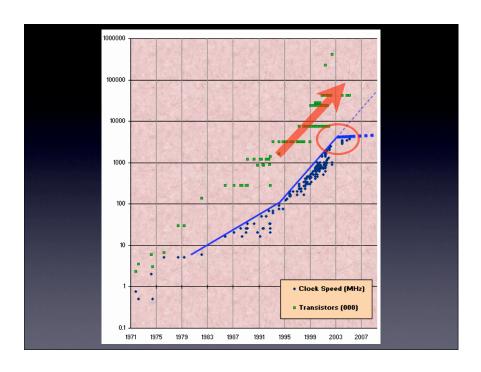
What the Enterprise Can Learn from Your Mom Jim Weirich Chief Scientist EdgeCase EdgeCase

Mom?

The Free Lunch Is Over A Fundamental Turn Toward Concurrency in Software By Herb Sutter • http://www.gotw.ca/publications/concurrency-ddj.htm • Published early 2005

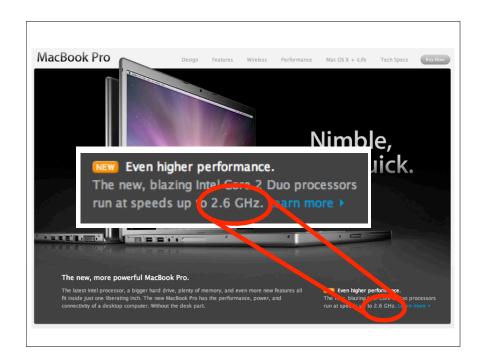
Moore's Law



Past Performance Gains Clock Speed Execution Optimization Cache

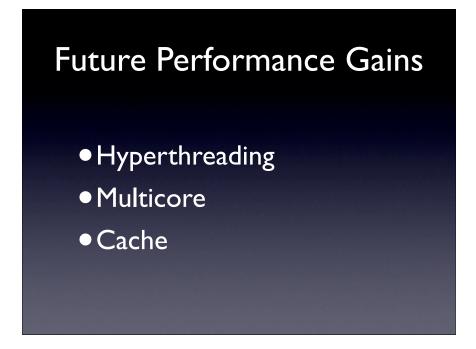


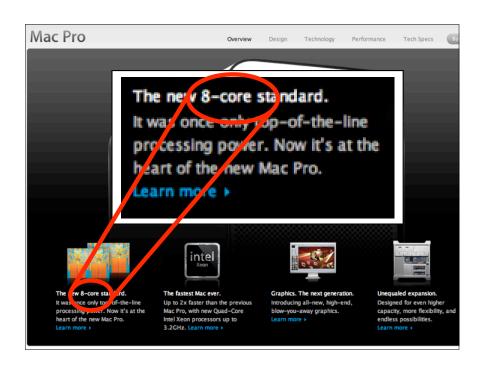


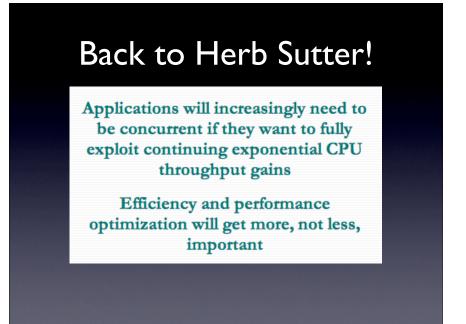


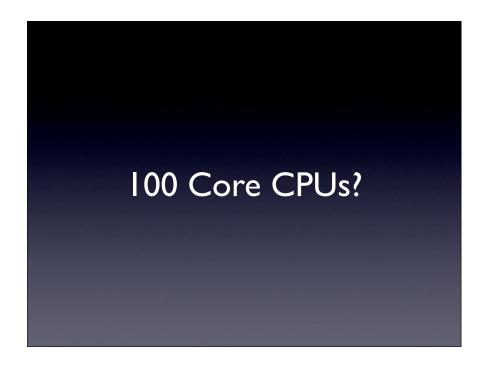




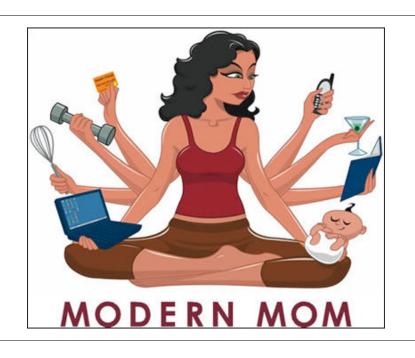








So, What Can the Enterprise Learn from Your Mom?





Charles Miller

In its place I would put Java Concurrency in Practice. Every new Atlassian developer gets handed this book and ordered to read it immediately. Writing multi-threaded code is hard, and a number of the things Java does under the hood to make its multi-threading more efficient makes it even harder. Unless you understand the subtleties described in this book of how Java shares data between threads, you will screw it up in some almost-impossible-to-debug way. [emphasis mine]

http://fishbowl.pastiche.org/2008/08/07/recommended_reading_for_java_d/

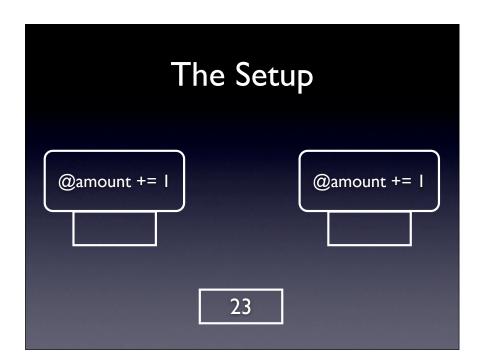


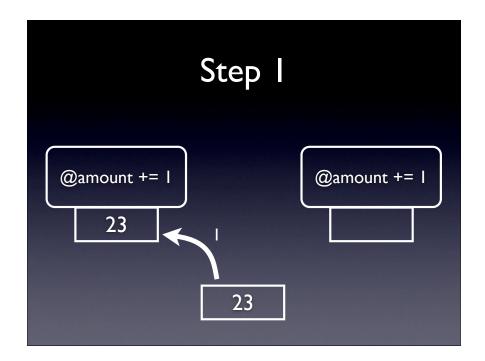


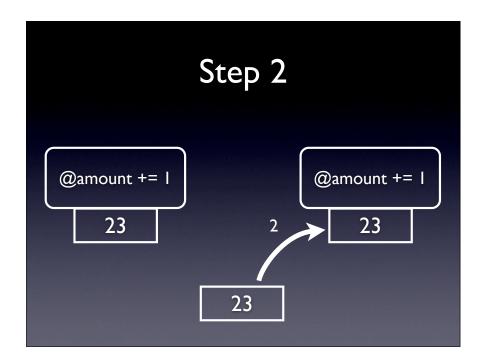
So, you want to write a concurrent program ...

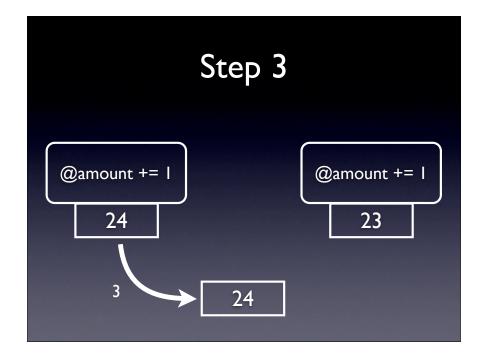
Demo: race I.rb

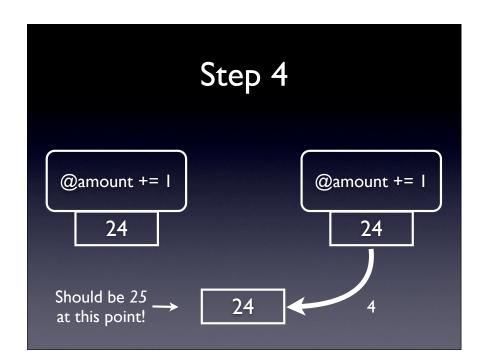
What happened?

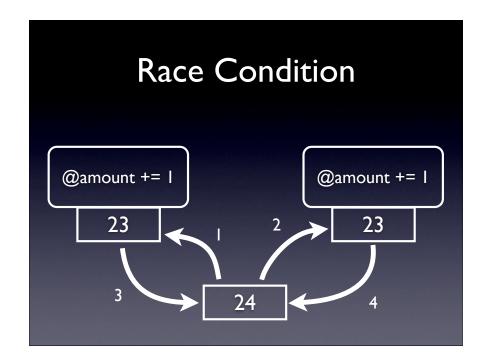


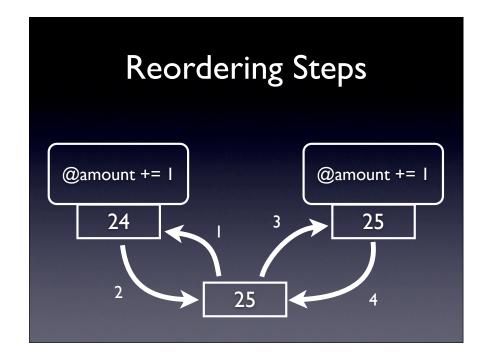


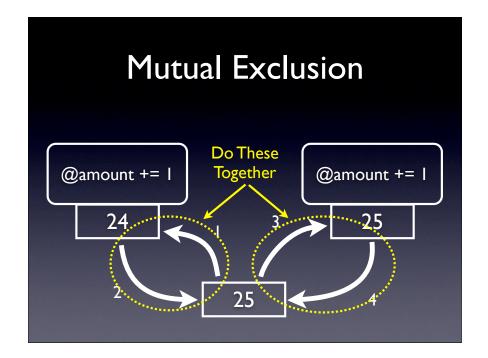












DISABLE CONTEXT SWITCHING
account.credit(1)
RE-ENABLE CONTEXT SWITCHING

require 'thread'
mutex = Mutex.new

...
mutex.synchonize do
 account.credit(1)
end

Demo: race4.rb

Demo: race7.rb

To Be Safe, You Must:

(I) Protect **every** shared memory access with a synchronizing lock.

Yes, **EVERY** access.

To Be Safe, You Must:

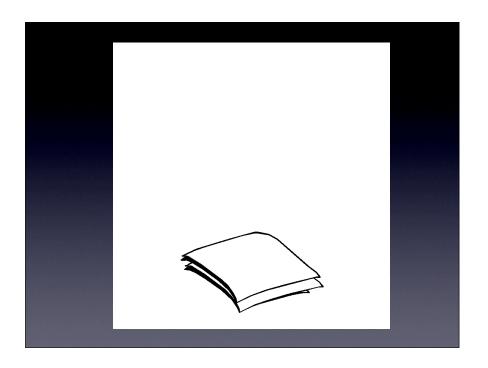
(2) Be aware of extended situations that need to be atomic.

To Be Safe, You Must:

(3) Have a strategy to avoid deadlock in the presence of multiple locks.

To Be Safe, You Must:

(4) Evaluate every single library used by your program to see if they also follow rules I - 3.





Horror Stories

- Real time data collection (1 in a million)
- Rake multitask dependencies
- Double Checked Lock

Double Checked Lock

```
public static Singleton getInstance()
{
  if (instance == null)
  {
    synchronized(Singleton.class) { //1
    if (instance == null) //2
       instance = new Singleton(); //3
    }
  }
  return instance;
}
```

The theory behind double-checked locking is perfect. Unfortunately, reality is entirely different.

Concurrent
Programming is
HARD

And it is hard because of Shared
Memory

What can we do to make concurrent programming easier?

Avoid
Shared
Memory

Are We Blug Programmers?

... Languages less powerful than Blub are obviously less powerful, because they're missing some feature he's used to. But when our hypothetical Blub programmer looks in the other direction, up the power continuum, he doesn't realize he's looking up. What he sees are merely weird languages... Blub is good enough for him, because he thinks in Blub.

-- Paul Graham, Beating the Averages



The Power of Messsaging

(Erlang)

Imagine a Language with

- No variables
- No assignment statements
- No explicit loops

Imagine a Language with

- No variables
- No assignment statements
- No explicit loops
- Only Constants
- Pattern Matching
- Recursion
 - (tail recursion)

Function Definitions

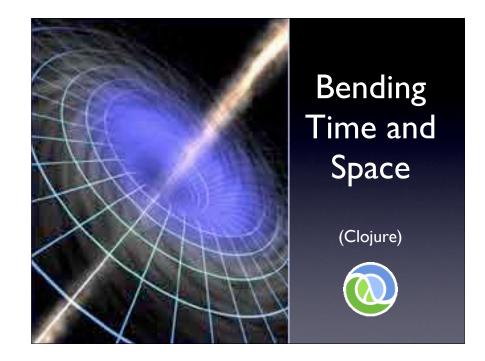
```
fact(0) ->
    1;
fact(N) ->
    N * fact(N-1).
```

Function Definitions

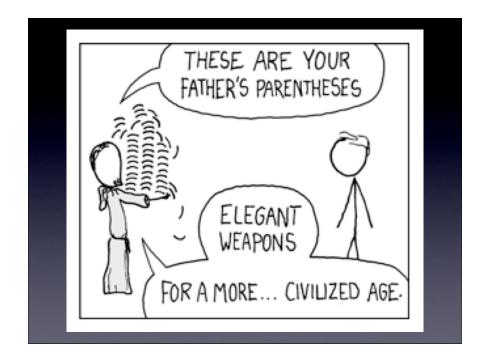
```
fact2(N) ->
    fact2(N, 1).

fact2(0, Acc) ->
    Acc;
fact2(N, Acc) ->
    fact2(N-1, N * Acc).
```

Erlang Demo







Quick Lisp Primer

Data Structures 1 321 ; Numbers a fido ; Names (1 2 3) ; Listss [1 2 3] ; Arrays

Calling Functions

```
(+ 2 4) ; => 6
(count '(a b c)) ; => 3
```

```
`(+ 2 4) ; => (+ 2 4)
```

Defining Functions

```
(defn factorial [n]
  (if (zero? i) 1
          (* (factorial (- n 1))) ))
```

Defining Functions

Sequences

Other Cool Stuff

Java Interface (import '(java.util.concurrent Executors)) (. Thread (sleep (* 1000 seconds)) Argument to sleep (static) Method to call (sleep) Class (or object) to get method call Do a Java call

No Modifiable State

(except ...)

 Vars - Thread Local Variables (impossible to share)

```
(def v 123)
(binding [v 321] v) # => 321
```

• v is bound to 321, but only for the current thread.

No Modifiable State

(except ...)

• Refs - STM (Sharable, but only in a transaction)

```
(def r (ref 0)
  (deref r) ; => 0

  (ref-set r 1) ; FAILS
  (dosync
        (ref-set r 1)) ; In Transaction
        (deref r) ; => 1
```

Clojure Demo

Summary

Concurrent Programming is Hard (primarily due to shared mutable state)

Hard enough that you probably want to avoid doing it in a traditional sequential programming language.

Functional Languages provide many advantages when dealing with concurrency.

Don't be a Blub Programmer.

Thank You! git://github.com/jimweirich/presentation_enterprise_mom.git

Contact Info

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