

The background of the slide is an abstract composition of various shades of blue, ranging from deep navy to light sky blue. It features a complex pattern of overlapping, semi-transparent geometric shapes, primarily triangles and polygons, which create a sense of depth and movement. A prominent white horizontal band runs across the middle of the image, serving as a backdrop for the title and author information.

Trees and Hierarchies in SQL

by Eduard Hildebrandt

**In most projects we have to deal with
some kind of trees or hierarchies!**

Think about...

- Categories
- Organisation
- Threads
- Folders
- Components
- ...

**How can we store and retrieve trees and hierarchies
efficiently from relational databases?**

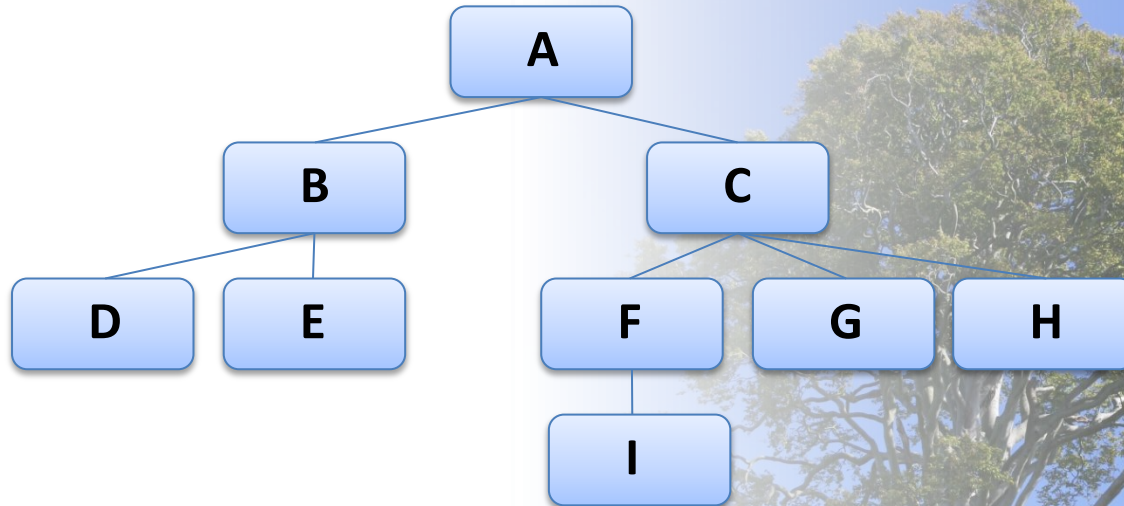
Agenda

- Adjacency List Model
- Closure Table Model
- Path Enumeration Model
- David Chandler Model
- Modified David Chandler Model
- Nested Set Model
- Hybrid Model
- Frequent Insertion in Nested Sets



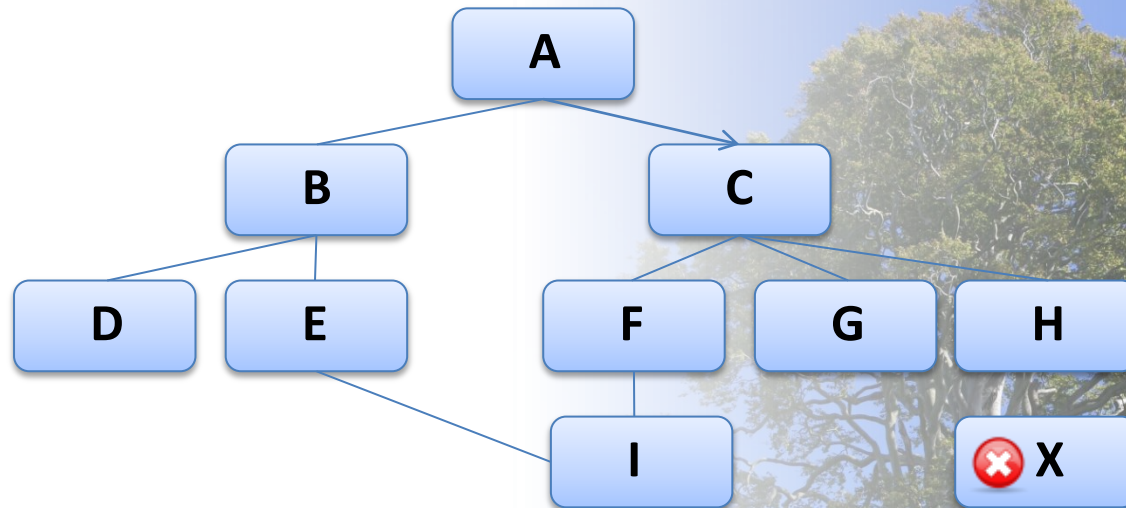
What is a „tree“?

A tree is a connected, undirected, acyclic grap.



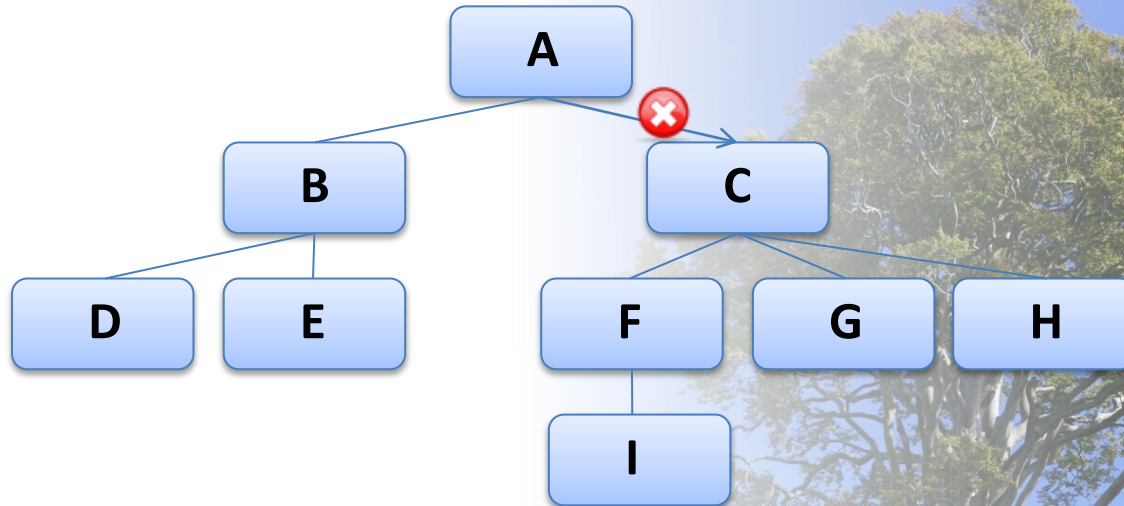
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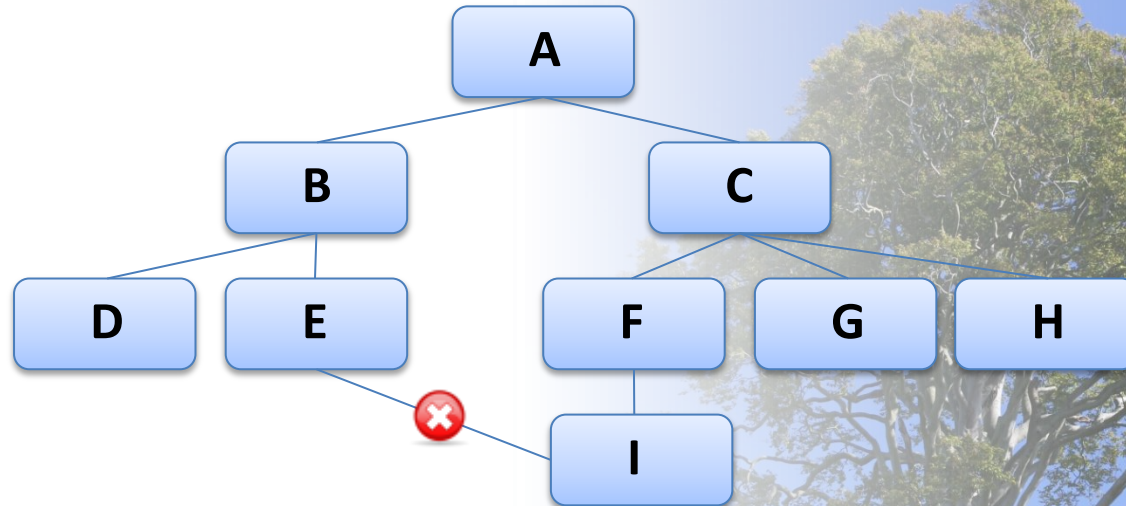
What is a „tree“?

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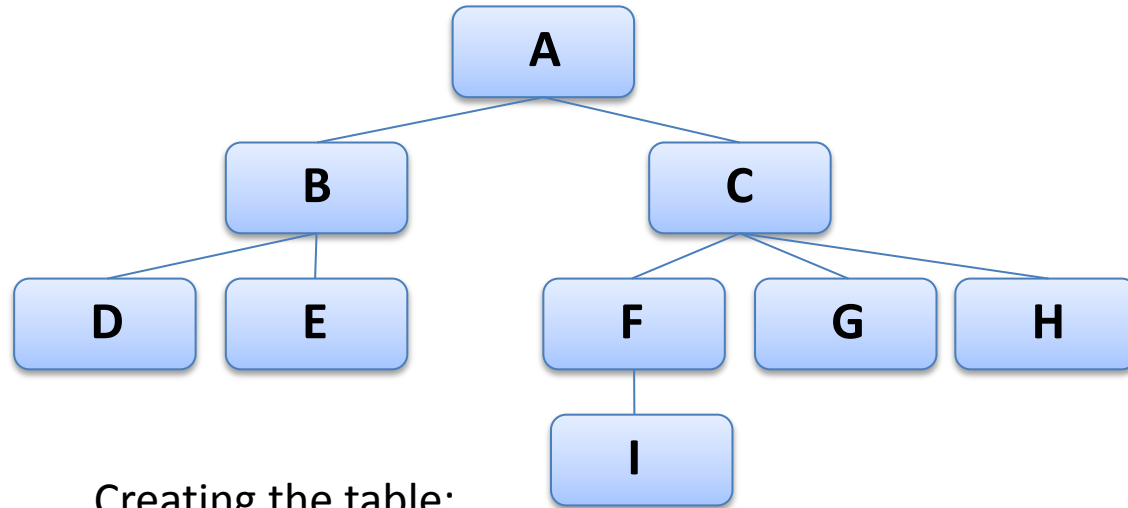


What is a „tree“?

A tree is a connected, undirected, **acyclic** grap.



Adjacency List Model



title	parent
A	null
B	A
C	A
D	B
E	B
F	C
G	C
H	C
I	F

Creating the table:

```
CREATE TABLE nodes
(title VARCHAR(255) NOT NULL PRIMARY KEY,
parent VARCHAR(255));
```

Finding the root node:

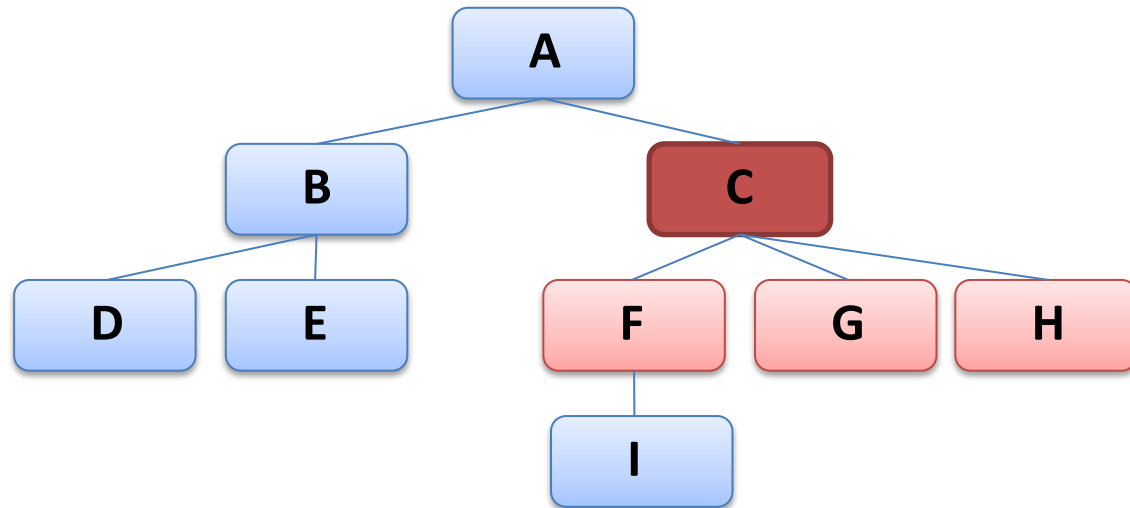
```
SELECT *
FROM nodes
WHERE parent IS NULL;
```

Finding the leaf nodes:

```
SELECT node1.*
FROM nodes AS node1
LEFT JOIN nodes AS node2 ON node1.title = node2.parent
WHERE node2.title = 'A';
```

Adjacency List Model

UPDATE anomalies



title	parent
A	null
B	A
C	A
D	B
E	B
F	C
G	C
H	C
I	F



```
UPDATE nodes
  SET title = 'X'
  WHERE title = 'C';
```



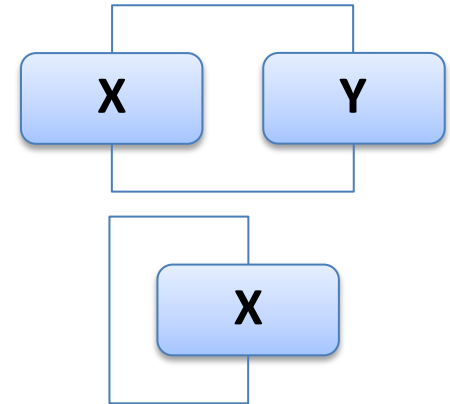
```
START TRANSACTION;
UPDATE nodes
  SET title = 'X'
  WHERE title = 'C';
UPDATE nodes
  SET parent = 'X'
  WHERE parent = 'C';
COMMIT;
```

Adjacency List Model

INSERT anomalies



```
INSERT INTO nodes (title, parent) VALUES  
('X', 'Y'),  
('Y', 'X');
```



```
INSERT INTO nodes (title, parent) VALUES  
('X', 'X');
```

Parent must exist:



```
ALTER TABLE nodes ADD CONSTRAINT parent_fk FOREIGN KEY (parent)  
REFERENCES nodes (id);
```

Node can not be it's own parent:



```
CHECK (title <> parent)
```

Prevent double connections:



```
UNIQUE (title, parent)
```

Connected graph (edges = nodes – 1):



```
CHECK ((SELECT COUNT(*) FROM nodes) - 1  
= SELECT COUNT(parent) FROM nodes);
```

One root only:



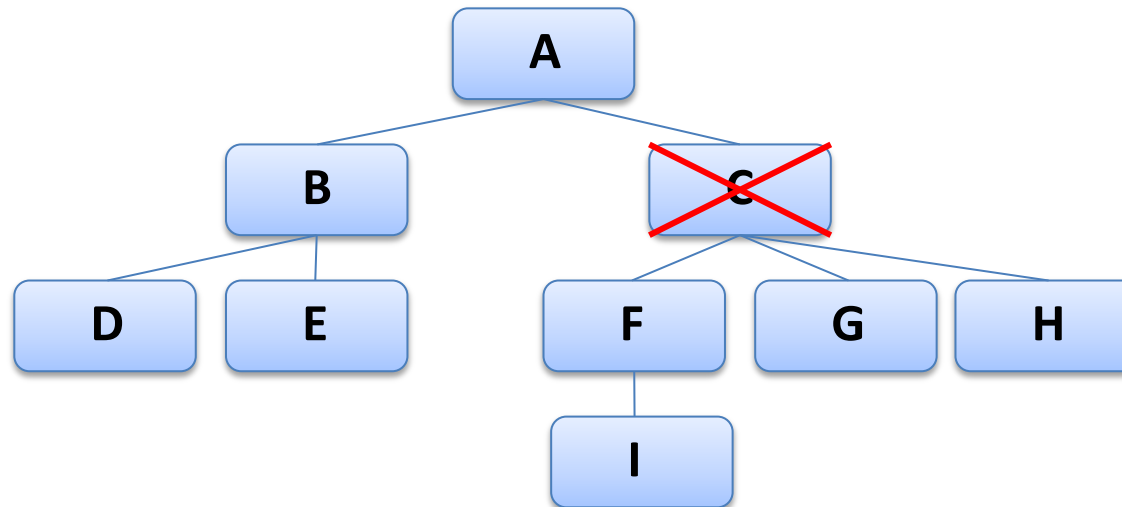
```
CHECK ((SELECT COUNT(*) FROM nodes WHERE parent IS NULL) = 1);
```

Adjacency List Model

DELETE anomalies



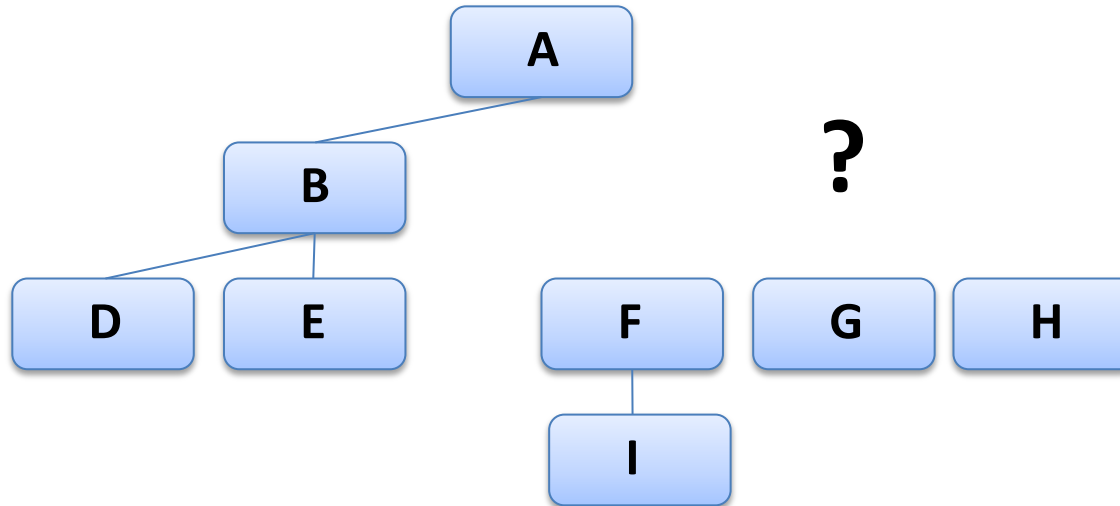
```
DELETE FROM nodes WHERE title = 'C' ;
```



Adjacency List Model

DELETE anomalies

What happens with nodes F, G, H and I?

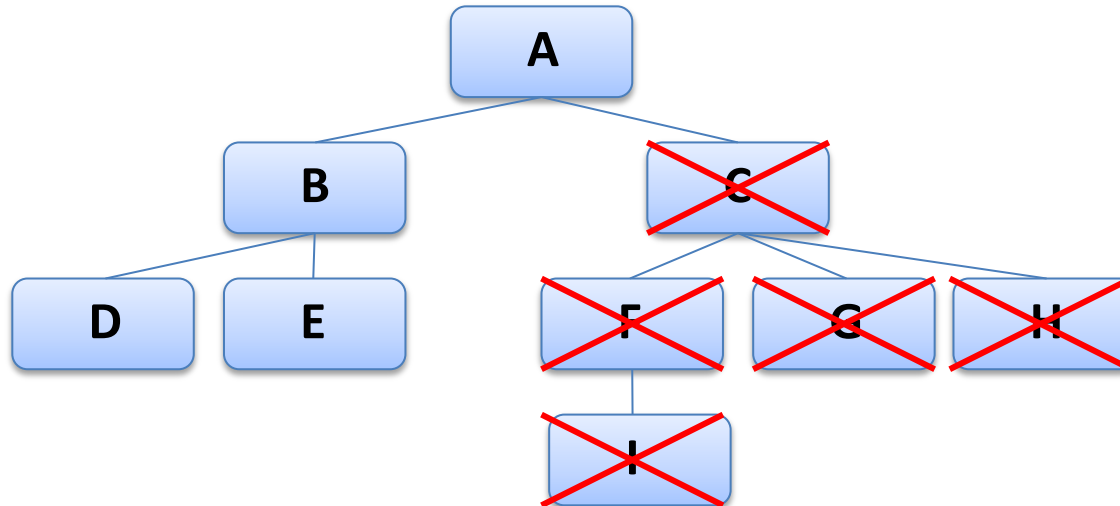


```
ALTER TABLE nodes ADD CONSTRAINT parent_fk FOREIGN KEY (parent)
REFERENCES nodes (id)
ON DELETE ...;
```

Adjacency List Model

DELETE anomalies

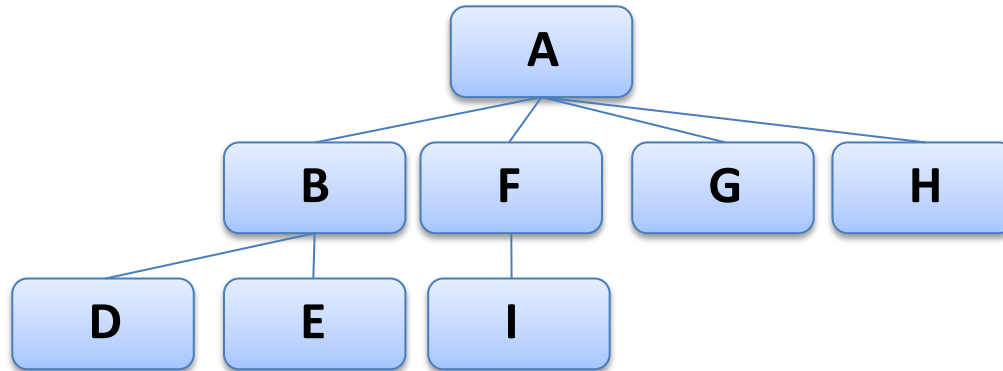
Solution 1: delete all children



Adjacency List Model

DELETE anomalies

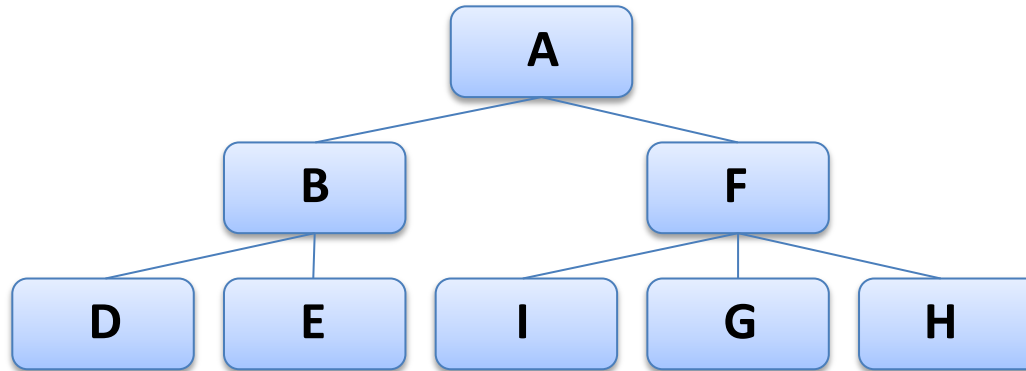
Solution 2: move children to the level of the parent



Adjacency List Model

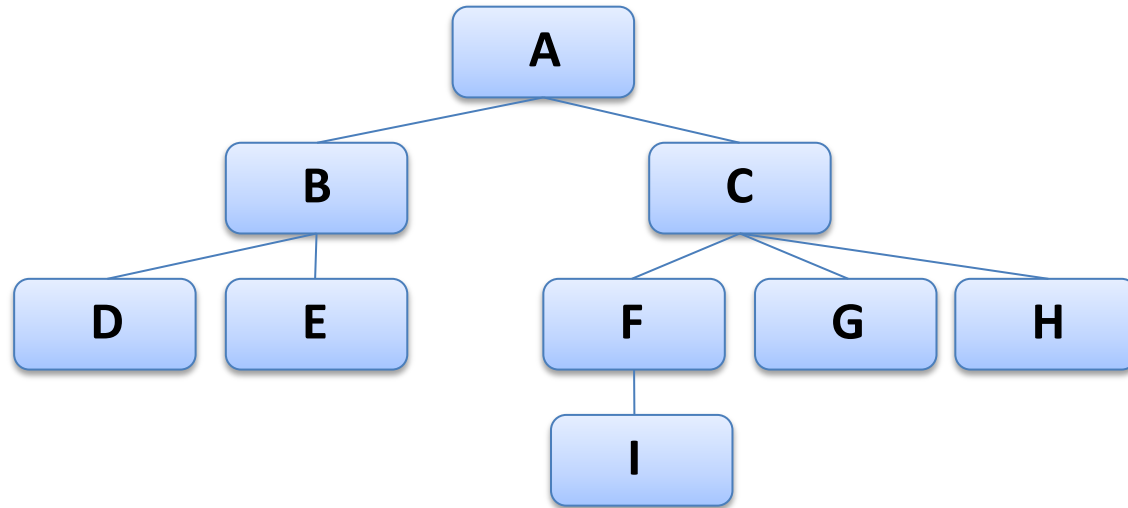
DELETE anomalies

Solution 3: replace parent with the first child



Adjacency List Model

with self -references



title	parent
A	A
B	A
C	A
D	B
E	B
F	C
G	C
H	C
I	F

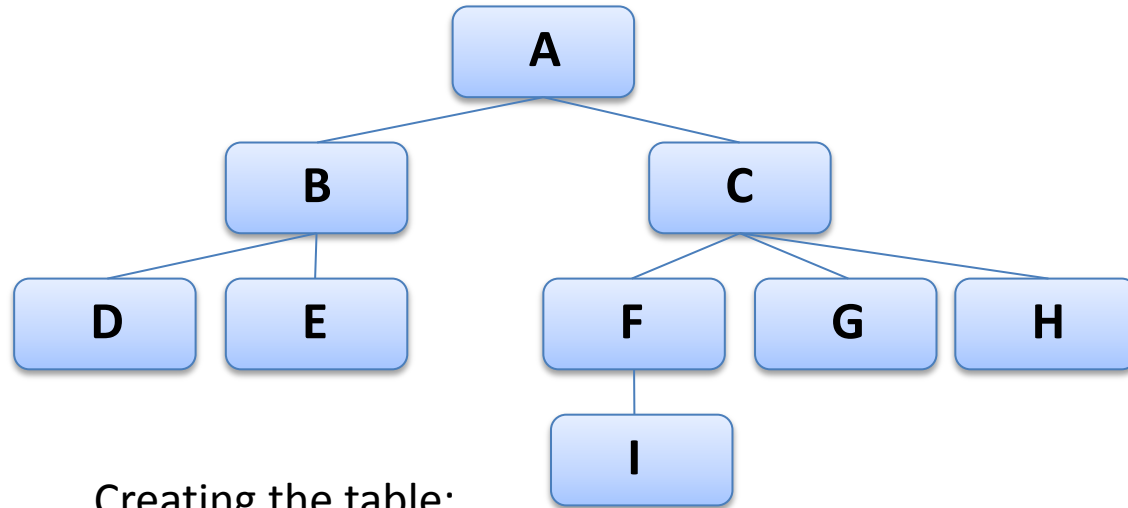
Creating the table:

```
CREATE TABLE nodes
(title VARCHAR(255) NOT NULL PRIMARY KEY,
parent VARCHAR(255) NOT NULL);
```



Avoids NULL in the parent column!

Path Enumeration Model

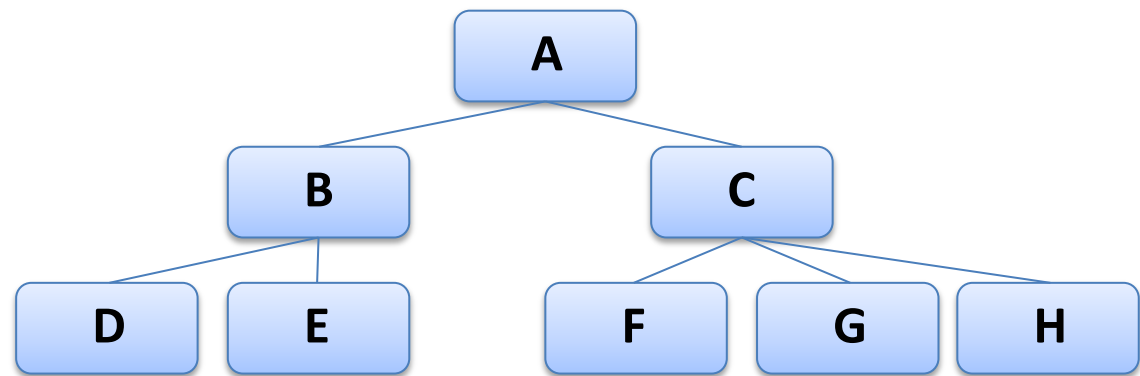


Creating the table:

```
CREATE TABLE nodes
(title VARCHAR(255) NOT NULL PRIMARY KEY
-- don't use a separator in primary key
CHECK (REPLACE (title, '/', '') = title,
path VARCHAR(1024) NOT NULL);
```

title	path
A	A
B	A/B
C	A/C
D	A/B/D
E	A/B/E
F	A/C/F
G	A/C/G
H	A/C/H
I	A/C/F/I

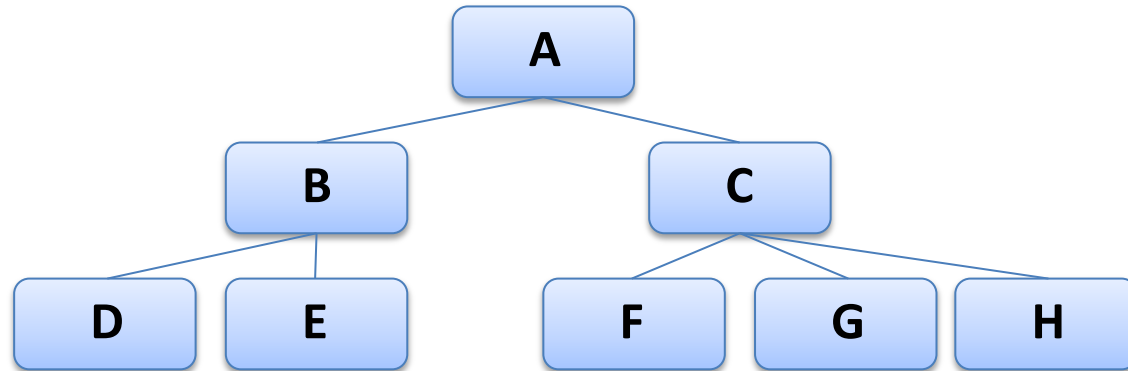
Closure Table Model



ID	parent
1	A
2	B
3	C
4	D
5	E
6	F
7	G
8	H

Ancestor	Descendant
1	2
1	4
1	5
2	4
2	5
1	3
1	6
1	7
1	8
3	6

David Chandler Model



ID	T	L1	L2	L3
1	A	1	0	0
2	B	1	1	0
3	C	1	2	0
4	D	1	1	1
5	E	1	1	2
6	F	1	2	1
7	G	1	2	2
8	H	1	2	3

Creating the table:

```
CREATE TABLE nodes
(id INTEGER NOT NULL PRIMARY KEY
 title VARCHAR(255) NOT NULL
 L1 INTEGER NOT NULL,
 L2 INTEGER NOT NULL,
 L3 INTEGER NOT NULL);
```

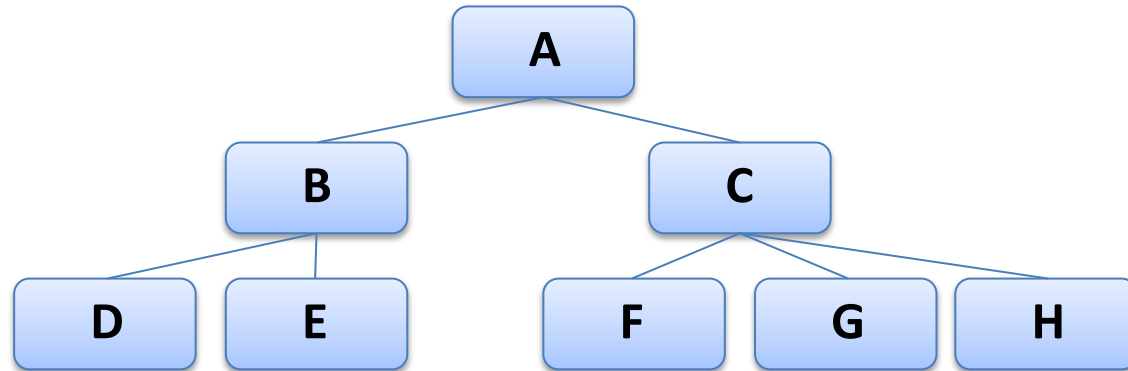


Algorithm is patented by David Chandler!



Hierarchy level is fixed!

Modified David Chandler Model



ID	T	L1	L2	L3
1	A	1	0	0
2	B	1	2	0
3	C	1	3	0
4	D	1	2	4
5	E	1	2	5
6	F	1	3	6
7	G	1	3	7
8	H	1	3	8

Creating the table:

```
CREATE TABLE nodes
(id INTEGER NOT NULL PRIMARY KEY
 title VARCHAR(255) NOT NULL
 L1 INTEGER NOT NULL,
 L2 INTEGER NOT NULL,
 L3 INTEGER NOT NULL);
```

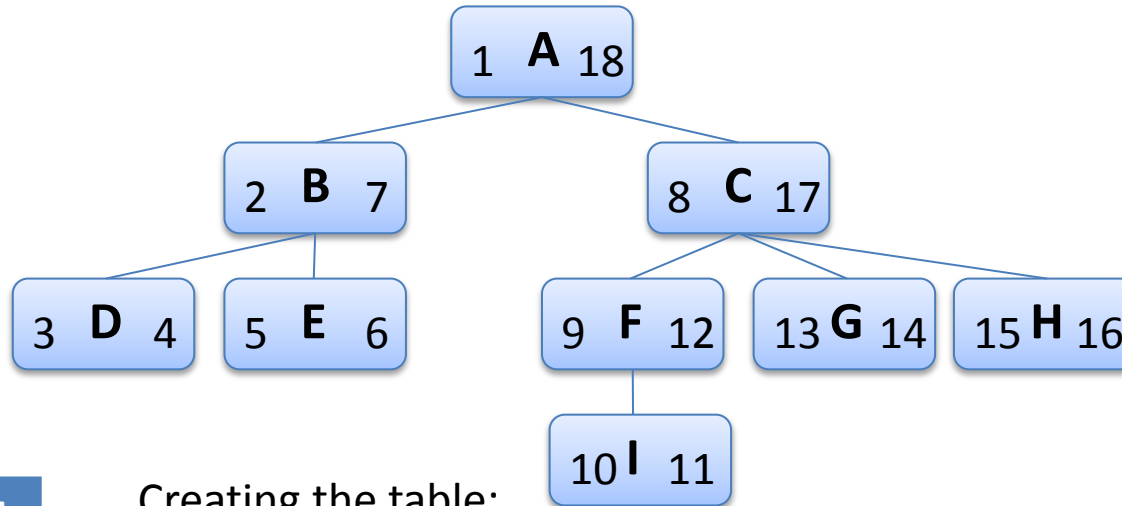


Not sure about implications of David Chandlers patent...



Hierarchy level is fixed!

Nested Set Model



title	lft	rgt
A	1	18
B	2	7
C	8	17
D	3	4
E	5	6
F	9	12
G	13	14
H	15	16
I	10	11

Creating the table:

```
CREATE TABLE nodes
  (title VARCHAR(255) NOT NULL PRIMARY KEY,
   lft  INTEGER NOT NULL,
   rgt  INTEGER NOT NULL);
```

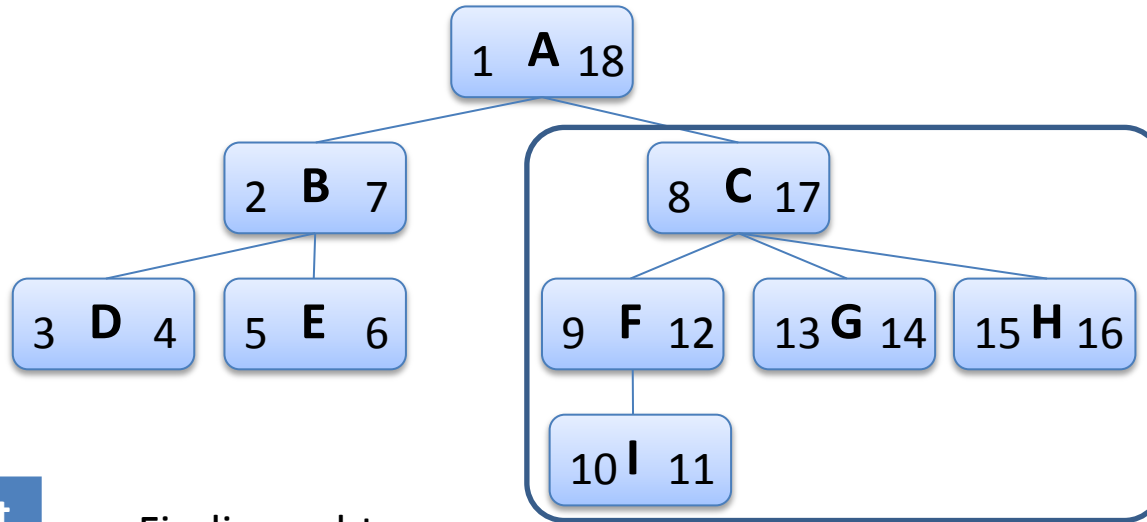
Finding the root node:

```
SELECT *
  FROM nodes
 WHERE lft = 1;
```

Finding the leaf nodes:

```
SELECT *
  FROM nodes
 WHERE lft = (MAX(rgt) - 1)
```

Nested Set Model



title	lft	rgt
A	1	18
B	2	7
C	8	17
D	3	4
E	5	6
F	9	12
G	13	14
H	15	16
I	10	11

Finding subtrees:

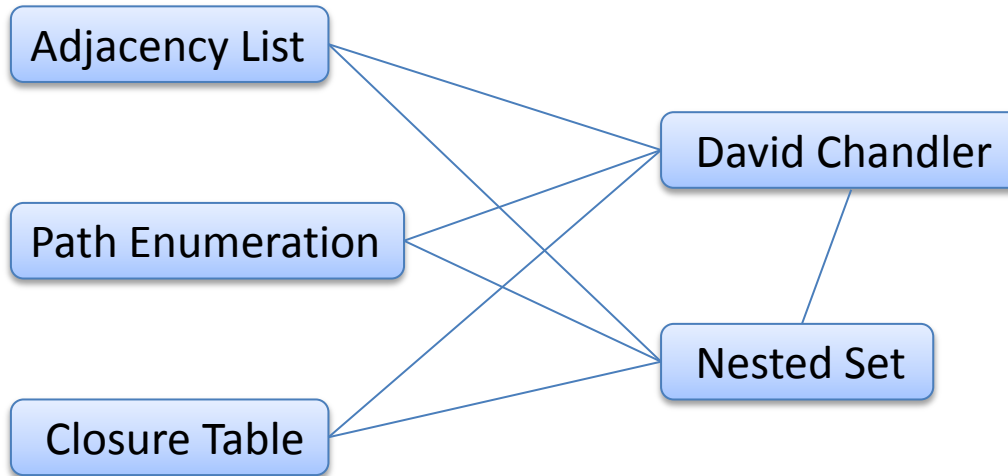
```
SELECT * FROM nodes
  WHERE lft BETWEEN lft AND rgt
 ORDER BY lft ASC;
```

Finding path to a node:

```
SELECT * FROM nodes
  WHERE lft < ? AND rgt > ?
 ORDER BY lft ASC;
```

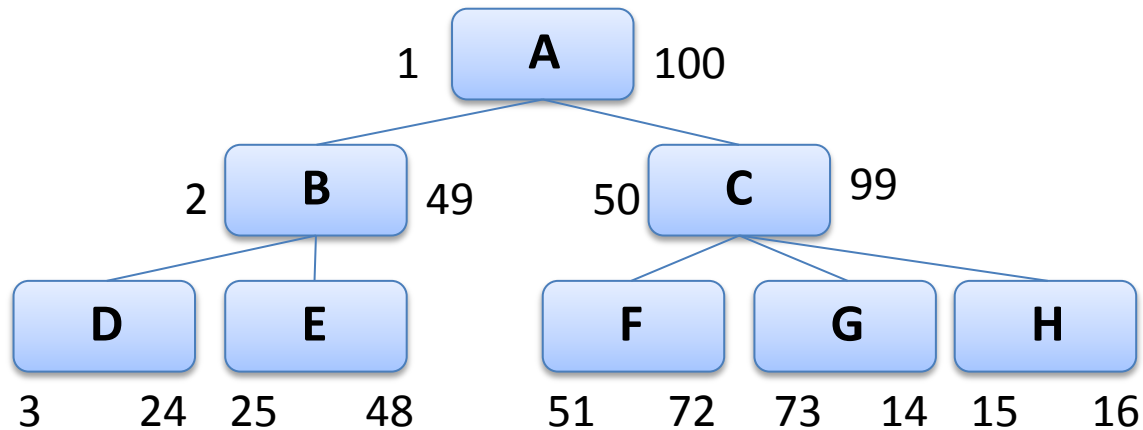
Hybrid Models

Combining models



Nested Set Model

Problem: Tree must be reorganized on every insertion (table lock).



Possible solution: Avoid overhead with larger spreads and bigger gaps

Nested Set Model

What datatype should I use for left and right values?



INTEGER



FLOAT / REAL / DOUBLE

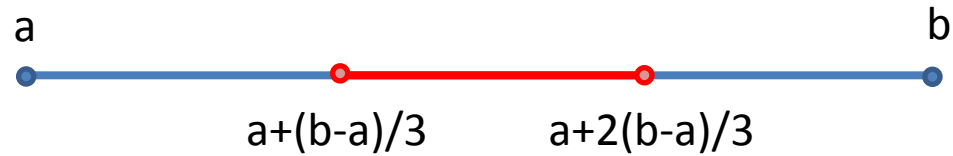
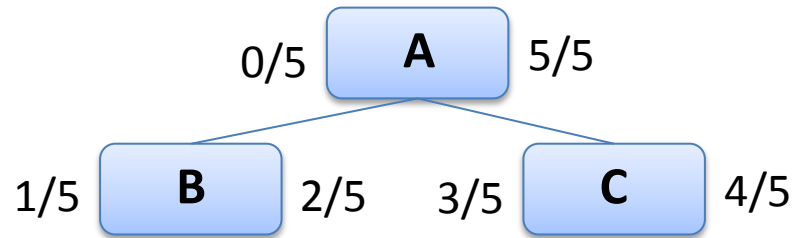


NUMERIC / DECIMAL



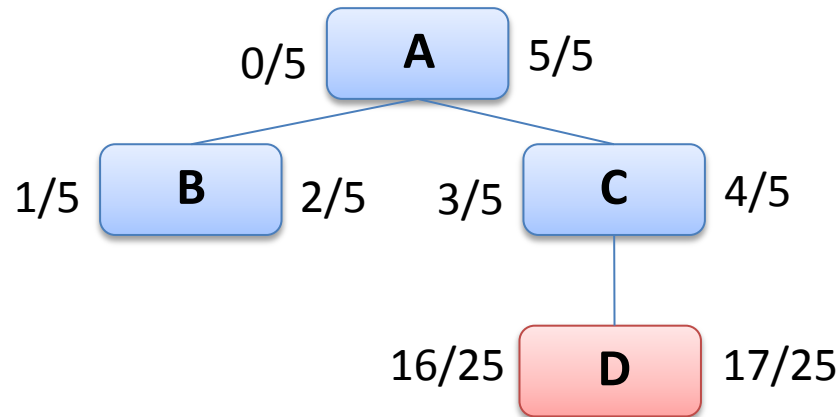
Nested Set Model

with rational numbers



Nested Set Model

with rational numbers



T	ln	ld	rn	rd
A	0	5	5	5
B	1	5	2	5
C	3	5	4	5
D	16	25	17	25

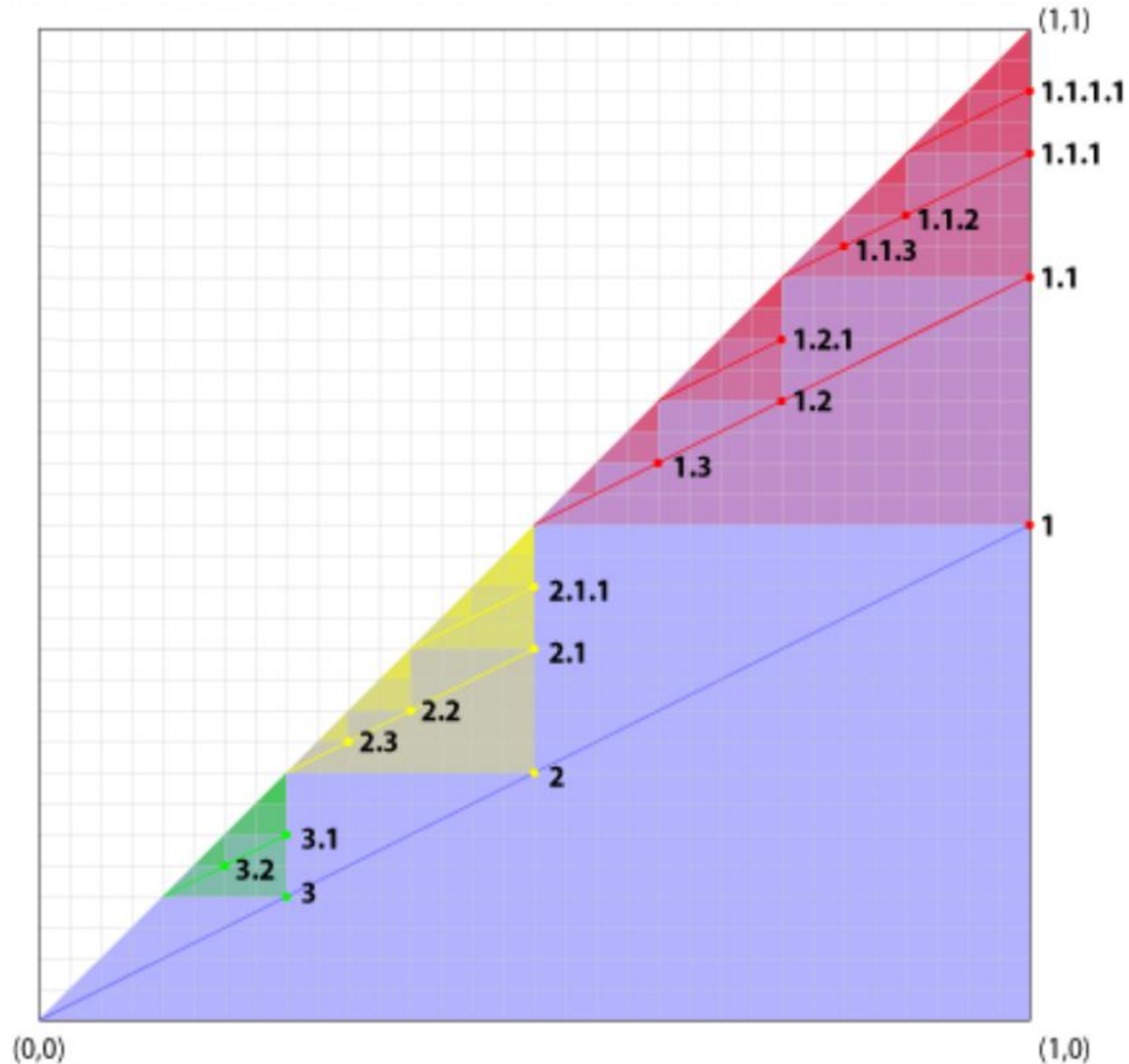


Adding an additional child node is always possible.
No need to reorganize the table!

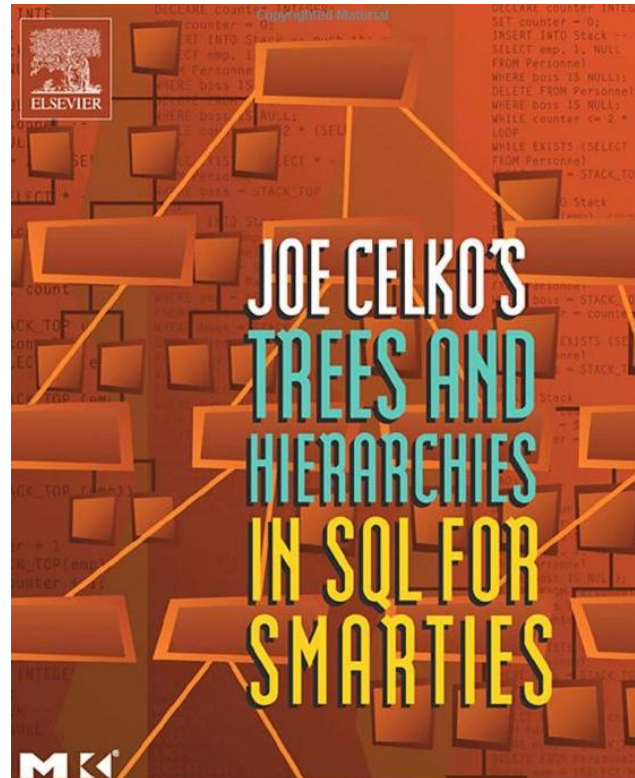
Nested Set Model

Binary Fraction

node	X	Y
1	1	1/2
1.1	1	3/4
1.1.1	1	7/8
1.1.1.1	1	15/16
1.1.2	7/8	13/16
1.1.3	13/16	25/32
1.2	3/4	5/8
1.2.1	3/4	11/16
1.3	5/8	9/16



Literature



Titel: Trees and Hierarchies in SQL


Autor: Joe Celko

Verlag: Morgan Kaufmann

ISBN: 1558609202

Q&A





THANK YOU!