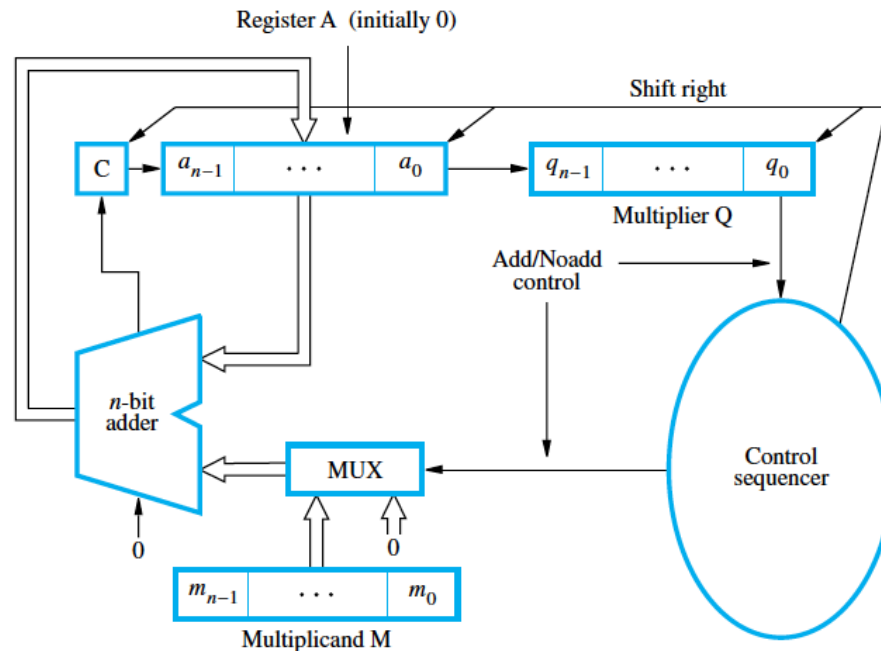


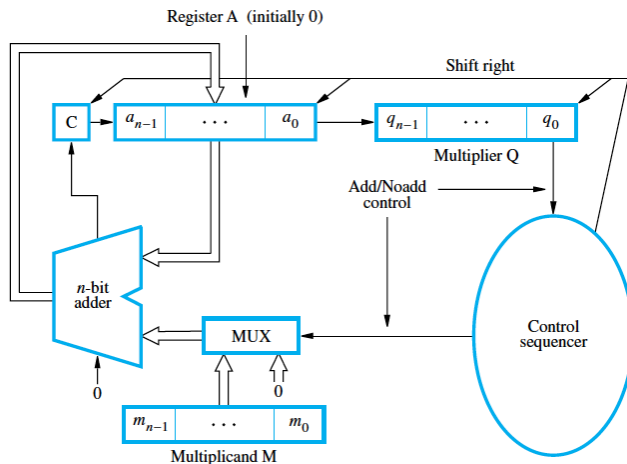
To multiply by  $2^n$ , just shift left  $n$  bits

In general, multiplication involves shifting and adding



The product of two  $n$ -bit numbers requires  $2n$  bits

Configuration for unsigned multiplication



4-bit numbers  $(-13 \times 11)$  unsigned multiply

|   |   |         |         |                            |
|---|---|---------|---------|----------------------------|
|   |   | M       |         |                            |
|   |   | 1 1 0 1 |         |                            |
|   |   | 0 0 0 0 | 1 0 1 1 | Initial configuration      |
| 0 | C | A       | Q       |                            |
| 0 |   | 1 1 0 1 | 1 0 1 1 | Add Shift } First cycle    |
| 0 |   | 0 1 1 0 | 1 1 0 1 |                            |
| 1 |   | 0 0 1 1 | 1 1 0 1 | Add Shift } Second cycle   |
| 0 |   | 1 0 0 1 | 1 1 1 0 |                            |
| 0 |   | 1 0 0 1 | 1 1 1 0 | No add Shift } Third cycle |
| 0 |   | 0 1 0 0 | 1 1 1 1 |                            |
| 1 |   | 0 0 0 1 | 1 1 1 1 | Add Shift } Fourth cycle   |
| 0 |   | 1 0 0 0 | 1 1 1 1 |                            |
|   |   | Product |         |                            |

Repeat N times (N= 4 here)

Add multiplicand to the A register only if the LSB of Q is 1

Shift for each cycle (C, A and Q together)

Generates an 8-bit product

Compute  $13 \times 11$  and negate the result to get -143



## A more general technique

Works for both positive and negative factors

Reduces the number of operations required

Example: suppose we want to multiply by  $30 = 0011110_2$   
30 can be regarded as

$$\begin{array}{r} 0100000 \quad (32) \\ - 0000010 \quad (2) \\ \hline 0011110 \quad (30) \end{array} \quad \text{i.e., multiply by } (32 - 2)$$

Add 32 x multiplicand plus -2 x multiplicand

| Multiplier |             | Version of multiplicand<br>selected by bit $i$ |
|------------|-------------|--|
| Bit $i$    | Bit $i - 1$ |  |
| 0          | 0           | $0 \times M$                                   |
| 0          | 1           | $+1 \times M$                                  |
| 1          | 0           | $-1 \times M$                                  |
| 1          | 1           | $0 \times M$                                   |

Scan multiplier bits from right to left

(assume initial 0 bit to right of LSB)

subtract multiplicand when moving from 0 to 1

add multiplicand when moving from 1 to 0

neither add nor subtract when moving from 0 to 0 or 1 to 1

shift multiplicand left one bit for each cycle



The original multiply is “*recoded*”

$30 = 0011110_2$  using 7 bits

0 0 1 1 1 1 1 0



0 +1 0 0 0 0 -1 0

Subtract multiplicand for each -1 in recoded multiplier

Add multiplicand for each +1 in recoded multiplier

Neither add nor subtract for each 0 in recoded multiplier


45 → 0101101

-45 → 1010011

|   |   |   |   |   |   |   |   |       |   |   |   |   |   |   |                                      |                         |                      |
|---|---|---|---|---|---|---|---|-------|---|---|---|---|---|---|--------------------------------------|-------------------------|----------------------|
|   |   |   |   |   |   |   |   | 0     | 1 | 0 | 1 | 1 | 0 | 1 | ← multiplicand                       |                         |                      |
|   |   |   |   |   |   |   |   | 0     | + | 1 | 0 | 0 | 0 | - | 1                                    | 0                       | ← Recoded multiplier |
|   |   |   |   |   |   |   |   | <hr/> |   |   |   |   |   |   |                                      |                         |                      |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0     | 0 | 0 | 0 | 0 | 0 | 0 |                                      |                         |                      |
| 1 | 1 | 1 | 1 | 1 | 1 | 1 | 0 | 1     | 0 | 0 | 1 | 1 |   |   | ← 2's complement of the multiplicand |                         |                      |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0     | 0 | 0 | 0 | 0 |   |   |                                      | Sign extended           |                      |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0     | 0 | 0 | 0 |   |   |   |                                      |                         |                      |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0     | 0 | 0 |   |   |   |   |                                      |                         |                      |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0     | 0 |   |   |   |   |   |                                      |                         |                      |
| 0 | 0 | 0 | 1 | 0 | 1 | 1 | 0 | 1     |   |   |   |   |   |   |                                      |                         |                      |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |       |   |   |   |   |   |   |                                      |                         |                      |
|   |   |   |   |   |   |   |   | <hr/> |   |   |   |   |   |   |                                      |                         |                      |
| 0 | 0 | 0 | 1 | 0 | 1 | 0 | 1 | 0     | 0 | 0 | 1 | 1 | 0 |   | →                                    | 0x0546 = 1350 = 45 x 30 |                      |

One addition and one subtraction is required to generate the product  
 Using 00111110 would have required 5 additions

Assume a 17-bit multiplier:

|   |    |    |    |   |    |   |    |   |   |    |    |    |    |   |    |   |   |  |  |
|---|----|----|----|---|----|---|----|---|---|----|----|----|----|---|----|---|---|--|--|
| 0 | 0  | 1  | 0  | 1 | 1  | 0 | 0  | 1 | 1   | 1  | 0  | 1  | 0  | 1 | 1  | 0 | 0 |  |  |
|   |    |    |    |   |    |   |    |   |  |    |    |    |    |   |    |   |   |  |  |
| 0 | +1 | -1 | +1 | 0 | -1 | 0 | +1 | 0 | 0   | -1 | +1 | -1 | +1 | 0 | -1 | 0 | 0 |  |  |