

System Software

System Software

- System software is a set of programs that support the operation of a computer
- Hides the complexities of the hardware from the user
- The user can focus on an application or other problem to be solved without needing to know the details of how the machine works internally.
- Controls the operation and extends the processing functionalities of a computer system.
- Makes the operation of a computer more fast, effective, and secure.
- Application Software is a program that does real work for the user. It is mostly created to perform a specific task for a user.

System Software

- Compiler
- Assembler
- Interpreter
- Text Editor
- Debugger
- Macro processor
- Loader
- Linker

Introduction

System Software

- General purpose software
- Intended to support the operation and use of computer
- Related to the architecture of the machine on which they are to run
- Machine dependent

Application Software

- Specific purpose software
- Primarily concerned with the solution of some problem using the computer as a tool
- Focus is on the application, not on the computing system
- Machine independent

System Software Vs Machine Architecture

- One characteristic in which most system software differ from application software is ***machine dependency***
 - assembler translate mnemonic instructions into machine code
 - compilers must generate machine language code
 - operating systems are directly concerned with the management of nearly all of the resources of a computing system
- There are some aspects of system software that do not directly depend upon the type of computing system
 - general design and logic of an assembler
 - general design and logic of a compiler
 - code optimization techniques

Simplified Instructional Computer(SIC) Machine

- Hypothetical computer that includes the hardware features most often found on real machines
- Two versions:
 - Standard model, SIC
 - XE version (SIC/XE) (extra equipment or extra expensive)
 - Upward compatible
 - An object program for the standard SIC will also execute properly on a SIC/XE system
- **SIC Machine Architecture**
- Memory
 - 8 bit bytes
 - 3 consecutive bytes from a word
 - All addresses are byte addresses
 - words addressed by the location of the lowest numbered byte
 - Address=15 bits
 - Memory size= 2^{15} bytes

SIC Machine Architecture

- Registers
- Five registers, 24 bit length

Mnemonic	Number	Special Use
A	0	Accumulator; used for arithmetic operations
X	1	Index register; used for addressing
L	2	Linkage register; the Jump to Subroutine (JSUB) instruction stores the return address in this register
PC	8	Program Counter; contains the address of the next instruction to be fetched for execution
SW	9	Status word; contains a variety of information, including a Condition Code(CC)

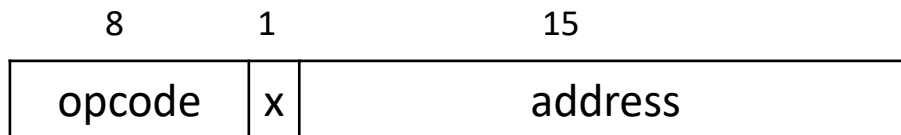
SIC Machine Architecture

- **Data Formats**

- Integers stored as 24 bit numbers
- 2's complement representation for negative numbers
- Characters stored using their 8 bit ASCII Code
- No floating point hardware

- **Instruction Formats**

- 24 bit format



- Flag bit-x indicates indexed addressing mode

Addressing Modes

- Two addressing modes –direct and indexed

Mode	Indication	Target Address Calculation
Direct	$x = 0$	TA = address
Indexed	$x = 1$	TA = address + (X)

- **Direct addressing mode**

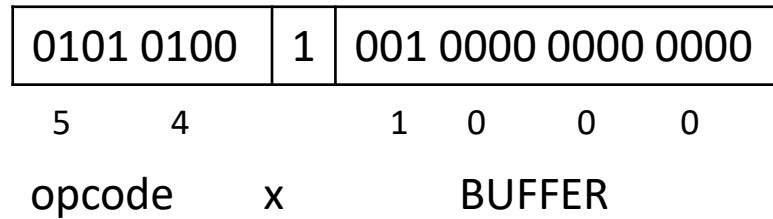
- LDA TEN

0000 0000	0	001 0000 0000 0000
opcode	x	TEN

Effective Address (EA) =1000

Addressing Modes

- **Indexed addressing mode**
 - STCH BUFFER, X



Effective Address (EA) = 1000 + [X]

Instruction Set

- Instructions that load and store registers
 - LDA, LDX, STA, STX
- Integer arithmetic operations
 - ADD, SUB, MUL, DIV
 - involve register A and a word in memory, result in the register
- Comparison, COMP
 - compares the value in register A with a word in memory and sets the condition code, CC.
- Conditional jump instructions JLT, JEQ, JGT
 - Test the setting of CC
- Subroutine linkage instructions, JSUB and RSUB
 - Return address placed in register L

Input and Output

- Performed by transferring 1 byte at a time to or from the rightmost eight bits of register A
- Each device is assigned a unique 8 bit code.
- Three I/O instructions
- Test Device (TD) checks whether the device is ready to send or receive data
 - Condition code is set to indicate the result of this test
 - < device is ready to send or receive, = device not ready
- Program needing to transfer data must wait until the device is ready, then execute a Read Data (RD) or Write Data (WD).
- Sequence repeated for each byte of data to be read or written

SIC/XE Machine Architecture

- Memory
 - 2^{20} bytes in the computer memory
- Registers
- Additional registers provided by SIC/XE

Mnemonic	Number	Special use
B	3	Base register; used for addressing
S	4	General working register
T	5	General working register
F	6	Floating-point accumulator (48bits)

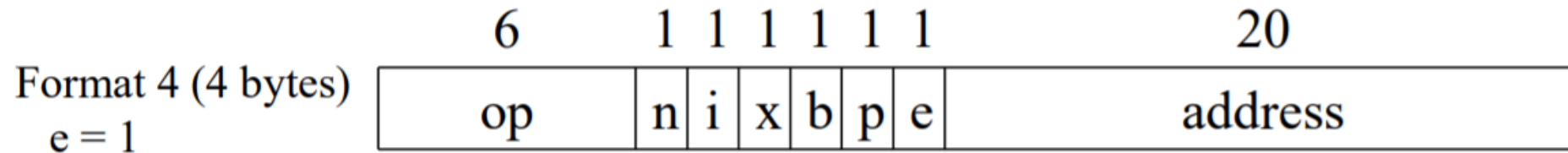
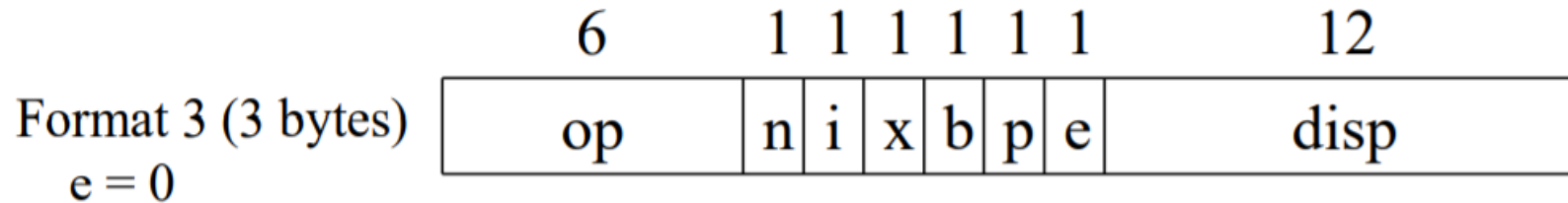
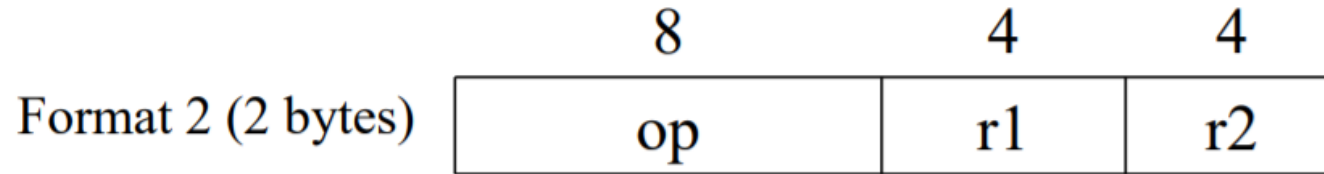
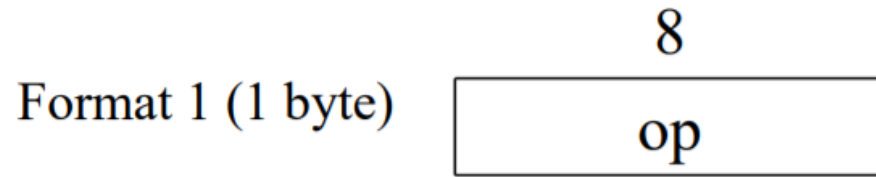
SIC/XE Machine Architecture

- Data Formats
 - Same data format as the standard version
 - 48 bit floating-point data type

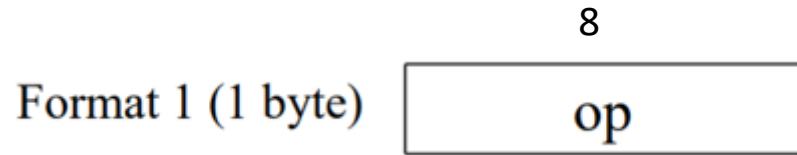


- Fraction is a value between 0 and 1, the assumed decimal point is immediately before the high order bit
- For normalized floating point numbers, the high order bit of the fraction must be 1
- Exponent is an unsigned number between 0 and 2047
- If the *exponent is e* and the *fraction has a value f*, the absolute value of the number is represented by
$$f * 2^{(e-1024)}$$
- s-represents the sign of the floating point number, s=0 positive, s=1 negative
- To represent zero, set all the bits to zero

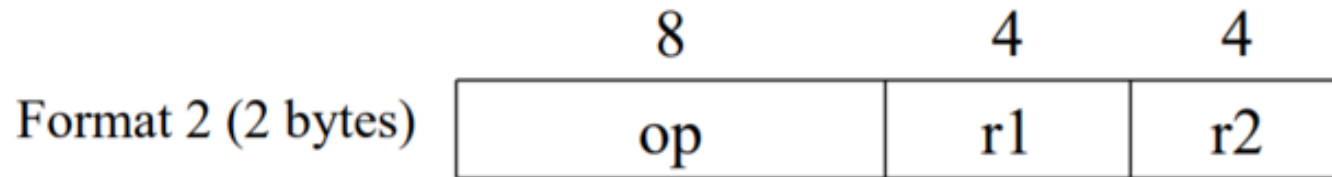
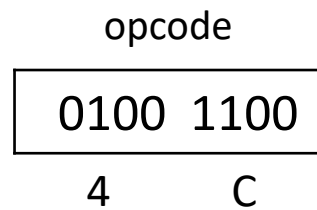
Instruction formats



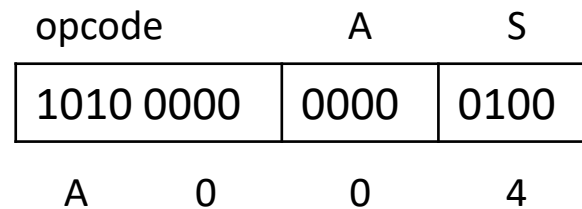
Instruction Format



Eg: RSUB

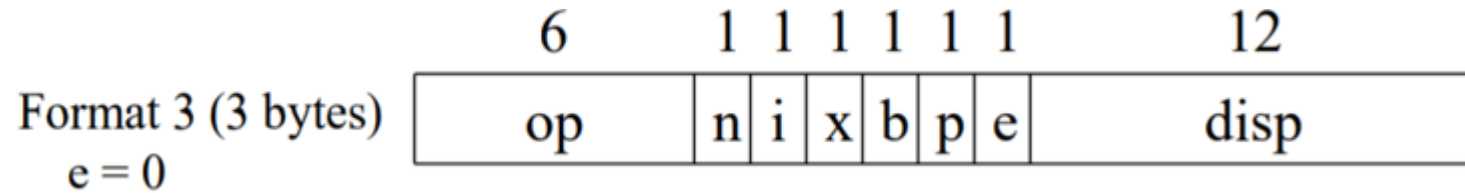


Eg: COMPR A,S

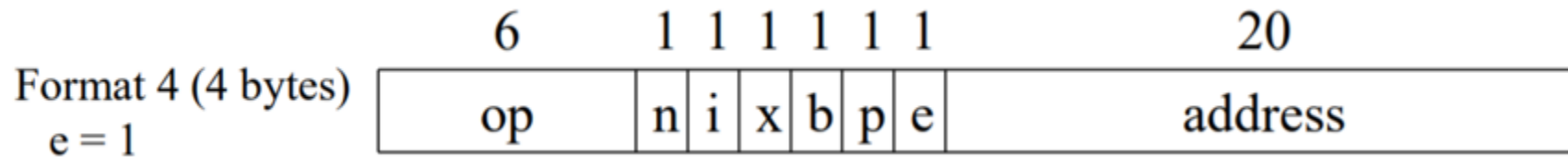
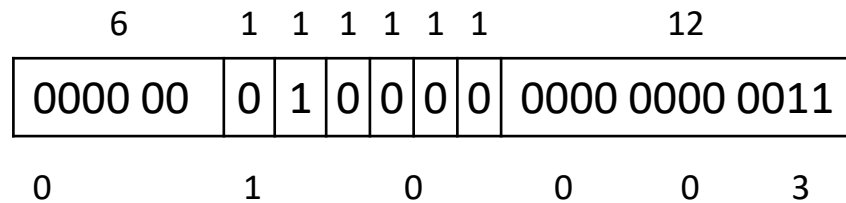


Formats 1 and 2 instructions do not reference memory at all

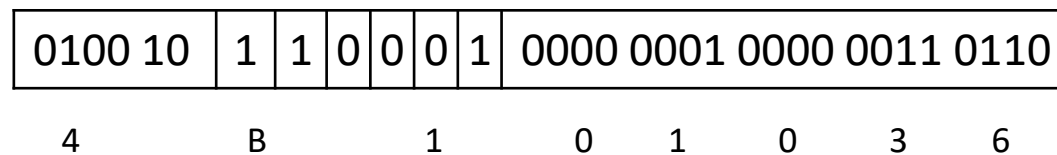
Instruction Format



- LDA #3 (Load 3 to Accumulator A)



- +JSUB RDREC



Addressing Modes

- **Format 3 Instruction**
- x – index bit
 - x is set to 1, X register value is added for target address calculation
- Bits b and p specify relative addressing.
 - b = 1 and p = 0: Relative to B.
 - b = 0 and p = 1: Relative to PC.
 - b = 0 and p = 0: Direct address.
- Bits i and n specify extended addressing modes.
 - i = 1 and n = 0: Immediate mode.
 - i = 0 and n = 1: Indirect mode.
 - i = n: Simple addressing.
- Bit e is always 0 for 3-byte instructions.

Addressing Modes

Mode	Indication	Target Address Calculation
Base relative	b=1, p=0	$TA = (B) + \text{disp}$ ($0 \leq \text{disp} \leq 4095$)
Program counter relative	b=0, p=1	$TA = (PC) + \text{disp}$ ($-2048 \leq \text{disp} \leq 2047$)
Immediate addressing	i=1, n=0	TA: TA is used as the operand value, no memory reference
Indirect addressing	i=0, n=1	((TA)): The word at the TA is fetched. Value of TA is taken as the address of the operand value
Simple addressing	i=0, n=0 or i=1, n=1	(TA): TA is taken as the address of the operand value
Base relative indexed addressing	b=1, p=0, x=1	$TA = (B) + (X) + \text{disp}$ ($0 \leq \text{disp} \leq 4095$)
Program counter relative indexed addressing	b=0, p=1, x=1	$TA = (PC) + (X) + \text{disp}$ ($-2048 \leq \text{disp} \leq 2047$)

Addressing Modes

- For base relative addressing, the displacement field *disp* in format 3 is interpreted as a 12 bit unsigned integer
- For program counter relative mode is used, *disp* is a 12-bit signed integer with negative numbers represented in 2's complement form
- b and p - both set to 0, *disp* field from format 3 instruction is taken to be the target address.
- Format 4 Instruction
 - bits b and p are normally set to 0, 20 bit address is the target address
 - Bit e is always 1

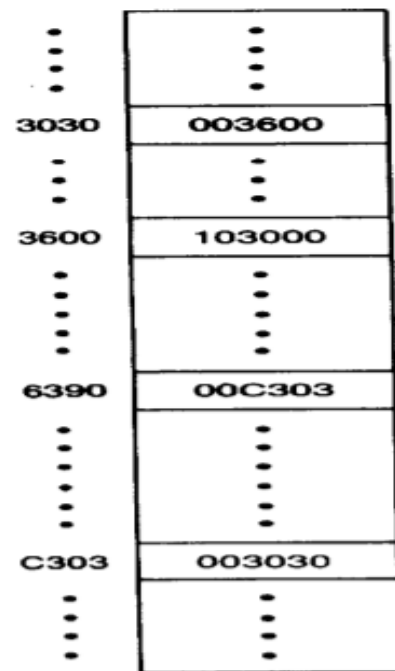
Addressing Modes Example

.	.	(B)=006000
.	.	(PC)=003000
.	.	(X)=000090
3030	003600	
.	.	
.	.	
.	.	
3600	103000	
.	.	
.	.	
.	.	
6390	00C303	
.	.	
.	.	
.	.	
C303	003030	
.	.	
.	.	
.	.	

Machine instruction										
Hex	Binary									
	op	n	l	x	b	p	e	disp/address		
032600	000000	1	1	0	0	1	0	0110	0000	0000
03C300	000000	1	1	1	1	0	0	0011	0000	0000
022030	000000	1	0	0	0	1	0	0000	0011	0000
010030	000000	0	1	0	0	0	0	0000	0011	0000
003600	000000	0	0	0	0	1	1	0110	0000	0000
0310C303	000000	1	1	0	0	0	1	0000	1100	0011 0000 0011

(b)

Examples of SIC/XE instructions and addressing modes.



(a)

(B) = 006000
(PC) = 003000
(X) = 000090

Machine instruction											Target address	Value loaded into register A
Hex	Binary											
	op	n	i	x	b	p	e	disp/address				
032600	000000	1	1	0	0	1	0	0110	0000	0000	3600	103000
03C300	000000	1	1	1	1	0	0	0011	0000	0000	6390	00C303
022030	000000	1	0	0	0	1	0	0000	0011	0000	3030	103000
010030	000000	0	1	0	0	0	0	0000	0011	0000	30	000030
003600	000000	0	0	0	0	1	1	0110	0000	0000	3600	103000
0310C303	000000	1	1	0	0	0	1	0000	1100	0011 0000 0011	C303	003030

(b)

Examples of SIC/XE instructions and addressing modes.

Instruction Set

- Instructions to load and store the new registers: LDB, STB etc.
- Floating-point arithmetic operations:
 - ADDF, SUBF, MULF, DIVF
- Register move instruction : RMO
- Register-to-register arithmetic operations
 - ADDR, SUBR, MULR, DIVR
- Supervisor call instruction : SVC
 - Executing this instruction generates an interrupt that can be used for communicating with the operating system
- Integer Arithmetic operations: ADD, SUB, MUL, DIV
- Compare instruction, COMP
- Subroutine linkage instructions: JSUB, RSUB
- Branching instructions (Conditional Jump Instruction): JLT, JEQ, JGT

Input and Output

- There are I/O channels that can be used to perform input and output while the CPU is executing other instructions
- Allows overlap of computing and I/O, resulting in more efficient system operation
- The instructions SIO, TIO, and HIO are used to start, test and halt the operation of I/O channels

Assembler Directives

- Assembler directives are instructions that direct the assembler to do something. The assembler directives in SIC are:

1. **START**

- Indicates the start of program
- Used to define program name and starting address
- Eg: COPY START 1000
 - This means the program name is COPY and the starting address is 1000

2. **END**

- Used to indicate the end of program.
- Optionally indicates first executable instruction.
- Eg: END ALPHA
 - This means the program ends here. And the first executable instruction is ALPHA

Assembler Directives

3. RESW

- Used to reserve specified words for a data area.
- Eg: ALPHA RESW 4
 - This is to reserve 4 words. (4 words means $4 * 3 = 12$ bytes)

4. RESB

- Used to reserve specified bytes for a data area.
- Eg: A RESB 5
 - This is to reserve 5 bytes

5. WORD

- Used to generate one word integer constant
- Eg: B WORD 6
 - This uses one word to store the integer constant 6 and the name B is assigned to the first location

6. BYTE

- Generate character constant using required number of bytes
- Eg: ALPHA BYTE C 'HAI'
 - This generates number of bytes needed to store HAI

Sample program SIC

- Data Movement Operations

	LDA	FIVE	LOAD CONSTANT 5 INTO REGISTER A
	STA	ALPHA	STORE IN ALPHA
	LDCH	CHARZ	LOAD CHARACTER 'Z' INTO REGISTER A
	STCH	C1	STORE IN CHARACTER VARIABLE C1
	...		
ALPHA	RESW	1	ONE WORD VARIABLE
FIVE	WORD	5	ONE WORD CONSTANT
CHARZ	BYTE	C'Z'	ONE BYTE CONSTANT
C1	REWB	1	ONE BYTE VARIABLE

Sample program SIC/XE

- Data Movement Operations

	LDA	#5	LOAD VALUE 5 INTO REGISTER A
	STA	ALPHA	STORE IN ALPHA
	LDA	#90	LOAD ASCII CODE FOR 'Z' INTO REGISTER A
	STCH	C1	STORE IN CHARACTER VARIABLE C1
	...		
ALPHA	RESW	1	ONE WORD VARIABLE
C1	REWB	1	ONE BYTE VARIABLE

Sample Arithmetic Operations- SIC

LDA	ALPHA	LOAD ALPHA INTO REGISTER A
ADD	INCR	ADD THE VALUE OF INCR
SUB	ONE	SUBTRACT 1
STA	BETA	STORE IN BETA
LDA	GAMMA	LOAD GAMMA INTO REGISTER A
ADD	INCR	ADD THE VALUE OF INCR
SUB	ONE	SUBTRACT 1
STA	DELTA	STORE IN DELTA

.....

ONE	WORD	1	ONE WORD CONSTANT
ALPHA	RESW	1	ONE WORD VARIABLE
BETA	RESW	1	ONE WORD VARIABLE
GAMMA	RESW	1	ONE WORD VARIABLE
DELTA	RESW	1	ONE WORD VARIABLE
INCR	RESW	1	

Sample Arithmetic Operations- SIC/XE

LDS	INCR	LOAD INCR INTO REGISTER S
LDA	ALPHA	LOAD ALPHA INTO REGISTER A
ADDR	S,A	ADD THE VALUE OF INCR
SUB	#1	SUBTRACT 1
STA	BETA	STORE IN BETA
LDA	GAMMA	LOAD GAMMA INTO REGISTER A
ADDR	S,A	ADD THE VALUE OF INCR
SUB	#1	SUBTRACT 1
STA	DELTA	STORE IN DELTA

.....

ALPHA	RESW	1	ONE WORD VARIABLE
BETA	RESW	1	
GAMMA	RESW	1	
DELTA	RESW	1	
INCR	RESW	1	

Sample Indexing and Looping- SIC

	LDX	ZERO	INITIALIZE INDEX REGISTER TO 0
MOVECH	LDCH	STR1,X	LOAD CHARACTER FROM STR1 INTO REG A
	STCH	STR2,X	STORE CHARACTER INTO STR2
	TIX	ELEVEN	ADD 1 TO INDEX, COMPARE RESULT TO 11
	JLT	MOVECH	LOOP IF INDEX IS LESS THAN 11
	.		
	.		
	.		
STR1	BYTE	C'TEST STRING'	11-BYTE STRING CONSTANT
STR2	RESB	11	11-BYTE VARIABLE
.			ONE-WORD CONSTANTS
ZERO	WORD	0	
ELEVEN	WORD	11	

Sample Indexing and Looping- SIC/XE

	LDT	#11	INITIALIZE REGISTER T TO 11
	LDX	#0	INITIALIZE INDEX REGISTER TO 0
MOVECH	LDCH	STR1,X	LOAD CHARACTER FROM STR1 INTO REG A
	STCH	STR2,X	STORE CHARACTER INTO STR2
	TIXR	T	ADD 1 TO INDEX, COMPARE RESULT TO 11
	JLT	MOVECH	LOOP IF INDEX IS LESS THAN 11
	.		
	.		
	.		
STR1	BYTE	C'TEST STRING'	11-BYTE STRING CONSTANT
STR2	RESB	11	11-BYTE VARIABLE

Sample Indexing and Looping- SIC

	LDA	ZERO	INITIALIZE INDEX VALUE TO 0
	STA	INDEX	
ADDLP	LDX	INDEX	LOAD INDEX VALUE INTO REGISTER X
	LDA	ALPHA, X	LOAD WORD FROM ALPHA INTO REGISTER A
	ADD	BETA, X	ADD WORD FROM BETA
	STA	GAMMA, X	STORE THE RESULT IN A WORD IN GAMMA
	LDA	INDEX	ADD 3 TO INDEX VALUE
	ADD	THREE	
	STA	INDEX	
	COMP	K300	COMPARE NEW INDEX VALUE TO 300
	JLT	ADDLP	LOOP IF INDEX IS LESS THAN 300
	.		
	.		
	.		
INDEX	RESW	1	ONE-WORD VARIABLE FOR INDEX VALUE
.			ARRAY VARIABLES--100 WORDS EACH
ALPHA	RESW	100	
BETA	RESW	100	
GAMMA	RESW	100	
.			ONE-WORD CONSTANTS
ZERO	WORD	0	
K300	WORD	300	
THREE	WORD	3	

Indexing and Looping- SIC/XE

	LDS	#3	INITIALIZE REGISTER S TO 3
	LDT	#300	INITIALIZE REGISTER T TO 300
	LDX	#0	INITIALIZE INDEX REGISTER TO 0
ADDLP	LDA	ALPHA,X	LOAD WORD FROM ALPHA INTO REGISTER A
	ADD	BETA,X	ADD WORD FROM BETA
	STA	GAMMA,X	STORE THE RESULT IN A WORD IN GAMMA
	ADDR	S,X	ADD 3 TO INDEX VALUE
	COMPR	X,T	COMPARE NEW INDEX VALUE TO 300
	JLT	ADDLP	LOOP IF INDEX VALUE IS LESS THAN 300
	.		
	.		
	.		
.			ARRAY VARIABLES--100 WORDS EACH
ALPHA	RESW	100	
BETA	RESW	100	
GAMMA	RESW	100	

Sample Input-Output Operations-SIC

INLOOP	TD	INDEV	TEST INPUT DEVICE
	JEQ	INLOOP	LOOP UNTIL DEVICE IS READY
	RD	INDEV	READ ONE BYTE INTO REGISTER A
	STCH	DATA	STORE BYTE THAT WAS READ
	.		
	.		
	.		
OUTLP	TD	OUTDEV	TEST OUTPUT DEVICE
	JEQ	OUTLP	LOOP UNTIL DEVICE IS READY
	LDCH	DATA	LOAD DATA BYTE INTO REGISTER A
	WD	OUTDEV	WRITE ONE BYTE TO OUTPUT DEVICE
	.		
	.		
	.		
INDEV	BYTE	X'F1'	INPUT DEVICE NUMBER
OUTDEV	BYTE	X'05'	OUTPUT DEVICE NUMBER
DATA	RESB	1	ONE-BYTE VARIABLE

Sample subroutine call and record input operations- SIC

	JSUB	READ	CALL READ SUBROUTINE
	.		
	.		
	.		
.			SUBROUTINE TO READ 100-BYTE RECORD
READ	LDX	ZERO	INITIALIZE INDEX REGISTER TO 0
RLOOP	TD	INDEV	TEST INPUT DEVICE
	JEQ	RLOOP	LOOP IF DEVICE IS BUSY
	RD	INDEV	READ ONE BYTE INTO REGISTER A
	STCH	RECORD,X	STORE DATA BYTE INTO RECORD
	TIX	K100	ADD 1 TO INDEX AND COMPARE TO 100
	JLT	RLOOP	LOOP IF INDEX IS LESS THAN 100
	RSUB		EXIT FROM SUBROUTINE
	.		
	.		
	.		
INDEV	BYTE	X'F1'	INPUT DEVICE NUMBER
RECORD	RESB	100	100-BYTE BUFFER FOR INPUT RECORD
.			ONE-WORD CONSTANTS
ZERO	WORD	0	
K100	WORD	100	

Sample subroutine call and record input operations-SIC/XE

	JSUB	READ	CALL READ SUBROUTINE
	.		
	.		
	.		
	.		SUBROUTINE TO READ 100-BYTE RECORD
READ	LDX	#0	INITIALIZE INDEX REGISTER TO 0
	LDT	#100	INITIALIZE REGISTER T TO 100
RLOOP	TD	INDEV	TEST INPUT DEVICE
	JEQ	RLOOP	LOOP IF DEVICE IS BUSY
	RD	INDEV	READ ONE BYTE INTO REGISTER A
	STCH	RECORD,X	STORE DATA BYTE INTO RECORD
	TIXR	T	ADD 1 TO INDEX AND COMPARE TO 100
	JLT	RLOOP	LOOP IF INDEX IS LESS THAN 100
	RSUB		EXIT FROM SUBROUTINE
	.		
	.		
	.		
INDEV	BYTE	X'F1'	INPUT DEVICE NUMBER
RECORD	RESB	100	100-BYTE BUFFER FOR INPUT RECORD