



Computer Organization

605.204

Module Three

Part Four

Language for the People



Module Three

- Part Four
- In this presentation, we are going to talk about :
-
- A second set of MIPS assembly language instructions



Previously

- Previously we talked about:
- A Language for the People - Part one
- Now: A Language for the People - Part two



At this point:

- MIPS
 - Loading words but addressing bytes
 - Arithmetic operands in registers only

<u>Instruction</u>	<u>Meaning</u>
add \$s1, \$s2, \$s3	$\$s1 = \$s2 + \$s3$
sub \$s1, \$s2, \$s3	$\$s1 = \$s2 - \$s3$
lw \$s1, 100(\$s2)	$\$s1 = \text{Memory}[\$s2+100]$
sw \$s1, 100(\$s2)	$\text{Memory}[\$s2+100] = \$s1$



Assembly Control Instructions

- Decision making instructions
 - alter the control flow,
 - select the "next" instruction to be executed
- **Branch**
- MIPS conditional branch instructions:

```
bne $t0, $t1, Label  
beq $t0, $t1, Label
```

- Example: if (i==j) h = i + j;

```
        bne $s0, $s1, Label  
        add $s3, $s0, $s1  
Label:      ....
```



Assembly Control Instruction

- **Jump**
- MIPS unconditional branch instructions:

```
j    label
```

- Example:

```
if (i!=j)                beq $s4, $s5, Lab1
    h=i+j;               add $s3, $s4, $s5
else                      j  Lab2
    h=i-j;               Lab1: sub $s3, $s4, $s5
                        Lab2: ...
```



A New Assembly Instruction

- We have: beq, bne, what about Branch-if-less-than?
- New instruction:

slt \$at, \$s1, \$s2

if $\$s1 < \$s2$ then
 \$at = 1
else
 \$at = 0

- Can use this instruction to build: "**blt \$s1, \$s2, Label**"
Can now build general control structures
- The assembler needs a register to do this, **\$at**
There is a policy of use convention for registers



Assembly Instructions for Constants

- Small constants are used quite frequently (50% of operands)

for example: $AB = A + 5;$
 $CQ = C - 18;$

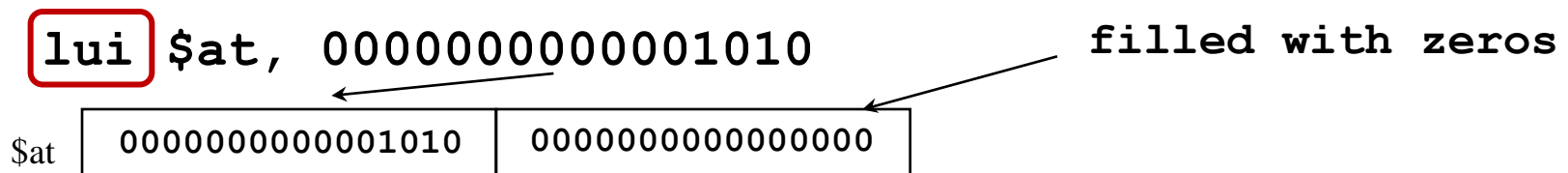
- Solutions ? ?
 - put 'typical constants' in memory and load them, or ...
 - create hard-wired registers (like \$zero) for constants, or ...
- MIPS Instructions:

```
addi $29, $29, 4
slti $8,  $18, 10
andi $29, $29, 6
ori  $29, $29, 4
```

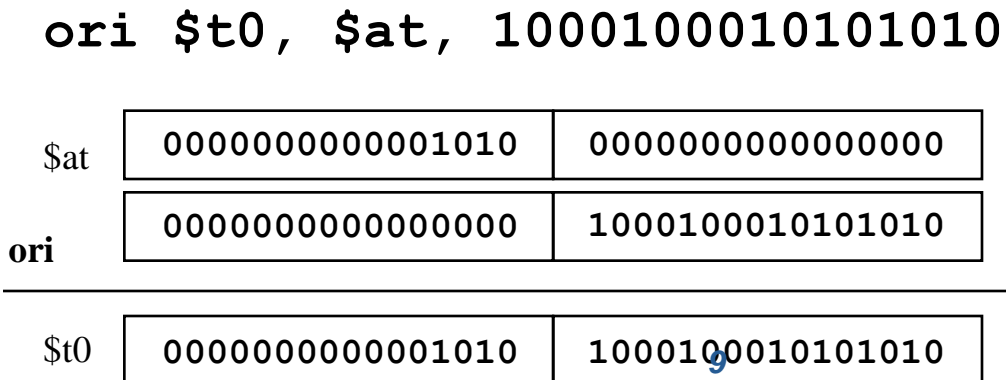



Large Assembly Constants

- We'd like to be able to load a big value constant into a register
- Must use two instructions, new "load upper immediate" instruction



- Then must get the lower order bits correct,





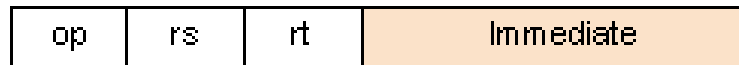
Assembly Addressing modes

- Immediate
- Register
- Base Relative
- PC Relative
- Pseudo Direct

Addressing Modes

- Immediate
 - ADDI

1. Immediate addressing



- Register
 - MULT

2. Register addressing

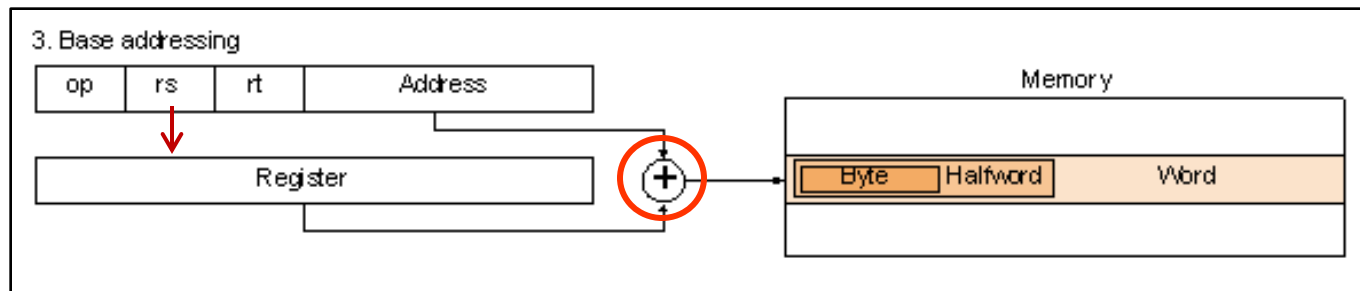


Registers

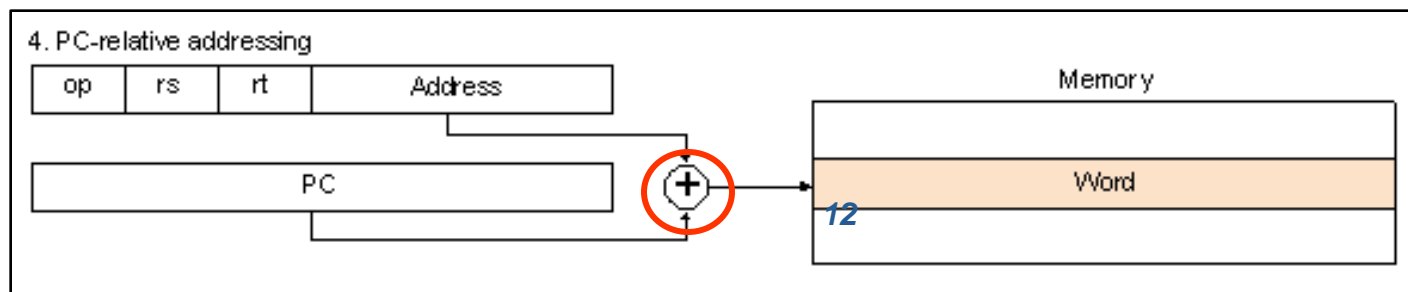
Register

Addressing Modes

- Base Relative
 - Memory access

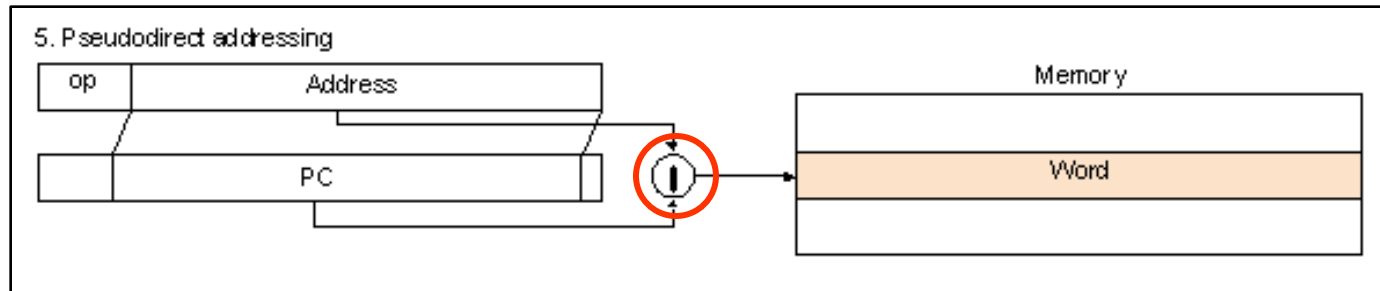


- PC Relative
 - Branch



Addressing Modes

- Pseudo Direct
 - Jump



concatenate

MIPS Assembly Summary

MIPS assembly language

Category	Instruction	Example	Meaning	Comments
Arithmetic	add	add \$s1, \$s2, \$s3	$\$s1 = \$s2 + \$s3$	Three operands; data in registers
	subtract	sub \$s1, \$s2, \$s3	$\$s1 = \$s2 - \$s3$	Three operands; data in registers
	add immediate	addi \$s1, \$s2, 100	$\$s1 = \$s2 + 100$	Used to add constants
Data transfer	load word	lw \$s1, 100(\$s2)	$\$s1 = \text{Memory}[\$s2 + 100]$	Word from memory to register
	store word	sw \$s1, 100(\$s2)	$\text{Memory}[\$s2 + 100] = \$s1$	Word from register to memory
	load byte	lb \$s1, 100(\$s2)	$\$s1 = \text{Memory}[\$s2 + 100]$	Byte from memory to register
	store byte	sb \$s1, 100(\$s2)	$\text{Memory}[\$s2 + 100] = \$s1$	Byte from register to memory
	load upper immediate	lui \$s1, 100	$\$s1 = 100 * 2^{16}$	Loads constant in upper 16 bits
Conditional branch	branch on equal	beq \$s1, \$s2, 25	if ($\$s1 == \$s2$) go to PC + 4 + 100	Equal test; PC-relative branch
	branch on not equal	bne \$s1, \$s2, 25	if ($\$s1 \neq \$s2$) go to PC + 4 + 100	Not equal test; PC-relative
	set on less than	slt \$s1, \$s2, \$s3	if ($\$s2 < \$s3$) $\$s1 = 1$; else $\$s1 = 0$	Compare less than; for beq, bne
	set less than immediate	slti \$s1, \$s2, 100	if ($\$s2 < 100$) $\$s1 = 1$; else $\$s1 = 0$	Compare less than constant
Unconditional jump	jump	j 2500	go to 10000	Jump to target address
	jump register	jr \$ra	go to \$ra	For switch, procedure return
	jump and link	jal 2500	$\$ra = \text{PC} + 4$; go to 10000	For procedure call



Review

- Design Principles:
 - **Simplicity favors regularity**
 - **Smaller is faster**
 - **Good design demands compromise**
 - **Make the common case fast**



Summary

- More of the MIPS Assembly Language

Next: The MIPS Machine Language