CSE107 Discrete Mathematics and Statistics

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1. Number Systems and Proof Techniques

1.1. Types of Numbers

Name	
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Name	Property	
Natural Numbers	Closed under addition	
Integer		
Relational Numbers	Can be expressed as $rac{x}{y}$ where $x,y\in Q$ and $y eq 0$	
Real Numbers		
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Prime Numbers *质数* You all know that.

1.2. Proof Techniques

1.2.1. Proof by contradiction

Example: Use proof by contracdition to prove $\sqrt{2}$ is a irrational number:

- ullet If $\sqrt{2}$ were rational, then it can be written as $\sqrt{2}=rac{x}{y}$ where $x,y\in Q$ and x,y
 eq 0
- By repeatedly cancelling common factors, we can make sure that x,y have no common factors so they are not both even.
- Since $2=rac{x^2}{y^2}$, so x^2 is even, then x is even.
- $\bullet \ \ \operatorname{Let} x = 2\overset{\circ}{w} \text{, } w \in N.$
- ullet Then $x^2=4w^2$ so $4w^2=2y^2$, $y^2=2w^2$ so y^2 is even then so is y .
- This contradicts the fact that x and y are not both even, so 2 being rational, must have been wrong.

1.2.2. Proof by induction

Normal steps:

- 1. Prove the first number holds the property.
- 2. Prove that if n=m holds the property, then so is n=m+1.

Examples:

ullet For every natural number n, $2^{n+2}+3^{2n+1}$ is divisible by 7

1.3. Tutorial 1

- Use counter-example for easy question.
- When it comes to rational numbers, use the form $\frac{x}{y}$ to represent it.
- ullet To prove \sqrt{n} is not rational, usually the parity on the sides of the equation is different.

2. Set Theory

2.1. Notation

$$B=\{1,3,5\}\Rightarrow 1\in B, 2\notin B$$

For inexplicit set: $S=\{x|P(x)\}$ e.g.

- $\begin{array}{l} \bullet \quad S = \{x | x \text{ is an odd positive integer}\} \\ \bullet \quad C = \{n^2 | n \text{ is an integer}\} \end{array}$

2.2. Important Sets

Name	Notation	Field
Natural Numbers	N	$\{0,1,2\ldots\}$
Integers	Z	$\{\ldots-1,0,1\ldots\}$
Positive Integers	Z^+	$\{1,2,3\ldots\}$
Rationals	Q	\$\lbrace\frac{x}{y}
Real numbers	R	Z+Q

2.3. Subset

ullet If **every** element of B is an element of A, then B is the subset of A. Notation: $B\subseteq A$.

2.4. Equality

If $B\subseteq A$ and $A\subseteq B$, then A=B.

2.5. Operations on sets

First of all, the sets to be operated must have the same data type.

Name	Notation	Definition
Union	$A \bigcup B$	$\{x\mid x\in A\bigvee x\in B\}$
Intersection	$A \bigcap B$	$\{x\mid x\in Aigwedge x\in B\}$
Relative complement	A - B	$\{x\mid x\in Aigwedge x otin B\}$
Complement	$\sim A$	$\{x otin A\}$
Symmetric difference	$A\triangle B$	$\{x\mid (x\in Aigwedge x otin B)igvee(x otin Aigwedge x\in B)\}$

2.6. Algebra of sets

Suppose $A,B,C\subseteq U$ in all the situations below.

Name	Content
Commutative laws	$A \bigcup B = B \bigcup A, A \cap B = B \cap A$
Associative laws 结 合律	$A \bigcup (B \bigcup C) = (A \bigcup B) \bigcup C, A \bigcap (B \bigcap C) = (A \bigcap B) \bigcap C$
Identity laws	$A igcup \emptyset = A, A igcup U = U, A igcap U = A, A igcap \emptyset = \emptyset$

Name	Content
Distributive laws分 配律	$A \cap (B \cup C) = (A \cap B) \cup (A \cap C), A \cup (B \cap C) = (A \cup B) \cap (A \cap C)$
Complement laws	$Aigcup_{}\sim A=U, \sim U=\emptyset, \sim (\sim A)=A, Aigcap_{}\sim A=\emptyset, \sim \emptyset=U$
De Morgan's laws	$\sim (A igcup B) = \sim A igcap \sim B, \sim (A igcap B) = \sim A igcup \sim B$

2.7. Power set

The power set Pow(A) is the set of all the subsets of A, denoted by $Pow(A) = \{C | C \subseteq A\}$.

2.7.1. Some laws

- $\forall A, B, Pow(A \cap B) = Pow(A) \cap Pow(B)$
- $s\exists A, B, Pow(A \cup B) \neq Pow(A) \cup Pow(B)$

2.8. Cardinality of sets

For a **finite** set S, the cardinality equals the number of the elements in S, denoted by |S|.

2.8.1. Computing the cardinality of sets

If A, B, C are sets then:

$$|A \bigcup B| = |A| + |B| - |A \cap B|$$

$$|A \bigcup B \bigcup C| = |A| + |B| + |C| - |A \cap B| - |A \cap C| - |B \cap C| + |A \cap B| + |A \cap C|$$

Those above are special cases of the principle of inclusion and exclusion.

2.9. Ordered pairs

Introduced **cartesian product** $A \times B$, the set of consisting all pairs (a,b) with $a \in A \bigwedge b \in B$.

2.9.1. Catesian plane

Use a cartesian coordinate system to represent pairs of numbers.

2.10. Bit strings of lengs n

Used to represent the **characteristic vector** of subsets of a set.

e.g.
$$S = \{1, 2, 3, 4, 5\}, A = \{1, 3, 5\}, B = \{3, 4\}$$

- the CV(characteristic vector) of A is {1,0,1,0,1}
- the CV of B is {0,0,1,1,0}

2.11. Tutorial 2

• The catesian product holds the **distrubutive laws** as well.

3. Relations

3.1. Definition (Binary relation 二元关系)

3.2. Representation

Туре	Example
Directed Graph(Digraph)	1 2 3 4 5 6
	7
Martix	$M = egin{bmatrix} T & T & F \ F & F & F \ T & T & T \end{bmatrix}$

3.3. Unary relation 单元关系

3.4. Infix noatation for binary relation

• Suppose R is a binary relation, then it can be write as xRy where $(x,y)\in R$.

3.5. Properties of binary relations

Suppose a binary relation R on a set A.

Name	Meaning
reflexive	$orall x \in A, xRx$
symmetric	$\forall x,y \in A, xRy \rightarrow yRx$
antisymmetric	$orall x,y\in A, xRyigwedge yRx o x=y$
transitive	$orall x,y,z\in A, xRyigwedge yRz o xRz$

• It can be reflected by digraph as well.

3.6. Transitive closure 传递闭包

• Example: Let A = $\{1, 2, 3\}$. Find the transitive closure of R = $\{(1, 1), (1, 2), (1, 3), (2, 3), (3, 1)\}$.

$$(3,1),(1,2) o (3,2)$$
 $(2,3),(3,1) o (2,1)$ $(2,1),(1,2) o (2,2)$ $(3,1),(1,3) o (3,3)$ so $t(R)=\{(1,1),(1,2),(1,3),(2,3),(3,1),(3,2),(2,1),(2,2),(3,3)\}$

• Obiviously, digraphs are more intuitive: check the graph of R, start at each node, find its end point, if there is no edge from that node to that end point, add that edge, and will get the relationship of t(R).

3.7. Equivalence Relations

- A binary relation being reflexive, transitive, and symmetric is called a equivalence relation.
- Equivalence class can be denoted by $E_x = \{y \mid yRx\}$ where $orall x \in A$.
 - $\circ \;\;$ e.g. $E_0=Z$ is the equivalence class of 0.

3.8. Partition of a set

Just remeber it has something to with equivalence relations.

3.9. Partial orders

- A binary relation being s reflexive, transitive and antisymmetric is called a **paritial order**.
 - \circ e.g. the relation \subseteq on Pow(A)

3.9.1. Predecessors (我也不知道翻译是啥)(好了我知道了,是前元,感谢Ramos)

- ullet x is a **predecessor** of y when R is a partial order on a set A and xRy, x
 eq y
- x is an **immediate predecessor** of y when x is a predecessor of y and there is no $z \notin \{x,y\}$ where xRz,zRy.

3.10. Hasse diagram

• The Hasse Diagram of a partial order is a digraph.

3.11. Total Order

- A binary relation R on a set A is a total order if it is a partial order such that $\forall x,y \in A, xRy \bigvee yRx$.
- The Hasse diagram of a total order is a **chain**.

3.12. Tutorial 3

- For matrices, figure out the right way to read it. The first number is the horizontal coordinate (up to down); the second is the vertical coordinate(left to right).
- A relation cannot be proved reflexive by combining transistive and symmetric, because there may not be all the elements in the relation compared with the set.9