

# Model-based software development

Lecture VIII.

Constraints, Syntax, Semantics

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# Syntax és Semantics

I. Constraints

II. Concrete and Abstract syntax

**III.** Editors

**IV. Semantics** 



# What are the components of a Domain-Specific Language?

- What do we need for a DSL?
  - > Language Structure (e.g. Metamodel)
  - > Constraints
  - > Views and Editors
  - > Meaning

- → Abstract Syntax
- → Concrete Syntax
- → Semantics

#### Constraints: Motivation

Challenge: How to describe complex structural patterns?

- Example: Object-oriented programming
  - > Each object is responsible for the consistency of itself
  - > "name is not empty"
  - > "age cannot be negative"
  - > "printedName on card needs to be the first part of name"
- Example: Domain-Specific Language
  - > Global constraints are more typical, responsibility is not on the object
  - > "validFrom and printedName needs to be unique"
  - > "Customer can have at most one valid card"

# Customer -name: string -age: int -isMale: bool -dateOfBirth: Date

-owner

\* -cards

CustomerCard

1

-valid: bool -validFrom: Date -goodThru: Date -printedName: string

#### About constraints

What is a constraint?

A restriction that filters unwanted models.

A transformation that maps models to a set of error messages.

- Sometimes can be enforced by metamodel, but:
  - > It can become difficult
  - > Or impssible

How to define constraints?

# Object Constraint Language (OCL)

- Allows precise UML/metamodel definitions
- OMG standard
- Features
  - > OCL constraints are declarative: they specify what is correct and not what should be done
  - > OCL constraints have no side effects: evaluating OCL expressions does not change the state of the system
  - > OCL constraints have formal syntax and semantics: their interpretation is unambiguous and can be automated
- Metamodel extension = more options
- Extending constraints = fewer options

#### Context

- Context: The model element for which the OCL term is defined
  - > class, interface, data type, component, operation, instance

```
context Customer inv: self.name <> ''
```

- Context Type: the type of model element for which the expression is evaluated
- If the context is a type, it is the same as the type of the context

```
context Customer inv:... \forall x: Customer(x) \Rightarrow \cdots
```

- Context instance: the specific model element for which the expression is evaluated
  - > Referenced by the keyword "self"

# Expression types by context

- Invariant
  - > "Every student has a Neptun code"
- Pre- and post-condition
  - > "It is dark before/after the sun rises/sets"
- Initial value
  - > "The car has run 0 km at the time of production"
- Derived value
  - > "The final grade is the average of the midterm and the final exam"
- Method body
  - > "The number of books in the library is the sum of the number of books on the shelves"

#### Invariants

- Constraint defined on a metamodel element
- A logical expression that must be true for all instances of a metamodel element at all times

```
context <metaelement>
inv [<contraint name>]: <logical expression>
```

# Invariant examples

context Customer inv:
 self.name <>''

context Customer inv:
 age >= 18

context CustomerCard inv
checkDates:
validFrom.isBefore(goodThru)

context Customer inv:
Customer.allInstances()->
forAll(c1, c2 | c1<>c2 implies
c1.name <> c2.name)

Date

+isBefore() : bool

#### Customer

-name : string

-age : int

-isMale : bool

-dateOfBirth : Date

1 -owner

\* -cards

#### CustomerCard

-valid : bool

-validFrom : Date -goodThru : Date

-printedName : string

# Pre- and post-conditions

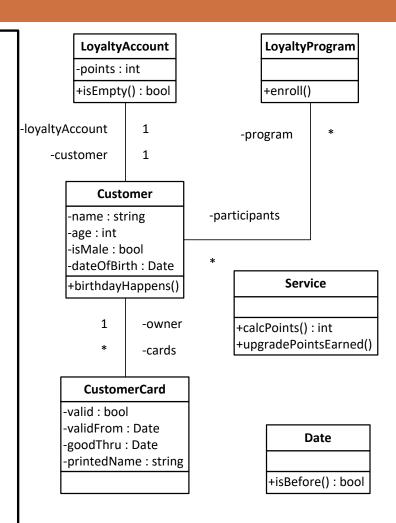
- Constraint defined for an operation
- Focuses on the effect of the operation regardless of algorithm or implementation
  - > Precondition: a condition true at the last moment before an operation is performed
  - > Postcondition: condition true at the first moment after an operation is performed

```
context <metaelement>::<operation> (<parameters>)
pre[<contraint name>]: <logical expression>

context <metaelement>::<operation> (<parameters>)
post[<contraint name>]: <logical expression>
```

# Pre- and post-condition examples

```
context LoyaltyAccount::
             isEmpty(): Boolean
post: result = (points = 0)
context Customer::birthdayHappens()
post: age = age@pre + 1
context Service::
upgradePointsEarned(amount: Integer)
post: calcPoints() = calcPoints@pre() +
amount
context LoyaltyProgram::
             enroll(c : Customer)
pre: c.name <> ' '
post: participants = participants@pre->
 including(c)
```



#### Initial value

- Constraint defined on an attribute or association
- A value taken by the attribute or association at the moment of context instance creation

```
context <metaelement>
init: <logical expression>
```

```
context CustomerCard::transactions : Set(Transaction)
init: Set{}

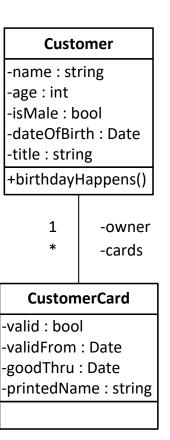
context CustomerCard::valid : Boolean
init: true
```

# CustomerCard -valid : bool -validFrom : Date -goodThru : Date -printedName : string 1 -card \* -transactions Transaction

#### Derived value

- Constraint defined on an attribute or association
- A derived element is not a value in itself, it is always defined by other elements

# +isBefore(): bool +addYears(): Date



# Query operations body

- Constraint defined for an operation
- There are operations that retrieve a specific value, with no other side effects
- A constraint can be used to specify what exactly they should return

#### Constraints and inheritance

- Constraints are also inherited
- The invariant is inherited by the derived class. The derived class may tighten the constraint, but it may not relax it.
  - > VIPCustomer needs to have a name as well. Name needs to start with "VIP-"

- The precondition can be relaxed in a redefined operation of the derived class. It cannot be tightened.
- The postcondition can be tightened in the redefined operation of the derived class. It cannot be relaxed.

# EXTRA: Types and Boole algebra in OCL

- All OCL expressions are typed
  - OclAny:
     The type that includes all others. E.g. x,
     y : OclAny
  - x = y x and y are the same object.
  - x <> y
    not (x = y).
  - x.oclType() The type of x.
  - x.isKindOf ( T ) True if T is a supertype (transitive) of the type of x.
  - T.allInstances() : CollectionAll the instances of type T.

- Boolean operators:
  - b and b2, b or b2, b xor b2, not b
     If any part of a Boolean expression fully determines the result, then it does not matter
     if some other parts of that expression have unknown or undefined results.
  - b implies b2
     True if b is false or if b is true and b2 is true.
  - if b then e1 else e2 endif
    If b is true the result is the value of e1;
    otherwise, the result is the value of e2.

#### EXTRA: Overview of Collection Valued Terms

- Size / aggregation:
  - c->size(): Integer Number of elements in the collection; for a bag or sequence, duplicates are counted as separate items.
  - c->sum(): Integer
     Sum of elements in the collection.
     Elements must be numbers
  - c->count(e): Integer
     The number of times that e is in c.
  - c->isEmpty(): Boolean
    Same as c->size() = 0.
  - c->notEmpty(): Boolean
    Same as not c->isEmpty().

- Equality
  - c = c2 : Boolean
- Collection membership
  - c->includes(e): Boolean; c->exists ( x | x = e ).
  - c->excludes(e): Boolean; not c->includes( e ).
  - c->includesAll(c2): Boolean; c includes all the elements in c2.
  - c->including(e): Collection
     The collection that includes all of c as well as e.
  - c->excluding(e): Collection
     The collection that includes all of c except e.

#### EXTRA: Overview of Collection Valued Terms

- Existential quantifier:
  - c->exists(x | P ): Boolean; there is at least one element in c, named x, for which predicate P is true.
  - Equivalent notation is: c->exists(P), c->exists(x:Type | P(x))
- Universal quantifier:
  - c->forAll(x | P): Boolean; for every element in c, named x, predicate P is true.
  - Equivalent notation is: c->forAll(P) c->forAll(x:Type | P)

#### Selection:

- c->select(x | P): Collection The collection of elements in c for which P is true.
- Equivalent is: c->select(P)
- Filtering:
  - c->reject(x | P): Collection
    c->select(x | not P).
  - Equivalent is: c->reject(P)
- Collection:
  - c->collect(x | E): Bag The bag obtained by applying E to each element of c, named x.
  - c.attribute : CollectionThe collection(of type of c)

# Syntax és Semantics

I. Constraints

**II.** Concrete and Abstract syntax

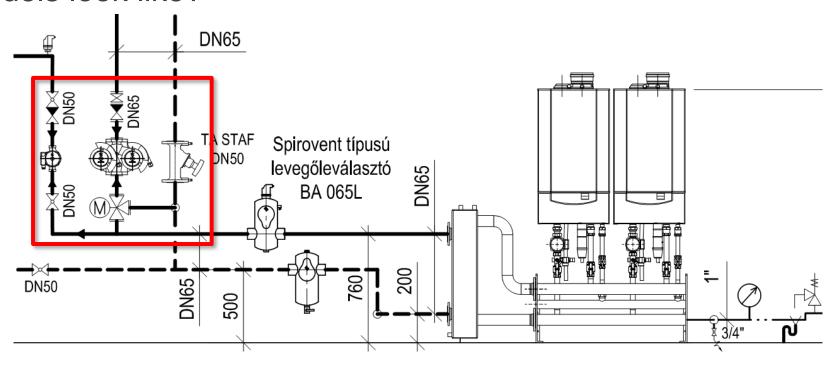
**III.** Editors

**IV. Semantics** 



# Concrete syntax

How models look like?



Honeywell keverőcsap DN50 Kvs 40 Spirovent típusú iszapleválasztó BE 065L

Remeha Quinta kaszkád rendszer hidrauliku

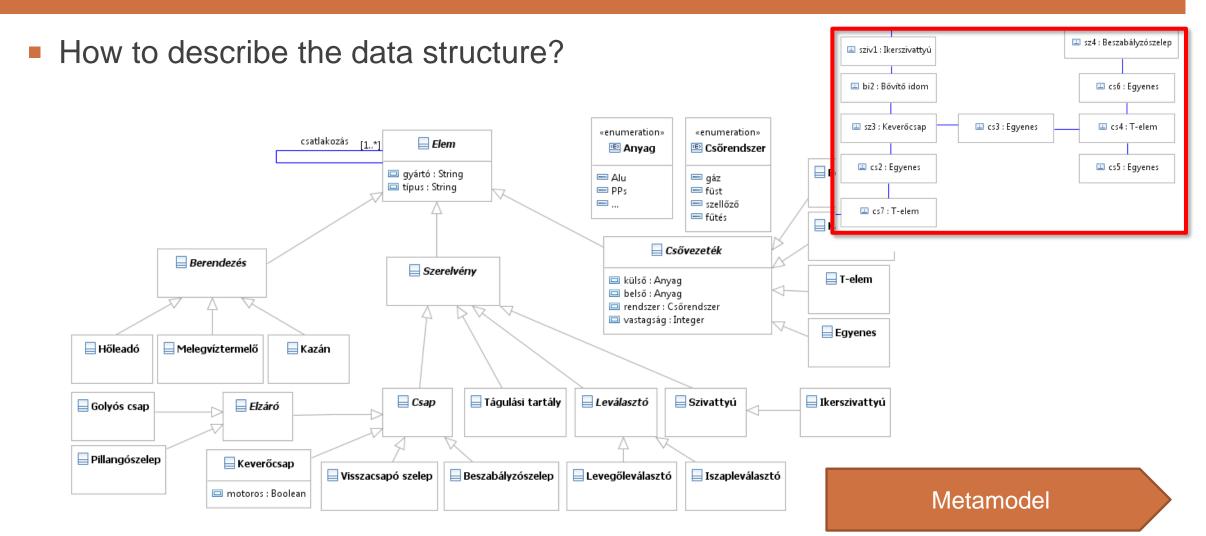
Concrete syntax

# Abstract syntax: model

How to represent models in a computer?



# Abstract syntax: description of the model



#### **Definitions**

Separate models to representation-independent and representation-specific parts.

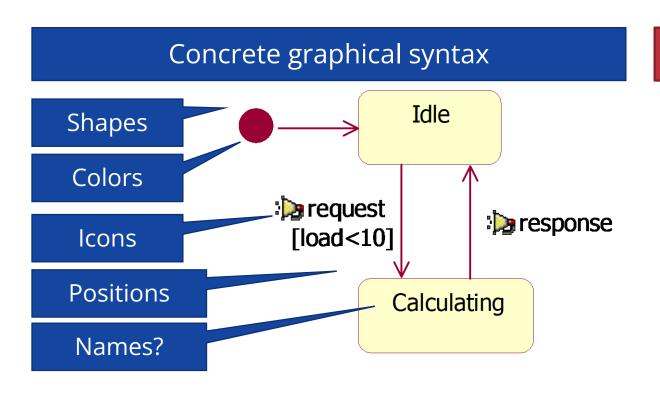
- Definition [Abstract syntax]: is the (abstract) data structure of a model, which excludes representation-specific details.
- Definition [Concrete syntax]: is the complete representation of a model (which includes representation-specific details).

 Abstract and concrete syntax may also denote the representation technique of a modeling language.

"In the concrete syntax of the Yakindu modeling language states are represented by rounded rectangles.."

# Concrete syntax

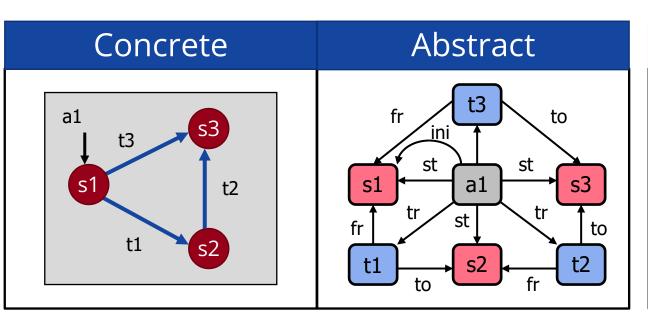
What is part of the concrete syntax, but not part of the abstract?

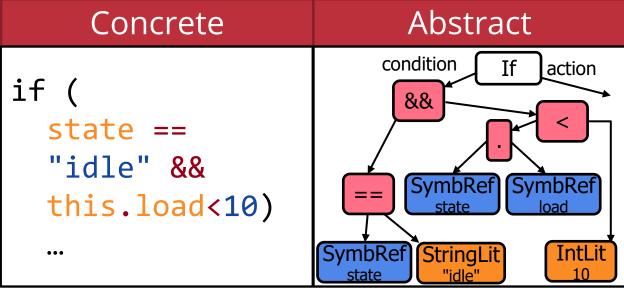


#### Concrete textual syntax

# Abstract syntax

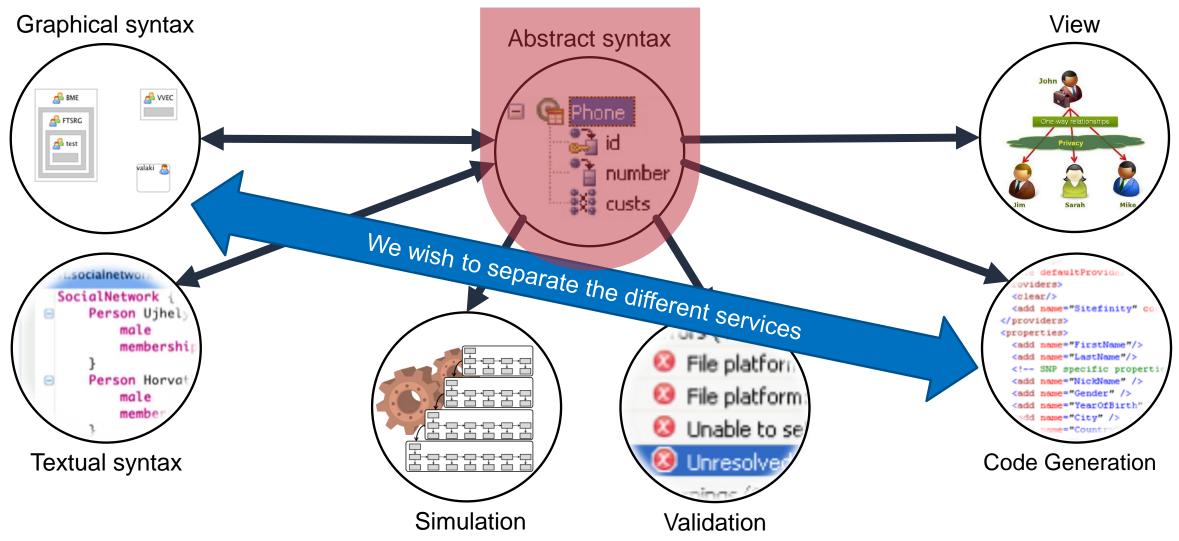
How to capture the abstract syntax of a model?



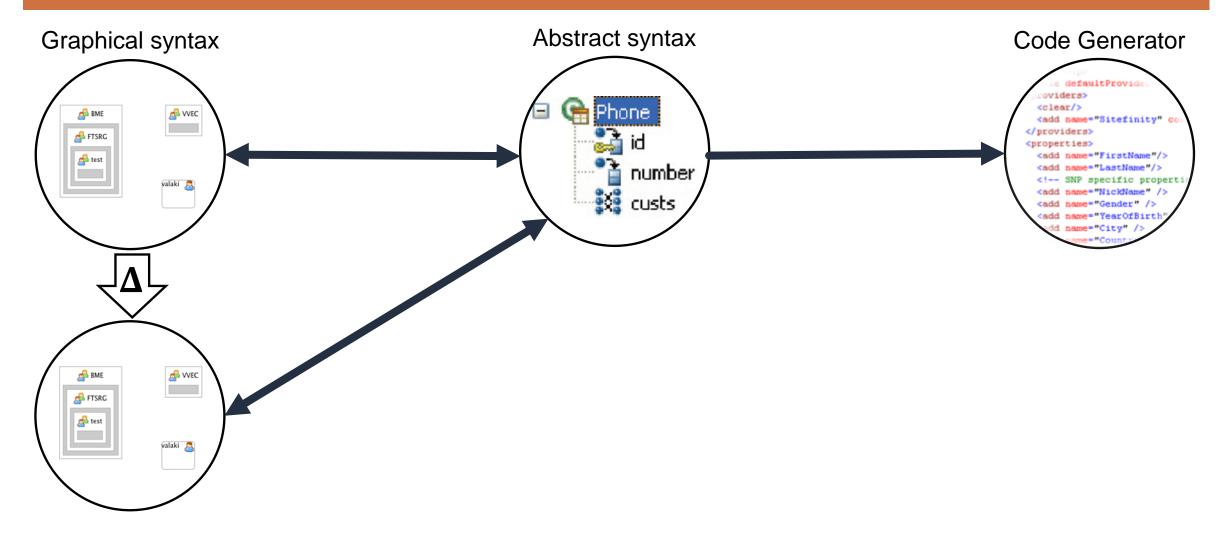


Abstract syntax: typically a graph-based structure.

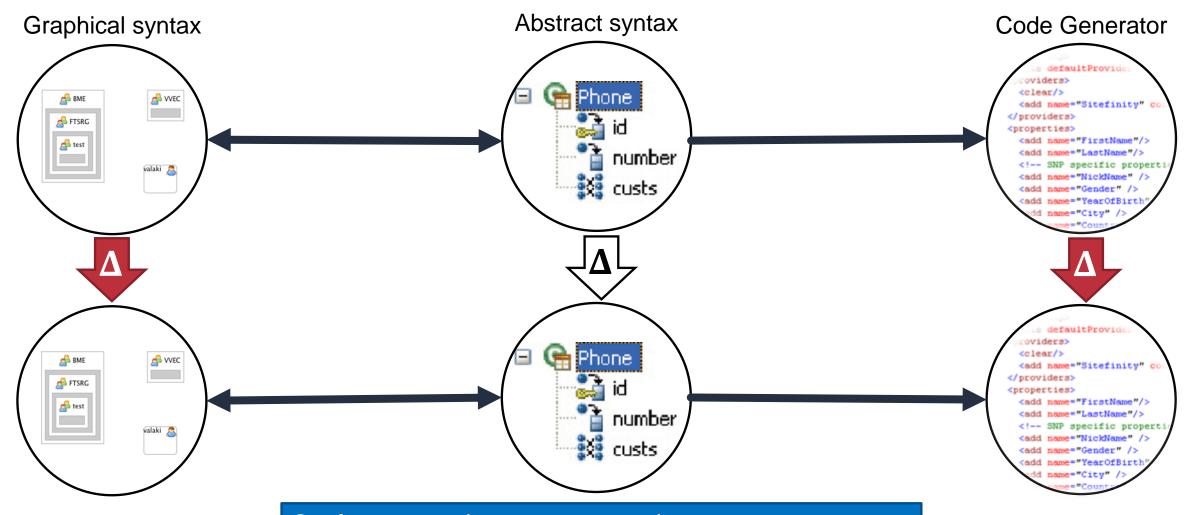
# How to separate concrete and abstract syntax?



# Change the concrete syntax, but not the abstract syntax



# Change the abstract syntax



**Goal:** separate the concrete ans abstract syntax without the modeling services interfering

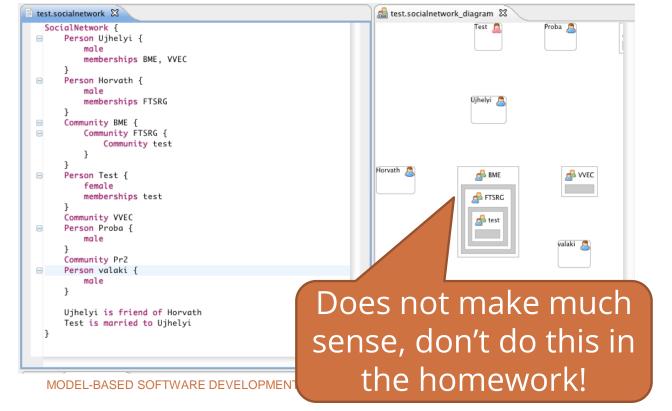
# Multiplicity of Notations

- 1 abstract syntax → many textual and visual notations
  - > Human-readable-writable textual or visual syntax
  - > Textual syntax for exchange or storage (typically XML)
  - > In case of UML, each diagram is only a partial view
- 1 abstract model → many concrete forms in 1 syntax!
  - > Whitespace, diagram layout
  - > Comments
  - > Syntactic sugar
- 1 semantic interpretation → many abstract models
  - > e.g. UML2 Attribute vs. one-way Association

# Textual + Graphical

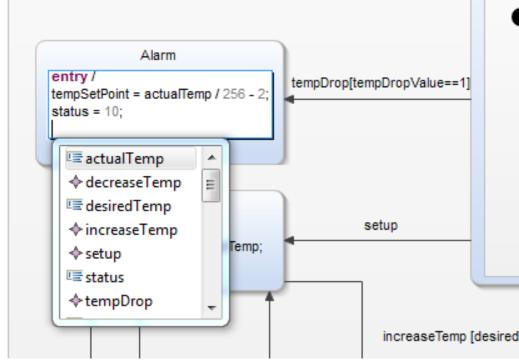
#### Same model, two syntaxes

- Text editor + graphical view □
- Xtext Generic Viewer



#### Different aspects of model

- Diagram with text fields
- Embedded Xtext support



# Syntax és Semantics

I. Constraints

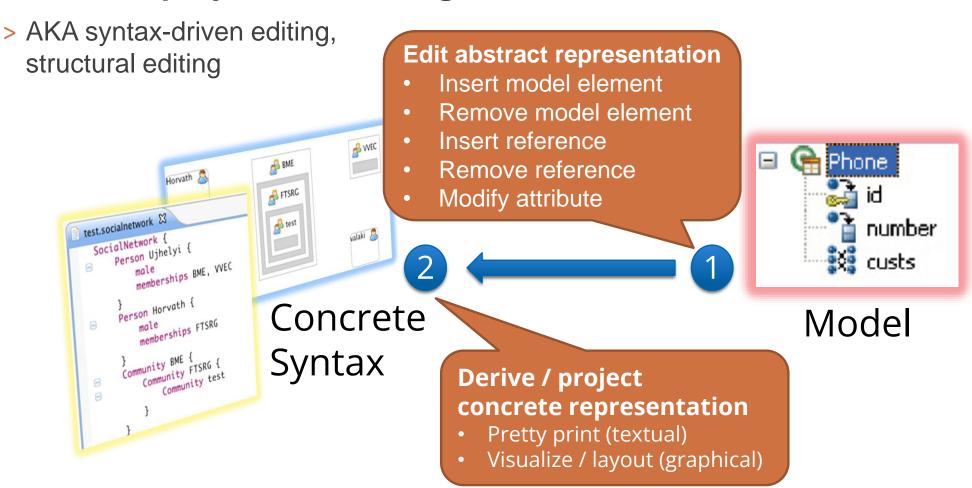
II. Concrete and Abstract syntax

**III.** Editors

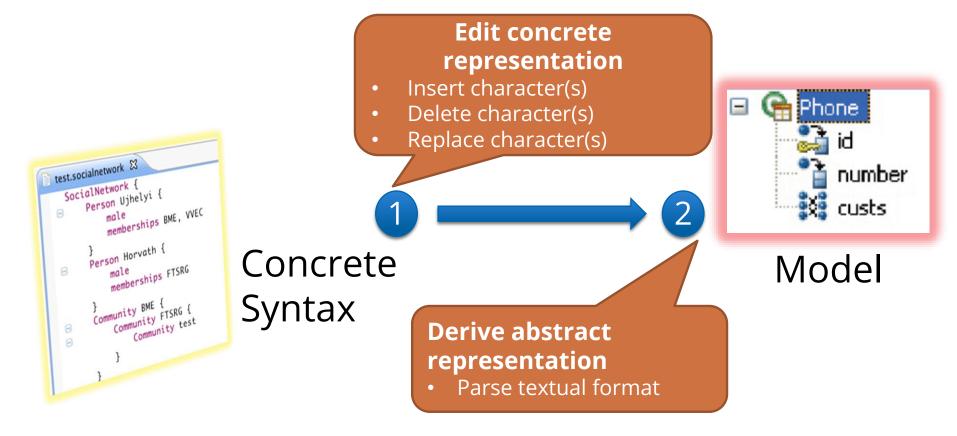
**IV. Semantics** 



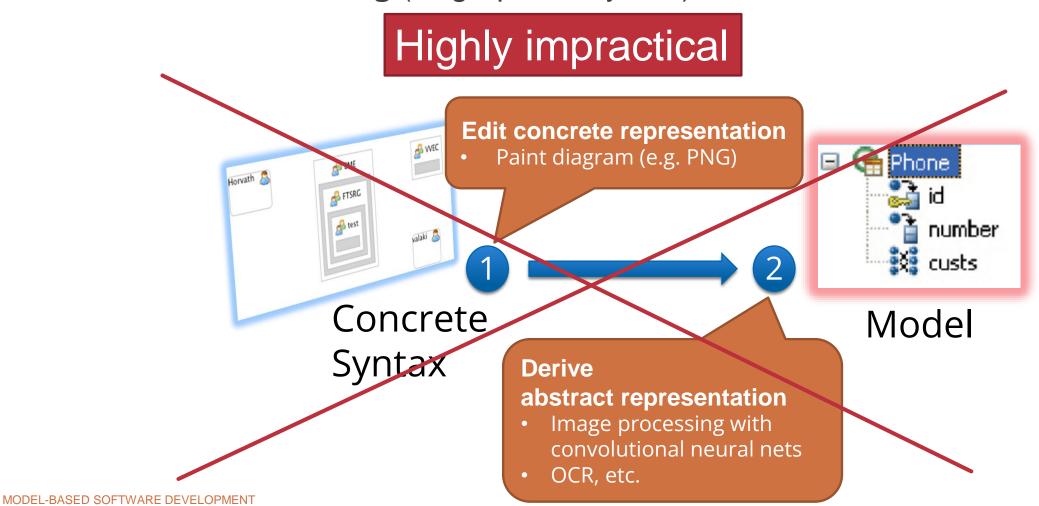
Workflow 1: projectional editing



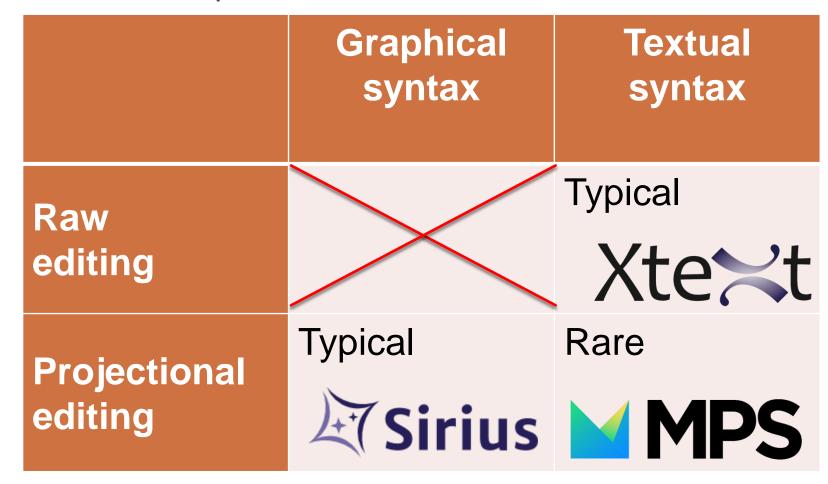
- Workflow 2: raw editing (w. textual syntax)
  - > AKA source editing



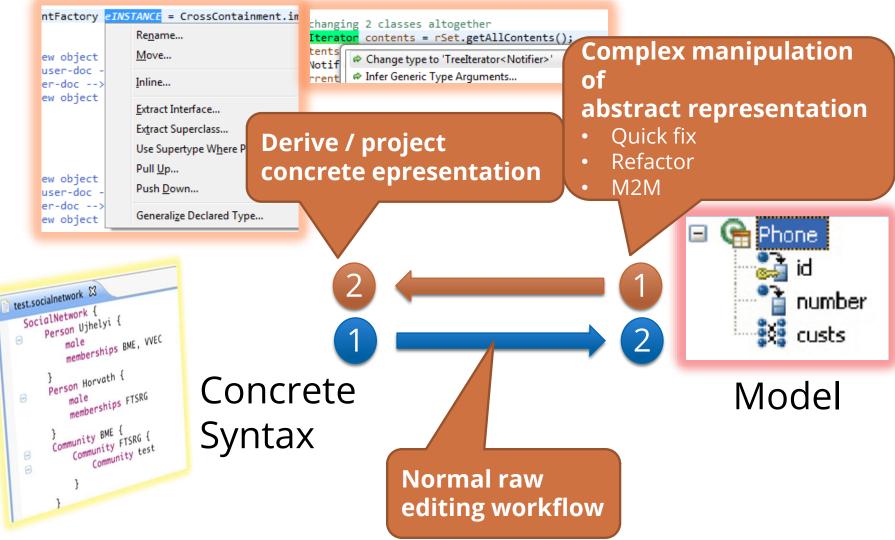
Workflow 2: raw editing (w. graphical syntax)



"Feature matrix" + examples



#### Mixed workflow



## Transactions in projectional editing

MODEL-BASED SOFTWARE DEVELOPMEN

Complex manipulation sequence as single action

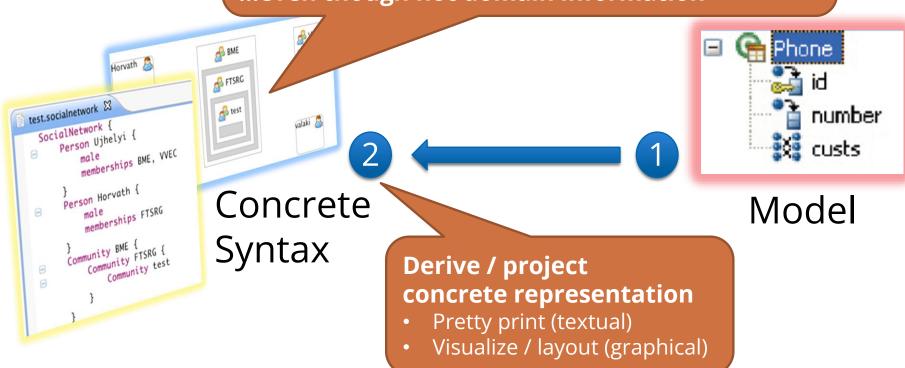
> "Extract subprocess", "Drag&drop attribute" etc. Transaction initialized ransaction Check for concurrent read or Begin Transaction write transactions **START** How to ensure? Revertibility **Declarative** Rollback commands Manipulation step 1 Record change Manual undo notifications Manipulation step 2... DO **Optional: check validity** Reject & roll back if violated Precommit Transaction finalized Postcommit **FINISH** *Issue change notifications* (if not earlier) Refresh projections

#### Superfluous notational parameters

Workflow 1: projectional editing



- Whitespace and comments, etc. (textual)
- · Layout, edge routing, size, shape, etc. (graphical)
- ...even though not domain information



### Deriving notational parameters

- Notational parameters can be...
  - > ... "baked into" projection code
    - e.g. all lines are black, all fonts are 10pt (graphical)
    - e.g. apply this code formatting template (textual)
  - ...derived from domain information
    - e.g. shape determined by type, color by visibility

#### **Problem 1**:

Editable parameters cannot be a function of the domain model, must be stored

#### **Problem 2:**

Providing sane values is difficult for some parameters e.g. position in diagram

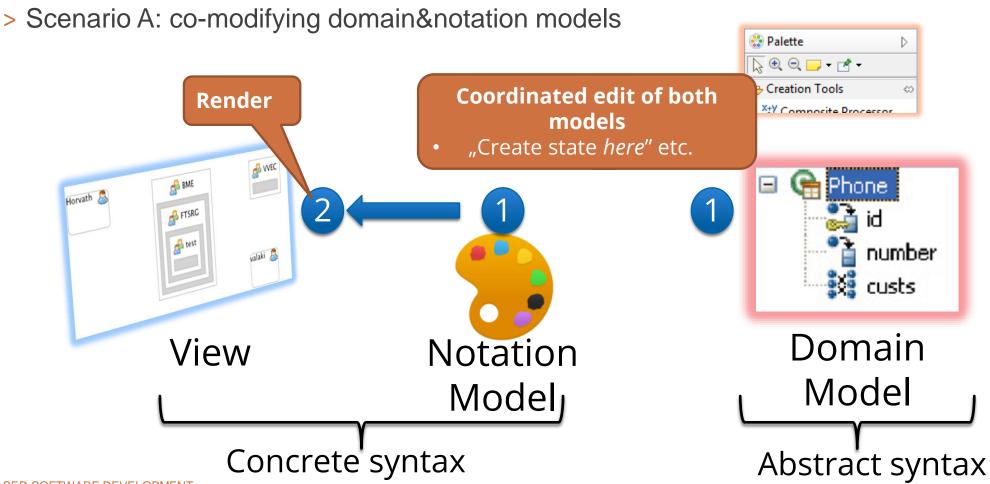
> ...stored in the model

#### Notation/view models

- Decompose model:
  - > Domain / Semantic model (abstract syntax)
  - > Notation model (view model): presentation state
    - may be editable by user
    - but still needs derivable defaults → see layouting
- Generic implementation in GMF and Graphiti
  - > Based on EMF, in fact
- Often stored in external files
  - > Separation of concerns
  - > E.g. code generator not interested in view information

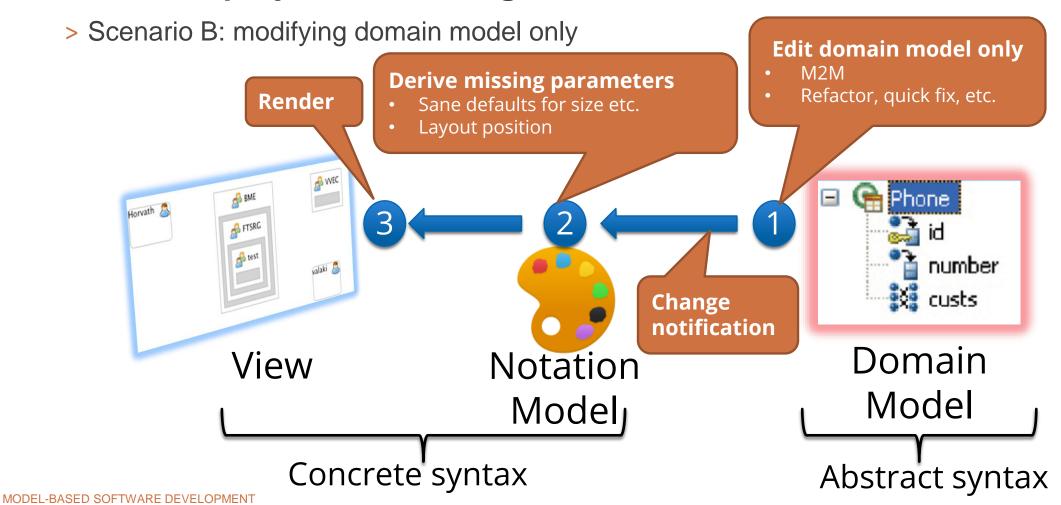
## Editing workflow with notation models

Workflow 1: projectional editing



## Editing workflow with notation models

Workflow 1: projectional editing



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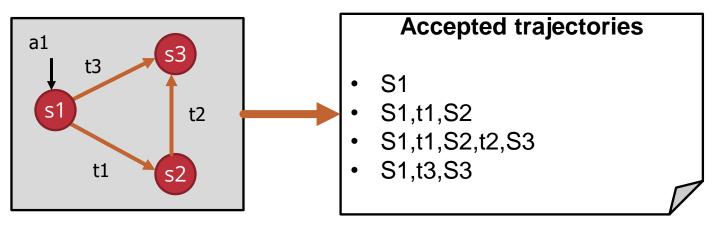
#### Semantics

- We talked a lot about syntax.
- Semantics: the meaning of concepts in a language
  - > Static: what does a snapshot of a model mean?
  - > Dynamic: how does the model change/evolve/behave?
- The semantics of a modeling language maps models to a real-world or a mathematical semantic domain.

## Types of semantics

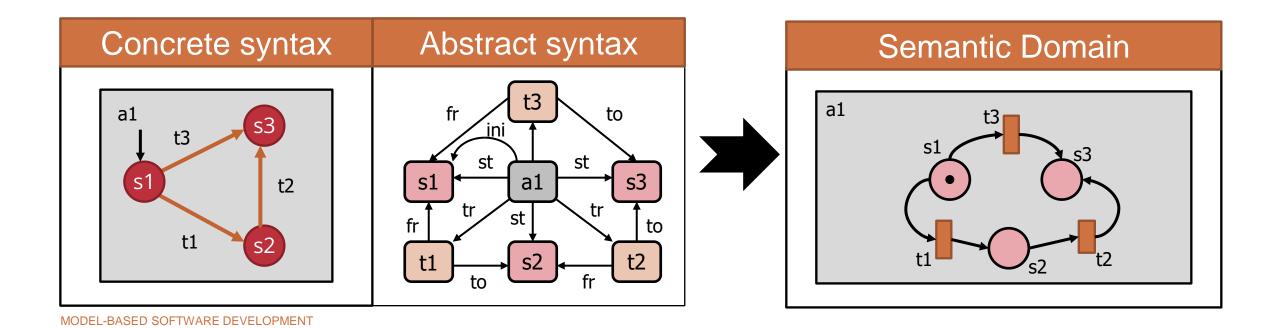
- Static Semantics
  - > Interpretation of metamodel elements: meaning of concepts in the abstract syntax

- Formal: mathematical statements about the interpretation
- Semantics is captured by mathematical axioms



## Example: denotational semantics

- Denotational (Translational)
  - > translating concepts in one language to another language (called **semantic domain**)
  - > "compiled"
  - > E.g. explaining state machines as Petri-net

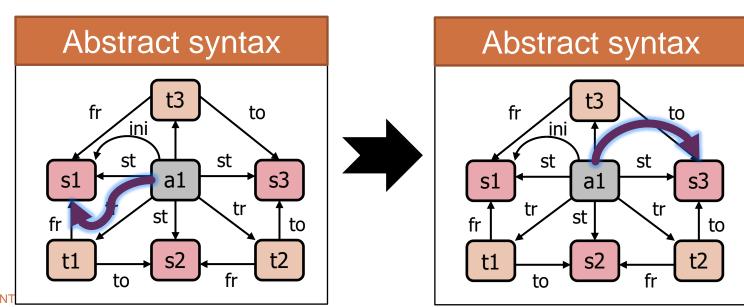


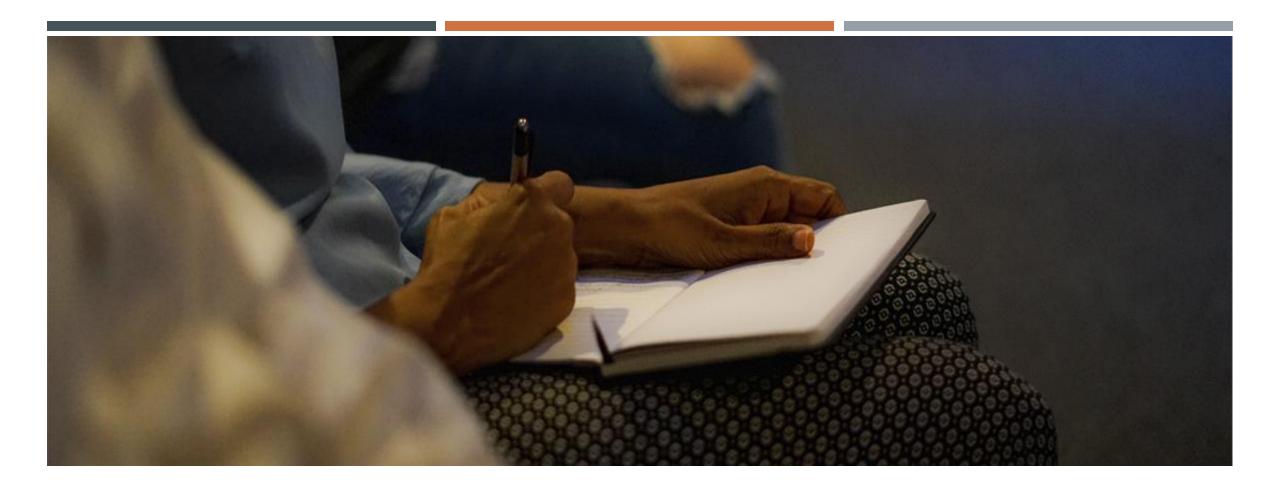
## Operational Semantics

#### Operational

- > Modeling the operational behavior of language concepts
- > "interpreted"
- > e.g. defining how the finite automaton may change state at run-time
- > Sometimes dynamic features are introduced only for formalizing dynamic sematics

Dynamic feature : current





# Thank you for your attention