Department of Computer Science Algorithms (CC4010)

FCUP 2021/2022

rugorumis (CC+010)	2021/2022
Test (26.11.2021)	duration: 2h
N.º Name	
1. [1.5] Is $\Theta(\log_2 n) = \Theta(\log_{100000} n)$? Explain.	
(G2) (G100000) I	
2. [2.5] Which is the difference between <i>quicksort</i> and <i>random</i> complexity of <i>quicksort</i> and the expected time complexity of <i>quicksort</i> no specific input elicits the worst-case behavior?	
3. [2.5] Explain why the asymptotic time complexity of Insis $\Theta(n^2)$. Start by presenting the main ideas of the two method	

Algorithms (CC4010)	DCC-FCUP	Test (26.11.2021)
N.º Name		
4. [6.0] Given an array x of n integers, Suppose that all elements of x are dis $b-a+1$. Assume that SELECT (x,a) deterministically, in-place, and in time	stinct, $x[1]$ is the first element, (b, k) is correct and MYPARTI	$1 \le a \le b \le n \text{ and } 1 \le k \le$
SELECT (x, a, b, k) 1. if $a = b$ then return $x[a]$ 2. $t = \text{MYPARTITION}(x, a, b)$ 3. $p = t - a + 1$ 4. if $p = k$ then return $x[t]$ 5. if $p < k$ then 6. return SELECT $(x, t + 1, k)$ 7. return SELECT $(x, a, t - 1, k)$	- /	
a) Let $x[i] = s_i$, for $1 \le i \le n$, be the of t ? Recall that SELECT (x, a, b, k) returns		
b) Could the time complexity of SELE plexity of SELECT (x, a, b, k) in the wor		
c) Under the condition stated on a), jus	tify the correctness of SELECT.	

Algorithms	(CC4010)	DCC-FCUP	Test (26.11.2021)
N.º	Name		
) Which are the	e differences to quicks	elect and median of five medians	?
order (CCW), w	ve want to check whet		ex polygon P , in counterclockwis rior of P . Assume v_1 is the vertexithm for solving the problem.
	-	. (32 // 6	<u> </u>
	why radix sort require stable sorting algorith		to be a stable sorting algorithm. I

Algorithms (CC4010)	DCC-FCUP	Test (26.11.2021)
N.° Name		
smallest y -value, P is sorted in str	$\{a_n\}$ be a set of n points in the plane. A ictly decreasing order of polar angle onvex hull in clockwise order (CW),	w.r.t. p_1 (there are no ties), and we
Using pseudocode, write the algor the time complexity in this case.	ithm (by adapting Graham-scan). Exp	plain why it is correct and which is

Master theorem:

Let $a \ge 1$ and b > 1 be constants, let f(n) be a function, and let T(n) be defined on the nonnegative integers by the recurrence T(n) = aT(n/b) + f(n), where we interpret n/b to mean either $\lfloor n/b \rfloor$ or $\lceil n/b \rceil$. Then T(n) has the following asymptotic bounds:

- $1. \ \ \text{If } f(n) = O(n^{\log_b a \varepsilon}) \text{ for some constant } \varepsilon > 0 \text{, then } T(n) = \Theta(n^{\log_b a}).$
- 2. If $f(n) = \Theta(n^{\log_b a})$, then $T(n) = \Theta(n^{\log_b a} \log_2 n)$.
- 3. If $f(n) = \Omega(n^{\log_b a + \varepsilon})$, for some constant $\varepsilon > 0$, and if $af(n/b) \le cf(n)$ for some constant c < 1 and all sufficiently large n, then $T(n) = \Theta(f(n))$.