

Encryption algorithm based on Rubik's cube principle

-An implementation on AD BF533 DSP-

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1. Introduction

The following paper presents an image encryption algorithm based on Rubik's cube principle as part of the project work for the Information Processing Technologies university course. The work is based entirely on the research published [1]. It is not my objective to present a new algorithm, but rather to implement the mathematical rationale behind it on AD BF533 Digital Signal Processor (DSP). The project firstly presents the logic behind the algorithm and its implementation in the C programming language, using Visual DSP as IDE. Results of both encrypted and decrypted images of different sizes show the diffusion and confusion elements of the algorithm. In addition, an Assembly programming language implementation has been done, in order to optimize the speed and memory characteristics of the algorithm. Last but not least, a speed comparison between the C and Assembly is displayed. The active work done by the student author of the project is creating a version of the Rubik's cube image encryption algorithm optimized for the BF553 DSP.

The paper is organized in the following manner: Section 2 describes the logical and mathematical operations of the algorithm, Section 3 presents the C implementation and results in Visual DSP, Section 4 shows the optimization done in Assembly programming language with the aim to increase performance.

2. Rubik's Cube image encryption algorithm

2.1. Encryption

1) As it has been stated above, the main source of inspiration for this project is [1]. The implemented code is entirely based on the following sequence of operations. The input is represented by an α -bit gray scale image with M rows and N columns called I_0 . The solution has three encryption keys:

- K_R and K_C randomly generated vectors of length M and N. Both vectors are populated with values ranging from $[0, 1, 2, \dots, 2^\alpha - 1]$, which must have non-constant values.
- The maximum number of iterations which will be done $ITER_{max}$.

After determining the encryption keys, which will be done in a setup phase, the encryption process can begin by initializing a counter $ITER = 0$.

2) The counter will be incremented $ITER = ITER + 1$.

3) For each row, an $a[i]$ vector will be calculated, representing the sum of all elements on the row i : $\sum_{j=0}^{N-1} I_0[i, j]$, $i = 0, 1, 2, \dots, M - 1$.

Compute Ma vector, where $Ma[i] = a[i] \% 2$.

Each row of the input image I_0 will be circular shifted by $K_R[i]$ positions according to the following rule:

If $Ma[i] = 0$ then right circular shift.

If $Ma[i] = 1$ then left circular shift.

4) For each column, a $b[i]$ vector will be calculated, representing the sum of all elements on the column j : $\sum_{i=0}^{M-1} I_0[i, j]$, $j = 0, 1, 2, \dots, N - 1$.

Compute Mb vector, where $Mb[i] = b[i] \% 2$.

Each column of the input image I_0 will be circular shifted by $K_C[i]$ positions according to the following rule:

If $Mb[i] = 0$ then up circular shift.

If $Mb[i] = 1$ then down circular shift.

Steps 4) and 5) will create a scrambled image, representing the diffusion element of the algorithm.

5) Using the K_C vector, the Bit wise operation XOR will be applied to each row of the scrambled image, using the following rule:

$$I_1(2i - 1, j) = I_{SCR}(2i - 1, j) \oplus K_C[j];$$

$I_1(2i, j) = I_{SCR}(2i, j) \oplus \text{rot180}K_C[j]$, where rot180 represents the 180 degree flip of vector K_C .

6) Using the K_R vector, the Bit wise operation XOR will be applied to each column of the scrambled image, using the following rule:

$$I_{ENC}(i, 2j - 1) = I_1(2i - 1, j) \oplus K_R[j];$$

$I_{ENC}(i, j) = I_1(2i, j) \oplus \text{rot180}K_R[j]$, where rot180 represents the 180 degree flip of vector K_R .

Steps 5) and 6) change the value of each pixel of the image, representing the confusion element of the algorithm.

7) If $ITER = ITER_{max}$, the encryption process is complete. Otherwise, the algorithm jumps back to step 2).

2.2. Decryption

The decrypted image is recovered from the encrypted image I_{ENC} , using the secret keys: vector K_R , vector K_C and $ITER_{max}$.

1) Initiate the counter $ITER = 0$.

2) Increment the counter $ITER = ITER + 1$.

3) Using the K_R vector, the Bit wise operation XOR will be applied to each column of the encrypted image, using the following rule:

$$I_1(i, 2j - 1) = I_{ENC}(2i - 1, j) \oplus K_R[j];$$

$I_1(i, j) = I_{SCR}(2i, j) \oplus \text{rot180}K_R[j]$, where rot180 represents the 180 degree flip of vector K_R .

4) Using the K_C vector, the Bit wise operation XOR will be applied to each row of the image I_1 , using the following rule:

$$I_{SCR}(2i - 1, j) = I_1(2i - 1, j) \oplus K_C[j];$$

$I_{SCR}(2i, j) = I_1(2i, j) \oplus \text{rot180}K_C[j]$, where rot180 represents the 180 degree flip of vector K_C .

Steps 3) and 4) will recreate the original scrambled image,

5) For each column, a b vector will be calculated, representing the sum of all elements on the column j : $\sum_{i=0}^{M-1} I_{SCR}[i, j]$, $j = 0, 1, 2, \dots, N - 1$.

Compute Mb vector, where $Mb[i] = b[i] \% 2$.

Each column of the scrambled image I_{SCR} will be circular shifted by $K_C[i]$ positions according to the following rule:

If $Mb[i] = 0$ then up circular shift.

If $Mb[i] = 1$ then down circular shift.

6) For each row, an $a[i]$ vector will be calculated, representing the sum of all elements on the row i : $\sum_{j=0}^{N-1} I_{SCR}[i, j]$, $i = 0, 1, 2, \dots, M - 1$.

Compute Ma vector, where $Ma[i] = a[i] \% 2$.

Each row of the scrambled image I_{SCR} will be circular shifted by $K_R[i]$ positions according to the following rule:

If $Ma[i] = 0$ then right circular shift.

If $Ma[i] = 1$ then left circular shift.

7) If $ITER = ITER_{max}$, the decryption process is complete. Otherwise, the algorithm jumps back to step 2).

3. Implementation in C and results

The first implementation for AD Blackfin 533 has been done using C programming language, in order to test the functionality of the algorithm in a Digital Signal Processor environment.

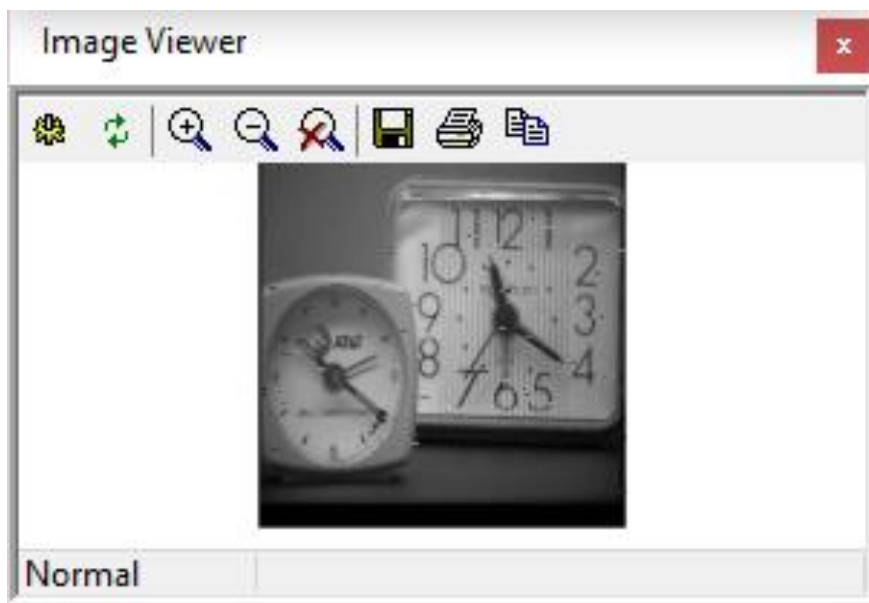
The C algorithm is based on the following key functions, which are based on the mathematical model presented in Section 2. They are used in both encryption and the decryption.

```
void Sum_Rows(unsigned char *input, int *a, unsigned char *Ma);  
void Rotate_Row(unsigned char *input, unsigned int pos, unsigned char dir);  
void Circular_Shift_Rows(unsigned char *input, unsigned char *Ma);  
void Sum_Columns(unsigned char *input, unsigned int *b, unsigned char *Mb);  
void Rotate_Column(unsigned char *input, unsigned int pos, unsigned char dir);  
void Circular_Shift_Columns(unsigned char *input, unsigned char *Mb);  
void Rows_XOR(unsigned char *input);  
void Columns_XOR(unsigned char *input);
```

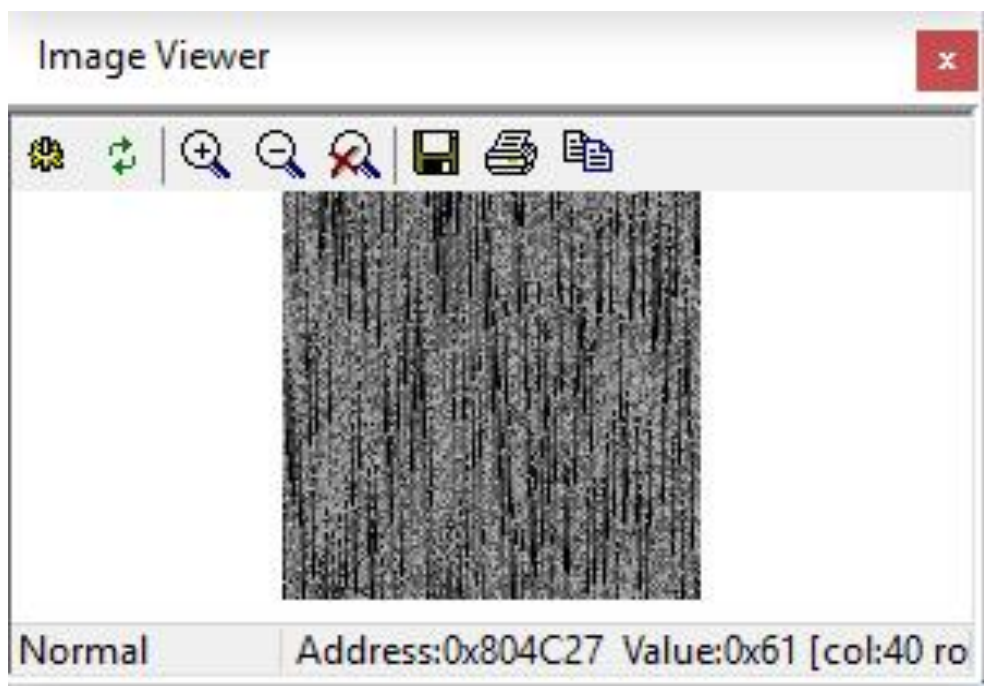
The following example illustrates the encryption of an 128x128 pixel gray scale image.

Example 1:

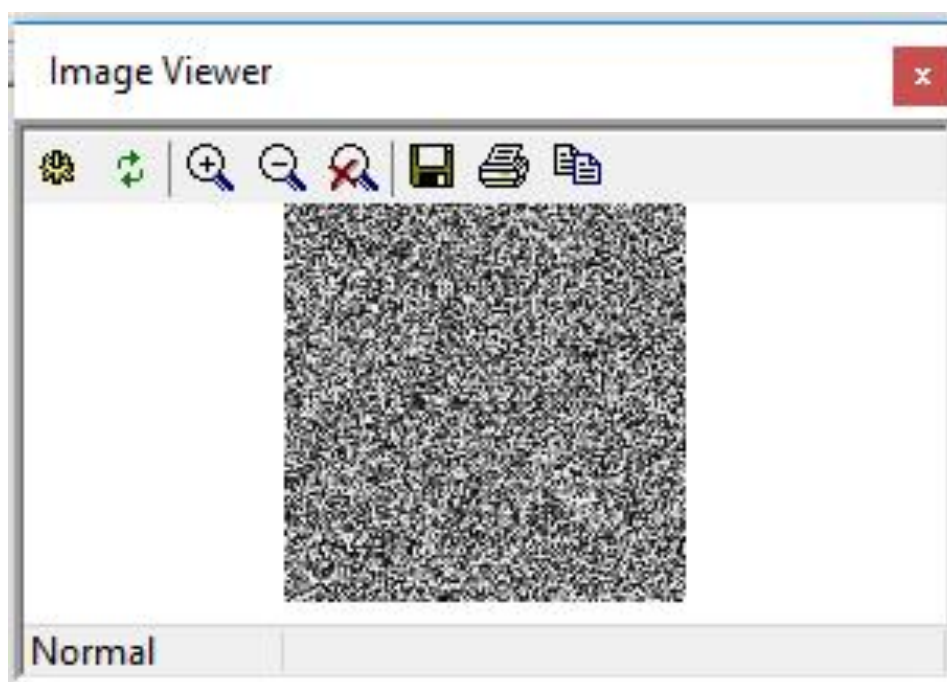
Original image:



Scrambled image:



Encrypted image:

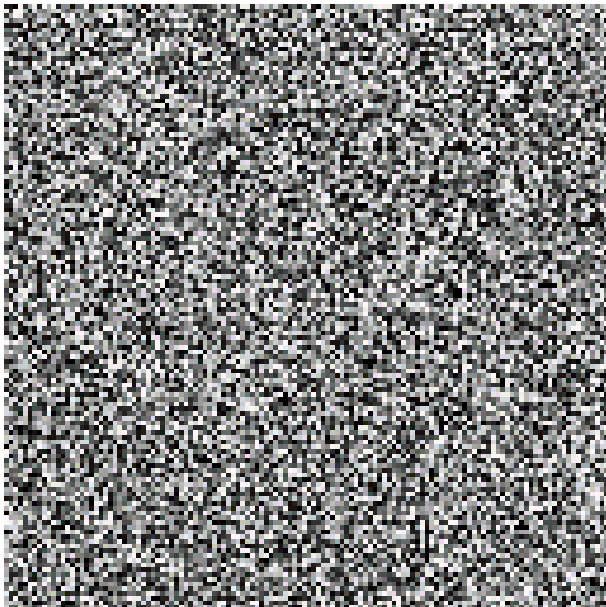


Example 2 presents the encryption of another 128x128 pixel gray scale image, “Lena”.

Original image:

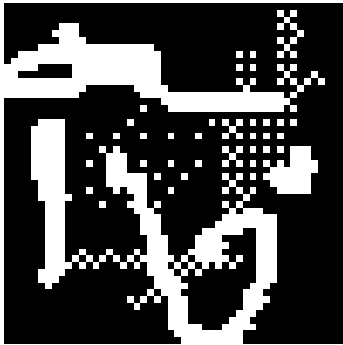


Encrypted image:

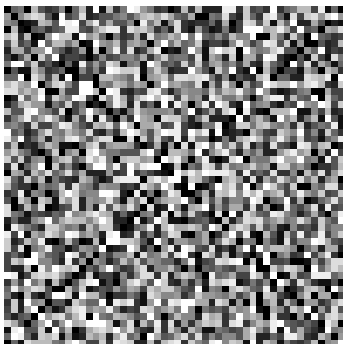


Example 3 presents the encryption of a 50x50 pixel gray scale image.

Original image:



Encrypted image:



4. Implementation in ASM and optimization

I have chosen to use Assembly language to optimize the algorithm for Blackfin 533. The following section firstly presents optimized functions, called in separately from the main project in order to illustrate speed improvement. The second part of the section analyses the global effect of the ASM written code.

The first functions that have been optimized have been **Sum_Rows** and **Sum_Columns**.

Profiling C implementation:

```
41
42 void Sum_Rows(unsigned char *input, unsigned char *a, unsigned char *Ma)
43 {
44     unsigned int sum;
45
46     unsigned char i, j;
47     for (i = 0; i < M; i++)
48     {
49         sum = 0;
50         for (j = 0; j < N; j++)
51         {
52             sum += input[i * M + j];
53         }
54     }
55 }
```

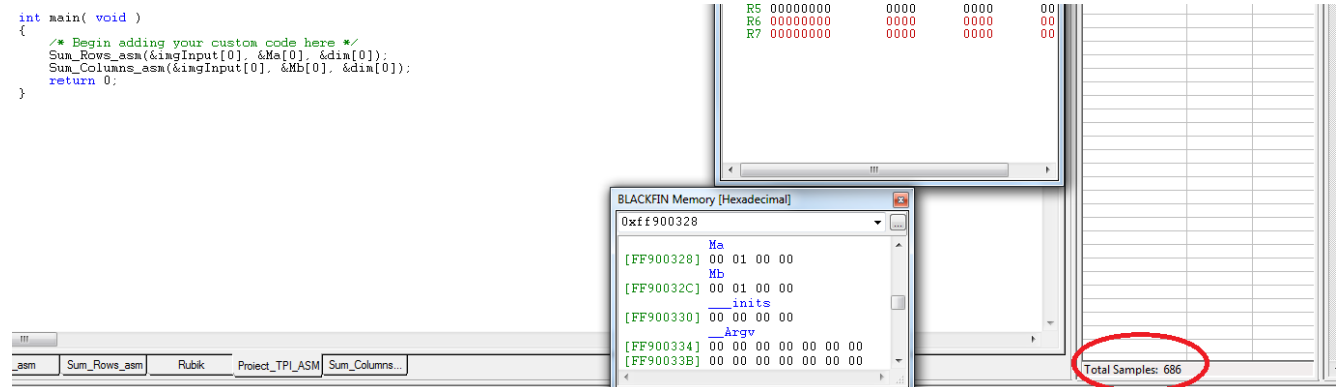
BLACKFIN Memory [Hexadecimal]

0xff900328	
[FF900328]	00 01 00 00
	Ma
[FF90032C]	00 01 00 00
	Ma
[FF900330]	06 0F 18 00
	a
[FF900334]	0C 0F 12 00
	b
	__init

Total Samples: 1259

Total samples 1259 after calling the C functions.

Profiling ASM implementation:

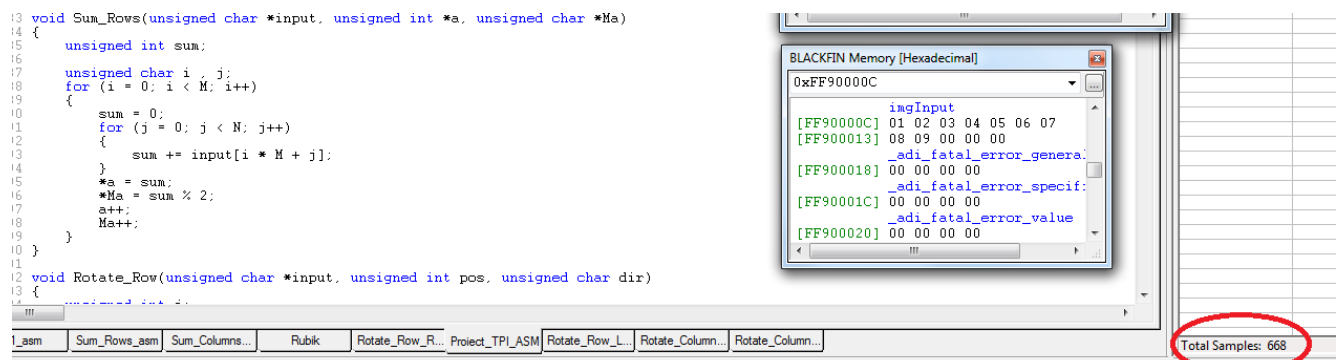


Total samples 686 after calling the ASM functions.

Rotate_Row() and **Rotate_Column()** functions have also been optimized for AD BF 533.

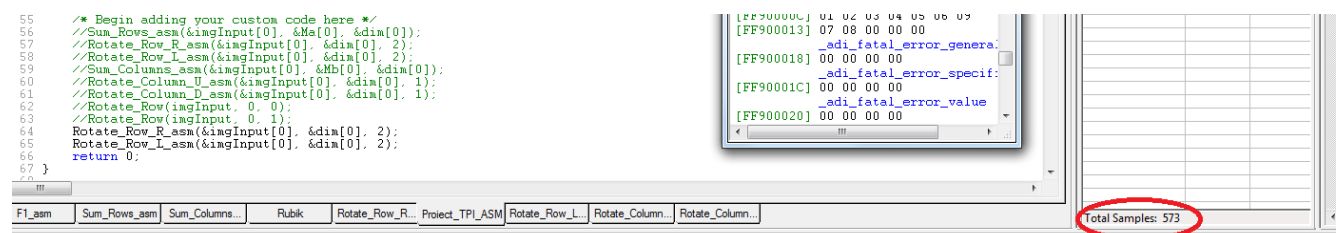
Rotate_Row() Profiling:

C implementation:



Total samples 668 after calling the C function.

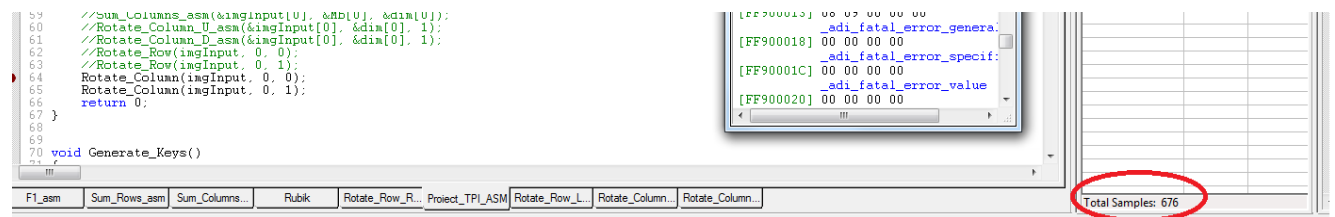
ASM implementation;



Total samples 573 after calling the ASM function.

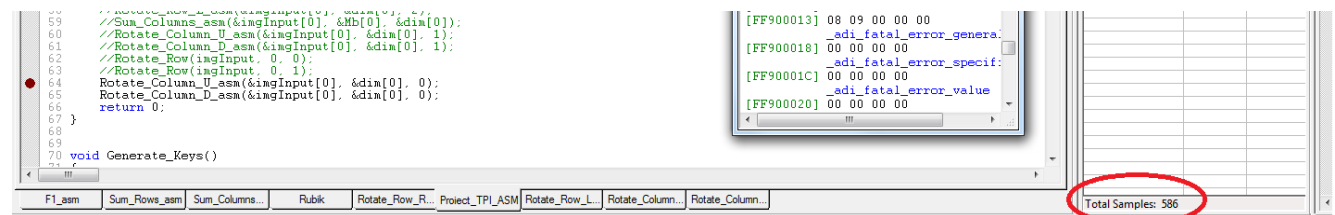
Rotate_Column() Profiling:

C implementation:



Total samples 676 after calling the C function.

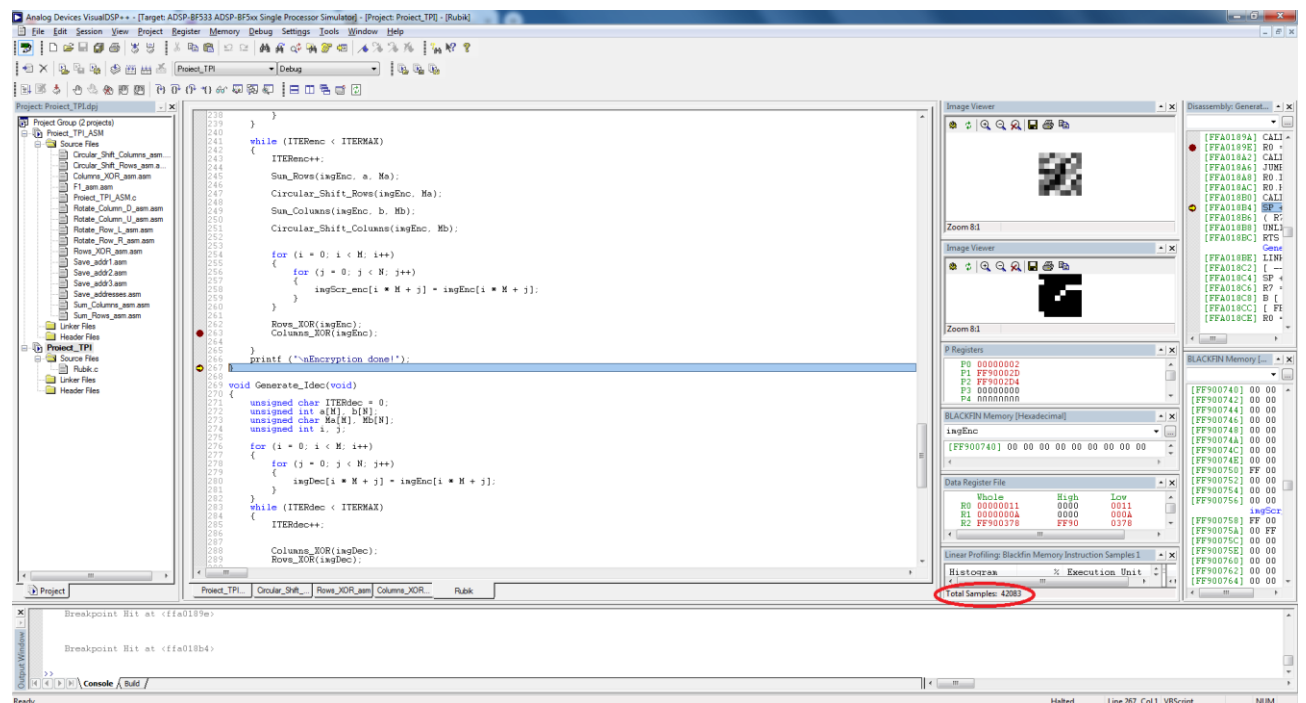
ASM implementation:



Total samples 586 after calling the ASM function.

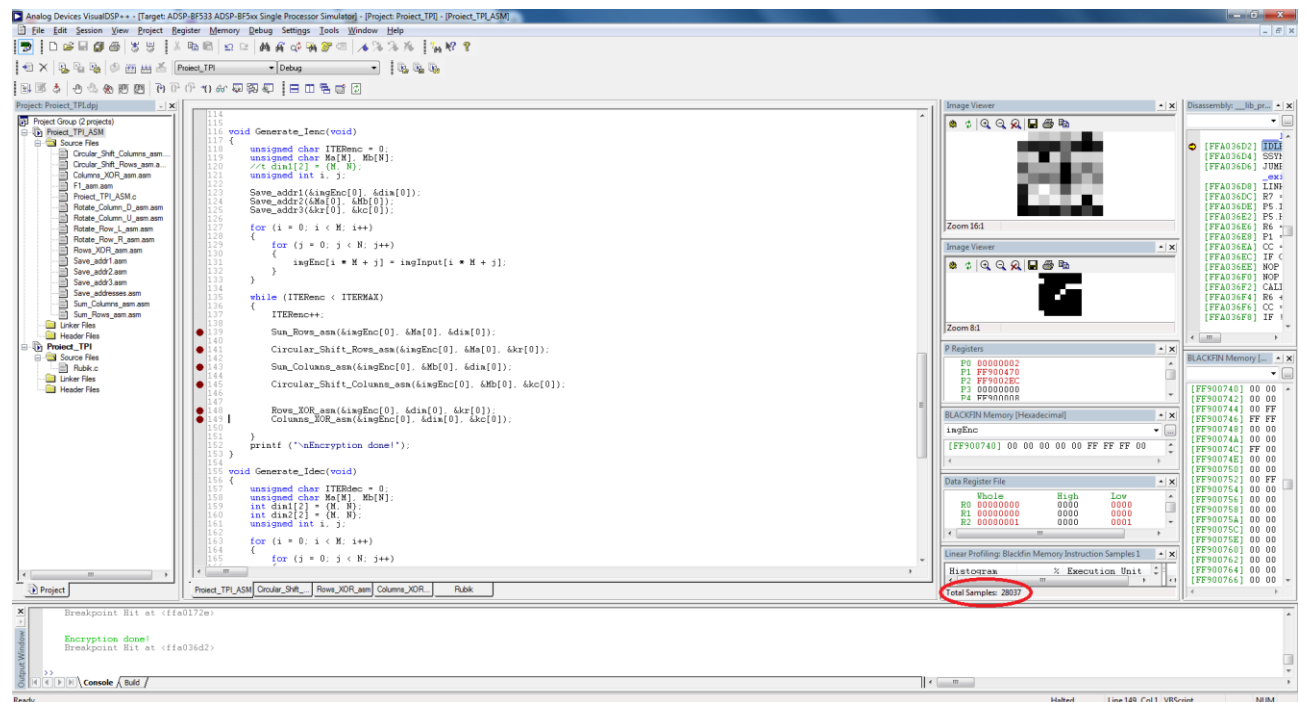
Additionally, **Rows_XOR()** and **Columns_XOR()** functions have also been implemented in ASM. The following examples present the global effects of the optimization.

Encryption C:



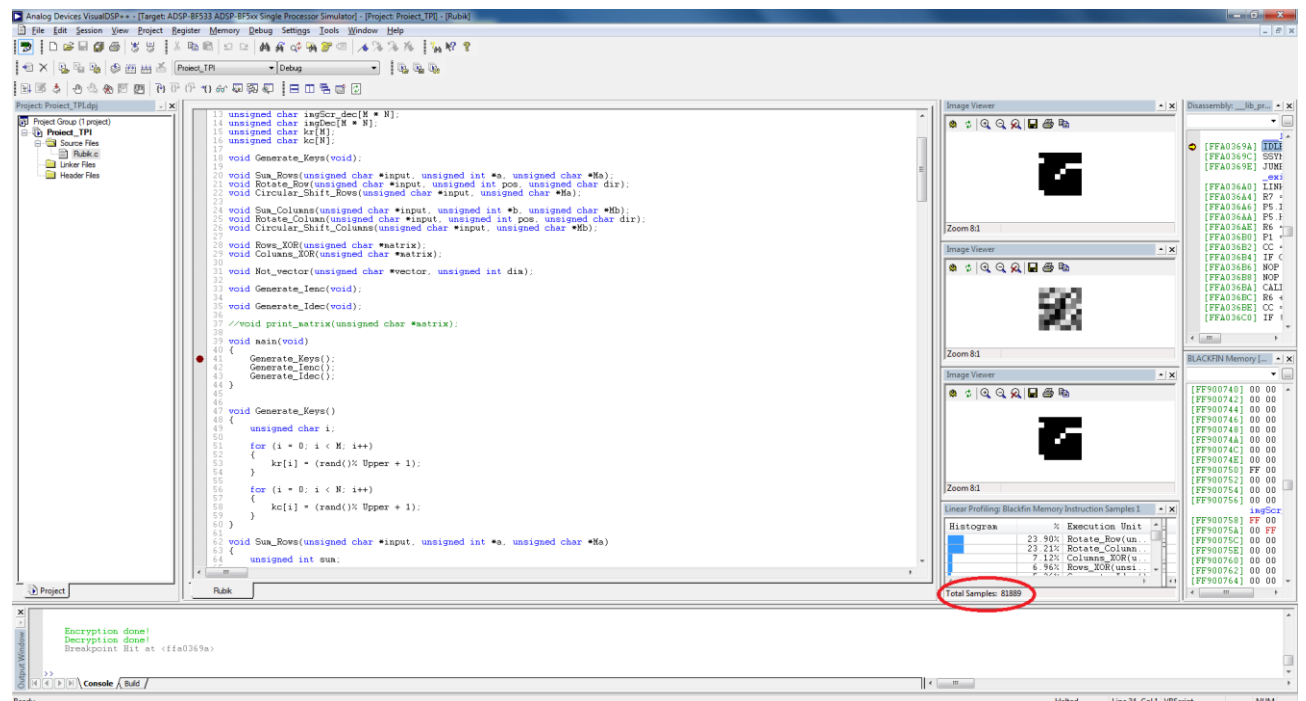
Total samples: 42083.

Encryption ASM:



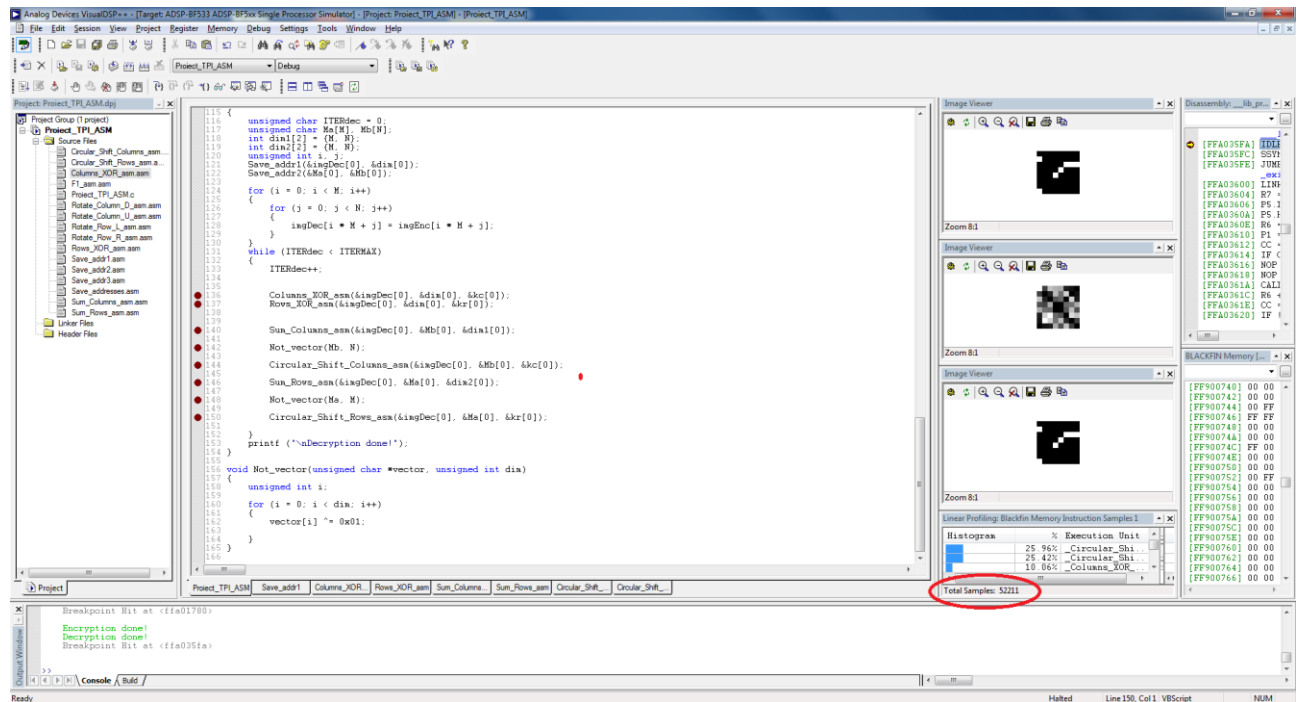
Total samples: 28037.

Full algorithm for an 8x8 gray scale image C:



Total samples: 81889.

Full algorithm for an 8x8 gray scale image ASM:



Total samples: 52211.

After running the project with the optimizations presented above, the algorithm runs with 36% faster than the C implementation.

5. Conclusions

The Rubik's Cube encryption algorithm is a real time image encryption algorithm, which runs even better on AD BF533 after optimization. It is a secure and robust algorithm, resistant to brute force attacks due to its large key space. Key sensibility is also present in our algorithm, a small change in the triplet of the secret keys (K_R , K_C , $ITER_{max}$) will create a big change in the output. In the same idea, for most images, the algorithm changes the value of 98% of the input image's pixels, making it resistant to statistical analysis attacks.

I would like to thank the researchers from source [1] for their contribution and for making the realization of this project possible. My contribution to this project was the following: I have read and understood the algorithm, I have implemented and tested it in C programming language, I optimized it for AD BF533 having good results overall.

6. References

[1] *Journal of Electrical and Computer Engineering* Volume 2012, Article ID 173931, 13 page; Authors: Khaled Loukhaoukha, Jean-Yves Chouinard, and Abdellah Berdai