

Fall CS 145 Final Review Guide.

📧 Prof: Gordon Cormack

📧 ISA: Ashish Mahto

📧 David Duan - j32duan@edu.uwaterloo.ca

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Topic. 8: Gen, RAM, CPU.

Subtopic. ★★☆☆ 8.2: RAM.

Remark.

Now we implement RAM using trie.

Well, basically just a wrapper...

Implementation.

<rkt>

```
(define-struct wrapper (w)) ;; a wrapper struct
(define ram (wrapper empty)) ;; the empty ram
```

</rkt>

Implementation. Store and fetch

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```
;; r: ram we are working with
;; addr: address we want to store the number at
;; num: the number we want to store in the ram
(define (ram-store r addr num)
  (make-wrapper
    (trie-store (wrapper-w r) addr
      (if (and (integer? num) (<= 0 num))
        num
        (error "attempt to store non-Nat value")))))
;; A wrapper contains a single object: w.
;; w is a trie.
;; We want to store num at address=addr inside w.
;; Also, before we storing the value inside w, we check if num is a natural number.

;; r: the ram we are working with
;; addr: the address we want to get value from
(define (ram-fetch r addr)
  (define val (trie-fetch (wrapper-w r) addr))
  (if (number? val)
    val
    (error "attempt to fetch undefined value")))
;; (wrapper-w r) unwraps the ram, gives us the trie.
;; Then we use trie-fetch to get the value at address=addr from the trie.
;; Next we test if val is a legit value (aka if it's a natural number).
;; If yes, we return val.
;; Otherwise we raise an error.
```

</rkt>

Implementation. core-dump

<rkt>

```

(define (core-dump ram)
  (define (core-dump ram start finish)
    (define (ram-fetch r addr)
      (define val (trie-fetch (wrapper-w r) addr))
      val)
    (cond
      [(> start finish) (void)] ;; base case: if start > finish, the job is finished.
      [else
       (define v (ram-fetch ram start)) ;; the value we get from ram-fetch from ram at address=start
       (if (number? v) ;; tests if it is a legit value or not
           (printf "~a: ~a\n" start v)
           (void))
       (core-dump ram (add1 start) finish)]))
    (core-dump ram 0 1000))

```

</rkt>

Remark.

This one should be pretty easy to understand.

(core-dump ram start finish): local helper.

(ram-fetch r addr): the ram-fetch function without the if-statement which checks if the value is legit.

v: the value we get from ram-fetch from ram at address=start

Think of this as a printing out elements of the list one by one.

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Example. (sumto n) using RAM

<rkt>

```

;; Recall sumto with accumulator helper
(define (sumto n)
  (define (sumacc i n a)
    (cond
      [(> i n) a]
      [else (sumacc (add1 i) n (+ a i))]))
  (sumacc 0 n 0))

;; Now we implement it using RAM
;; Rules:
;; 1. tail recursion only
;; 2. one parameter only -- must be RAM
;; 3. result must be RAM

(define (sumto-ram n)
  (define (sumacc ram)
    ;; really (sumacc i n a)
    ;; this calculates a + (sum from i to n)
    ;; ram[0]: i
    ;; ram[1]: n
    ;; ram[2]: a
    (define i (ram-fetch ram 0))
    (define n (ram-fetch ram 1))
    (define a (ram-store ram 2))
    (cond

```

```

[(> i n) ram]
[else
  (define r1 (int-add ram 2 i))    ;; a += i
  (define r2 (int-add1 r1 0))      ;; i += 1
  (sumacc r2))])
(define r1 (ram-store ram 0 0))
(define r2 (ram-store r1 1 n))
(define r3 (ram-store r2 2 0))
(ram-fetch (sumacc r3 2))

```

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Example. PairRAM

A pair-ram can store ints or cons pairs.

pairram[0]: where the list starts

pairram[1..99]: open to interpretation.

<rkt>

```

(define listram (ram-store ram 0 100)) ;; list-start = 100.
(define (list-cons r head-value tail-addr)
  (let* ([tail-value (ram-fetch r tail-addr)]
        [list-start (ram-fetch r 0)]
        [r1 (ram-store r list-start head-value)] ;; we store head-value at list-start
        [r2 (ram-store r1 (add1 list-start) tail-value)] ;; we store tail-value at (add1 list-start)
        [r3 (ram-store r2 tail-addr list-start)] ;; now tail-addr points to list-start
        [r4 (ram-store r3 0 (+ 2 list-start))]) ;; we modify list-start (by adding 2 to it)
    r4))

```

</rkt>

<stepper>

```

~> (define listram (ram-store ram 0 100))
~> (define p1 listram)
~> (define p2 (ram-store p1 77 0))    ;; [77] = 0
~> (define p3 (list-cons p2 10 77))  ;; [77] = '(10)

```

```
=> head-value = 10; tail-addr = 77
```

```
=> tail-value = (ram-fetch p2 77) = 0 (which is empty)
```

```
list-start = (ram-fetch p2 0) = 100
```

```
r1 = (ram-store p2 100 10) = [100]: 10
```

```
r2 = (ram-store r1 101 0) = [101]: 0
```

```
r3 = (ram-store r2 77 100) = [77] : 100
```

```
r4 = (ram-store r3 0 102) = [0] : 102
```

```

;; Right now: [0]    : 102    ;; address 0 points to 102, which is our next empty spot to add values.
              [77]    : 100    ;; address 77 points to a list, which starts at 100.
              [100]   : 10     ;; address 100 stores 10, which is the first value of our list.
              [101]   : 0      ;; address 101 stores 0, indicates there is no more value in the list.

```

```
~> (define p4 (list-cons p3 20 77)) ;; [77] = '(20 10)
```

```
=> head-value = 20; tail-addr = 77
```

```
=> tail-value = (ram-fetch p3 77) = 100
list-start = (ram-fetch p3 0) = 102
r1 = (ram-store p3 102 20) = [102]: 20
r2 = (ram-store r1 103 100) = [103]: 100
r3 = (ram-store r2 77 102) = [77]: 102
r4 = (ram-store r3 0 104) = [0]: 104

;; Right now: [0] : 104      ;; address 0 points to 104, which is our next empty spot to add values.
              [77] : 102    ;; address 77 points to a list, which starts at 102.
              [100] : 10    ;; address 100 stores 10
              [101] : 0     ;; address 101 stores 0, indicates there is no more value in the list.
              [102] : 20    ;; address 102 stores 20
              [103] : 100   ;; address 103 points to 100, which is the address storing the element
                          ;; after 20.
```

```
~> (define p5 (list-cons p4 30 77)) ;; [77] = '(30 20 10)
```

```
=> head-value = 30; tail-addr = 77
=> tail-value = (ram-fetch p4 77) = 102
list-start = (ram-fetch p4 0) = 104
r1 = (ram-store p4 104 30) = [104]: 30
r2 = (ram-store r1 105 102) = [105]: 102
r3 = (ram-store r2 77 104) = [77]: 104
r4 = (ram-store r3 0 106) = [0]: 106

;; Right now: [0] : 106      ;; address 0 points to 104, which is our next empty spot to add values.
              [77] : 104    ;; address 77 points to a list, which starts at 104.
              [100] : 10    ;; address 100 stores 10
              [101] : 0     ;; address 101 stores 0, indicates there is no more value in the list.
              [102] : 20    ;; address 102 stores 20
              [103] : 100   ;; address 103 points to 100, which is the address storing the element
                          ;; after 20.
              [104] : 30    ;; address 104 stores 30
              [105] : 102   ;; address 105 points to 102, which is the address storing the element
                          ;; after 30.
```

```
;; At this point, you should realize that tail-addr (in this case 77) always points to the address
;; of first element in the list.
```

```
;; Thus, following the pointers we can retrieve the entire list.
```

```
</stepper>
```

```
<rkt>
```

```
(define (list-empty? r list-addr)
  (zero? (ram-fetch r list-addr)))

;; get the first element of the list
;; r: ram
;; list-addr: where does our list start?
(define (list-first r list-addr)
  (let* ([addr-of-car (ram-fetch r list-addr)]
         [val-of-car (ram-fetch r addr-of-car)])
    val-of-car))
```

```

;; drop the first element of the list, aka replace the first element with the second element
;; and since this is a linked list, we don't need to worry about anything else.
;; r: ram
;; list-addr: where does our list start?
(define (list-dropfirst r list-addr)
  (let* {[addr-of-car (ram-fetch r list-addr)]
         [addr-of-cdr (add1 addr-of-car)] ;; make sure you understand this step
         [val-of-second (ram-fetch addr-of-cdr)]}
    ;; now we want to store value of second at the address which originally stores first.
    (ram-store r addr-of-car val-of-second)))

;; store val as the first element.
;; r: ram
;; list-addr: where does our list start?
;; mew-val: our new value for first element.
(define (list-setfirst r list-addr new-val)
  (let* {[addr-of-car (ram-fetch r list-addr)]
         [new-ram (ram-store r addr-of-car new-val)]}
    new-ram))

;; This is explained in CPU.pdf, I'm too lazy to do it again.
(define (ram-copy ram from-addr to-addr)
  (ram-store ram to-addr (ram-fetch ram from-addr)))

;; listasRAM->list
(define (outstream r list-addr)
  (define r1 (ram-copy r list-addr 55));; 55 is arbitrary. We copy [list-addr] to [55]
  (define (iterate r)
    (cond
      [(list-empty? r 55) empty]
      [else (cons (list-first r 55) (iterate (list-dropfirst r 55)))]))
  (iterate r1))

;; list->listasRAM
(define (instream ram list-addr lst)
  (define r1 (ram-store ram list-addr 0))
  (define (helper r lst)
    (cond
      [(empty? lst) r]
      [else (helper (list-cons r (first lst) list-addr) (cdr-lst))]))
  (helper r1 (reverse lst))) ;; make sure you understand why we need to reverse the list.
</rkt>

```

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END

