



# Computer Organization

COMP2120

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## Number Representation Part I



## Number system in computer

- We understand decimal number, base-10
- Computer: Binary, not decimal  
Easy to be represented, stored, and transmitted.  
E.g. hole in a punched card, high or low voltage
- There are 10 types of people in this world, those who understand binary and those who don't
- NOT the same as in mathematics.  
E.g., the product of two positive integers may be negative  
Associative law

$$(3.14 + 1e20) - 1e20 \neq 3.14 + (1e20 - 1e20)$$



# Number system in computer

## Finite Precision Numbers

- the size of memory available for storing a number is fixed. e.g. 32-bit integers
- The number of possible representation is fixed, i.e.  $2^{32}$  for 32-bit
- Overflow (number too large to be represented) and Underflow (number too small to be represented)
- Negative numbers
- Fractions
- Irrational Number\*
- Complex Numbers\*



# Radix Number Systems

- Positional number system: a string of digits

$$(\dots a_3 a_2 a_1 a_0 . a_{-1} a_{-2} a_{-3} \dots)_r$$

- position  $i$ , an associated weight  $r^i$ ,  $r$  is radix, or base
- $0 \leq a_i < r$ . Decimal:  $r = 10$ ,  $a_i \in \{0, 1, \dots, 9\}$ ; Binary:  $r = 2$ ,  $a_i \in \{0, 1\}$
- Dot: radix point
- Value:

$$\begin{aligned} & \dots + a_3 r^3 + a_2 r^2 + a_1 r + a_0 r^0 + a_{-1} r^{-1} + a_{-2} r^{-2} + a_{-3} r^{-3} + \dots \\ &= \sum_i (a_i \times r^i) \end{aligned}$$



# Radix Number Systems

- The formula applies to all radices, binary (base 2), octal (base 8), hexadecimal (base 16)
- Hexadecimal: 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, A(10), B(11), C(12), D(13), E(14), F(15)
- 0x or 0X also mean Hexadecimal  
 $FA1D3_{16} = 0xFA1D3$
- What is 10 (hexadecimal) ?
- The conversion between octal/hexadecimal and binary is very easy.
- Octal and Hexadecimal representation is to facilitate human only. Computer still use binary representation.



# Radix Number Systems

Example:

$$1984_{10} = 1 \times 10^3 + 9 \times 10^2 + 8 \times 10 + 4$$

$$1B84_{16} = 1 \times 16^3 + 11 \times 16^2 + 8 \times 16 + 4$$

Note that the principle behind decimal representation and other radix representation is the same.

We may use decimal numbers as analogy for explanation.



## Conversion between different radices-to radix 10

$$\begin{aligned} & (\dots a_3 a_2 a_1 a_0 . a_{-1} a_{-2} a_{-3} \dots)_r \\ &= \dots + a_3 r^3 + a_2 r^2 + a_1 r + a_0 r^0 + a_{-1} r^{-1} + a_{-2} r^{-2} + a_{-3} r^{-3} + \dots \\ &= \sum_i (a_i \times r^i) \end{aligned}$$

- Radix 2

$$\begin{aligned} 10010011_2 &= 1 \times 2^7 + 0 \times 2^6 + 0 \times 2^5 + 1 \times 2^4 + 0 \times 2^3 \\ &\quad + 0 \times 2^2 + 1 \times 2^1 + 1 \times 2^0 = 147 \end{aligned}$$

- Radix 8

$$223_8 = 2 \times 8^2 + 2 \times 8^1 + 3 \times 2^0 = 147$$

- Radix 16

$$93_{16} = 9 \times 16^1 + 3 \times 16^0 = 147$$



## Conversion between different radices-to radix 10

- Fractional numbers  $1A6.BE_{16}$

$$\begin{aligned} 1A6.BE_{16} &= 1 \times 16^2 + 10 \times 16^1 + 6 \times 16^0 + 11 \times 16^{-1} + 14 \times 16^{-2} \\ &= 422.7421875 \end{aligned}$$

- How many operations ?  $D = \sum_{i=0}^n (a_i \times r^i)$

$$Oper \leq n + \sum_{i=0}^n i = \mathcal{O}(n^2)$$

- A faster way

$$\begin{aligned} D &= \sum_{i=0}^n a_i \times r^i = r \left( \sum_{i=1}^n a_i \times r^{i-1} \right) + a_0 \\ &= r \left( r \left( \sum_{i=2}^n a_i \times r^{i-2} \right) + a_1 \right) + a_0 = \dots \end{aligned}$$





## Conversion between different radices-to radix 10

- Example

$$\begin{aligned} 1A6_{16} &= 16 \times (16 \times 1 + 10) + 6 \\ &= 16 \times (16 + 10) + 6 = 422_{10}. \end{aligned}$$

- $Oper = \mathcal{O}(n)$



## Conversion between different radices-From radix 10 to radix 2

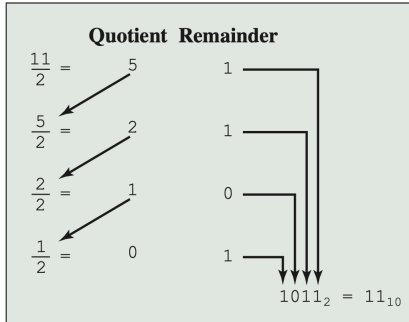
- Example:

$$\begin{aligned} 1A6_{16} &= 16 \times (16 \times 1 + 10) + 6 \\ &= 16 \times (16 + 10) + 6 = 422_{10}. \end{aligned}$$

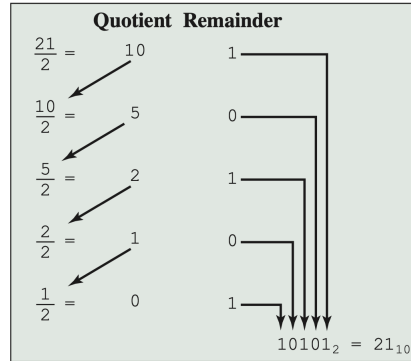
- Find  $a_0, a_1, \dots$  by repeatedly divided by  $r$  and find the quotient and remainder.
- $\frac{422}{16}$ , Quotient: 26, Remainder: 6.  $a_0 = 6$ .
- Divide quotient by  $r$ ,
- $\frac{26}{16}$ , Quotient: 1, Remainder: 10.  $a_1 = A, a_2 = 1$ .
- Find all digits by repeatedly dividing the quotient by  $r$ .



# Conversion between different radices-From radix 10 to radix 2



(a) 11<sub>10</sub>



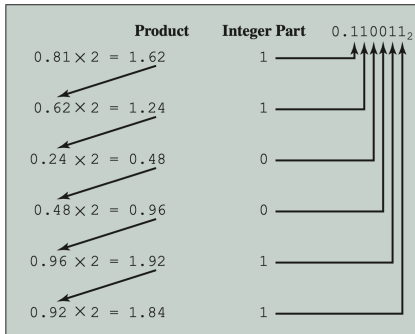
(b) 21<sub>10</sub>

**Figure 9.1** Examples of Converting from Decimal Notation to Binary Notation for Integers

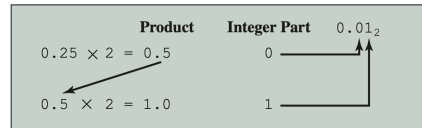


# Conversion between different radices-From radix 10 to radix 2

- Fractional numbers: multiply by  $r$  and get the integer part and repeat this process



(a)  $0.81_{10} = 0.110011_2$  (approximately)



(b)  $0.25_{10} = 0.01_2$  (exactly)

**Figure 9.2** Examples of Converting from Decimal Notation to Binary Notation for Fractions



## Example

Convert 19.7421875 to octal:

- Integer part:  $\frac{19}{8} = 2 + \frac{3}{8}$ ,  $19 = 23_8$ .
- Fractional part:
  - $0.7421875 \times 8 = 5.9375$ , First digit  $a_{-1} = 5$
  - $0.9375 \times 8 = 7.5$ , Second digit  $a_{-2} = 7$
  - $0.5 \times 8 = 4$ , Third digit  $a_{-3} = 4$

Thus,  $0.7421875 = 0.574_8$

- $19.7421875 = 23.574_8$



## Between radix 2, 8 and 16

- there exist convenient ways to convert between radix 2 and octal (hexadecimal).
- Binary digits are grouped into sets of four bits.

0000 = 0	0100 = 4	1000 = 8	1100 = C
0001 = 1	0101 = 5	1001 = 9	1101 = D
0010 = 2	0110 = 6	1010 = A	1110 = E
0011 = 3	0111 = 7	1011 = B	1111 = F

- Starting from the decimal point, group 4 binary digits to the left (to the right, after decimal point), and convert the 4 digits to hex digit
- Example

$$0001\ 1001\ 1000\ 0100 . 1011\ 0110 = 1984.B6_{16}$$

- Similar to octal

$$001\ 100\ 110\ 000\ 100 . 101\ 101\ 100 = 14604.554_8$$



## To radix 2

- Example:

$$7BA3.BC5_{16} = 0111\ 1011\ 1010\ 0011.1011\ 1100\ 0101$$

$$56743.5704_8 = 101\ 110\ 111\ 100\ 011.101\ 111\ 000\ 100$$

0000 = 0	0100 = 4	1000 = 8	1100 = C
0001 = 1	0101 = 5	1001 = 9	1101 = D
0010 = 2	0110 = 6	1010 = A	1110 = E
0011 = 3	0111 = 7	1011 = B	1111 = F



## Sign and Magnitude Representation

- Radix number system: unsigned encoding, representing positive numbers.
- How to deal with negative numbers
- One solution: represent the sign using the leftmost bit, and the rest of the bits for magnitude.
- Example, 1-bit sign, 7-bit magnitude

$$+17 = 00010001$$

$$-17 = 10010001$$

- How to represent 0, we have two zeros 00000000 or 10000000 ?
- Difficult to do arithmetic.





## Excess $K$

- Also referred to as offset binary, excess code or biased representation
- $n$ -bit Excess  $K$  notation: range from  $0 - K$  to  $2^n - 1 - K$ .
- 3 bit is from 0 to 7. Middle is 3 or 4

	$K = 4$	$K = 3$
111	3	4
110	2	3
101	1	2
100	0	1
011	-1	0
010	-2	-1
001	-3	-2
000	-4	-3

- Excess- $2^{n-1}$ , Excess- $2^{n-1} - 1$ , usually the second one
- The lower half is negative, the upper half positive.
- useful in floating point exponent representation



## Representation of Signed Numbers (Positive and Negative)

- The representation scheme is a mapping between the bit pattern (unsigned integer value) and the actual value it represents,

$$f(\text{bit pattern}) = \text{actual value}$$

- The bit pattern can be treated as a positive number. The mapping is from a positive number to a number (data) to be represented.
- E.g., excess  $K$ ,

$$f(\text{bit pattern}) = \text{value} = \text{bit pattern} - K$$



## One's complement

- Positive numbers: same as an unsigned bit pattern
- Negative numbers: representing  $-N$  by inverting all bits of binary representation of  $N$
- inverting a bit  $x$  is  $1 - x$

$$\begin{aligned}7 &= 0111(4 - \text{bit}) \\ -7 &= \text{inverting}(0111) \\ &= 1000\end{aligned}$$

$$\begin{aligned}7 &= 0000\ 0111(8 - \text{bit}) \\ -7 &= \text{inverting}(0000\ 0111) \\ &= 1111\ 1000\end{aligned}$$

- What is mapping  $f(\text{bit pattern})$



## One's complement

- The bit pattern of  $-7$  is

$$1111\ 1111 - 0000\ 0111 = 1111\ 1000$$

$$1111\ 1111 - N = (2^n - 1) - N$$

- bit pattern value  $= (2^n - 1) - N$ , represented value  $= -N$

$$f(\text{bit pattern}) = \text{bit pattern} - (2^n - 1)$$

$$\text{bit pattern} = 2^n - 1 - N$$

- Negative number: Most significant bit (MSB, the leftmost bit)=1
- What about positive numbers?

$$f(\text{bit pattern}) = \text{bit pattern}$$



## One's complement

- What 1101 1010 represents?

$$f = 1101\ 1010 - (2^8 - 1) = 2^7 + 2^6 + 2^4 + 2^3 + 2 - (2^8 - 1) = -37$$

- How to represent 0? Two zeros, 000..., 0, 111..., 1
- Problem with arithmetic: cannot just add them together as if they are simple binary numbers,



## Two's complement

- Positive numbers, same as an unsigned bit pattern
- Negative numbers, representing by One's complement +1,

$$-7(8 \text{ bit}) = 1111 \ 1000 \text{ One's}$$

$$-7(8 \text{ bit}) = 1111 \ 1001 \text{ Two's}$$

- the bit pattern of  $-N$  has a value of  $2^n - N$ .
- Negative number: Most significant bit (MSB, the leftmost bit)=1
- Only one zero, 000 . . . 0 .
- Simple arithmetic



## Example: (8-bit numbers)

Scheme	Bit pattern of $-7$	Bit pattern of $9$	$9 + (-7) = 2$
One's complement	1111 1000	0000 1001	0000 0010
Two's complement	1111 1001	0000 1001	0000 0010
Sign and magnitude	1000 0111	0000 1001	0000 0010
Excess $127(2^{8-1} - 1)$	0111 1000	1000 1000	1000 0001

- Two's complement: just add them together as if they are simple binary numbers, and discard any carry from the MSB,  $9 + (-7)$

$$9 + 2^n - 7 = 2^n(\text{carry out from MSB}) + 2$$

•

$$\begin{array}{r}
 1111\ 1001 \\
 +\ 0000\ 1001 \\
 \hline
 =\ 1\ 0000\ 0010
 \end{array}$$



## Range Extension

- Take an  $n$ -bit integer and store it in  $m$  bits,  $m > n$
- Example: Sign magnitude, just move the sign bit and add zeros

+18	=	00010010	(sign magnitude, 8 bits)
+18	=	0000000000010010	(sign magnitude, 16 bits)
-18	=	10010010	(sign magnitude, 8 bits)
-18	=	1000000000010010	(sign magnitude, 16 bits)

- Two's complement: adding zeros does not work

+18	=	00010010	(twos complement, 8 bits)
+18	=	0000000000010010	(twos complement, 16 bits)
-18	=	11101110	(twos complement, 8 bits)
-32,658	=	1000000001101110	(twos complement, 16 bits)





## Range Extension

- Move the sign bit to the new leftmost position and fill with copies of the sign bit

$$\begin{aligned} -18 &= 11101110 && \text{(twos complement, 8 bits)} \\ -18 &= 1111111111101110 && \text{(twos complement, 16 bits)} \end{aligned}$$

- Two's complement: the value of  $n$  bit representation of  $-|A|$  is  $2^n - |A|$
- After extension, the value of  $m$  bit representation is

$$2^n - |A| + 2^{m-1} + 2^{m-2} + \dots + 2^n = 2^n - |A| + 2^m - 2^n = 2^m - |A|$$

- After range extension, the value represented is still  $|A|$ .



## Two's complement

**Table 10.1** Characteristics of Twos Complement Representation and Arithmetic

<b>Range</b>	$-2^{n-1}$ through $2^{n-1} - 1$
<b>Number of Representations of Zero</b>	One
<b>Negation</b>	Take the Boolean complement of each bit of the corresponding positive number, then add 1 to the resulting bit pattern viewed as an unsigned integer.
<b>Expansion of Bit Length</b>	Add additional bit positions to the left and fill in with the value of the original sign bit.
<b>Overflow Rule</b>	If two numbers with the same sign (both positive or both negative) are added, then overflow occurs if and only if the result has the opposite sign.
<b>Subtraction Rule</b>	To subtract $B$ from $A$ , take the twos complement of $B$ and add it to $A$ .