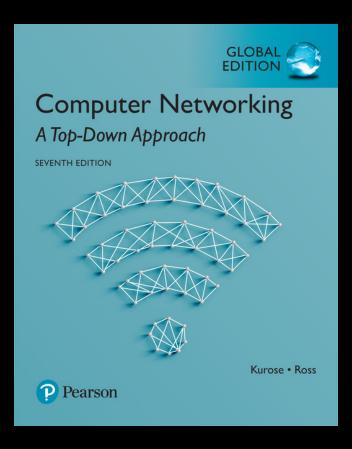
Computer / Network Introduction

School of Electric and Computer Engineering

Pusan National University, KOREA

Younghwan Yoo





Computer Networking

A Top-Down Approach

7th edition

Jim Kurose, Keith Ross

Pearson

April 2016

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01. What is the Internet?

What is the Internet?

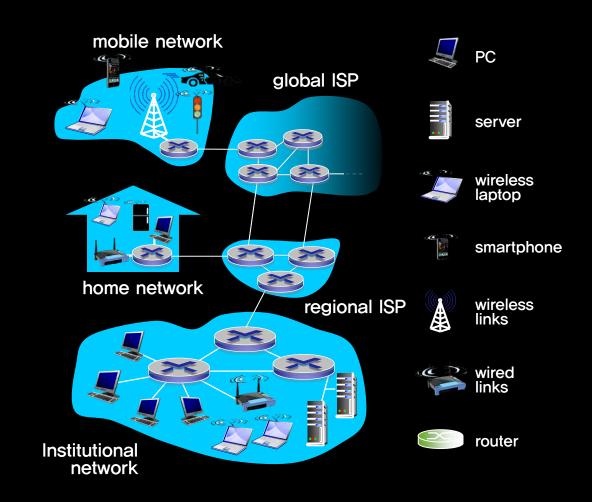


What is the Internet?



- Internet = Inter- + net (work)
 - "network of networks"

Various types of networks



What is the Internet?

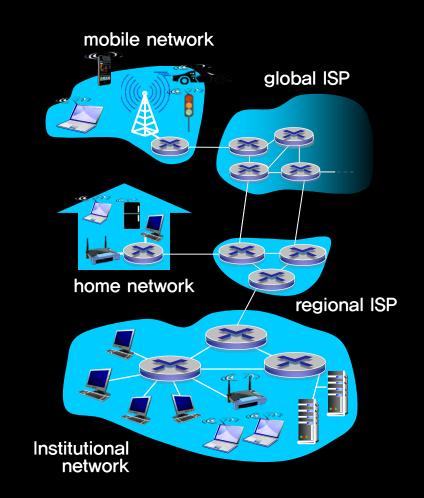


HW components

- end hosts
- links
 - copper, fiber, radio, satellite
- interconnection devices
 - router, switch, repeater

SW components

- operating software
- application programs
- protocols





Protocols

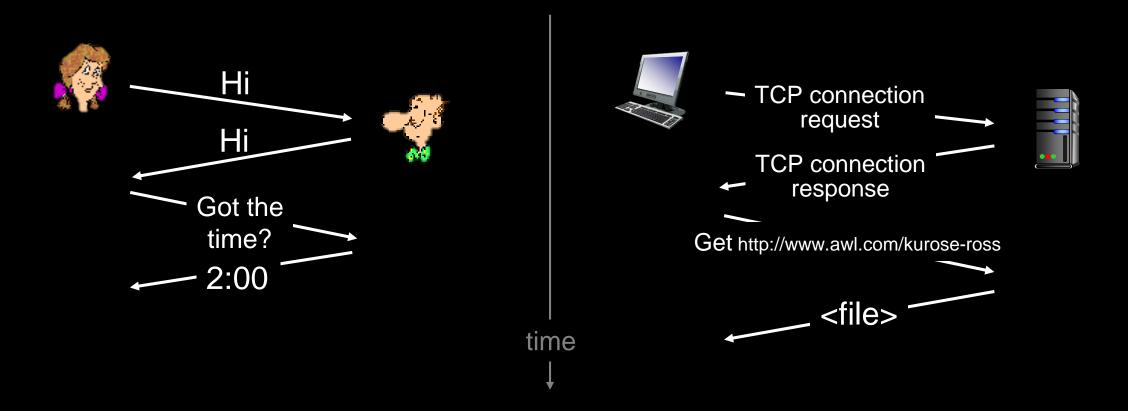
 a defined set of rules and regulations that determine how data is transmitted in telecommunications and computer networking (from Wikipedia)



출처 – http://rtrfm.com.au/story/talk-the-talk-why-its-hard-to-learn-a-language/



Human protocol vs. Comm. protocol





- All communication activity in Internet governed by protocols
- Protocols define
 - message format
 - order of messages sent and received among network entities
 - actions taken on message transmission, receipt

02. Network Edge

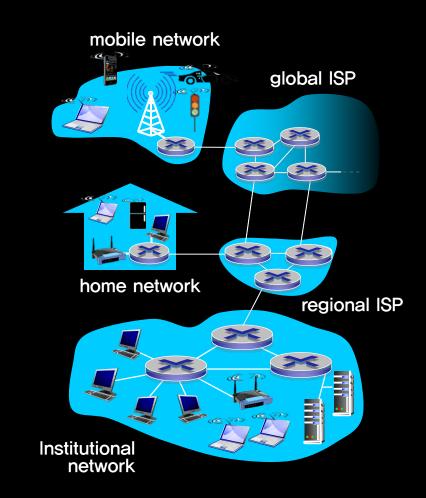


Network edge

- hosts: clients and servers
- servers often in data center

Network core

Interconnected routers or switches

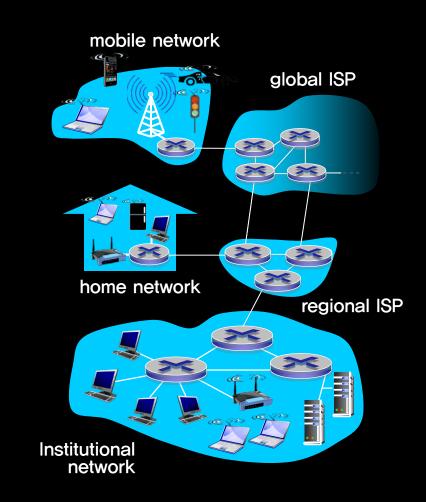




Network connecting end systems to edge router?

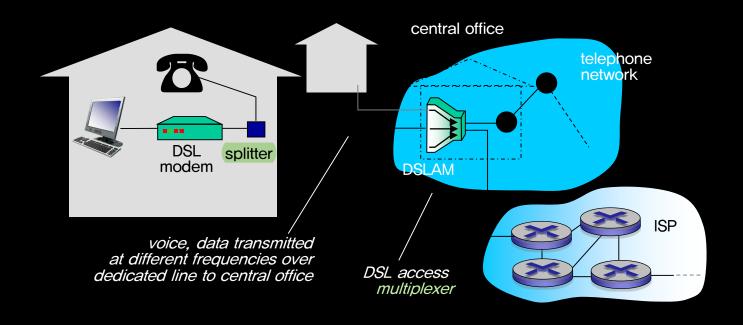
Access network

- residential access nets
- institutional acc. nets
- wireless access nets

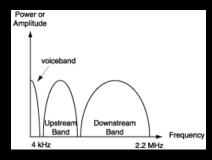




Digital subscriber line (DSL)





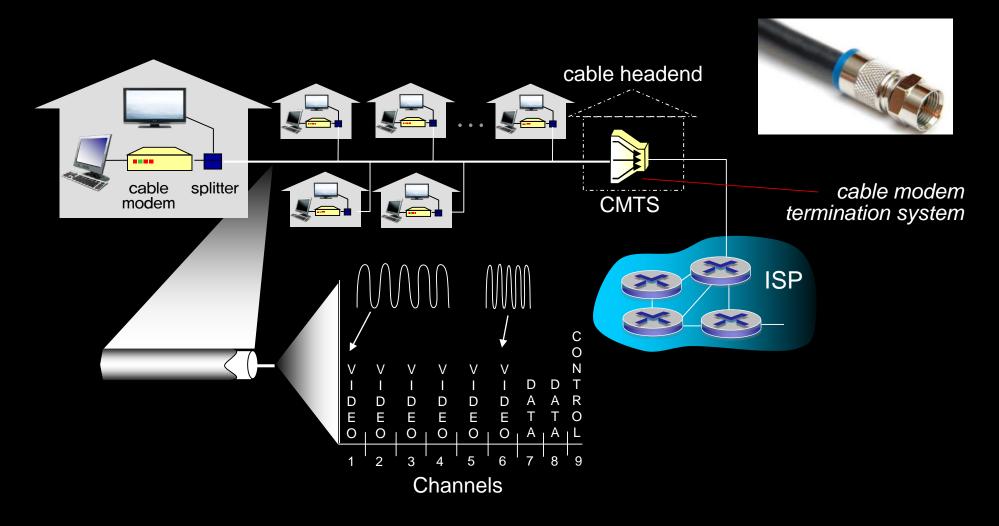


출처 - http://sunsite.uakom.sk/sunworldonline/swol-02-1998/swol-02-connectivity.html/

- voice over a phone line to telephone net
- data over a phone line to the Internet



Cable network

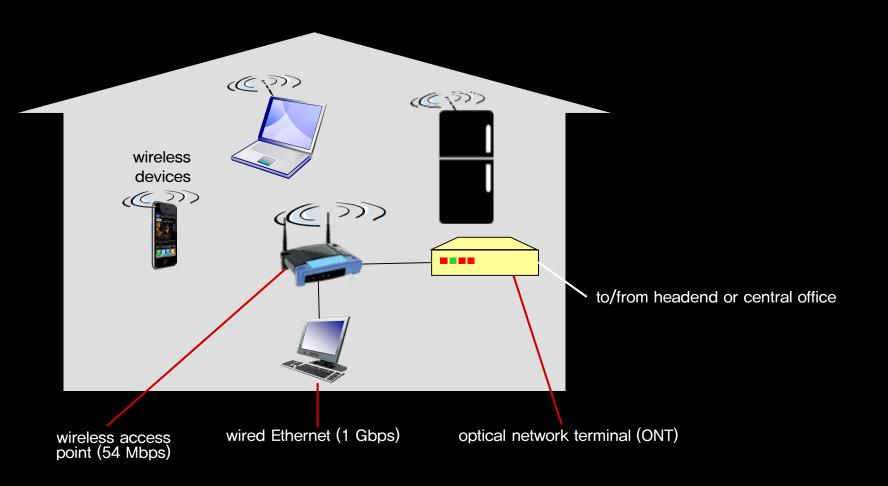




Fiber—to—the—Home (FTTH)

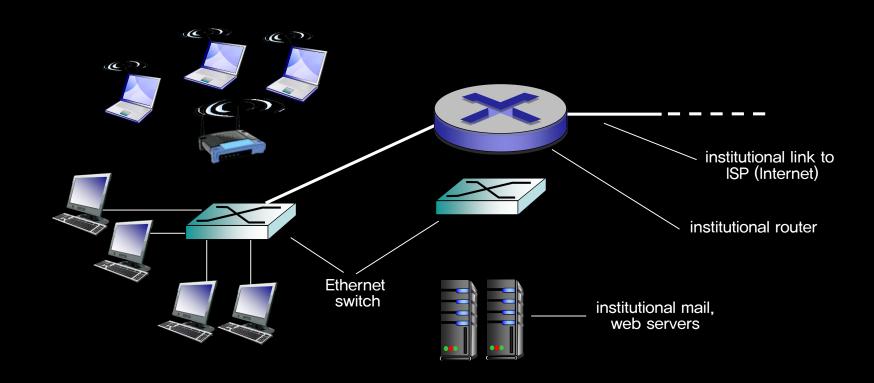


is-optical-fiber-definition-lesson-quiz,html/





- Typically used in companies, universities
- 10Mbps, 100Mbps, 1Gbps, 10Gbps transmission rates





Wireless LAN

■ 802.11g: 54 Mbps

■ 802.11n: 600 Mbps

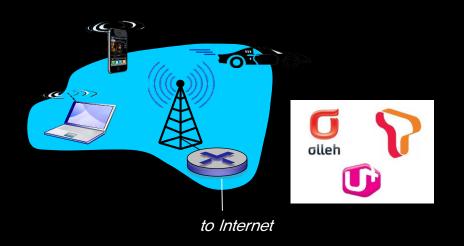
■ 802.11ac: 2.6 Gbps (three 160MHz chs)

Cellular network

■ 3G, 4G LTE, 5G

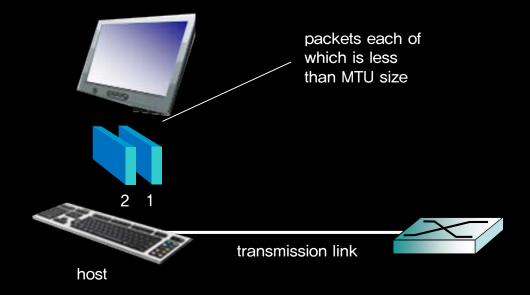
■ 1, 10, 100 Mbps per person







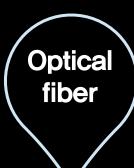
- Takes an application message
- Breaks it into smaller chunks, known as packets, of length less than MTU
- Transmits packet into access network
 - link transmission
 rate, a.k.a. link capacity or link
 bandwidth







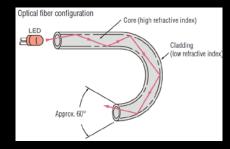




- Two insulated copper wires
- 100 Mbps \sim 10 Gbps
- Two concentric copper conductors
- Broadband
- Multiple channels on cable
 - d Dielectric Jacket

출처 - http://www.standard-wire.com/coax_cable_theory_and_application.html/

- High speed, low error
- Less attenuation
- Immune to electromagnetic noise



출처 – https://www.keyence.co.uk/ss/products/sensor/ sensorbasics/fiber/info//



출처 - https://commons.wikimedia.org/wiki/File:USB_ Twisted_Pair.svg/



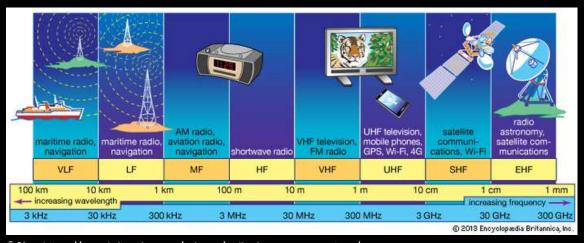
Bandwidth for each service

• WiFi: 2.4 GHz, 5 GHz

■ IMT-Advanced (4G LTE): about 2 GHz

■ IMT-2020 (5G): 30 GHz~300 GHz

■ Satellite communication: 1~40 GHz



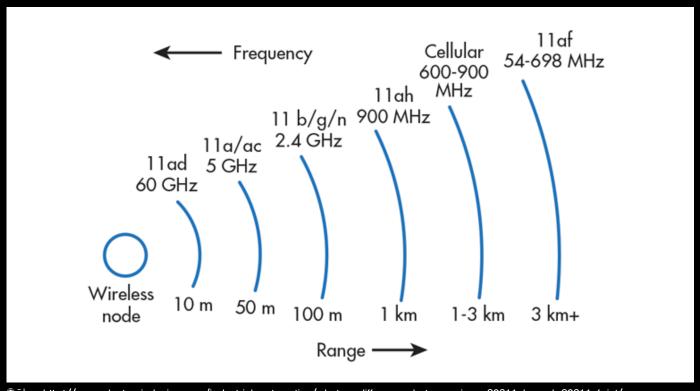
Frequency Spectrum Visible Light Microwaves Service Frequency Designation Designation 0 Hertz 30 - 300 MHz Direct Current (DC) Very High Frequency (VHF) Very Low Frequency (VLF) 3 - 30 KHz Ultra High Frequency (UHF) 300 - 3000 MHz 30 - 300 KHz Super High Frequency (SHF) 3 - 30 GHz Low Frequency (LF) Medium Frequency (MF) 300 - 3000 KHz Extreme High Frequency (EHF) 30 - 300 GHz High Frequency (HF) 3 - 30 MHz Microwave Frequencies 500 MHz - 300 GHz

출처 - https://www.britannica.com/science/radio-frequency-spectrum/

출처 - https://mars.nasa.gov/MPF/rovercom/radio.html/



■ The higher the frequency, the stronger the linearity and the faster the attenuation

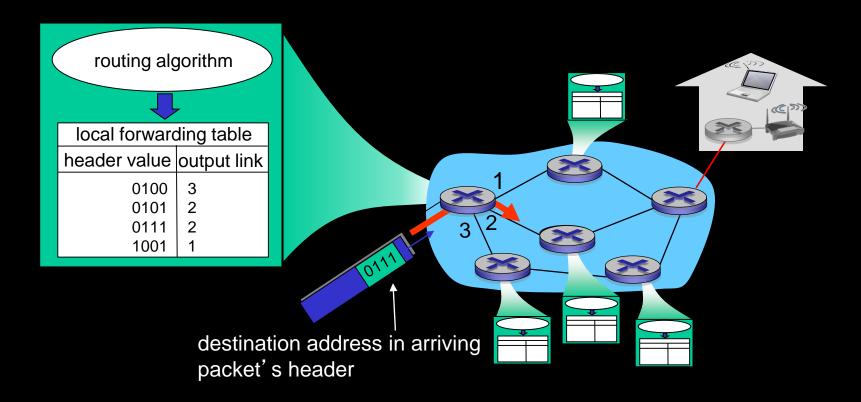


출처 - http://www.electronicdesign.com/industrial-automation/what-s-difference-between-ieee-80211ah-and-80211af-iot/

03. Network Core



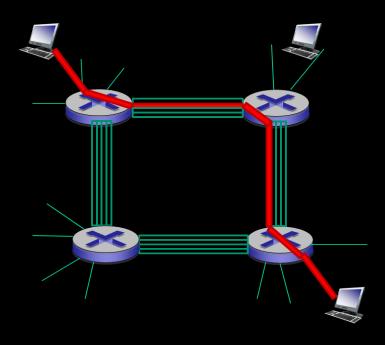
- Mesh of interconnected routers
- Function
 - packet forwarding from one router (or switch) to the next along the path from source to dest.





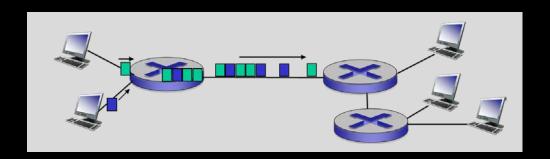
Circuit Switching

- End—end resources reserved for "call" between src. & dest.
- Entire data flew along the path like water



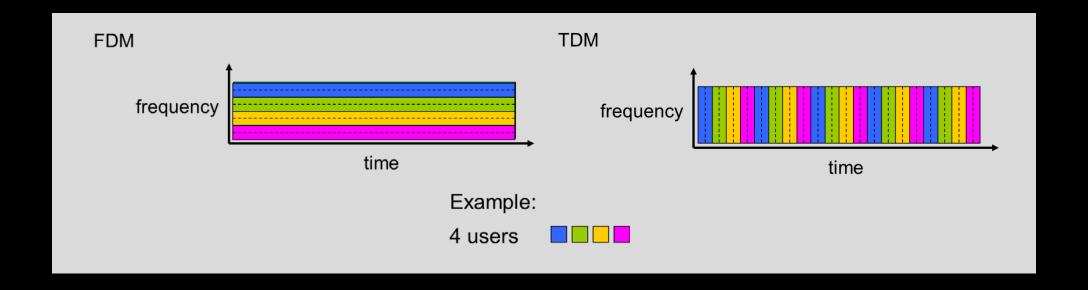
Packet Switching

- Entire data broken into small packets
- Each packet has its destination address
- Each packet handled independently





- Resources dedicated to each call
 - circuit segment idle if not used by the call (no sharing)
- Commonly used in traditional telephone networks
- Channel allocation methods: FDM, TDM

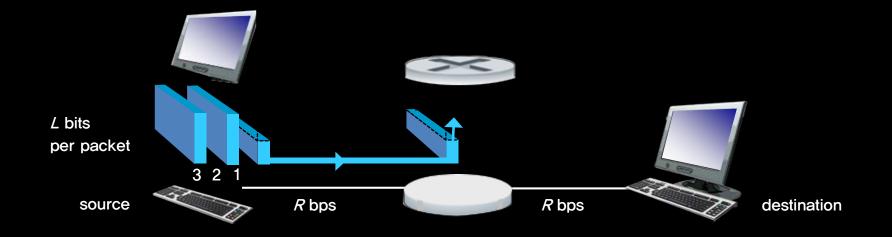




- Packet transmitted at full link capacity
- Takes L/R seconds to transmit L—bit packet into link at R bps
- Store—and—forward: entire packet must arrive at router before forwarded to the next
- End-end delay ≈ 2L/R

numerical example:

- L = 7.5 Mbits
- R = 1.5 Mbps
- One-hop transmission delay = 5s
- End−end delay ≈ 10s



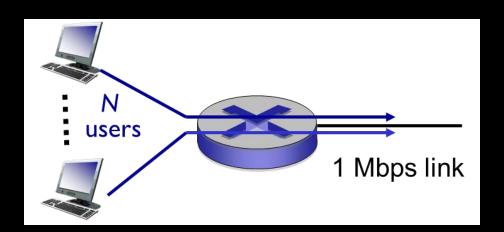
Switching: Comparison



Packet switching allows more users to use network!

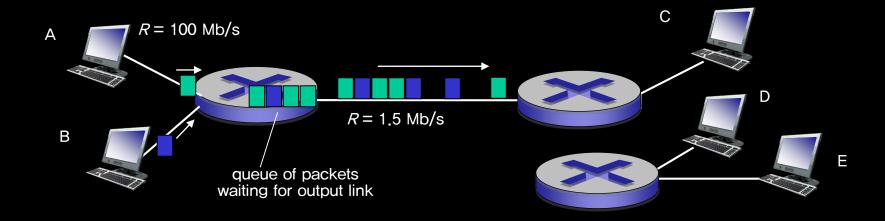
- Example
 - 1 Mbps link
 - each user:
 - 100 kbps when "active"
 - only 10% of time active
 - circuit switching
 - 10 users
 - packet switching
 - with 35 users, probability of 10 users being active at the same time is about .0004

$$1 - \sum_{k=0}^{10} {35 \choose k} \left(\frac{1}{10}\right)^k \left(\frac{9}{10}\right)^{35-k} = 0.0004243$$





- Circuit switching guarantees the quality of service for each call!
- On the other hand, packet switching may suffer from queuing and loss
 - if arrival rate exceeds transmission rate of link
 - packet queued before being transmitted
 - can be dropped (lost) if memory fills up



04. Internet Structure



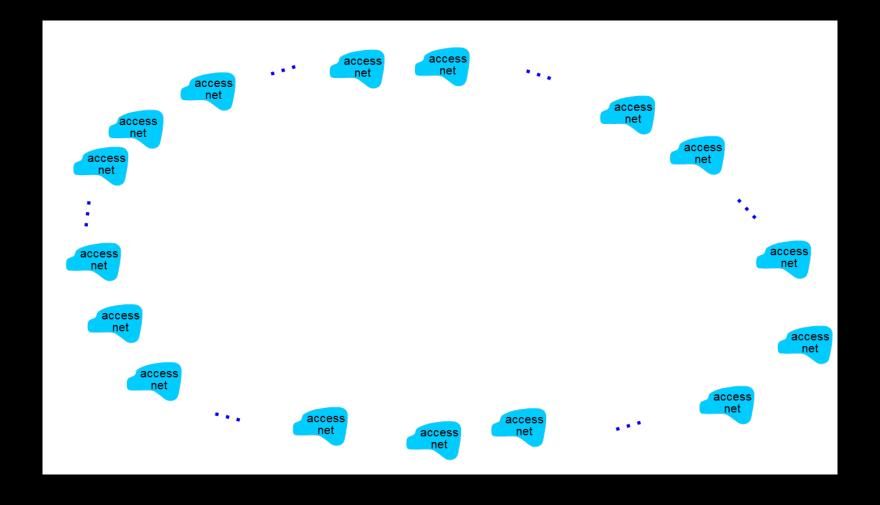
Who's in charge of the Internet?

Nobody! Or Everybody!

- independently operated networks
- End systems connect to Internet via access ISPs (Internet Service Providers)
 - residential, company, and university ISPs
- Access ISPs in turn must be interconnected
 - so that any two hosts can be communicated
- Resulting network of networks is very complex
 - evolution driven by economics and national policy

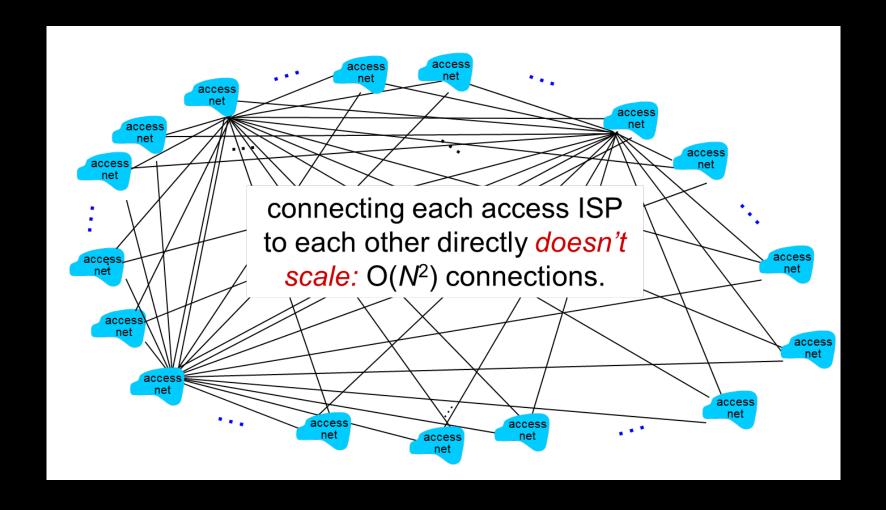


• Question: given millions of access ISPs, how to connect them together?



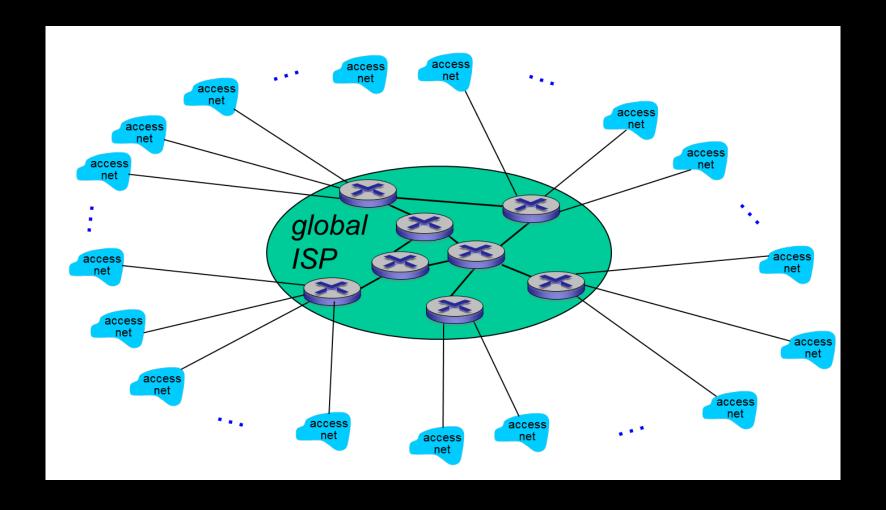


■ Naïve method: connect each access ISP to every other access ISP



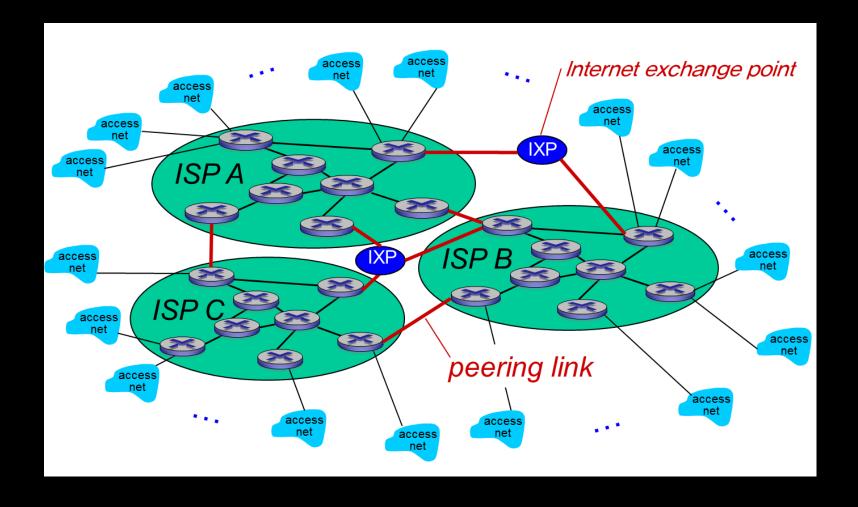


Scalable method: connect each access ISP to one global transit ISP



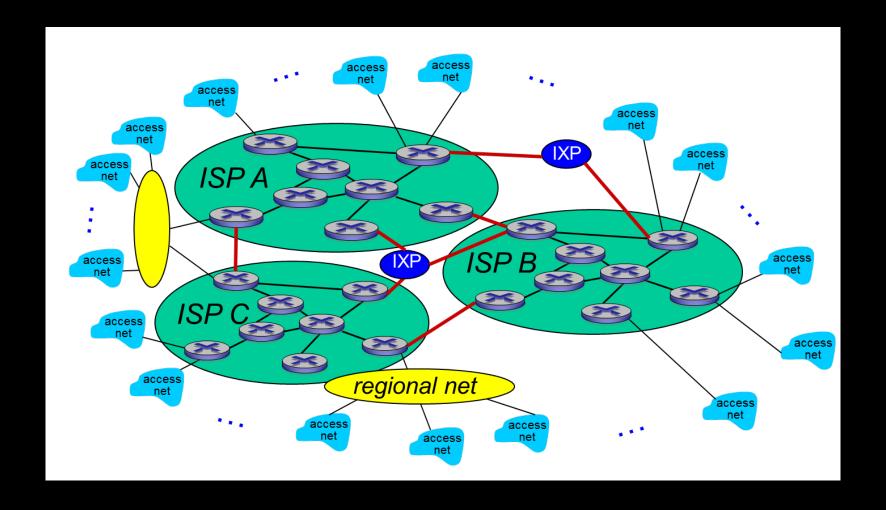


Competing ISPs appear which must be interconnected



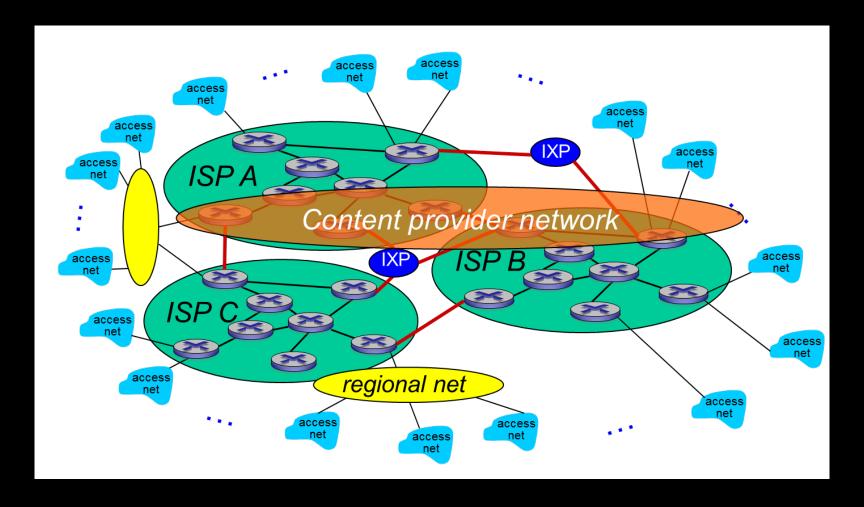


Regional networks arise to connect access networks to ISPs



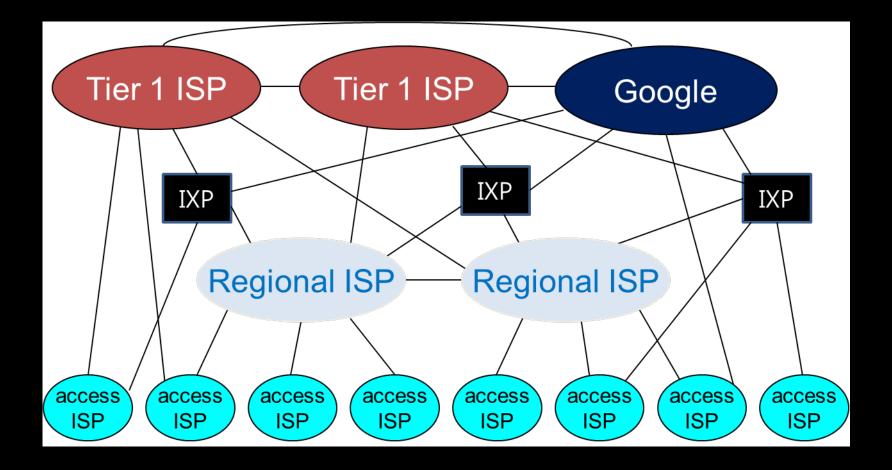


 Contend provider (e.g., Google, Microsoft, Akamai) run their own network to bring services and content close to end users



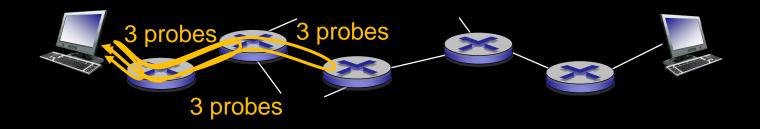


Seen vertically as:





- Nobody in charge of the entire structure…
- Can we see the entire structure?
- Traceroute
 - sends three packets that will reach router i on path towards destination
 - router i will return packets to sender
 - shows interval between transmission and reply





```
- - X
C:\Windows\system32\cmd.exe
             >tracert cisco.com
C:\Users\
Tracing route to cisco.com [72.163.4.161]
over a maximum of 30 hops:
                        <1 ms 192.168.1.1
       <1 ms
               <1 ms
                        32 ms 194.146.109.226
       49 ms
               37 ms
               29 ms
                        53 ms cpe-188-129-0-253.dynamic.amis.hr [188.129.0.253]
       40 ms
      41 ms
               45 ms
                        37 ms ljubljana9-ge-2-5.amis.net [212.18.39.113]
  4
      50 ms
               47 ms
                        81 ms mx-lj1-te-1-2-0.amis.net [212.18.44.137]
      103 ms
                        60 ms mx-vi1-te-0-0-0.amis.net [212.18.44.142]
               72 ms
      53 ms
               53 ms
                        61 ms xe-0-0-0-300.vie20.ip4.tinet.net [77.67.75.93]
     169 ms
              145 ms
                       150 ms xe-10-3-2.was14.ip4.tinet.net [141.136.110.217]
  8
      330 ms
              225 ms
                       303 ms te-7-2.car4.Washington1.Level3.net [4.68.110.97]
     217 ms
                       209 ms vlan60.csw1.Washington1.Level3.net [4.69.149.62]
10
     205 ms
              208 ms
11
                       200 ms ae-61-61.ebr1.Washington1.Level3.net [4.69.134.129]
     209 ms
              185 ms
                       204 ms ae-2-2.ebr3.Atlanta2.Level3.net [4.69.132.85]
     204 ms
                       202 ms ae-7-7.ebr3.Dallas1.Level3.net [4.69.134.21]
13
              204 ms
14
     282 ms
              197 ms
                      210 ms ae-63-63.csw1.Dallas1.Level3.net [4.69.151.133]
     200 ms
              219 ms 230 ms ae-1-60.edge9.Dallas1.Level3.net [4.69.145.16]
     210 ms
              197 ms
                       213 ms CISCO-SYSTE.edge9.Dallas1.Level3.net [4.30.74.46]
16
17
                               Request timed out.
     322 ms
              310 ms
                      329 ms rcdn9-cd2-dmzdcc-gw2-por1.cisco.com [72.163.0.182]
     319 ms
                       315 ms rcdn9-14a-dcz05n-gw1-ten5-5.cisco.com [72.163.0.238]
     324 ms
             299 ms 309 ms www1.cisco.com [72.163.4.161]
20
Trace complete.
C:\Users\
```

05. Performance Metrics



Delay

Packet delivering time from source to destination

Packet Loss

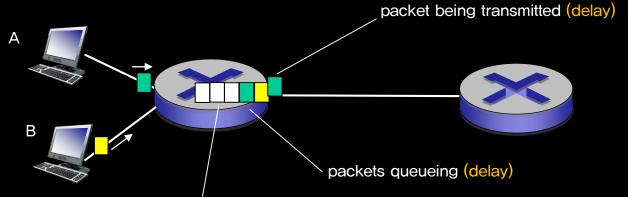
Ratio of lost packets to total sent packets

* PDR (Packet Delivery Ratio): Raio of packets successfully delivered to the total number of packets

Throughput

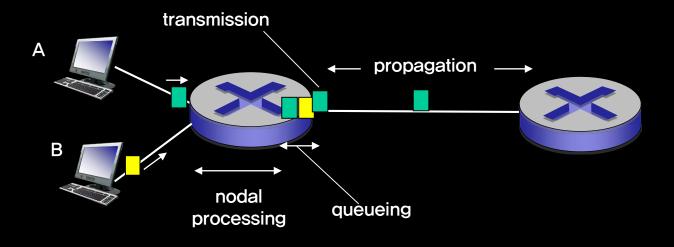
The amount of traffic delivered per unit time

if packet arrival rate to link (temporarily) exceeds output link capacity ...



free (available) buffers: arriving packets dropped (loss) if no free buffers

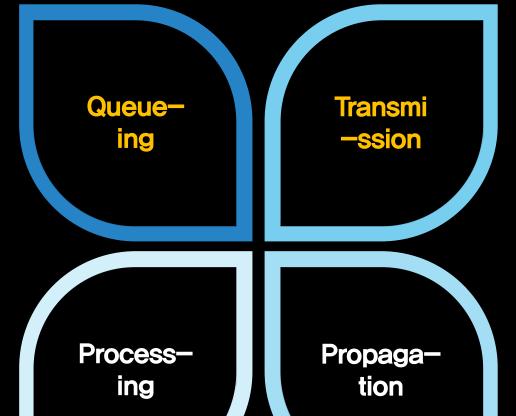




$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$



- time waiting at output buffer for transmission
- congestion dependent



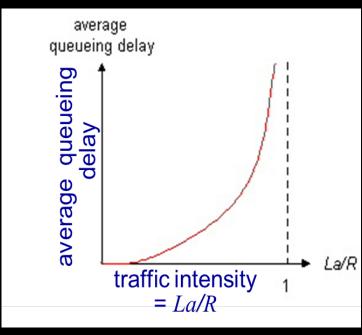
- L: packet length (bits)
- R: link bandwidth (bps)
- $d_{trans} = L/R$

- bit error checking
- decision of output link
- typically < msec

- d: length of physical link
- s: signal speed
 (~2x10⁸ m/s)
- $d_{prop} = d/s$

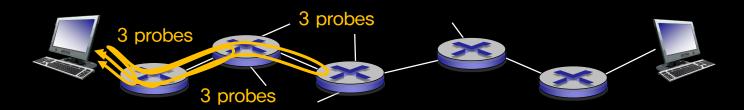
- R: link bandwidth (bps)
- L: packet length (bits)
- a: avg. packet arrival rate

- La/R~0: small queueing delay
- La/R->1: large queueing delay
- La/R~1: increase to the infinite!









traceroute: gaia.cs.umass.edu to www.eurecom.fr

■ 1 cs-gw (128.119.240.254) 1 ms 1 ms 2 ms

2 border1-rt-fa5-1-0.gw,umass,edu (128.119.3.145) 1 ms 1 ms 2 ms

3 cht-vbns.gw.umass.edu (128,119,3,130) 6 ms 5 ms 5 ms

• 4 jn1-at1-0-0-19.wor.vbns.net (204.147.132.129) 16 ms 11 ms 13 ms

■ 5 jn1-so7-0-0-0,wae.vbns.net (204,147,136,136) 21 ms 18 ms 18 ms

6 abilene-vbns.abilene.ucaid.edu (198.32.11.9) 22 ms 18 ms 22 ms

■ 7 nycm-wash,abilene,ucaid.edu (198,32,8,46) 22 ms 22 ms 22 ms

8 62.40.103,253 (62,40.103,253) 104 ms 109 ms 106 ms

de2-1.de1.de.geant.net (62.40.96.129) 109 ms 102 ms 104 ms

...

15 eurecom-valbonne,r3t2,ft,net (193,48,50,54) 135 ms 128 ms 133 ms

■ 16 194,214,211,25 (194,214,211,25) 126 ms 128 ms 126 ms

* means no response (probe lost, router not replying)

18 * * *

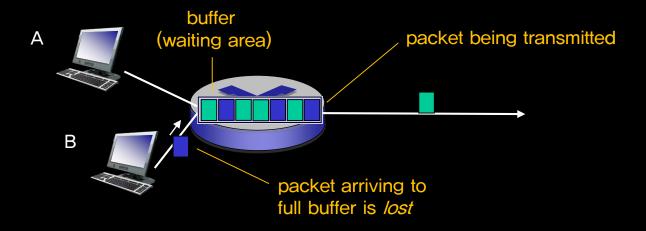
19 fantasia.eurecom.fr (193.55.113.142) 132 ms 128 ms 136 ms

3 delay measurements from gaia.cs.umass.edu to cs-gw.cs.umass.edu

trans-oceanic link

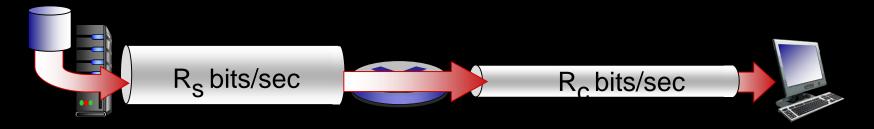


- Packet arriving to full queue dropped
- Lost packet may be retransmitted by previous node, by source end system, or not at all





- Rate at which bits transferred b_w s-d
 - instantaneous: rate at given point in time
 - average: rate over longer period of time



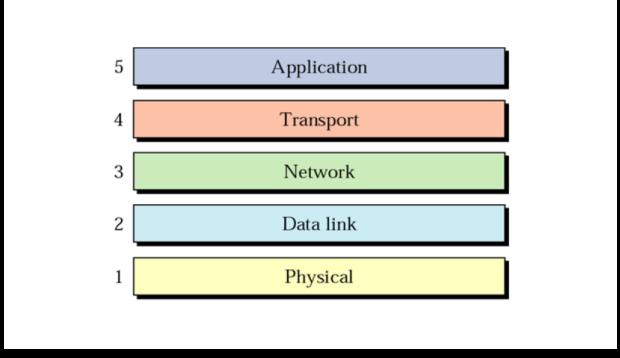
bottleneck link -

link on end-end path that constrains end-end throughput

06. Protocol Stack

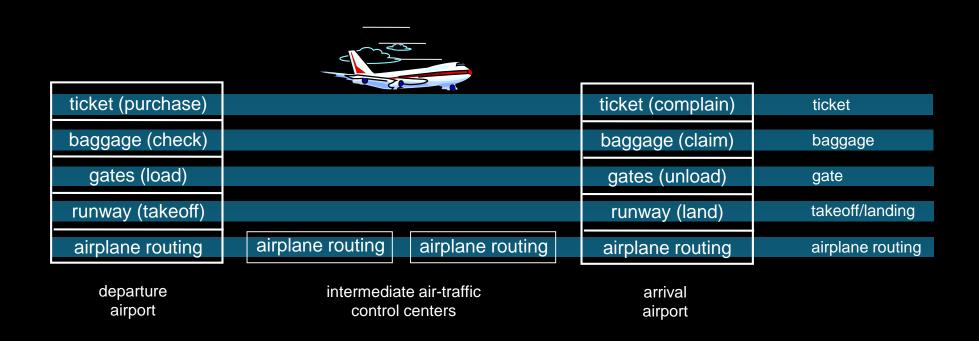


A communication protocol stack is composed of several layers



출처 - http://www.cse.yorku.ca/course_archive/2012-13/F/2041/slides/networks.html/





- layers: each layer implements a service
 - via its own internal—layer actions
 - relying on services provided by layer below



- Modularization eases development, maintenance, and updating of system
 - change of implementation of layer's service transparent to rest of system
 - e.g., change in gate procedure doesn't affect rest of system

Layering considered harmful?



- application: supporting network applications
 - FTP, SMTP, HTTP
- transport: process—process data transfer
 - TCP, UDP
- network: routing of datagrams from source to destination
 - IP, routing protocols
- link: data transfer between neighboring network elements
 - Ethernet, 802.11 (WiFi), PPP
- physical: bits "on the wire"

application

transport

network

link

physical



- presentation: allow applications to interpret meaning of data,
 e.g., encryption, compression, machine—specific
 conventions
- session: synchronization, connection management, recovery of data exchange
- Internet stack "missing" these layers!
 - these services, if needed, must be implemented in application

application

presentation

session

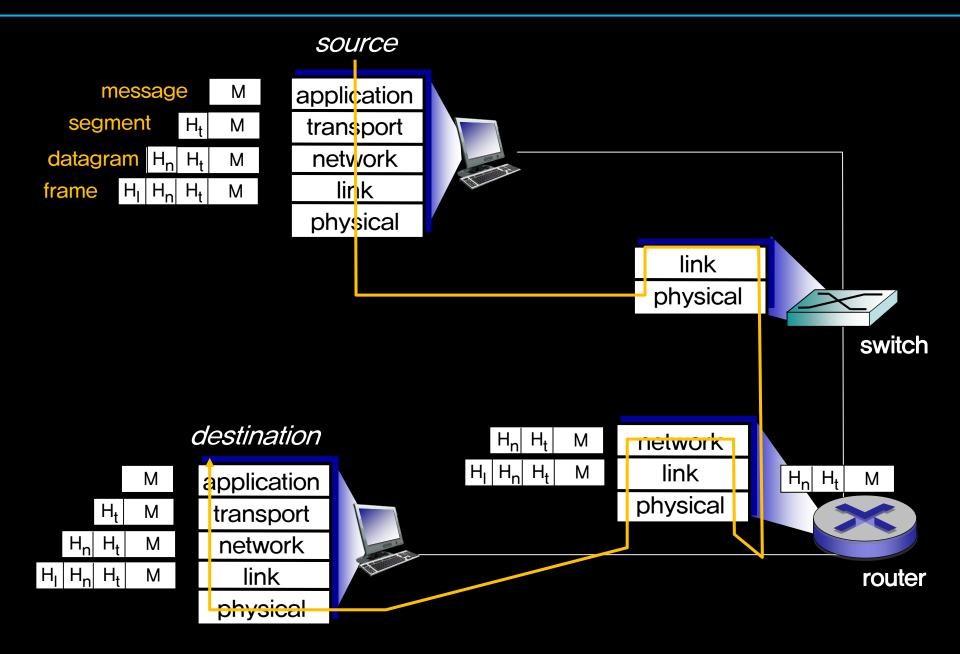
transport

network

link

physical





07. Network Security



Field of network security

- how bad guys can attack computer networks
- how we can defend networks against attacks
- how to design architectures that are immune to attacks

Internet not originally designed with (much) security in mind

- original vision: "a group of mutually trusting users attached to a transparent network" ©
- Internet protocol designers playing "catch—up"
- security considerations in all layers!



Malware can get in host from:

- virus: self-replicating infection by receiving/executing object (e.g., e-mail attachment)
- worm: self-replicating infection by passively receiving object that gets itself executed

Malware can

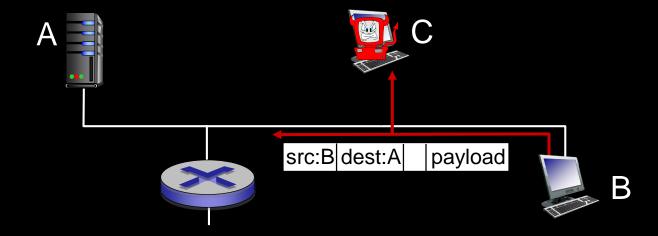
- record keystrokes, web sites visited, upload info to collection site (spyware)
- require ransom or destroy system (ransomware)



출저 - https://epatientfinder.com/worried-ransomware-need-backup-plan-literally/

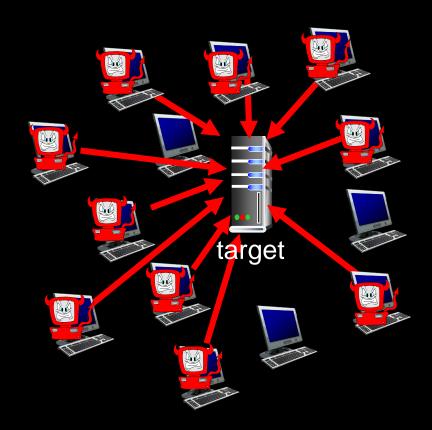


- Broadcast media (shared Ethernet, wireless)
- Promiscuous network interface reads/records all packets (e.g., including passwords!)
 passing by



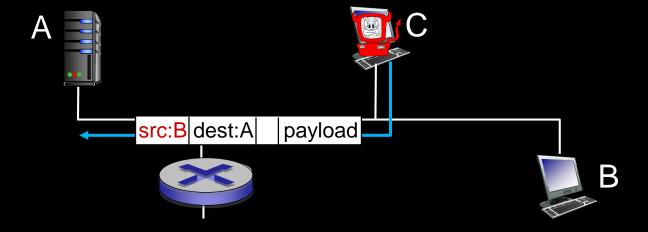


Attackers make resources (server, bandwidth)
 unavailable by sending the huge amount of
 bogus traffic through botnets





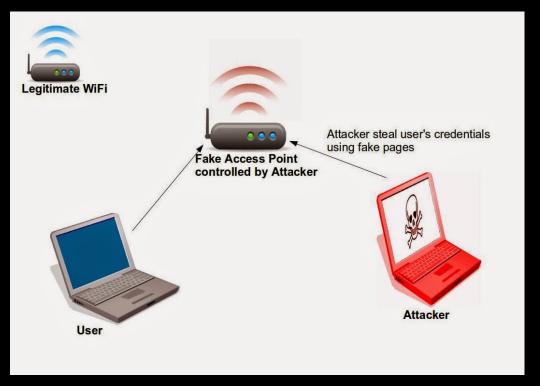
- Send packets with fake source address
- a.k.a. "Sybil" attack







Steal user's credentials using fake AP



출저 – https://developertz.wordpress.com/2017/03/14/evil-twin-and-fake-wireless-access-point-hacks-what-they-are-how-to-defend//

08. History of the Internet

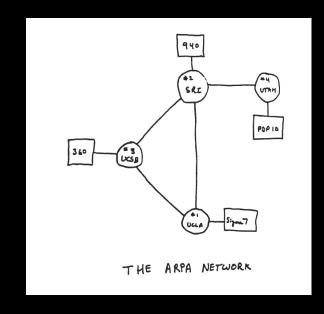


Early packet—switching principles

- **1961**
 - Kleinrock queueing theory shows effectiveness of packet– switching
- **1**967
 - ARPAnet conceived by Advanced Research Projects Agency
- **1969**
 - first ARPAnet node operational
- **1972**
 - ARPAnet public demo
 - NCP (Network Control Protocol) first host-host protocol
 - first e-mail program
 - ARPAnet has 15 nodes



Leonard Kleinrock





- Internetworking, new and proprietary nets
 - **1974**
 - Cerf and Kahn architecture for interconnecting networks



출저 - http://www.amongtech.com/unsung-heroes-internet-pioneers-youve-never-heard/

- **1976**
 - Ethernet at Xerox PARC
- **1979**
 - ARPAnet has 200 nodes

Cerf and Kahn's internetworking principles:

- minimalism, autonomy no internal changes required to interconnect networks
- best effort service model
- stateless routers
- decentralized control

define today's Internet architecture



- New protocols, a proliferation of networks
 - **1**983
 - deployment of TCP/IP
 - **1982**
 - smtp e-mail protocol defined
 - **1983**
 - DNS defined for name—to—IP—address translation
 - **1985**
 - ftp protocol defined
 - **1988**
 - TCP congestion control



Bill Joy 출저 – http://fortune.com/2011/04/06/bill-joy-better-batteries-key-to-green-power-adoption/

Internet History (1990~2000)



Commercialization, the Web, new apps

- **1991**
 - NSF lifts restrictions on commercial use of NSFnet (decommissioned, 1995)
- Early 1990's: Web
 - HTML, HTTP: Tim Berners—Lee
 - 1994: Mosaic, later Netscape
 - Commercialization of the Web
- Late 1990's-2000's
 - more killer apps: instant messaging, P2P file sharing
 - network security to forefront
 - est. 50 million host, 100 million+ users



Tim Berners-Lee 출저 - http://techland.time.com/2012/09/05/10-auestions-for-sir-tim-berners-lee/



- 5B devices attached to Internet (2016)
 - smartphones and tablets
- Aggressive deployment of broadband access
- Increasing ubiquity of high—speed wireless access
- Emergence of online social networks:
 - Facebook: ~ one billion users
- Service providers (Google, Microsoft) create their own networks
 - bypass Internet, providing "instantaneous" access to search, video content, email, etc.
- e-commerce, universities, enterprises running their services in "cloud" (e.g., Amazon EC2)

Summary

01

Internet overview

- Internet: "network of networks"
- Internet is composed of network edge, access network, core network, and protocol

02

What's a protocol?

- rules and regulations to transfer data in computer networks
- message format, order of messages, actions

03

Network edge, core, access network

- access network: a network connect edge and core
- core network techniques: packet switching vs. circuit switching

04

Internet structure

- independently operated networks
- hierarchical structure of access ISPs, regional ISPs, global ISPs, and so on

05

Network performance

- Delay (transmission, queueing, processing, propagation)
- Packet loss, packet delivery rate
- Throughput

06

Protocol stack

- Layering: advantages of modularization
- Internet protocol stack
- ISO/OSI reference model

07

Network security

- Internet not originally designed with (much) security in mind
- Internet protocol designers playing "catch-up"

08

History of the Internet

- Firstly developed as ARPAnet
- Internetworking architecture = autonomy + minimalism
- TCP/IP, World Wide Web