

■ Memory mapped Input/Output

- The code below is an example illustrating the use of mapped I/O
- Polls keyboard to echo characters until CR (carriage return) is hit

	LEA	EBX, LOC	Initialize register EBX to point to the address of the first location in main memory where the characters are to be stored.
READ:	BT	KBD_STATUS, 1	Wait for a character to be entered in the keyboard buffer KBD_DATA.
	JNC	READ	
	MOV	AL, KBD_DATA	Transfer character into AL (this clears KIN to 0).
	MOV	[EBX], AL	Store the character in memory
	INC	EBX	and increment pointer.
ECHO:	BT	DISP_STATUS, 2	Wait for the display to become ready.
	JNC	ECHO	
	MOV	DISP_DATA, AL	Move the character just read to the display buffer register (this clears DOUT to 0).
	CMP	AL, CR	If it is not CR, then
	JNE	READ	branch back and read another character.

- Port Mapped I/O (isolated I/O) is common with IA-32

- Requires special Input/Output instructions

- Examples:

`IN AL,DX` ; reads one byte from the devices whose port number is in DX

`IN AX,DX` ; reads two bytes of data from the port

`IN EAX,DX` ; reads 4 bytes

`OUT 61h,AX` ; sends two bytes to output port number 61h

`OUT DX,AL` ; sends one byte to the port whose number is in DX

Port mapped scheme uses I/O address space separate from the program address space.

- IA-32 Interrupts and Exceptions
 - Two interrupt lines
 - NMI non-maskable is always accepted by processor
 - INTR user interrupt is maskable
 - Enabled/disabled using IF flag within status register
 - Accepted if priority > privilege level of currently running program
 - Events other than external interrupts cause exceptions
 - Vector number is assigned to each interrupt or exception
 - This index identifies an IDT (interrupt descriptor table) entry
 - Table entry contains starting address of corresponding handler
 - I/O devices are connected through an APIC
 - Advanced Programmable Interrupt Controller
 - Controller implements priority scheme and sends vector number to CPU

■ Actions taken for Interrupts or Exceptions

1. Push status register, code segment register and EIP onto processor stack
2. For exceptions due to abnormal condition, push cause of exception
3. Disable further interrupts of the same type
4. Use vector number to load appropriate handler address from IDT into EIP

When finished, the handler uses the IRET instruction to resume the program.

IRET pops EIP, code segment register and status register from stack

Handler must undo any changes it made to stack pointer before executing IRET