



Computer Organization

605.204

Module Three

Part Three

Language for the People



Module Three

- Part Three
- In this presentation, we are going to talk about :
- Assembly Language
- The first set of MIPS assembly language instructions



Previously

- Previously we talked about:
- Language for the Machine
- Now: A Language for the People



People Language

- People read and write ‘natural’ languages
 - English
 - Spanish
 - Chinese
- People use words, phrases, nouns, pronouns, and verbs
- Natural languages are ambiguous (English is the worst)
 - “The chicken is ready to eat.”



People Language

- People prefer to write computer programs in languages that most easily allow them to describe the problems they need to solve.
- High level programming languages were developed for this very reason.
- But ...first ...
- People understand letter patterns better than number patterns

00

LDA

12

SUB

28

COMP



Hardware Assembly Language

Symbolic language,

- almost 1 to 1 with the machine language.

- Fixed Format Fixed or Variable Length
- One instruction per line

Label: Operation Fixed number of operands #Comments

loop:	LW	\$v1, 0(\$a0)	# get the next source word
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100011 00100 00010 0000000000000000



Who Uses Assembly Language

- The machine designer
 - Must implement and trade off instruction functionality
- The compiler writer
 - Must generate machine language from a High Level Languages
- The writer of time or space critical code
 - Performance goals may force program-specific optimizations of the assembly language
- Special purpose or imbedded processor programmers
 - Special functions and heavy dependence on unique devices can make High Level Languages useless
- Students



MIPS Design Overview

All processing operands are in registers
to promote speed

Load and Store access to memory
load operand data value into register from memory
store operand data value from register to memory

Design goals: **Maximize performance,**
Minimize cost,
and Reduce design time



Operands in Registers

- To allow the MIPS processor to run faster
- MIPS Registers are 32 bits in size
 - Integer values are +/- 2 billion
 - Floating point values are +/- 1.70×10^{38}
 - 4 ASCII characters
- MIPS has 32 registers



Load and Store

- Move operand values from and to the main memory
- LOAD - copy value from memory to a register
 - LW – Load Word - copy all 32 bits to a register
- STORE - copy value to memory from a register
 - SW – Store Word - copy all 32 bits to memory



MIPS Assembly Instructions

- Language of this Machine
- More primitive than higher level languages
 - No sophisticated control flow: Branch and Jump only
- Very restrictive
 - MIPS Arithmetic Instructions
 - ADD \$t3, \$t1, \$t2 ($\$t3 = \$t1 + \$t2$)
- We'll be working with the MIPS instruction set architecture
 - similar to other architectures developed since the 1980's
 - used by NEC, Nintendo, Silicon Graphics, Sony

Design goals: **Maximize performance, Minimize cost, and Reduce design time**



MIPS Assembly Language

- Assembly language provides convenient symbolic representation
 - much easier than writing down numbers
 - destination first
- Machine language is the underlying reality
- Assembly language can provide 'pseudo-instructions'
 - “**move** \$t0, \$t1” exists only in Assembler
 - would be implemented using “**add** \$t0,\$t1,\$zero”
- When considering performance you should count real instructions



MIPS Assembly Arithmetic

- Instructions have 3 operands.
- Operand values are in registers.
- Operand order is fixed - destination first

Example:

C code: $A = B + C$

MIPS code: `add $s0, $s1, $s2`
 A B C

A red arrow originates from the third operand, `$s2`, in the MIPS code line and points diagonally upwards and to the left towards the first operand, `$s0`, illustrating the fixed operand order where the destination register is first.

(registers associated with variables by the compiler)



MIPS Assembly Arithmetic

- Design Principle: **Simplicity favors regularity.**
- Of course this complicates some things...

C code: $A = B + C + D;$
 $E = F - A;$

MIPS code: `add $t0, $s1, $s2`
 `add $s0, $t0, $s3`
 `sub $s4, $s5, $s0`

- Operands must be in registers, only 32 registers provided
- Design Principle: **Smaller is faster.**



Load and store Instructions

- Example:

C code: $A[8] = h + A[8];$

MIPS code: `lw $t0, 32($s3) ; content of A[8]`
 `add $t0, $s2, $t0; add h`
 `sw $t0, 32($s3) ; save to A[8]`

- Machine process: Memory address = `rs` + 16 bit value



Load and store Instructions

- Consider the load-word and store-word instructions,
 - What would the regularity principle have us do?
 - Design principle: **Good design demands a compromise**
- Example:

lw \$s1, 100(\$s2) $\$s1 = \text{Memory}[\$s2+100]$

sw \$s1, 100(\$s2) $\text{Memory}[\$s2+100] = \$s1$

- Where's the compromise?

Store word has destination last.



At this point:

- MIPS
 - Loading words but addressing bytes
 - Arithmetic operands in registers only

<u>Instruction</u>	<u>Meaning</u>
add \$s1, \$s2, \$s3	$\$s1 = \$s2 + \$s3$
sub \$s1, \$s2, \$s3	$\$s1 = \$s2 - \$s3$
lw \$s1, 100(\$s2)	$\$s1 = \text{Memory}[\$s2+100]$
sw \$s1, 100(\$s2)	$\text{Memory}[\$s2+100] = \$s1$



Summary

- People Languages
- MIPS Assembly Language
- All processing operands are in registers
- Load and Store access to memory

Next: More of the MIPS Assembly Language