



Model-based software development

Lecture IX.

Graph pattern
matching,
Graph transformation

Dr. Semeráth Oszkár

Graph pattern matching, Graph transformation

Definitions

Graph pattern matching

Model transformations

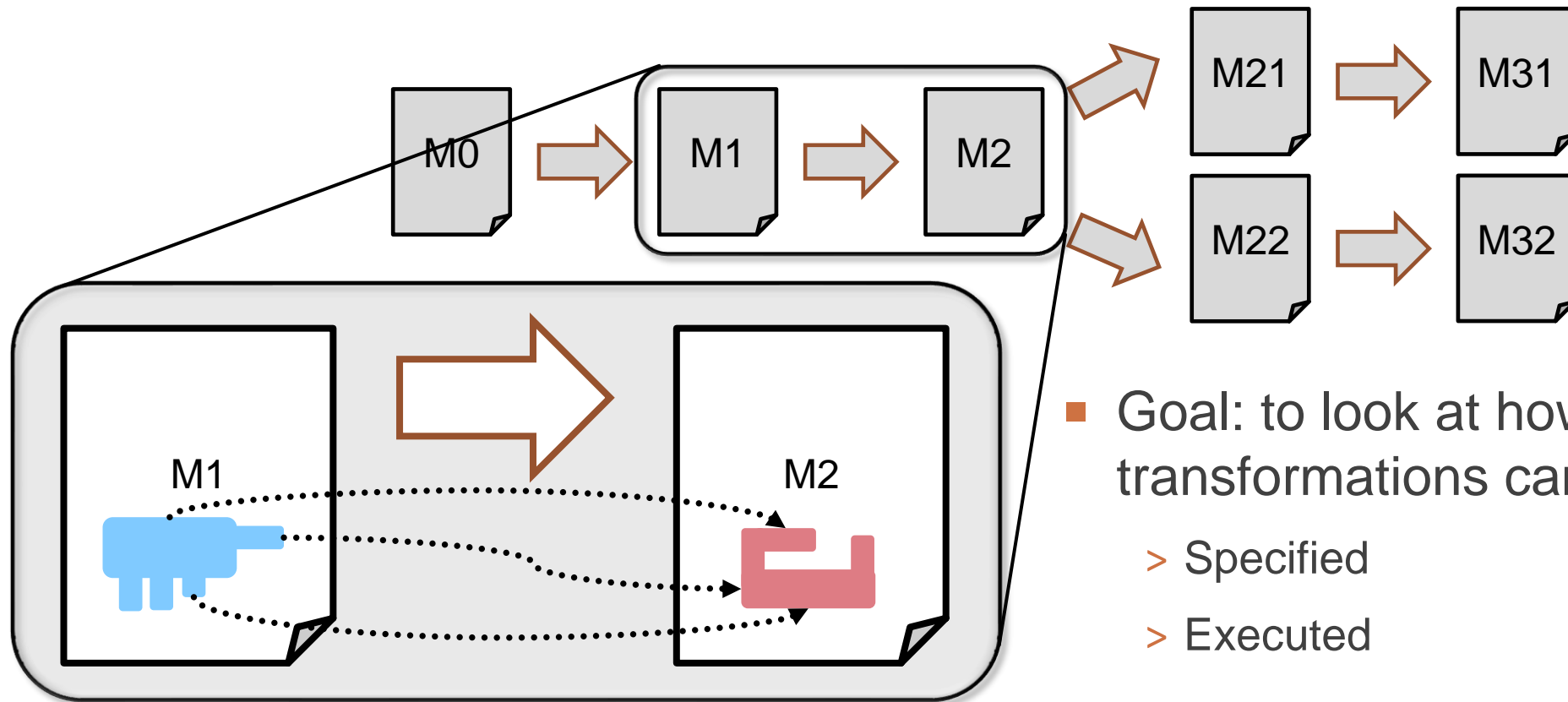
Incremental transformations

Design space exploration



Motivation: Transformation of models

- **Model-based development:** Models as primary documents
- Developing models, automating model processing

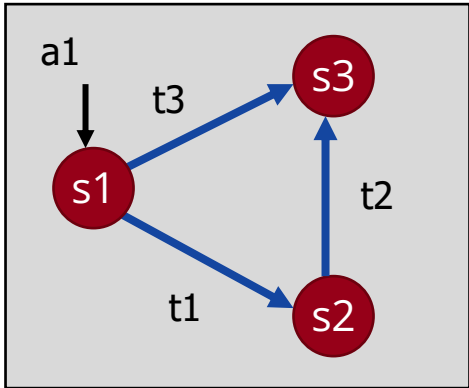


- Goal: to look at how model transformations can be
 - > Specified
 - > Executed

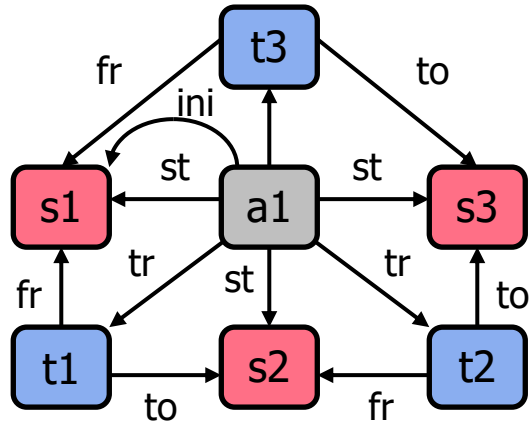
Abstract syntax

- How to modify the models?
- **Idea:** modify the representation of models directly → **Abstract syntax**

Concrete



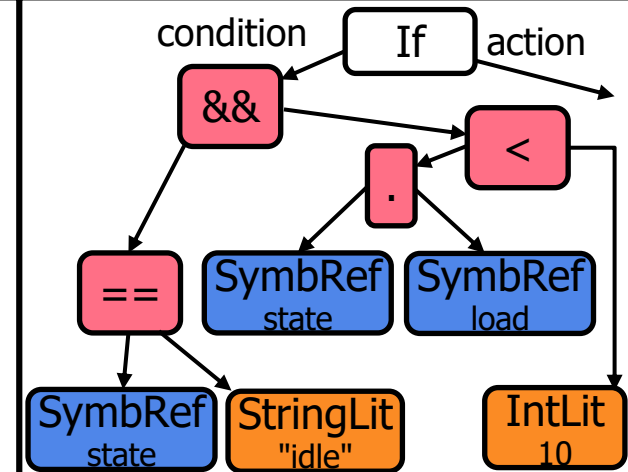
Abstract



Concrete

```
if (  
  state ==  
  "idle" &&  
  this.load < 10)  
  ...
```

Abstract



- **Task:** method to modify graphs!

Graph pattern matching, Graph transformation

Definitions

Graph pattern matching

Model transformations

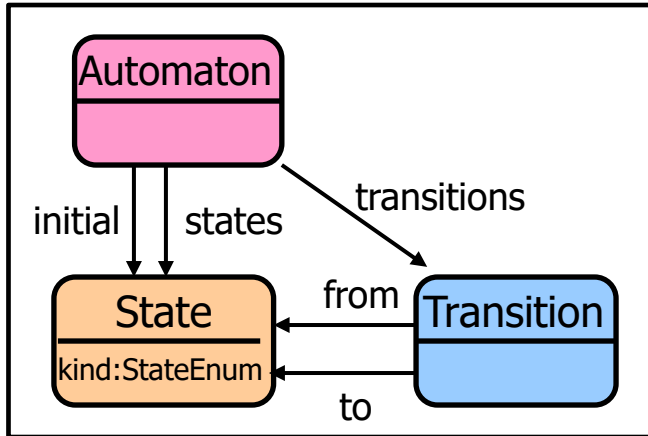
Incremental transformations

Design space exploration

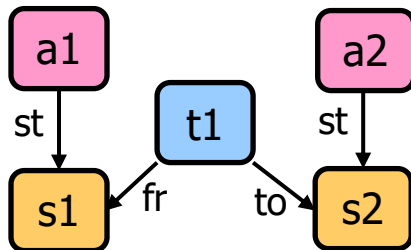


A simple example

Metamodel



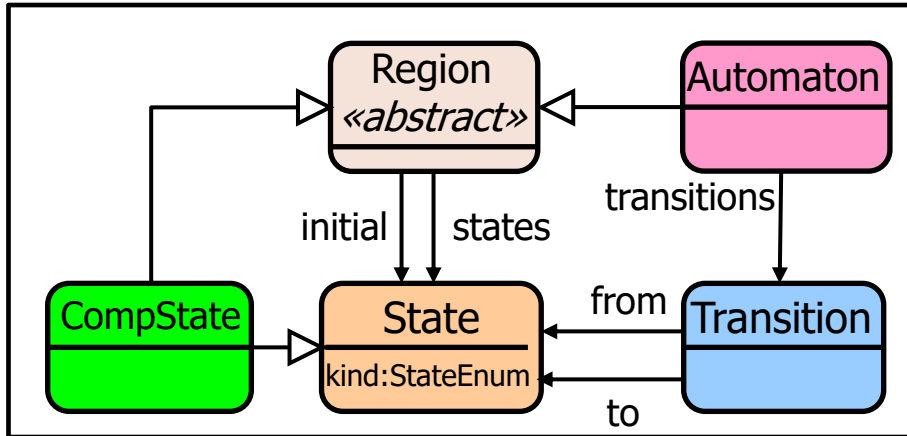
Violation example



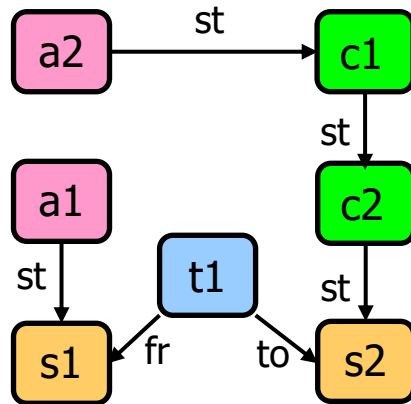
- Well-formedness constraint:
 - > Transition source & target states must be owned by same automaton
- Goal: to find violations...
 - > A violation is a *Transition*, whose „*from*” link points to a *State* *s1*, and „*to*” link points to a *State* *s2*, where the automaton of *s1* is not the automaton of *s2*

A more complex example

Metamodel



Violation example



- Well-formedness constraint:
 - > Transition source & target states must be owned by same automaton
- Goal: to find violations...
 - > A violation is a *Transition*, whose „from” link points to a *State* *s1*, and „to” link points to a *State* *s2*, where the automaton of *s1* is not the automaton of *s2*

Programmatic traversal vs. queries

- **Goal:** find constraint violations in model

Traverse model in general-purpose language

```
for (Automaton automaton : automatons) {  
    for (Transition transition : automaton.getTransitions()) {  
        State sourceState = transition.from;  
        // which automaton defines this state?  
        Automaton sourceAutomaton = null;  
        for (Automaton candidate : automatons) {  
            if (candidate.getStates().contains(sourceState)) {  
                sourceAutomaton = candidate;  
                break;  
            }  
        }  
        // ... do the same for targetState, then  
        if (sourceAutomaton != targetAutomaton)  
            // report violation  
    }  
}
```

„simple
example”

Programmatic traversal vs. queries

- Goal: find constraint violations in model
 - > Traverse model in general-purpose language
 - > Use a **Query DSL**
 - More concise
 - **Declarative** functional specification of the query
 - Freely interpreted by **query engine** (e.g. optimization)
 - Can be platform-independent
- Validation is just one use cases for **model queries**
 - > Derived features
 - > M2M/M2T Transformation, Simulation
 - > ...

Query Language Styles

- SQL-like (relational algebra)

- > Example: EMF Query
- > 😊 Good for attribute restrictions
- > ☹ Not very concise for relationships (many joins)

- Functional style

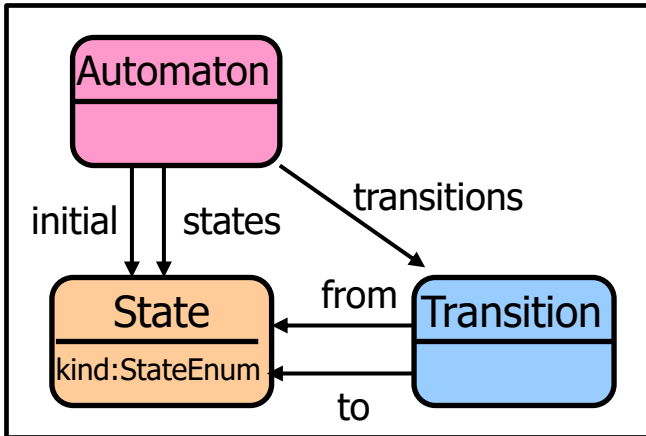
- > Example: OCL
- > Somewhat declarative

```
context Transition inv:  
  Automaton.allInstances()->forAll(a |  
    a.states->includes(self.from) =  
    a.states->includes(self.to)  
  );
```

- Logic style

- > Domain relational calculus / graph patterns / Datalog
- > Even more declarative

Metamodel

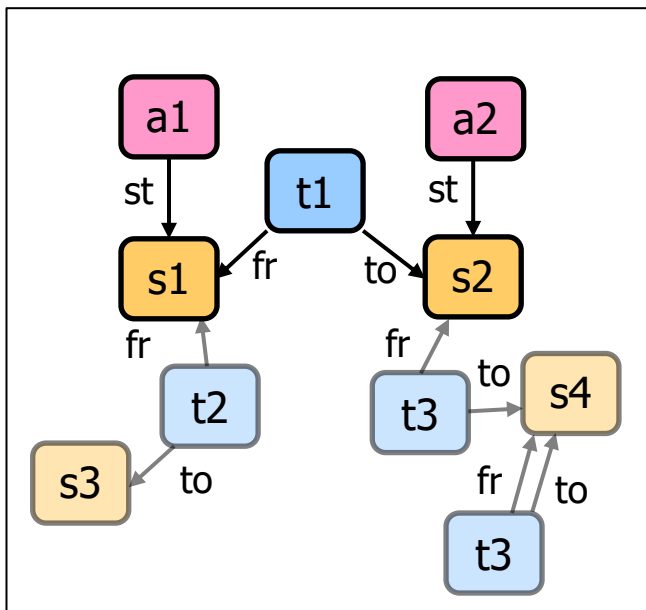


A violation is a *Transition*, whose „*from*” link points to a *State* *s1* and „*to*” link points to a *State* *s2*, where the automaton of *s1* is not the automaton of *s2*

in formal logic
(Domain
Relational
Calculus)

$$\{t \mid \exists s_1, s_2, a_1, a_2: Transition(t) \wedge from(t, s_1) \wedge to(t, s_2) \wedge states(a_1, s_1) \wedge states(a_2, s_2) \wedge a_1 \neq a_2\}$$

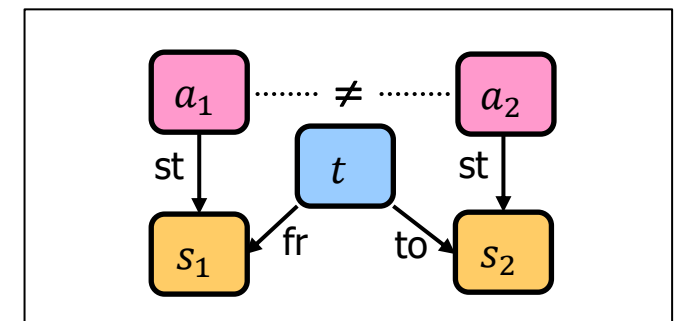
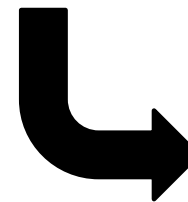
Violation example



violates(t) :-
Transition(t), from(t, s1), to(t, s2),
states(a1, s1), states(a2, s2), a1 ≠ a2

Datalog-like
query
languages

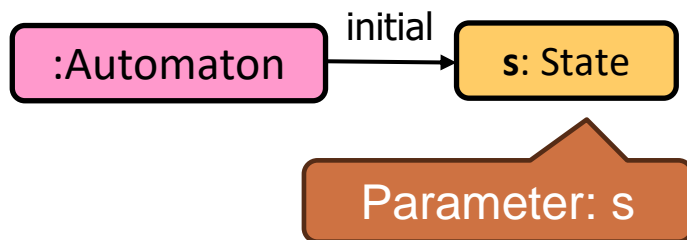
Query engine



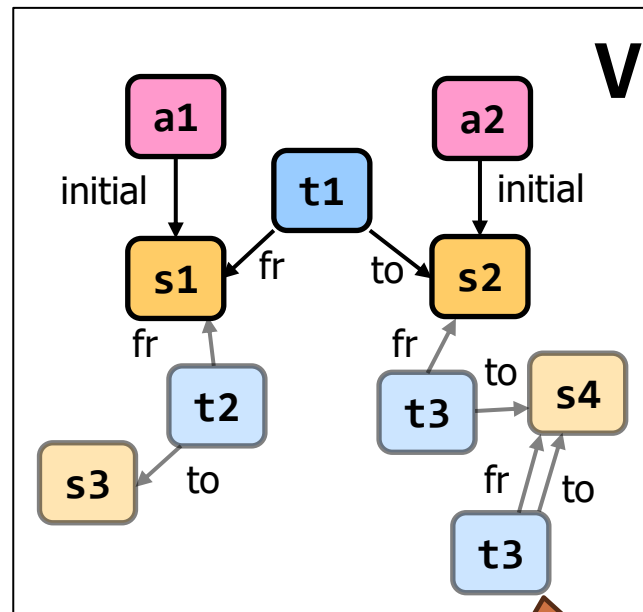
Pattern

Pattern matching

Pattern



Example model



Matches

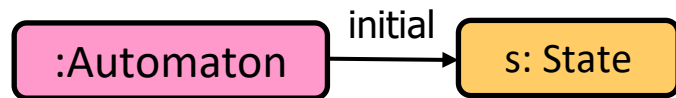
s
s1
s2

Parameter: s

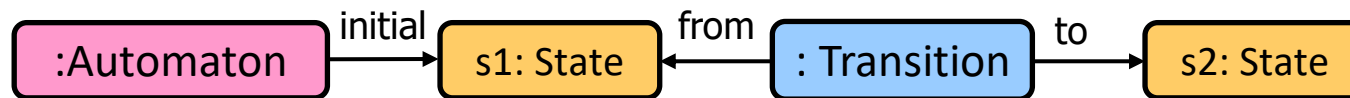
Model elements

Examples

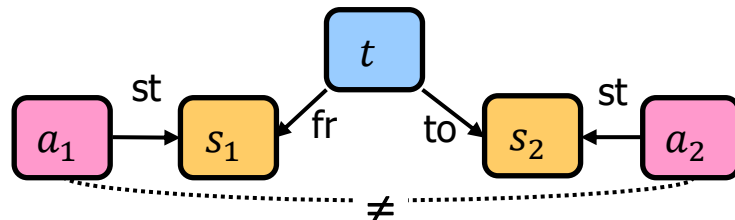
- **Simple example:** Initial states in the model



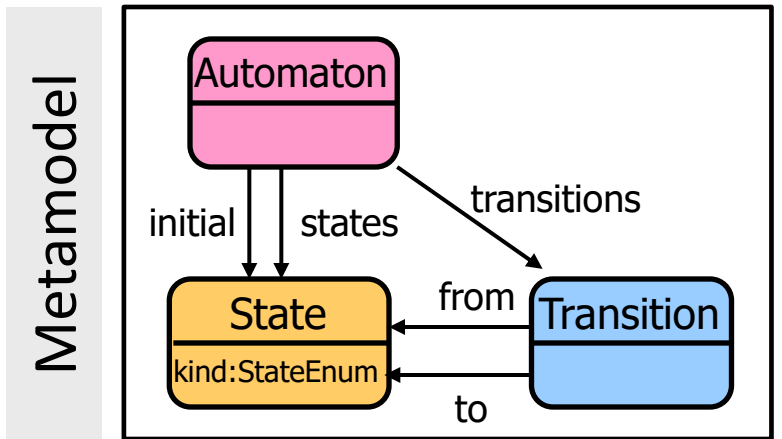
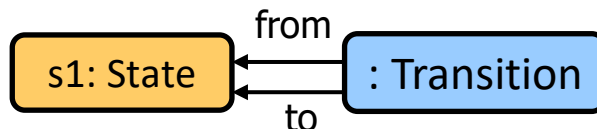
- **Chain (\wedge):** Second states in the model



- \neq : Transition across automata

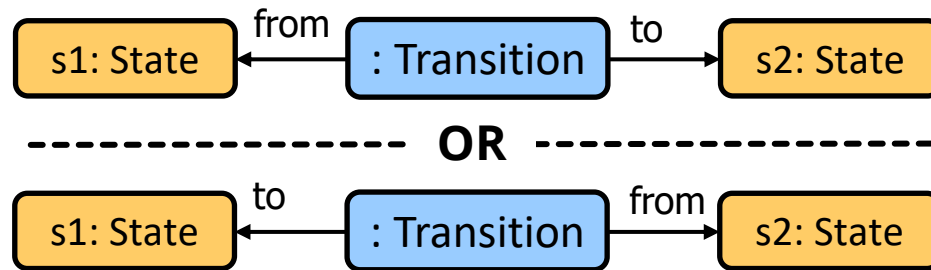


- $=$: Loop edge



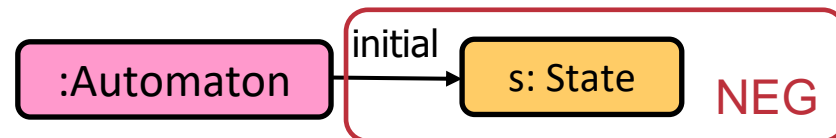
Examples 2

- (v): Two states are connected

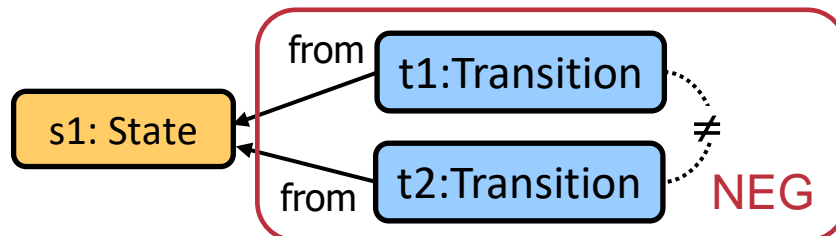


- (\neg , Negative Application Condition):

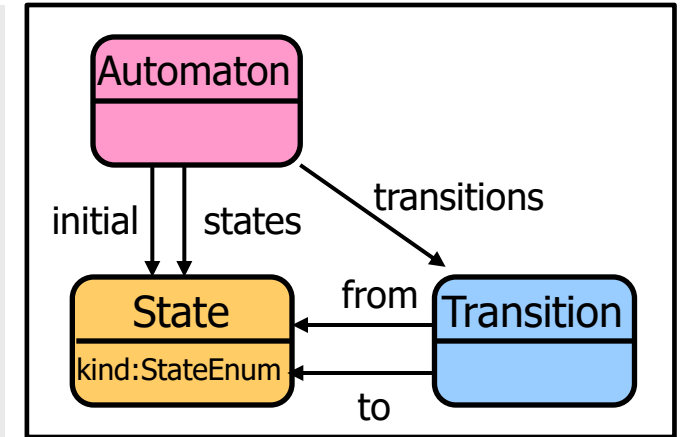
> automaton without initial state



> a state from whose initial state no two transactions leave (deterministic)



Metamodel



Graph pattern matching, Graph transformation

Definitions

Graph pattern matching

Model transformations

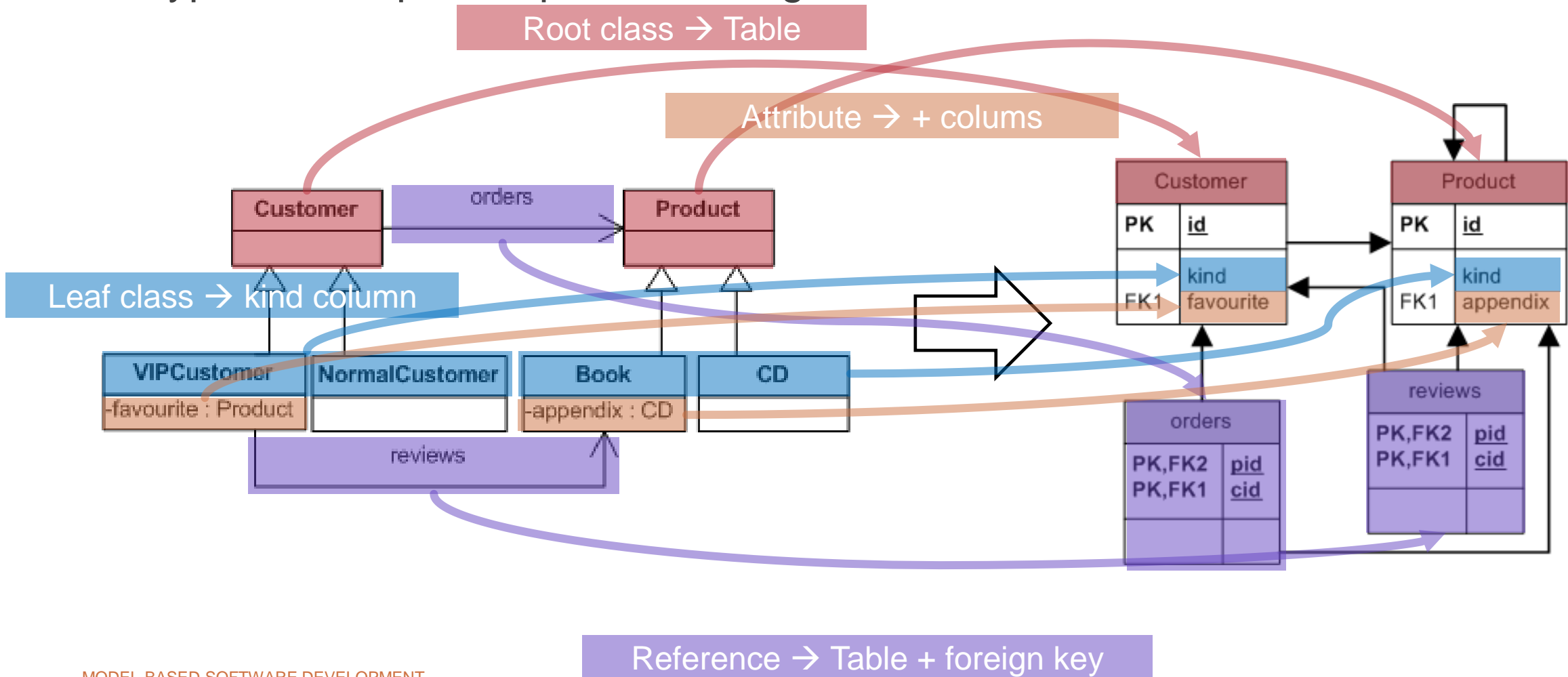
Incremental transformations

Design space exploration



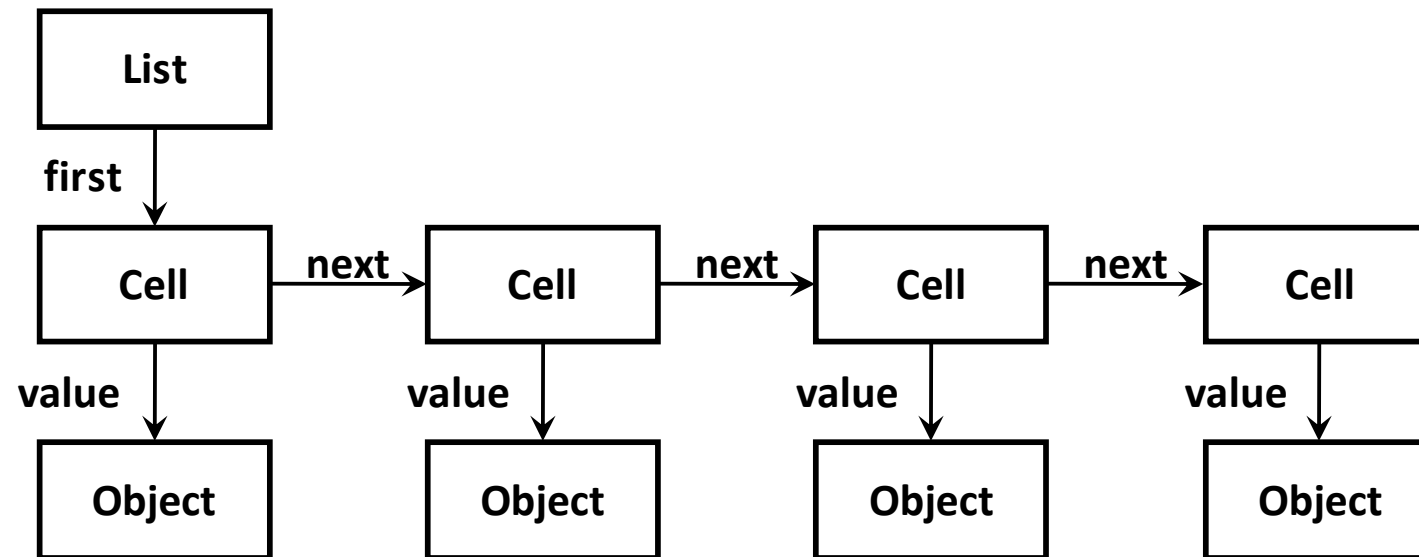
Example Transformation

- Typical example: map a class diagram to database tables!



Graph transformation

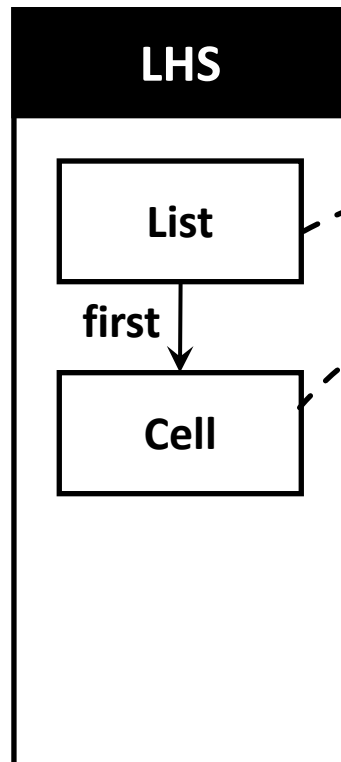
- Model = Labelled Graph



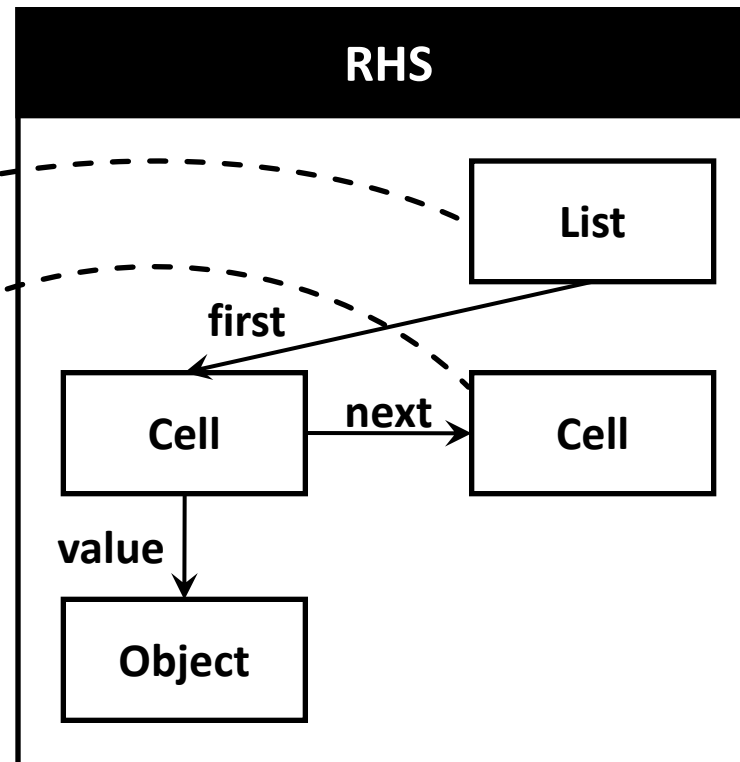
Graph Transformation rule

- Graph rewriting rule, defined with two graphs

Left Hand Side

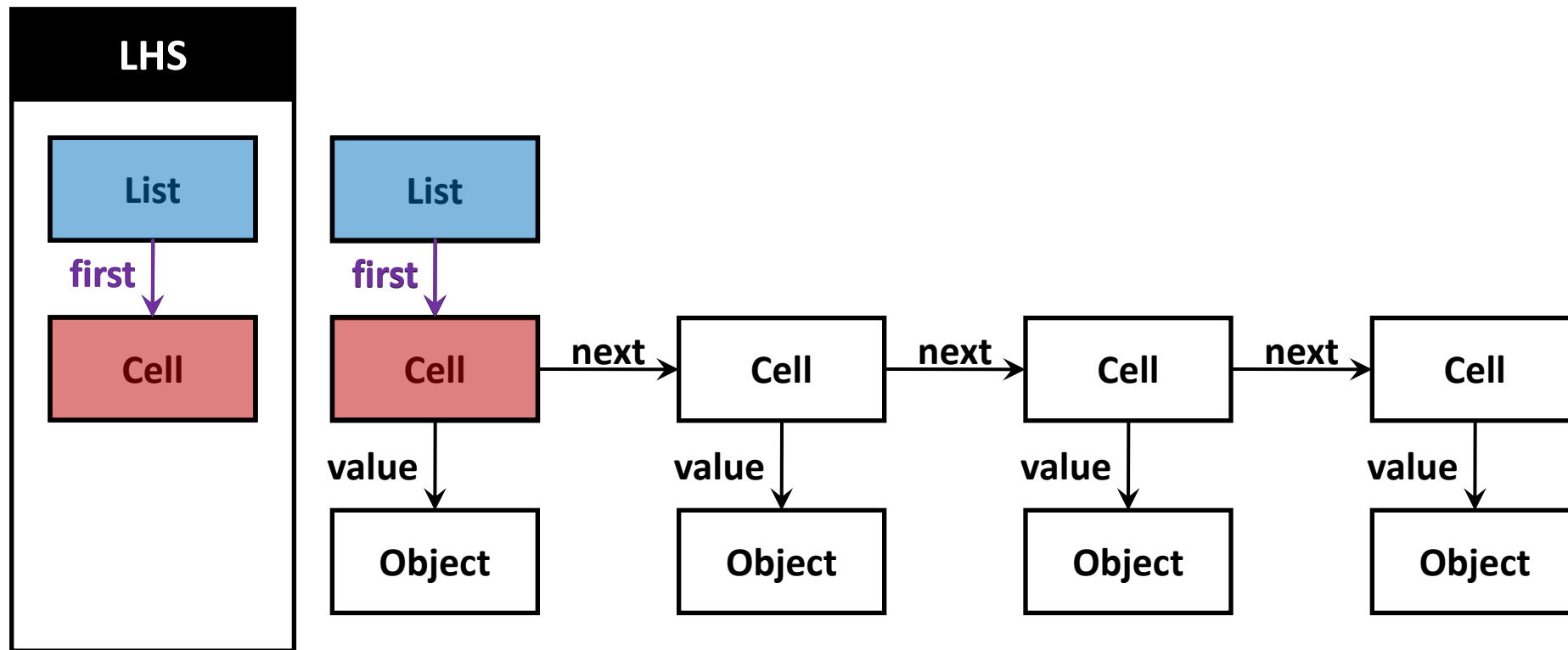


Right Hand Side



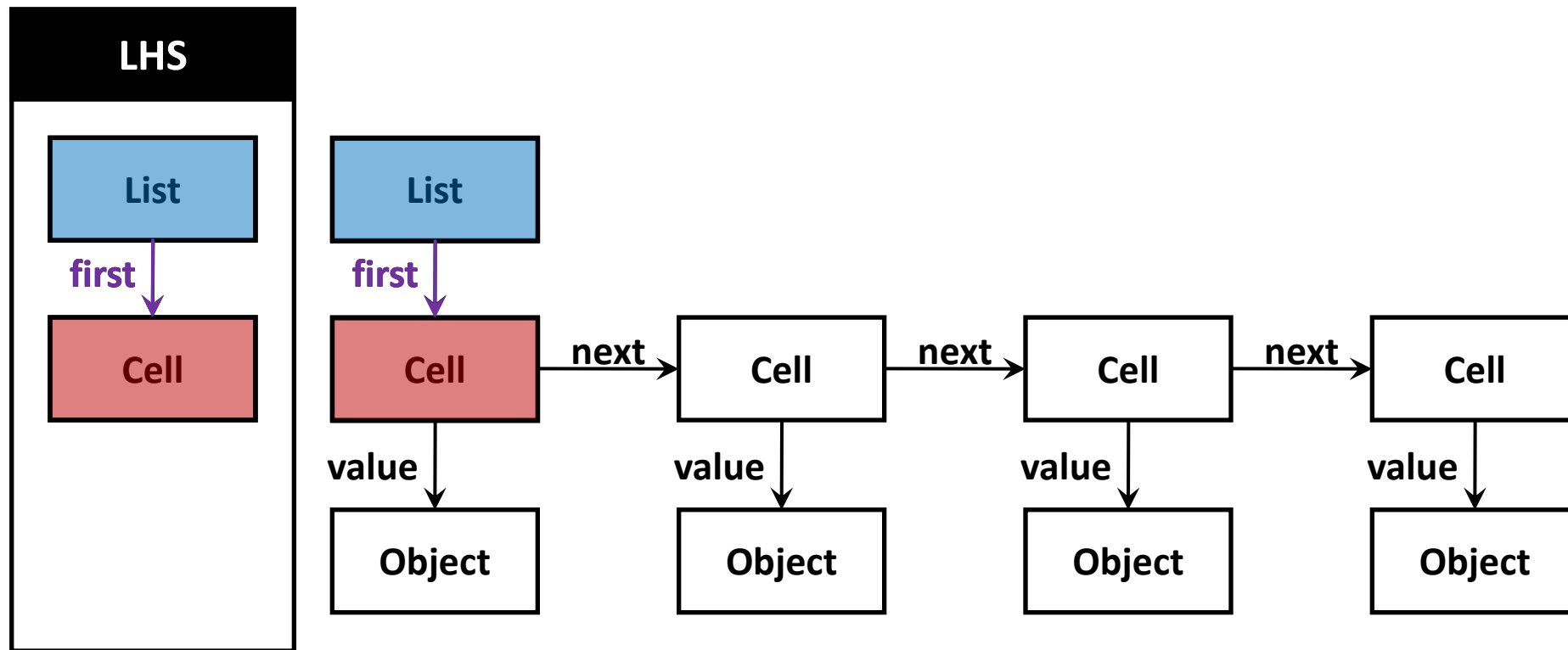
Graph Transformation: Pattern matching

- **Matching:** find the subgraphs containing LHS in the source graph



Graph Transformation: Pattern matching

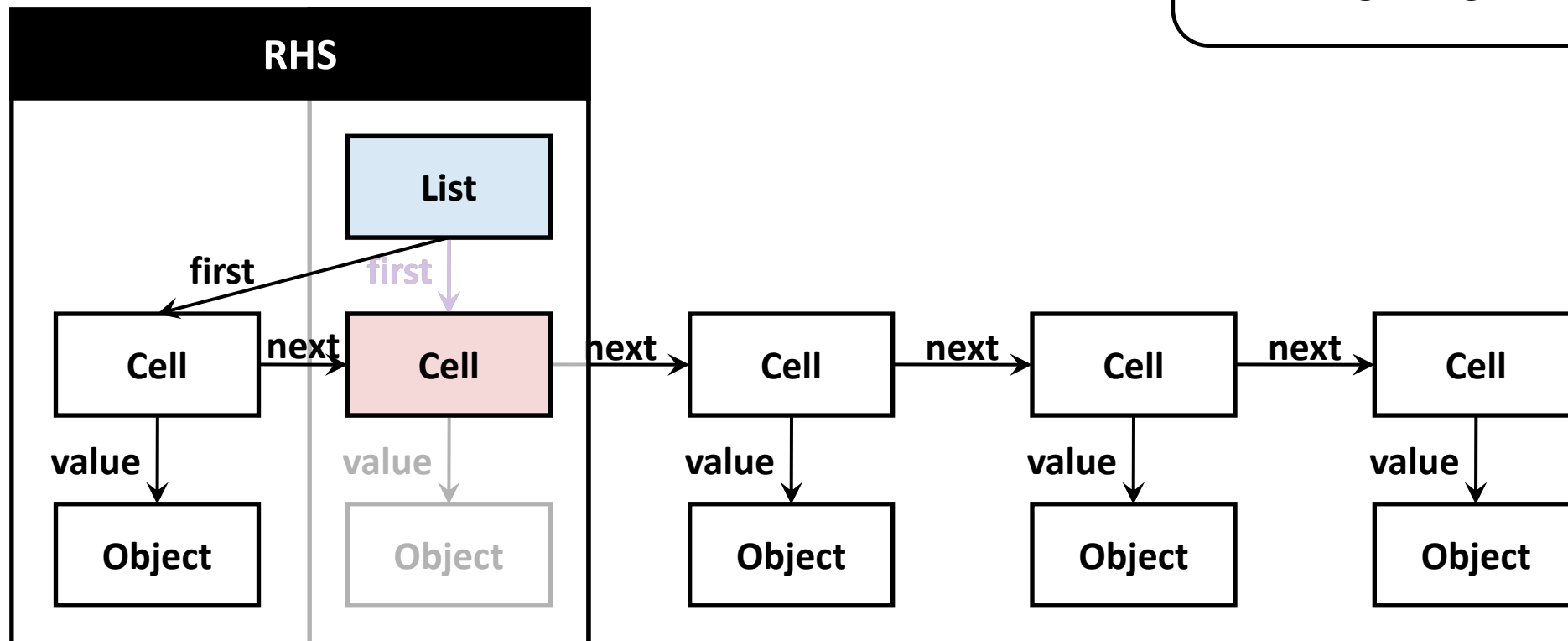
- **Matching:** find the subgraphs containing LHS in the source graph



Graph Transformation: Execution of rewriting

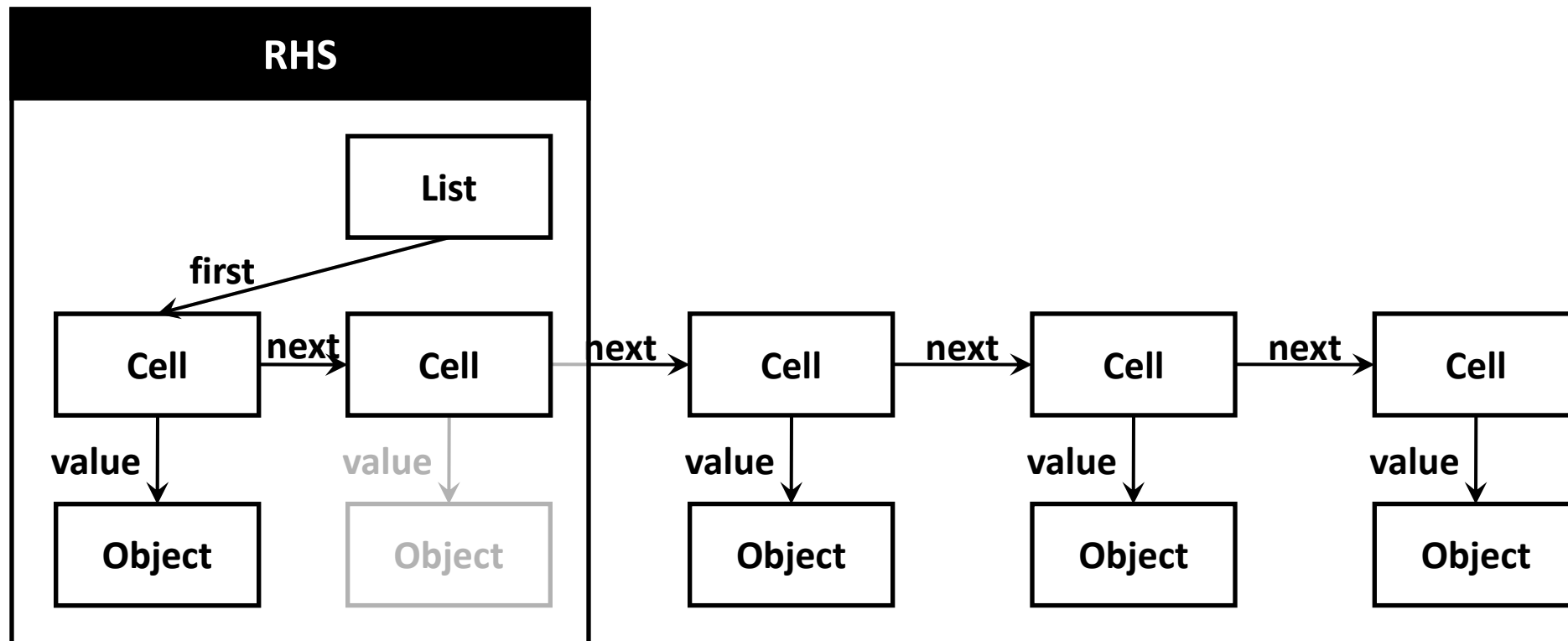
- Rewriting the graph by the match: replace LHS with RHS.

LHS\RHS → Delete
RHS\LHS → Insert
RHSnLHS → Leave it



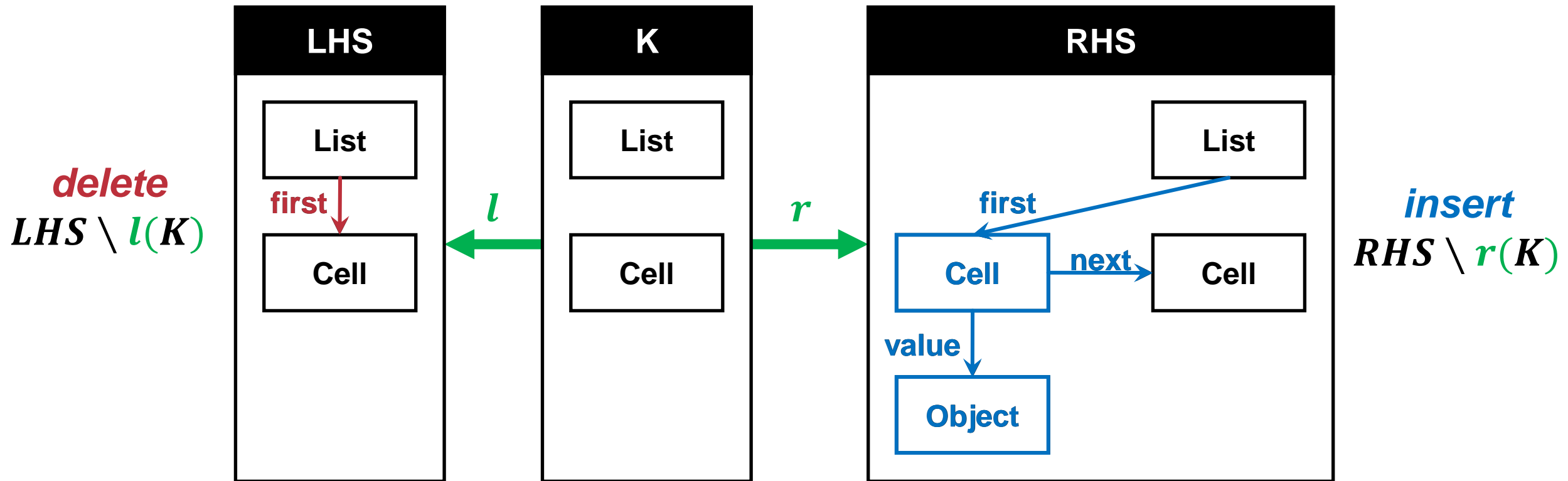
Graph Transformation: Execution of rewriting

- We get a new graph

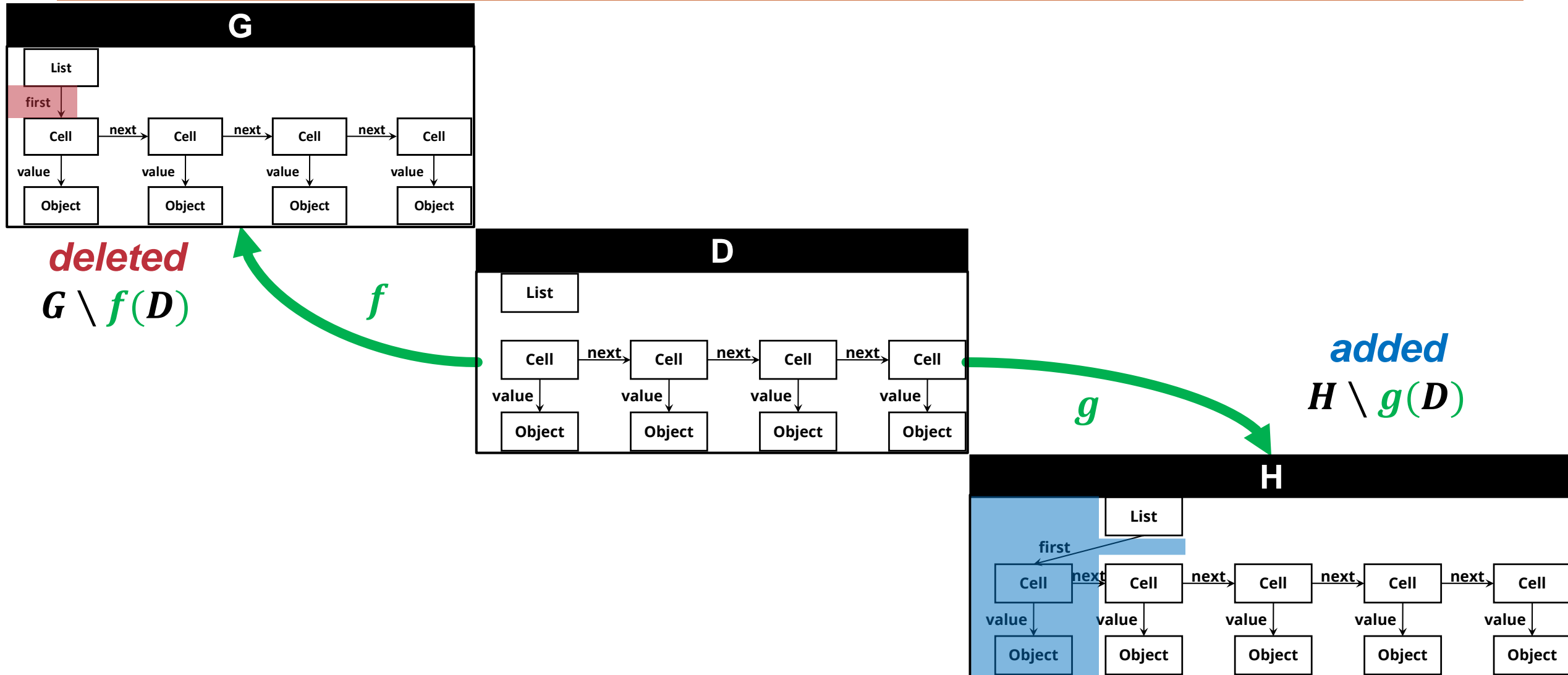


Anatomy of graph transformation

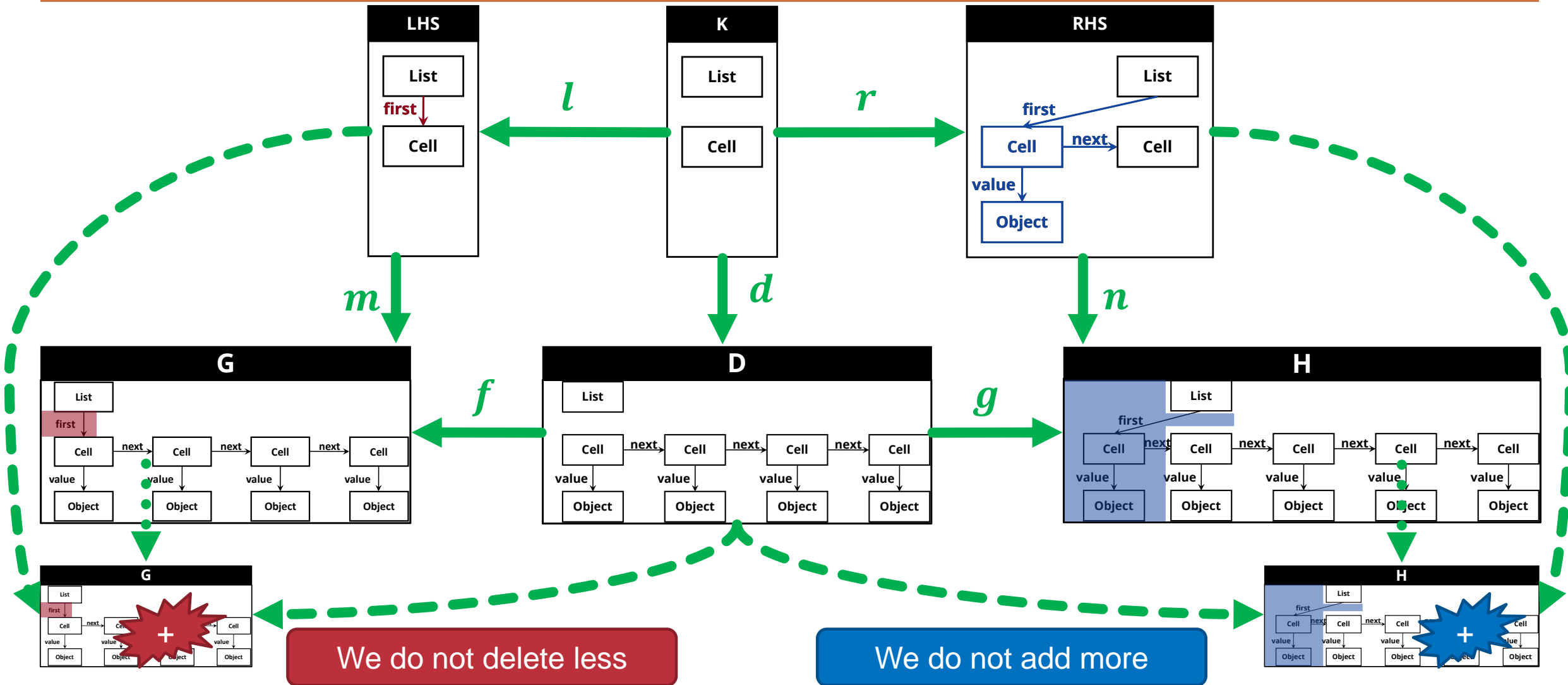
- Let us examine which graph can be matched to which graph during the transformation.



Anatomy of transformed models



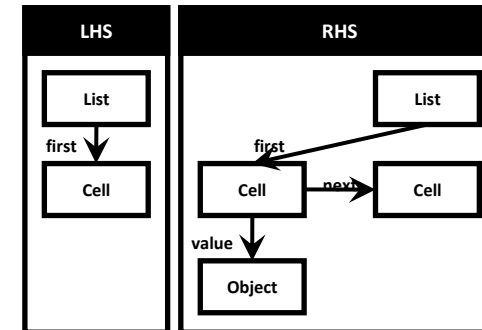
Complete anatomy



Graph Transformation

- Formulating rules for rewriting models
- Extending grammar rules

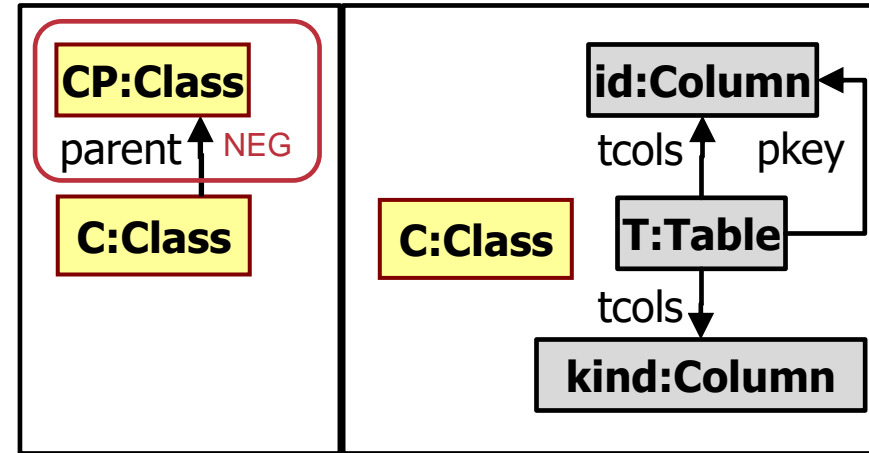
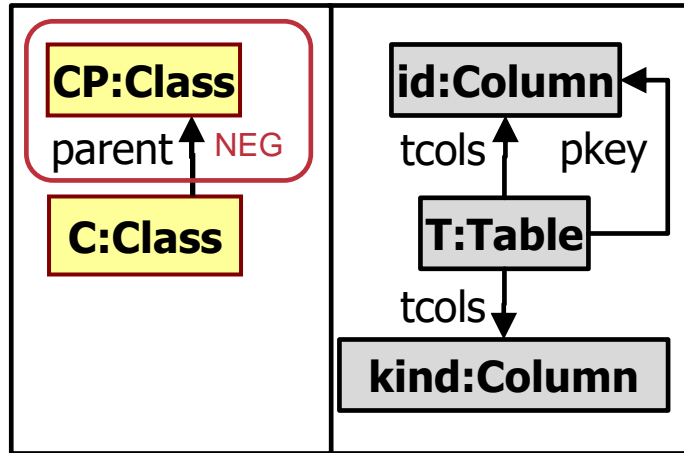
$List \rightarrow List, Cell$ **vs**



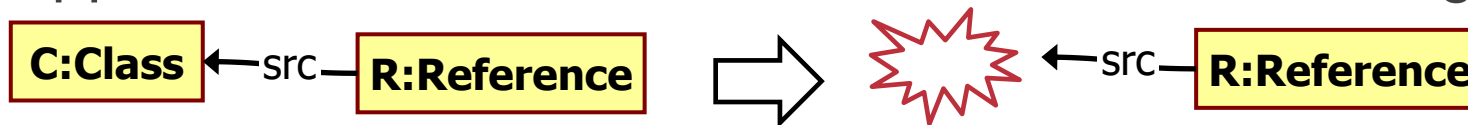
- Clear but mathematically precise formalism
(Termination, Ordering, Confluence, ...)
- Tool support (see previous practice)

“Dangling edge” problem

- Mapping of ancestor classes (with and without deletion)



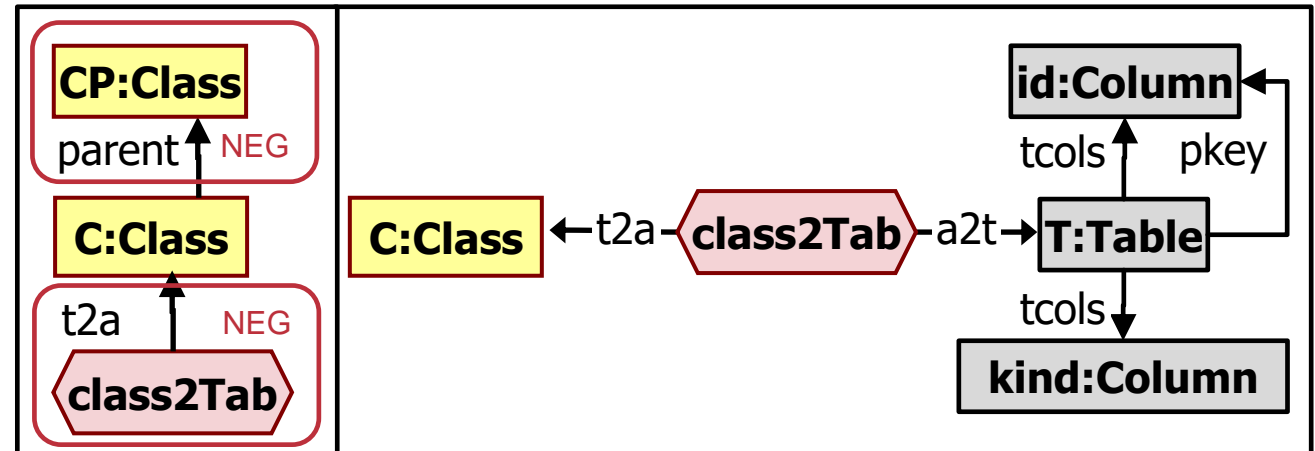
- What happens if we delete an element that still has an edge pointing to it?



- Resolving „dangling edges”: Delete edges / Undo transformation

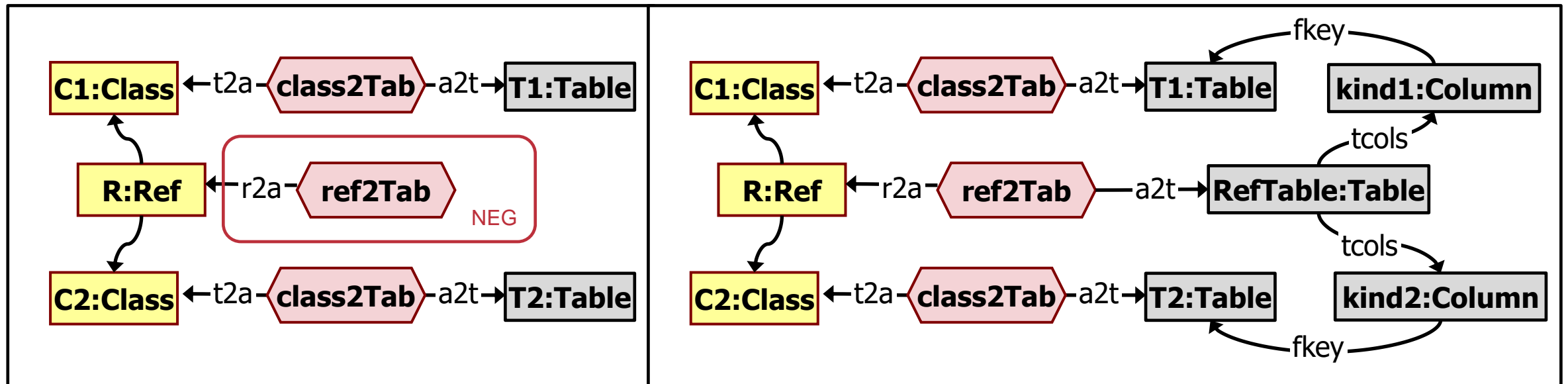
Examples

- Mapping of ancestor classes with traceability :
 - > Find an ancestor class
 - > that has not yet been mapped,
 - > then map it.



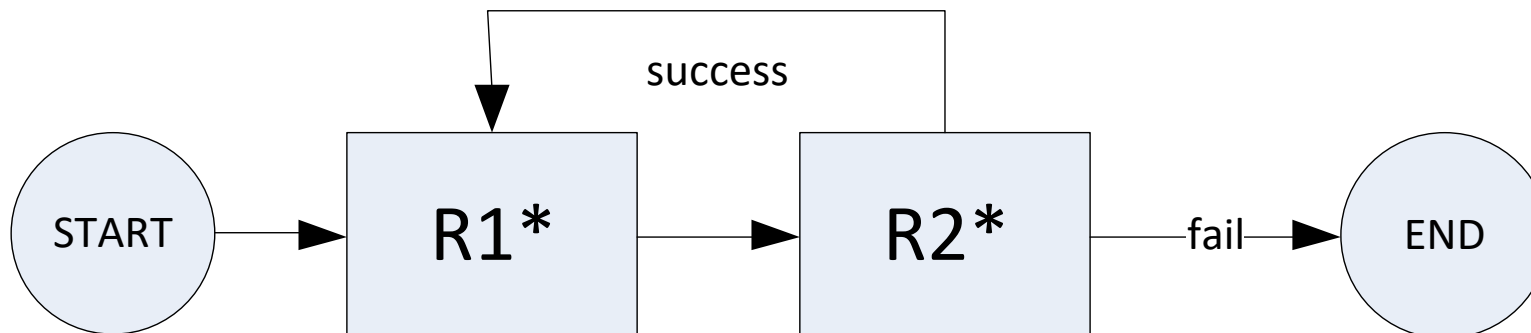
Examples

- Mapping references

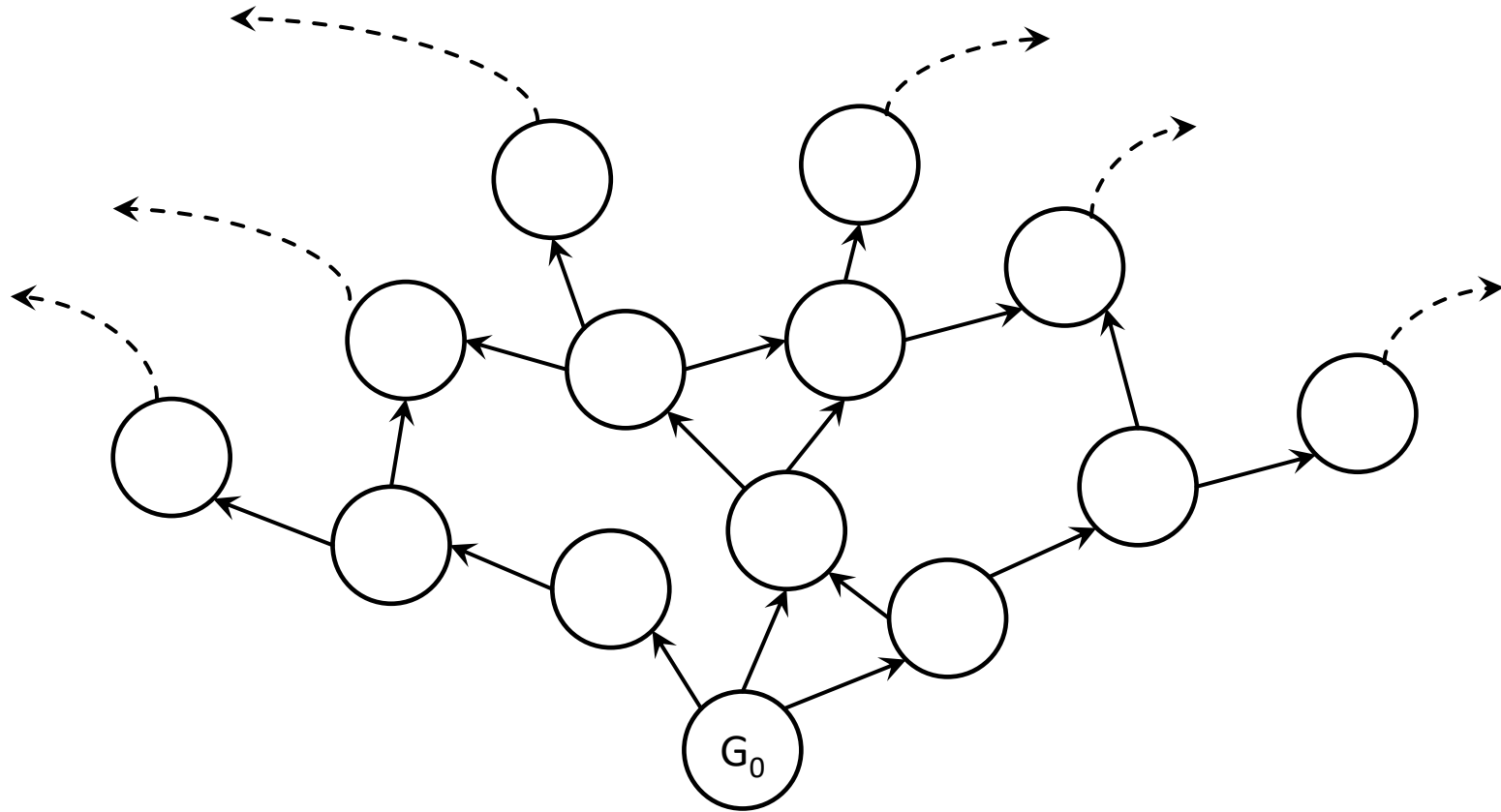


Control mechanisms

- In what order should the rules be executed?
- Multiple options, see previous lecture.
- But for example:
 - > Fire arbitrary transformations as long as possible (~ default)
 - > Fire all possible transformations once
 - > control graph (explicit control)

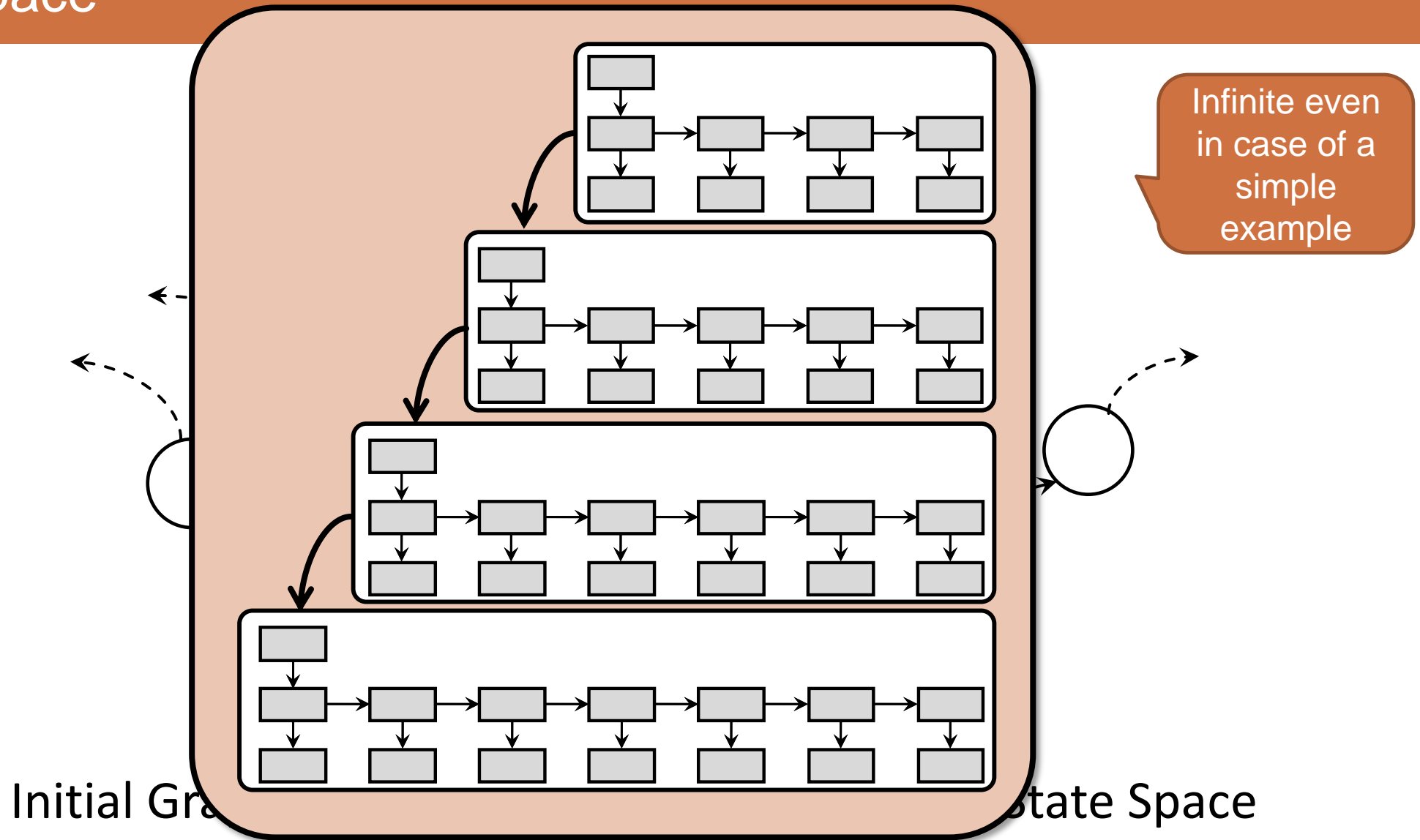


State Space



Initial Graph + GT rules \rightarrow **(Typically infinite)** State Space

State Space



Types of model transformations

- By number of inputs and outputs
(In-place vs out-place)
- By the language
(Endogenous vs exogenous)
- By the direction
(unidirectional vs bidirectional)

Graph pattern matching, Graph transformation

Definitions

Graph pattern matching

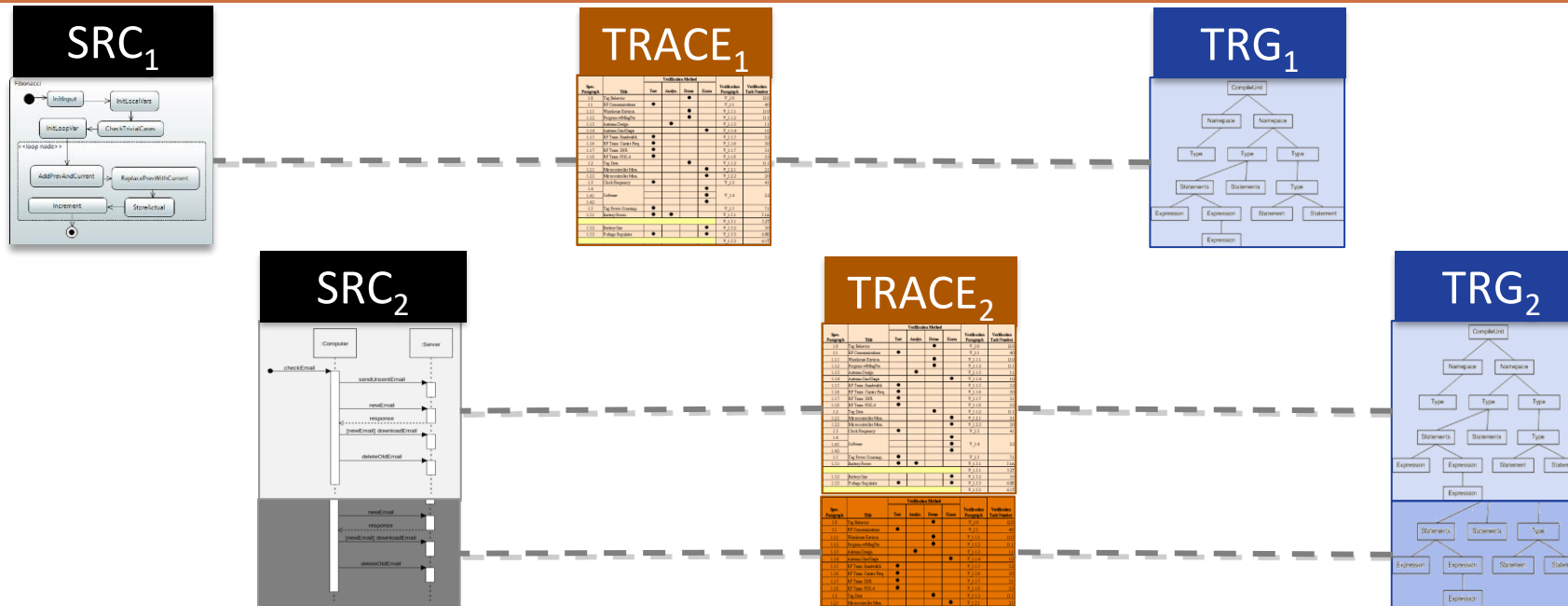
Model transformations

Incremental transformations

Design space exploration



No Incrementality: Batch Transformations

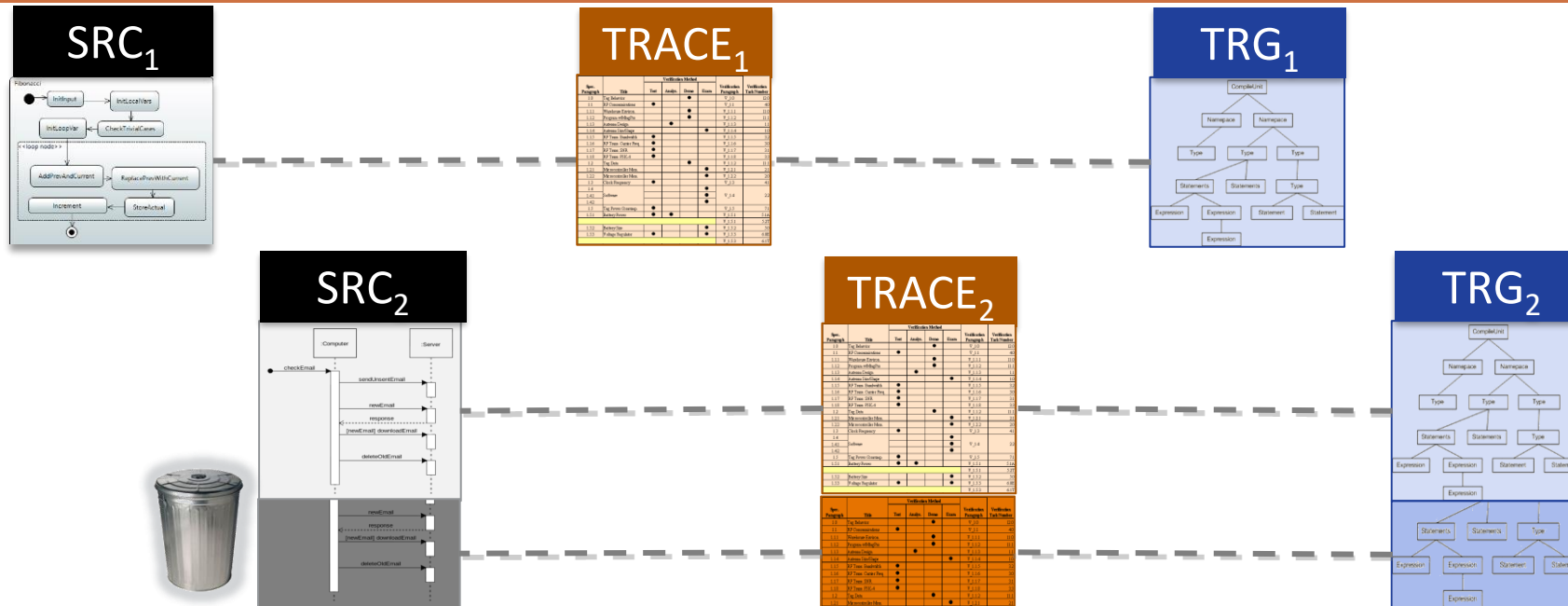


1. First transformation

2. Source model changes

3. Re-execute from scratch for all source models

Dirty Incrementality



Pros:

- Large-step incrementality
- Avoids continuous execution

Cons:

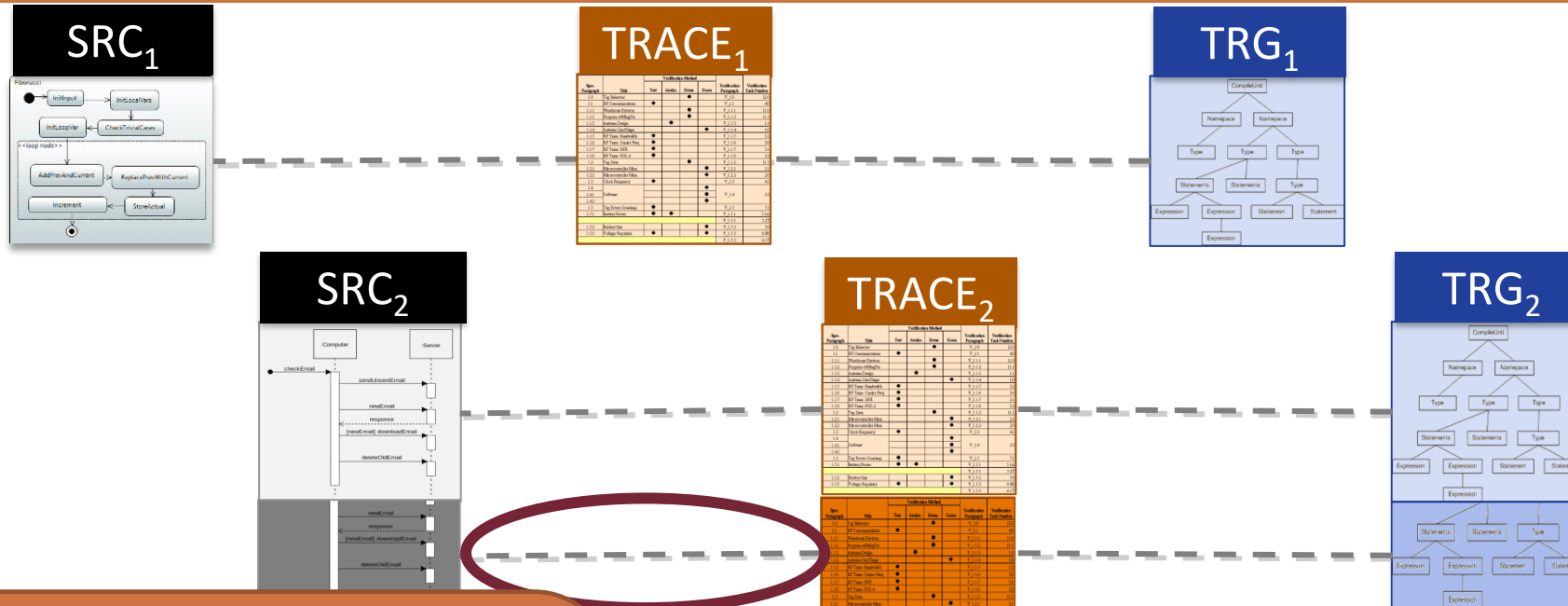
- Complex MT can be slow
- Cleanup (after an error)?
- Chaining?

1. First transformation

2. Source model changes

3. Re-execute from scratch only for changed models

Incrementality by Traceability



Pros:

- Small-step incrementality
- Better performance

Cons:

- Highly depends on traceability links
- Smart matcher needed

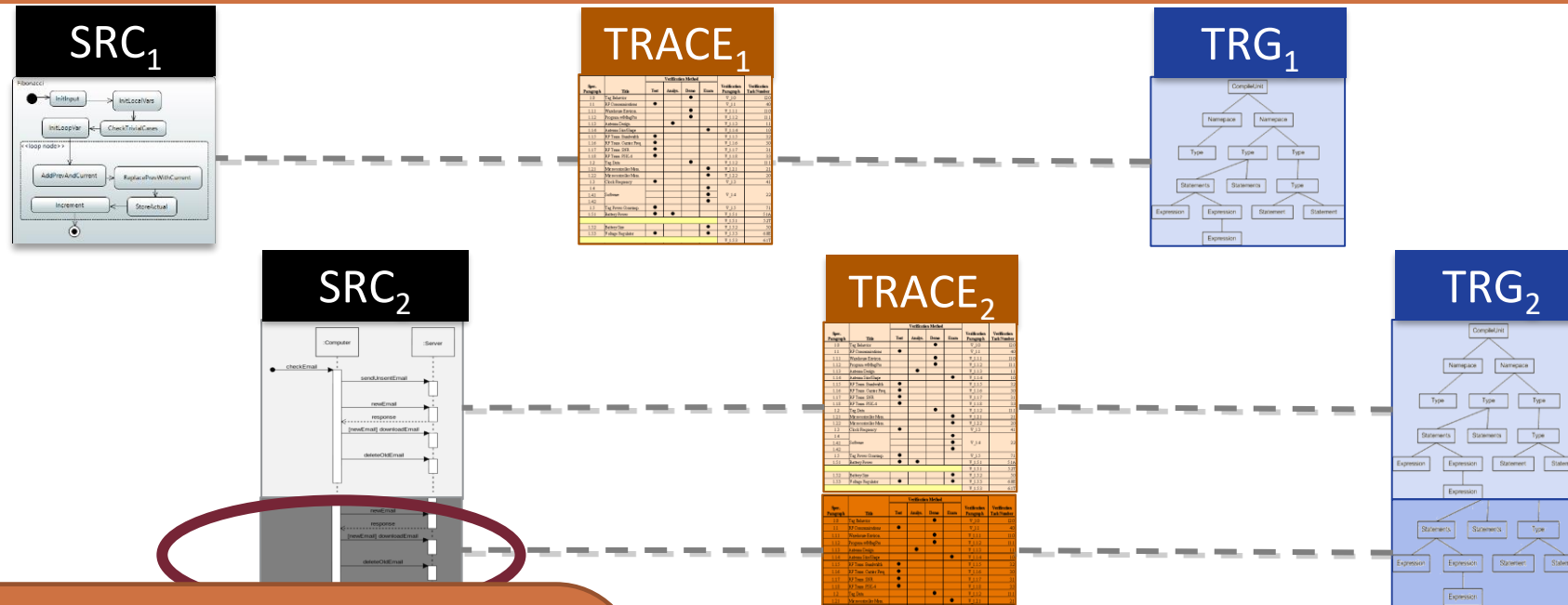
1. First transformation

2. Source model changes

3. Detect missing trace links

4. Re-execute MT only for untraceable elements

Event Driven Transformations



Pros:

- Refined context: driven by changes of query result set
- Chaining
- Avoids continuous comp.

Cons:

- Language-level restrictions
- Must "listen" live

1. First transformation

2. Source model changes

3. Process change notification

4. Propagate change

Incremental Forward Transformation

- Goals: reuse computations

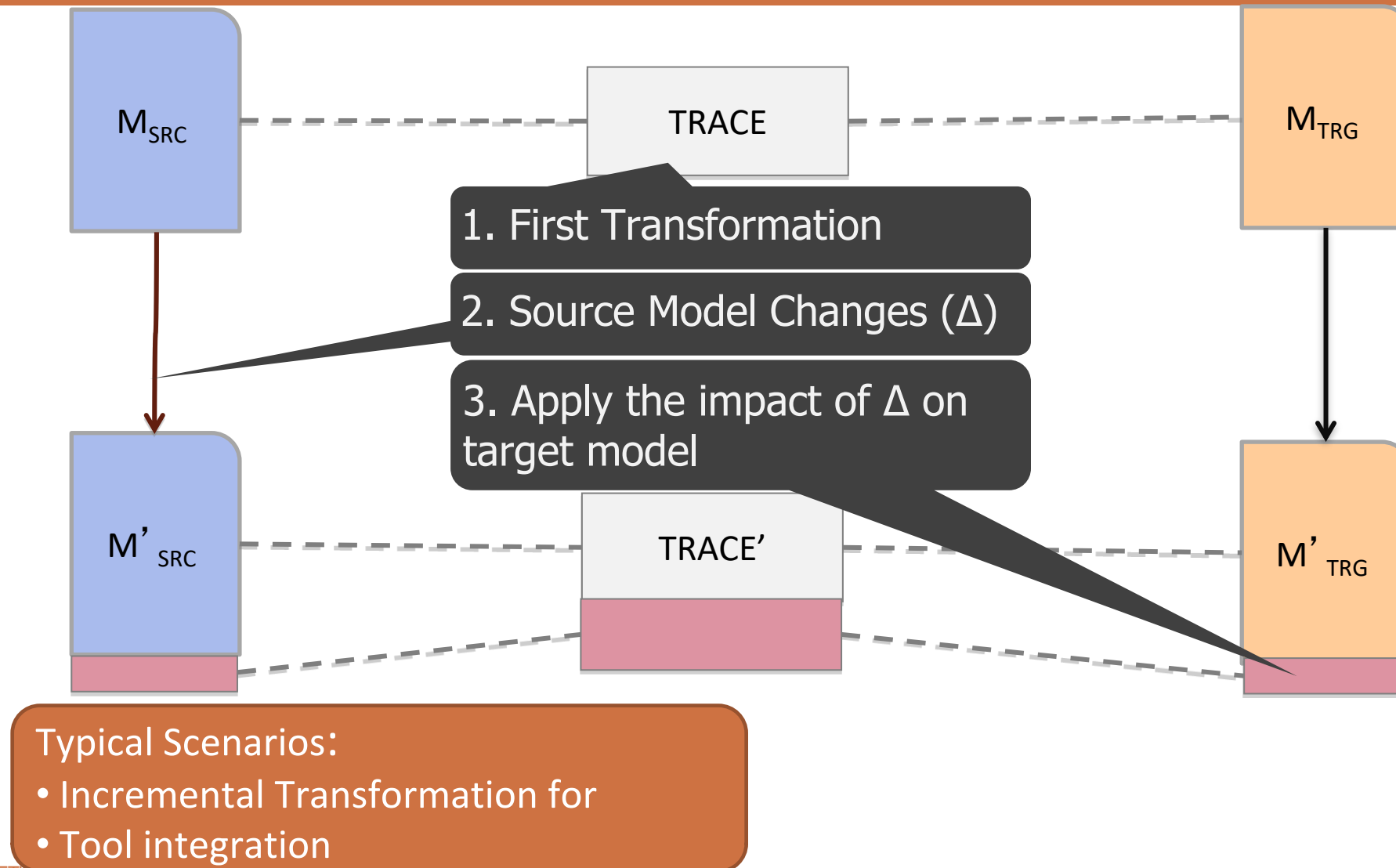
- > **Target Incrementality**

- ...by reusing the unchanged parts of the target model
 - Unchanged model elements does not need to be modified
 - Anything derived from the unchanged part does not need to be regenerated (e.g., code)
 - Change does not propagate further

- > **Source Incrementality**

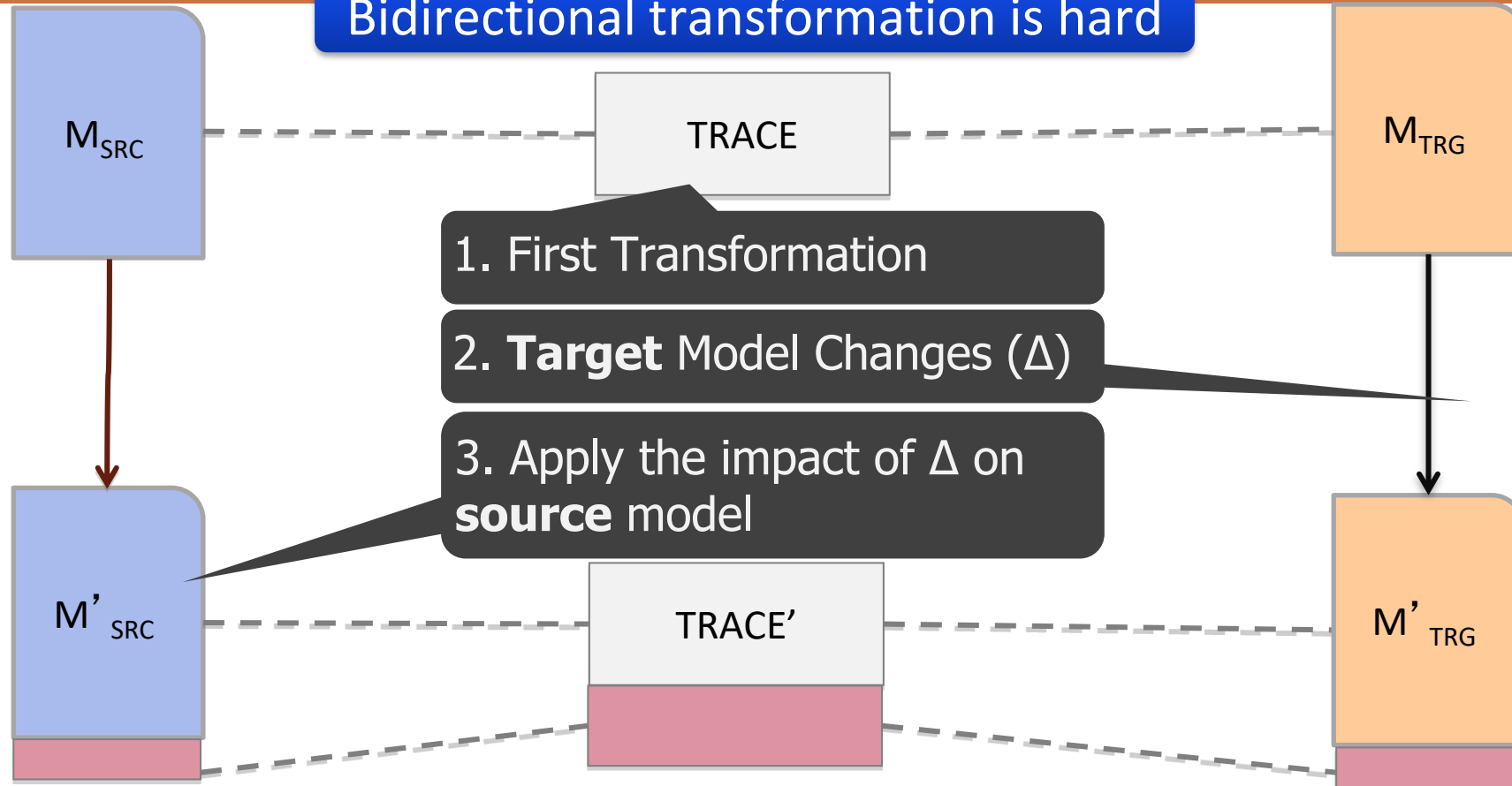
- ... by ignoring the unchanged parts of the source model
 - Incremental pattern matchers.

Incremental Forward Transformation Setup



Incremental backward transformation?

Bidirectional transformation is hard



Some related work:

A. Schürr, P. Stevens, N. Foster, T. Hettel,
Cicchetti&Pierantonio, Czarnecki&Diskin

Graph pattern matching, Graph transformation

Definitions

Graph pattern matching

Model transformations

Incremental transformations

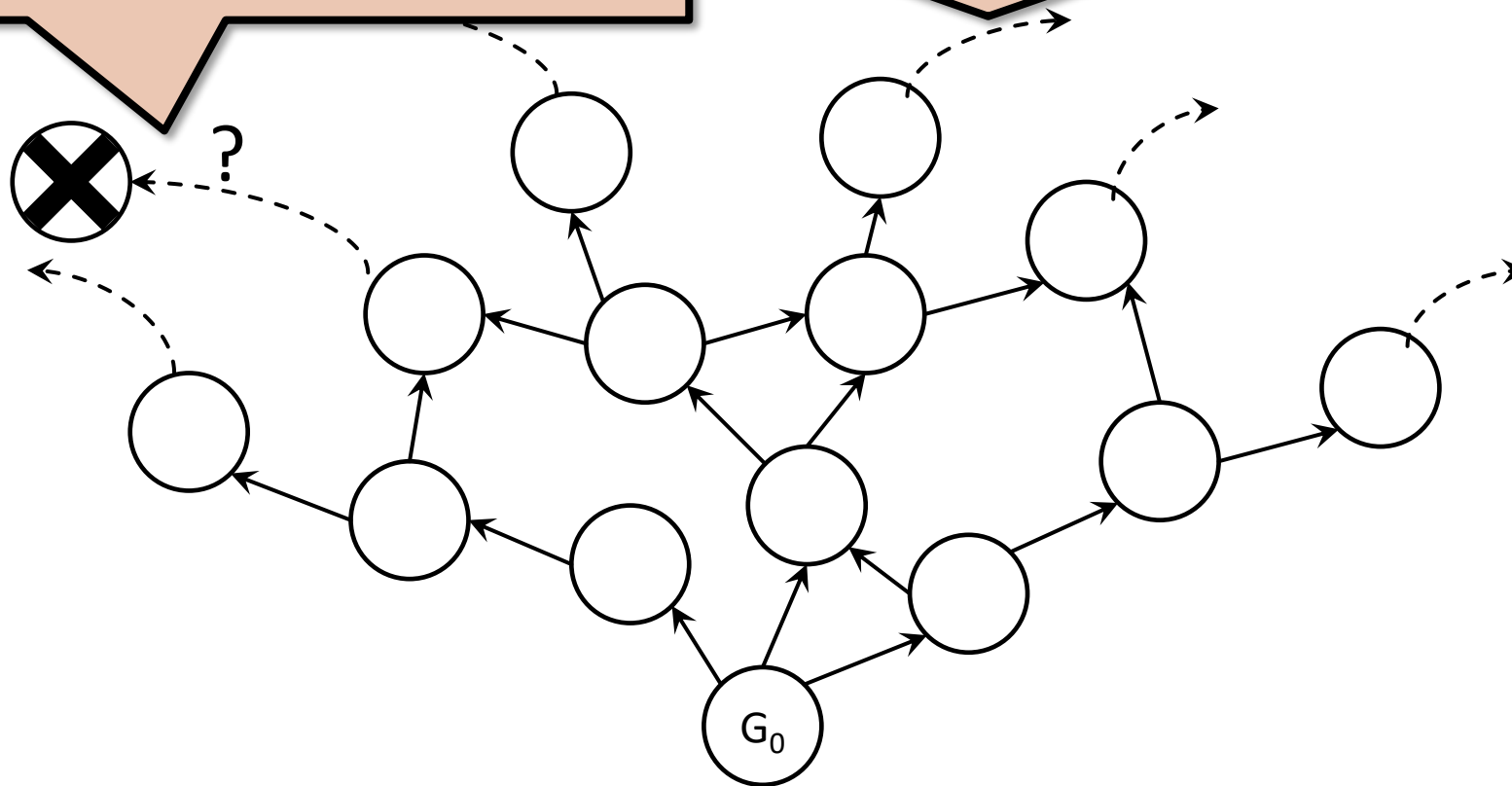
Design space exploration



Revisit: state space of GT system

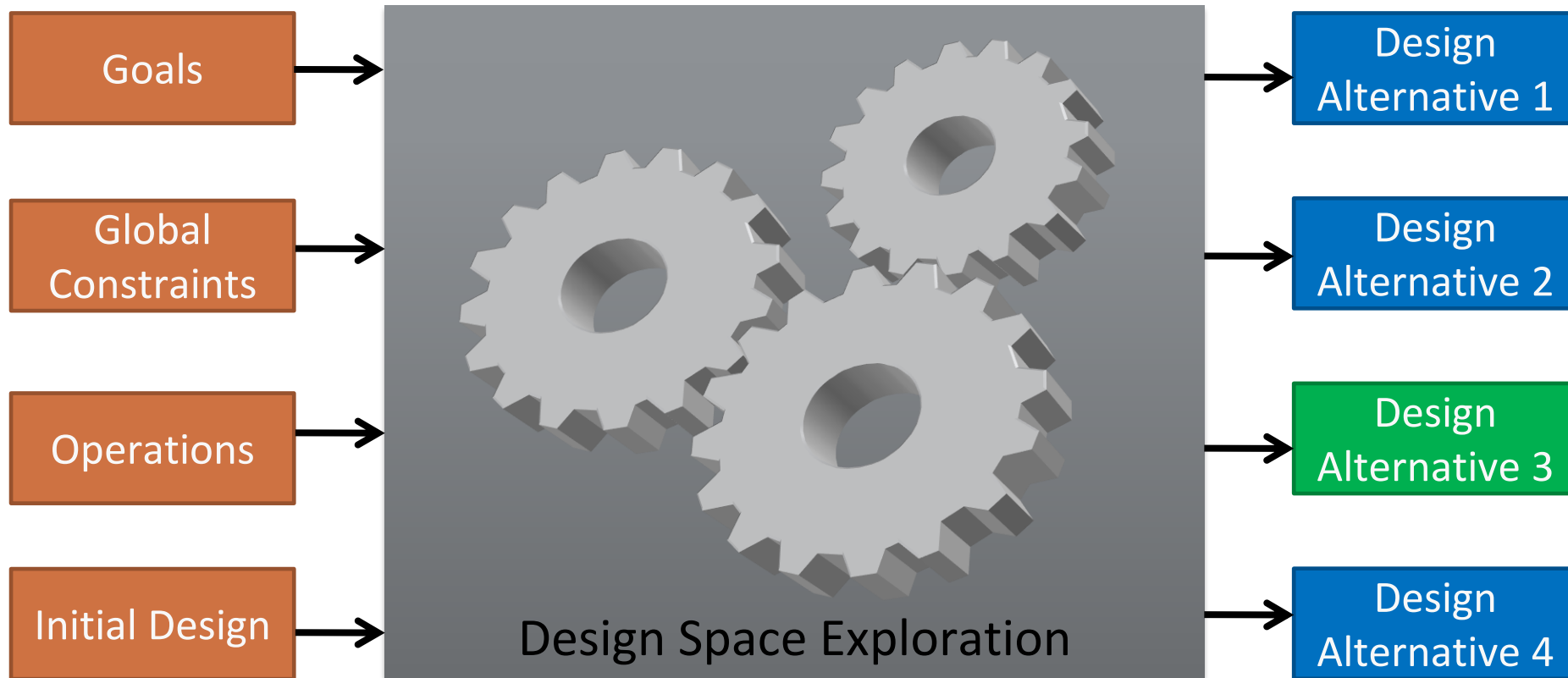
Solutions are
in the state space

Potentially infinite state space



Initial Graph + GT rules \rightarrow State Space

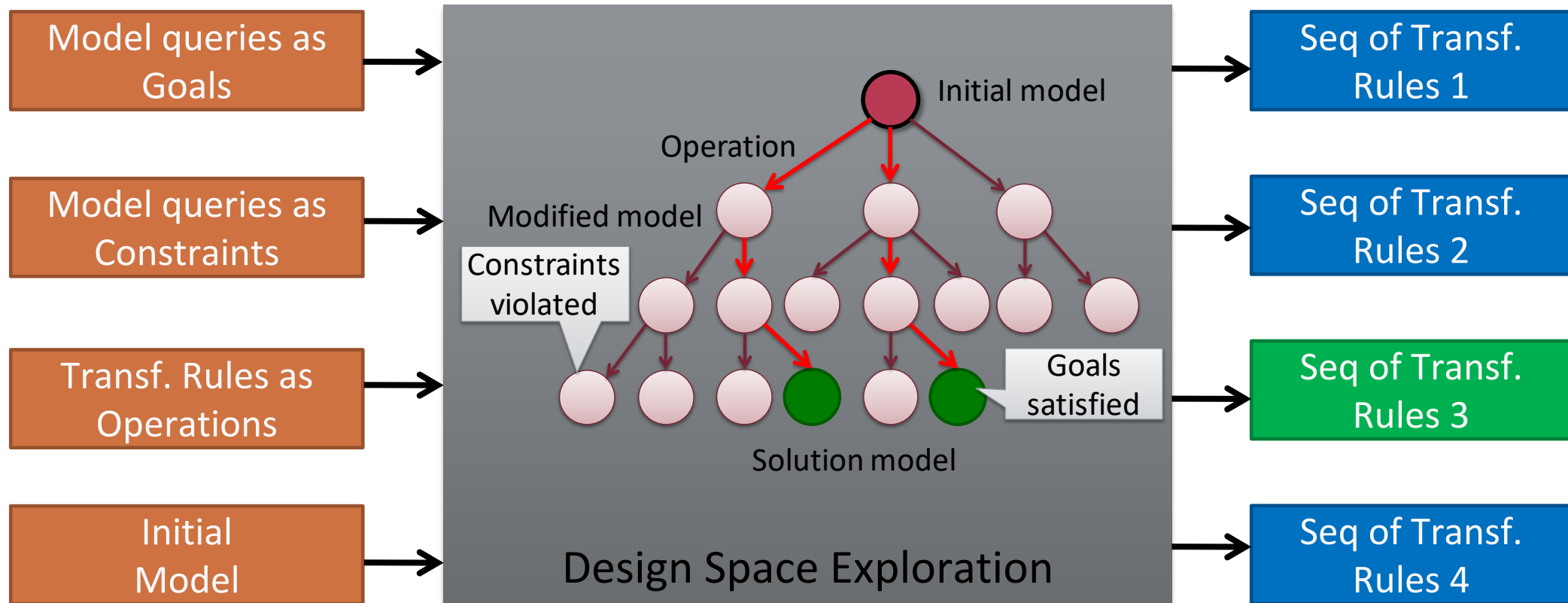
Design Space Exploration



Special state space exploration

- potentially infinite state space
- „dense” solution space

Model Driven Guided Design Space Exploration

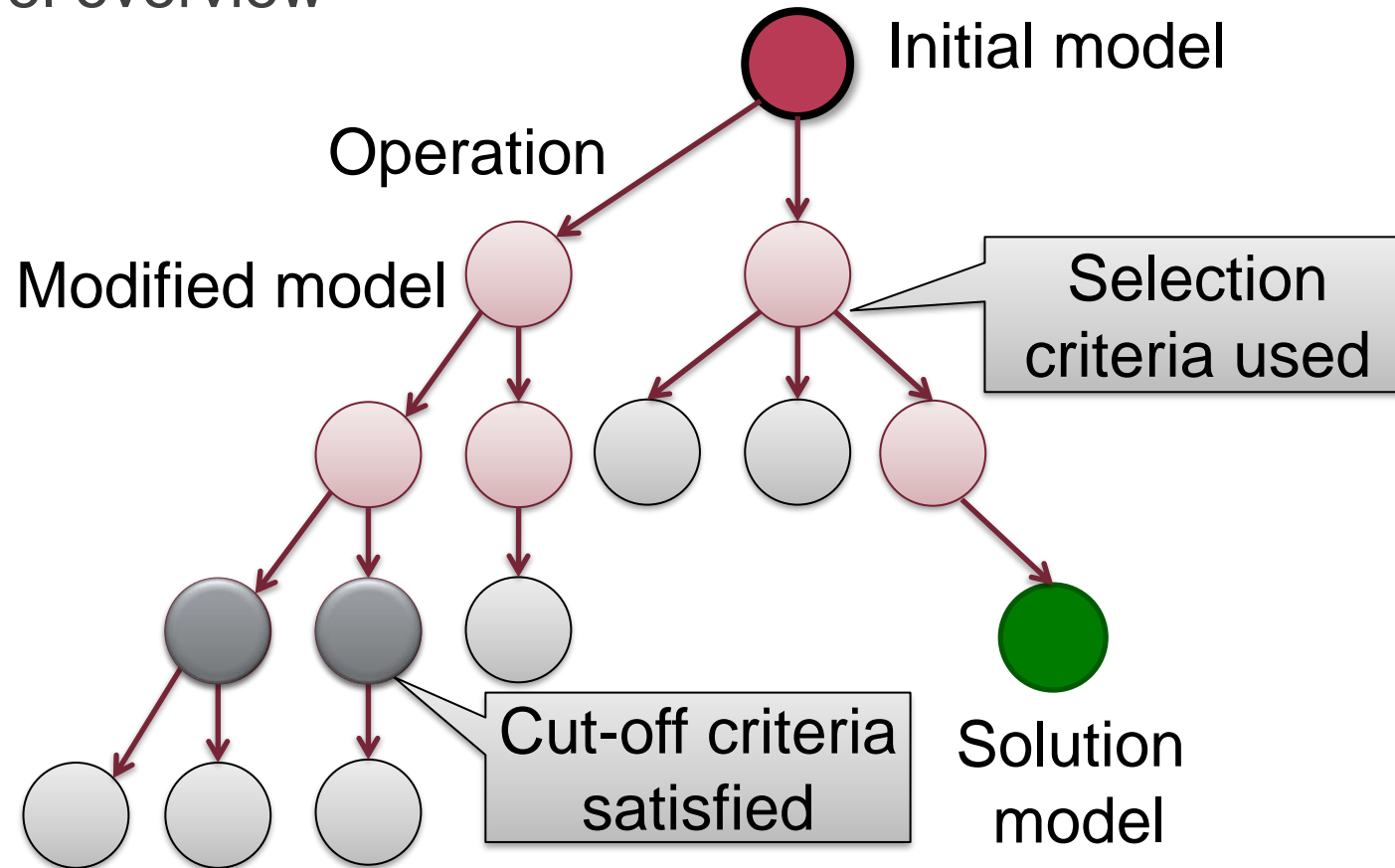


Guidance for exploration: Hints

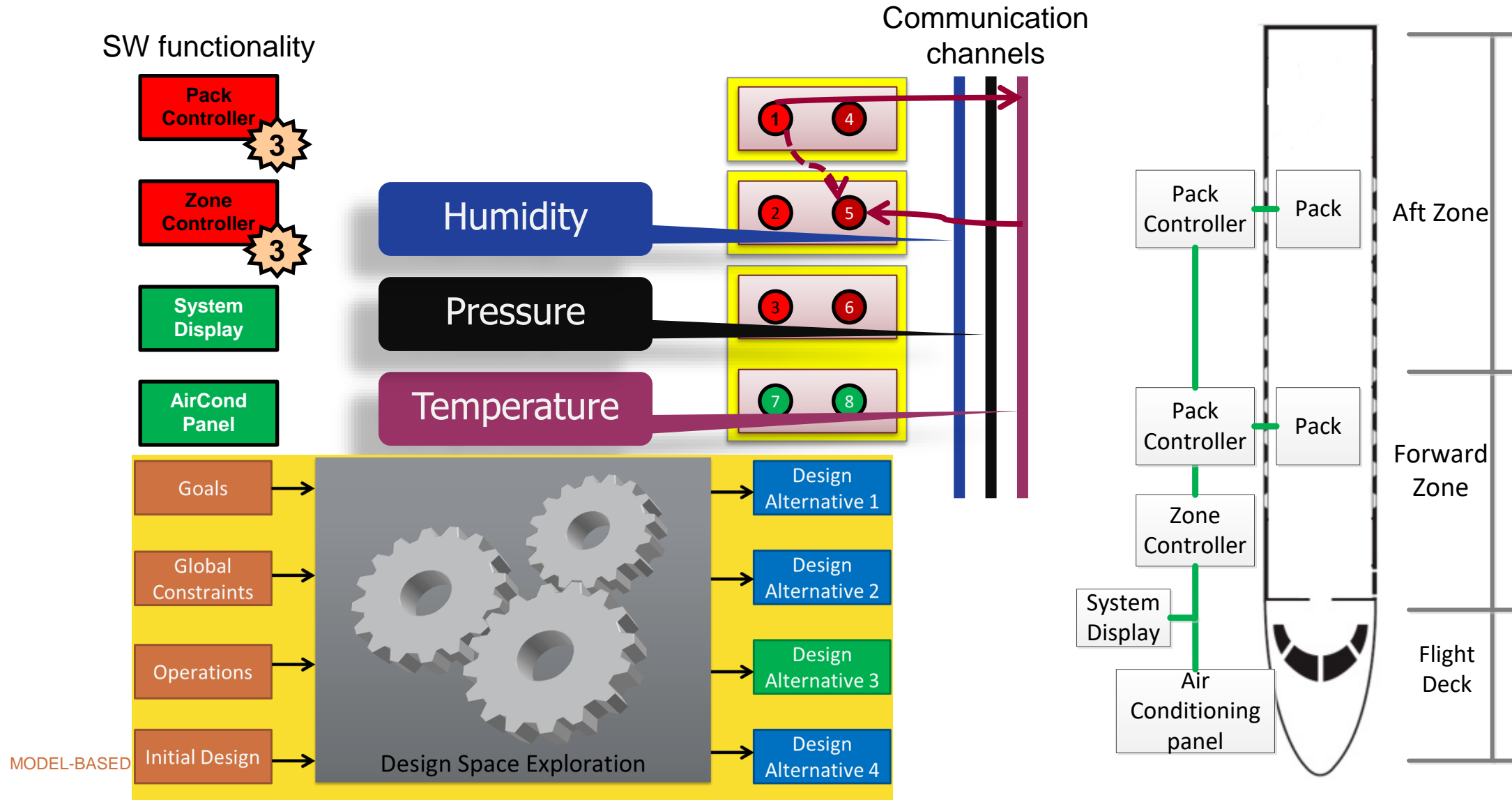
- designer / end user
- formal analysis

Guided Design Space Exploration

- High-level overview

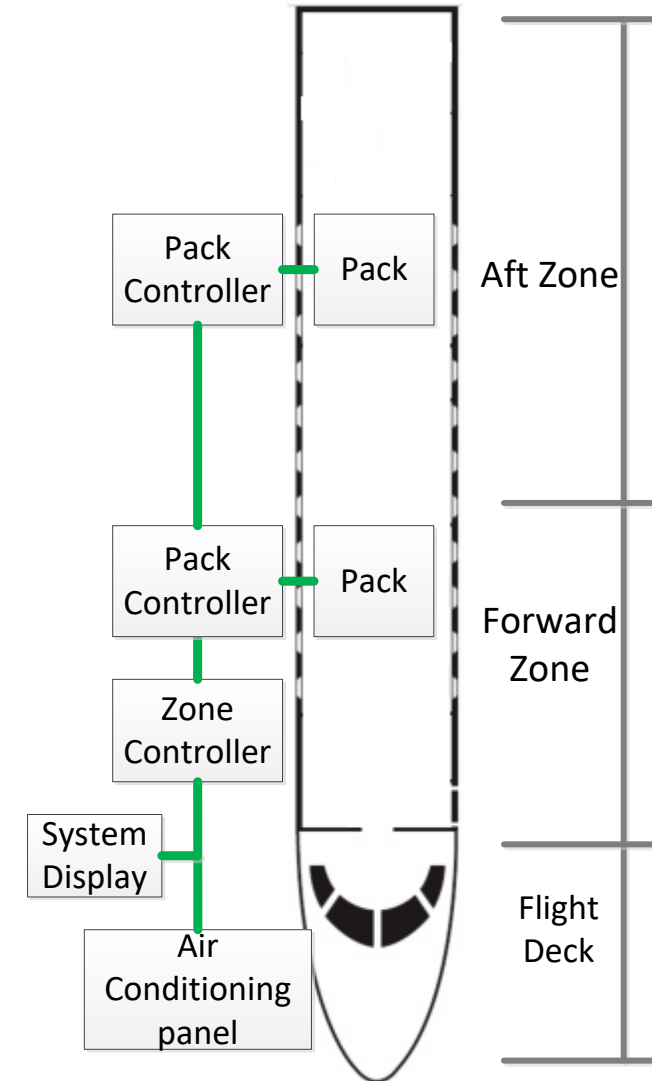
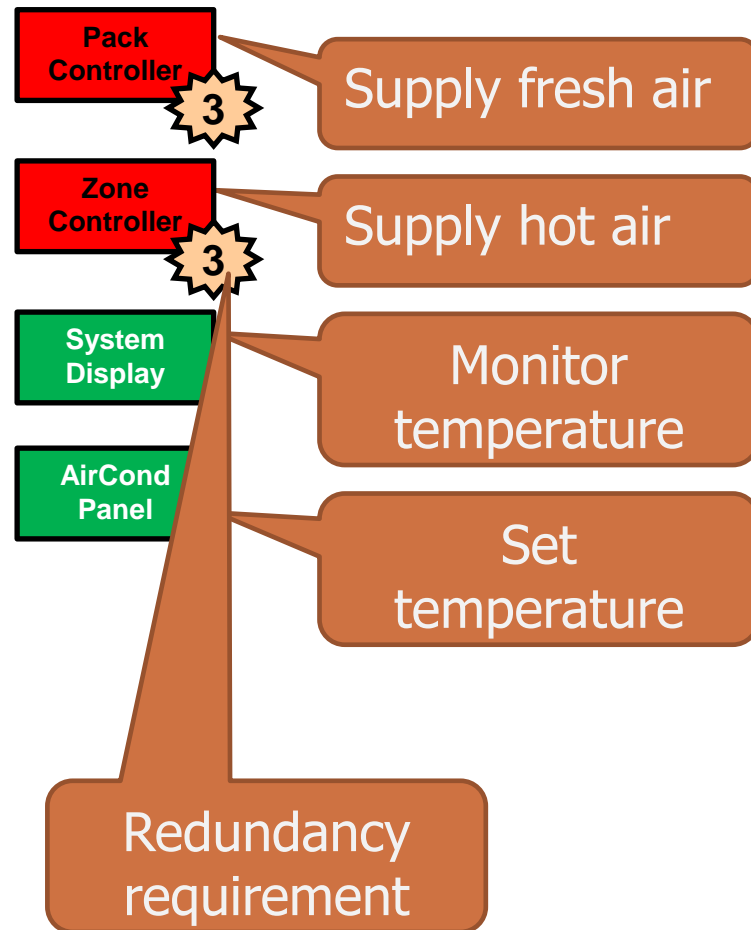


Design Space Exploration for IMA Configuration Design



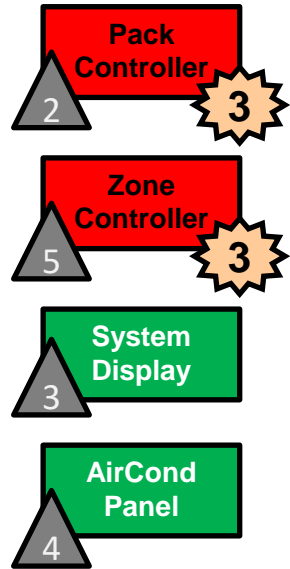
Designing ARINC653 configurations

SW functionality
(critical + non-critical)

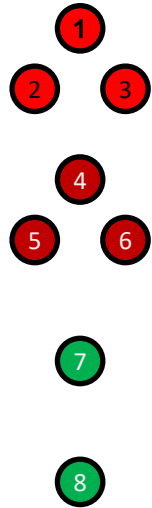


Job instances, Partitions, Modules

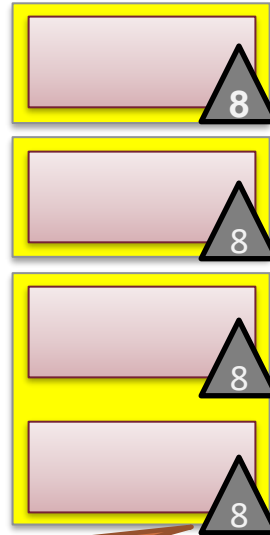
SW functionality
(critical + non-critical)



Job instances



Partitions

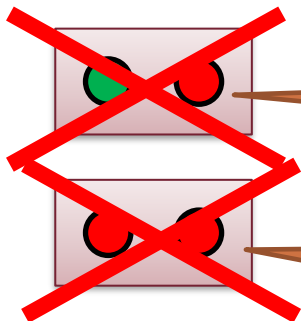


Modules

Additional constraints

- WCET,
- scheduling, etc.
- interfaces
- datatypes

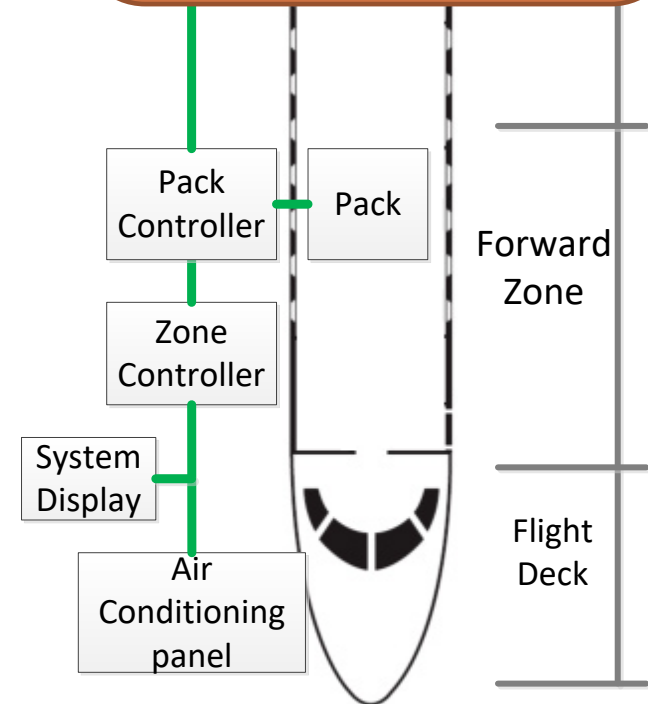
Constraints



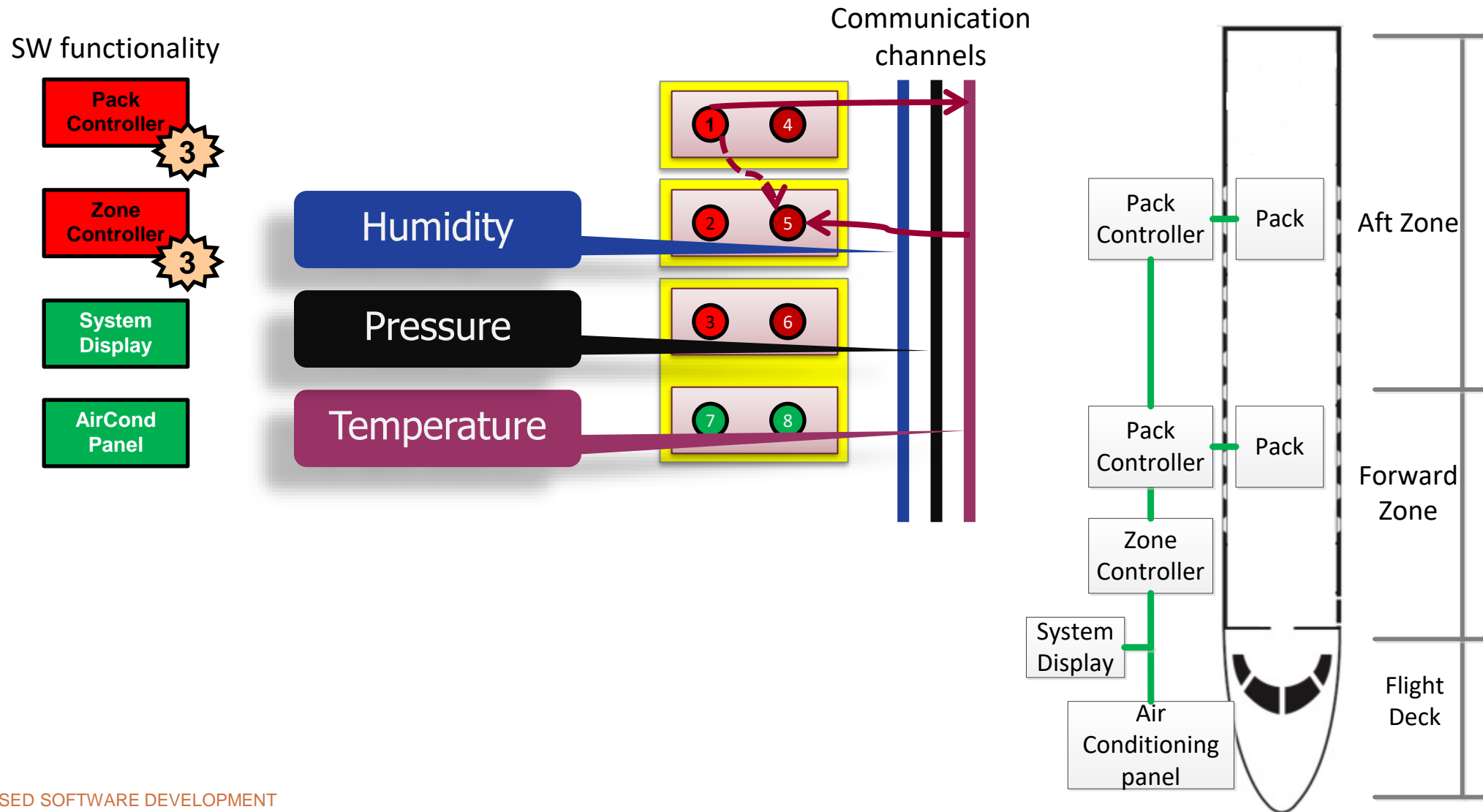
Memory needs
+ constraints

Do not mix critical
and non-crit. jobs

Do not mix instances
of the same critical job



Allocating communication channels



Graph pattern matching, Graph transformation

Definitions

Graph pattern matching

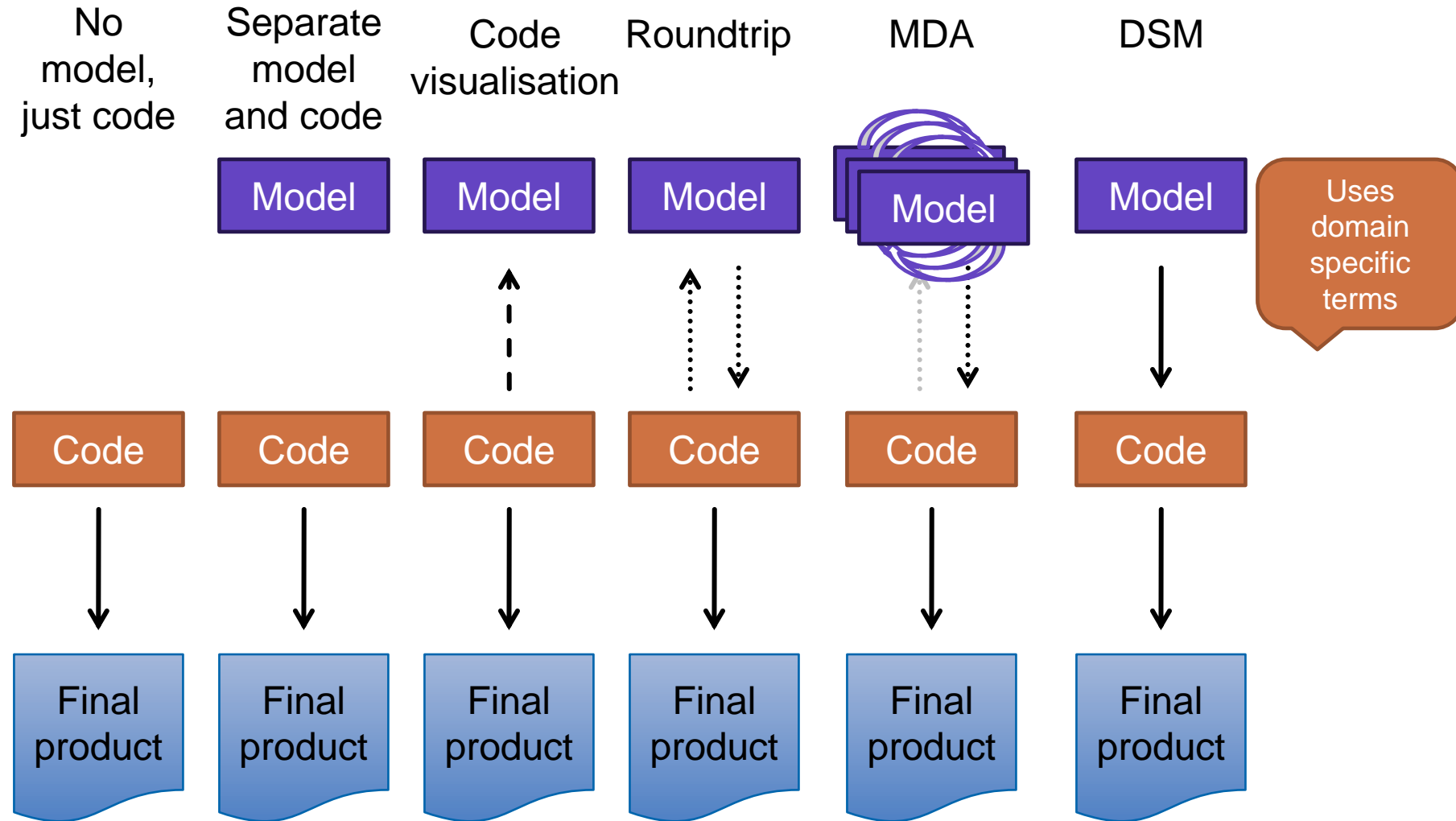
Model transformations

Incremental transformations

Design space exploration



How do we use models?





Thank you for your attention