#### Hilbert Transform Pairs of Wavelets

September 19, 2007



#### Outline

#### The Challenge

Hilbert Transform Pairs of Wavelets

#### The Reason

The Dual-Tree Complex Wavelet Transform

#### The Dissapointment

Background

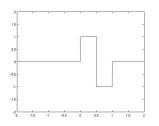
Characterization by means of Frequency Responses

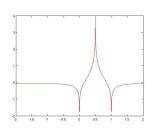
#### A Workaround

Approximate Hilbert Transform Pairs The Design Problem



#### Hilbert Transform Pairs





$$\mathcal{H}f(x) = \lim_{\varepsilon \to 0^+} \frac{1}{\pi} \int_{\{|y| > \varepsilon\}} \frac{f(x-y)}{y} \, dy.$$

equivalently, if  $\widehat{f}(\xi)$  is defined,

$$\mathcal{F}\{\mathcal{H}f\}(\xi) = -i\operatorname{sign}(\xi)\widehat{f}(\xi).$$



#### Hilber Pairs of Wavelets

Given two mother wavelets  $\psi_1, \psi_2 \colon \mathbb{R} \to \mathbb{R}$ , find necessary and sufficient conditions so that  $\psi_2 = \mathcal{H}\psi_1$ .

- ▶ If this is possible, find an example.
- ▶ If not, characterize "how close we can get" (and give an example).



We are interested in approximants with the following properties:

- ▶ Fast algorithms
- ▶ Perfect Reconstruction
- ▶ Modeling geometric features
- ▶ Shift Invariance



We are interested in approximants with the following properties:

- ▶ Fast algorithms
   ▶ Perfect Reconstruction

  Orthogonal real-valued wavelets (MRA)
- ► Modeling geometric features
- ► Shift Invariance



We are interested in approximants with the following properties:

- Fast algorithms
   Perfect Reconstruction

  Orthogonal real-valued wavelets (MRA)
- ► Modeling geometric features Shift Invariance  $\begin{cases} \text{Complex valued wavelets} \\ \psi = u + i\mathcal{H}u. \end{cases}$



We are interested in approximants with the following properties:

- ► Fast algorithms
- Perfect ReconstructionModeling geometric features
- ► Shift Invariance



$$\begin{array}{cccc}
\cdots V_2 \subset V_1 \subset V_0 \subset V_{-1} \subset \cdots \\
\phi \in V_0
\end{array}$$

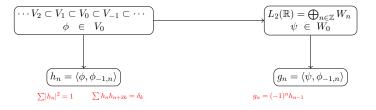


$$\begin{array}{cccc}
 & & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 & & \\
 &$$

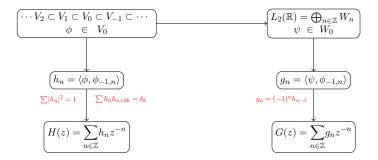




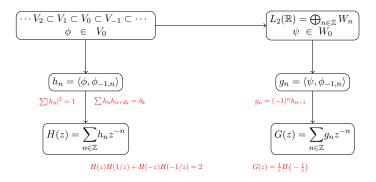




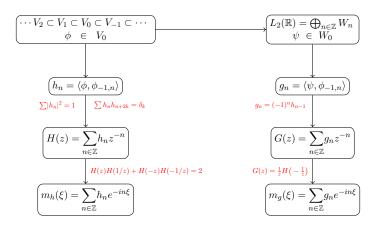




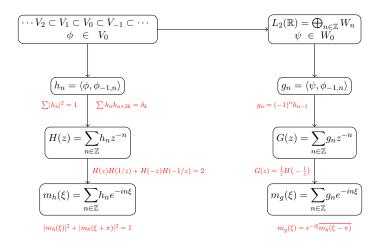




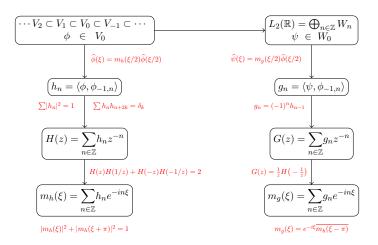














#### The Phase Condition

Assume there exists a  $2\pi$ -periodic function  $\theta \colon \mathbb{R} \to \mathbb{R}$  so that  $m_{h_2}(\xi) = e^{-i\theta(\xi)} m_{h_1}(\xi)$ .

$$\widehat{\phi}_2(\xi) = e^{-i\sum_{k=1}^{\infty} \theta(\xi/2^k)} \widehat{\phi}_1(\xi).$$

• 
$$m_{g_2}(\xi) = e^{i\theta(\xi-\pi)} m_{g_1}(\xi)$$
.

$$\hat{\psi}_2(\xi) = e^{i \{ \theta(\xi/2 - \pi) - \sum_{k=1}^{\infty} \theta(\xi/2^{k+1}) \}} \hat{\psi}_1(\xi).$$

Therefore,

$$-i\operatorname{sign}(\xi) = e^{i\left\{\theta(\xi/2-\pi) - \sum_{k=2}^{\infty} \theta(\xi/2^k)\right\}} \text{ implies } \psi_2 = \mathcal{H}\psi_1.$$

$$\theta(\xi/2+\pi) - \sum_{k=1}^{\infty} \theta(\xi/2^{k+1}) = -\frac{\pi}{2}\operatorname{sign}(\xi).$$

$$\operatorname{e.g. } \theta(\xi) = \xi/2 \text{ for } -\pi < \xi < \pi.$$



### Approximate Hilbert Transform Pairs

The Phase condition is equivalent to a "half-sample delay":

$$h_n^{(2)} = h_{n+1/2}^{(1)}.$$

Design approximate Hilbert transform pairs, by making error function small near  $\xi = 0$ :

$$E(\xi) = m_{h_2}(2\xi) - e^{-i\xi} m_{h_1}(2\xi).$$

Equivalently, near z = 1,

$$\mathbb{E}(z) = G(z^2) - \frac{1}{z}H(z^2).$$



### The Design Problem

Construct shortest FIR filters  $m_{h_1}$ ,  $m_{h_2}$ , with specified number of zero moments  $K_1$ ,  $K_2$ ; and  $m_{h_2}(\xi) \approx e^{-i\xi/2} m_{h_1}(\xi)$ .



### The Design Problem

Construct shortest FIR filters  $m_{h_1}$ ,  $m_{h_2}$ , with specified number of zero moments  $K_1$ ,  $K_2$ ; and  $m_{h_2}(\xi) \approx e^{-i\xi/2} m_{h_1}(\xi)$ .

- ► The CoCoA Method
- ► The "Flat-Delay Allpass Filter" Method.



## Computational Commutative Algebra (CoCoA)

$$\begin{split} \delta_k &= \sum h_n^{(1)} h_{n+2k}^{(1)} & \lfloor N_1/2 \rfloor \text{ conditions} \\ \delta_k &= \sum h_{n+2k}^{(2)} & \lfloor N_2/2 \rfloor \text{ conditions} \\ H_1(z) &= Q(z) \left(1 + \frac{1}{z}\right)^{K_1} & K_1 \text{ conditions} \\ H_2(z) &= Q(z) \left(1 + \frac{1}{z}\right)^{K_2} & K_2 \text{ conditions} \\ H_2(z^2) &- \frac{1}{z} H_1(z) &= Q(z) \left(1 - \frac{1}{z}\right)^L & L \text{ conditions} \end{split}$$



## Computational Commutative Algebra (CoCoA)

▶ Each condition gives a polynomial

$$p_k \in \mathbb{Q}[X_0, \dots, X_{N_1}, Y_0, \dots, Y_{N_2}].$$

Find Gröbner basis of  $I = (p_1, \ldots, p_s) \subset Q[X, Y]$  with Lex order.

$$(s = \lfloor N_1/2 \rfloor + \lfloor N_2/2 \rfloor + K_1 + K_2 + L)$$

▶ From Gröbner basis, easier to choose a set of solutions.



### Computational Commutative Algebra (CoCoA)

```
))) ---- CoCoA 3.7
                -- by L. Robbiano, A. Capani, G. Niesi,
                -- J. Abbott. A. Bigatti. M. Caboara.
               --- M. Kreuzer, D. Perkinson
             ---- online help : type "Man():"
                    release notes : type "RelNotes():"
-- Current ring is R = Q[t,x,y,z]
Use S::=0[abcdefahl, Lex:
A:=Ideal(a^2+b^2+c^2+d^2-1, e^2+f^2+g^2+h^2-1, ac+bd, eg+fh, a+b+c+d-e-f-g-h);
B:=GBasis(A);
B[1];
a + b + c + d - e - f - a - h
B[2];
ea + fh
B[3];
e^2 + f^2 + g^2 + h^2 - 1
B[4];
bc - bd + c^2 + cd - ce - cf - ca - ch
-2b^2 - 4bd + 2be + 2bf + 2ba + 2bh - 2d^2 + 2de + 2df + 2da + 2dh - 2ef - 2eh - 2fa - 2ah
B[6];
efh - f^2g - g^3 - gh^2 + g
B[7];
2bd^2 - bde - bdf - bdg - bdh - c^3 - c^2d + c^2e + c^2f + c^2g + c^2h - cd^2 + cde + cdf + cdg + cdh - cef - ceh -
cfg - cgh + d^3 - d^2e - d^2f - d^2g - d^2h + def + deh + dfg + dgh
```



# Flat-Delay Allpass Filter $N_1 = N_2$ , $K_1 = K_2$

Design filters with z-transforms  $H_1$  and  $H_2$  satisfying

$$H_1(z) = F(z)D(z),$$
  $H_2(z) = F(z)z^{-L}D(z).$ 



Flat-Delay Allpass Filter 
$$N_1 = N_2$$
,  $K_1 = K_2$ 

Design filters with z-transforms  $H_1$  and  $H_2$  satisfying

$$H_1(z) = F(z)D(z),$$
  $H_2(z) = F(z)z^{-L}D(z).$ 

- $\triangleright$  D is chosen to achieve the (approximate) half-sample delay.
- $\triangleright$  F is chosen to achieve the "vanishing moments" condition.



## Flat-Delay Allpass Filter

 $N_1 = N_2, K_1 = K_2$ 

Design of D:

$$\begin{aligned} \frac{d_n}{d_n} &= (-1)^n \binom{L}{n} \frac{(\frac{1}{2} - L) \cdots (\frac{1}{2} - L + n)}{(\frac{1}{2} + 1) \cdots (\frac{1}{2} + n + 1)}.\\ (n &= 0, 1, \dots, L) \end{aligned}$$



## Flat-Delay Allpass Filter

 $N_1 = N_2, K_1 = K_2$ 

Design of D:

$$d_n = (-1)^n \binom{L}{n} \frac{(\frac{1}{2} - L) \cdots (\frac{1}{2} - L + n)}{(\frac{1}{2} + 1) \cdots (\frac{1}{2} + n + 1)}.$$
  
$$(n = 0, 1, \dots, L)$$

Design of F:

$$F(z) = Q(z) \left(1 + \frac{1}{z}\right)^K$$



# Flat-Delay Allpass Filter $N_1 = N_2$ , $K_1 = K_2$

Design of D:

$$d_n = (-1)^n \binom{L}{n} \frac{(\frac{1}{2} - L) \cdots (\frac{1}{2} - L + n)}{(\frac{1}{2} + 1) \cdots (\frac{1}{2} + n + 1)}.$$
$$(n = 0, 1, \dots, L)$$

Design of F:

$$F(z) = \frac{Q(z)(1 + \frac{1}{z})^K}{2}$$

Design of Q:  $\blacktriangleright$  Generate  $(r_n)$  symmetric, minimal length such that

$$R(z)(z+2+\frac{1}{z})^K D(z)D(1/z)$$
 is halfband.

ightharpoonup Q is a spectral factor of R:

$$R(z) = Q(z)Q(1/z).$$