

Statement

↳ Given an **array**, which contains a combination of the following three elements:

- 0 (Representing red)
- 1 (Representing white)
- 2 (Representing blue)

Sort the array **in place** so that the elements of the same colour are adjacent, and the final order is: red (0), white (1), blue (2).

Naive

↳ Use conventional sorting algorithm, i.e., Merge sort gives $O(n \log n)$.

Optimal

↳ Order elements in single pass

- Initialize pointers
 - **start** and **current**: initially point to the **first** index
 - **end**: initially point to the **last** index
- Iterate over array until **current** pointer is greater than **end** pointer.
- During iteration, there are **3** conditions:
 - If **colours[current]** is 0
 - ↳ The **current** pointer points to red. **Swap** **colours[current]** and **colours[start]**. This places the red element at the start of the array.
 - ↳ Increment both **start** and **current** pointers by **one**. Moving **start** ensures that the next red element will occupy the correct position.
 - If **colours[current]** is 1
 - ↳ The white element is already in its correct section, **no swapping** required.

↳ Increment the **current** pointer by **one**.

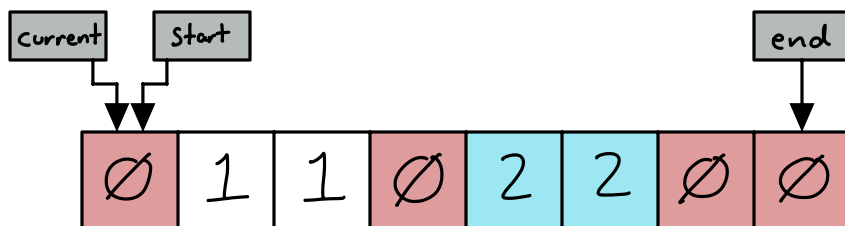
• If **colours[current]** is **2**

↳ The **current** pointer points to blue. **Swap** **colours[current]** and **colours[end]**. This pushes the blue element to the end of the array.

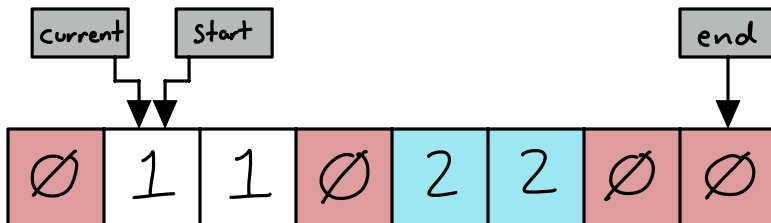
↳ Decrement **end** pointer by **one**.

Visualization

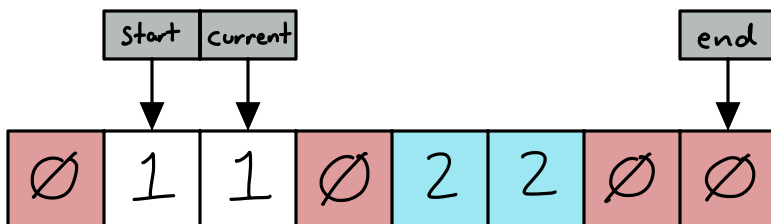
i) **colours[current] == \emptyset** so swap the elements of **colours[current]** and **colours[start]**.
Increment both the **current** and **start** pointers.



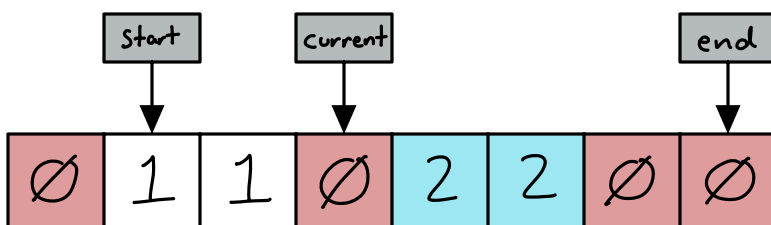
ii) **colours[current] == 1**
Increment **current** pointer



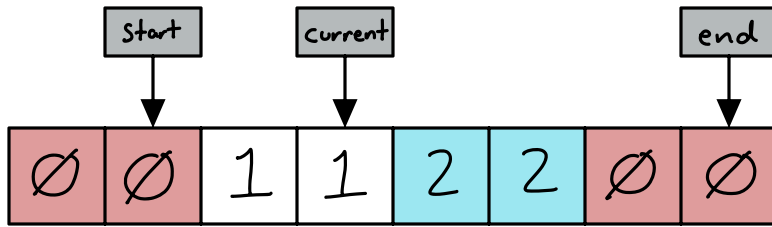
iii) **colours[current] == 1**
Increment **current** pointer



iv) **colours[current] == \emptyset**
Swap the elements of **colours[current]** and **colours[start]**.

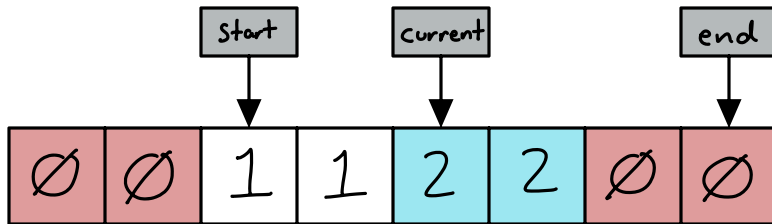


v) Increment both **current** and **start** pointers.



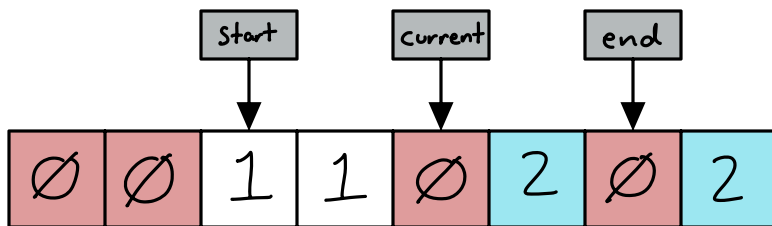
vi) **colours[current] == 2**

Swap the elements of **colours[current]** and **colours[end]** and decrement the **end** pointer.

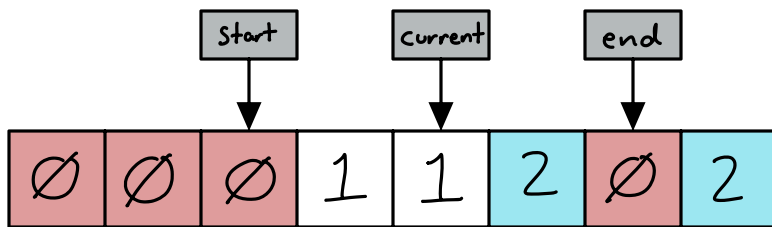


vii) **colours[current] == \emptyset**

Swap the elements of **colours[current]** and **colours[start]**.

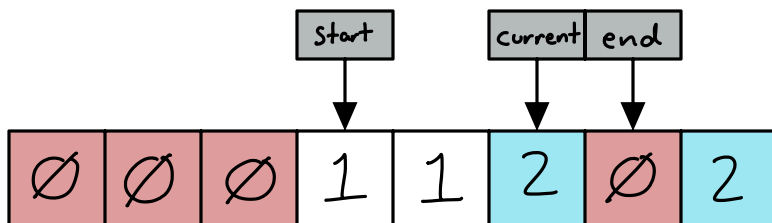


viii) Increment **start** and **current** pointers



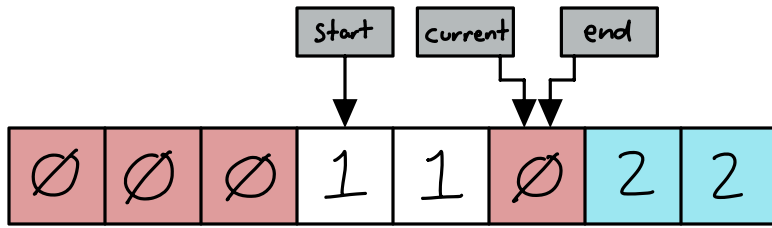
ix) **colours[current] == 2**

Swap the elements of **colours[current]** and **colours[end]** and decrement the **end** pointer.

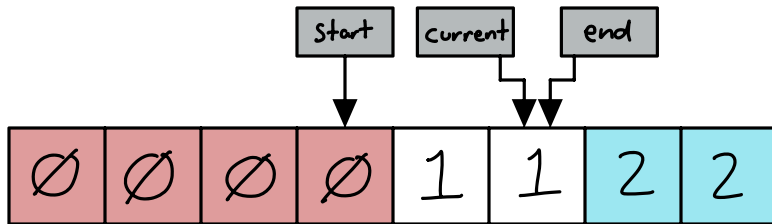


x) $colours[current] == \emptyset$

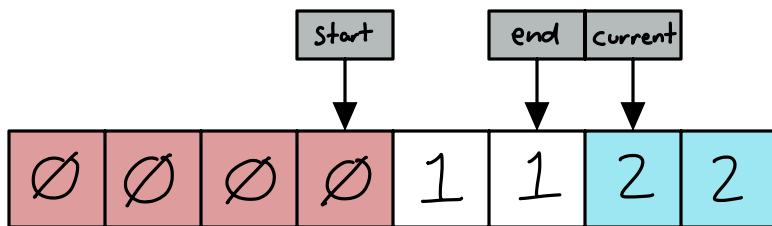
Swap $colours[current]$ and $colours[start]$



xi) Increment both $current$ and $start$ pointers.



xii) The end pointer is less than the $current$ pointer, no further swapping can be performed, return $colours$.



Code

```
vector<int> SortColors(vector<int>& colors) {  
    int start = 0, current = 0, end = colors.size() - 1;  
  
    while (current <= end) {  
        if (colors[current] == 0) {  
            swap(colors[start], colors[current]);  
            current++;  
            start++;  
        } else if (colors[current] == 1) {  
            current++;  
        } else {  
            swap(colors[current], colors[end]);  
            end--;  
        }  
    }  
    return colors;  
}
```

Time Complexity

↳ Only traverse array once, thus $O(n)$.

Space Complexity

↳ No extra space used, thus $O(1)$.