Notes and exercises from Set Theory

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Introduction

This document contains notes and exercises from [1].

Chapter 1

Exercise (2). There is no set *X* such that $P(X) \subset X$.

Proof. By the axiom of regularity (1.8), X is \in -minimal in $\{X\}$, so $X \not\in X$ and hence $P(X) \not\subset X$.

Exercise (3). If *X* is inductive, then the set $\{x \in X \mid x \subset X\}$ is inductive. Hence *N* is transitive and for each $n \in N$, $n = \{m \in N \mid m < n\}$.

Proof. Let $S = \{x \in X \mid x \subset X\}$. By inductivity of X, $\emptyset \in S$, and if $x \in S$, then $x \cup \{x\} \in S$, so S is inductive. Taking X = N, it follows that S = N since N is the smallest inductive set. Hence $n \in N$ implies $n \subset N$, so N is transitive and $n = \{m \in N \mid m < n\}$.

Remark. We proved transitivity of N "by induction" on N: $0 \subseteq N$ and if $n \subseteq N$ then $n+1 \subseteq N$, so $n \subseteq N$ for all $n \in N$. The following exercises are similar.

Exercise (4). If X is inductive, then the set $\{x \in X \mid x \text{ is transitive}\}$ is inductive. Hence every $n \in N$ is transitive.

Proof. The class C of transitive sets is inductive. Indeed, \emptyset is transitive, and if x is transitive then $x \cup \{x\}$ is transitive since $y \in x \cup \{x\}$ implies $y \subset x \subset x \cup \{x\}$. It follows that $\{x \in X \mid x \text{ is transitive}\} = X \cap C$ is inductive since the intersection of two inductive classes is inductive. Taking X = N, it follows as above that every $n \in N$ is transitive.

Exercise (5). If *X* is inductive, then the set $\{x \in X \mid x \text{ is transitive and } x \notin x\}$ is inductive. Hence $n \notin n$ and $n \neq n+1$ for all $n \in N$.

Proof. The class $C = \{x \mid x \text{ is transitive and } x \notin x\}$ is inductive. Indeed, $\emptyset \in C$. If $x \in C$, then $x \cup \{x\}$ is transitive (by inductivity of the class of transitive sets). Also $x \cup \{x\} \notin x$, lest $x \cup \{x\} \subset x$ by transitivity of x and hence $x \in x$ —contradicting $x \notin x$. Similarly $x \cup \{x\} \neq x$. Therefore $x \cup \{x\} \notin x \cup \{x\}$. So $x \cup \{x\} \in C$, and C is inductive. It follows as above that $X \cap C$ is inductive, and taking X = N that $n \notin n$ and hence $n \neq n+1$ for all $n \in N$. □

Remark. In order to prove that $n \notin n$ for all $n \in N$ by induction on N, we "loaded the induction hypothesis" with transitivity.

Exercise (6). If X is inductive, then the set $\{x \in X \mid x \text{ is transitive and regular}\}$ is inductive, where a set x is called *regular* if every nonempty subset of x has an ϵ -minimal element.

Proof. The class $C = \{x \mid x \text{ is transitive and regular}\}$ is inductive. Indeed, $\emptyset \in C$. If $x \in C$, then $x \cup \{x\}$ is transitive. If $y \subset x \cup \{x\}$ is nonempty, let $z = y - \{x\} \subset x$. If $z = \emptyset$, then $y = \{x\}$, and x is ∈-minimal in y by regularity of x. If $z \neq \emptyset$ and t is ∈-minimal in x, then $x \notin t$ by transitivity and regularity of x, so t is ∈-minimal in y. Therefore $x \cup \{x\}$ is regular. So $x \cup \{x\} \in C$, and C is inductive. It follows as above that $X \cap C$ is inductive.

Remark. Taking X = N it follows as above that every $n \in N$ is regular.

Exercise (7). Every nonempty $X \subset N$ has an \in -minimal element.

Proof. Choose $n \in X$. If n is not ϵ -minimal in X, then $n \cap X$ is a nonempty subset of n, which has an ϵ -minimal element m (Exercise 6). By transitivity of n (Exercise 4), m is ϵ -minimal in X.

Exercise (8). If *X* is inductive, then the set $\{x \in X \mid x = \emptyset \text{ or } \exists y(x = y \cup \{y\})\}$ is inductive. Hence for all $n \in \mathbb{N}$, n = 0 or n = m + 1 for some $m \in \mathbb{N}$.

Proof. The class $C = \{x \mid x = \emptyset \text{ or } \exists y(x = y \cup \{y\})\}$ is obviously inductive, so $X \cap C$ is inductive. Taking X = N, the rest follows, with $m \in N$ by transitivity of N (Exercise 3). □

Exercise (9). Let $A \subset N$ be such that $0 \in A$, and if $n \in A$ then $n + 1 \in A$. Then A = N.

Proof. A is inductive, so A = N since N is the smallest inductive set.

Alternately, if $A \neq N$, let m be \in -minimal in N - A (Exercise 7). Then $m \neq 0$, and $m \neq n + 1$ for any $n \in N$, which is impossible (Exercise 8).

Exercise (10). Each $n \in N$ is T-finite.

Proof. By induction on N. First, n = 0 is T-finite since $P(\emptyset) = \{\emptyset\}$ and \emptyset is \subseteq maximal in $\{\emptyset\}$. If n is T-finite, suppose $X \subseteq P(n \cup \{n\})$ is nonempty. Let

$$X' = \{ u - \{ n \} \mid u \in X \}$$

Then $X' \subset P(n)$ is nonempty and has a \subset -maximal element $u - \{n\}$ for some $u \in X$ by T-finiteness of n, where we may assume $n \in u$ if $u \cup \{n\} \in X$. We claim u is \subset -maximal in X. Indeed, if $v \in X$ and $u \subset v$, then $u - \{n\} \subset v - \{n\}$, so $u - \{n\} = v - \{n\}$ by \subset -maximality of $u - \{n\}$ in X', so u = v by the assumption about u. Therefore n + 1 is T-finite.

Exercise (11). *N* is T-infinite. In fact, $N \subset P(N)$ has no \subset -maximal element.

Proof. We know $N \subset P(N)$ by transitivity of N (Exercise 3), and for all $n \in N$ we have $n \subset n+1 \in N$ but $n \neq n+1$ (Exercise 5).

Remark. If A is T-finite and $\pi: A \to B$ is surjective (onto), then B is T-finite.

Proof. By pullback. If $X \subset P(B)$ is nonempty, let

$$X_{-1} = \pi_{-1}(X) = \{ u_{-1} = \pi_{-1}(u) \mid u \in X \}$$

Then $X_{-1} \subset P(A)$ is nonempty and has a \subset -maximal element u_{-1} for some $u \in X$ by T-finiteness of A. We claim u is \subset -maximal in X. Indeed, if $v \in X$ and $u \subset v$, then $u_{-1} \subset v_{-1}$, so $u_{-1} = v_{-1}$ by \subset -maximality of u_{-1} in u_{-1} , so u = v by surjectivity of u. Therefore u is T-finite.

Exercise (12). Every finite set is T-finite.

Proof. By definition, every finite set is the image of a T-finite natural number (Exercise 10), so is T-finite by the previous remark.

Alternately, if S is T-infinite, choose $X \subset P(S)$ nonempty with no \subset -maximal element. By induction on N, for each $n \in N$ there is a properly ascending chain $u_0 \subset \cdots \subset u_n \subset S$ of length n+1 in X, and hence an injection $n \to S$ sending $m \in n$ into $u_{m+1} - u_m$. If S had k elements for some $k \in N$, there would then be an injection $k+1 \to k$, which is impossible by an easy induction on k. Therefore S is infinite.

Exercise (13). Every infinite set is T-infinite.

Proof. If *S* is infinite, let *X* be the set of finite subsets of *S*. Clearly *X* is nonempty since $\emptyset \in X$. Also *X* has no ⊂-maximal element. Indeed, if $u \in X$ then there is $s \in S - u$ since *S* is infinite, and $u \cup \{s\}$ is finite (if *u* has *k* elements, then $u \cup \{s\}$ has k + 1 elements), so $u \subset u \cup \{s\} \in X$ where the inclusion is proper. Therefore *S* is T-infinite.

Remark. The previous two exercises show that (Cantor) finiteness is equivalent to Tarski finiteness in ZF.

Chapter 2

Remark. In (2.1), < is actually a *well-ordering* on *Ord.* Indeed, if $C \subset Ord$ is a nonempty class of ordinals and $\alpha \in C$ is not \in -minimal in C, then $\alpha \cap C$ is a nonempty subset of α which has an \in -minimal element β . By transitivity of α , β is \in -minimal in C.

This observation provides an alternative proof that Ord is a proper class: if Ord were a set, then because it is transitive and strictly well-ordered by \in , it would be an ordinal, and hence $Ord \in Ord$ —contradicting strictness.

References

[1] Jech, Thomas. Set Theory, 3rd ed. Springer, 2002.