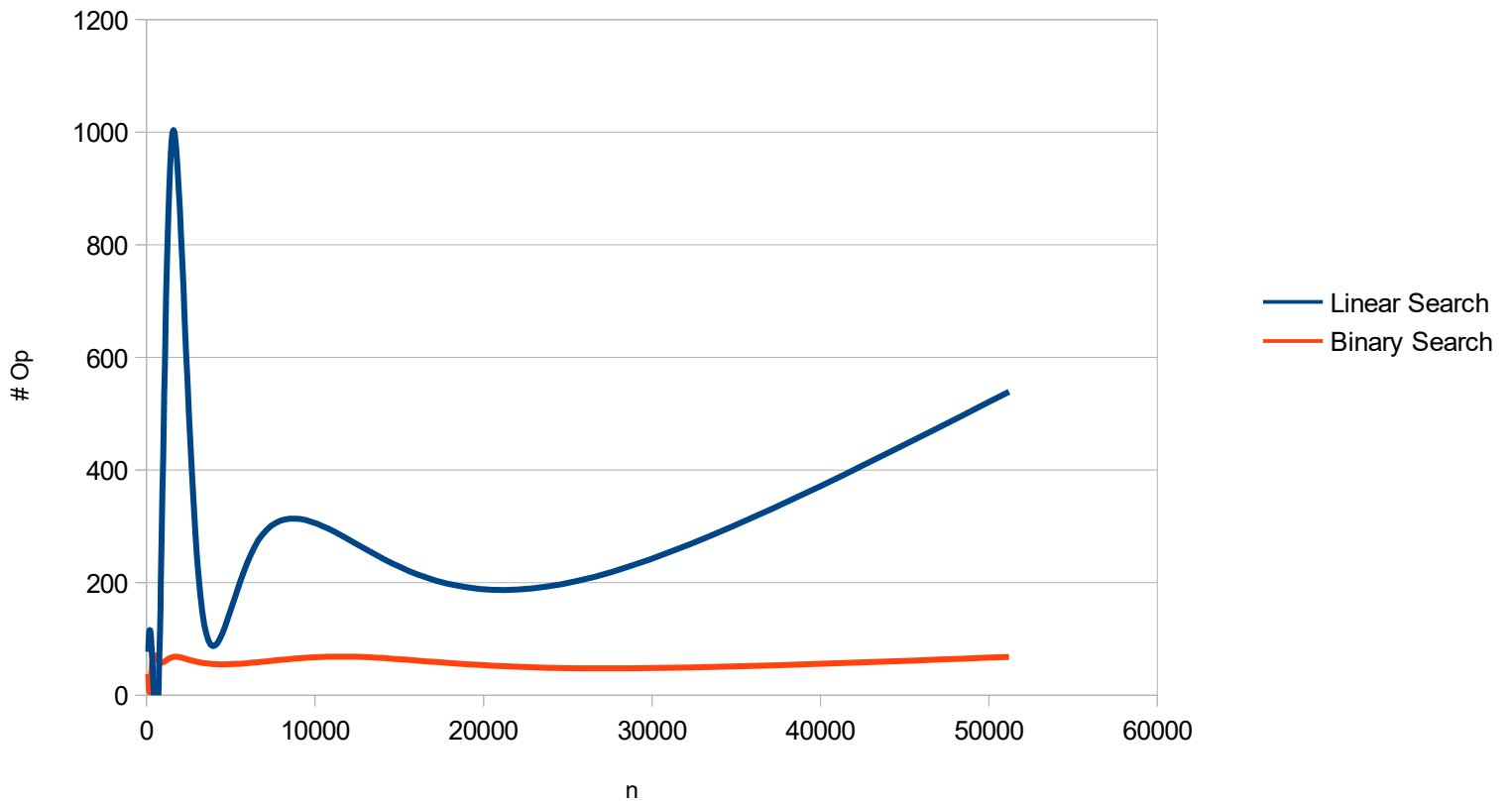
Time Comparison

Linear Search vs. Binary Search



Operation Comparison

Linear Search vs. Binary Search



Problem2)

1. Selection Sort

$$O_b + \sum_{i=0}^{n-1} (O_i + PO_s) + O_a$$

Before the loop Inner loop If statement after loop the

$$= O_b + \frac{n(n-1)}{2} O' + O_a$$

$$= \frac{n^2}{2} O' - O' \frac{n}{2} + O_a + O_b$$

$$\Rightarrow \boxed{O(n^2)}$$

2. Bubble Sort

① W_n : the number of iterations of dowhile loop in a worst-case execution

② for loop. $\sum_{i=0}^{n-1} PO_s = (n) PO_s.$

③ $\Rightarrow \underline{(n-2)n \cdot PO_s.}$

∴ $O(n^2)$

Time Comparison

Type 1 Bubble Sort
Type 2 Selection Sort

(Unit: micro second)

n	Type 1	Type 2
100	46	19
200	180	86
400	706	290
800	2,244	1,218
1600	10,391	4,814
3200	29,876	12,512
6400	169,078	47,740
12800	621,179	229,346
25600	2,697,879	827,399
51200	10,730,047	3,255,863

Operation Comparison

Type 1

Bubble Sort

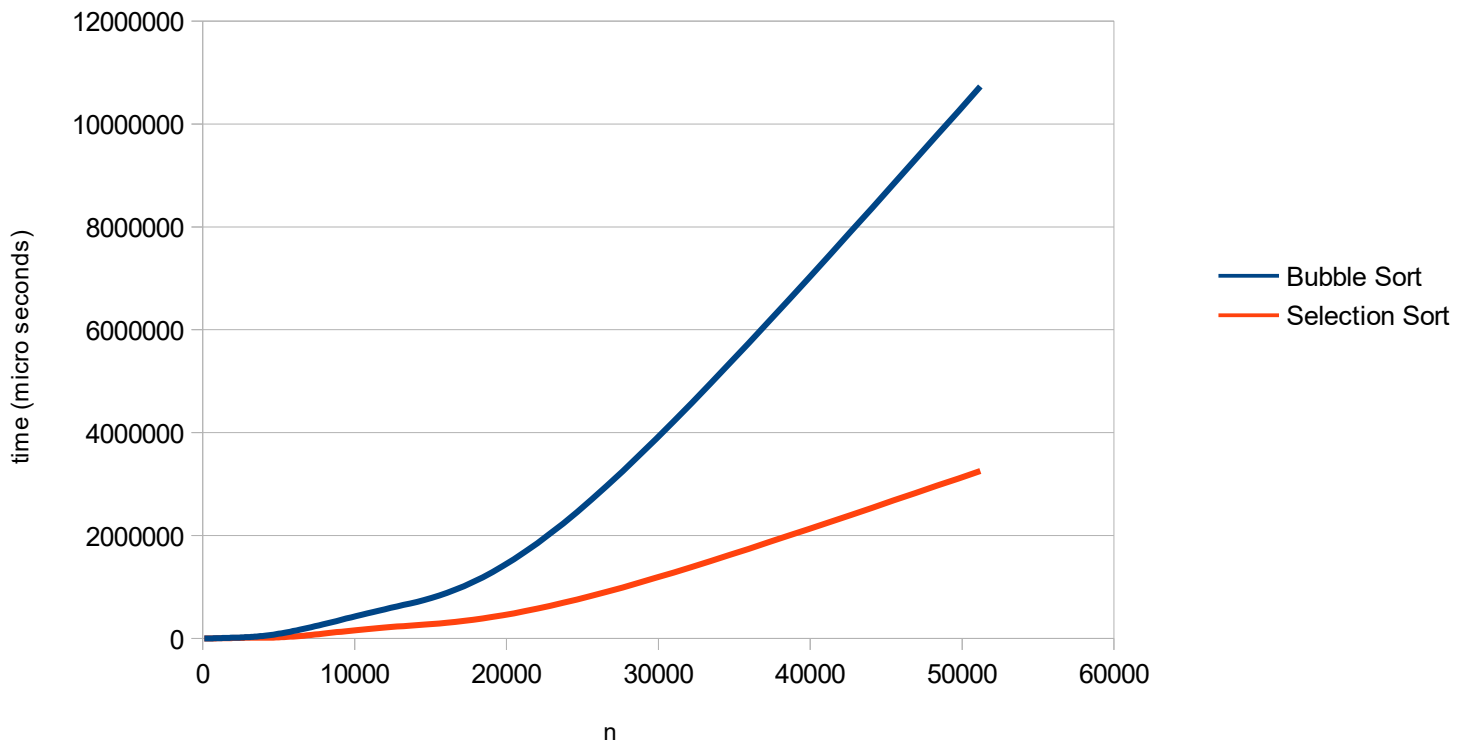
Type 2

Selection Sort

n	Type 1	Type 2
100	80,248	16,933
200	366,304	64,030
400	1,495,626	248,785
800	5,874,208	977,713
1600	23,489,446	3,876,397
3200	94,893,812	15,433,258
6400	383,158,458	61,585,522
12800	1,541,414,186	246,057,775
25600	6,167,223,228	983,628,193
51200	24,598,212,394	3,933,343,522

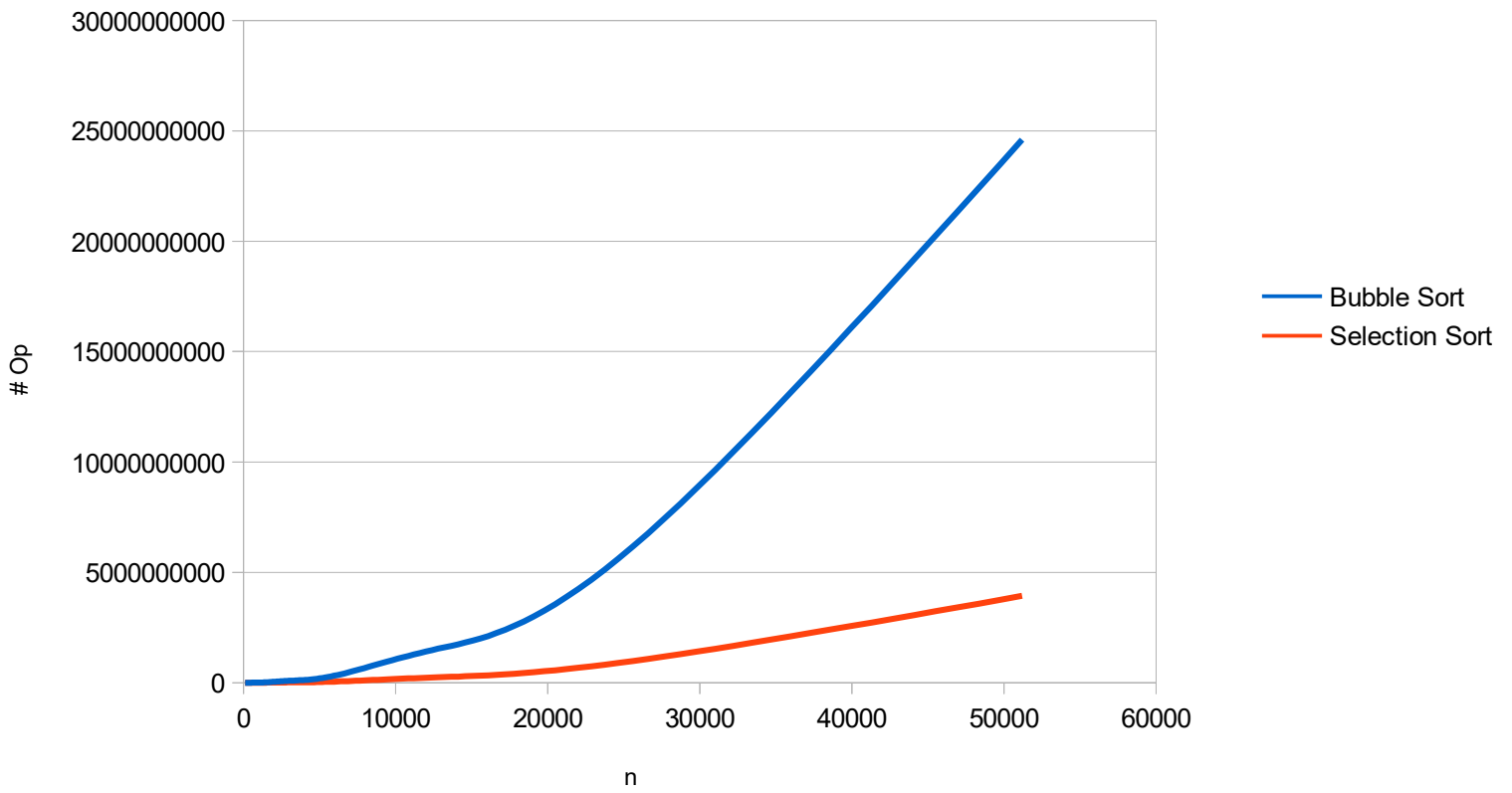
Time Comparison

Bubble Sort vs. Selection Sort



Operation Comparison

Bubble Sort vs. Selection Sort



Problem3)

1. Simple Vector.

$$\underbrace{O_b}_{\text{before the loop}} + \underbrace{\sum_{i=0}^{n-1} O_i}_{\text{for loop}} + \underbrace{O_a}_{\text{after the loop}}$$

$$= \frac{n(n-1)}{2} + O_a + O_b.$$

$$\Rightarrow O(n^2)$$

2. Simple Vector Push Efficient

$$\textcircled{1} \text{ if } (\text{maxSize} == \text{arraySize})$$
$$\sum_{i=0}^{\text{arraySize}} O_i$$

$$\textcircled{2} \text{ else } O(1)$$

In the beginning, since the array size is relatively small it doesn't affect to the efficiency and as the array size becomes larger, the case to create additional memory space, which uses for-loop, rarely happens. In this reason, its graph is $y=c$ where c is constant in the simulation.

3. Simple Vector Using linked list

① addList

$$W_n = \underbrace{n}_{\text{w}} \quad \text{(w = worst)}$$

(worst-case of while loop)

$$\Rightarrow \boxed{O(n)}$$

Time

Time Comparison

Type 1 Simple Vector using arrays
 Type 2 Optimized Simple Vector using arrays
 Type 3 Simple Vector with Linked list

(Unit: micro second)

n in for-loop in main function	n	Type 1	Type 2	Type 3
100	464	14,690		
	446		434	
	439			145
200	915	53,588		
	898		743	
	906			310
300	1347	114,943		
	1346		1,347	
	1442			812
400	1754	190,953		
	1845		1,400	
	1760			912
500	2277	280,202		
	2235		2,285	
	2269			988
600	2856	424,997		
	2628		3,179	
	2752			1,967
700	3059	538,281		
	3109		3,521	
	3078			2,101
800	3521	793,849		
	3695		2,862	
	3528			3,037
900	4083	903,452		
	4006		3,865	
	3972			2,996
1000	4393	1,120,424		
	4594		3,183	
	4398			4,640
1100	4993	1,365,880		
	4848		5,740	
	4930			4,092
1200	5387	1,846,003		
	5499		5,426	
	5520			6,543
	5731	1,918,461		
	5981		6,165	

Time

Type 1

Simple Vector using arrays

Type 2

Optimized Simple Vector using arrays

Type 3

Simple Vector with Linked list

(Unit: micro second)

n in for-loop in main function	n	Type 1	Type 2	Type 3
1300	5829			6,525
1400	6238	2,275,923		
	6363		6,551	
	6264			7,113
1500	6812	2,583,400		
	6867		6,837	
	6552			9,580
1600	7109	2,904,343		
	7232		7,582	
	7281			8,978
1700	7513	3,286,578		
	7940		8,407	
	7515			11,209
1800	8072	3,609,546		
	8306		7,332	
	7944			12,441
1900	8460	3,937,719		
	8491		6,668	
	8691			11,277
2000	9167	4,405,148		
	9091		6,278	
	8949			16,604
5000	22335	27,253,067		
10000	45105		52,451	
	44556			534,780
100000	225852		188,903	
	224950			12,650,044

Op

Operation Comparison

Type 1 Simple Vector using arrays
 Type 2 Optimized Simple Vector using arrays
 Type 3 Simple Vector with Linked list

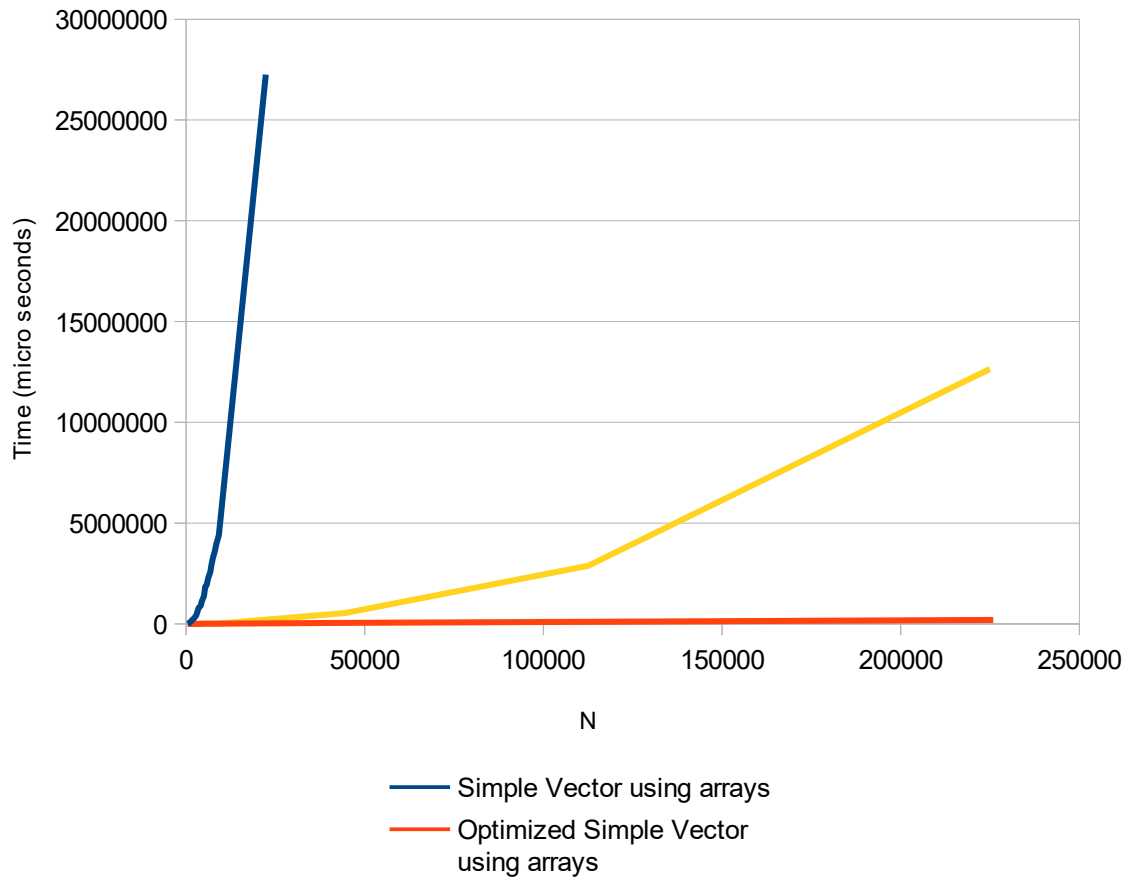
(Unit: micro second)

n in for-loop in main function	n	Type 1	Type 2	Type 3
100	405	26,664		
	421		994	
	467			15,650
200	919	103,314		
	885		1,942	
	883			61,300
300	1342	229,964		
	1364		3,530	
	1391			136,950
400	1878	406,614		
	1855		3,830	
	1920			242,600
500	2129	633,264		
	2237		4,130	
	2286			378,250
600	2691	909,914		
	2740		6,998	
	2665			543,900
700	3295	1,236,564		
	3226		7,298	
	3149			739,550
800	3667	1,613,214		
	3605		7,598	
	3635			965,200
900	3949	2,039,864		
	4086		7,898	
	3925			1,220,850
1000	4523	2,516,514		
	4515		8,198	
	4371			1,506,500

plot_Time

Timing Comparison

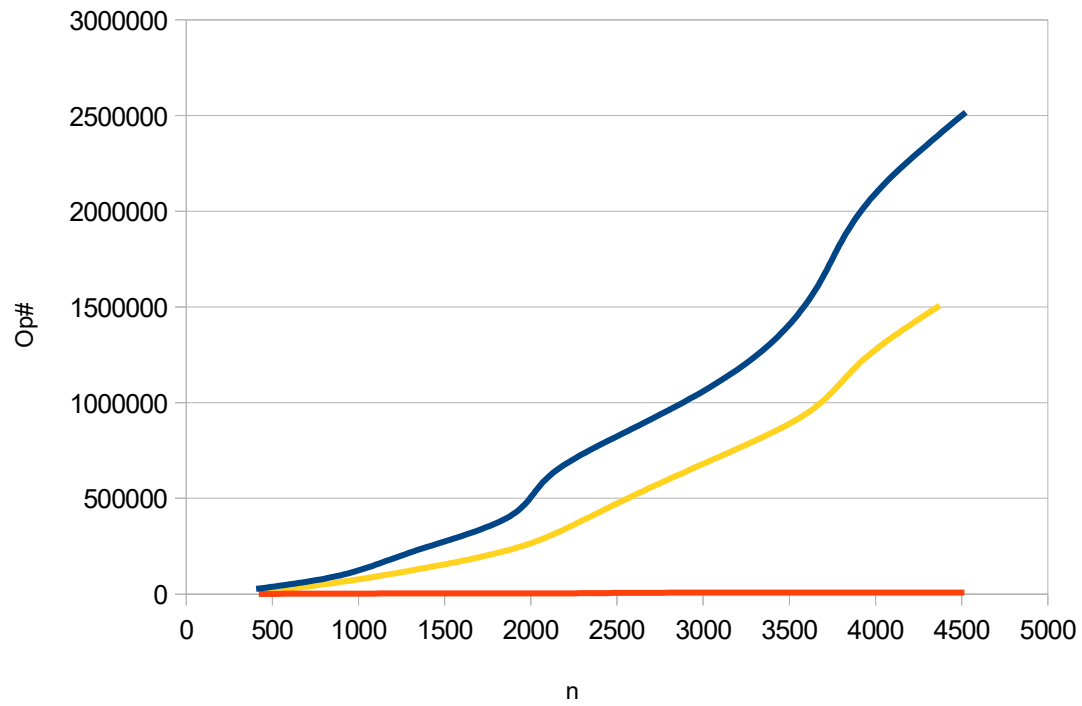
push method



plot_Op

Operation Comparison

push method



- Simple Vector using arrays
- Optimized Simple Vector using arrays
- Simple Vector with Linked list

Problem 4)

① Left Hand Side or LHS of Equation

$$\textcircled{2} \log_B X = \frac{\log X}{\log B}$$

② Right Hand Side or RHS of Equation

$$\begin{aligned} \frac{\log_H X}{\log_H B} &= \frac{\frac{\log X}{\log H}}{\frac{\log B}{\log H}} \\ &= \frac{\log X}{\log B} \end{aligned}$$

③ LHS $\stackrel{?}{=}$ RHS

$$\frac{\log X}{\log B} = \frac{\log X}{\log B} \quad \checkmark$$

$$\text{Therefore, } \log_B X = \frac{\log_H X}{\log_H B}$$

Problem5)

$$\sum_{i=1}^n i = 1 + 2 + 3 + \dots + n-1 + n$$

$$+ \left| \sum_{i=1}^n i = n + n-1 + n-2 + \dots + 2 + 1 \text{ (Reverse order)} \right|$$

$$2 \sum_{i=1}^n i = \underbrace{(1+n) + (1+n) + (1+n) + \dots + (1+n) + (1+n)}_n$$

$$= n(n+1)$$

$$\text{Therefore, } \sum_{i=1}^n i = \frac{n(n+1)}{2}$$

Problem6)

$$e^x = \sum_{i=0}^{\infty} \frac{x^i}{i!}$$

$$= \frac{x^0}{0!} + \frac{x^1}{1!} + \frac{x^2}{2!} + \dots$$

$$= 1 + x + \frac{x^2}{2} + \frac{x^3}{6} + \dots$$

$$\text{So, } e^{-\frac{1}{n}} = 1 - \frac{1}{n} + \frac{\left(-\frac{1}{n}\right)^2}{2} + \frac{\left(-\frac{1}{n}\right)^3}{6} + \dots$$

$$= 1 - \frac{1}{n} + \frac{1}{2n^2} - \frac{1}{6n^3} + \dots$$

$$= 1 - \frac{1}{n} \underbrace{\left(1 - \frac{1}{2n} \left(1 - \frac{1}{3n} \left(1 - \frac{1}{4n} \left(1 - \frac{1}{5n} \dots\right)\right)\right)\right)}_{\approx 1}$$

$$\approx 1 - \frac{1}{n}$$

Problem7)

In the simulation,

$$\begin{aligned} &\text{The number of calling fiboRec}(0) + \text{The number of calling fiboRec}(1) \\ &= \text{fiboRec}(n+1) \end{aligned}$$

So, it takes Fibonacci time when we execute the recursive Fibonacci.

Fibonacci Recursion

n	fiboRec(n)	#0 Call func(0) if(n==0)	#1 Call func(1) Else if(n==1)	#0 + #1 Total # of Call
0	0	1	0	1
1	1	0	1	1
2	1	1	1	2
3	2	1	2	3
4	3	2	3	5
5	5	3	5	8
6	8	5	8	13
7	13	8	13	21
8	21	13	21	34
9	34	21	34	55
10	55	34	55	89
11	89	55	89	144
12	144	89	144	233
13	233	144	233	377
14	377	233	377	610
15	610	377	610	987
16	987	610	987	1597
17	1597	987	1597	2584
18	2584	1597	2584	4181
19	4181	2584	4181	6765
20	6765	4181	6765	10946
21	10946	6765	10946	17711
22	17711	10946	17711	28657
23	28657	17711	28657	46368
24	46368	28657	46368	75025
25	75025	46368	75025	121393
26	121393	75025	121393	196418
27	196418	121393	196418	317811
28	317811	196418	317811	514229
29	514229	317811	514229	832040
30	832040	514229	832040	1346269
31	1.34627e+06	832040	1346269	2178309
32	2.17831e+06	1346269	2178309	3524578
33	3.52458e+06	2178309	3524578	5702887
34	5.70289e+06	3524578	5702887	9227465
35	9.22746e+06	5702887	9227465	14930352
36	1.49304e+07	9227465	14930352	24157817
37	2.41578e+07	14930352	24157817	39088169
38	3.90882e+07	24157817	39088169	63245986
39	6.3246e+07	39088169	63245986	102334155
40	1.02334e+08	63245986	102334155	165580141
41	1.6558e+08	102334155	165580141	267914296

Sheet1

42	2.67914e+08	165580141	267914296	433494437
43	4.33494e+08	267914296	433494437	701408733
44	7.01409e+08	433494437	701408733	1134903170
45	1.1349e+09	701408733	1134903170	1836311903
46	1.83631e+09	1134903170	1836311903	2971215073

Problem 8)

1. Calculation:

$$1. m(\Omega) = nP_m^r \quad \text{where } n=6, m=4$$
$$= n^m = 6^4 = 1296$$

2. a pair.

$$\textcircled{1} \# \text{ of a pair} = {}_6P_1 = \frac{6!}{5!1!} = 6$$

$$\textcircled{2} \# \text{ of a pair combination} = {}_4C_2 = 6$$

$$\textcircled{3} \# \text{ of combination of other dice}$$

$$= {}_5P_2 = \frac{5!}{3!} = 20$$

$$\Rightarrow \textcircled{1} \times \textcircled{2} \times \textcircled{3} = 6P_1 \cdot {}_4C_2 \cdot {}_5P_2 = 120.$$

3. two pairs.

$$\textcircled{1} \# \text{ of two pairs} = {}_6C_2 = 15$$

$$\textcircled{2} \# \text{ of first pair combination} = {}_4C_2 = 6.$$

$$\textcircled{3} \# \text{ of second pair combination} = {}_2C_2 = 1.$$

$$\Rightarrow \textcircled{1} \times \textcircled{2} \times \textcircled{3} = {}_6C_2 \cdot {}_4C_2 \cdot {}_2C_2 = 90$$

4. three of a kind.

① # of three of a kind = $6P_1 = 6$

② # of three of a kind combination = $4C_3 = 4$

③ # of other dice = $5P_1 = 5$

$\Rightarrow ① \times ② \times ③ = 6P_1 \cdot 4C_3 \cdot 5P_1 = 120$

5. four of a kind

① # of four of a kind = $6P_1 = 6$

② # of four of a kind combination = $4C_4 = 1$

③ No other dice

$\Rightarrow ① \times ② = 6P_1 = 6$

2. Simulation:

Pair	number of event	%	Theoretical
a Pair	585	58.5	55.5556
two Pairs	67	6.7	6.94444
three of a kind	81	8.1	9.25926
four of a kind	7	0.7	0.462963

RUN SUCCESSFUL (total time: 1s)

3. Sumarry

	Theoretical		Simulation		Diff
	#	Probability	#	Probability	
a pair	720	56%	585	59%	-3%
two pair	90	7%	67	7%	0%
three of a kind	120	9%	81	8%	1%
four of a kind	6	0%	7	1%	0%

Problem9)

- a) (1) p = Probability getting Head = $3/10$, q = Probability getting tail = $7/10$
(2) $(p+q)^n = (\text{sigma sign}) \text{ from } i=0 \text{ to } i=n \binom{n}{i} p^i q^{n-i} = 1$
 $\binom{n}{i}$ is n choose i combination
(3) So, if we substitute $i=5$, $n=5$, $p=0.3$, $q=0.7$
 then $(0.3)^5 * (0.7)^0 = 0.00243$, that is, 0.243%
b) Simulation result: 0.243%

RUN SUCCESSFUL (total time: 700ms)

Problem10, 11, 12)

1. Calculation (Problem10):

$$Z = x^n.$$

① R_n : the number of ~~iterations~~ recursive call

② If $2^i \leq n < 2^{i+1}$, then $R_n = i + 1$

$$\rightarrow i = \log_2 n, R_n = \log_2 n + 1$$

③ $O(\log_2 n)$

2. Simulation Result(Problme10, 11, 12)

Problem10: (1) $x=30.5$, $n=5$ (2) $x^n = 2.63936e+07$

Problem11: (1) $x=0.5$ (2) $g(x) = 0.244919$

Problem12: (1) $x=PI/4$ (2) $C(x) = 1.41421$, $S(x) = 1.41421$

RUN SUCCESSFUL (total time: 5s)

Problem13)

When arySize=50 and modNum=12, simulation result:

```
0 1 2 3 4 5 6 7 8 9
10 11 0 1 2 3 4 5 6 7
8 9 10 11 0 1 2 3 4 5
6 7 8 9 10 11 0 1 2 3
4 5 6 7 8 9 10 11 0 1
```

Mode Freq = 5

Number of modes = 2

The number of modes = 2

The max Frequency = 5

The mode set = {0,1}

RUN SUCCESSFUL (total time: 103ms)