

**GCXX**

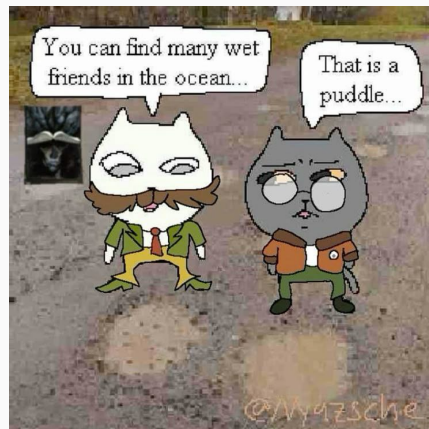
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## GC: what is needed?

- **Allocator**
- **Mark**
- **Sweep**

## The whole API:

```
std::map<void*, size_t> allocated;  
void* heap;  
static constexpr size_t MIN_SIZE =  
    sizeof(void*) * 2;  
  
[[nodiscard]] void* allocate(size_t size);  
void deallocate(void* p);
```



Find all reachable objects.

Search begins with GC roots, i.e.

- Globals
- Registers
- Stack



— [Logic – Challenging 12] I know how to find the roots.



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**LOGIC** [Challenging: Failure] — You didn't think that scanning the globals and the stack would be so easy, right?

## Globals and stack

Turns out, finding the location of static memory is quite hard. There's no easy way to do it, only parsing the current ELF file and asking all of the loaded dynamic libraries.





**AUTHORITY** [Trivial: Success] — Who cares about globals after all?  
Check only the registers and the stack.





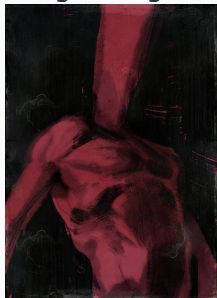
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**LOGIC** [Easy: Success] — In GC constructor save the current `rsp` value  
as the stack base and then iterate over the whole stack, finding the  
roots.

## Mark phase

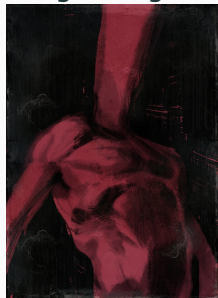
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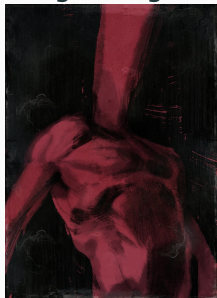
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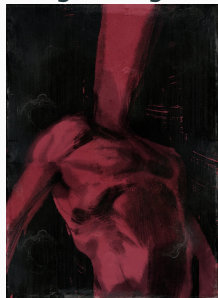


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**PHYSICAL INSTRUMENT** [Heroic: Success] — I'll say it again. The %!@& do you need memory for? Write a bitset for the whole heap.

## Sweep phase

Three lines:

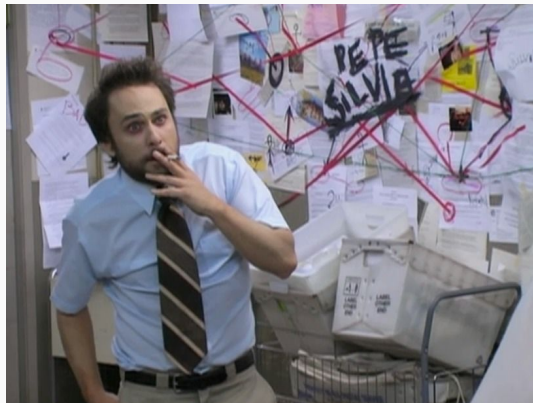
```
for (auto&& [ptr, _] : allocator.allocated) {  
    if (!reachable[...]) {  
        allocator.deallocate(ptr);  
    }  
}
```

Just find all unreachable blocks and free them.

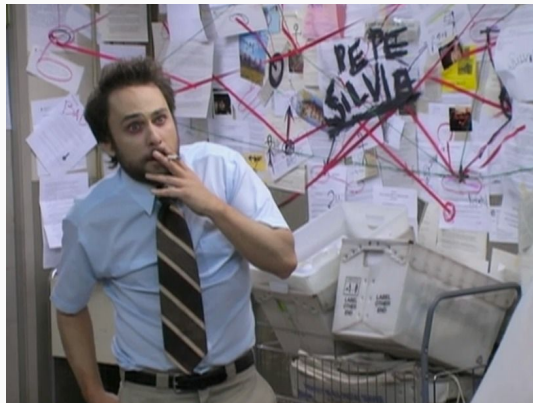
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That led to some funny cosequences while testing: 'std::println' uses 'std::format', which, in turn, allocates a string, and its constructor is using 'new'.  
But we used 'println' for logging in our allocator. We got to the point where doing literally anything with dynamic memory will cause a stack overflow.



# IT WORKS

