



Longhand division in binary is similar to that in decimal

$$\begin{array}{r} 21 \\ 13 \overline{) 274} \\ \underline{26} \\ 14 \\ \underline{13} \\ 1 \end{array}$$

$$\begin{array}{r} 10101 \\ 1101 \overline{) 100010010} \\ \underline{1101} \\ 10000 \\ \underline{1101} \\ 1110 \\ \underline{1101} \\ 1 \end{array}$$

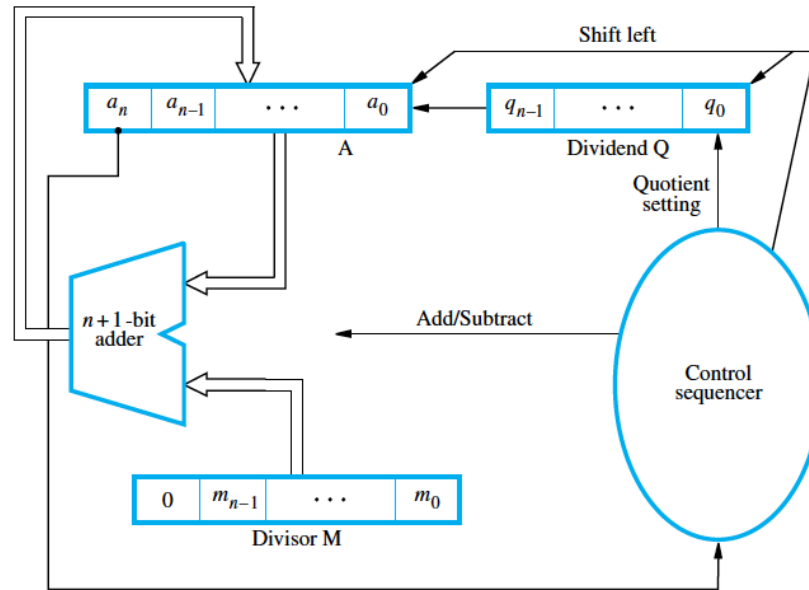
Both a quotient and a remainder are produced

Handle signed numbers by dividing the positive equivalents

Sign of remainder = sign of dividend (dividend rule)

Quotient is negative if the divisor and dividend differ in sign

Set A to 0  
Put divisor in M  
Put dividend in Q

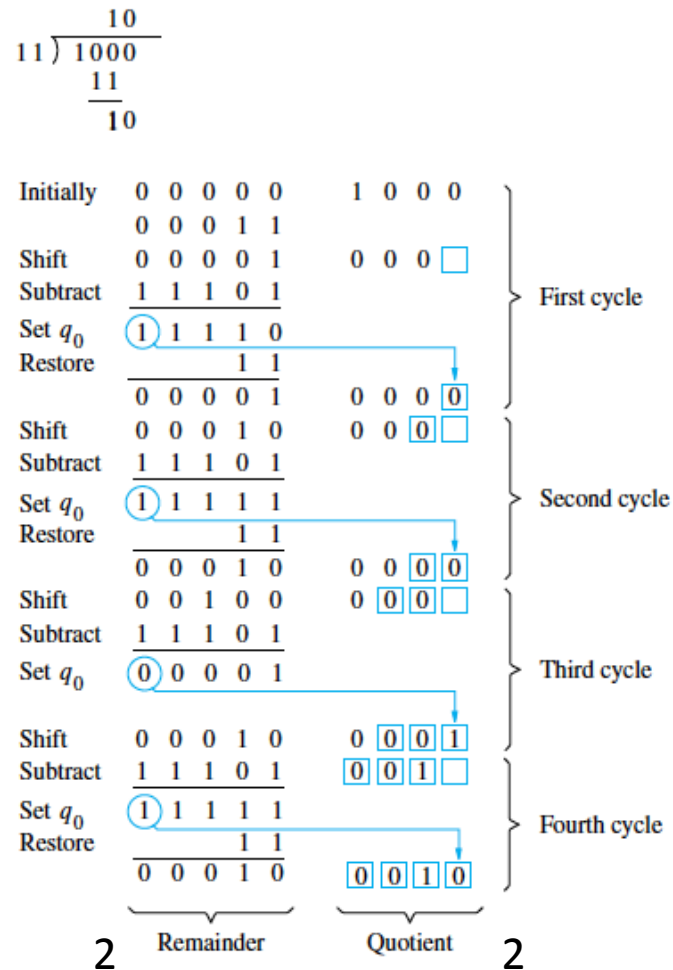


Repeat the following 3 steps  $n$  times:

1. Shift A and Q left one bit position
2. Subtract M from A
3. If A is negative, set  $q_0$  to 0 and add M back to A (i.e. restore A); otherwise set  $q_0$  to 1

When done, Q contains quotient and A contains the remainder

Assume 4-bit values. Division of 8 by 3



Avoids having to add M back when subtracting makes  $A < 0$

Restoring division computes  $A-M$  and if negative adds A back before shifting left 1 bit for the next cycle

Shifting  $A-M$  left first and then adding M gives:

$$2(A-M)+M = 2A - M \text{ (which is needed in the next cycle)}$$

This avoids the restore step

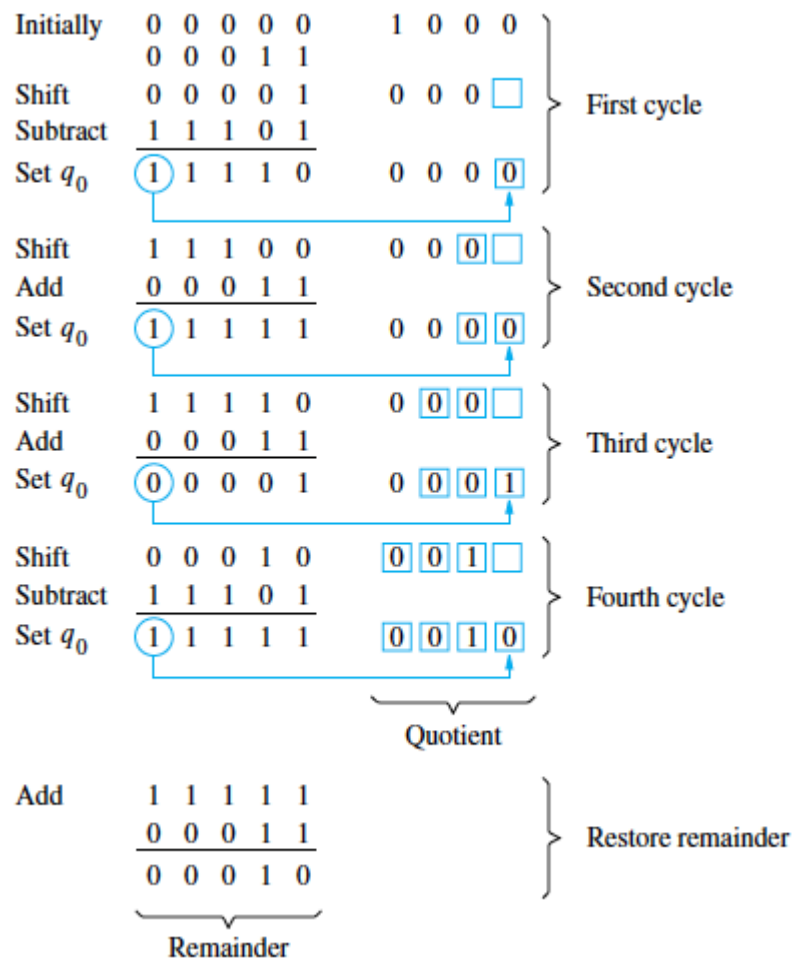
Stage 1: Repeat the following 2 steps  $n$  times:

1. If  $A < 0$ , shift  $A$  and  $Q$  left 1 bit and add  $M$  to  $A$   
Else shift  $A$  and  $Q$  left 1 bit and subtract  $M$  from  $A$
2. If  $A < 0$ , set  $q_0 = 0$  else set  $q_0 = 1$

Stage 2: If  $A < 0$ , add  $M$  to  $A$

Stage 2 is needed to leave the proper positive remainder in  $A$

Assume 4-bit values. Division of 8 by 3



Use absolute values of the dividend and divisor

Quotient is negative if the signs of divisor & dividend differ

Sign of remainder should match the sign of the dividend  
this is called the dividend rule