



# Module 4

MIPS Assembly Language;  
Control Structures



# Module Four

- MIPS Assembly Language; Control Structures - Part Two
- In this presentation, we are going to talk about :
- MIPS Assembly Language support for Subroutines



# Overview

- Previously we talked about:
  - Process Flow
  - Control Structures
    - IF – THEN – ELSE
    - Loops
- Now: Subroutines



# Subroutines

- Sub-programs
- Program code that is used multiple places in a program.
- Program code that is common and copied from a library.
- Main program jumps to the subroutine code, and subroutine jumps back to the main program.

- MIPS

JAL - Jump and Link

- Accepts argument values
- Functions

Subset of subroutines that return a result value.



# Subroutines

- Main Program
  - Sets argument values
  - Jumps to the subroutine code
- Subroutine
  - Allocates memory space
  - Runs the procedure task
  - Sets the return values
  - Frees memory space
  - Jumps back to Main Program



# MIPS

- Call a subroutine
- Jump and Link instructions **JAL** routine-name
- Address of subroutine being called
- Return from subroutine
- Return Jump instructions **JR** \$ra

# Jump and Link

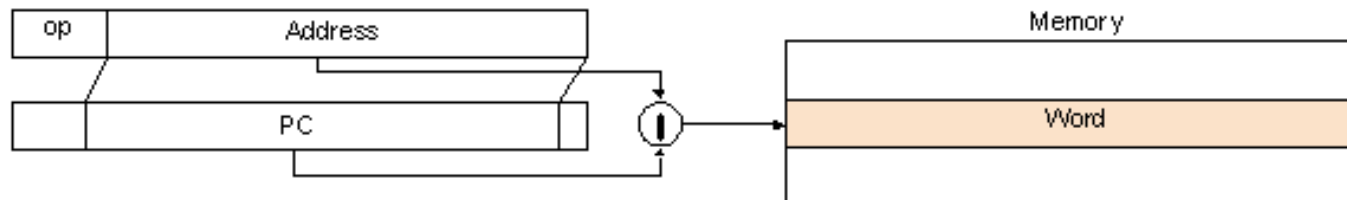
## Jump and link

`jal target`

3	target
6	26

Unconditionally jump to the instruction at target. Save the address of the next instruction in register `$ra`.

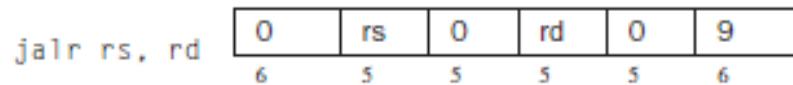
## Pseudodirect addressing



$PC + 4 \longrightarrow$  `$ra` Register

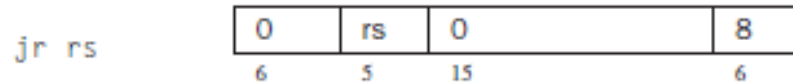
# Jump Register

## Jump and link register



Unconditionally jump to the instruction whose address is in register `rs`. Save the address of the next instruction in register `rd` (which defaults to 31).

## Jump register



Unconditionally jump to the instruction whose address is in register `rs`.





# Arguments

- Argument values passed to the subroutine
- Several methods
- In registers

As values - copy of value - very safe

As addresses - address - can cause 'side effects'

- In memory
- On the Stack



# MIPS Registers

- General agreement of MIPS programmers
- \$a0 - \$a3 first four values input to the subroutine  
if more - then on Stack
- \$v0 - \$v1 returned result values  
if more - then on Stack
- \$ra - the return address
- \$t0 - \$t9 registers may be used by subroutine
- \$s0 - \$s7 values must be saved if registers used
- \$sp - Stack pointer



# A MIPS example

- Can you figure out the code?

```
swap(int v[], int k);  
{  int temp;
```

```
    temp = v[k];  
    v[k] = v[k+1];  
    v[k+1] = temp;
```

```
}
```

swap:

```
mulh $v0, $a1, 4    ; k * 4  
add  $v0, $a0, $v0  ; v + k
```

```
lw   $t6, 0($v0)    ; v[k]  
lw   $t7, 4($v0)    ; v[k+1]  
sw   $t6, 4($v0)    ; v[k+1]  
sw   $t7, 0($v0)    ; v[k]
```

```
jr   $ra
```



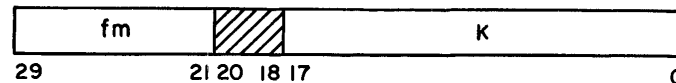
# Another Method

- The CDC 6500 was an early super computer. 1965
- It used direct addresses in the instructions.
- Some instructions were 30 bits.
- The word size is 60 bits.
- The subroutine call instruction is the Return Jump.

# CDC 6500 Return Jump

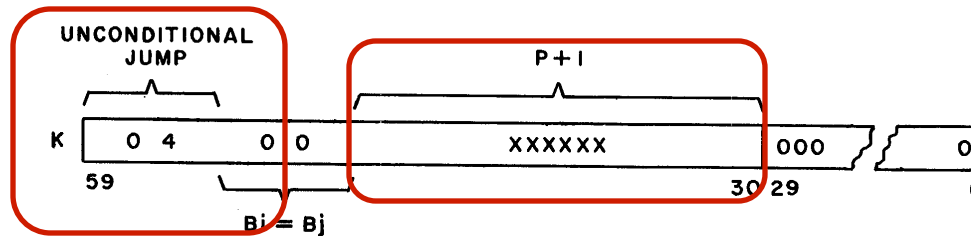
Branch

**010**      **RJ**      **K**      *Return jump to K*      (30 Bits)



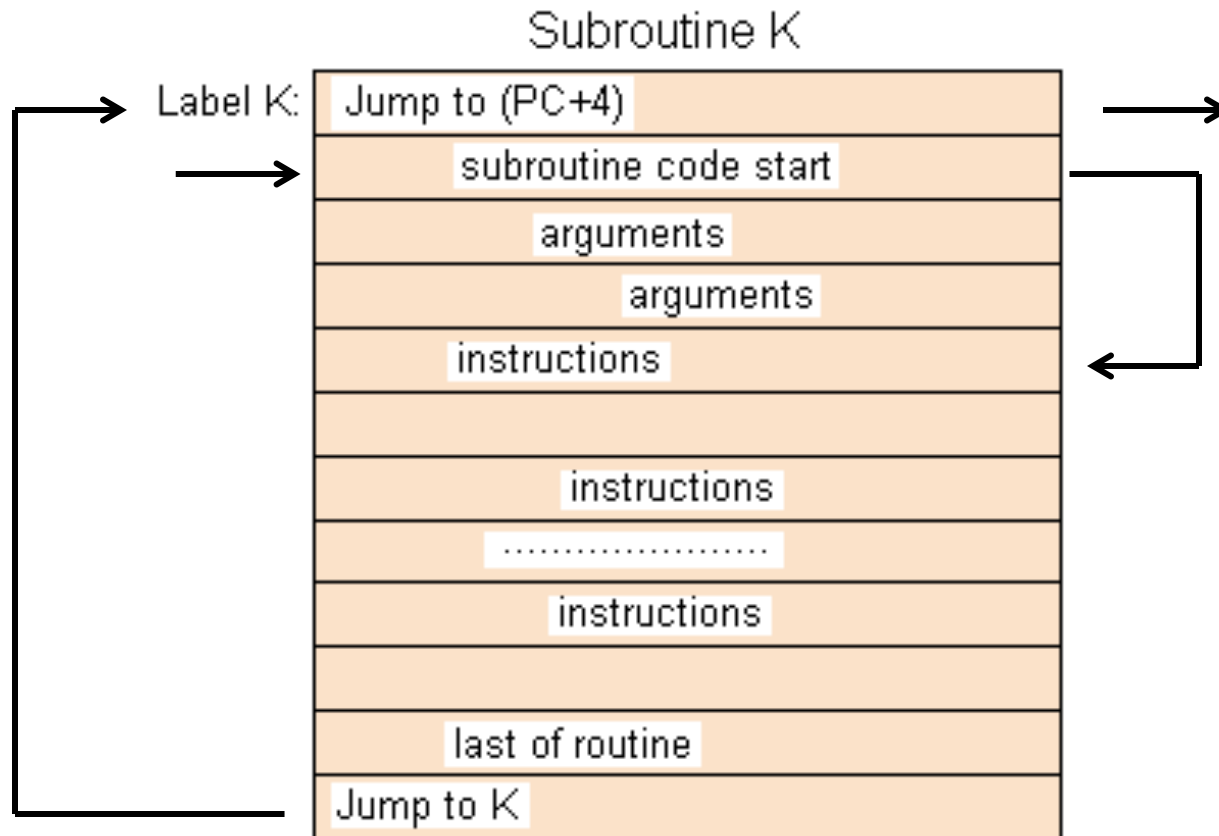
The instruction stores an 04 unconditional jump and the current address plus one  $[P + 1]$  in the upper half of address K, then branches to  $K + 1$  for the next instruction. Note that this instruction is always out of the instruction stack, thus voiding the stack.

The octal word at K after the instruction appears as follows:



A jump to address K at the end of the branch routine returns the program to the original sequence.

# Return Jump in action





# Types of Subroutines

- Leaf - subroutine that processes and returns
- Nested - subroutines that call other subroutines before returning  
Need to save the return address value in \$ra
- Recursive - subroutines that call themselves  
Can be direct or via additional subroutines
- Re-entrant - allocates stack and memory space each time called



# Summary

- MIPS Assembly Language and control structures
  - Subroutines
- Next: Data structure support for Subroutines