

VLSI CAD: Logic to Layout

Rob A. Rutenbar
University of Illinois

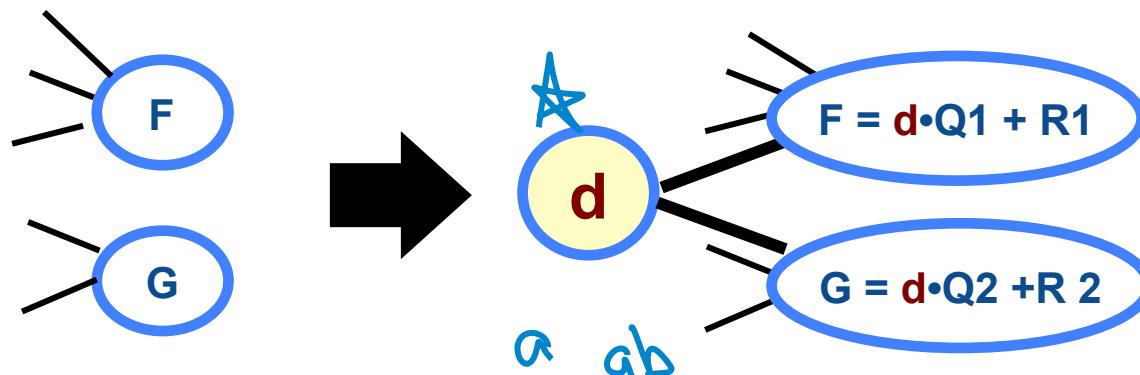
Lecture 7.1 Logic Synthesis: Multilevel Logic and Divisor Extraction – Single Cube Case



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How Do We Find Good Divisors?

- Called: **Extraction**. How do we extract good divisors?
 - Extraction of a **single-cube divisor** from multiple expressions?
 - Extraction of a **multiple-cube divisor** from multiple expressions?



- Remember: if we want a single-cube divisor: **Go look for co-kernels** ↩
- Remember: if we want a multiple-cube divisor: **Go look for kernels** ↩

$a + bc$



How We Do It

- **For single cube extraction**
 - Build a very large matrix of 0s and 1s
 - Heuristically look for “prime rectangles” in this matrix
 - Each such “prime” is a good common single-cube divisor
- **For multiple cube extraction**
 - Build a (different) very large matrix of 0s and 1s
 - Heuristically look for “prime rectangles” in this matrix
 - Each such “prime” is a good multiple-cube divisor
- **Surprisingly, a lot like Kmaps(!)**
 - Except we do it all automatically algorithmically. Still, unifying and familiar idea...



Single Cube Extract: Matrix Representation

- Given: a set of SOP algebraic model Boolean equations (P, Q, R)
- Construct the cube-literal matrix as follows
 - 1 row for each unique product term; 1 column for each unique literal
 - A “1” in the matrix if this product term uses this literal, else a “.”
 - Number each row and column for convenience in referring to them, later

$$\begin{aligned}P &= \textcircled{abc} + abd + eg \\Q &= abfg \\R &= bd + ef\end{aligned}$$

	a	b	c	d	e	f	g
1	1	2	3	4	5	6	7
abc	1
abd	2	1	1	.	1	.	.
eg	3	1	.
abfg	4	1	1	.	.	1	1
bd	5	.	1	.	1	.	.
ef	6	1	1

Example from: Richard Rudell, Logic Synthesis for VLSI Design, PhD Thesis, Dept of EECS, University of California at Berkeley, 1989.



Covering this Matrix: Prime Rectangles

- A **rectangle** of this cube-literal matrix...
 - Set of rows **R** and columns (cols) **C** that has a '1' in every row/column intersection
 - Need **not** be contiguous rows or cols in matrix. **Any set** of rows or columns
- A **prime rectangle**:
 - A rectangle that **cannot** be made any bigger by adding another row or a column

Prime rectangle example

	a	b	c	d	e	f	g
abc	1	1	1	1	.	.	.
abd	2	1	1	.	1	.	.
eg	3	1	.
abfq	4	1	1	.	.	1	1
bd	5	.	1	.	1	.	.
ef	6	.	.	.	1	1	.



Result: Prime Rectangle Columns = Divisor!

- **Primes are “biggest possible” common single-cube divisors**
 - **Makes sense:** columns of (R,C) tell you the literals in the product term divisor (ie. If you AND these columns), rows tell you which product terms you can divide!

$$\begin{aligned}P &= abc + abd + eg \\Q &= abfg \\R &= bd + ef\end{aligned}$$

	a 1	b 2	c 3	d 4	e 5	f 6	g 7
abc 1	1	1	1
abd 2	1	1	.	1	.	.	.
eg 3	1	.	1
abfg 4	1	1	.	.	.	1	1
bd 5	.	1	.	1	.	.	.
ef 6	1	1	.

1 cube divisor

$$X = ab$$

$$P = Xc + Xd + eg$$

$$Q = Xfg$$

$$R = bd + ef$$

$$X = ab$$

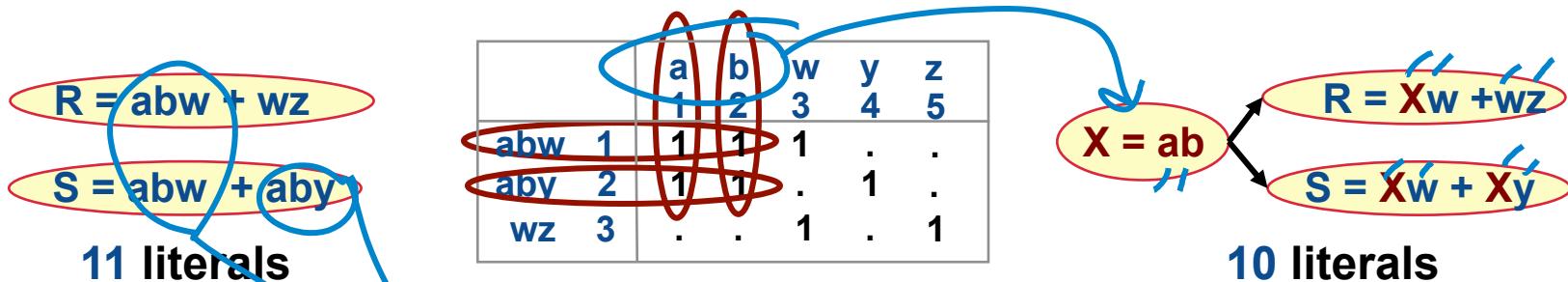


Simple Bookkeeping to Track Literals

- Recall: we factor – extract– to **reduce literals** in logic network
 - Would be nice if there was a simple formula to compute this. (There is!)
 - Start with a prime rectangle
 - Let **C** = number of ("#") columns in rectangle
 - For each **row r** in rectangle: let **Weight(r)** = # times this product appears in network
 - Compute **L** = $(C-1) \times [\sum_{\text{rows } r} \text{Weight}(r)] - C$
- Nice result: for a prime rectangle, **L** = # literals saved
 - To be precise: if you count literals *before* extracting this divisor cube, and *after*, **L** is how many literals are saved



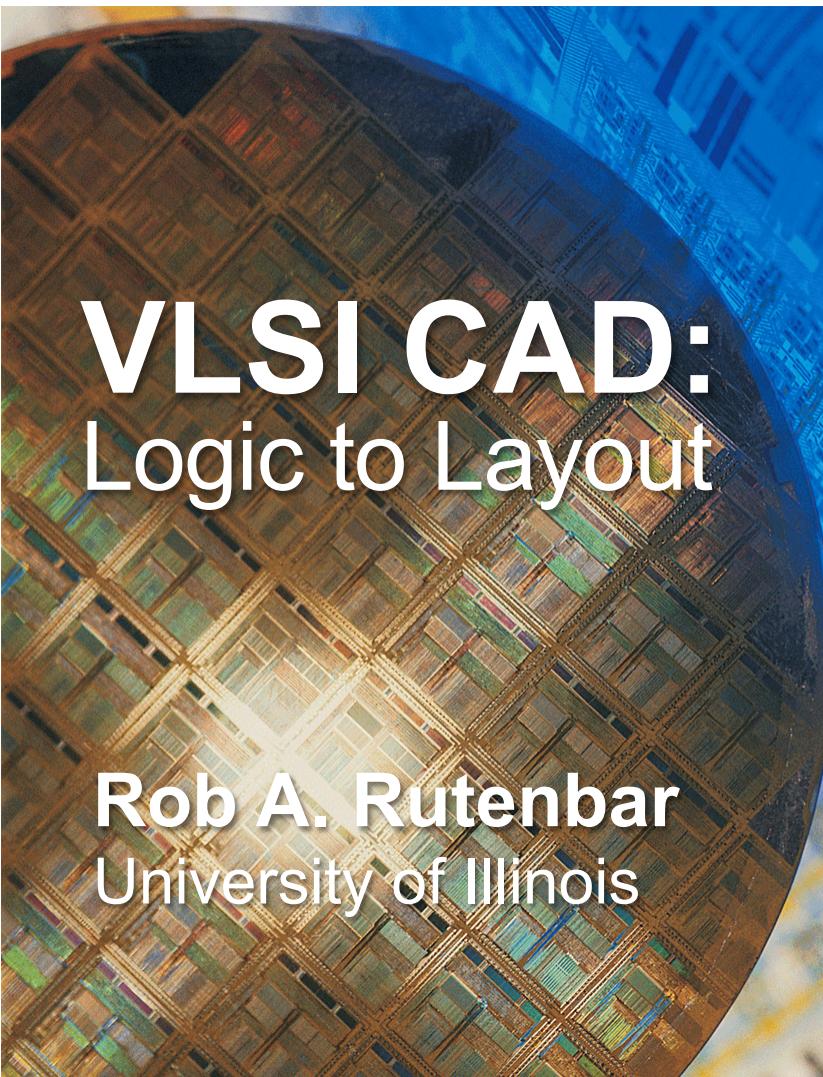
Compute Saved Literals, Single-Cube Extract



- Literals saved formula: Compute $L = (C-1) \times [\sum_{\text{rows } r} \text{Weight}(r)] - C$
 - $C = 2 \text{ columns}$
 - $\text{Weight}(abw) = 2$ (appears twice in network). $\text{Weight}(aby) = 1$ (appears once)
 - $L = (2-1)(2+1) - 2 = 1$ it works! We did, in fact, **save just 1 literal.**

Counts all places **ab** appeared (6 literals), the savings to replace each with **X** (3 literals), and new overhead of adding **X=ab** term (+2 literals). $-6+3+2 = 1 \text{ saved.}$





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Lecture 7.2

Logic Synthesis: Multilevel Logic and Divisor Extraction – Multiple Cube Case



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Multiple-Cube Factors...?

- Remarkably, a very *similar* matrix-rectangle-prime concept
 - Make an appropriate **matrix**. Find **prime rectangle**. Do literal count **bookkeeping** with numbers associated with rows/columns. It all just works! Nice, **unifying** idea.
 - Given: A set of Boolean functions (nodes in a network)

$$P = af + bf + ag + cg + ade + bde + cde$$

$$Q = af + bf + ace + bce$$

$$R = ade + cde$$

- First: find **co-kernels and kernels** of each of these functions
 - Why? Brayton-McMullen says: Multiple-cube factors are **intersections of the product terms in the kernels** for each of these functions



Kernels / Co-Kernels of P,Q,R Example

- **P = af + bf + ag + cg + ade + bde + cde**
 - Co-kernel **(a)** kernel **(de + f + g)** ✓
 - Co-kernel **(b)** kernel **(de + f)** ✓
 - Co-kernel **(de)** kernel **(a + b + c)** ✓
 - Co-kernel **(f)** kernel **(a + b)**
 - Co-kernel **(c)** kernel **(de + g)**
 - Co-kernel **(g)** kernel **(a+c)**
 - [Co-kernel **(1)** kernel **(af + bf + ag + cg + ade + bde + cde)** trivial, **ignore**] ▲
- **Q = af + bf + ace + bce**
 - Co-kernel **(a)** or co-kernel **(b)** kernel **(ce + f)** (can have > 1 co-kernel) ▲
 - Co-kernel **(f)** or co-kernel **(ce)** kernel **(a + b)** (ditto) ▲
 - Co-kernel **(1)** kernel **(af + bf + ace + bce)** trivial, **ignore**
- **R = ade + cde**
 - Co-kernel **(de)** kernel **(a + c)**
 - (Note, **R** is not its own kernel; why not?)

Example from: Richard Rudell, Logic Synthesis
for VLSI Design, PhD Thesis, Dept of EECS,
University of California at Berkeley, 1989.



New Matrix: Co-Kernel--Cube Matrix

P	co-kernel (a)	kernel	$(de + f + g)$
P	co-kernel (b)	kernel	$(de + f)$
P	co-kernel (de)	kernel	$(a + b + c)$
P	co-kernel (f)	kernel	$(a + b)$
P	co-kernel (c)	kernel	$(de + g)$
P	co-kernel (g)	kernel	$(a+c)$
Q	co-kernel (a)	kernel	$(ce + f)$
Q	co-kernel (b)	kernel	$(ce + f)$
Q	co-kernel (f)	kernel	$(a + b)$
Q	co-kernel (ce)	kernel	$(a + b)$
R	co-kernel (de)	kernel	$(a + c)$

One row for each **unique (Function, Co-kernel) pair** in problem

One column for each **unique cube** among all kernels in problem

	a	b	c	ce	de	f	g
	1	2	3	4	5	6	7
P	a	1					
P	b	2					
P	de	3					
P	f	4					
P	c	5					
P	g	6					
Q	a	7					
Q	b	8					
Q	ce	9					
Q	f	10					
R	de	11					

What goes in the individual (row, column) entries here in the matrix?



Entries in the Co-Kernel--Cube Matrix

- From the **row**, take the **co-kernel**, go look at the associated kernel
- Look at **cubes** in this kernel; put “1” in **columns** that are cubes in this kernel; else put “.”

$$P = af + bf + ag + cg + ade + bde + cde$$

P co-kernel (a) kernel $(de + f + g)$

P co-kernel (g) kernel $(a+c)$

In effect, each row says
“here is a function, if I divide it
by **this** co-kernel cube, **that**
is the result (kernel) I get...”

	a 1	b 2	c 3	ce 4	de 5	f 6	g 7
P a 1	1	1	1
P b 2	1	1	.
P de 3	1	1	1
P f 4	1	1
P c 5	1	.
P g 6	1	.	1
Q a 7	.	.	.	1	.	1	.
Q b 8	.	.	.	1	.	1	.
Q ce 9	1	1
Q f 10	1	1
R de 11	1	.	1



Prime Rectangles in Co-Kernel--Cube Matrix

- **Prime rectangle** is again a good divisor: now multiple cube(!)
 - Create the multiple cube divisor as **sum** (OR) of cubes of prime rectangle **columns**

	a	b	c	ce	de	f	g
P a	1	1	.	.	1	1	1
P b	2	.	.	.	1	1	.
P de	3	1	1	1	.	.	.
P f	4	1	1	.	1	.	.
P c	5	.	.	.	1	.	1
P g	6	1	.	1	.	.	.
Q a	7	.	.	1	.	1	.
Q b	8	.	.	1	1	.	.
Q ce	9	1	1	1	.	.	.
Q f	10	1	1	.	1	.	.
R de	11	1

Interpretation

$$P = (\text{co-kernel } de) \cdot (a + b)$$

$$P = (\text{co-kernel } f) \cdot (a + b)$$

$$Q = (\text{co-kernel } ce) \cdot (a + b)$$

$$Q = (\text{co-kernel } f) \cdot (a + b)$$

+ stuff1 + rem1

+ ~~stuff2~~ + rem2

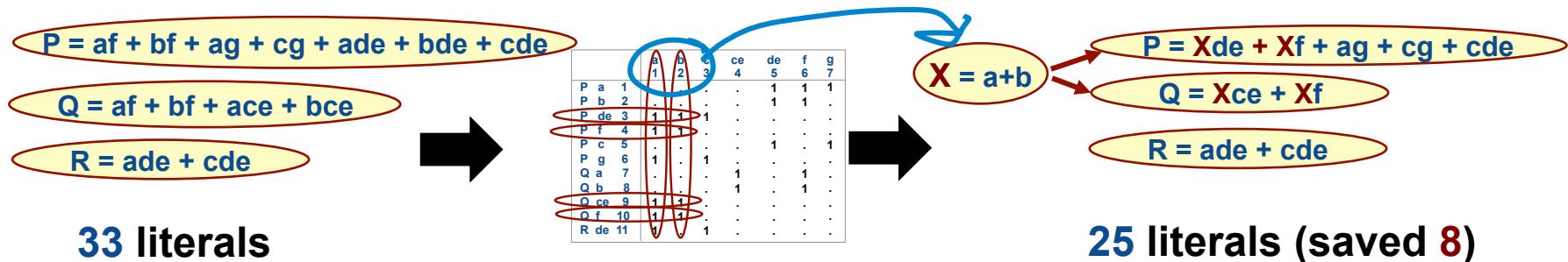
+ ~~stuff3~~ + rem3

+ ~~stuff4~~ + rem4

OR together cubes of the **columns** of the prime rectangle – this is the divisor **(a+b)**!



Multiple-Cube Extraction: # Literals Saved



- **Similar formula**

- For each **column c** in rectangle: let **Weight(c)** = # literals in column cube
- For each **row r** in rectangle: let **Weight(r)** = 1 + # literals in co-kernel label
- For each “1” covered at **row r, column c**: AND row co-kernel and column cube. Let **Value(r,c)** = # literals in this new ANDed product
- Compute $L = \sum_{\text{rows } r, \text{columns } c} \text{Value}(r,c) - \sum_{\text{rows } r} \text{Weight}(r) - \sum_{\text{column } c} \text{Weight}(c)$
- **Result:** L is the number of **saved literals** from this extraction



Multiple-Cube Extraction: # Literals Saved

		a	b	c	ce	de	f	g
Row weights		1	2	3	4	5	6	7
P	a	1	1	1
P	b	2	1	.
P	de	3	1	1	1	.	.	.
P	f	4	1	1
P	c	5	.	.	1	.	.	1
P	g	6	1
Q	a	7	1	.
Q	b	8	1	.
Q	ce	9	1	1	1	.	.	.
Q	f	10	1	1	1	.	.	.
R	de	11	1	.	1	.	.	.

\sum element values = 20

\sum row weights = 10

\sum column weights = 2

$$(\text{value sum}) - (\text{row sum}) - (\text{column sum}) = 20 - 10 - 2 = 8$$

Before

33 literals

$$P = af + bf + ag + cg + ade + bde + cde$$

$$Q = af + bf + ace + bce$$

$$R = ade + cde$$

After

25 literals

$$X = a+b \rightarrow P = Xde + Xf + ag + cg + cde$$

$$X = a+b \rightarrow Q = Xce + Xf$$

$$X = a+b \rightarrow R = ade + cde$$

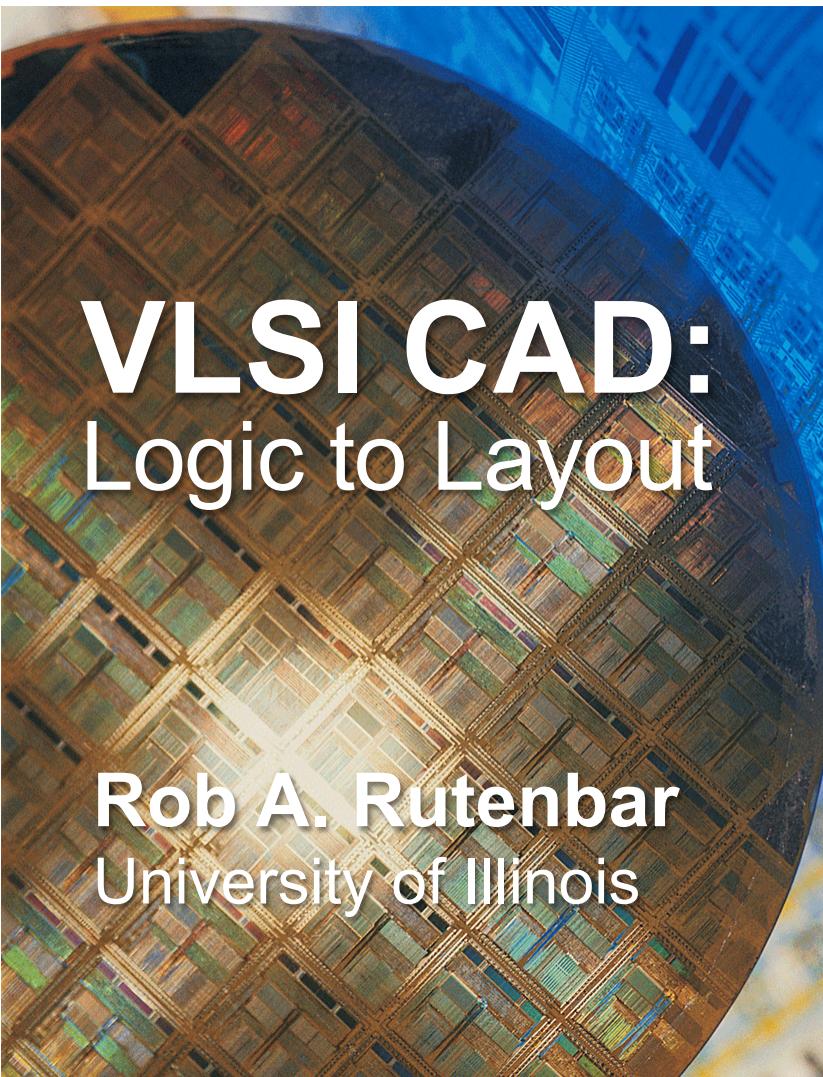
Change in #literals = 33 – 25 = 8 !



Details for Both Single/Multiple Cube Extract

- You can extract a **second, third, etc, divisor using same matrix**
 - Works for the single-cube divisors
 - Works for the multiple-cube divisors
- ...but must **update matrix** to reflect new Boolean logic network
 - Because the node contents are different, and there is a new divisor node in network
 - Like adding “don’t cares” to a Kmap: some things no longer essential to cover
 - And for multiple-cube case, must kernel new divisor nodes to update matrix
 - All mechanical. A bit tedious, to be honest. Just skip it for this presentation
 - **For us:** just know how to **extract first good divisor**, this is good enough





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Lecture 7.3

Logic Synthesis: Multilevel Logic and Divisor Extraction – Finding Prime Rectangles & Summary



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How to *Find* a Prime Rectangle in Matrix?

- **Greedy heuristics** work well for this **rectangle covering** problem
 - Start with a single row rectangle with “good #literal savings”
 - Grow the rectangle alternatively by adding more rows, more columnss
- **Example: Rudell’s Ping Pong heuristic (from his ‘89 Berkeley PhD)**
 - To grow a big prime rectangle, from a good 1-row rectangle
 1. Pick the **best single row** (the 1-row rectangle with best #literals saved)
 2. Look at **other rows with 1s** in same places (and more!) that you can add, one at a time, that each maximize #literals saved. Add them until can’t find any more.
 3. Look at **other columns with 1s** in same places (and more!) that you can add, one at a time, that maximize # literals saved. Add them until you can’t find any more.
 4. Goto 2.
 5. Quit when cannot grow rectangle any more in any direction.



Ping Pong Heuristic -- How Good?

- **Very. Results from Rudell's 1989 Berkeley PhD**
 - **PING PONG:** this heuristic that grows a good prime alternatively by rows, columns
 - **BEST RECT:** uses an aggressive search to find the ideal **best rectangle**
 - Examples are industrial circuits, goal is single cube extraction

Ex	Inputs	Outputs	Initial #Literals	Using PINGPONG		Using BESTRECT	
				#Literals	CPU	#Literals	CPU
apex1	45	45	2,887	1,314	14	1,304	858
apex2	39	3	15,531	3,996	193	3,901	3,408
apex5	117	88	7,369	3,798	43	3,779	343

- Relevant result: within **<1% of optimal, 10-100x faster** than brute force



Extraction: Summary

- **Single cube extraction**
 - Build the cube-literal matrix; prime rectangle is a good **single cube divisor**
 - Simple bookkeeping lets us **estimate savings in #literals** in Boolean network
- **Multiple cube extraction**
 - Kernel all the expressions in network; build the co-kernel--cube matrix
 - Each prime rectangle is a good **multiple cube divisor**
 - Simple bookkeeping lets us **estimate savings in #literals** in Boolean network
- **Mechanically, both rectangle cover problems (very like Kmaps!)**
 - Good **heuristics** for solving for a good prime rectangle, fast and effective
 - Ways to extract *more* than one divisor from network (but we didn't cover these)



Aside: How to We Really Do This?

- ...**not** with rectangle covering on ***all* kernels/co-kernels**
 - Too expensive to do rectangle problem on big circuits (> 20K gates)
 - Too expensive to go compute **complete** set of kernels, co-kernels
- **Often use heuristics to find a “quick” set of likely divisors**
 - Don’t fully kernel each node of network: too many cubes to consider here
 - Instead, can extract a ***subset*** of useful kernels quickly
 - Then, can either do rectangle cover on these smaller problems (smaller since fewer things to consider in covering problem)...
 - ...or, try to do simpler overall network restructuring, eg, try all pairwise substitutions of one node into another node: keep good ones, continue in a greedy way



Notes:

- **Good references to read about all these ideas**
 - R.K. Brayton, R. Rudell, A. Sangiovanni-Vincentelli, A.R. Wang, "MIS: A Multiple-Level Logic Optimization System," *IEEE Transactions on CAD of ICs*, vol. CAD-6, no. 6, November 1987, pp. 1062-1081.
 - Giovanni De Micheli, *Synthesis and Optimization of Digital Circuits*, McGraw-Hill, 1994.
 - R.K. Brayton, R. Rudell, A.S. Vincentelli, and A. Wang, "Multi-Level Logic Optimization and the Rectangular Covering Problem," Proceedings of the International Conference on Computer Aided Design, pp. 66-69, 1987.
 - Richard Rudell, *Logic Synthesis for VLSI Design*, PhD Thesis, Dept of EECS, University of California at Berkeley, 1989.
 - Srinivas Devadas, Abhijit Gosh, Kurt Keutzer, *Logic Synthesis*, McGraw Hill Inc., 1994.

