Approximation as Differentiation notes

1 OVERVIEW

Some notes on the key constructions:

• Categories of families, special case of Grothendieck construction

Set will often be Setoid in the Agda code but we'll ignore that here for now.

2 **DEFINITIONS**

2.1 Indexed families of objects

For I a set and C a category, write Fam(I, C) for the set of I-indexed families of objects of C.

Suppose C a category and I, J sets and x an I-indexed family of objects of C and y a J-indexed family of objects of C. Define y_f for any $f: I \to J$ to be the I-indexed family of objects of C.

2.2 Category of families

For a functor $F: C \rightarrow Set$, we have the category where:

- objects are pairs (I, x) with I an object of C and $x \in FI$
- morphisms from (I, x) to (J, y) are morphisms $f: I \to J$ in C where (Ff)(x) = y

This is the Grothendieck construction for a functor $F: C \to \mathbf{Cat}$, in the special case where $F: C \to \mathbf{Set}$. (Where we read a Set-valued functor as a Cat-valued functor restricted to discrete categories.)

In our code we give a slightly different construction. For a category C, define the category where:

- objects are pairs (I, x) with I a set and x an I-indexed family of objects of C
- $\bullet \,$ morphisms from (I,x) to (J,y) are functions $f:I\to J$ paired with morphisms

Author's address:

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