

Model-based software development

Lecture X.

Graph pattern matching,
Graph transformation

Dr. Semeráth Oszkár

Graph pattern matching, Graph transformation

Definitions

Graph pattern matching

Model transformations

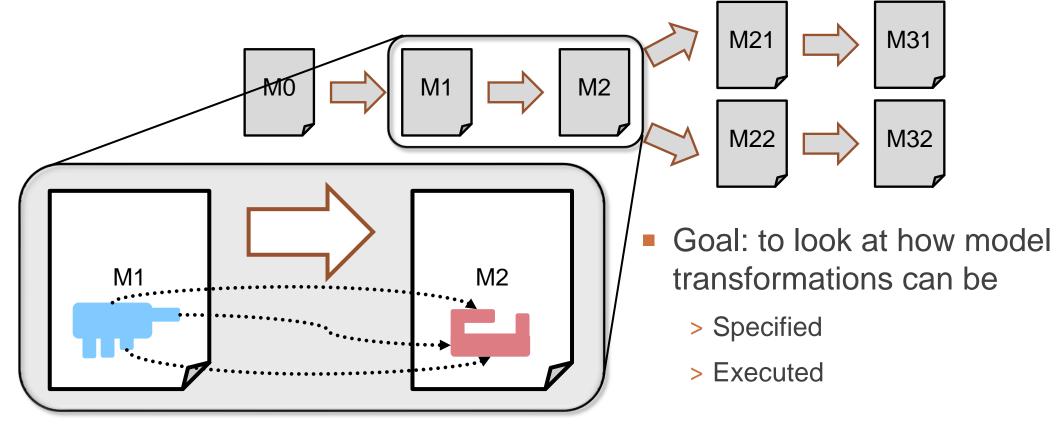
Incremental transformations

Design space exploration



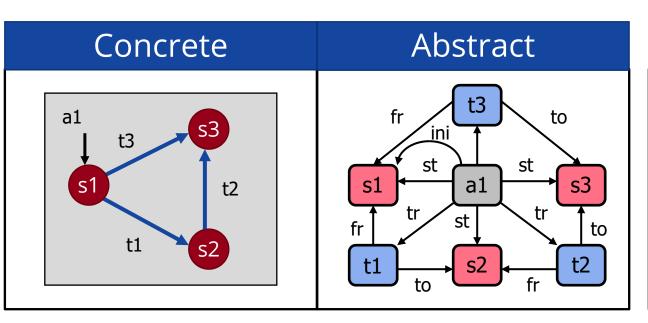
Motivation: Transformation of models

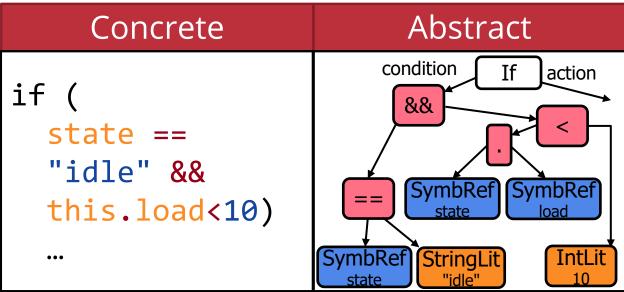
- Model-based development: Models as primary documents
- Developing models, automating model processing



Abstract syntax

- How to modify the models?
- Idea: modify the representation of models directly → Abstract syntax





Task: method to modify graphs!

Graph pattern matching, Graph transformation

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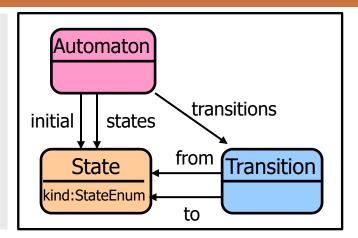
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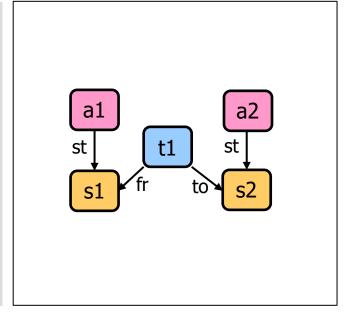


A simple example

Metamodel







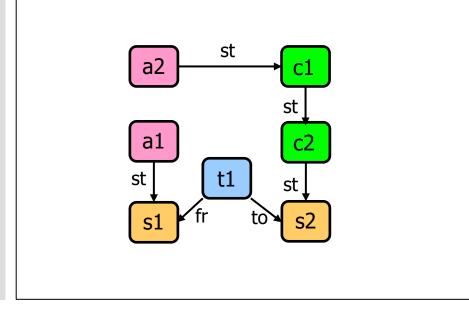
- Well-formedness constraint:
 - > Transition source & target states must be owned by same automaton
- Goal: to find violations...
 - > A violation is a *Transition*, whose "*from*" link points to a *State s1*, and "*to*" link points to a *State s2*, where the automaton of *s1* is not the automaton of *s2*

A more complex example

Metamodel

Region Automaton wabstract transitions initial states from Transition kind:StateEnum to

Violation example



- Well-formedness constraint:
 - > Transition source & target states must be owned by same automaton
- Goal: to find violations...
 - > A violation is a *Transition*, whose "*from*" link points to a *State s1*, and "*to*" link points to a *State s2*, where the automaton of *s1* is not the automaton of *s2*

Programmatic traversal vs. queries

Goal: find constraint violations in model

Traverse model in general-purpose language

```
for (Automaton automaton : automatons) {
  for (Transition transition : automaton.getTransitions()) {
    State sourceState = transition.from;
    // which automaton defines this state?
                                                                      "simple
   Automaton sourceAutomaton = null;
   for (Automaton candidate : automatons) {
                                                                     example"
      if (candidate.getStates().contains(sourceState)) {
        sourceAutomaton = candidate;
        break;
    // ... do the same for targetState, then
    if (sourceAutomaton != targetAutomaton)
      // report violation
```

Programmatic traversal vs. queries

- Goal: find constraint violations in model
 - > Traverse model in general-purpose language
 - > Use a Query DSL
 - More concise
 - Declarative functional specification of the query
 - Freely interpreted by query engine (e.g. optimization)
 - Can be platform-independent
- Validation is just one use cases for model queries
 - > Derived features
 - > M2M/M2T Transformation, Simulation

> ...

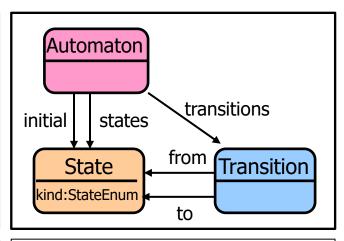
Query Language Styles

- SQL-like (relational algebra)
 - > Example: EMF Query
 - > © Good for attribute restrictions
 - >

 Not very concise for relationships (many joins)
- Functional style
 - > Example: OCL
 - > Somewhat declarative

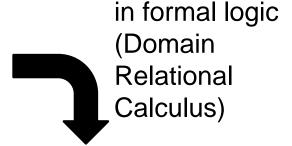
```
context Transition inv:
   Automaton.allInstances()->forAll(a |
       a.states->includes(self.from) =
       a.states->includes(self.to)
   );
```

- Logic style
 - > Domain relational calculus / graph patterns / Datalog
 - > Even more declarative



a1 a2 st t1 st st st to s2 fr to s4 to fr to t3

A violation is a *Transition*, whose "*from*" link points to a *State s1* and "*to*" link points to a *State s2*, where the automaton of *s1* is not the automaton of *s2*

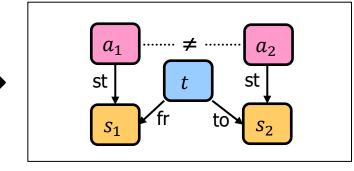


 $\{t \mid \exists s_1, s_2, a_1, a_2 : Transition(t) \land from(t, s_1) \land to(t, s_2) \land states(a_1, s_1) \land states(a_2, s_2) \land a_1 \neq a_2\}$

violates(t):Transition(t),from(t,s1),to(t,s2),
states(a1,s1),states(a2,s2),a1=\=a2

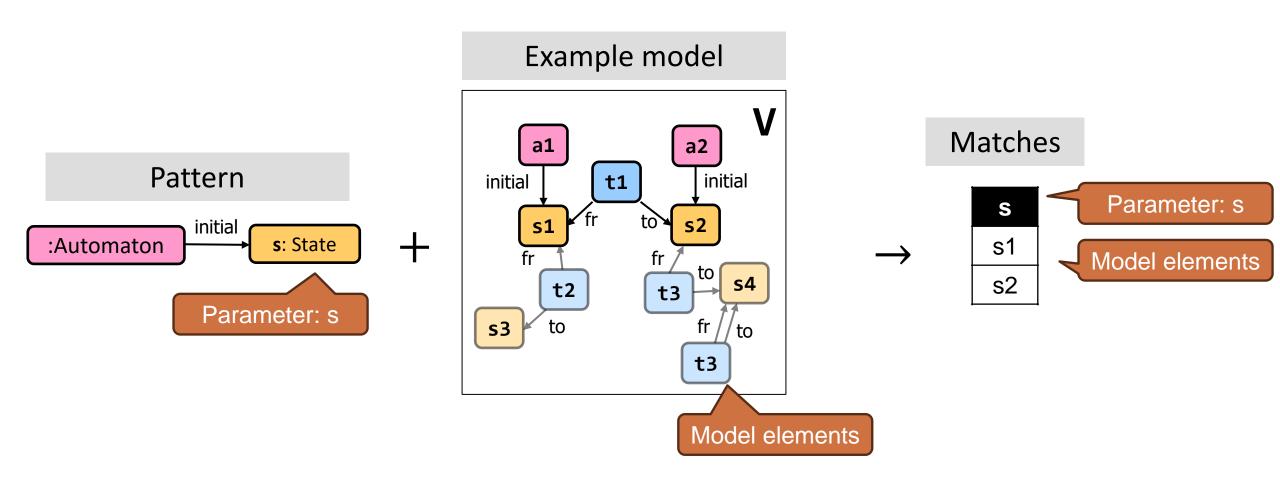






Pattern

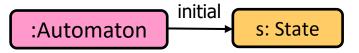
Pattern matching



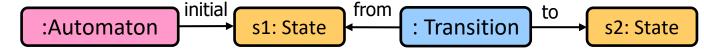
MODELLALAPÚ SZOFTVERFEJLESZTÉS 18

Examples

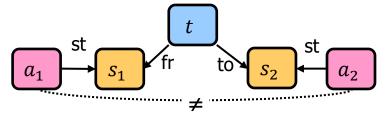
Simple example: Initial states in the model



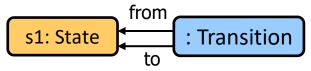
■ Chain (∧): Second states in the model



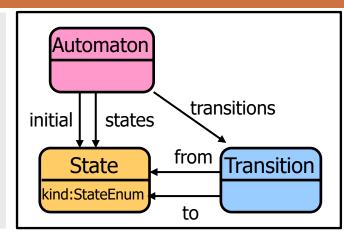
≠: Transition across automata



=: Loop edge

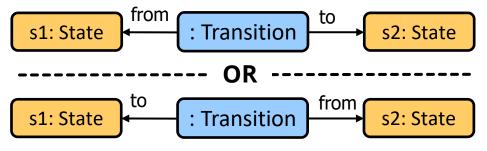


Metamodel

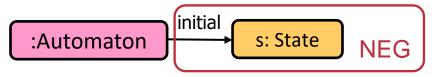


Examples 2

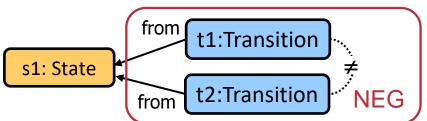
(v): Two states are connected



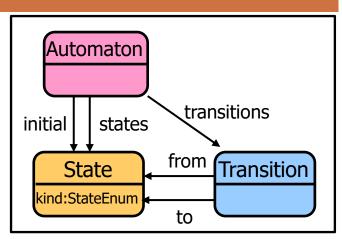
- (¬, Negative Application Condition):
 - > automaton without initial state



> a state from whose initial state no two transactions leave (deterministic)



Metamodel



Graph pattern matching, Graph transformation

Definitions

Graph pattern matching

Model transformations

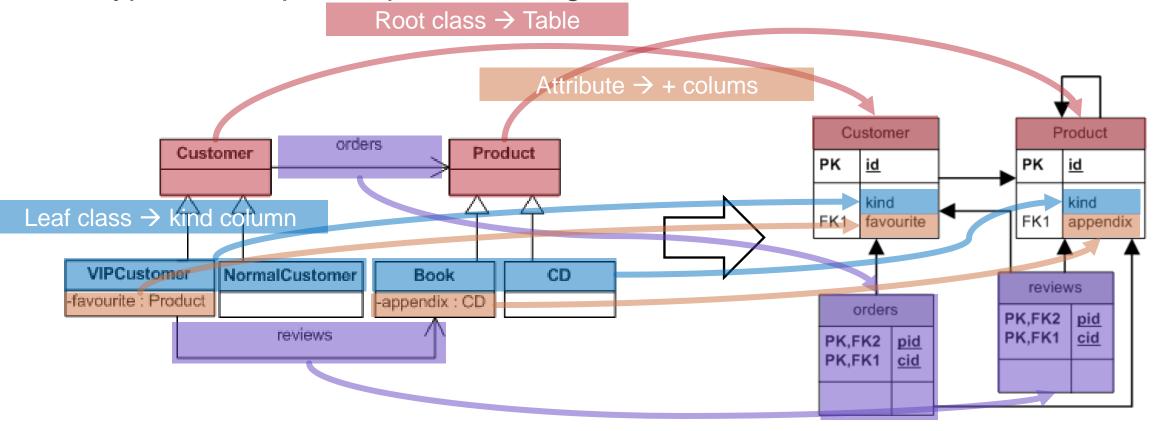
Incremental transformations

Design space exploration



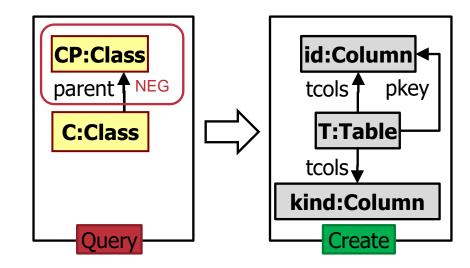
Example Transformation

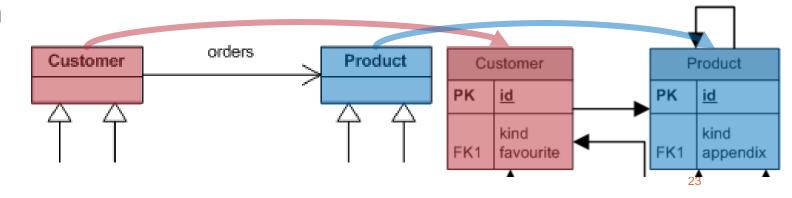
Typical example: map a class diagram to database tables!



Example Transformation

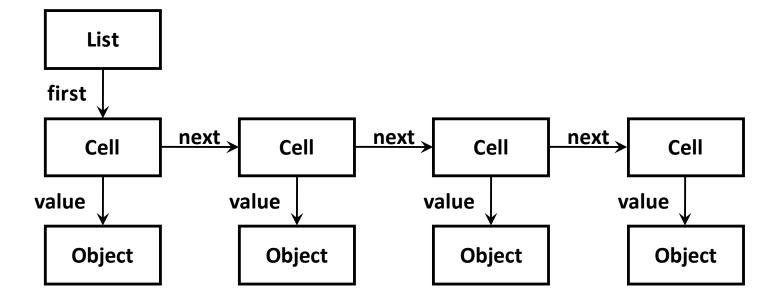
- How would we solve the problem of creating tables representing root classes?
- Query the root classes (class that has no ancestor)
- 2. Create the tables and with them the necessary columns
- 3. Repeat as long as we can
- Goal: To formulate the whole transformation with similar rules





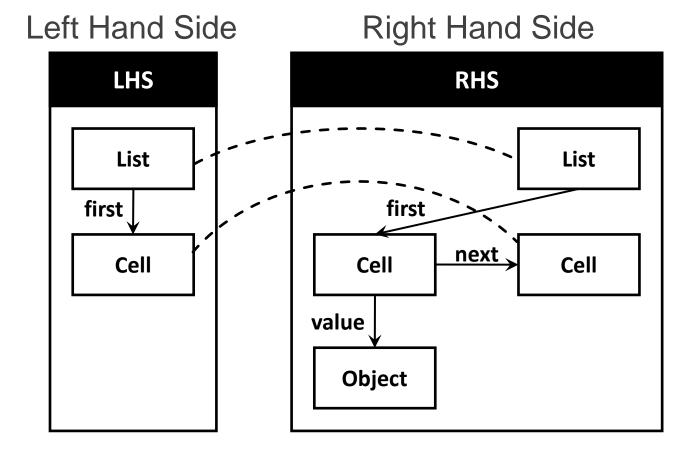
Graph transformation

Model = Labelled Graph



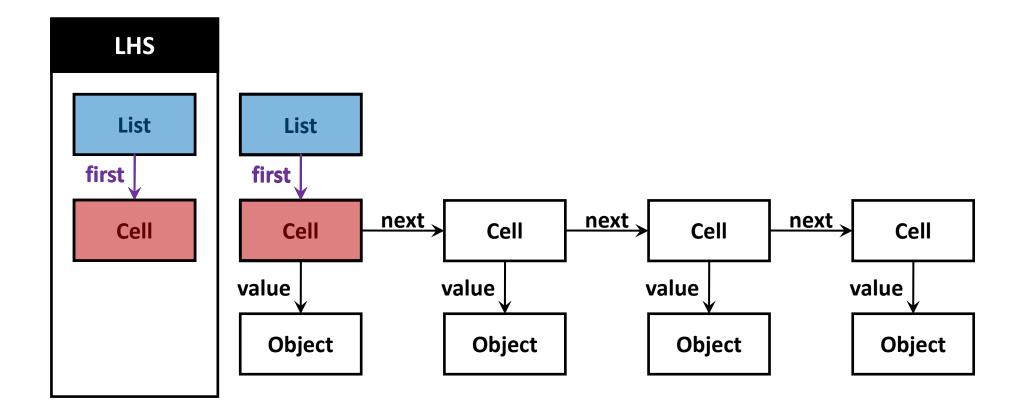
Graph Transformation rule

Graph rewriting rule, defined with two graphs



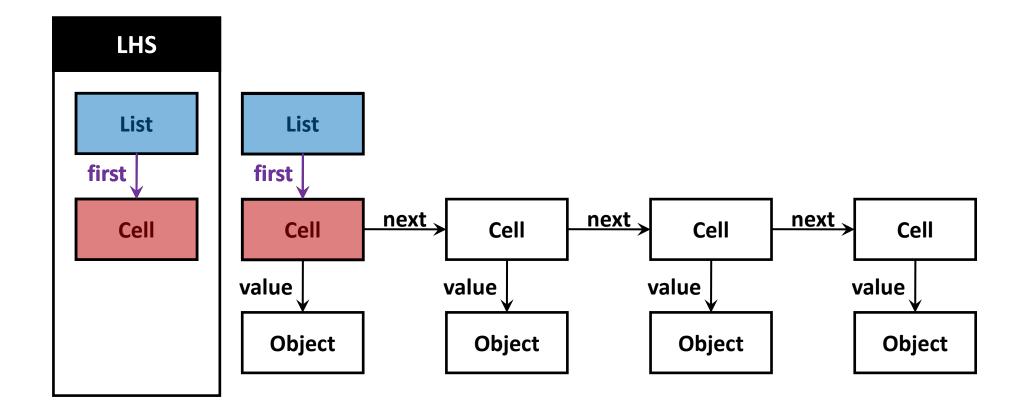
Graph Transformation: Pattern matching

Matching: find the subgraphs containing LHS in the source graph



Graph Transformation: Pattern matching

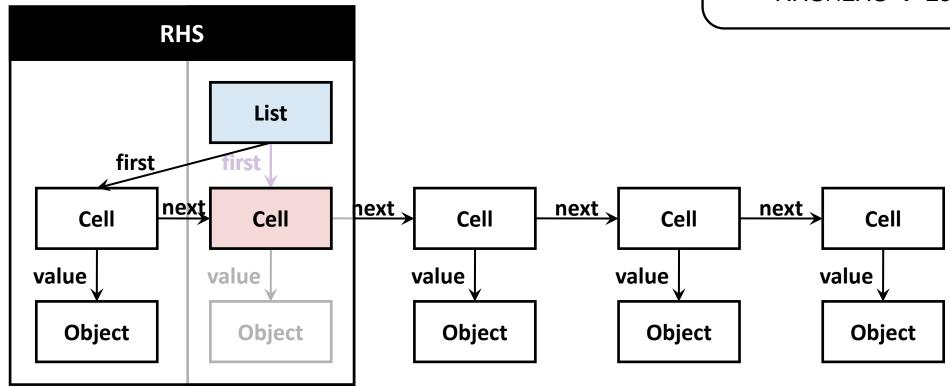
Matching: find the subgraphs containing LHS in the source graph



Graph Transformation: Execution of rewriting

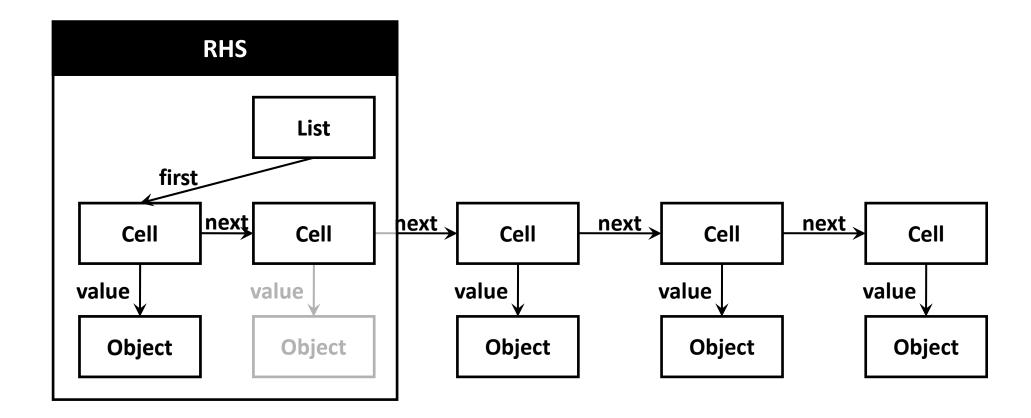
Rewriting the graph by the match: replace LHS with RHS.

LHS\RHS → Delete RHS\LHS → Insert RHS∩LHS → Leave it



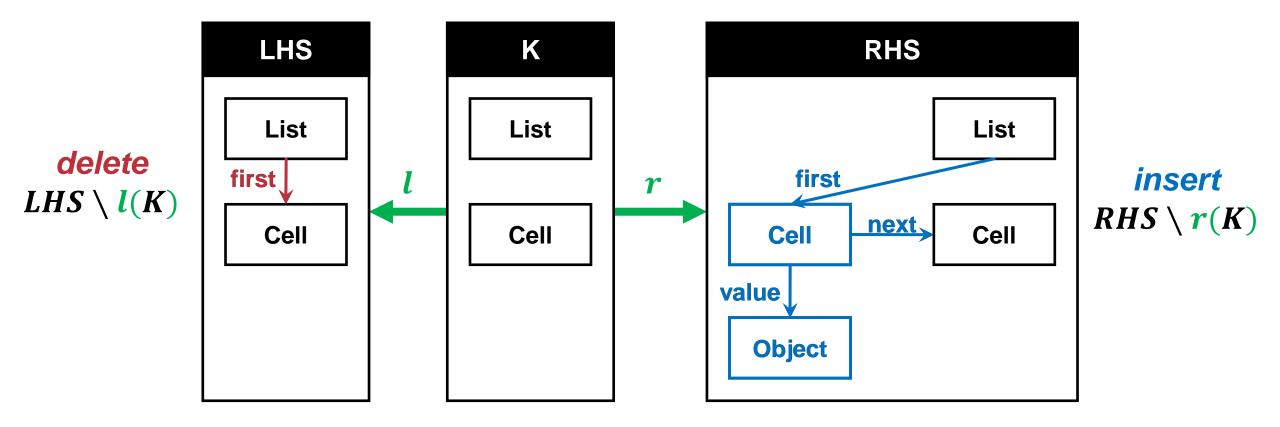
Graph Transformation: Execution of rewriting

We get a new graph

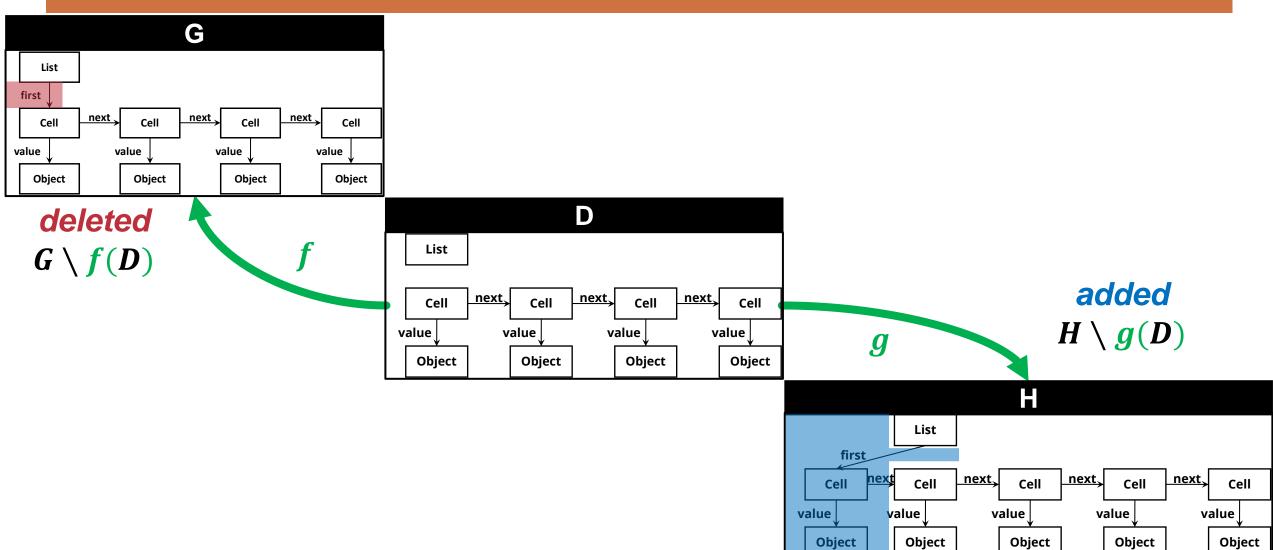


Anatomy of graph transformation

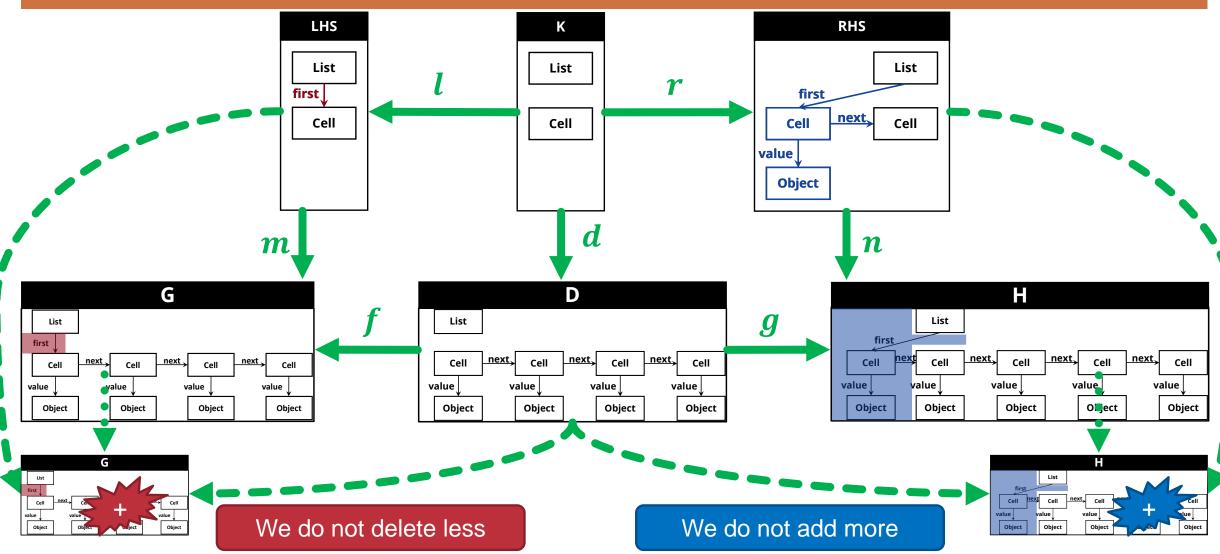
 Let us examine which graph can be matched to which graph during the transformation.



Anatomy of transformed models



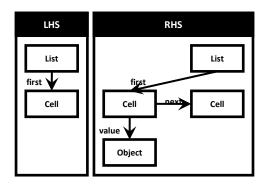
Complete anatomy



Graph Transformation

- Formulating rules for rewriting models
- Extending grammar rules

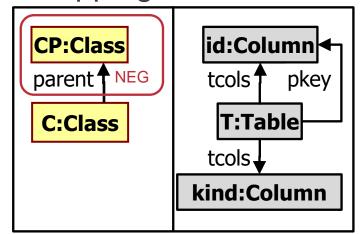
List → List, Cell vs

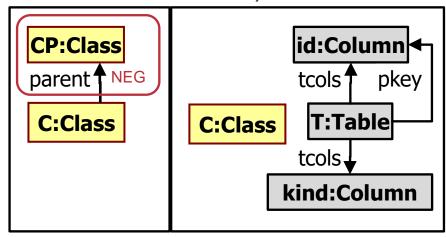


- Clear but mathematically precise formalism
 - (Termination, Ordering, Confluence, ...)
- Tool support (see previous practice)

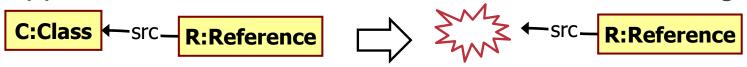
"Dangling edge" problem

Mapping of ancestor classes (with and without deletion)





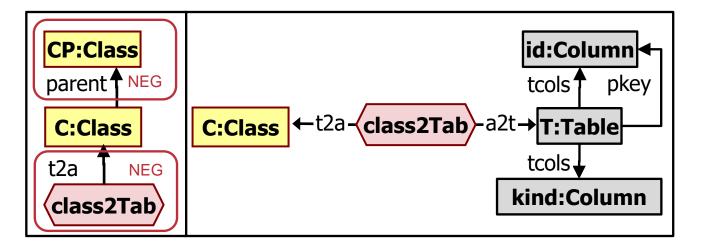
What happens if we delete an element that still has an edge pointing to it?



Resolving "dangling edges": Delete edges / Undo transformation

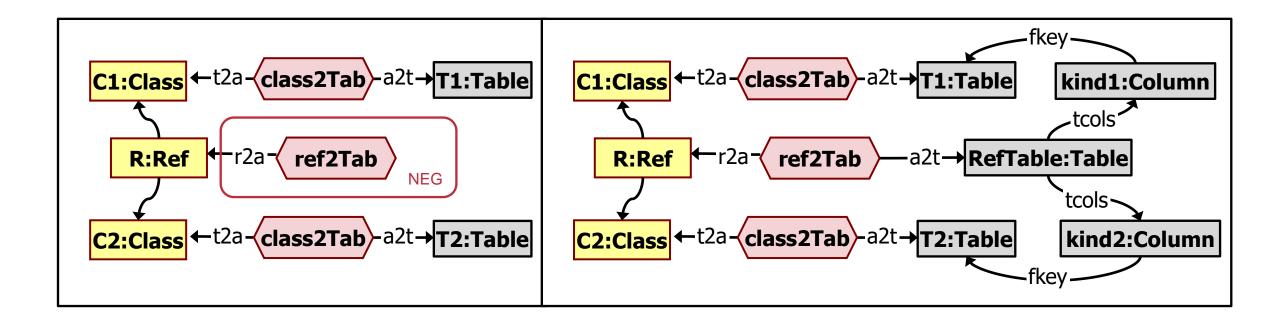
Examples

- Mapping of ancestor classes with traceability :
 - > Find an ancestor class
 - > that has not yet been mapped,
 - > then map it.



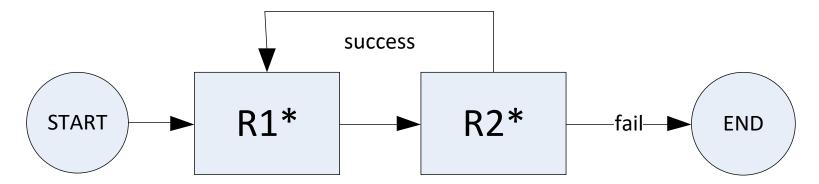
Examples

Mapping references

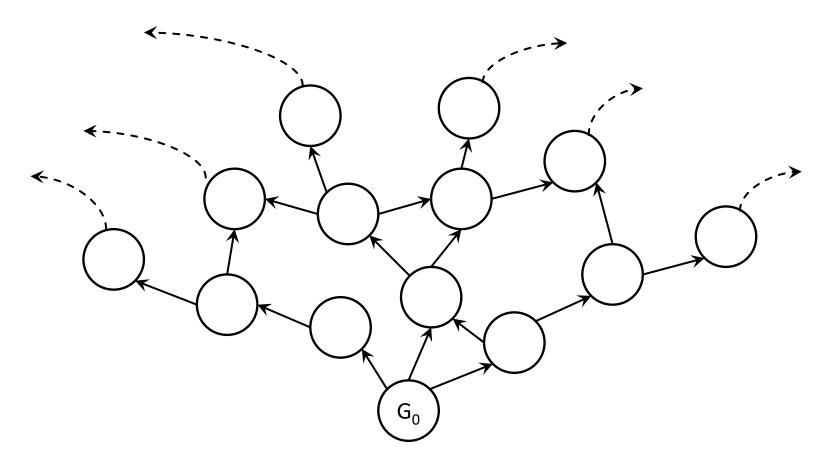


Control mechanisms

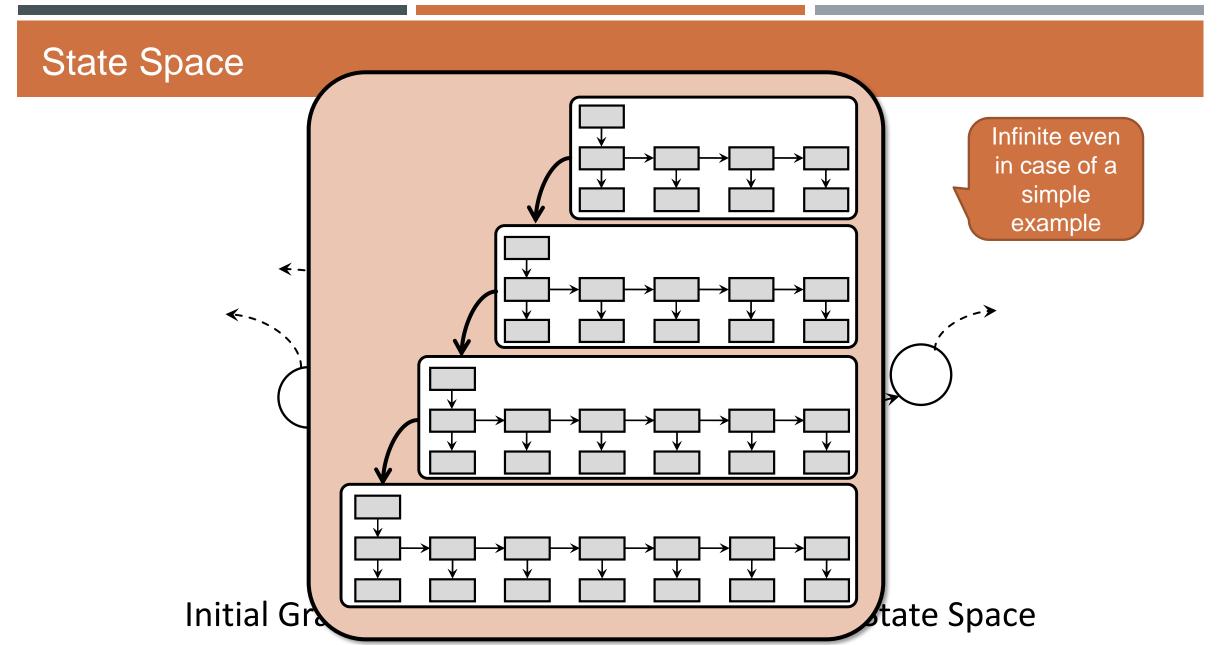
- In what order should the rules be executed?
- Multiple options, see previous lecture.
- But for example:
 - > Fire arbitrary transformations as long as possible (~ default)
 - > Fire all possible transformations once
 - > control graph (explicit control)



State Space



Initial Graph + GT rules → (Tipically infinite) State Space



Types of model transformations

- By number of inputs and outputs (In-place vs out-place)
- By the language (Endogenous vs exogenous)
- By the direction (unidirectional vs bidirectional)

Graph pattern matching, Graph transformation

Definitions

Graph pattern matching

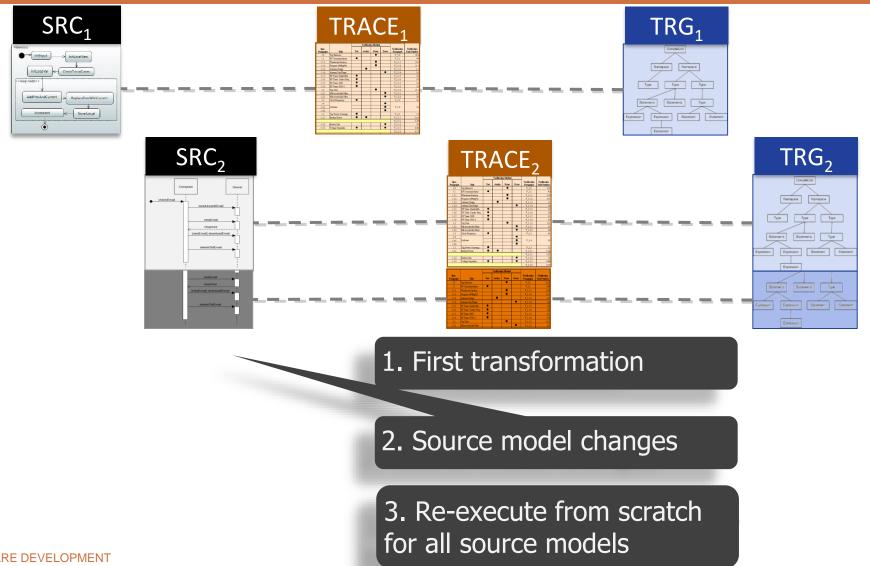
Model transformations

Incremental transzformations

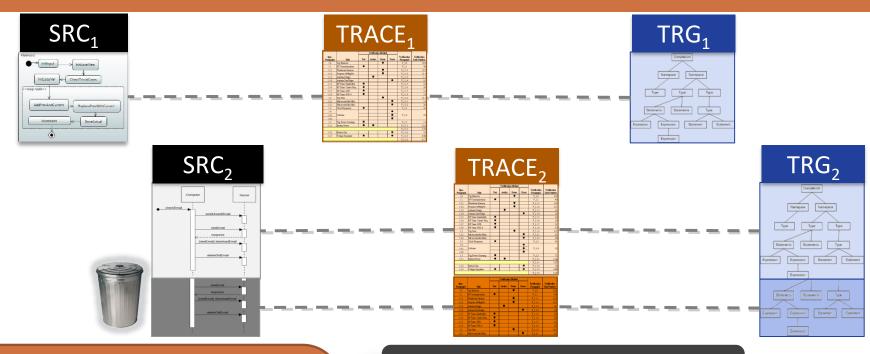
Design space exploration



No Incrementality: Batch Transformations



Dirty Incrementality



Pros:

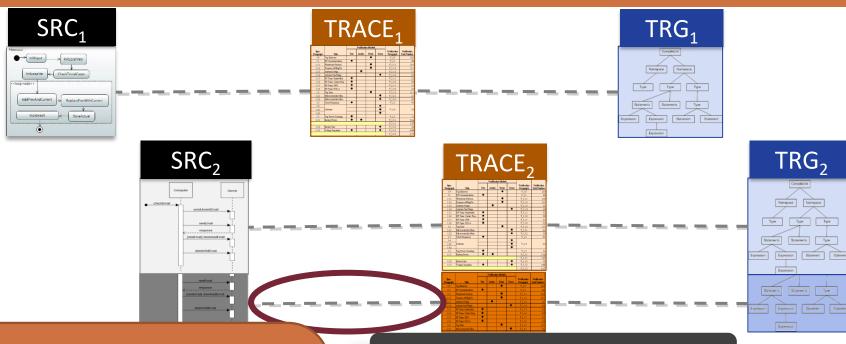
- Large-step incrementality
- Avoids continuous execution

Cons:

- Complex MT can be slow
- Cleanup (after an error)?
- Chaining?

- 1. First transformation
- 2. Source model changes
- 3. Re-execute from scratch only for changed models

Incrementality by Traceability



Pros:

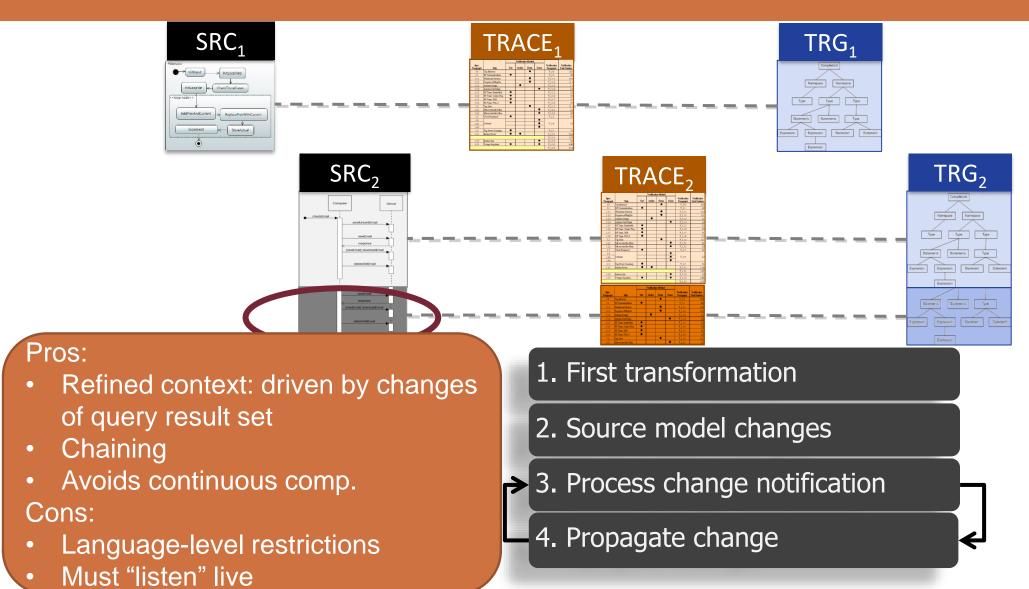
- Small-step incrementality
- Better performance

Cons:

- Highly depends on traceability links
- Smart matcher needed

- 1. First transformation
- 2. Source model changes
- 3. Detect missing trace links
- 4. Re-execute MT only for untraceable elements

Event Driven Transformations



Incremental Forward Transformation

Goals: reuse computations

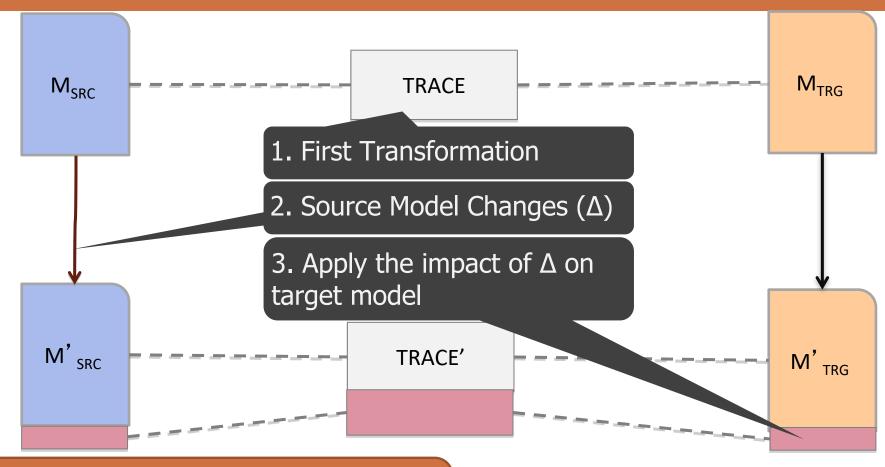
> Target Incrementality

- ...by reusing the unchanged parts of the target model
- Unchanged model elements does not need to be modified
- Anything derived from the unchanged part does not need to be regenerated (e.g., code)
- Change does not propagate further

> Source Incrementality

- by ignoring the unchanged parts of the source model
- Incremental pattern matchers.

Incremental Forward Transformation Setup

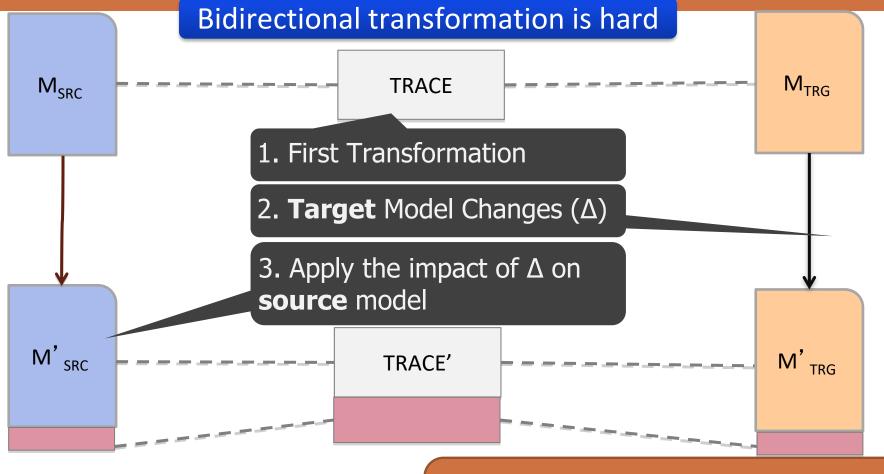


Typical Scenarios:

- Incremental Transformation for
- Tool integration

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Incremental backward transformation?



Some related work:

A. Schürr, P. Stevens, N. Foster, T. Hettel, Cicchetti&Pierantonio, Czarnecki&Diskin

Graph pattern matching, Graph transformation

Definitions

Graph pattern matching

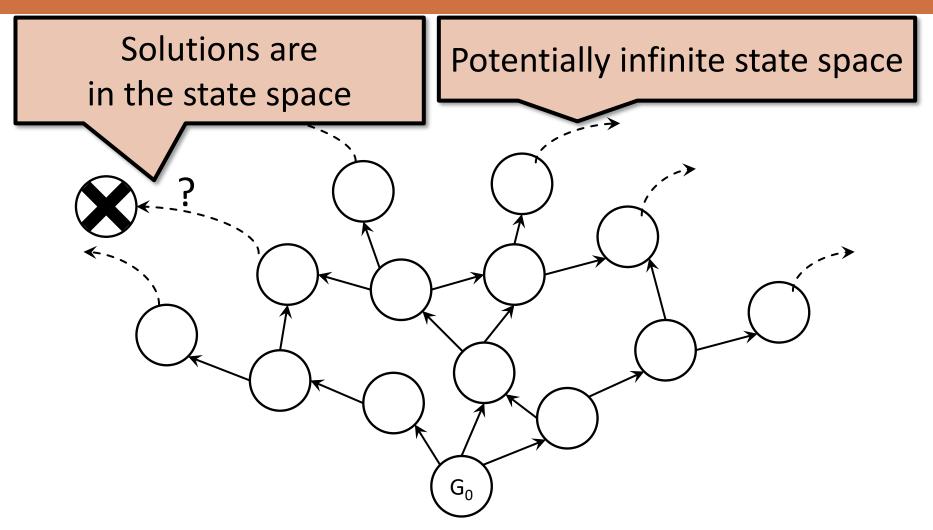
Model transformations

Incremental transzformations

Design space exploration

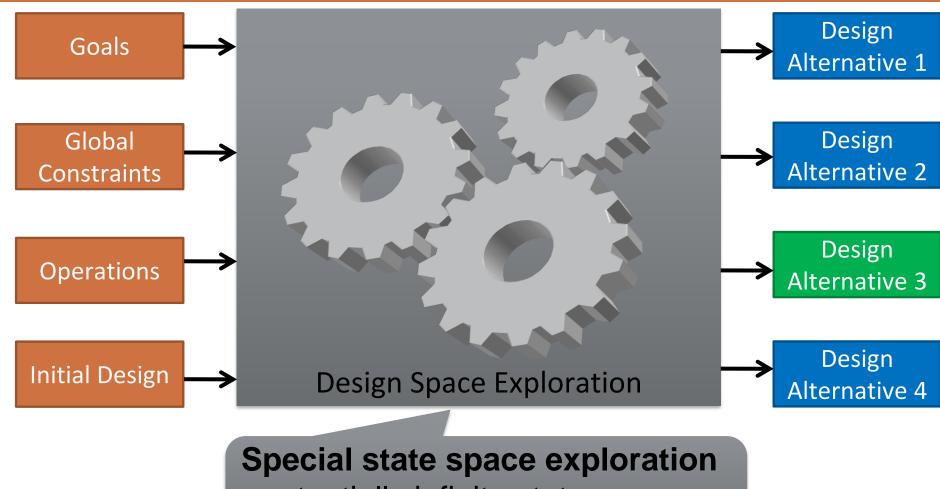


Revisit: state space of GT system



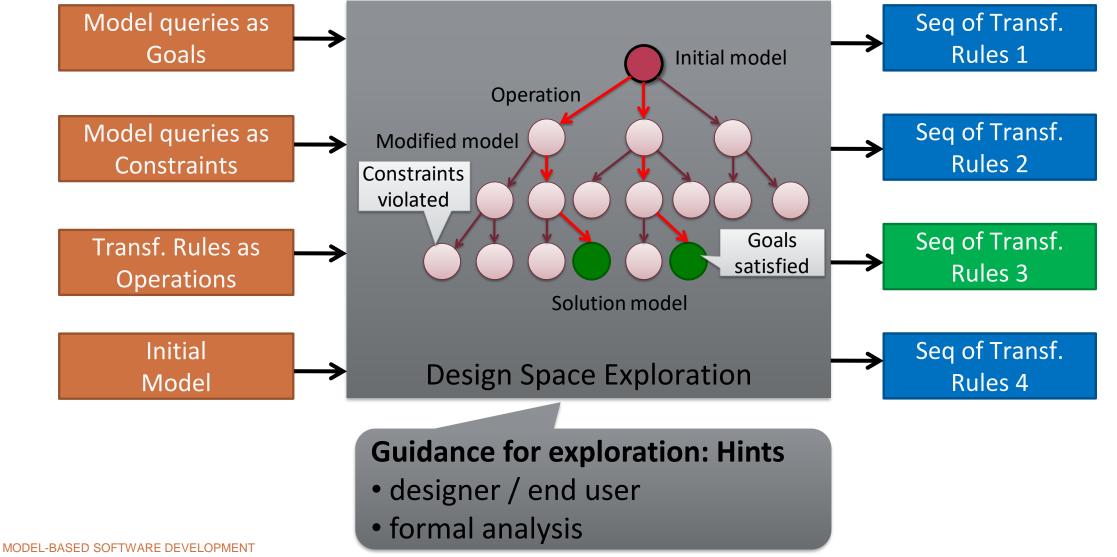
Initial Graph + GT rules → State Space

Design Space Exploration

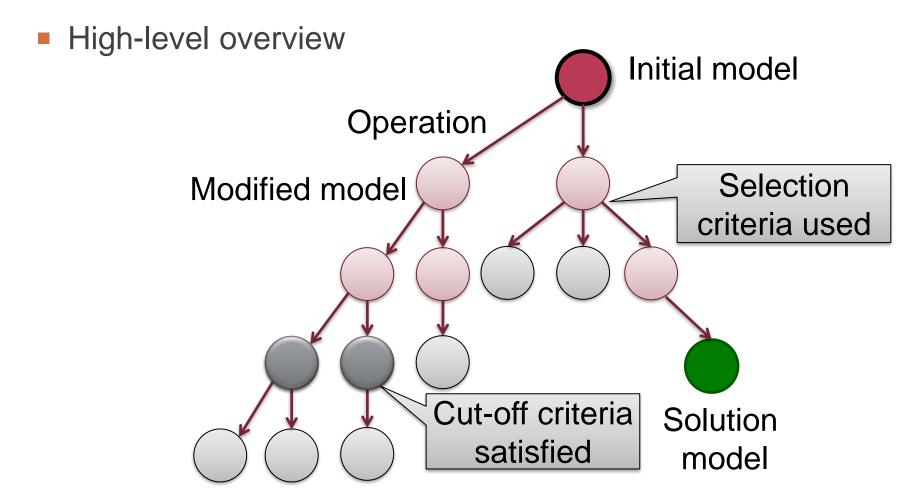


- potentially infinite state space
- "dense" solution space

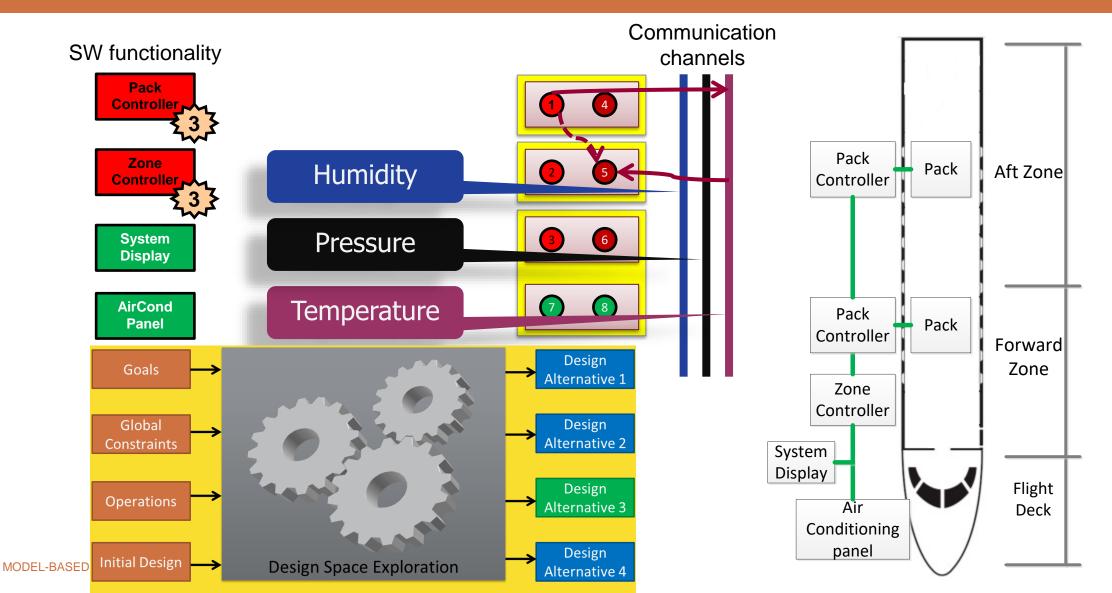
Model Driven Guided Design Space Exploration



Guided Design Space Exploration

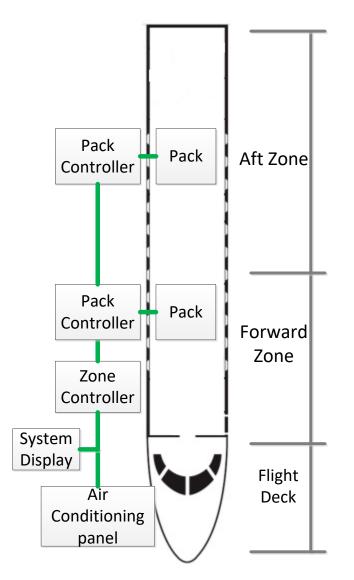


Design Space Exploration for IMA Configuration Design

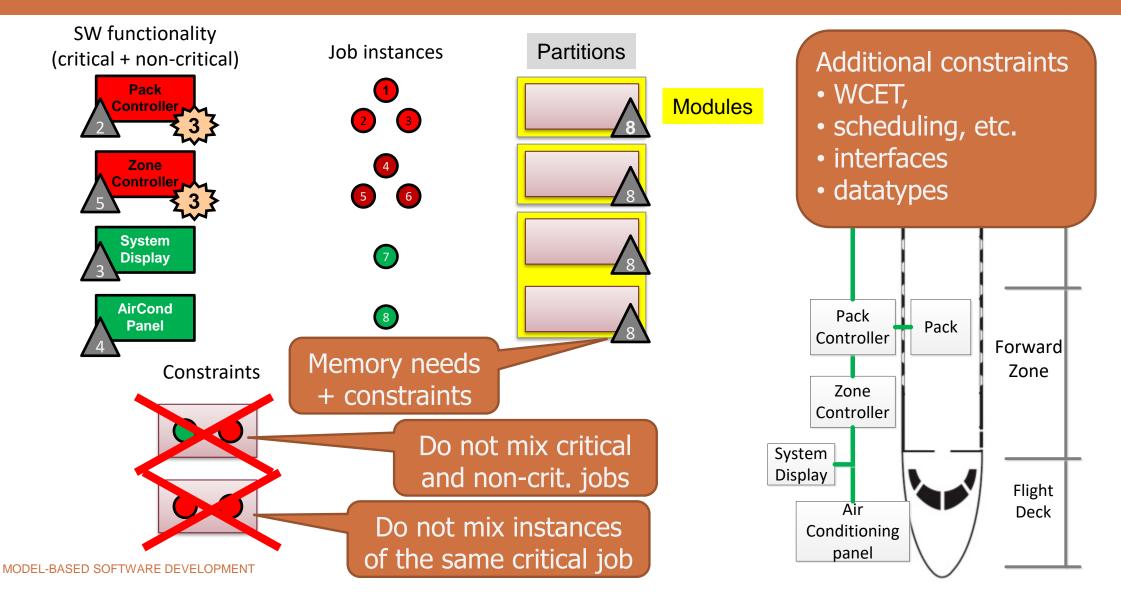


Designing ARINC653 configurations

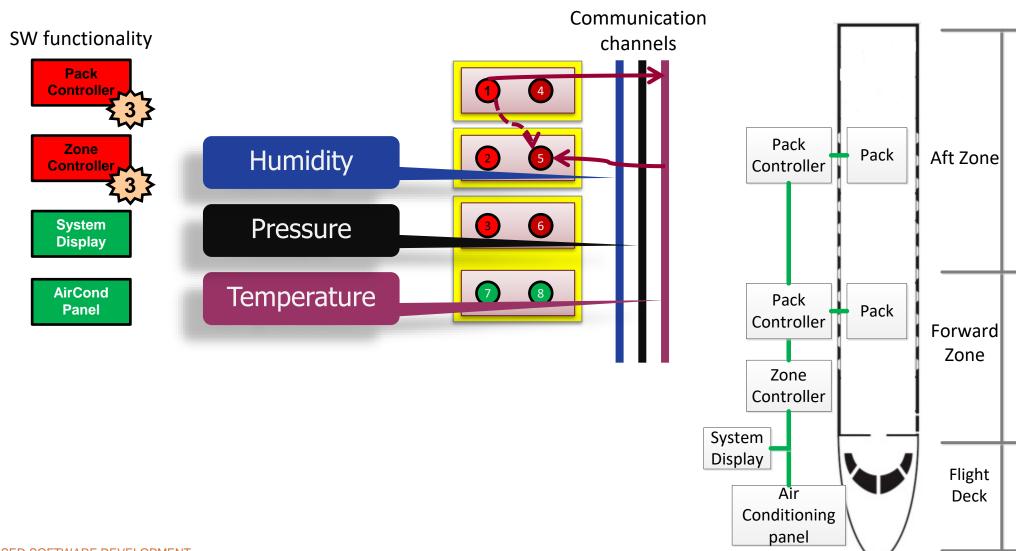
SW functionality (critical + non-critical) **Pack** Controller Supply fresh air Zone Supply hot air Controller **System** Monitor Display temperature **AirCond Panel** Set temperature Redundancy requirement



Job instances, Partitions, Modules



Allocating communication channels



Graph pattern matching, Graph transformation

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Graph pattern matching

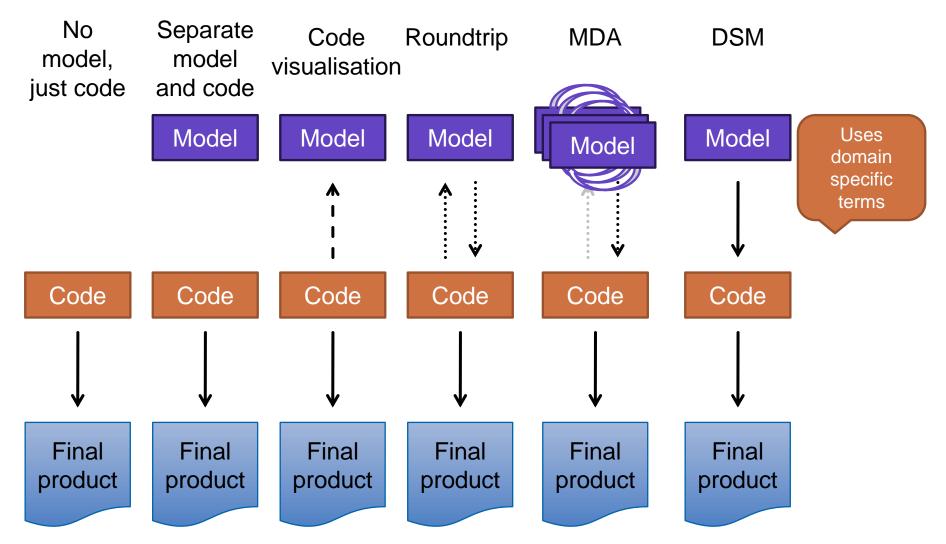
Model transformations

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How do we use models?





Thank you for your attention