

Live-Verification While Programming:

Dafny

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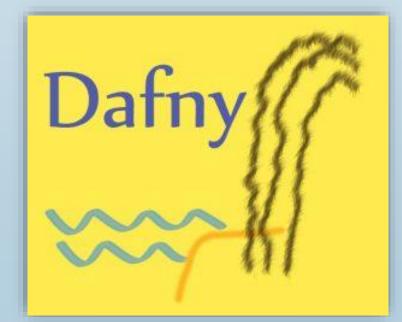


Image Source: https://www.microsoft.com/

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Introduction



- Dafny was developed by K. Rustan M. Leino at Microsoft in 2008
- Part of Microsoft Research RiSE (Research in Software Engineering)
- Live Verification Language
- Cross Language Compiler to
 - JavaScript, Go, .NET Languages
- Supports
 - Generic Classes
 - Dynamic Allocation
 - Specification Constructs
 - Pre- and Post-Conditions
 - Ghost State
 - Dynamic Frames
 - Proof of Termination







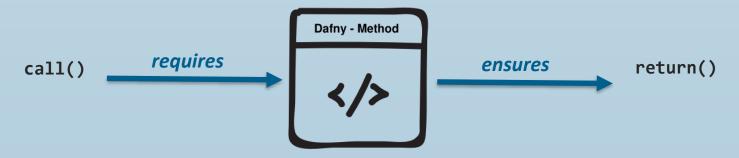
Image Sources:

http://ase-conferences.org/ase/past/ase2007/images/RustanLeino.png https://www.microsoft.com/

Pre- and Post- Conditions



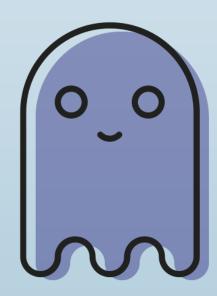
- Pre- and Post- Conditions (requires and ensures) are checked before compilation
 - Works similar to Typing constraints in other programming languages like C# or Java
 - Dafny analyzes possible code paths, to validate the solvability of the given specifications
- This assures that certain specifications are complied before and after method execution
 - Preconditions need to be ensured before the Methods Body
 - Postconditions need to be ensured after the Methods Body



Ghost State



- Everything that is specified as ghost is only considered by the Verifier,
 and will be completely <u>ignored</u> by the Compiler
- Variables can be initialized as Ghost Variables
- Ghost Methods are called Functions
- Ghost Variables and Functions are prohibited to influence normal Variables and Methods
 - This is also checked by the Verifier



Dynamic Frames



- Dynamic Frames assure that only the selected objects can be accessed by the current Method
- The 'modifies' statement assures, that only the selected objects can be altered by the current Method
- The 'reads' statement assures, that only the selected objects can be read by the current Method
- The selection is recursive complete, which means every object that is a parameter of the selected object (and so on) will also be accessible.



Proof of Termination



- Dafny can check if a Method with Loops and / or Recursion will terminate eventually
- The termination proof will only take seconds, even if the programs termination itself could need years
 - E.g. recursive Ackermann or Fibonacci



Proof of Termination

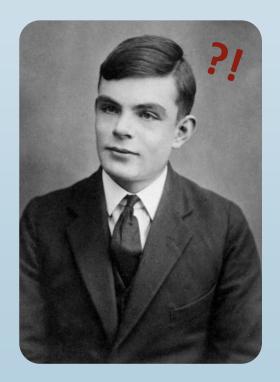


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WAIT A MINUTE!?

(Pun not intended)

Isn't that the Halting-Problem?
Why can Dafny solve NP Complete??



Proof of Termination



 Dafny can check if a Method with Loops and / or Recursion will terminate eventually

- The termination proof will only take seconds, even if the programs termination itself could need years
 - E.g. recursive Ackermann or Fibonacci
- Dafny will analyze the termination with the help of the 'decreases' variable and tries to figure out if the selected variable decreases towards it's halting anchor
- This works for Loops and Recursion, but there could still be algorithms that <u>can not</u> be proven to terminate
 - This will result in a verification timeout and a warning

Is this program correct? 1 function Ackermann(m: int, n: int): int 2 decreases m, n 3 { 4 if m <= 0 then 5 n + 1 6 else if n <= 0 then 7 Ackermann(m - 1, 1) 8 else 9 Ackermann(m - 1, Ackermann(m, n - 1)) 10 }

How? Magic?



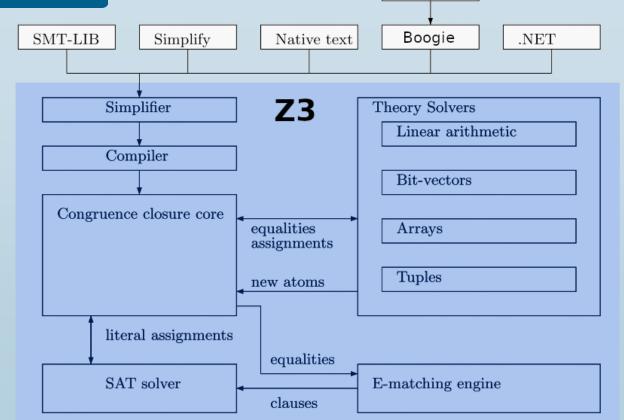
Result

- How can Dafny verify all those things?
- Dafny uses an Intermediate Verification
 Language called Boogie (Microsoft)
- Dafny Program Program Program SMT Z3
- Boogie is able to translate specifications defined by Dafny into Mathematical Formulas, the so called 'Theories'
- Those Mathematical Formulas will then be solved by the SMT-Solver Z3 (Microsoft)
- This will result in a feedback if (e.g.) a Pre- or Post- Condition can be fulfilled or not

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SMT Solver





Dafny

Satisfiability Modulo Theories



Specification - Classes



- Abstract Classes are called Traits
- Classes may extend other classes or Traits
- Consist of Constructor, Variables and Methods
- Variables can be initialized as Ghost Variables

```
class Account {

var balance : real;

constructors - constructor (balance: real) {this.balance := balance; }

method deposit(amount: real) modifies this { balance := balance + amount ;}

methods - method withdraw(amount: real) modifies this { balance := balance - amount ;}

method getBalance() returns (res : real) { return balance; }

}
```

Specification - Methods



- Has multiple Input and Output Parameters
- Uses Pre- and Post- Conditions
- Uses Dynamic Frames to manage Memory Access
- Proofs Termination of Loops and Recursive Methods
- Ghost Methods are called Functions

```
type
                                           in-parameters
                                                                           out-parameters
         attributes
                         name params
method {:att1}{:att2} M<T1, T2>(a: A, b: B, c: C) returns (x: X, y: Y, z: Z)
                      precondition (boolean expression)
  requires Pre
                      objects whose fields may be updated by the method
  modifies Frame
                      postcondition (boolean expression)
  ensures Post
  decreases Rank
                      variant function (to prove termination of recursive methods)
  Body
          imperative style (statement or sequence of statements)
```

Example: Queue Datastructure



```
method Dequeue () Not Empty
                                    Dynamic Frame
     requires Valid();
     requires 0 < |contents|
                                        Swinging Pivot Restriction
     modifies footprint;
     ensures Valid() \( \) fresh(footprint - old(footprint));
5
6
     ensures contents = old(contents)[1..];
                                      Old Queue but WITHOUT
     var n := head.next;
     head := n:
                                     the first entry
     contents := n.tailContents;
10
                                  Function Body
```



Literature



Introduction and Theory:

- 'Live-Verification while Programming: Dafny' (Philippe Heiler, 2020)
- 'Dafny: An automatic program verifier for functional correctness' (K. Rustan M. Leino, 2010)
- 'Specification and Verification of Object-Oriented Software' (K. Rustan M. Leino, 2010)

Programming Documentation:

- 'Dafny Quick Reference' (J. Pascoal Faria, 2020)
- 'Dafny Reference Manual' (Richard L. Ford et al. 2017)

Try out Dafny:

- 'Dafny GitHub Page'

 (https://github.com/dafny-lang/dafny/releases)
- 'Dafny Online Verifier' (https://rise4fun.com/Dafny/)

