2021SCUCTF新生赛-逆向题WriteUp

时间: 2021/11/19-2021/11/23

2021SCUCTF新生赛-逆向题WriteUp

[RE1]Welcome

[RE2]Xor

[RE3]DebugMe

[RE4]upxed

[RE5]ezpak

[RE6]EzLinearEquation

[RE7]VirtualMachine

解法一:分析虚拟机解法二:angr梭哈

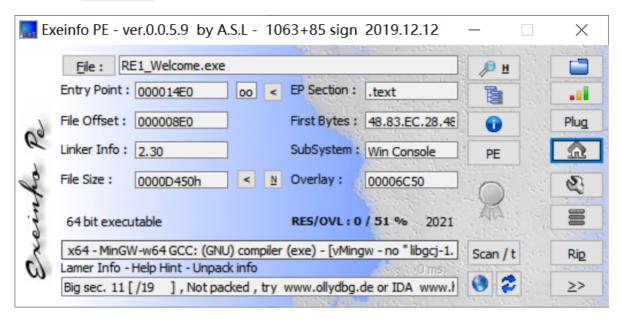
[RE8]PyCode

[RE9]Plants_Vs_Zombies

[HappyRE]ezjava

[RE1]Welcome

将文件拖入 Exeinfo PE 检查文件格式和加壳情况。64位PE文件,没有加壳:



拖入IDA后查看伪代码,分析程序逻辑,可以得到,只需要输入1919810即可输出正确的flag:

```
1int __cdecl main(int argc, const char **argv, const char **envp)
  2{
     int v4[3]; // [rsp+2Ch] [rbp-Ch] BYREF
     _main(argc, argv, envp);
  6 puts("Input a correct number to rescue the Fuhrer.");
    if ( scanf("%d", v4) != 1 )
     puts("??");
10
      exit(0);
 11
    if ( v4[0] != 1919810 )
12
 13 {
14
      puts("Wrong!");
15
      exit(0);
17 atexit(getflag);
18 return 0;
19}
```

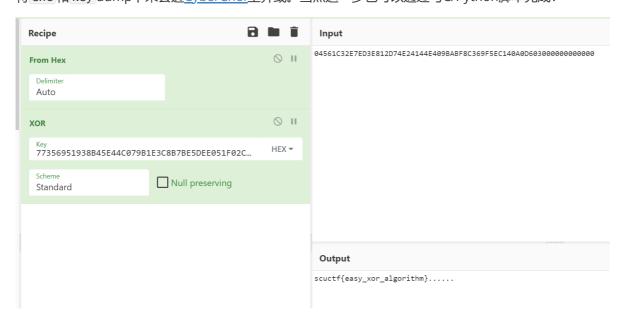
```
D:\CTF\SCUCTF-XSS-2021\Review\Release\1>RE1_Welcome.exe
Input a correct number to rescue the Fuhrer.
1919810
Correct.Flag to enter Fuhrerbunker:
scuctf{w31c0m3_to_the_world_of_r3!}
```

[RE2]Xor

直接定位关键代码。中间有一段很奇怪的代码我们可能不用去仔细分析,而是通过关键代码猜测其作用。第一个if判断一个很复杂的式子是不是等于26,结合下面for循环的次数也是26可以确定第一个if是判断输入的长度是否为26。第二个if指的是将输入和key异或之后与密文enc比较,比对通过则说明输入为正确的flag:

```
_main(argc, argv, envp);
      puts("Input your flag, and i will check it:");
scanf("%40s", v8);
      v8[39] = 0;
 13
      v3 = v8;
  15
      {
        v4 = *(DWORD *)v3;
        v3 += 4;
        v5 = \sim v4 \& (v4 - 16843009) \& 0 \times 8080808080;
  19
20
      while ( !v5 );
      if ( (~v4 & (v4 - 16843009) & 0x8080) == 0 )
 21
        v5 >>= 16;
      if ( (~v4 & (v4 - 16843009) & 0x8080) == 0 )
2
      if ( &v3[-__CFADD__((_BYTE)v5, (_BYTE)v5) - 3] - v8 != 26 )
        puts("Nope.");
28
        exit(1);
  29
      for ( i = 0i64; i != 26; ++i )
        if ( (key[i] ^ (unsigned __int8)v8[i]) != enc[i] )
 32
           puts("Nope.");
 34
          exit(0);
      puts("Correct! Input is your flag.");
```

将 enc 和 key dump下来丢进CyberChef里异或。当然这一步也可以通过写C/Python脚本完成:



[RE3]DebugMe

这题输入一个正确的数字 v22 即可输出flag,但是 v22 的生成过程非常复杂,手动计算很困难,所以这题我们可以考虑动态调试:

首先你需要一个Linux环境,比如说Ubuntu虚拟机。并在上面运行IDA的linux server64:

```
root@ubuntu:~# cd re/
root@ubuntu:~/re# ls

10 11 12 access.log capstone chacha20 ###### example-1_recovered exp.py gghdl xiangyunbei
root@ubuntu:~/re# cd dbgsrv/
root@ubuntu:~/re/dbgsrv# ls
21111961978038e80fb admin_panel babyre babyre.orig csb linux_server linux_server64 ps sob
root@ubuntu:~/re/dbgsrv# ./linux_server64
IDA Linux 64-bit remote debug server(ST) v7.5.26. Hex-Rays (c) 2004–2020
Listening on 0.0.0.0:23946...
```

搭建好动态调试环境之后就可以通过动调来查看v22的值了,是0x5AD,即1453:

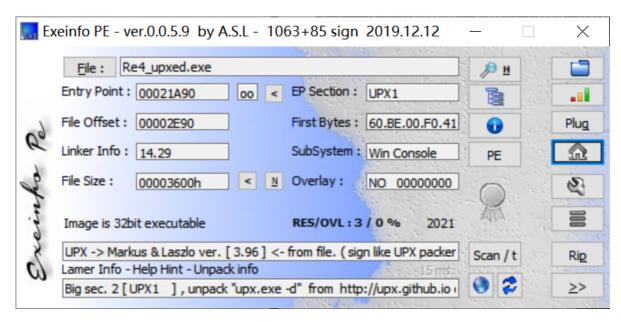
```
| Description |
```

将输入改为1453即可得到正确的flag:

```
Correct! Here is your flag:
scuctf{just_ask_you_to_have_Linux}
```

[RE4]upxed

查壳,发现有upx壳:



这题说是不能直接脱壳,但直接脱壳又确实可以,挺奇怪的:

用upx直接脱壳:

```
D:\CTF\SCUCTF-XSS-2021\Review\Release\4>upx -d Re4_upxed.exe -o dump_.exe

Ultimate Packer for eXecutables

Copyright (C) 1996 - 2020

UPX 3.96w Markus Oberhumer, Laszlo Molnar & John Reiser Jan 23rd 2020

File size Ratio Format Name

41984 <- 13824 32.93% win32/pe dump_.exe

Unpacked 1 file.
```

找到关键代码,应该是把输入加密之后与密文进行比较:

```
qmemcpy(v8, "THIS_IS_THE_KEYa", 16);
• 10
• 11
       v6 = 0x5618435D;
      v7[0] = -1608884855;
v7[1] = 1372647817;
v7[2] = 1872257453;
v7[3] = 835943406;
• 12
•
•
• 15
• 16
      v7[4] = -1137280198:
      puts(Buffer);
puts(off_41A004);
puts(off_41A008);
17
•
•
20 puts(off_41A00C);
21
     puts(off_41A010);
22
     j_memset(Str, 0, 0x64u);
23 puts("Input your flag:");24 sub_41127B("%100s", (char)Str);
• 25 if ( j_strlen((const char *)Str) != 24 )
27
         puts("Wrong!");
28
         exit(0);
30
         sub_41100A(&Str[i], v8);
32
                                       (&v6 + 2 * i) || HIDWORD(Str[i]) != v7[2 * i] )
            puts("Wrong!");
35
36
            exit(0);
      puts("Correct! flag is your input!");
39
• 40
41}
    00001142 main 0:32 (411D42)
```

这样看起来还有点抽象,稍作些修改,可以看到密文很明显就是v6了:

```
QWORD v6[4]; // [esp+154h] [ebp-38h]
int v7[5]; // [esp+174h] [ebp-18h] BYREF

CheckForDebuggerJustMyCode(&unk_41C015);
;memcpy(v7, "IHIS_IS_IHE_KEYa", 16);
ODWORD(v6[0]) = 0x5618435D;
1 HIDWORD(v6[0]) = 0x51087189;
1 HIDWORD(v6[1]) = 0x51087189;
1 HIDWORD(v6[1]) = 0x51087189;
1 HIDWORD(v6[1]) = 0x31037BEE;
1 HIDWORD(v6[2]) = 0x31037BEE;
1 HIDWORD(v6[2]) = 0x8C367B3A;
toto(suffean);
puts(off_41A004);
puts(off_41A000);
puts(off_41A010);
j_memset(Str, 0, 0x64u);
j_uts(off_41A010);
j_i_memset(Str, 0, 0x64u);
if (j_strlen((const char *)Str) != 24 )
{
   puts("Mrong!");
   exit(0);
}

for (i = 0; i < 3; ++i )
{
   sub_41100AR((int)&Str[i], (int)v7);
   if ( LODWORD(Str[i]) != HIDWORD(v6[i]) != HIDWORD(v6[i]) )</pre>
```

加密是很明显的TEA,密钥是之前就找到的"THIS_IS_THE_KEYa":

```
2{
     int result; // eax
  4 unsigned int i; // [esp+10Ch] [ebp-2Ch]
  5 int v4; // [esp+118h] [ebp-20h]
     unsigned int v5; // [esp+124h] [ebp-14h] unsigned int v6; // [esp+130h] [ebp-8h]
     CheckForDebuggerJustMyCode(&unk_41C015);
10 v6 = *a1;
\bullet 11 v5 = a1[1];
     v4 = 0;
13 for ( i = 0; i < 0x20; ++i )</p>
 14 {
15
       v4 -= 0x61C88647;
       v6 += (a2[1] + (v5 >> 5)) ^ (v4 + v5) ^ (*a2 + 16 * v5);
16
17
      v5 += (a2[3] + (v6 >> 5)) ^ (v4 + v6) ^ (a2[2] + 16 * v6);
19
     *a1 = v6;
20 result = 4;
21 a1[1] = v5;
22 return result;
23 }
```

直接去WikiPedia抄一份TEA源码,dump密文和密钥进行解密即可:

```
#include <stdint.h>
#include <cstdio>
void encrypt (uint32_t* v, uint32_t* k) {
    uint32_t v0=v[0], v1=v[1], sum=0, i;
                                                  /* set up */
                                                    /* a key schedule constant */
    uint32_t delta=0x9e3779b9;
    uint32_t k0=k[0], k1=k[1], k2=k[2], k3=k[3]; /* cache key */
    for (i=0; i < 32; i++) {
                                                     /* basic cycle start */
        sum += delta;
        v0 += ((v1 << 4) + k0) \land (v1 + sum) \land ((v1 >> 5) + k1);
        v1 += ((v0 << 4) + k2) \land (v0 + sum) \land ((v0 >> 5) + k3);
                                                     /* end cycle */
    v[0]=v0; v[1]=v1;
}
void decrypt (uint32_t* v, uint32_t* k) {
    uint32_t v0=v[0], v1=v[1], sum=0xC6EF3720, i; /* set up */
    uint32_t delta=0x9e3779b9;
                                                    /* a key schedule constant */
    uint32_t k0=k[0], k1=k[1], k2=k[2], k3=k[3]; /* cache key */
    for (i=0; i<32; i++) {
                                                     /* basic cycle start */
        v1 = ((v0 << 4) + k2) \land (v0 + sum) \land ((v0 >> 5) + k3);
        v0 = ((v1 << 4) + k0) \land (v1 + sum) \land ((v1 >> 5) + k1);
        sum -= delta;
                                                     /* end cycle */
    v[0]=v0; v[1]=v1;
}
int main(){
    uint32_t enc[6];
    enc[0] = 0x5618435D;
    enc[1] = 0xA01A5D89;
    enc[2] = 0x51D0F189;
```

```
enc[3] = 0x6F9861AD;
enc[4] = 0x31D37BEE;
enc[5] = 0xBC367B3A;
uint8_t key[20] = "THIS_IS_THE_KEYA";
for(int i = 0;i < 3;i ++){
    decrypt(enc + 2 * i, (uint32_t*)key);
}
puts((char*)enc);
}</pre>
```

输出:

```
scuctf{Upx_with_TEA_Alg}
```

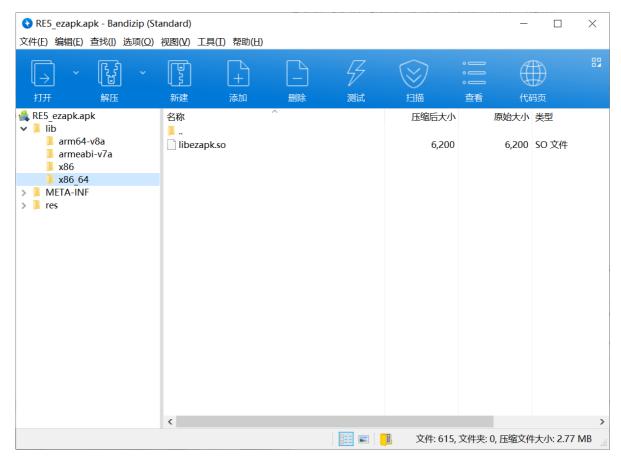
当然也可以手动脱壳,CTF-Wiki的脱壳教程: https://ctf-wiki.org/reverse/windows/unpack/esp/

[RE5]ezpak

APK文件可以用jadx打开,找到MainActivity类。可以看到这里是调用了一个native函数 check ,native函数是用C语言写的,编译成共享库(后缀为.so的文件)后塞到APK里供Java代码调用:

```
static {
35
           System.loadLibrary("ezapk");
}
      /* access modifiers changed from: protected */
43p
     public void onCreate(Bundle bundle) {
          super.onCreate(bundle);
           setContentView((int) R.layout.activity_main);
          this.password = (EditText) findViewById(R.id.password);
         this.result = (TextView) findViewById(R.id.result);
          Button button = (Button) findViewById(R.id.submit);
48
          this.submit = button;
49
          button.setOnClickListener(new MainActivity$$ExternalSyntheticLambda@(this));
}
      /* renamed from: lambda$onCreate$0$cn-bluesadi-ezapk-MainActivity reason: not valid java name */
49p
                                                               pdiezapkMainActivity(View view) {
           String obj = this.password.getText().toString();
50
51
           if (check(obj))
                                                   uccess:\nscuctf{%s}", new Object[]{MD5(obj)}));
52
               tnis.result.setlext(String.format(
              this.result.setTextColor(getColor(R.color.green));
53
56
              return:
           this.result.setText("Wrong password.");
55
56
           this.result.setTextColor(getColor(R.color.red));
```

APK本质上是一个压缩包,我们用解压缩软件打开APK,可以看到lib文件夹内有各种架构下的共享库文件:



不同架构的so文件反编译结果会有差异,可以每种架构都试一遍,选最容易看的一种。这里我选的是x86_64架构,先还原参数类型(非ARM架构需要手动导入jni.h头文件):

```
IDA View-A A Ham Pseudocode-A A Ham Hex View-1 A MA Structures A Ham

1 bool __fastcall Java_cn_bluesadi_ezapk_MainActivity_check(JNIEnv *a1, jobject a2, jstring a3)

2 {
3 const char *v3; // rbx
4 m128i v4: // xmm1
```

第一步是将输入转存到buffer,注意这里输入的字符顺序发生了变化:

```
v8 = __readfsqword(0x28u);
      input = (*a1)->GetStringUTFChars(a1, a3, &v6);
• 10
• 11
     if ( strlen(input) != 22 )
12
       return 0;
13
      *(_OWORD *)buffer = 0LL;
14
      *(_QWORD *)&buffer[14] = 0LL;
      buffer[0] = *input;
15
16
      buffer[2] = input[1];
17
      buffer[4] = input[2];
18
      buffer[6] = input[3];
      buffer[8] = input[4];
19
20
      buffer[10] = input[5];
21
      buffer[12] = input[6];
22
      *(_QWORD *)&buffer[14] = *((unsigned __int8 *)input + 7);
     *(_DWORD *)&buffer[16] = *((unsigned __int8 *)input + 8);
23
      *(_WORD *)&buffer[18] = *((unsigned __int8 *)input + 9);
*(_WORD *)&buffer[20] = *((unsigned __int8 *)input + 10);
24
25
26
      buffer[1] = input[11];
27
      buffer[3] = input[12];
28
      buffer[5] = input[13];
• 29
      buffer[7] = input[14];
30
      buffer[9] = input[15];
31
      buffer[11] = input[16];
32
      buffer[13] = input[17];
      buffer[15] = input[18];
33
34
      buffer[17] = input[19];
35
      buffer[19] = input[20];
      buffer[21] = input[21];
36
```

这一部分很容易总结出规律,可以看做是把输入分成两个部分,然后穿插:

```
for(int i = 0;i < 2;i += 1){
   for(int j = 0;j < 22;j += 2){
      buf[i + j] = input[ptr++];
   }
}</pre>
```

后面则是一个异或加密和比较的过程:

虽然比较零散,但可以看出来规律,归纳出来是一个这样的加密过程:

```
for(int i = 0;i < 22;i ++){
   buf[i] += i;
   buf[i] ^= i;
}</pre>
```

写出exp:

```
#include <cstdio>
int main(){
    unsigned char enc[23] =
"\x79\x72\x73\x71\x7d\x6b\x63\x6c\x61\x61\x76\x79\x7d\x7f\x63\x72\x61\x6b\x92\x9
b\x6c\x9d";
    for(int i = 0; i < 22; i ++){
        enc[i] \wedge= i;
        enc[i] -= i;
    }
    int ptr = 0;
    unsigned char passwd[23] = {0};
    for(int i = 0; i < 2; i += 1){
        for(int j = 0; j < 22; j += 2){
            passwd[ptr++] = enc[i + j];
    }
    printf("%s\n", passwd);
}
```

[RE6]EzLinearEquation

这一段本质上是构造了一个线性方程组,参数存储在mat数组里,这个线性方程组的解即是flag:

```
scanf("%40s", Str);
if ( strlen(Str) == 27 && !strncmp(Str, "scuctf{", 7ui64) && BYTE2(v7) == 125 )

{
    v9 = &Str[7];
    v12 = 0i64;
    for ( i = 0; i <= 18; ++i )
    {
        v12 = 0i64;
        for ( j = 0; j <= 18; ++j )
            v12 += mat[19 * i + j] * (unsigned __int8)v9[j];
        if ( v12 != enc[i] )
            goto LABEL_10;
    }
    puts("Correct!");
    result = 0;
}
else
{</pre>
```

求解方程组使用的库是z3,我之前写过一个<u>教程</u>。dump参数和密文,编写Python脚本求解方程即可:

```
from z3 import *
```

```
mat = [0x0000000000000010, 0x00000000000051, 0x00000000000000002,
0x00000000000032, 0x000000000001B, 0x0000000000064, 0x00000000034,
0x0000000000003E, 0x0000000000002E, 0x0000000000017, 0x00000000000A,
0x00000000000062, 0x000000000000040, 0x000000000003F, 0x00000000005D,
0x0000000000002E, 0x00000000000013, 0x0000000000049, 0x000000000022,
0 \times 000000000000038, 0 \times 00000000000000000, 0 \times 0000000000000032, 0 \times 000000000000043,
0x0000000000002E, 0x00000000000002A, 0x000000000002D, 0x00000000004A,
0 \times 000000000000040, 0 \times 0000000000000003, 0 \times 00000000000004, 0 \times 0000000000000000,
0 \times 000000000000056, 0 \times 00000000000000007, 0 \times 000000000000062, 0 \times 000000000000004B,
0x00000000000000, 0x000000000002f, 0x0000000000024, 0x00000000043,
0x000000000005c, 0x000000000004A, 0x000000000003A, 0x00000000001c,
0 \times 000000000000030, 0 \times 000000000000000025, 0 \times 000000000000022, 0 \times 00000000000000005,
0 \times 000000000000049, 0 \times 000000000000005A, 0 \times 0000000000045, 0 \times 00000000000005,
0x00000000000026, 0x000000000000000000, 0x00000000001B, 0x000000000001c,
0x0000000000003B, 0x0000000000001F, 0x000000000005E, 0x000000000051,
0x0000000000000, 0x0000000000003c, 0x000000000001E, 0x000000000047,
0x00000000000039, 0x00000000000021, 0x000000000005E, 0x00000000002F,
0 \times 000000000000049, 0 \times 00000000000001A, 0 \times 0000000000058, 0 \times 000000000000000
0 \times 00000000000003c, 0 \times 000000000000019, 0 \times 000000000000017, 0 \times 00000000000036,
0x0000000000002f, 0x00000000000016, 0x0000000000052, 0x00000000004B,
0 \times 00000000000003, 0 \times 00000000000004, 0 \times 00000000000002, 0 \times 00000000000000,
0x000000000001E, 0x000000000000040, 0x000000000003B, 0x0000000000000
0 \times 0 0 0 0 0 0 0 0 0 0 0 0 2 F, 0 \times 0 0 0 0 0 0 0 0 0 0 0 1, 0 \times 0 0 0 0 0 0 0 0 0 0 13, 0 \times 0 0 0 0 0 0 0 0 0 0 3,
0x0000000000002D, 0x00000000000023, 0x000000000001D, 0x000000000051,
0 \times 0000000000000047, 0 \times 00000000000000058, 0 \times 00000000000035, 0 \times 000000000000039,
0 \times 000000000000060, 0 \times 00000000000000000, 0 \times 00000000000006, 0 \times 0000000000000000
0x0000000000005E, 0x0000000000034, 0x0000000000028, 0x0000000001B,
0 \times 000000000000017, 0 \times 0000000000000036, 0 \times 00000000000005, 0 \times 000000000000048,
0x0000000000005B, 0x000000000000D, 0x000000000004E, 0x000000000031,
0 \times 00000000000003B, 0 \times 0000000000000005D, 0 \times 0000000000003F, 0 \times 00000000000000A,
0 \times 000000000000019, 0 \times 0000000000000042, 0 \times 000000000000051, 0 \times 000000000000034,
0x0000000000004F, 0x0000000000004c, 0x000000000003E, 0x00000000016,
0 \times 00000000000012, 0 \times 000000000000019, 0 \times 00000000000003, 0 \times 0000000000000032,
0 \times 000000000000046, 0 \times 0000000000000023, 0 \times 0000000000000, 0 \times 000000000000000
0 \times 000000000000049, 0 \times 000000000000004, 0 \times 00000000000003, 0 \times 00000000000008,
0x00000000000064, 0x00000000000030, 0x0000000000002F, 0x000000000037,
0x000000000000025, 0x000000000001E, 0x000000000005A, 0x000000000017,
0 \times 000000000000004, 0 \times 00000000000004F, 0 \times 000000000000032, 0 \times 00000000000000D,
0x000000000001D, 0x0000000000053, 0x000000000003c, 0x00000000005,
0 \times 000000000000055, 0 \times 000000000000002F, 0 \times 00000000000055, 0 \times 000000000000017,
0 \times 000000000000003, 0 \times 000000000000004A, 0 \times 000000000000026, 0 \times 0000000000000033,
0 \times 000000000000004D, 0 \times 000000000000005A, 0 \times 000000000000029, 0 \times 000000000000036,
0 \times 00000000000037, 0 \times 000000000000031, 0 \times 000000000000059, 0 \times 00000000000004c,
0x000000000005A, 0x000000000001c, 0x0000000000055, 0x000000000052,
0 \times 00000000000005A, 0 \times 000000000000002B, 0 \times 00000000000017, 0 \times 0000000000000048,
```

```
0 \times 00000000000004B, 0 \times 0000000000000000, 0 \times 000000000000018, 0 \times 00000000000000
0 \times 0000000000000046, 0 \times 000000000000000, 0 \times 00000000000039, 0 \times 0000000000000042,
0 \times 00000000000003E, 0 \times 0000000000000061, 0 \times 0000000000000042, 0 \times 000000000000006,
0 \times 00000000000004E, 0 \times 000000000000038, 0 \times 00000000000002A, 0 \times 00000000000058,
0x0000000000003B, 0x00000000000051, 0x000000000001E, 0x0000000001A,
0 \times 000000000000038, 0 \times 000000000000005, 0 \times 00000000000017, 0 \times 00000000000000,
0x00000000000003, 0x000000000000057, 0x0000000000010, 0x000000000002E,
0x00000000000052, 0x00000000000017, 0x000000000005B, 0x000000000042,
0 \times 000000000000014, 0 \times 00000000000000036, 0 \times 00000000000004, 0 \times 000000000000019,
0 \times 000000000000030, 0 \times 000000000000003c, 0 \times 00000000000063, 0 \times 000000000000047,
0 \times 00000000000002c, 0 \times 00000000000000026, 0 \times 000000000000042, 0 \times 0000000000000019,
0 \times 00000000000003c, 0 \times 0000000000000005, 0 \times 000000000000044, 0 \times 00000000000006,
0 \times 00000000000004A, 0 \times 00000000000000051, 0 \times 0000000000000062, 0 \times 000000000000033,
0 \times 000000000000036, 0 \times 00000000000000058, 0 \times 0000000000038, 0 \times 0000000000000042,
0x00000000000011, 0x00000000000014, 0x0000000000004B, 0x00000000005A,
0 \times 000000000000032, 0 \times 000000000000001D, 0 \times 0000000000055, 0 \times 00000000000000D,
0 \times 000000000000052, 0 \times 0000000000000034, 0 \times 000000000000060, 0 \times 00000000000001E,
0 \times 000000000000040, 0 \times 000000000000004B, 0 \times 0000000000011, 0 \times 0000000000000009,
0 \times 000000000000046, 0 \times 0000000000000056, 0 \times 0000000000000, 0 \times 000000000000
0 \times 0000000000000040, 0 \times 000000000000000, 0 \times 000000000000004, 0 \times 000000000000004,
0 \times 000000000000005, 0 \times 00000000000000057, 0 \times 00000000000008, 0 \times 0000000000000035,
0x0000000000002A, 0x00000000000012, 0x000000000002E, 0x00000000005B,
0 \times 000000000000015, 0 \times 000000000000002A, 0 \times 00000000000046, 0 \times 00000000000000000,
0 \times 000000000000031, 0 \times 000000000000000000, 0 \times 00000000000034, 0 \times 000000000000018,
0 \times 000000000000047, 0 \times 000000000000005, 0 \times 00000000000003, 0 \times 000000000003,
0 \times 000000000000021, 0 \times 0000000000000056, 0 \times 00000000000014, 0 \times 000000000000004,
0x0000000000053, 0x000000000005F, 0x0000000000004C, 0x0000000000040,
0 \times 000000000000044, 0 \times 000000000000005F, 0 \times 0000000000000, 0 \times 0000000000000000,
0x000000000000000001
enc = [0x00000000001424A, 0x0000000001AA02, 0x00000000001962c,
0x00000000013FF4, 0x0000000000183E3, 0x000000000016B57, 0x000000000017BEC,
0x000000000133E2, 0x0000000000162AA, 0x000000000195AC, 0x000000000014AF1,
0x0000000001d1d0, 0x000000000019EE2, 0x000000000019135, 0x000000000014E8B,
0x00000000019ADF, 0x000000000015600, 0x000000000013D5A, 0x00000000001827F,
0x00000000000000001
X = [Int('x\%d'\% i) \text{ for } i \text{ in } range(19)]
solver = Solver()
for i in range(19):
    equation = 0
    for j in range(19):
        equation += mat[i * 19 + j] * X[j]
    solver.add(equation == enc[i])
print(solver.check())
model = solver.model()
print(''.join([chr(model[X[i]].as_long()) for i in range(19)]))
```

[RE7]VirtualMachine

解法一: 分析虚拟机

一道典型的虚拟机逆向题,虽然没有设置什么陷阱或者与其他技术综合,但还是有一定难度,需要有一 定的汇编基础以及对操作系统的抽象架构有所了解。

首先是VM的初始化:

```
1int cdecl VM init(void **a1, int a2)
  2{
     DWORD *v3; // [esp+D0h] [ebp-8h]
      CheckForDebuggerJustMyCode(&unk 41D02F);
     *a1 = malloc(0x2Cu);
  7 if (!*a1)
  8 {
       puts("Error while mallocing for VM");
• 9
10
       exit(1);
 11
12 v3 = *a1;
\bullet 13 *v3 = 0:
014 v3[1] = 1;
015 v3[2] = 2;
\bullet 16 v3[3] = 3;
017 v3[4] = a2;
\bullet 18 v3[5] = a2;
19 v3[6] = malloc(0x1000u);
0 20 if (!v3[6])
 21
       puts("Error while mallocing for SP");
22
23
       exit(1);
 24
0 25 v3[7] = malloc(0x100u);
0 26 if (!v3[7])
 27
       puts("Error while mallocing for PWRITE");
28
29
       exit(1);
 30
\bullet 31 v3[8] = v3[7];
32 return 0;
33 }
```

然后是VM正式开始运行,从红框中可以看出这题的VM使用了结构体,所以逆向起来会更加复杂,为了方便分析VM架构以及dump出VM指令,需要首先还原VM结构体:

```
1void __cdecl VM_start(int a1)
       int v1; // ecx
      int v2; // ecx
      int i; // [esp+D4h] [ebp-44h]
      int v4; // [esp+E0h] [ebp-38h]
    7 unsigned __int8 v5; // [esp+113h] [ebp-5h]
8 unsigned __int8 v6; // [esp+113h] [ebp-5h]
        __CheckForDebuggerJustMyCode(&unk 41D02F);
 11 \quad v4 = 0;
      while (1)
• 12
         switch ( **( BYTE **)(a1 + 16) )
           case 0xA:
             sub 411398(
               a1 + 4 * *(unsigned __int8 *)(*(_DWORD *)(a1 + 16) + 1),
            a1 + 4 * *(unsigned __int8 *)(*(_DWORD *)(a1 + 16) + 2));
           *(_DWORD *)(a1 + 16) += 3;
break;
           case 0xB:
            sub_41137F(
               a1 + 4 * *(unsigned __int8 *)(*(_DWORD *)(a1 + 16) + 1),
               a1 + 4 * *(unsigned __int8 *)(*(_DWORD *)(a1 + 16) + 2));
           *(_DWORD *)(a1 + 16) += 3;
break;
26
          case 0xC:
            sub_411069(
              a1 + 4 * *(unsigned __int8 *)(*(_DWORD *)(a1 + 16) + 1),
               a1 + 4 * *(unsigned __int8 *)(*(_DWORD *)(a1 + 16) + 2));
             *(DWORD *)(a1 + 16) += 3;
             break;
```

还原结构体的方法是在 Structures 中插入自定义的新结构体,并手动设置每个结构体变量的大小、类型等,这里我们自定义了一个VM结构体。还原需要根据代码推断结构体变量的作用,具体的过程就不细说了,总之是个苦差事,需要耐心和时间:

```
X
                   × o
                                        × A
                             Hex View-1
       Pseudocode-A
                                                Structures
00000000 ; Ins/Del : create/delete structure
00000000 ; D/A/* : create structure member (data/ascii/array)
00000000 ; N
                   : rename structure or structure member
00000000 ; U : delete structure member
00000000 ; [00000010 BYTES. COLLAPSED STRUCT GUID. PRESS CTRL-NUMPAD+
0000000
0000000 O VM
                          struc ; (sizeof=0x2C, mappedto_51)
000000000 regs
                          dd 4 dup(?)
0000001<mark>0 PC</mark>
                          dd?
                                                    ; offset
00000014 PC base
                          dd?
00000018 sp_
0000001C mem
                          dd?
                                                    ; offset
                          dd?
00000010 mem_ptr
                          dd ?
00000014 len
                          dd ?
                          dd?
0000002<mark>8 flag</mark>
0000002C VM
                           ends
999999
```

还原结构体后的代码如图所示,还是比较清晰的:

```
case 0xFu:
         rsh((int)&a1->regs[a1->PC[1]], (int)&a1->regs[a1->PC[2]]);
         a1->PC += 3;
         break;
       case 0x10u:
         rol((int)&a1->regs[a1->PC[1]], (int)&a1->regs[a1->PC[2]]);
         a1->PC += 3;
         break;
       case 0x11u:
         ror((int)&a1->regs[a1->PC[1]], (int)&a1->regs[a1->PC[2]]);
         a1->PC += 3;
         break;
       case 0x12u:
         a1->regs[a1->PC[1]] = a1->regs[a1->PC[2]];
         a1->PC += 3;
         break:
       case 0x13u:
         a1->regs[a1->PC[1]] = _byteswap_ulong(*(_DWORD *)(a1->PC + 2));// mov_i2r
         a1->PC += 6;
         break;
       case 0x14u:
         *a1->sp_ = a1->regs[a1->PC[1]];
        a1->sp_ += 4;
         a1->PC += 2;
50
         break;
          *a1->sp_ = _byteswap_ulong(*(_DWORD *)(a1->PC + 1));
         a1->sp_ += 4;
a1->PC += 5;
         break;
       case 0x16u:
         v5 = a1->PC[1];
```

分析出每个opcode对应的功能后,dump字节码为一种类似汇编的形式,dump脚本如下:

```
#include <cstdio>
#include <cstdint>
#include <stdlib.h>
#define READ_WORD(PC) ((code[PC] << 8) | code[PC + 1])</pre>
#define READ_DWORD(PC) __builtin_bswap32(*(uint32_t*)&code[PC])
uint8_t code[1176] = {
    0x23, 0x00, 0x65, 0x1B, 0x1E, 0x00, 0x13, 0x00, 0x00, 0x00, 0x00, 0x0A,
0x1A, 0x00, 0x1C, 0x01,
    0x13, 0x00, 0x00, 0x00, 0x00, 0x57, 0x1A, 0x00, 0x1C, 0x01, 0x13, 0x00,
0x00, 0x00, 0x00, 0x72,
    0x1A, 0x00, 0x1C, 0x01, 0x13, 0x00, 0x00, 0x00, 0x00, 0x6F, 0x1A, 0x00,
0x1c, 0x01, 0x13, 0x00,
    0x00, 0x00, 0x00, 0x6E, 0x1A, 0x00, 0x1C, 0x01, 0x13, 0x00, 0x00, 0x00,
0x00, 0x67, 0x1A, 0x00,
    0x1c, 0x01, 0x13, 0x00, 0x00, 0x00, 0x00, 0x21, 0x1A, 0x00, 0x1c, 0x01,
0x13, 0x00, 0x00, 0x00,
    0x00, 0x20, 0x1A, 0x00, 0x1C, 0x01, 0x13, 0x00, 0x00, 0x00, 0x00, 0x0A,
0x1A, 0x00, 0x1C, 0x01,
    0x1E, 0x00, 0x18, 0x08, 0xFF, 0x13, 0x00, 0x00, 0x00, 0x00, 0x48, 0x1A,
0x00, 0x1C, 0x01, 0x13,
    0x00, 0x00, 0x00, 0x00, 0x65, 0x1A, 0x00, 0x1C, 0x01, 0x13, 0x00, 0x00,
0x00, 0x00, 0x6C, 0x1A,
    0x00, 0x1c, 0x01, 0x13, 0x00, 0x00, 0x00, 0x00, 0x6c, 0x1A, 0x00, 0x1c,
0x01, 0x13, 0x00, 0x00,
    0x00, 0x00, 0x6F, 0x1A, 0x00, 0x1C, 0x01, 0x13, 0x00, 0x00, 0x00, 0x00,
0x21, 0x1A, 0x00, 0x1C,
    0x01, 0x13, 0x00, 0x00, 0x00, 0x00, 0x20, 0x1A, 0x00, 0x1C, 0x01, 0x13,
0x00, 0x00, 0x00, 0x00,
```

```
0x57, 0x1A, 0x00, 0x1C, 0x01, 0x13, 0x00, 0x00, 0x00, 0x00, 0x65, 0x1A,
0x00, 0x1c, 0x01, 0x13,
   0x00, 0x00, 0x00, 0x00, 0x6c, 0x1A, 0x00, 0x1c, 0x01, 0x13, 0x00, 0x00,
0x00, 0x00, 0x63, 0x1A,
    0x00, 0x1c, 0x01, 0x13, 0x00, 0x00, 0x00, 0x00, 0x6F, 0x1A, 0x00, 0x1c,
0x01, 0x13, 0x00, 0x00,
   0x00, 0x00, 0x6D, 0x1A, 0x00, 0x1C, 0x01, 0x13, 0x00, 0x00, 0x00, 0x00,
0x65, 0x1A, 0x00, 0x1C,
   0x01, 0x13, 0x00, 0x00, 0x00, 0x00, 0x20, 0x1A, 0x00, 0x1C, 0x01, 0x13,
0x00, 0x00, 0x00, 0x00,
   0x74, 0x1A, 0x00, 0x1C, 0x01, 0x13, 0x00, 0x00, 0x00, 0x00, 0x6F, 0x1A,
0x00, 0x1C, 0x01, 0x13,
    0x00, 0x00, 0x00, 0x00, 0x20, 0x1A, 0x00, 0x1C, 0x01, 0x13, 0x00, 0x00,
0x00, 0x00, 0x73, 0x1A,
    0x00, 0x1c, 0x01, 0x13, 0x00, 0x00, 0x00, 0x00, 0x63, 0x1A, 0x00, 0x1c,
0x01, 0x13, 0x00, 0x00,
   0x00, 0x00, 0x75, 0x1A, 0x00, 0x1C, 0x01, 0x13, 0x00, 0x00, 0x00, 0x00,
0x63, 0x1A, 0x00, 0x1C,
   0x01, 0x13, 0x00, 0x00, 0x00, 0x00, 0x74, 0x1A, 0x00, 0x1C, 0x01, 0x13,
0x00, 0x00, 0x00, 0x00,
   0x66, 0x1A, 0x00, 0x1c, 0x01, 0x13, 0x00, 0x00, 0x00, 0x00, 0x20, 0x1A,
0x00, 0x1C, 0x01, 0x13,
    0x00, 0x00, 0x00, 0x00, 0x63, 0x1A, 0x00, 0x1C, 0x01, 0x13, 0x00, 0x00,
0x00, 0x00, 0x68, 0x1A,
    0x00, 0x1c, 0x01, 0x13, 0x00, 0x00, 0x00, 0x00, 0x61, 0x1A, 0x00, 0x1c,
0x01, 0x13, 0x00, 0x00,
   0x00, 0x00, 0x6c, 0x1A, 0x00, 0x1c, 0x01, 0x13, 0x00, 0x00, 0x00, 0x00,
0x6C, 0x1A, 0x00, 0x1C,
   0x01, 0x13, 0x00, 0x00, 0x00, 0x00, 0x65, 0x1A, 0x00, 0x1C, 0x01, 0x13,
0x00, 0x00, 0x00, 0x00,
   0x6E, 0x1A, 0x00, 0x1C, 0x01, 0x13, 0x00, 0x00, 0x00, 0x00, 0x67, 0x1A,
0x00, 0x1c, 0x01, 0x13,
    0x00, 0x00, 0x00, 0x00, 0x65, 0x1A, 0x00, 0x1C, 0x01, 0x13, 0x00, 0x00,
0x00, 0x00, 0x21, 0x1A,
    0x00, 0x1c, 0x01, 0x1b, 0x23, 0x18, 0x23, 0x13, 0x00, 0x00, 0x00, 0x00,
0x0A, 0x1A, 0x00, 0x1C,
    0x01, 0x13, 0x00, 0x00, 0x00, 0x00, 0x49, 0x1A, 0x00, 0x1C, 0x01, 0x13,
0x00, 0x00, 0x00, 0x00,
   0x6E, 0x1A, 0x00, 0x1C, 0x01, 0x13, 0x00, 0x00, 0x00, 0x00, 0x70, 0x1A,
0x00, 0x1c, 0x01, 0x13,
   0x00, 0x00, 0x00, 0x00, 0x75, 0x1A, 0x00, 0x1C, 0x01, 0x13, 0x00, 0x00,
0x00, 0x00, 0x74, 0x1A,
   0x00, 0x1c, 0x01, 0x13, 0x00, 0x00, 0x00, 0x00, 0x3A, 0x1A, 0x00, 0x1c,
0x01, 0x1D, 0x07, 0x18,
   0x07, 0x1B, 0x17, 0x1E, 0x00, 0x1F, 0x13, 0x01, 0x00, 0x00, 0x00, 0x1A,
0x20, 0x00, 0x01, 0x22,
    0x00, 0x03, 0x19, 0x00, 0x1c, 0x01, 0x19, 0x01, 0x1c, 0x01, 0x13, 0x02,
0x00, 0x00, 0x00, 0x73,
   0x20, 0x00, 0x02, 0x22, 0x00, 0x03, 0x13, 0x03, 0x00, 0x00, 0x00, 0x63,
0x20, 0x01, 0x03, 0x22,
   0x00, 0x03, 0x19, 0x00, 0x1c, 0x01, 0x19, 0x01, 0x1c, 0x01, 0x13, 0x02,
0x00, 0x00, 0x00, 0x75,
    0x20, 0x00, 0x02, 0x22, 0x00, 0x03, 0x13, 0x03, 0x00, 0x00, 0x00, 0x63,
0x20, 0x01, 0x03, 0x22,
   0x00, 0x03, 0x19, 0x00, 0x1c, 0x01, 0x19, 0x01, 0x1c, 0x01, 0x13, 0x02,
0x00, 0x00, 0x00, 0x74,
    0x20, 0x00, 0x02, 0x22, 0x00, 0x03, 0x13, 0x03, 0x00, 0x00, 0x00, 0x66,
0x20, 0x01, 0x03, 0x22,
```

```
0x00, 0x03, 0x19, 0x00, 0x1c, 0x01, 0x13, 0x02, 0x00, 0x00, 0x00, 0x7B,
0x20, 0x00, 0x02, 0x22,
    0x00, 0x03, 0x1E, 0x00, 0x1C, 0x19, 0x19, 0x00, 0x13, 0x02, 0x00, 0x00,
0x00, 0x7D, 0x20, 0x00,
    0x02, 0x22, 0x00, 0x03, 0x1E, 0x00, 0x1C, 0x07, 0x15, 0x00, 0x00, 0x00,
0x12, 0x15, 0x00, 0x00,
    0x00, 0x01, 0x19, 0x00, 0x16, 0x01, 0x0D, 0x00, 0x01, 0x0A, 0x00, 0x01,
0x1A, 0x00, 0x14, 0x01,
    0x1c, 0x01, 0x16, 0x02, 0x16, 0x03, 0x20, 0x02, 0x03, 0x24, 0x16, 0x14,
0x00, 0x13, 0x00, 0x00,
    0x00, 0x00, 0x01, 0x0A, 0x02, 0x00, 0x16, 0x00, 0x14, 0x03, 0x14, 0x02,
0x23, 0x02, 0xB2, 0x1E,
    0x07, 0x19, 0x02, 0x13, 0x03, 0x00, 0x00, 0x00, 0x45, 0x20, 0x02, 0x03,
0x22, 0x00, 0x03, 0x1c,
    0x01, 0x19, 0x02, 0x13, 0x03, 0x00, 0x00, 0x00, 0x7A, 0x20, 0x02, 0x03,
0x22, 0x00, 0x03, 0x1c,
    0x01, 0x19, 0x02, 0x13, 0x03, 0x00, 0x00, 0x00, 0x5F, 0x20, 0x02, 0x03,
0x22, 0x00, 0x03, 0x1c,
    0x01, 0x19, 0x02, 0x13, 0x03, 0x00, 0x00, 0x00, 0x56, 0x20, 0x02, 0x03,
0x22, 0x00, 0x03, 0x1c,
    0x01, 0x19, 0x02, 0x13, 0x03, 0x00, 0x00, 0x00, 0x71, 0x20, 0x02, 0x03,
0x22, 0x00, 0x03, 0x1c,
    0x01, 0x19, 0x02, 0x13, 0x03, 0x00, 0x00, 0x00, 0x7A, 0x20, 0x02, 0x03,
0x22, 0x00, 0x03, 0x1c,
    0x01, 0x19, 0x02, 0x13, 0x03, 0x00, 0x00, 0x00, 0x7A, 0x20, 0x02, 0x03,
0x22, 0x00, 0x03, 0x1c,
    0x01, 0x19, 0x02, 0x13, 0x03, 0x00, 0x00, 0x00, 0x85, 0x20, 0x02, 0x03,
0x22, 0x00, 0x03, 0x1c,
   0x01, 0x19, 0x02, 0x13, 0x03, 0x00, 0x00, 0x00, 0x71, 0x20, 0x02, 0x03,
0x22, 0x00, 0x03, 0x1c,
    0x01, 0x19, 0x02, 0x13, 0x03, 0x00, 0x00, 0x00, 0x70, 0x20, 0x02, 0x03,
0x22, 0x00, 0x03, 0x1c,
    0x01, 0x19, 0x02, 0x13, 0x03, 0x00, 0x00, 0x00, 0x5F, 0x20, 0x02, 0x03,
0x22, 0x00, 0x03, 0x1c,
    0x01, 0x19, 0x02, 0x13, 0x03, 0x00, 0x00, 0x00, 0x4D, 0x20, 0x02, 0x03,
0x22, 0x00, 0x03, 0x1c,
    0x01, 0x19, 0x02, 0x13, 0x03, 0x00, 0x00, 0x00, 0x79, 0x20, 0x02, 0x03,
0x22, 0x00, 0x03, 0x1c,
    0x01, 0x19, 0x02, 0x13, 0x03, 0x00, 0x00, 0x00, 0x7B, 0x20, 0x02, 0x03,
0x22, 0x00, 0x03, 0x1C,
    0x01, 0x19, 0x02, 0x13, 0x03, 0x00, 0x00, 0x00, 0x76, 0x20, 0x02, 0x03,
0x22, 0x00, 0x03, 0x1c,
    0x01, 0x19, 0x02, 0x13, 0x03, 0x00, 0x00, 0x00, 0x89, 0x20, 0x02, 0x03,
0x22, 0x00, 0x03, 0x1c,
    0x01, 0x19, 0x02, 0x13, 0x03, 0x00, 0x00, 0x00, 0x90, 0x20, 0x02, 0x03,
0x22, 0x00, 0x03, 0x1c,
    0x01, 0x19, 0x02, 0x13, 0x03, 0x00, 0x00, 0x00, 0x89, 0x20, 0x02, 0x03,
0x22, 0x00, 0x03, 0x1c,
    0x01, 0x1B, 0x1E, 0x00, 0x13, 0x00, 0x00, 0x00, 0x00, 0x0A, 0x1A, 0x00,
0x1c, 0x01, 0x13, 0x00,
    0x00, 0x00, 0x00, 0x43, 0x1A, 0x00, 0x1C, 0x01, 0x13, 0x00, 0x00, 0x00,
0x00, 0x6F, 0x1A, 0x00,
    0x1c, 0x01, 0x13, 0x00, 0x00, 0x00, 0x00, 0x72, 0x1A, 0x00, 0x1c, 0x01,
0x13, 0x00, 0x00, 0x00,
    0x00, 0x72, 0x1A, 0x00, 0x1C, 0x01, 0x13, 0x00, 0x00, 0x00, 0x00, 0x63,
0x1A, 0x00, 0x1C, 0x01,
    0x13, 0x00, 0x00, 0x00, 0x00, 0x74, 0x1A, 0x00, 0x1C, 0x01, 0x13, 0x00,
0x00, 0x00, 0x00, 0x21,
```

```
0x1A, 0x00, 0x1C, 0x01, 0x13, 0x00, 0x00, 0x00, 0x00, 0x0A, 0x1A, 0x00,
0x1C, 0x01, 0x1E, 0x00,
           0x18, 0x08, 0xff, 0xff, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00,
0x00, 0x00, 0x00, 0x00,
           0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00,
0x00, 0x00, 0x00, 0x00,
           0 \times 00, 0 \times 
0x00, 0x00, 0x00, 0x00,
           0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00, 0x00
};
uint32_t PC = 0;
uint32_t mem_ptr = 0;
uint32_t imm;
int main(){
            freopen("dump.txt", "w", stdout);
           while(code[PC]){
                        printf("%d: ", PC);
                        switch (code[PC]){
                                    case 0xA:
                                                printf("add r%d, r%d\n", code[PC + 1], code[PC + 2]);
                                                PC += 3;
                                               break;
                                    case 0xB:
                                                printf("sub r%d, r%d\n", code[PC + 1], code[PC + 2]);
                                                PC += 3;
                                                break;
                                    case 0xC:
                                                printf("mul r%d, r%d\n", code[PC + 1], code[PC + 2]);
                                                PC += 3;
                                                break;
                                    case 0xD:
                                                printf("xor r%d, r%d\n", code[PC + 1], code[PC + 2]);
                                                PC += 3;
                                                break;
                                    case 0xE:
                                                printf("lsh r%d, r%d n", code[PC + 1], code[PC + 2]);
                                                PC += 3;
                                                break;
                                    case 0xF:
                                                printf("rsh r%d, r%d\n", code[PC + 1], code[PC + 2]);
                                                PC += 3;
                                                break;
                                    case 0x10:
                                                printf("rol r%d, r%d\n", code[PC + 1], code[PC + 2]);
                                                PC += 3;
                                                break;
                                    case 0x11:
                                                printf("ror r%d, r%d\n", code[PC + 1], code[PC + 2]);
                                                PC += 3;
                                                break;
                                    case 0x12:
                                                printf("mov r%d, r%d\n", code[PC + 1], code[PC + 2]);
                                                PC += 3;
                                                break;
                                    case 0x13:
                                                imm = READ_DWORD(PC + 2);
```

```
if(imm >= 32 && imm <= 126){
        printf("mov r%d, %d(\'%c\')\n", code[PC + 1], imm, imm);
    }else{
        printf("mov r%d, %d\n", code[PC + 1], imm);
    PC += 6;
    break;
case 0x14:
    printf("push r%d\n", code[PC + 1]);
    PC += 2;
    break;
case 0x15:
    printf("push %x\n", READ_DWORD(PC + 1));
    break;
case 0x16:
    printf("pop r%d\n", code[PC + 1]);
    PC += 2;
   break;
case 0x17:
    printf("read\n");
    PC ++;
    break;
case 0x18:
    printf("write %d\n", code[PC + 1]);
    PC += 2;
    break;
case 0x19:
    printf("mov r%d, mem[%d]\n", code[PC + 1], mem_ptr);
    PC += 2;
    break;
case 0x1A:
    printf("mov mem[%d], r%d\n", mem_ptr, code[PC + 1]);
    PC += 2;
    break;
case 0x1B:
    printf("clear_mem\n");
    PC ++;
    break;
case 0x1c:
    printf("add_mem_ptr %d\n", code[PC + 1]);
    mem_ptr += code[PC + 1];
    PC += 2;
   break;
case 0x1D:
    printf("sub_mem_ptr %d\n", code[PC + 1]);
    mem_ptr -= code[PC + 1];
    PC += 2;
   break;
case 0x1E:
    printf("set_mem_ptr %d\n", code[PC + 1]);
    mem_ptr = code[PC + 1];
    PC += 2;
   break;
case 0x1F:
    printf("mov r0, len\n");
    PC ++;
    break;
```

```
case 0x20:
                printf("cmp r%d, r%d\n", code[PC + 1], code[PC + 2]);
                PC += 3;
                break;
            case 0x21:
                printf("je %d\n", READ_WORD(PC + 1));
                PC += 3;
                break:
            case 0x22:
                printf("jne %d\n", READ_WORD(PC + 1));
                PC += 3;
                break;
            case 0x23:
                printf("jmp %d\n", READ_WORD(PC + 1));
                PC += 3;
                break;
            case 0x24:
                printf("je_rel %d\n", PC + code[PC + 1]);
                PC += 2;
                break;
            case 0x25:
                printf("jne_rel %d\n", PC + code[PC + 1]);
                PC += 2;
                break;
            case 0x26:
                printf("jmp_rel %d\n", PC + code[PC + 1]);
                PC += 2;
                break;
            case 0x27:
                printf("debug\n");
                PC ++;
                break;
            case 0x28:
                printf("nop\n");
                PC ++;
                break;
            case OxFF:
                printf("end\n");
                PC ++;
                break;
            default:
                printf("unknown code %d\n", code[PC]);
                PC ++;
                break;
        }
    }
}
```

分析dump下来的VM指令,首先直接跳转到101位置处的指令:

```
dump.txt ×

1 0: jmp 101
2 3: clear_mem
3 4: set_mem_ptr 0
```

```
101: mov r0, 7!('H')
107: mov mem[0], r0
109: add_mem_p:r 1
111: mov r0, 1 1('e')
117: mov mem[1, r0
119: add_mem_prr 1
121: mov r0, 108('l')
127: mov mem[2], r0
129: add_mem_p:r 1
131: mov r0, 1 8('l')
137: mov mem[3], r0
139: add_mem_p r 1
141: mov r0, 1 1('o')
147: mov mem[4], r0
149: add mem p:r 1
151: mov r0, 3 ('!')
157: mov mem[5], r0
159: add mem p:r 1
161: mov r0, 3!(' ')
167: mov mem[6], r0
169: add_mem_p:r 1
171: mov r0, 8 ('W')
177: mov mem[7], r0
179: add_mem_p:r 1
181: mov r0, 1 1('e')
187: mov mem[8 , r0
189: add_mem_p r 1
191: mov r0, 108('l')
```

读取输入,判断输入的长度是不是26,以及格式是不是 scuctf{xxx}:

```
164
165
      530: read
166
     531: set mem ptr 0
      533: mov r0, len
167
      534: mov r1, 26
168
      540: cmp r0, r1
169
170
      543: jne 3
      546: mov r0, mem[0]
171
172
      548: add mem ptr 1
      550: mov r1, mem[1]
173
      552: add_mem_rt. 1
174
      554: mov r2, 115('s')
175
176
      560: cmp r0, r2
177
      563: jne 3
     566: mov r3, 🗐('c')
178
      572: cmp r1,
179
180
      575: jne 3
      578: mov r0, mem[2]
181
182
      580: add_mem_ptr 1
      582: mov r1, mem[3]
183
      584: add_mem_rtr 1
184
      586: mov r2, 117('u')
185
      592: cmp r0, r2
186
187
      595: jne 3
      598: mov r3, 99('c')
188
189
      604: cmp r1,
190
      607: jne 3
      610: mov r0, mem[4]
191
      612: add_mem_rtr 1
192
      614: mov r1, mem[5]
193
      616: add_mem_rtr 1
194
      618: mov r2, 116('t')
195
196
      624: cmp r0,
197
      627: jne 3
      630: mov r3, 102('f')
198
199
      636: cmp r1,
      639: jne 3
200
      642: mov r0, nem[6]
201
      644: add mem ptr 1
202
```

```
203 646: mov r2, 123('{')
204 652: cmp r0, r2
```

接下来对 scuctf{} 内的输入进行循环加密,加密算法非常简单:

```
676: set_mem_ptr 0
212
     678: add mem ptr 7
213
    680: push 12
214
215 685: push 1
    690: mov r0, mem[7]
216
21
     692: pop r1
21
     694: xor r0, r1
     697: add r0. r1
219
220
     700: mov mem[7], r0
221
     702: push r1
222
     704: add_mem_ptr 1
223
     706: pop r2
224
     708: pop r3
225
     710: cmp r2, r3
226 713: je rel 735
227
     715: push r0
     717: mov r0, 1
228
229
     723: add r2, r0
     726: pop r0
230
231
     728: push r3
     730: push r2
232
233 732: jmp 690
     735: set_mem_ptr 7
234
235
     737: mov r2, mem[7]
236 739: mov r3, 69('E')
237
   745: cmp r2, r3
```

最后与密文做比较:

```
739: mov r3, 6 ('E')
236
237
      745: cmp r2, r
      748: jne 3
238
      751: add_mem_p r 1
239
      753: mov r2, mem[8]
240
      755: mov r3, 112('z')
241
      761: cmp r2, r
242
      764: jne 3
243
      767: add_mem_p r 1
244
245
      769: mov r2, mem[9]
      771: mov r3, 9!('_')
246
247
      777: cmp r2, r
      780: jne 3
248
      783: add_mem_p r 1
249
250
      785: mov r2, mem[10]
      787: mov r3, 8 ('V')
251
      793: cmp r2, r
252
      796: jne 3
253
     799: add mem p r 1
254
255
      801: mov r2, m m[11]
      803: mov r3, 1:3('q')
256
      809: cmp r2, r
257
258
      812: jne 3
      815: add mem_p r 1
259
      817: mov r2, m m[12]
260
      819: mov r3, 1 2('z')
261
      825: cmp r2, r
262
           jne 3
263
      828:
```

写出exp:

```
#include <cstdio>

char flag[100] = "Ez_Vqzz\x85qp_My{v\x89\x90\x89";

int main(){
   for(int i = 1;flag[i - 1];i ++){
      flag[i - 1] = (flag[i - 1] - i) ^ i;
   }
   printf("scuctf{%s}", flag);
}
```

解法二: angr梭哈

之前写的一个angr教程: https://bluesadi.github.io/0x401RevTrain-Tools/angr/11 %E5%88%A9%E 7%94%A8angr%E7%AC%A6%E5%8F%B7%E6%89%A7%E8%A1%8C%E6%A2%AD%E5%93%88VM%E 7%B1%BBCTF%E8%B5%9B%E9%A2%98/

话不多说, 先贴exp:

```
import angr
proj = angr.Project('RE7_VirtualMachine.exe', load_options={'auto_load_libs':
False})
proj.hook(addr=0x00413210, hook=angr.SIM_PROCEDURES['stubs']['Nop']())
__CheckForDebuggerJustMyCode
proj.hook(addr=0x004131E0, hook=angr.SIM_PROCEDURES['stubs']['Nop']())
j___RTC_CheckEsp
proj.hook(addr=0x00411037, hook=angr.SIM_PROCEDURES['libc']['scanf']())
proj.hook(addr=0x004110E6, hook=angr.SIM_PROCEDURES['libc']['printf']())
proj.hook_symbol('putchar', angr.SIM_PROCEDURES['libc']['putchar']())
proj.hook_symbol('exit', angr.SIM_PROCEDURES['libc']['exit']())
proj.hook_symbol('puts', angr.SIM_PROCEDURES['libc']['puts']())
proj.hook_symbol('memset', angr.SIM_PROCEDURES['libc']['memset']())
proj.hook_symbol('strlen', angr.SIM_PROCEDURES['libc']['strlen']())
proj.hook_symbol('malloc', angr.SIM_PROCEDURES['libc']['malloc']())
state = proj.factory.blank_state(addr=proj.loader.find_symbol('main_0'))
state.options.add(angr.options.ZERO_FILL_UNCONSTRAINED_MEMORY)
state.options.add(angr.options.ZERO_FILL_UNCONSTRAINED_REGISTERS)
simgr = proj.factory.simgr(state)
while len(simgr.active):
    print(simgr)
    for active in simgr.active:
        print(active.posix.dumps(1))
        if b'Correct' in active.posix.dumps(1):
            print(active.posix.dumps(0))
            exit(0)
    simgr.step()
```

有几个要注意的地方,第一是angr对PE文件的库函数识别效果很差,很多库函数都识别不出来,所以需要手动hook:

```
proj.hook(addr=0x00413210, hook=angr.SIM_PROCEDURES['stubs']['Nop']()) #
    __CheckForDebuggerJustMyCode
proj.hook(addr=0x004131E0, hook=angr.SIM_PROCEDURES['stubs']['Nop']()) #
j___RTC_CheckEsp
proj.hook(addr=0x00411037, hook=angr.SIM_PROCEDURES['libc']['scanf']())
proj.hook(addr=0x004110E6, hook=angr.SIM_PROCEDURES['libc']['printf']())
proj.hook_symbol('putchar', angr.SIM_PROCEDURES['libc']['putchar']())
proj.hook_symbol('exit', angr.SIM_PROCEDURES['libc']['exit']())
proj.hook_symbol('puts', angr.SIM_PROCEDURES['libc']['memset']())
proj.hook_symbol('memset', angr.SIM_PROCEDURES['libc']['memset']())
proj.hook_symbol('strlen', angr.SIM_PROCEDURES['libc']['strlen']())
proj.hook_symbol('malloc', angr.SIM_PROCEDURES['libc']['malloc']())
```

```
state.options.add(angr.options.ZERO_FILL_UNCONSTRAINED_MEMORY)
```

因为程序在读取输入之前没有初始化VM结构体的 len 变量(一个小疏忽),所以在程序的一开始打印"Hello…"语句时 len 是处于未约束状态的,这样会导致angr在下面的for循环中产生路径爆炸。为了避免这种情况,我们需要用0来填充未约束的内存:

[RE8]PyCode

这题需要逆向一个用Python3.9编译的pyc:



传统的pyc反编译工具uncompyle6只支持到了Python3.8,在线反编译也只能出一部分代码:

uncomputed translates Bython bytecode back into equivalent Python source code. It accepts bytecodes from Python version 1.0 to version 3.8, panning over 24 years of Python releases. We include Dropbox's Python 2.5 bytecode and some PyPy bytecodes.

Google查一下,发现有一个叫做pycdc的工具,据说可以反编译Python3.9的字节码:

3 Answers





Sadly enough, it's currently impossible. Decompile 3 has the latest pyc to py methods (decompilation), but it hasn't updated for python 3.9 yet as that update takes a very long time to create



And it will most likely never happen for 3.9 (the developer of decompyle3 said that he is focusing more on his main job and that he doesn't have time to create this update as the 3.9 python update really changed the workflow, so it will be very hard and time-consuming).



So for now, the only solution is to wait, but if you want to speed things up, you can always sponsor the creator of decompile 3 (https://github.com/sponsors/rocky) (as he said that if you would get enough money to work more on this project, he will)

Edit:

I have recently found out that there is an alternative
I haven't used it myself, but its meant to decompile the compiled
python file (.pyc) back to humanly readable code (.pv). For any python version!
You can check it out here: https://github.com/zrax/pycdc

这个工具既可以反编译出Python字节码也可以反编译出Python源码:

To run pycdas, the PYC Disassembler: ./pycdas [PATH TO PYC FILE] The byte-code disassembly is printed to stdout.

To run pycdc, the PYC Decompiler: ./pycdc [PATH TO PYC FILE] The decompiled Python source is printed to stdout. Any errors are printed to stderr.

用pycdc反编译出来的源码跟在线反编译的结果是一样的,有所残缺:

```
5676,
5520,
3504,
7550]

flag = input('Input:')

if len(flag) != 26:
    print('Wrong!')
    exit(1)

flag_arr = (lambda .0: [ ord(i) for i in .0 ])(flag)

flag_arr = (lambda .0: [ flag_arr[i] * 2 * i for i in .0 ])(range(26))

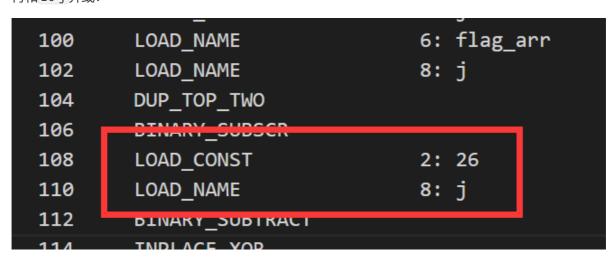
print('Success!')

print('Input is your flag!')
```

于是尝试用pycdas反编译出Python字节码,可以分析出加密流程。首先是将输入的每一个字符加上 j:

```
GET ITER
64
                                  20 (to 88)
66
        FOR ITER
        STODE NAME
68
                                  6: flag_arr
70
        LOAD NAME
72
        LOAD NAME
                                  8: j
74
        DUP TOP TWO
        BINARY SUBSCR
76
78
        LOAD NAME
                                  8: j
        INPLACE ADD
80
        ROT THREE
82
84
        STORE SUBSCR
```

再和 26-j 异或:



最后调用一个列表生成表达式,对应pycdc反编译出来的那一部分代码:

```
POP_TOP

LOAD_CONST 5: <CODE> listcomp>
LOAD_CONST 6: '<listcomp>'

MAKE_FUNCTION 0

LOAD_NAME 2: flag

GET_ITER

CALL_FUNCTION 1
```

```
0: flag_arr
        LOAD_GLOBAL
8
                                   1: i
10
        LOAD FAST
12
        BINARY SUBSCR
14
        LOAD CONST
                                   0: 2
        BINARY MULTIPLY
16
18
        LOAD FAST
                                   1: i
        BINARY MULTIPLY
20
        LIST APPEND
22
                                   2
24
        JUMP ABSOLUTE
                                   4
26
        RETURN VALUE
```

根据加密流程写出exp:

```
enc = [0, 250, 444, 678, 880, 1260, 1788, 952, 2352, 1944, 1960, 1144, 2784,
2522, 2576, 3450, 3712, 4182, 5040, 5282, 4680, 3906, 5676, 5520, 3504, 7550]

flag = 's'

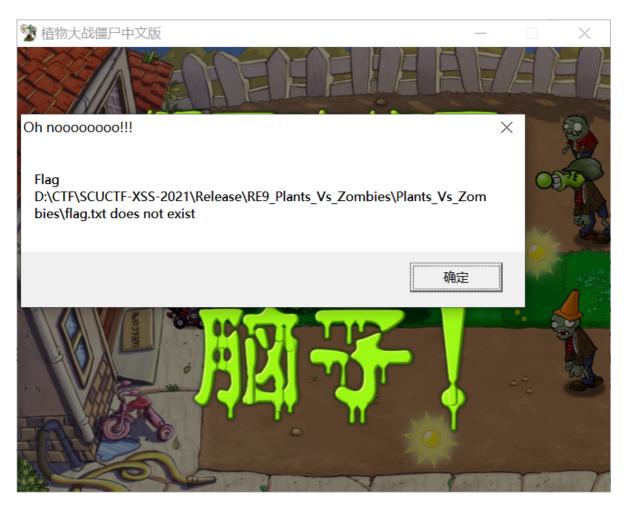
for i in range(1, 26):
    flag += chr(((enc[i] // i // 2) ^ (26 - i)) - i)

print(flag)
```

[RE9]Plants_Vs_Zombies

一道很有意思的题,这题改编自我之前用DLL注入的方法写的一个PVZ补丁: https://github.com/bluesadi/PVZPatch

这题会在游戏失败加密PVZ目录下的flag.txt文件,加密后的文件保存为flag.enc,如果没有flag.txt文件也会提示。所以此题的关键是找出加密代码,并对flag.enc文件解密得到flag:



首先将PlantsVsZombies.exe拖入IDA,可以发现在start函数的开头就调用了LoadLibrary函数,加载DBGHELP.DLL,start函数在正常情况下不可能出现这样的内容,说明DBGHELP.DLL是有问题的:

```
; CODE XREF: start-18F↑j
start
                proc near
; FUNCTION CHUNK AT .text:0061EA12 SIZE 0000018D BYTES
; FUNCTION
                         offset aDbghelpDll ; "DBGHELP.DLL"
                push
loc_61EBEC:
                                          ; CODE XREF: start-10↑j
                         ds:LoadLibraryA
                call
                call
                        sub_62E8F1
                         loc_61EA12
                 jmp
start
                endp
```

分析DBGHELP.DLL,发现有个叫做GameOver的导出函数,并且该函数中也出现了flag.txt这样的字符串,基本可以断定这个函数就是加密函数:

```
| Company | Comp
```

这里用到了天堂之门技术,在32位环境下获取64位API函数的句柄,并调用。有关天堂之门技术,我在看雪论坛上发了个<u>帖子</u>:

```
v7 = sub_10001690();
v6 = sub_100014A0(v7, HIDWORD(v7), "CreateFileA");
sub_100014A0(v7, HIDWORD(v7), "WriteFile");
v4 = sub_100014A0(v7, HIDWORD(v7), "ReadFile");
v5 = sub_100014A0(v7, HIDWORD(v7), "CloseHandle");
v2 = sub_100014A0(v7, HIDWORD(v7), "VirtualAlloc");
v3 = sub_100014A0(v7, HIDWORD(v7), "DeleteFileA");
momset(Ruffer, A sizeof(Ruffer));
```

主要加密过程如下,先读取flag.txt文件的内容,然后删除文件;加载一段shellcode到内存中,并执行:

```
else
{
    sub_100015A0(ReadFile_, HIDWORD(ReadFile_), 5, v8, (__int64)(int)v12, 100, 0, 0, 0, 0, 0);
    sub_100015A0(CloseHandle, HIDWORD(CloseHandle), 1, v8);
    sub_100015A0(DeleteFile, HIDWORD(DeleteFile), 1, (__int64)(int)Buffer);
    if ( !*(_QWORD *)&shellcode)
    {
        *(_QWORD *)&shellcode = sub_100015A0(VirtualAlloca, HIDWORD(VirtualAlloca), 4, 0, 0, 32, 0, 12288, 0, 64, 0);
        sub_10001020(shellcode, *(&shellcode + 1), &unk_100040E8, (int)&unk_100040E8 >> 31, 32, 0);
    }
    result = MK_FP(v1, shellcode)(&v12[strlen(v12) + 1] - &v12[1], &loc_10001D3B);
}
```

查看shellcode的内容,这段shellcode非常简单,本来是为了降低题目难度,没想到有同学直接用密码学的手段推出了加密流程(非预期解):

```
.data:100040E8 ;
.data:100040E8 loc_100040E8:
.data:100040E8
.data:100040E8
                               dec
                                        eax
                               dec
                                        ecx
                                        al, [esi+ecx]
                               mov
                               ror
                               add
                               dec
                                        eax
                               cmp
                                        short loc_100040FB
                               inz
                                        al, 64h
               loc_100040FB:
                               mov
                                        [esi+ecx], al
                               dec
                                        eax
                               test
                                        short loc_100040E8
                               jnz
                               push
                               push
                                        edx
                               dec
                                        eax
                               retf
```

根据加密流程和密文写出exp:

```
#include <cstdio>
#include <cstdint>
// scuctf{knockin_on_heavens_gate}
//
x6e\x6d\xb0\x6f\x92\xd1\x11\x74\xd5\xf6\x76\x78\x39\xda\xf9\xfc\xdd\xfc\x1f\xbf
\x40\xe3\xc2\xe4\x86\x04\x06\x47\xaa\xc9\xcd
int main(){
   uint8_t flag[] =
f\x40\xe3\xc2\xe4\x86\x04\x06\x47\xaa\xc9\xcd";
   for(int i = 0; i < 31; i ++){
       if(i == 6){
          flag[i] \wedge = 0x64;
       }
       flag[i] -= i;
       flag[i] = (flag[i] \ll 3) \mid (flag[i] \gg 5);
   printf("%s\n", flag);
}
```

[HappyRE]ezjava

hgg出的一道Java题,先用jd-gui打开看看,关键的CrackLicense类完全被混淆了:

```
| Description | Condition | Description | De
```

这题通过常规手段逆向有一定困难,因此可以考虑另辟蹊径。首先运行一下这个jar文件,发现会一直输出这一段话:

```
D:\CTF\SCUCTF-XSS-2021\Release\HappyRE_ezjava>java -jar HappyRE_ezjava.jar
[2021-12-06 21:12:01] Your license has expired, please purchase a new license from y4tacker!
[2021-12-06 21:12:06] Your license has expired, please purchase a new license from y4tacker!
```

对应的应该是这一部分的内容:

我们可以发现,如果校验失败会输出刚刚截屏中的那一段话,如果校验成功则应该输出正确的flag。可上面的if语句中并没有进行任何的加解密操作,而是直接输出一个字符串,所以可以认为其实flag已经提前解密出来在内存中的某个位置了,因此可以考虑用CE扫描内存,果然找到了了flag:

