

| Algorithmics | Student information | Date | Number of session |
|--------------|------------------------|------|-------------------|
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Activity 1. Bubble Algorithm

| n | T ordered | T reverse | T random |
|------------|-----------|-----------|----------|
| 10000 | 331 | 1532 | 1028 |
| 2*10000 | 1306 | 5988 | 4129 |
| 2**2*10000 | 5251 | 24914 | 17469 |
| 2**3*10000 | 21250 | OoT | OoT |
| 2**4*10000 | OoT | OoT | OoT |

Bubble Sort has a worst-case and average-case time complexity of $O(n^2)$, which is evident in the exponential growth of execution times as the input size doubles. For **ordered arrays**, Bubble Sort performs slightly better ($O(n)$) because it can detect early that no swaps are needed. However, for reverse-ordered arrays, it performs the worst, requiring the maximum number of swaps and comparisons. In random arrays, the performance lies between the two extremes, as the number of swaps depends on the initial arrangement.

Activity 2. Selection algorithm

| n | T ordered | T reverse | T random |
|------------|-----------|-----------|----------|
| 10000 | 306 | 325 | 327 |
| 2*10000 | 1222 | 1281 | 1292 |
| 2**2*10000 | 5106 | 4903 | 5202 |
| 2**3*10000 | 20784 | 20954 | 20318 |
| 2**4*10000 | OoT | OoT | OoT |

Selection sort has a consistent time complexity of $O(n^2)$ across all cases, as evidenced by the times increasing quadratically with input size. For smaller inputs, the algorithm completes in a reasonable time, but as the input size grows, it results in "Out of Time"

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(OoT) errors, highlighting its inefficiency for large numbers. The ordered, reversed, and random cases perform similarly, as it always performs the same number of comparisons and swaps.

Activity 3. Insertion algorithm

| n | Tordered | Treversed | Trandom |
|-------------|----------|-----------|---------|
| 10000 | LoR | 151 | 149 |
| 2*10000 | LoR | 573 | 569 |
| 2**2*10000 | LoR | 2331 | 2315 |
| 2**3*10000 | LoR | 9766 | 9281 |
| 2**4*10000 | LoR | 37825 | 36957 |
| 2**5*10000 | LoR | OoT | OoT |
| 2**6*10000 | LoR | OoT | OoT |
| 2**7*10000 | LoR | OoT | OoT |
| 2**8*10000 | LoR | OoT | OoT |
| 2**9*10000 | 91 | OoT | OoT |
| 2**10*10000 | 184 | OoT | OoT |
| 2**11*10000 | 361 | OoT | OoT |
| 2**12*10000 | 722 | OoT | OoT |
| 2**13*10000 | 1477 | OoT | OoT |

For smaller input sizes, the algorithm performs efficiently in all cases, but as the input size grows exponentially, the ordered and reversed cases show a sharp increase in time complexity, eventually leading to "Out of Time" errors for larger inputs. The algorithm struggles with nearly sorted or reverse-sorted data, likely due to its quadratic time complexity $O(n^2)$ in such cases. Interestingly, the random case performs significantly better, completing in constant time for larger inputs, which implies the algorithm handles randomness more efficiently, possibly due to fewer required comparisons or swaps. This

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highlights the algorithm's sensitivity to input order and its inefficiency for worst-case scenarios.

Activity 4. Quicksort algorithm

| n | Tordered | Treversed | Trandom |
|------------|----------|-----------|---------|
| 250000 | LoR | 97 | 94 |
| 2*25000 | 64 | 190 | 190 |
| 2**2*25000 | 123 | 404 | 402 |
| 2**3*25000 | 256 | 869 | 870 |
| 2**4*25000 | 525 | 1879 | 1884 |
| 2**5*25000 | 1084 | 4233 | 4255 |
| 2**6*25000 | 2231 | 10327 | 103332 |

Quicksort performs efficiently in the random case, as expected, with times increasing linearly or near-linearly, reflecting its average-case complexity of $O(n \log n)$. However, in the ordered and reversed cases, the times grow significantly faster, particularly for larger inputs, indicating a degradation to its worst-case complexity of $O(n^2)$. This occurs because quicksort's performance heavily depends on the choice of pivot; in ordered or reversed data, poor pivot selection leads to unbalanced partitions.

Activity 5. QuicksortInsertion algorithm

| n | T random |
|---------------------------|----------|
| QuickSort | |
| QuickSort+Insertion(k=5) | 11699 |
| QuickSort+Insertion(k=10) | 11215 |
| QuickSort+Insertion(k=20) | 11685 |
| QuickSort+Insertion(k=30) | 10521 |

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| | |
|-----------------------------|-------|
| QuickSort+Insertion(k=50) | 12031 |
| QuickSort+Insertion(k=100) | 9957 |
| QuickSort+Insertion(k=200) | 8314 |
| QuickSort+Insertion(k=500) | 12379 |
| QuickSort+Insertion(k=1000) | 22440 |

The results obtained for the combined QuickSort and Insertion Sort algorithm on a random vector of size 16,000,000 do not align with the expected behavior. Ideally, for small values of k (e.g., $k = 5$ or $k = 10$), the algorithm should perform efficiently, as Insertion Sort is well-suited for small subarrays. However, the execution times show inconsistencies, with some larger k values (e.g., $k = 100$ or $k = 200$) performing better than smaller ones, which is unexpected. This suggests a potential flaw in the implementation, such as incorrect pivot selection in QuickSort or improper handling of subarray sizes.