

# Adrenaline

Pinpointing and Reining in Tail Queries  
with Quick Voltage Boosting

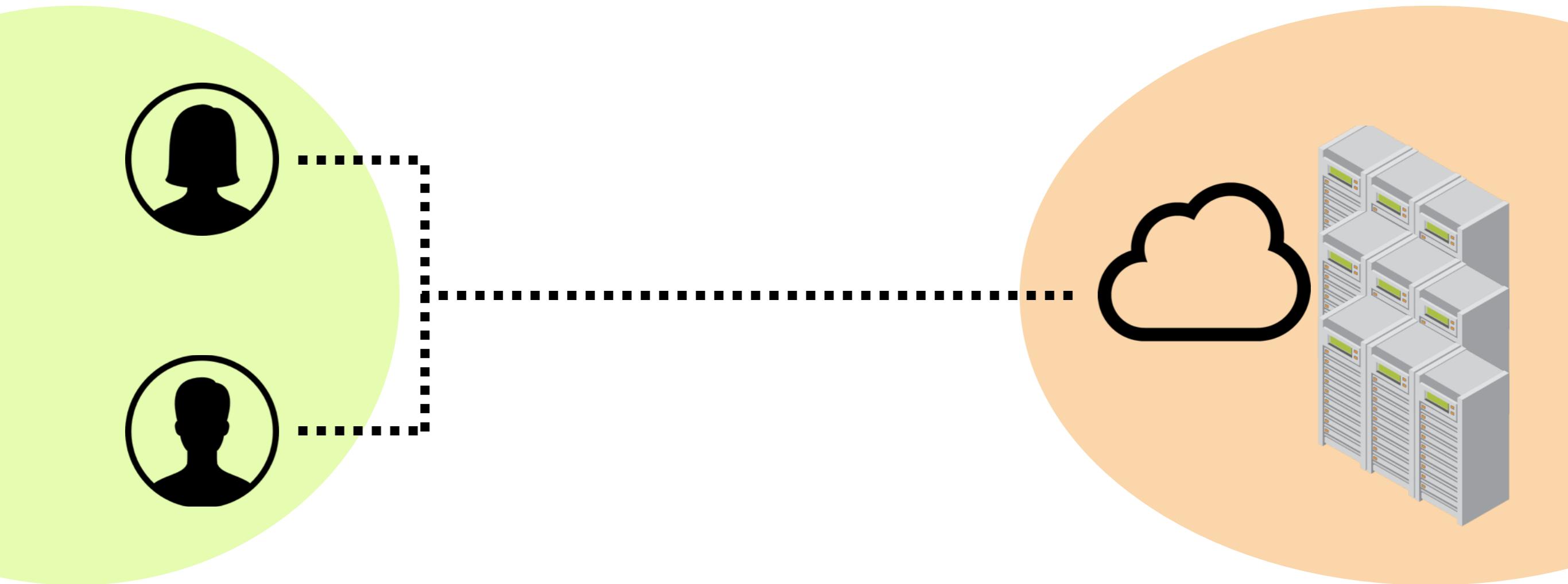
Chang-Hong Hsu, Yunqi Zhang, Michael A. Laurenzano,  
David Meisner, Thomas Wenisch, Jason Mars,  
Lingjia Tang, Ronald G. Dreslinski



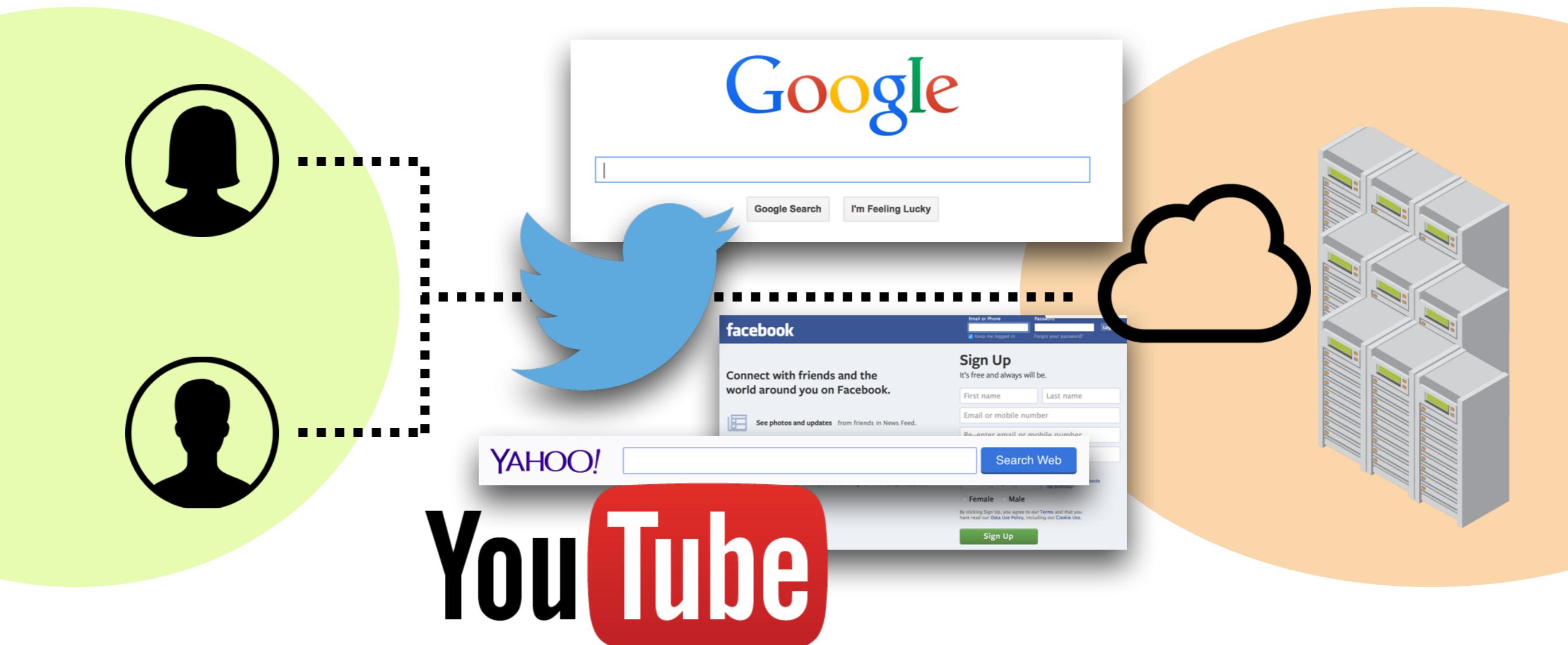
ClarityLab



# When datacenters meet users



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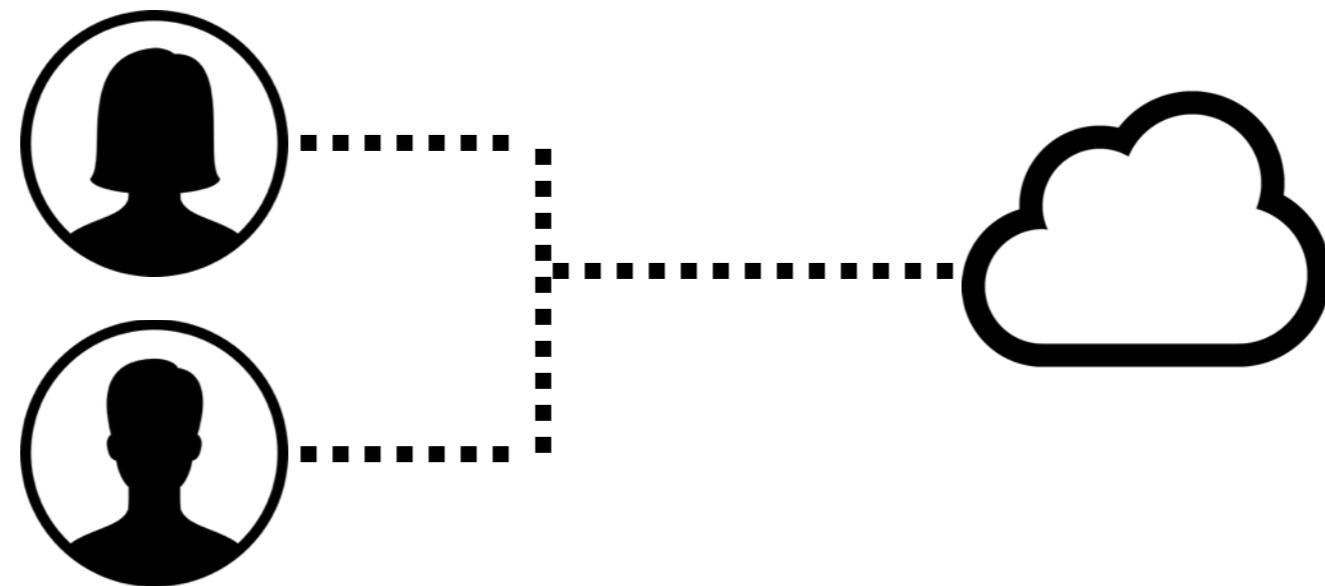


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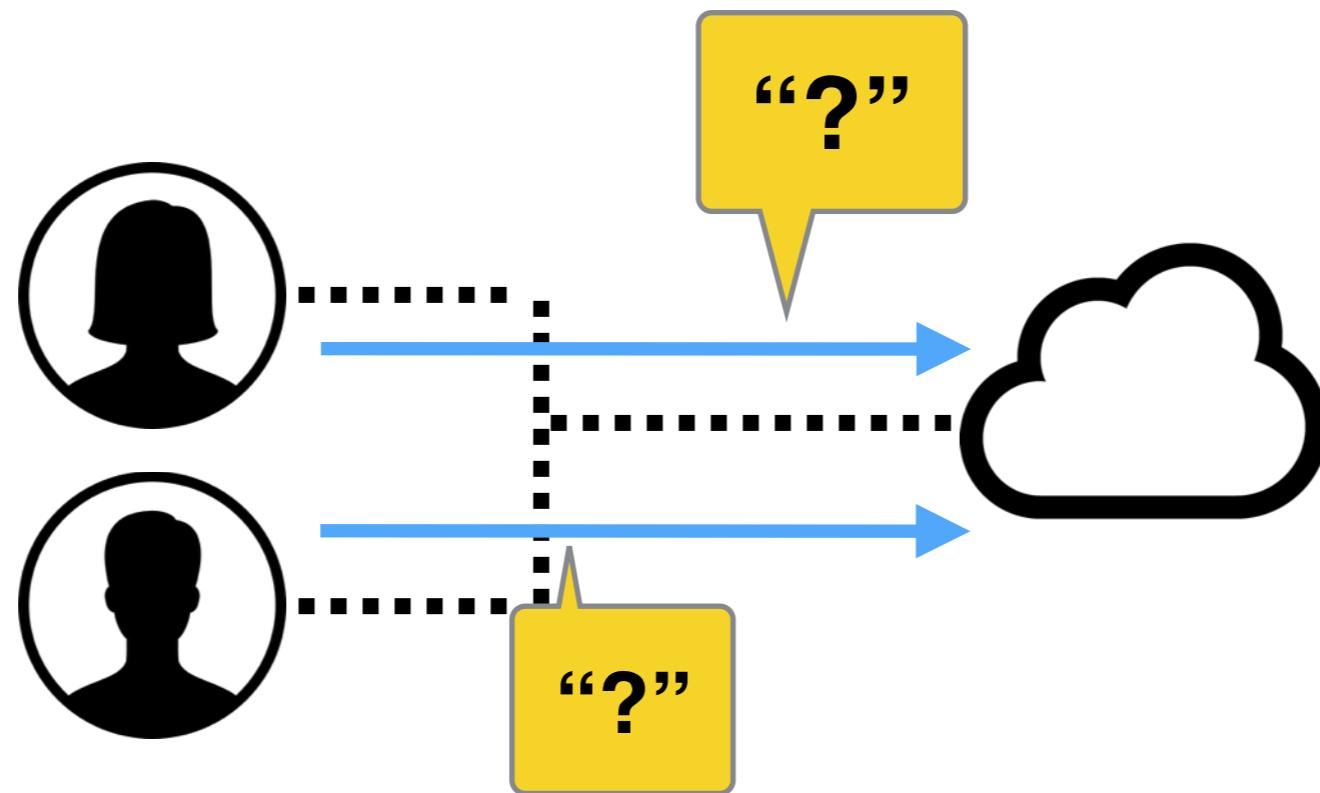
**User-facing services are  
everywhere!**

# User-facing services need to be FAST!



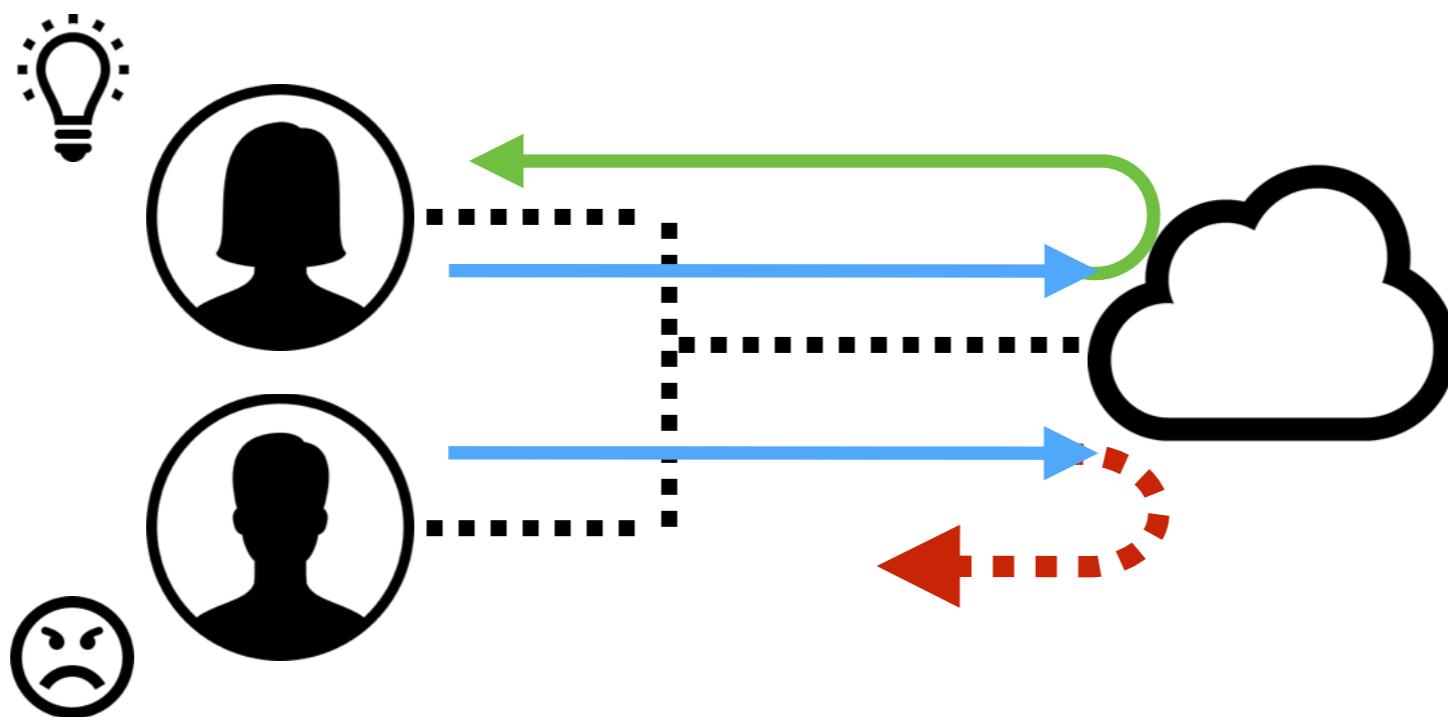
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- When users send queries,



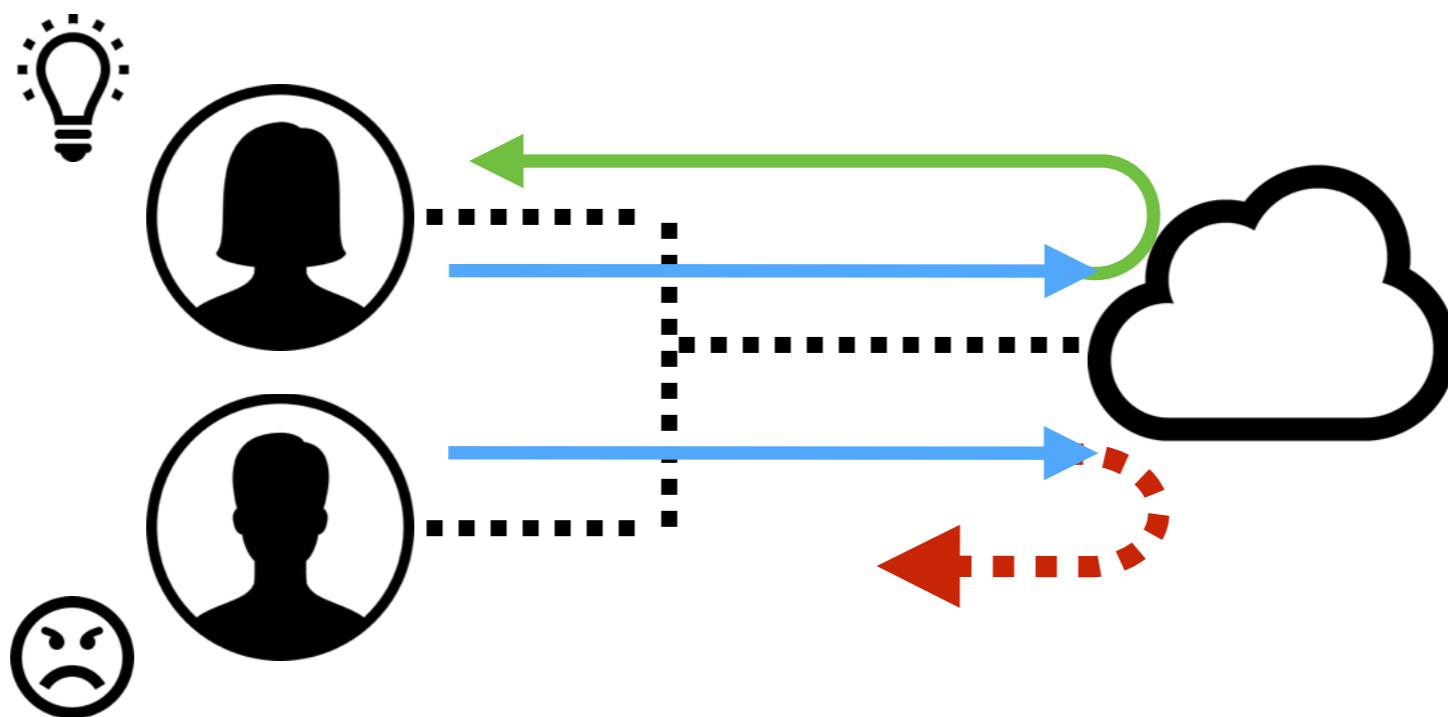
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**Latency is the key to the high quality of user-facing services**

# How fast is “FAST”?

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For user-facing services, “fast” means **low tail latency**

# “Tail latency?”

- Tail latency is the latency of the **slowest queries** of the entire distribution
- Tail latency represents the “**worst-case**” quality of a service (worst-case QoS)
- Tail latency directly relates to **user experience**



# “Fast” means low tail latency

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**Reduce the long TAIL latency makes a lot of sense!**

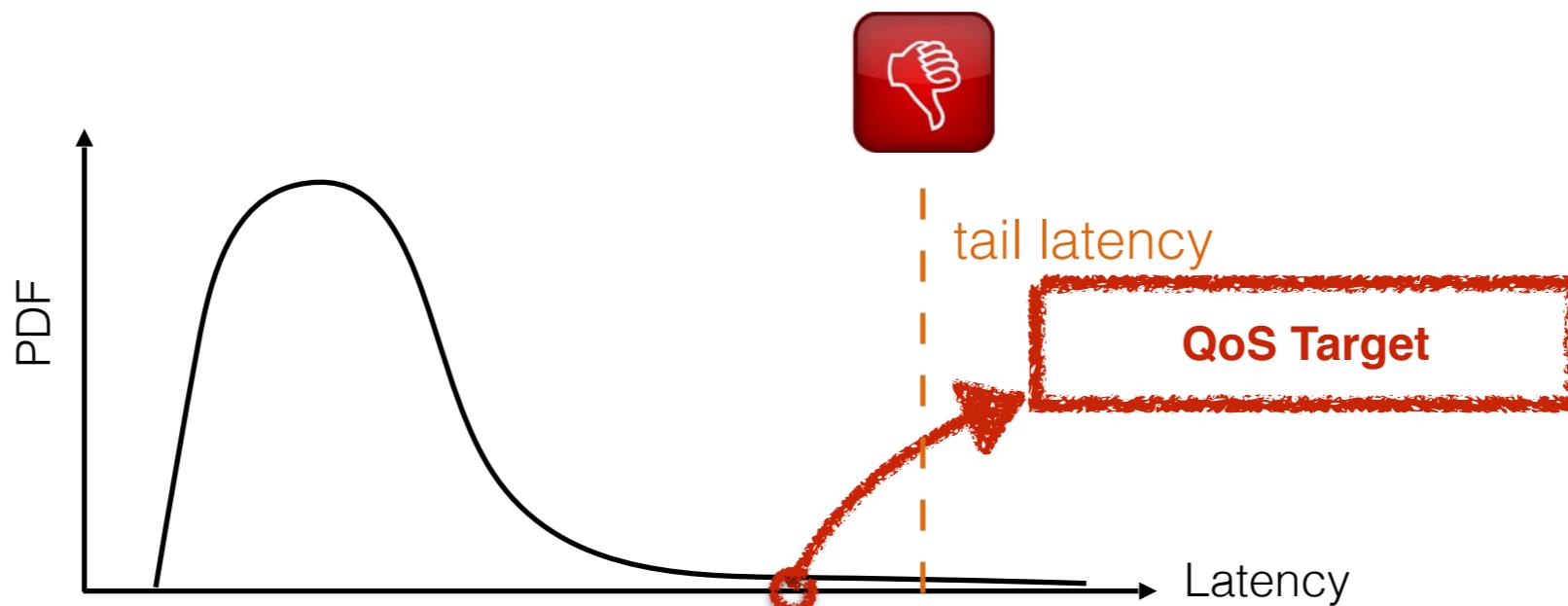
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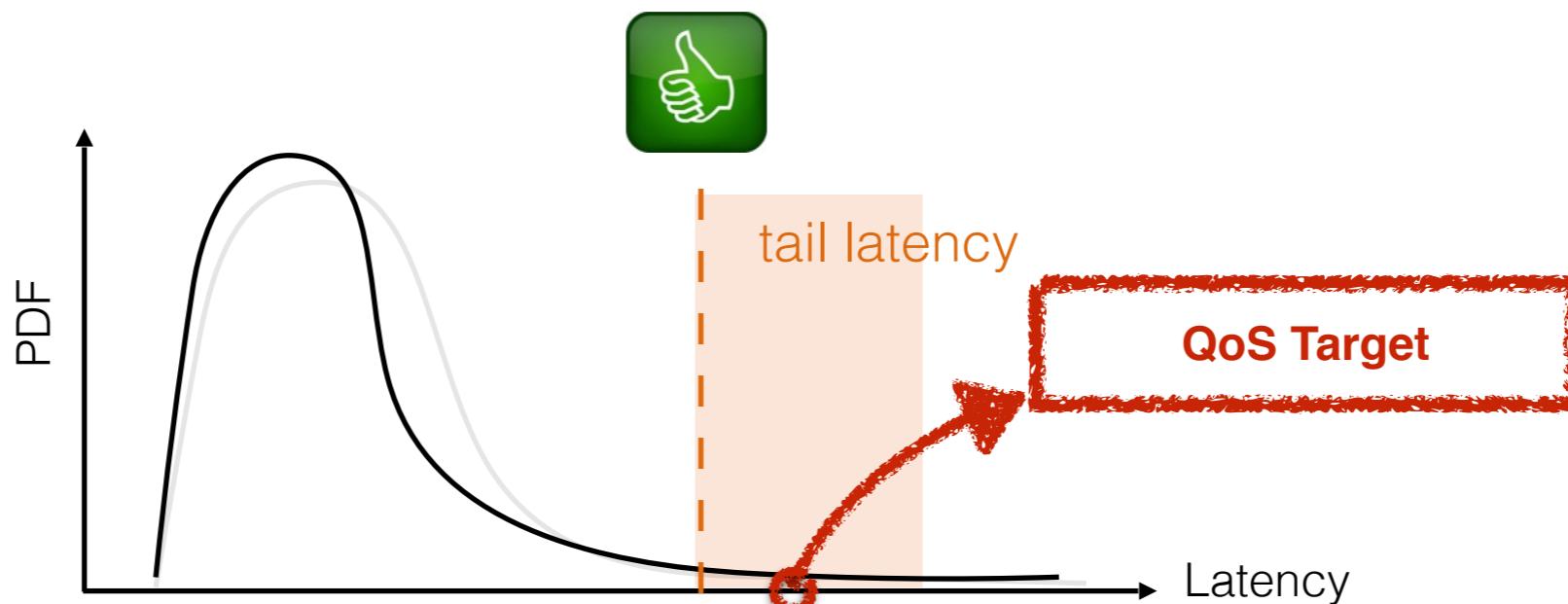
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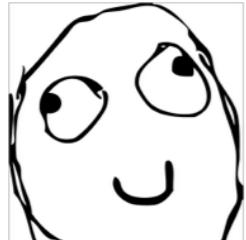
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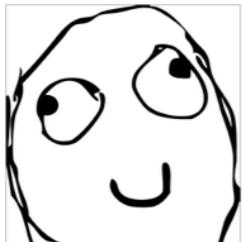
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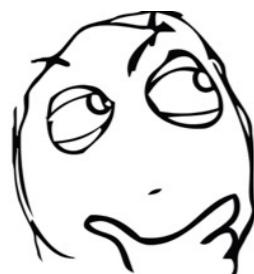




OK, I'll **boost the voltage/frequency of my cores** when the load is high, making the service run fast

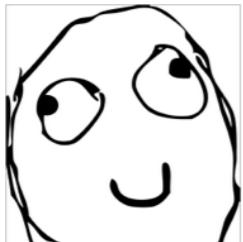


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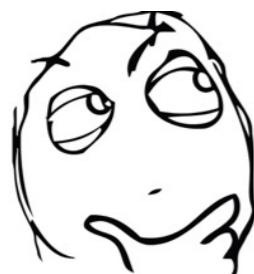


Wait a sec...I do want my service fast, but I remember  $P_{dyn} \propto f^3$

That means **A LOT of extra energy!** No!!



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Can I have a **fast service** with only **little energy overhead?**





**Yes!** We can achieve that goal if we can

1. **pinpoint those queries in the tail**
2. **boost these tail queries** specifically



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## Adrenaline introduces query-level boosting

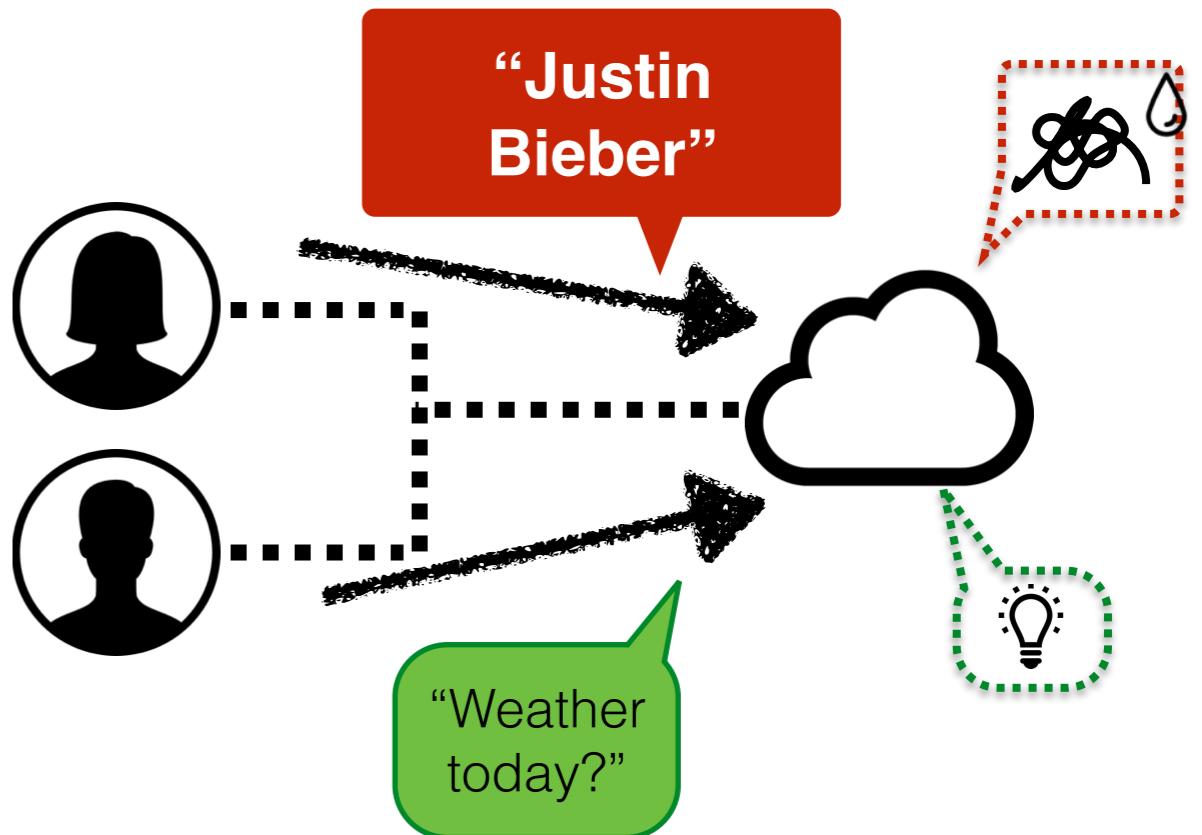
- ✓ Identify and take advantages of the differences among queries
- ✓ Switching core's voltage/frequency quickly to boost for the slow-running queries

# Adrenaline insights

| Adrenaline |

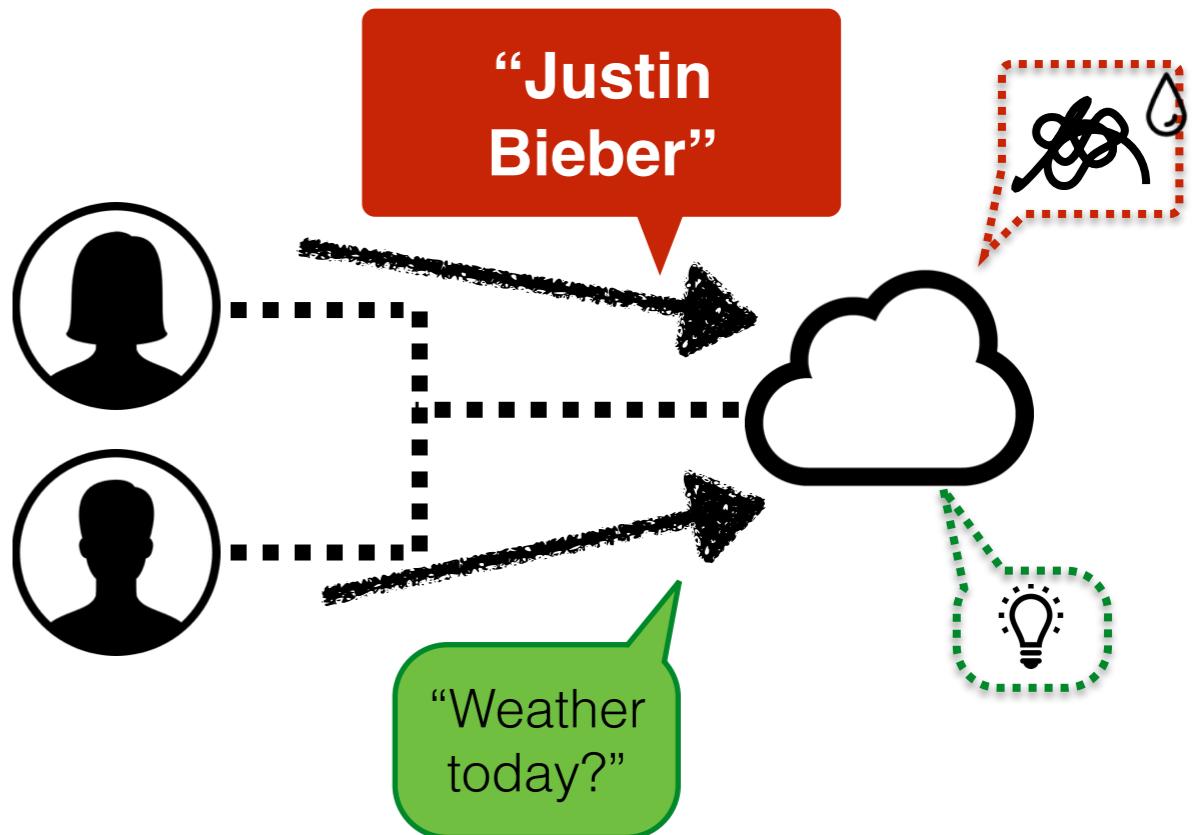
# Queries are highly variable

- Even for the same web service, queries
  - **come from different users**
  - **have different contents**
  - **require different actions**



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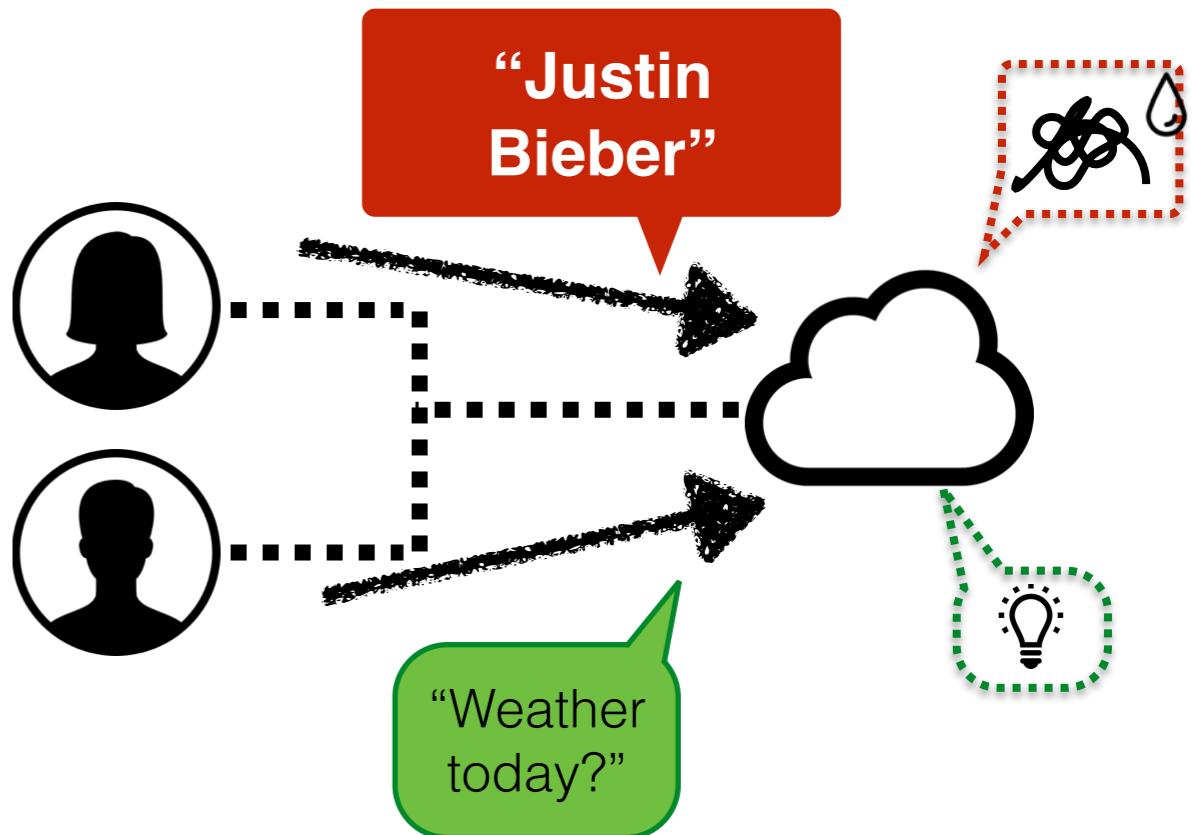
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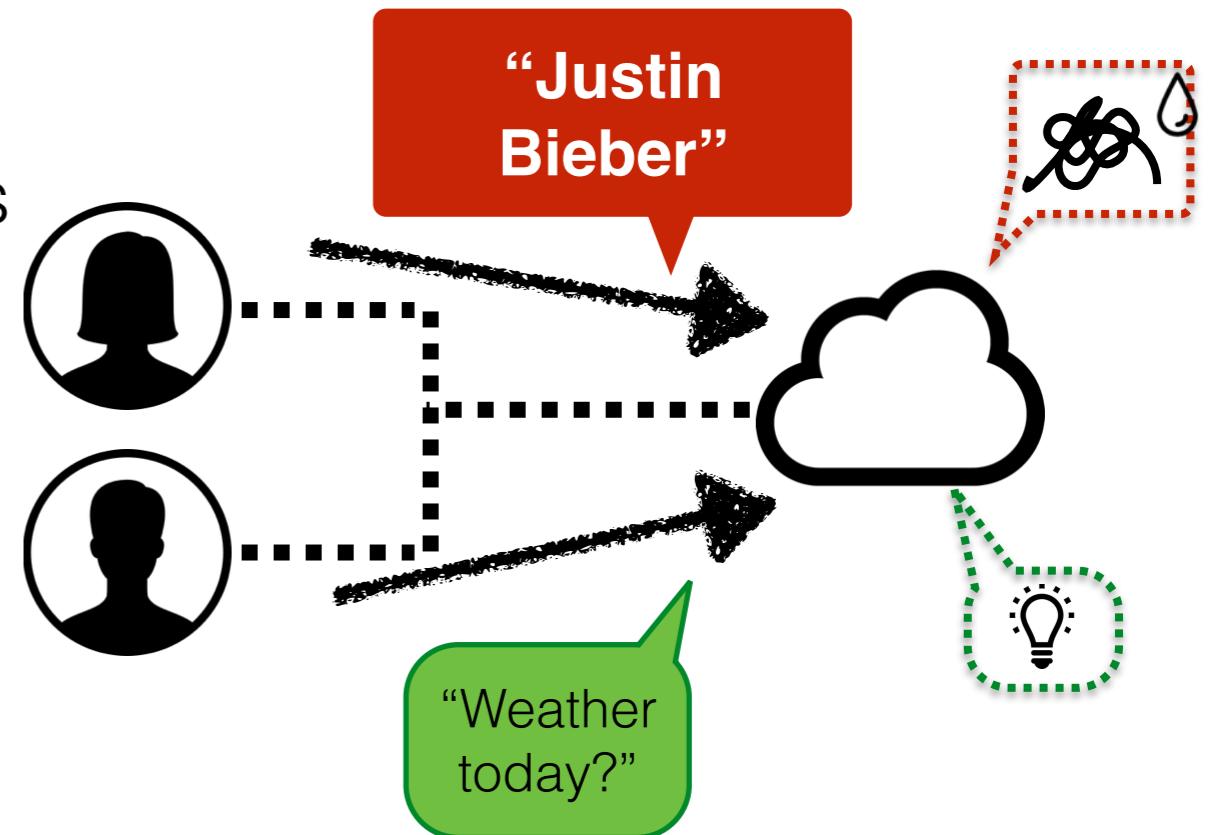


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**We observe that...**

- they need different amount of time to process
- they react differently to core's V/f scaling

# Variability matters!

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## Insight:

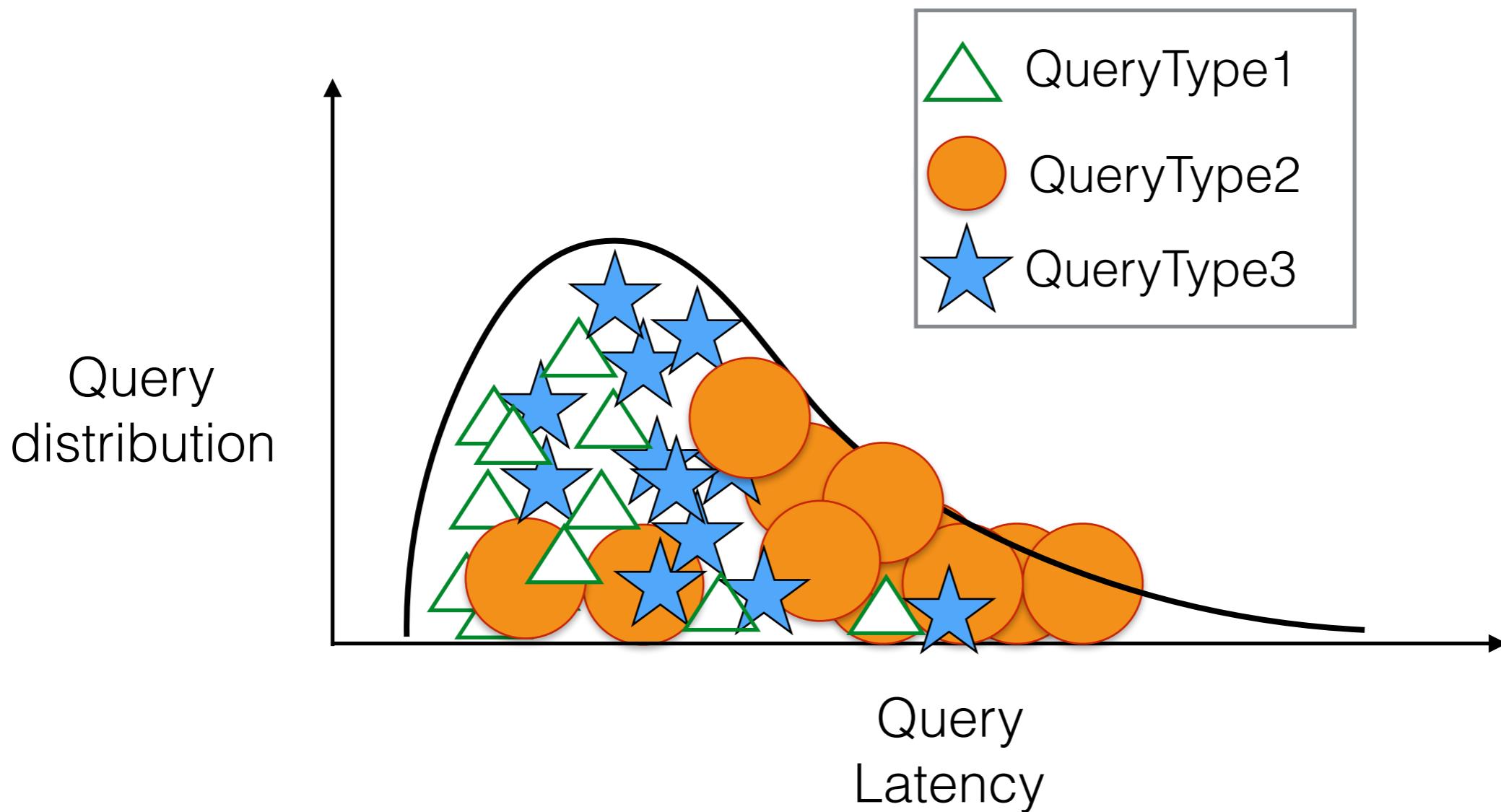
For a user-facing service, if we know

1. **“what query to target”**
2. **“how much I can boost”**

we can boost cores' V/f accordingly and intelligently

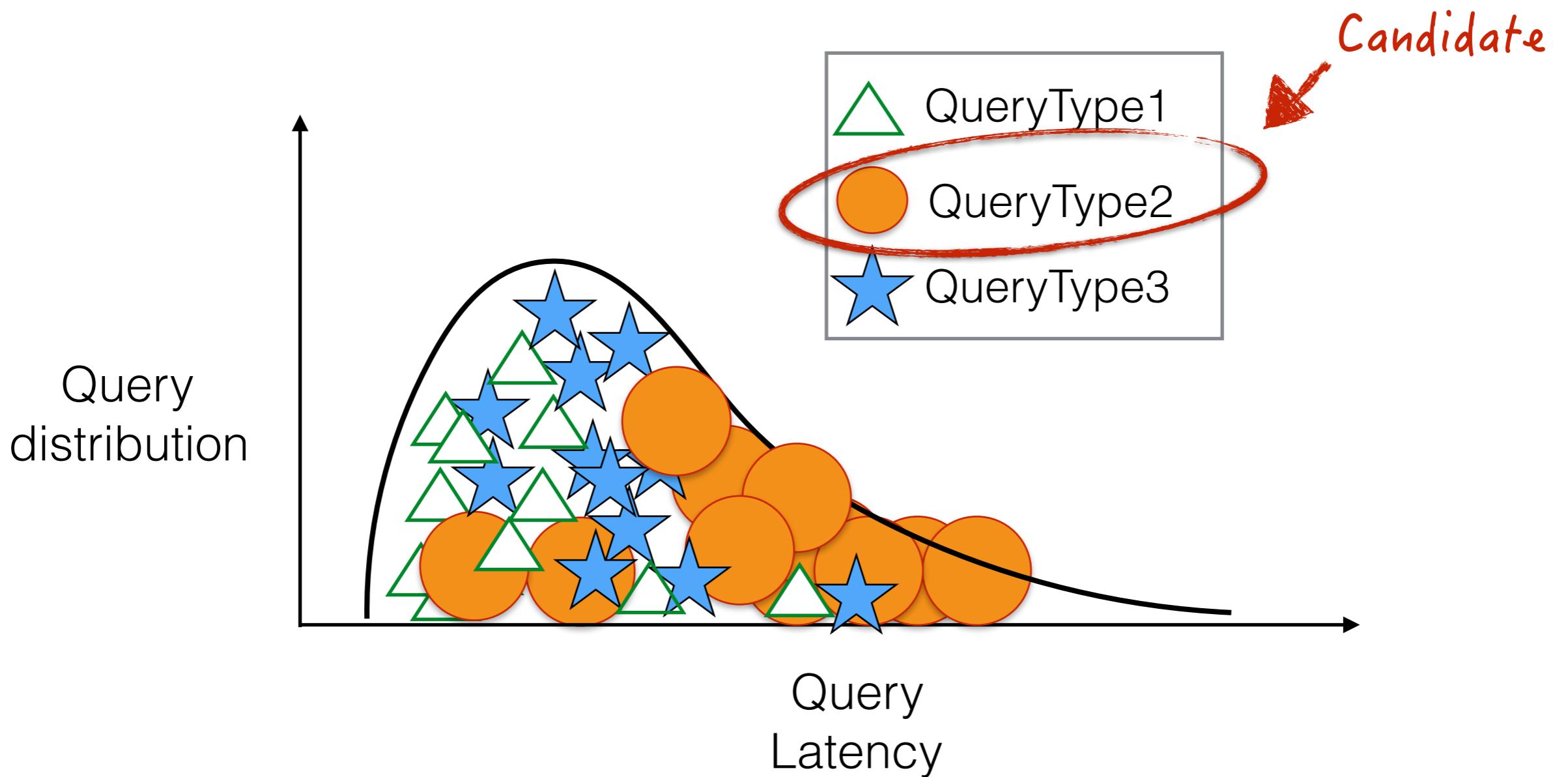
# “What query to target?”

Characteristic 1: **high contribution to the tail**

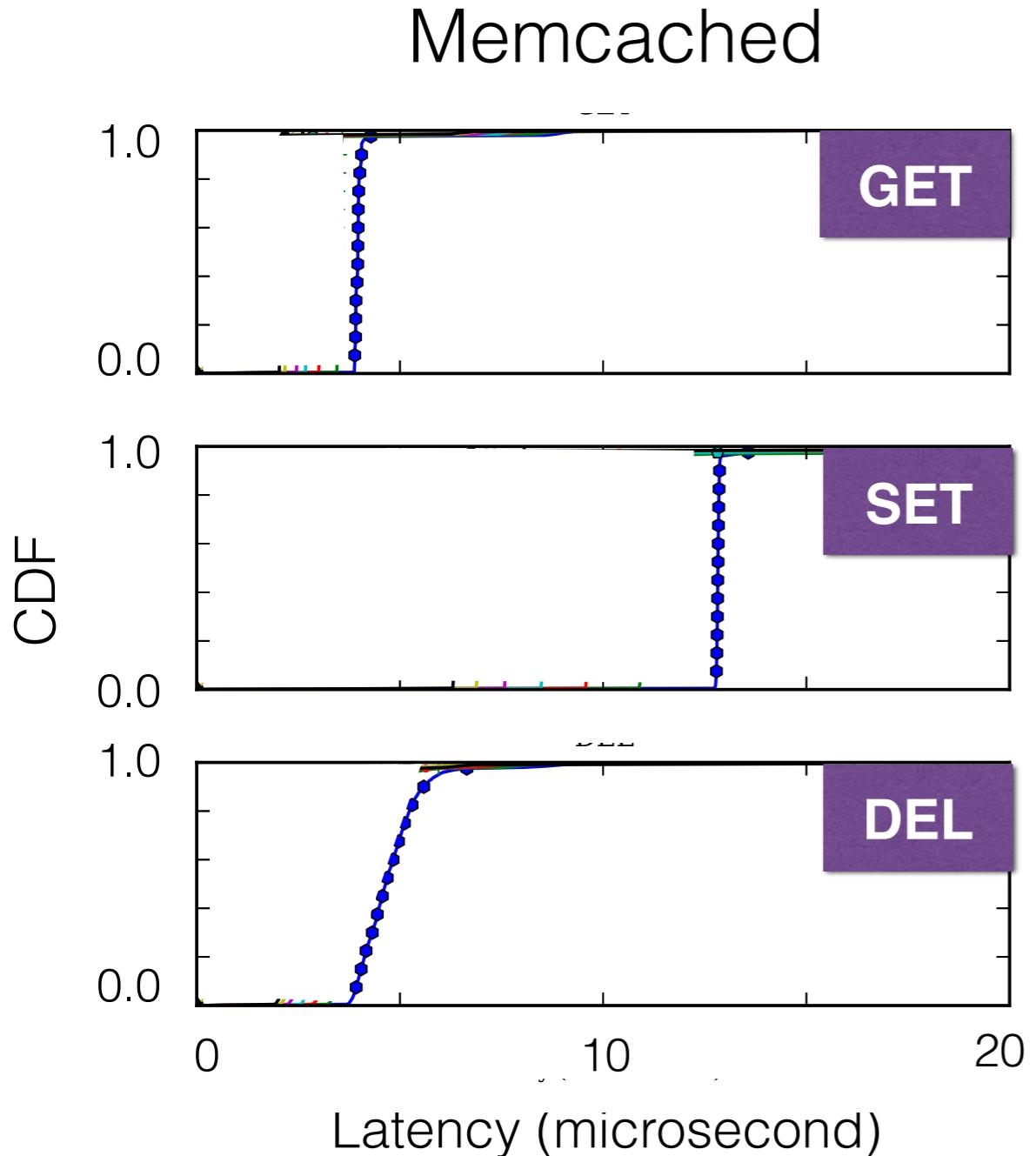


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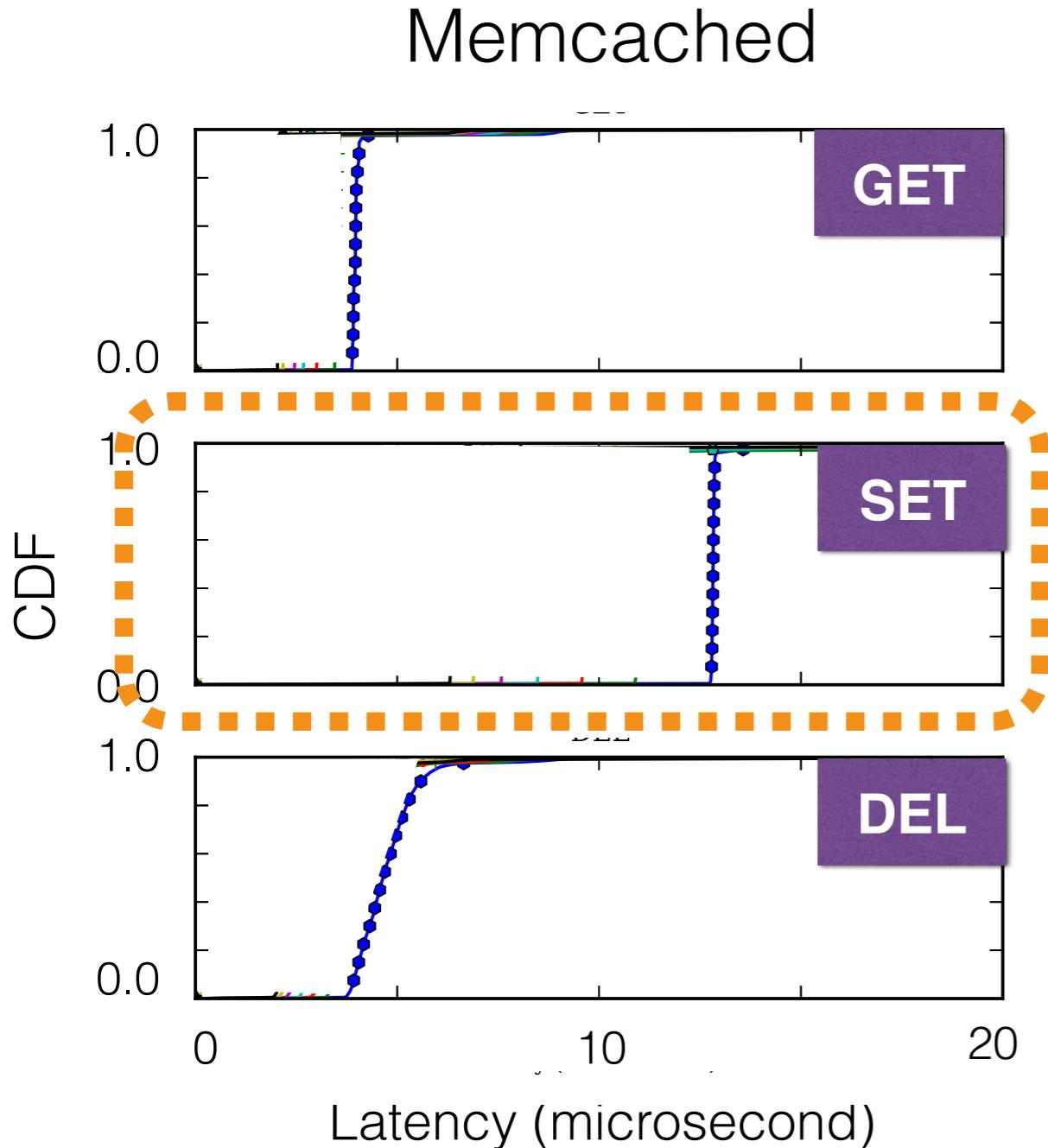
# High tail-contribution example in Memcached



## What is in this graph

- Latency distributions (CDFs) of different types of requests
- Running with a fixed core frequency

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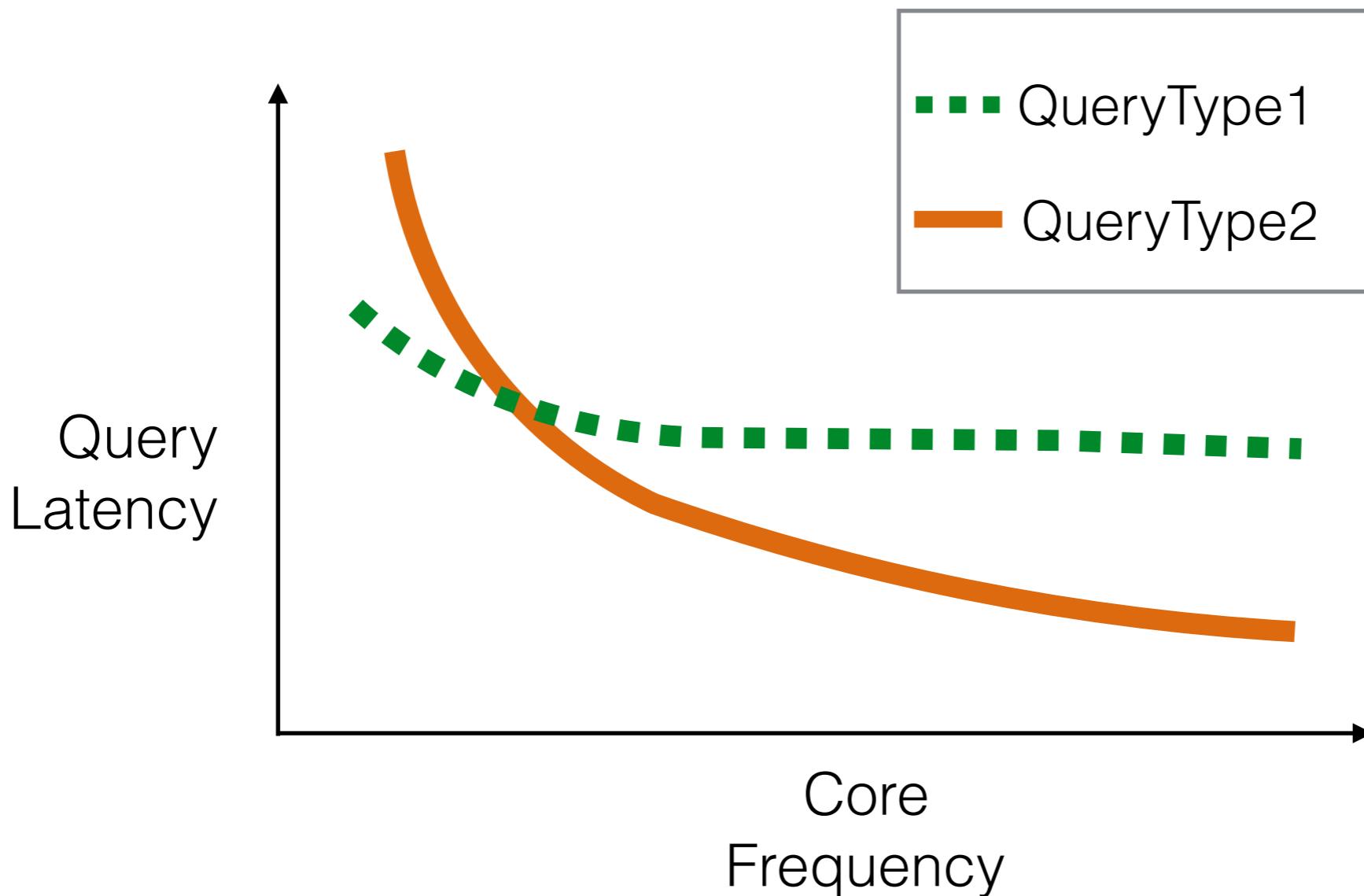
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## Observation

- **SETs have 2-3x longer tail latency**

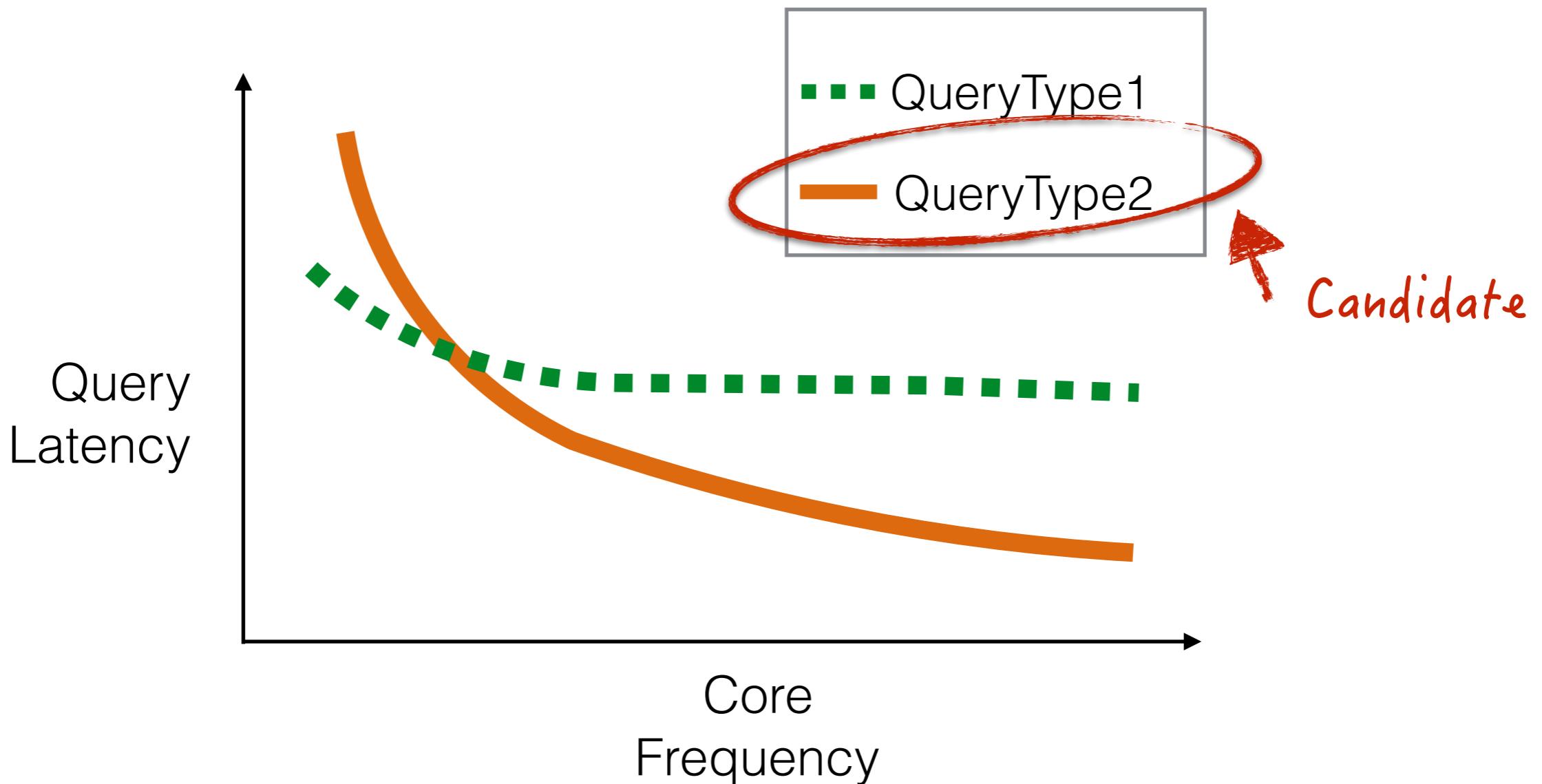
# How much can we boost?

Characteristic 2: **High boost-ability**



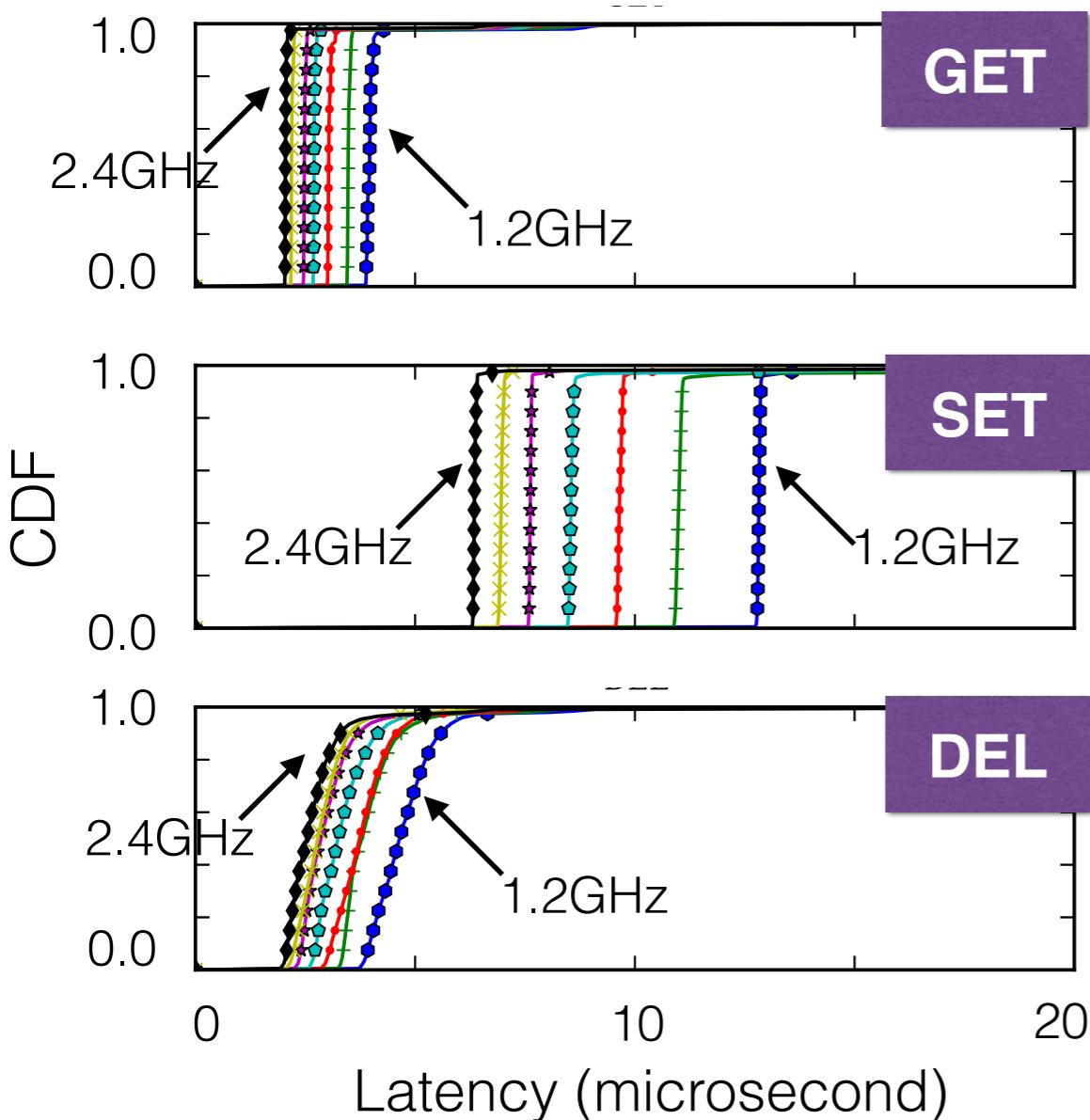
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# High boost-ability example in Memcached

Memcached

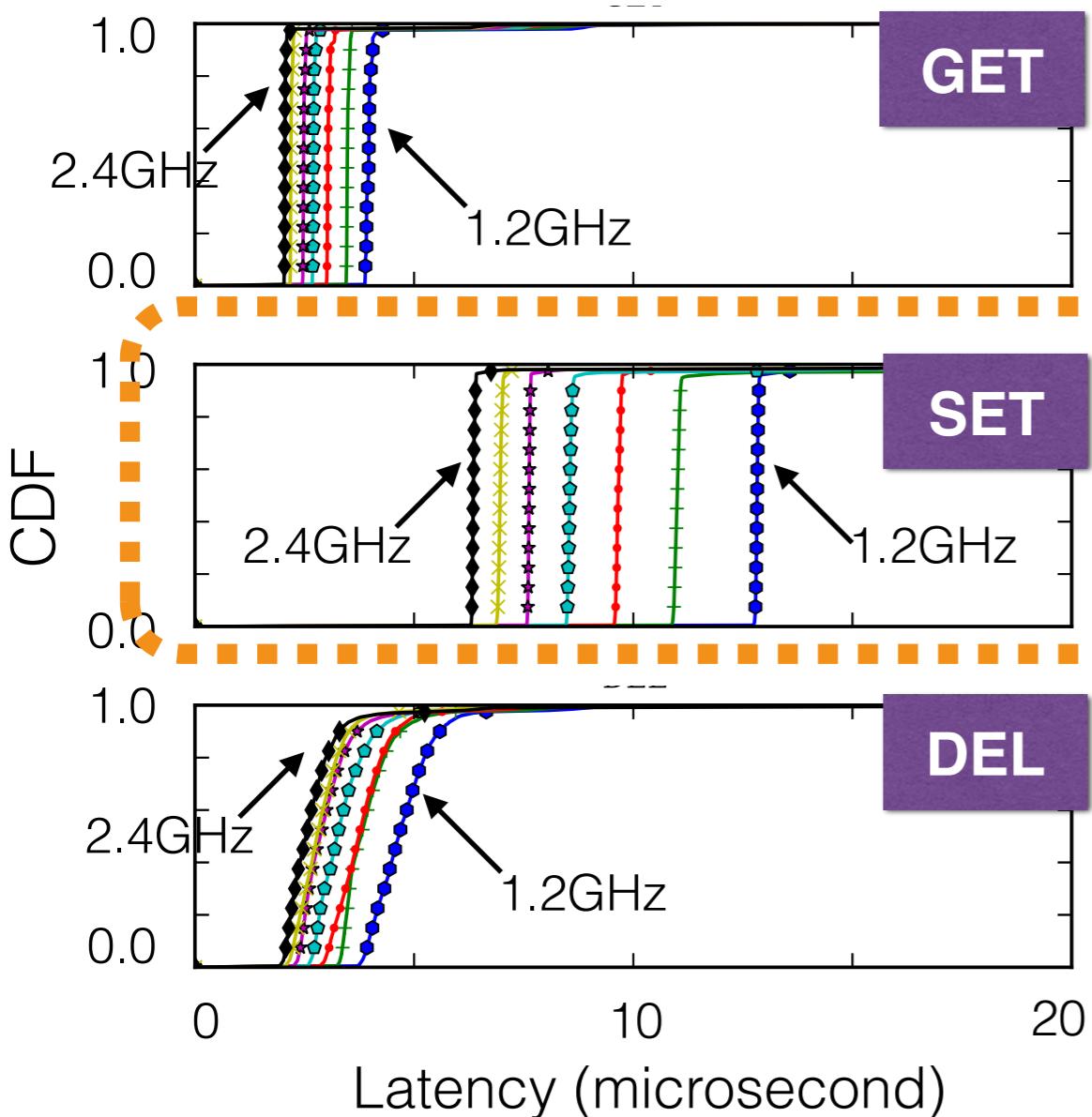


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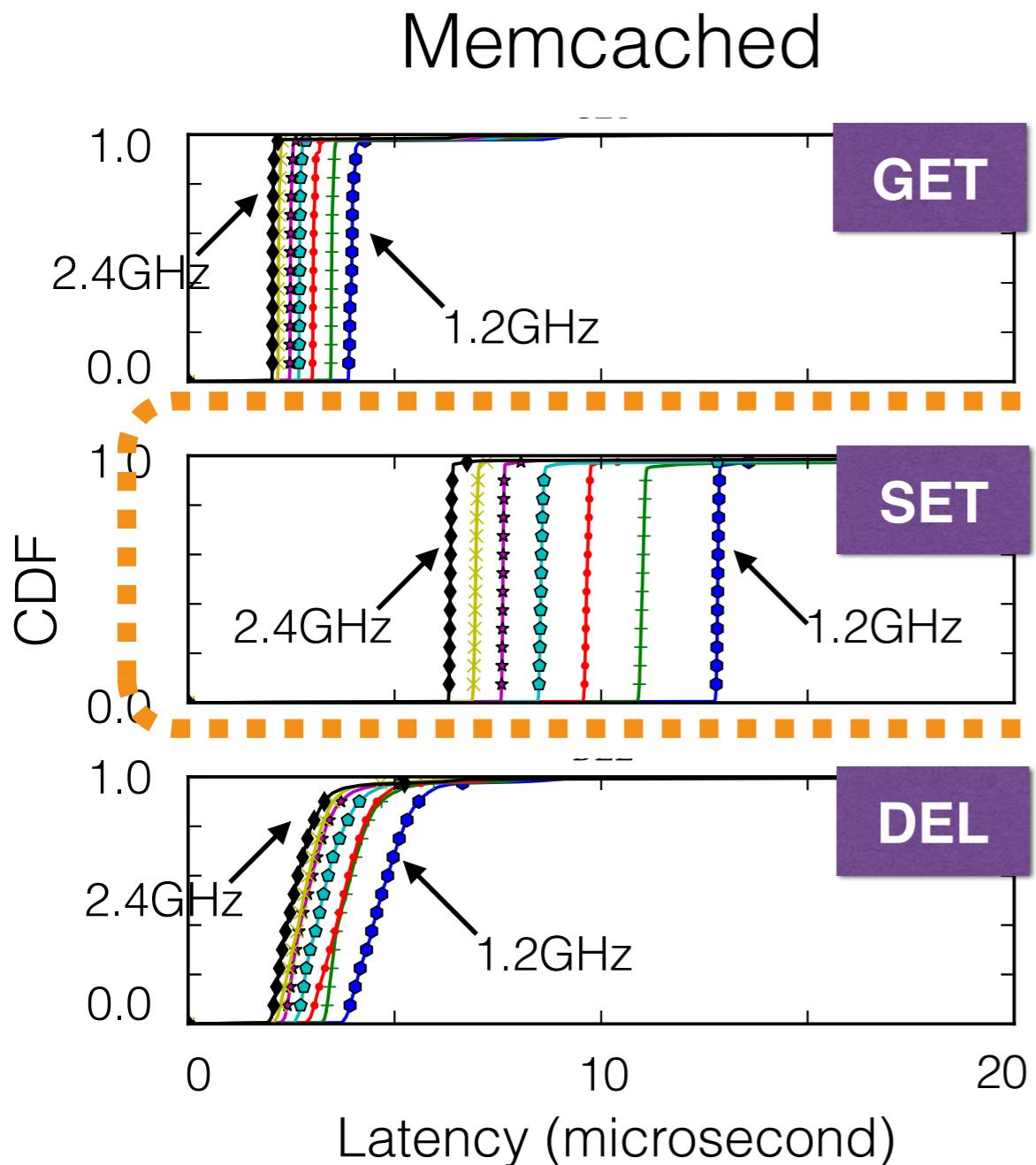
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## Observation

- Boosting the core from the lowest to the highest frequency improves SET's tail latency from **13 $\mu$ s to 7 $\mu$ s**

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**SET requests are  
good candidates  
for boosting**

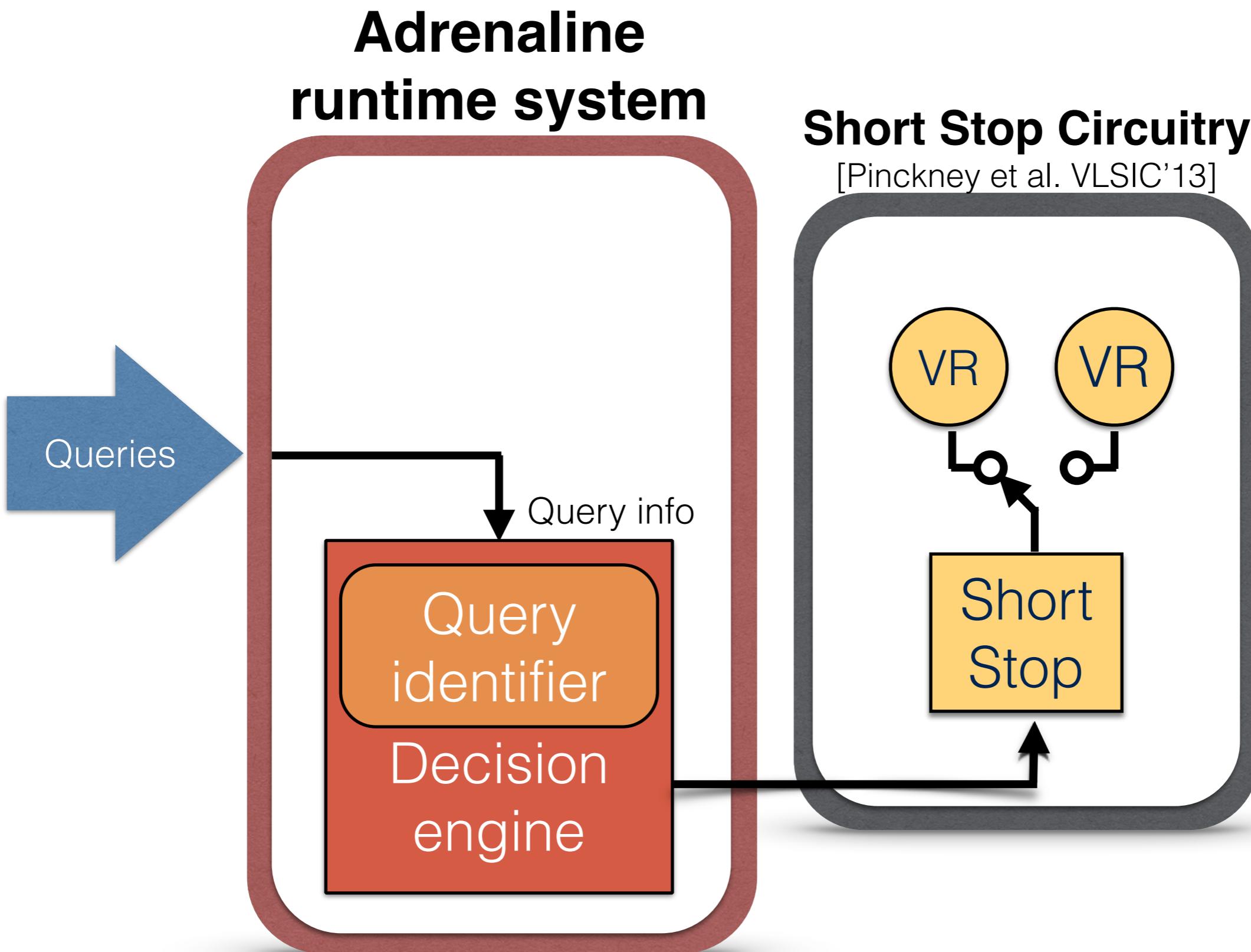
## **Query-level boosting with Adrenaline**

- ✓ Pinpoint queries that are highly likely to contribute to the tail
- ✓ Quickly boost the core via ultra-fast switching circuitry

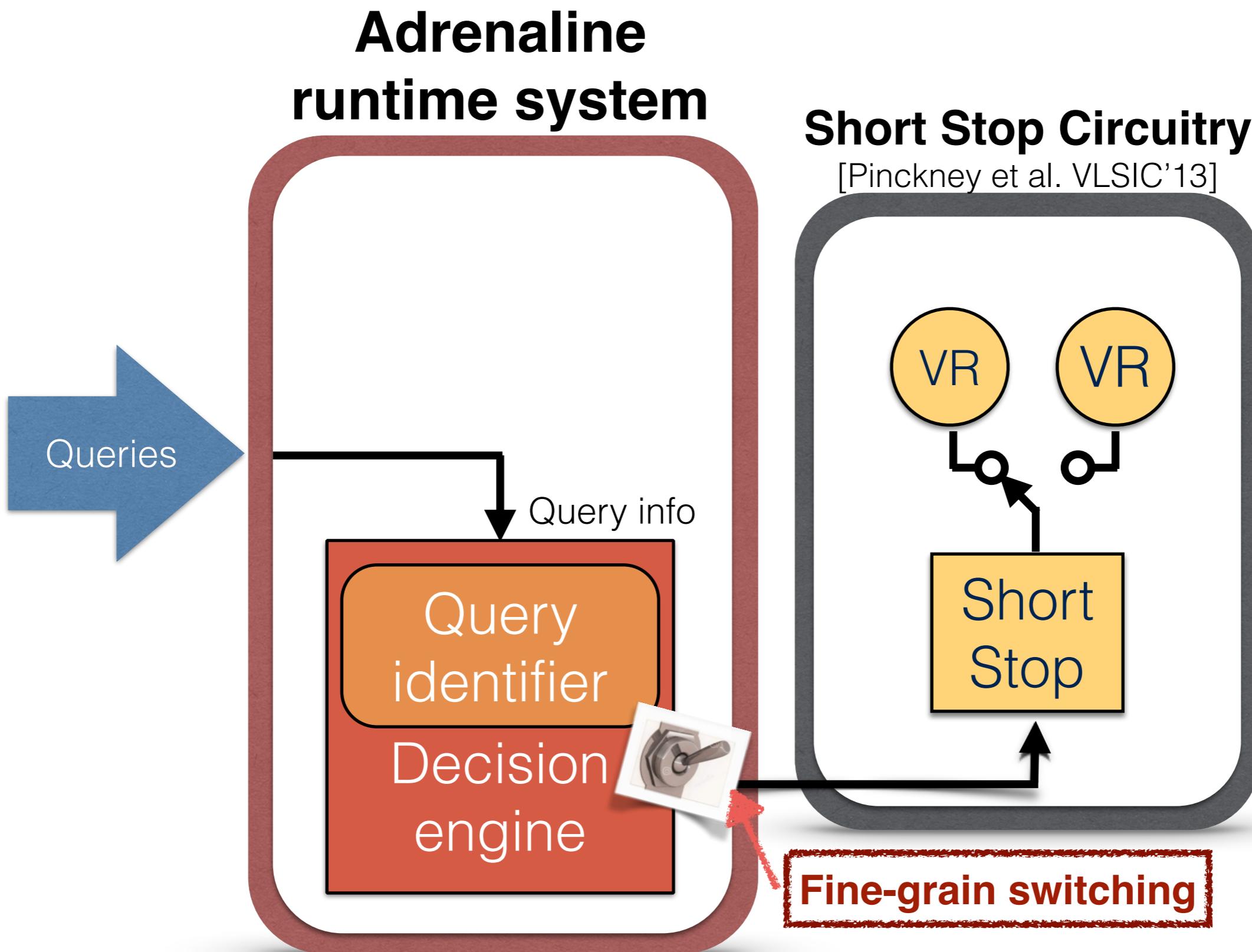
# Adrenaline Design

| Adrenaline |

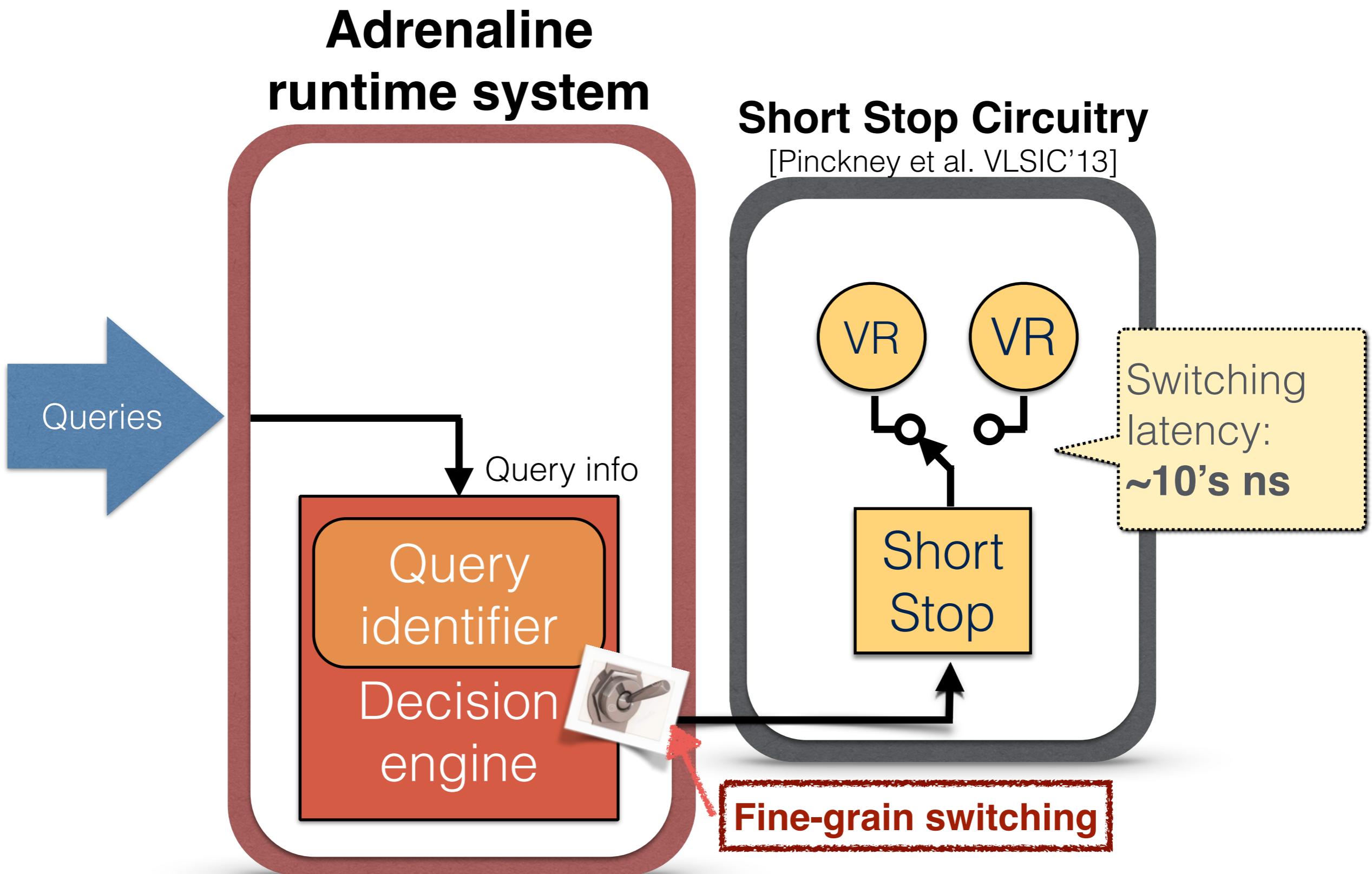
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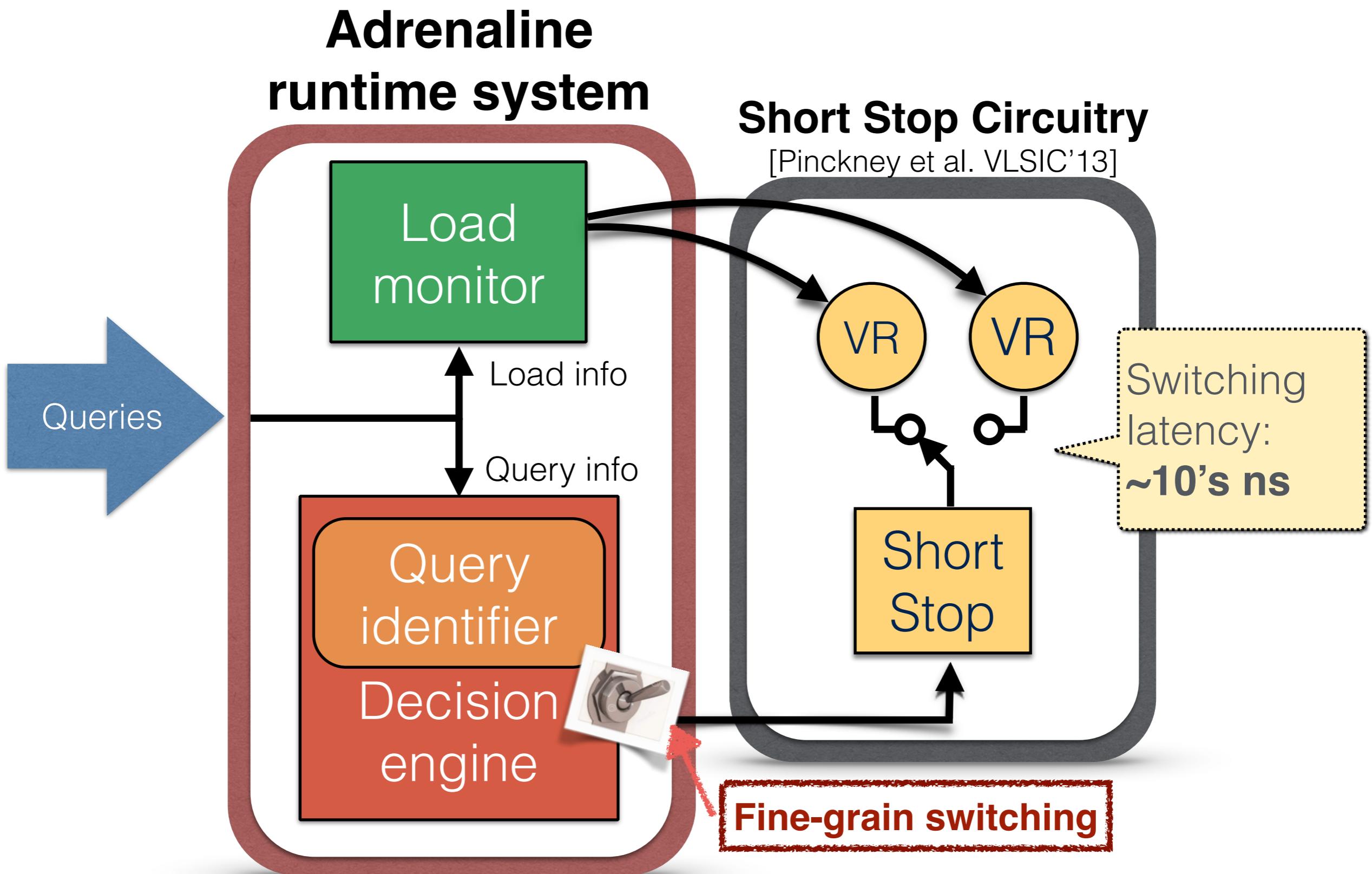
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# Boosting based on the query-level indicators

Query  
identifier

## Rapidly identifying query types

- Candidate vs. non-candidate type
- Needs to be simple to achieve low overhead

Decision  
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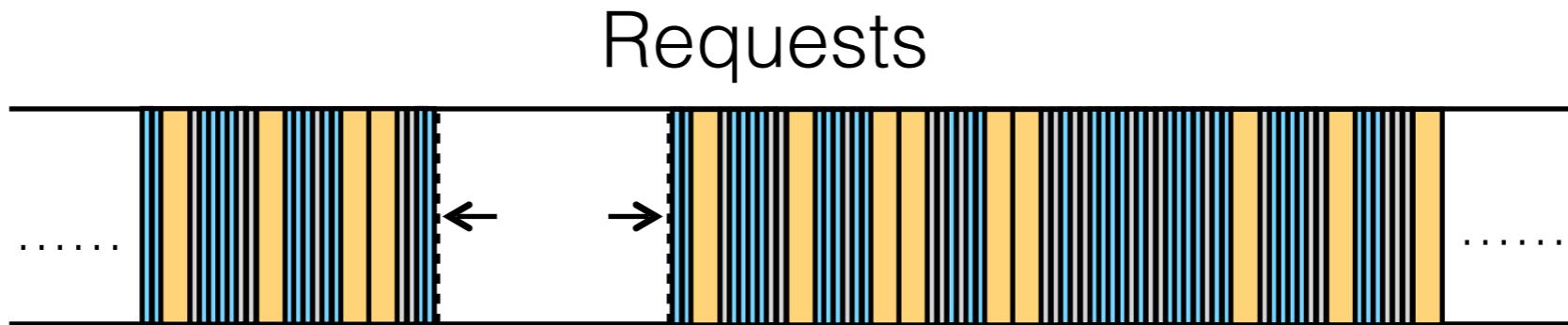
### **Candidate** type:

- Boost this query as soon as possible

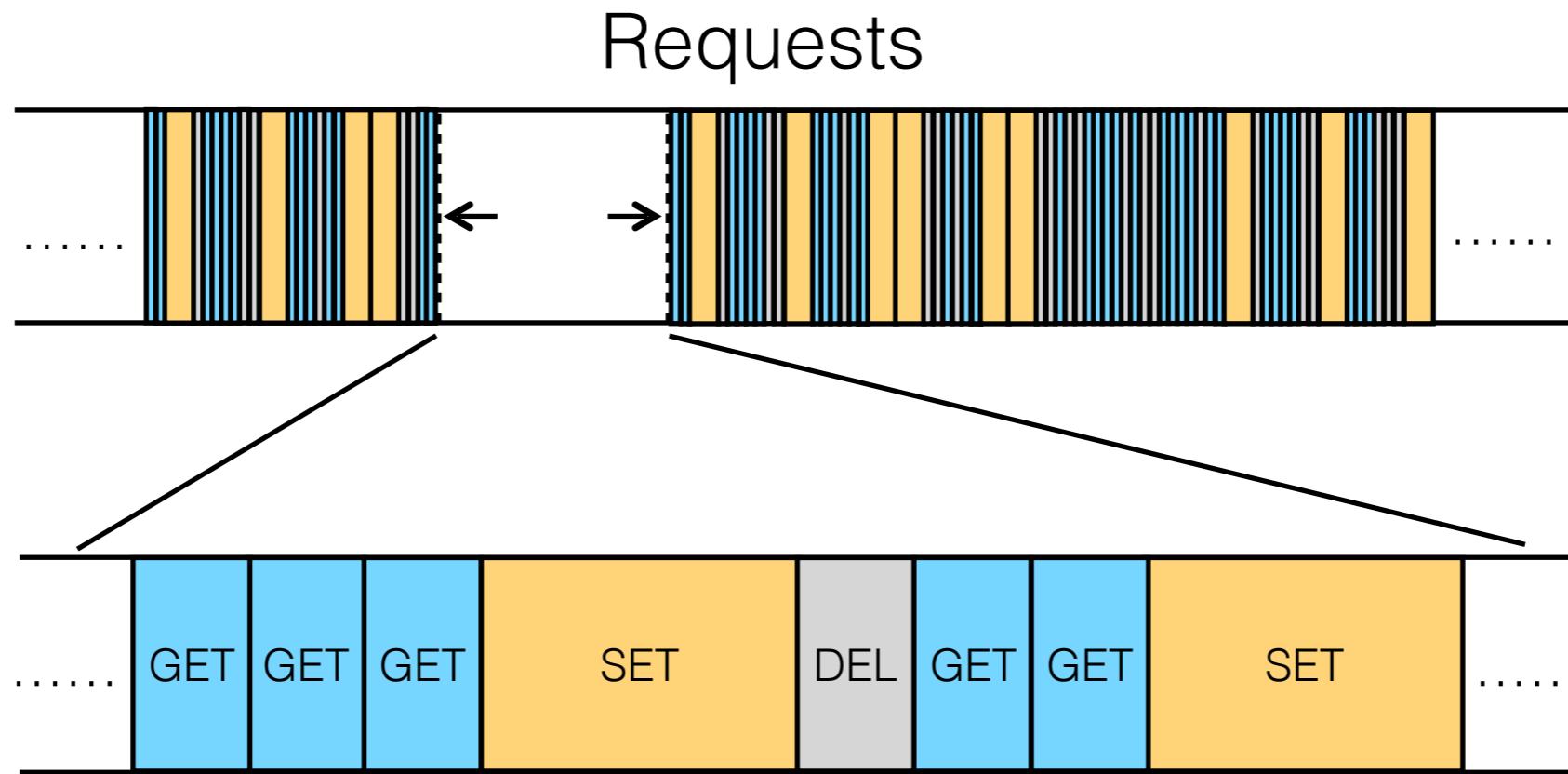
### **Non-candidate** type:

- Boost when query runtime exceeds half of QoS target

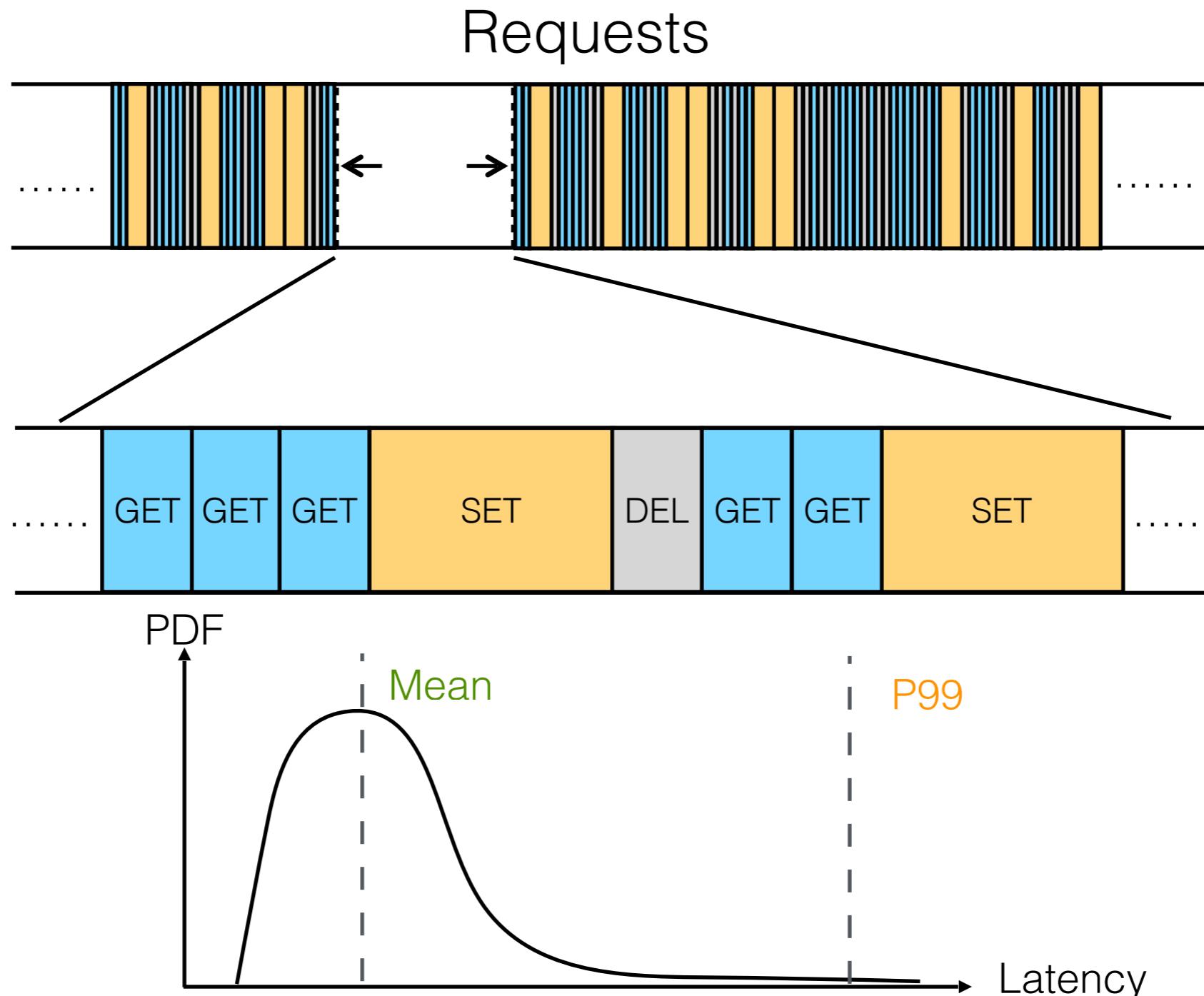
# Adrenaline: a closer look



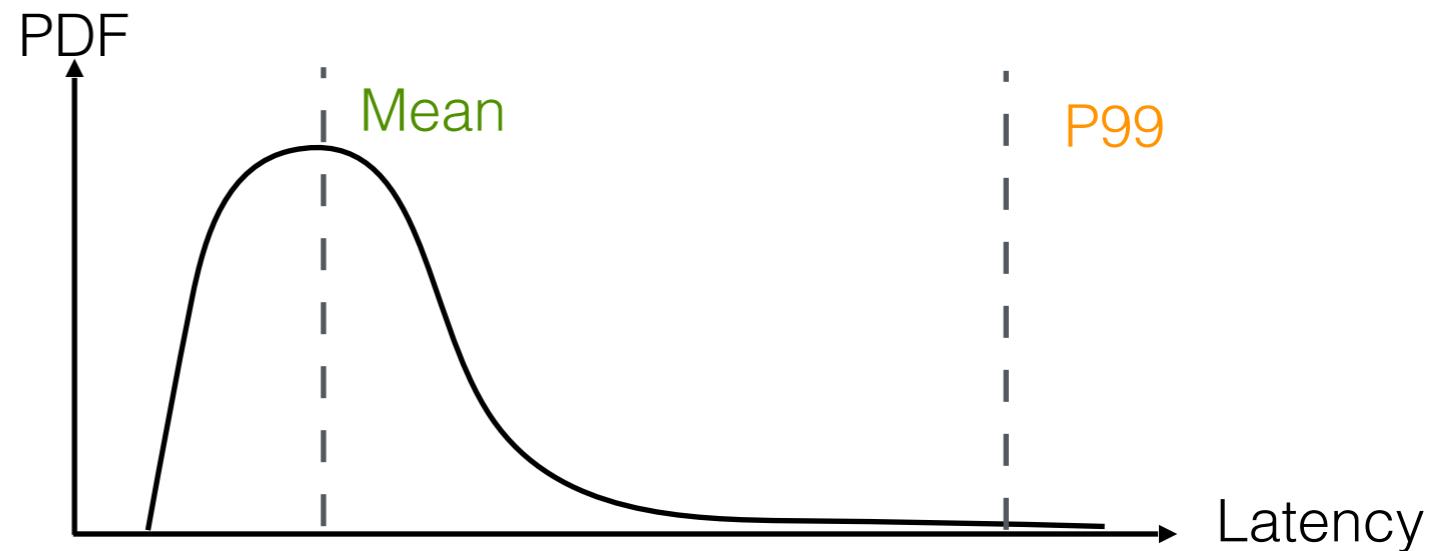
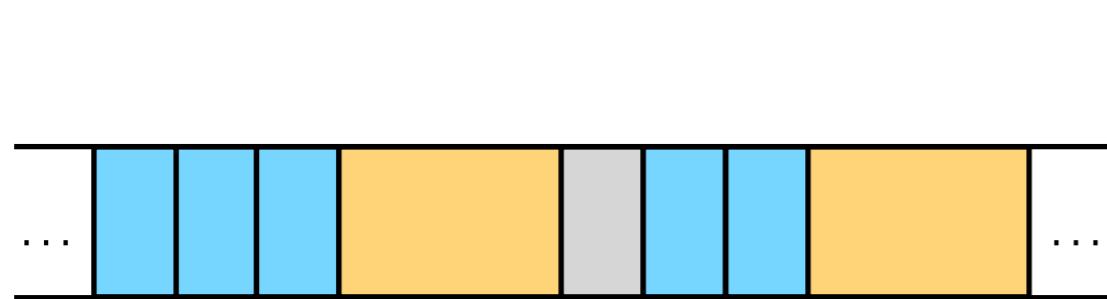
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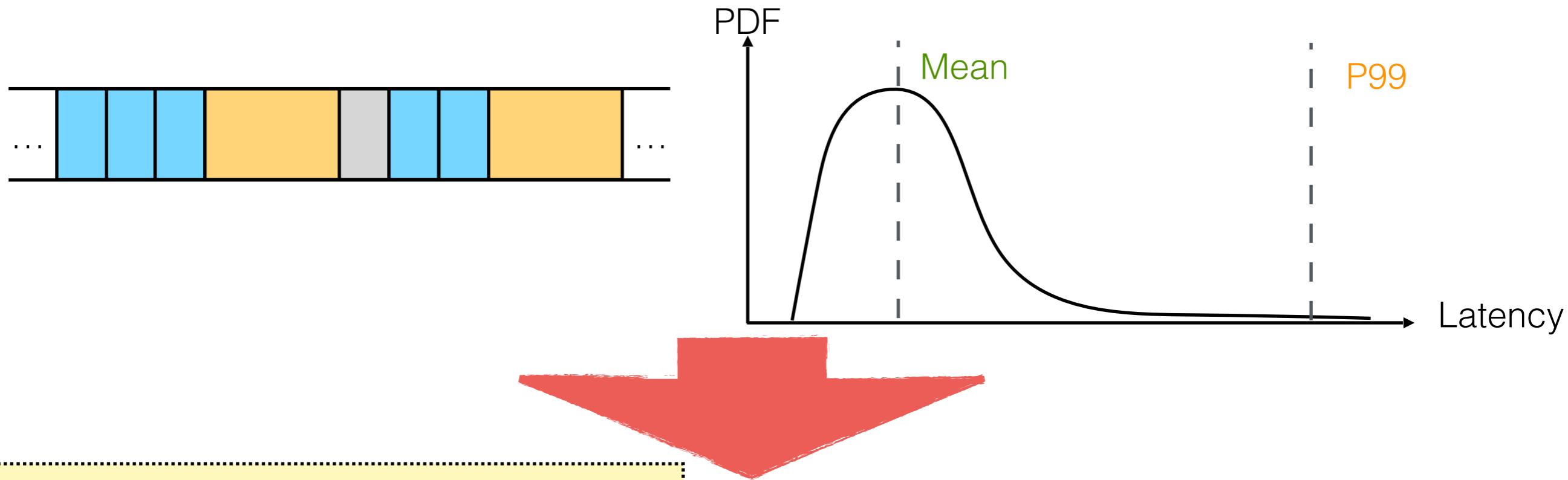
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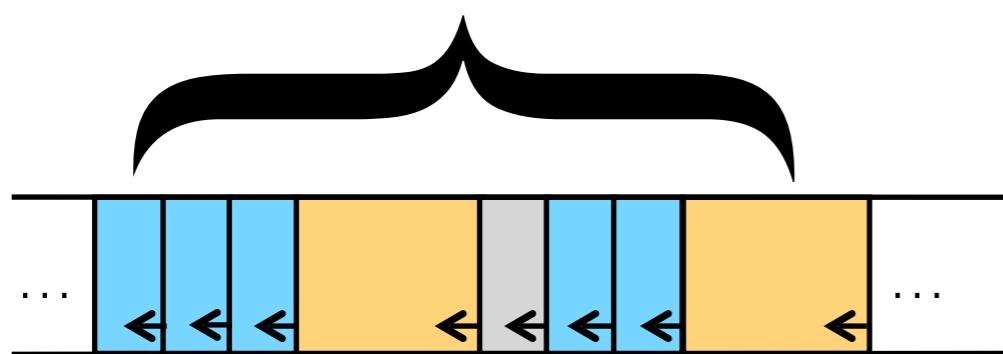
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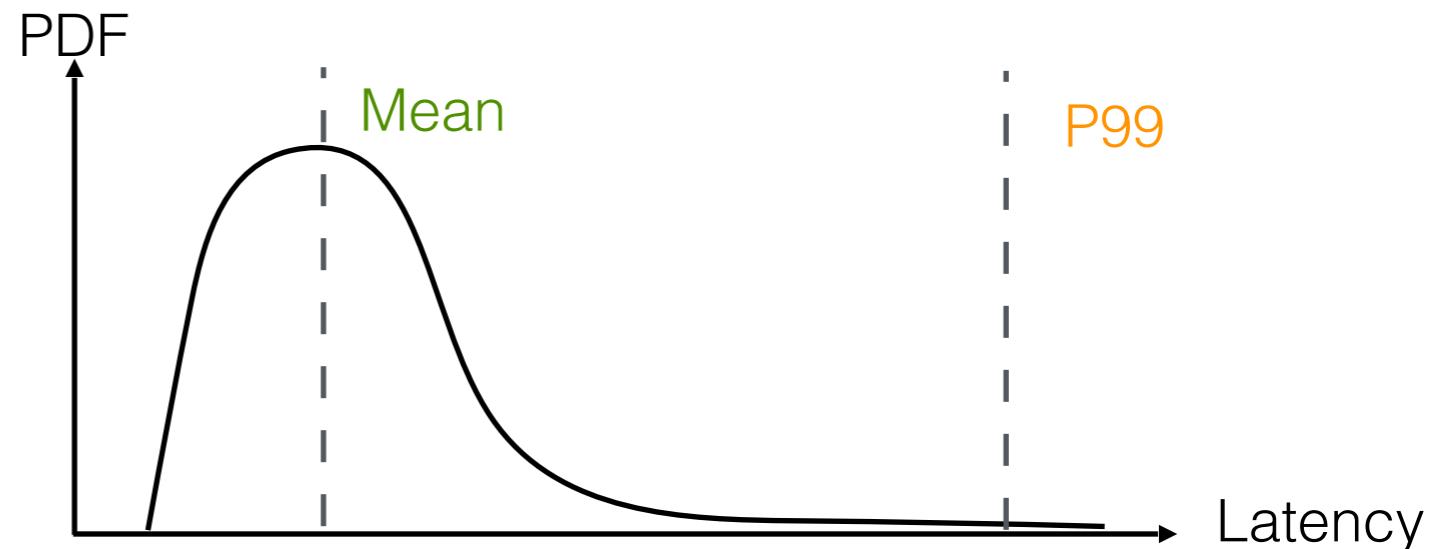
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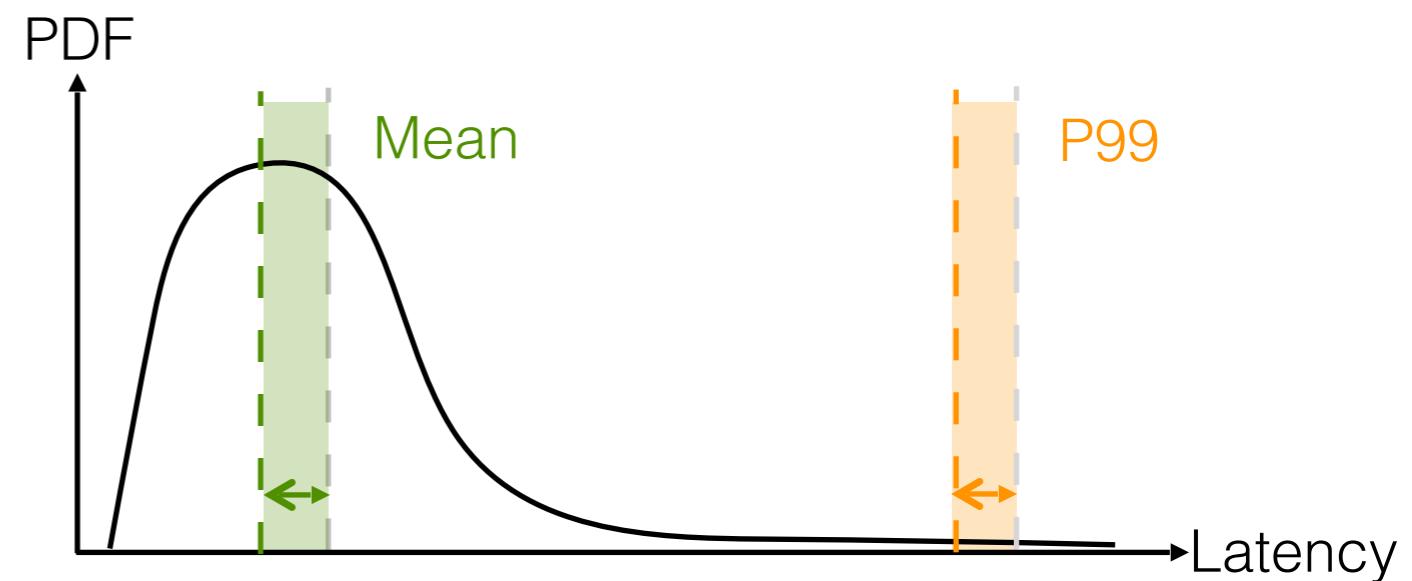
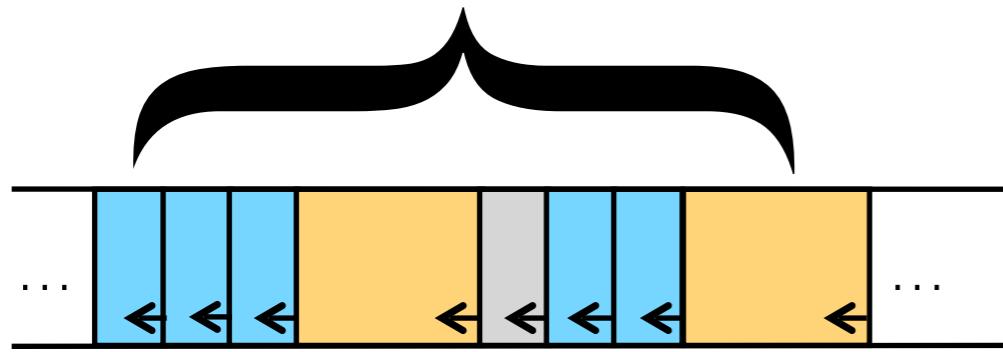
Coarse-grain solutions boost every query in that time interval



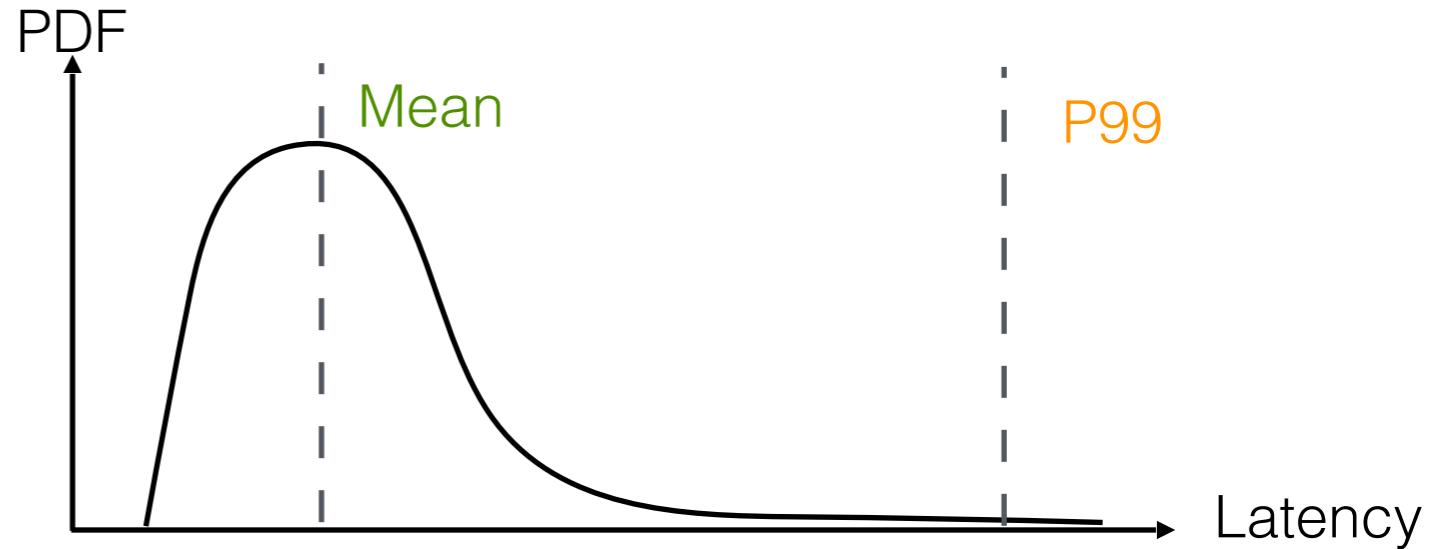
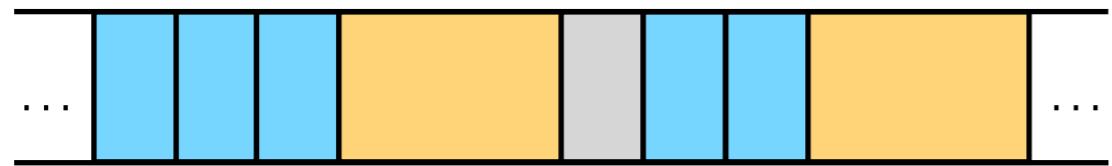
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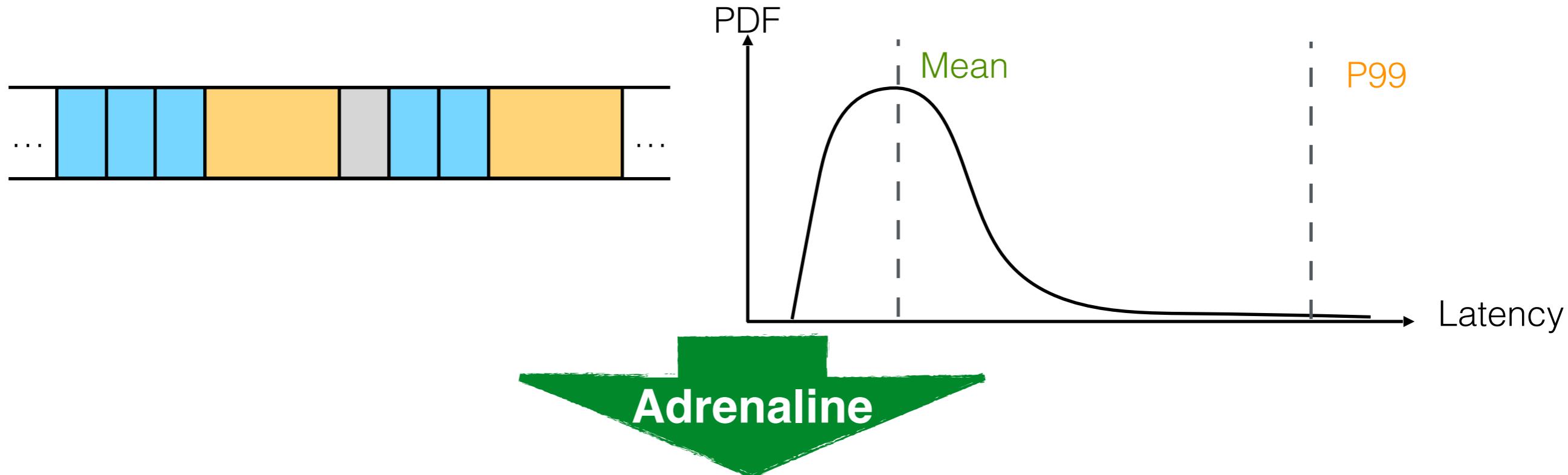
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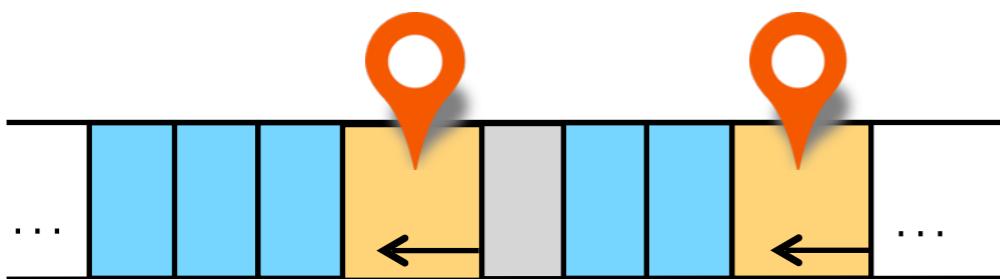
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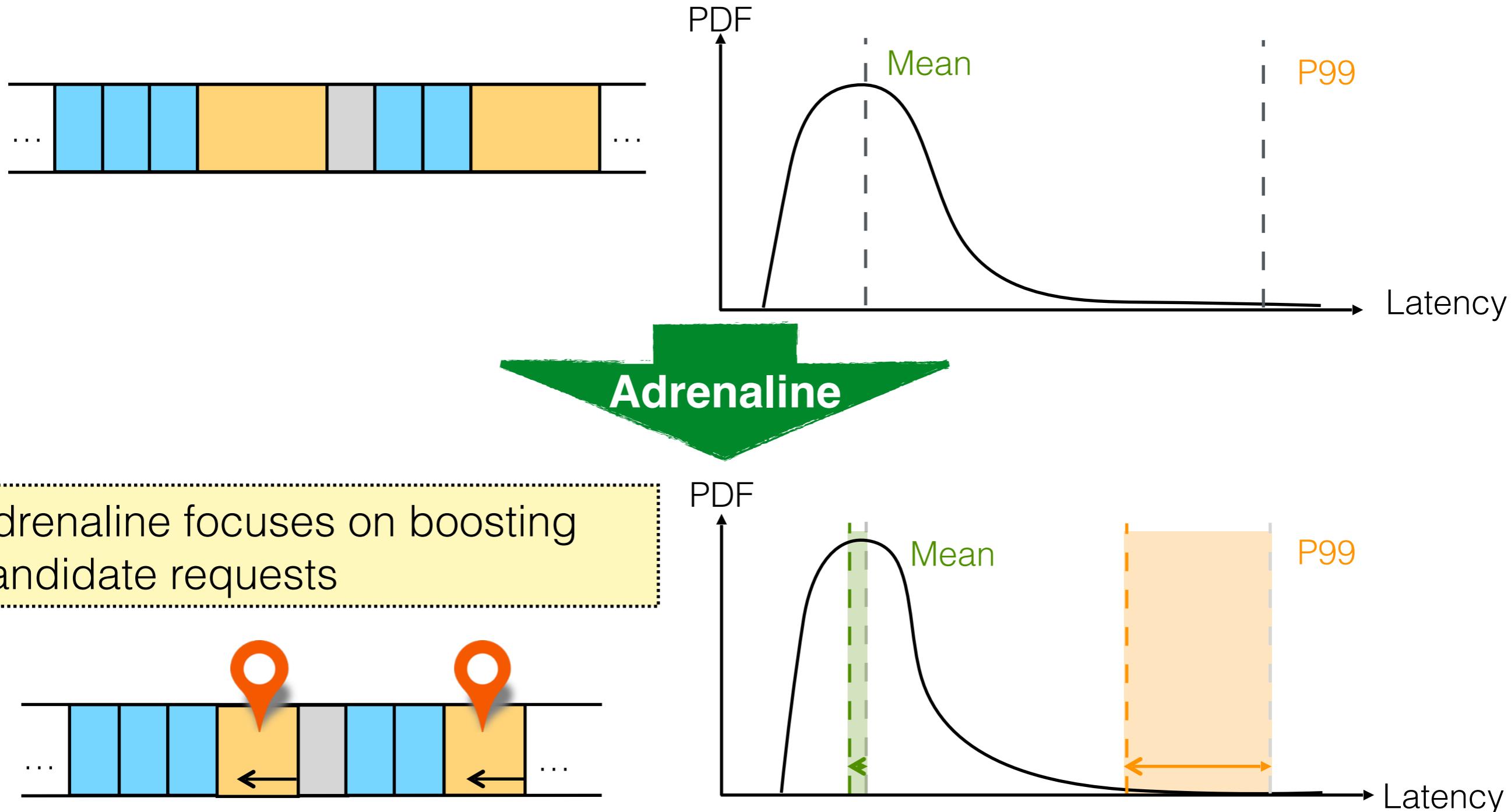
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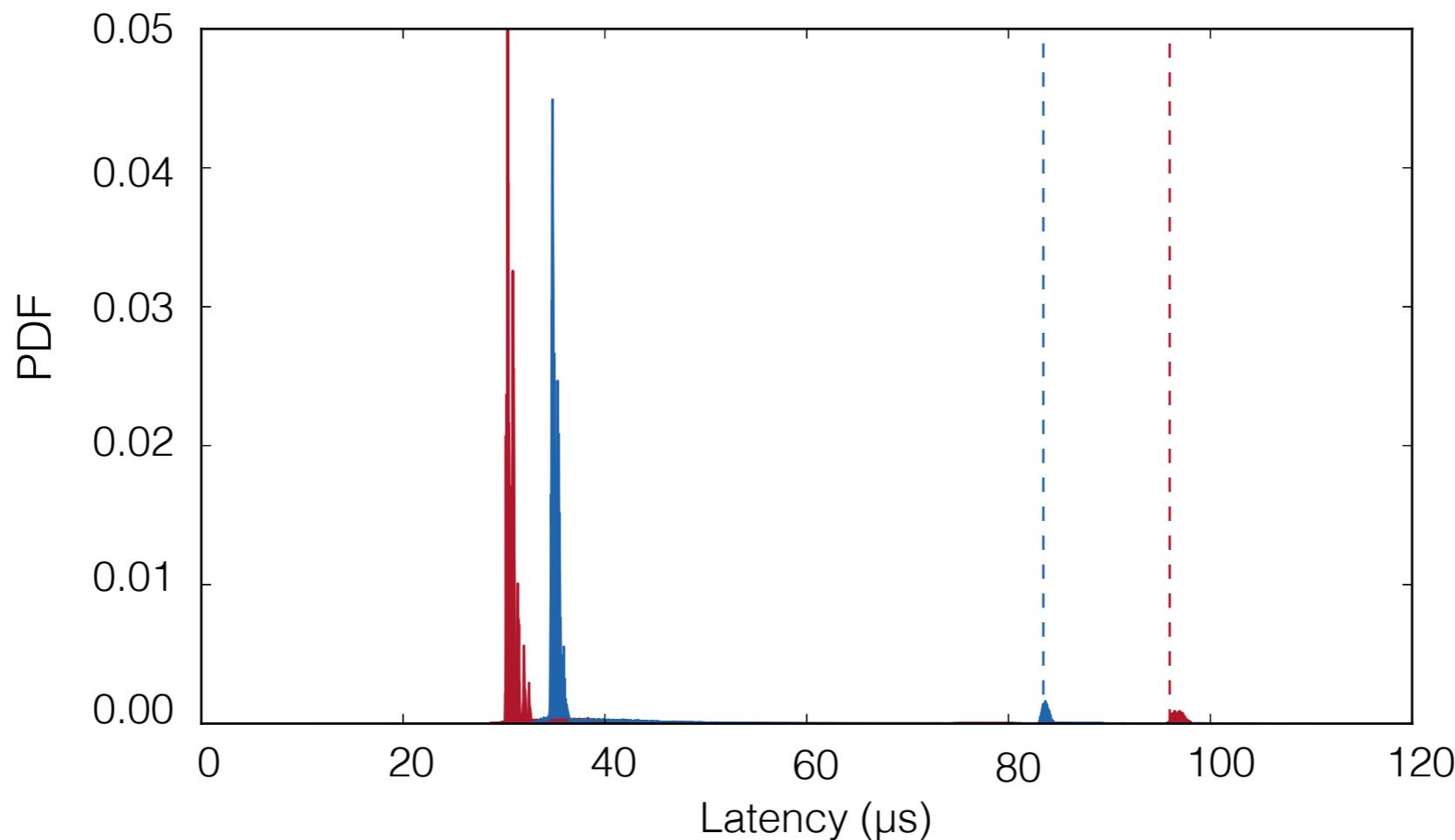
Adrenaline focuses on boosting candidate requests



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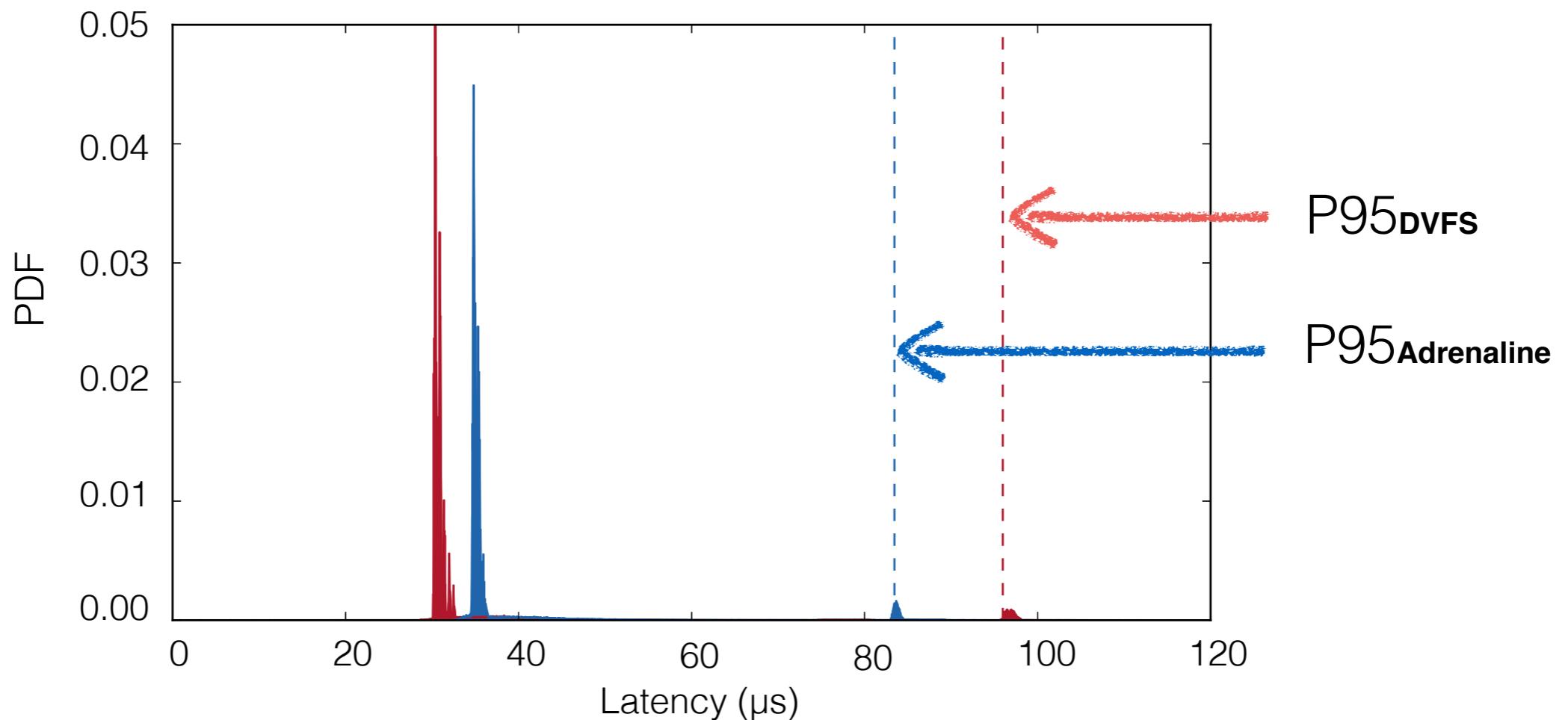


# Measured latency distributions



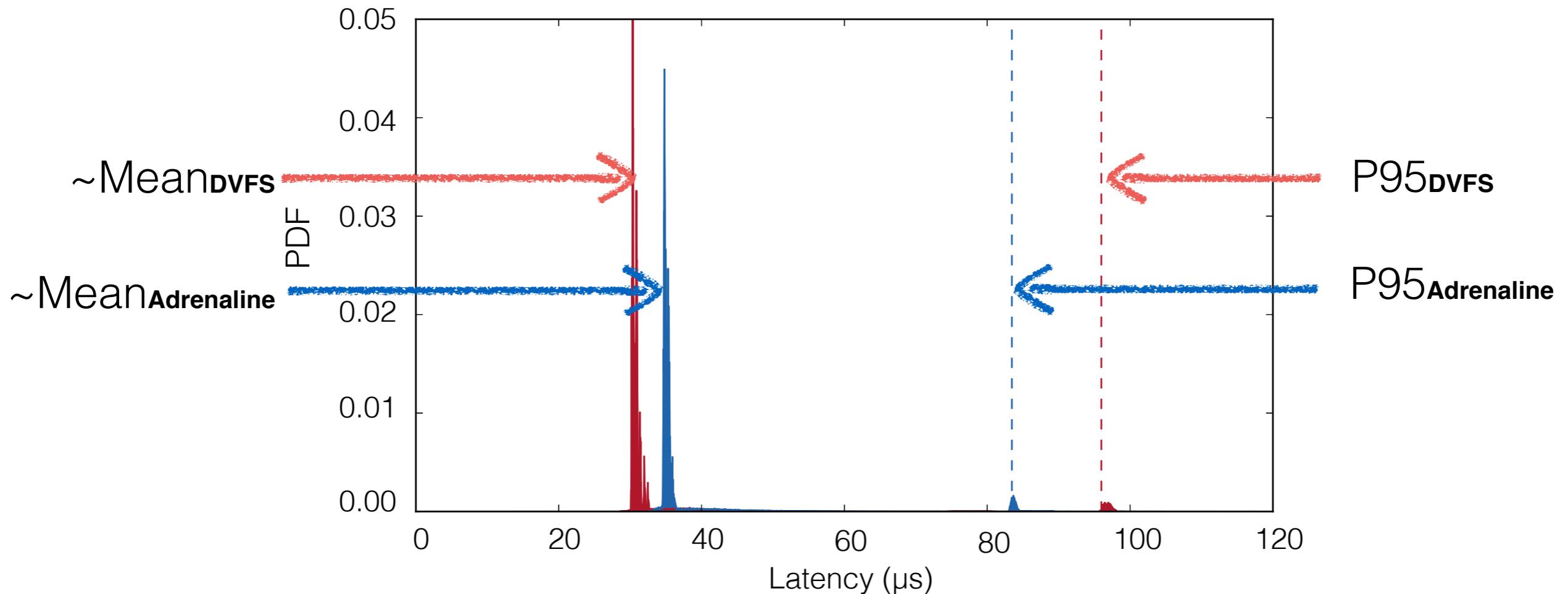
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- Compared to coarse-grain DVFS, Adrenaline...
  - ✓ reduces the tail latency significantly, achieving superior QoS
  - ✓ trades the mean latency for energy saving

# Evaluation

| Adrenaline |

# Evaluation methodology

- Measure request latency distributions on a real system
  - Intel Ivy Bridge + 136 GB RAM
  - Analyze Adrenaline & coarse-grain DVFS in BigHouse  
[Meisner et al. ISPASS'12]

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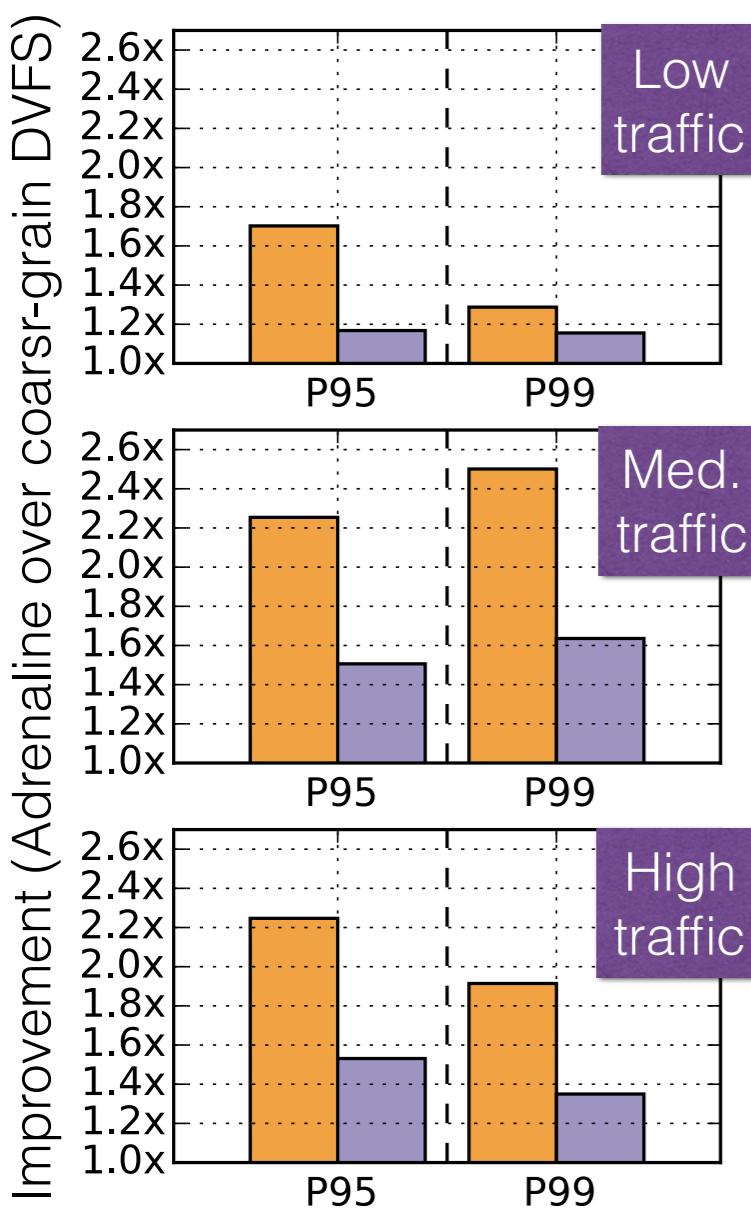
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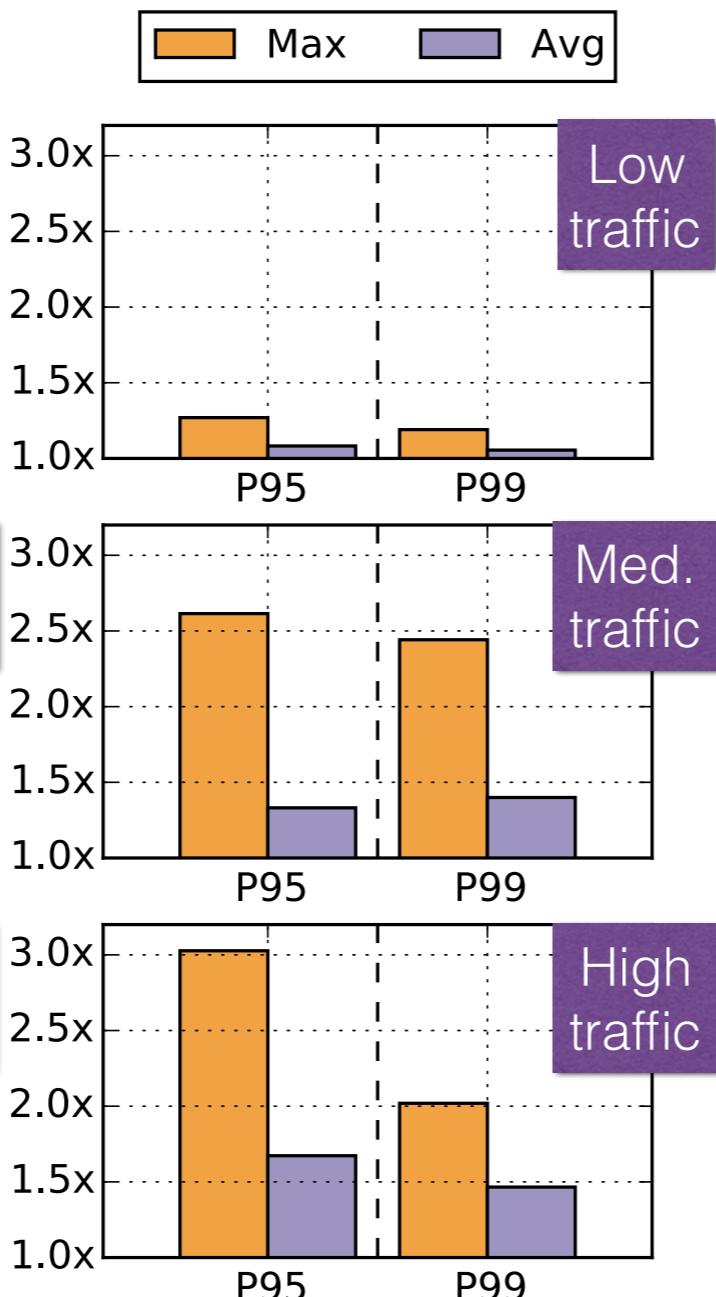
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- We generate various workload compositions using Facebook's published result [Atikoglu et al. SIGMETRIC'12]

# Rein in the tail with Adrenaline

Memcached



Web Search

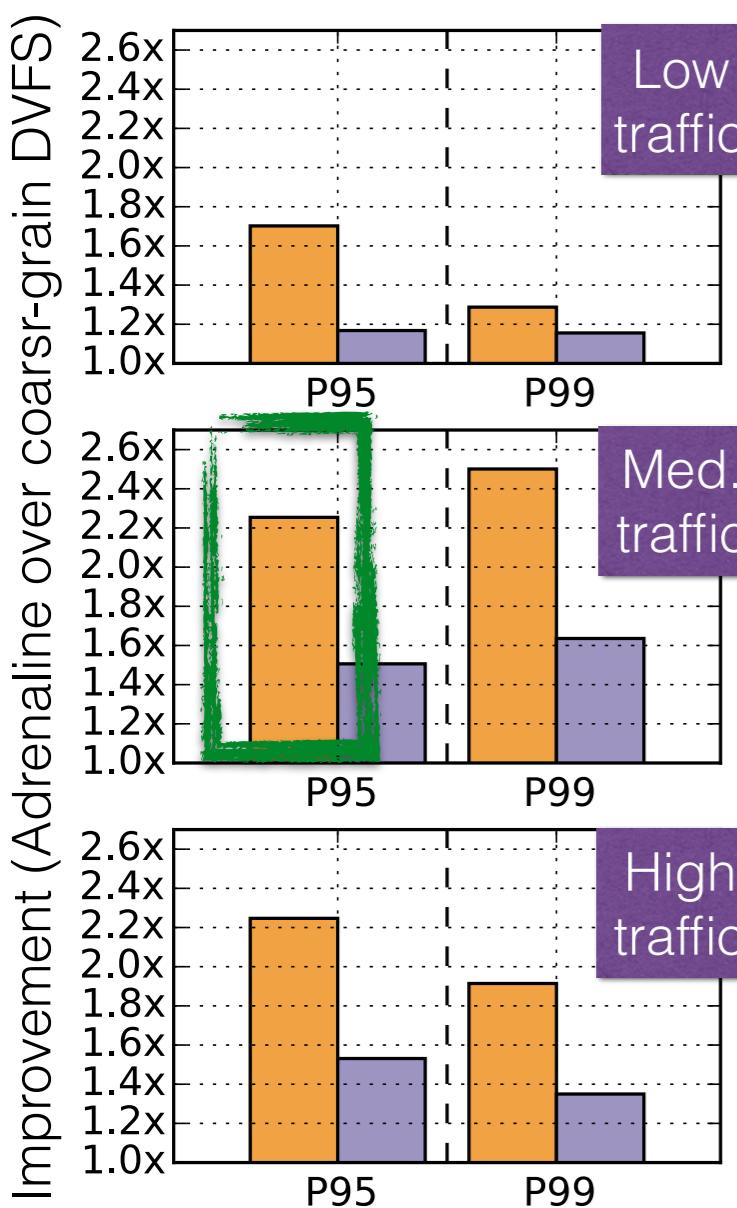


## Optimizing for tail latency:

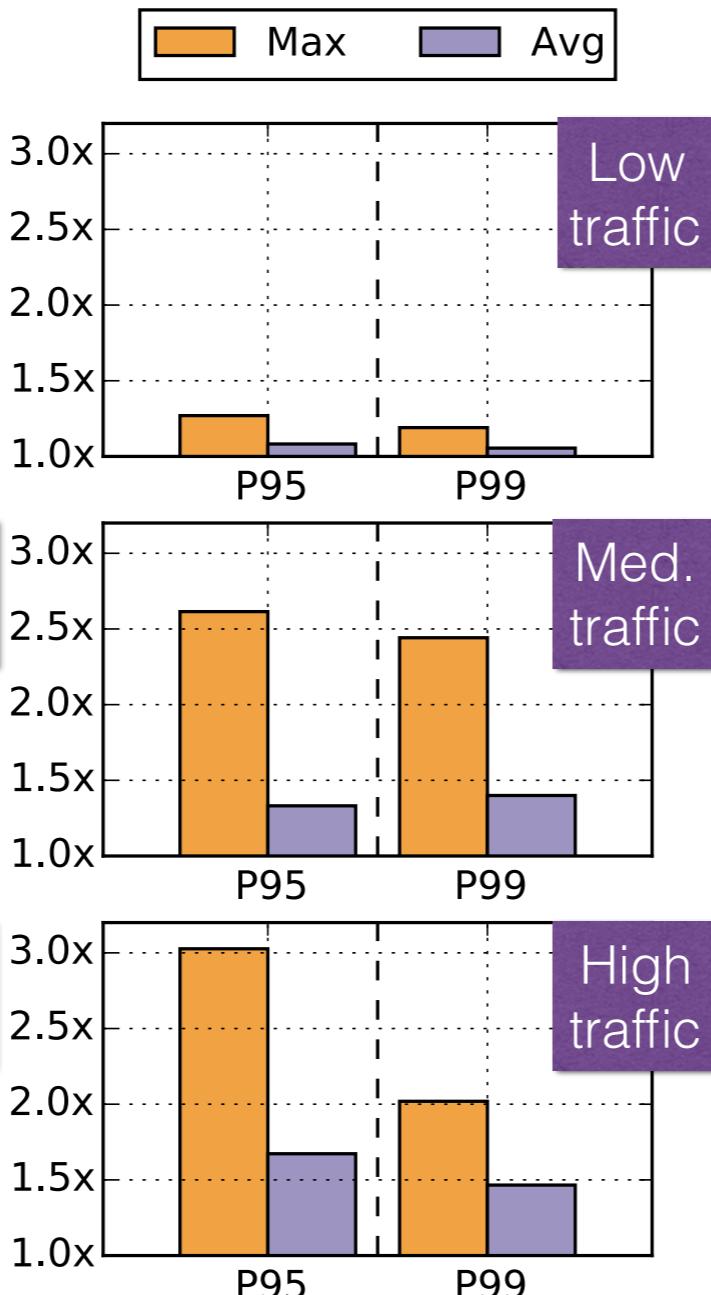
- Same energy budget
- Memcached:
  - Max: up to 2.5x
  - Avg: up to 1.6x
- Websearch
  - Max: up to 3x
  - Avg: up to 1.7x

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Memcached



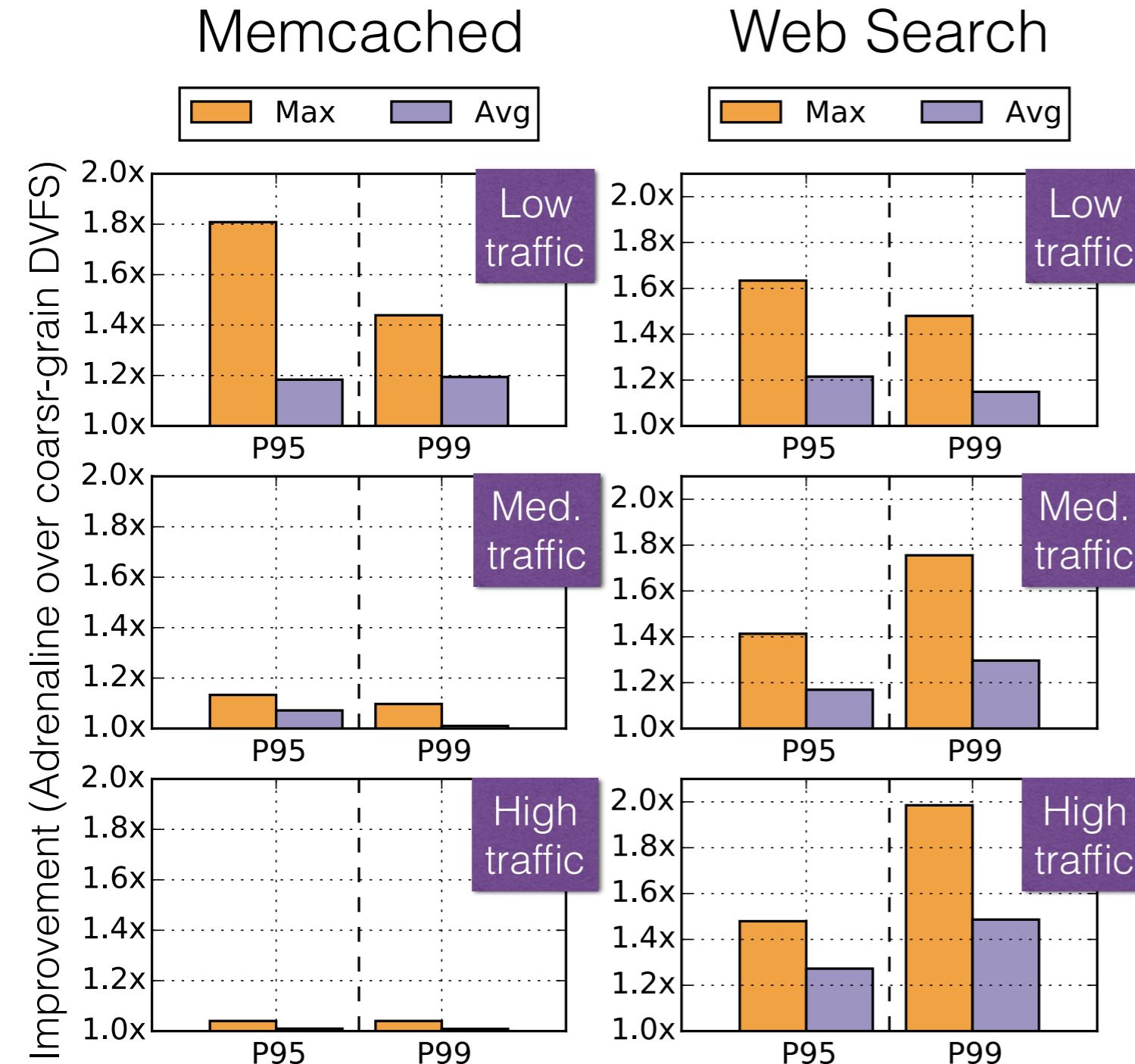
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# When we want low energy



## Optimizing for energy

- Same tail latency target
- Memcached:
  - Max: up to 1.8x
  - Avg: up to 1.2x
- Websearch
  - Max: up to 2x
  - Avg: up to 1.5x

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- User-facing services requires responsiveness and energy efficiency
- Pinpointing and boosting for tail queries gives the best of both worlds
- Adrenaline outperforms coarse-grain DVFS
  - ✓ Up to 3x tail latency improvement
  - ✓ Up to 2x energy saving improvement

