

A Mathematical Description of the CSS Cipher

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Given:

- a 40-bit scrambling key expressed as a vector of five 8-bit bytes in $k[0:4]$,
- a DVD sector expressed as a vector of 2^{11} (2048) 8-bit bytes in $s[0:2047]$,
- a function $r()$ which reverses the bits in an 8-bit byte,
- two linear feedback shift registers $l[]$ and $m[]$, defined as follows,
- the sum of $l[]$ and $m[]$, called $x[]$,
- an 8-bit PAL $p[]$, defined as follows, and
- the temporary variables $k'[]$, $l'[]$ and $m'[]$,

Then:

$$k'[n] = r(k[n] \wedge s[84+n])$$

$$l[127] = k'[1] \mid (1 \ll 8) \mid (k'[0] \ll 9)$$

$$m[127] = k'[4] \mid (k'[3] \ll 8) \mid ((k'[2] \& 00011111) \ll 16) \mid (1 \ll 21) \mid ((k'[2] \& 11100000) \ll 17)$$

$$x[127] = 0$$

$$l'[n] = l[n] \wedge (l[n] \gg 14)$$

$$l[n=128:2047] = (l[n-1] \gg 8) \wedge (l'[n-1] \ll 9) \wedge (l'[n-1] \ll 12) \wedge (l'[n-1] \ll 15)$$

$$m'[n] = m[n] \wedge (m[n] \gg 3) \wedge (m[n] \gg 4) \wedge (m[n] \gg 12)$$

$$m[n=128:2047] = (m[n-1] \gg 8) \wedge (m'[n-1] \ll 17)$$

$$x[n=128:2047] = !(l[n] \gg 9) + (m[n] \gg 17) + (x[n-1] \gg 8)$$

The PAL $p[]$ is defined with the following equations, where a is the lowest input bit and A is the lowest output bit:

$$A = !(a \& b) \wedge d$$

$$B = !(e \& f) \wedge g$$

$$C = !(A \& B) \wedge f$$

$$D = !(A \& B) \wedge b$$

$$E = !(a \& b) \wedge c$$

$$F = !(e \mid f) \wedge h$$

$$G = !(E \& F) \wedge a$$

$$H = !(E \mid F) \wedge e$$

And the descrambled sector corresponding to $k[]$ and $s[]$ is:

$$s'[n=0:127] = s[n]$$

$$s'[n=128:2047] = p[s[n]] \wedge x[n]$$

A translation of the previous into C code:

```
#include #include
int
main()
{
    unsigned char k[5], s[2048], kp[5], sp[2048];
    unsigned int l[2048], m[2048], x[2048], lp[2048], mp[2048];
    unsigned char r[256], p[256];
    int n;

    for (n = 0; n <= 255; n++) {
        int a, b, c, d, e, f, g, h, A, B, C, D, E, F, G, H;

        a = (n >> 0) & 1;
        b = (n >> 1) & 1;
        c = (n >> 2) & 1;
        d = (n >> 3) & 1;
        e = (n >> 4) & 1;
        f = (n >> 5) & 1;
        g = (n >> 6) & 1;
        h = (n >> 7) & 1;

        r[n] = (a << 7) | (b << 6) | (c << 5) | (d << 4) | (e << 3) |
            (f << 2) | (g << 1) | (h << 0);

        A = (a & b) ^ d ^ 1;
        B = (e & f) ^ g ^ 1;
        C = (A & B) ^ f ^ 1;
        D = (A & B) ^ b ^ 1;
        E = (a & b) ^ c ^ 1;
        F = (e | f) ^ h ^ 1;
        G = (E & F) ^ a ^ 1;
        H = (E | F) ^ e ^ 1;

        p[n] = (H << 7) | (G << 6) | (F << 5) | (E << 4) | (D << 3) |
            (C << 2) | (B << 1) | (A << 0);
    }

    for (read(0, k, 5); read(0, s, 2048); write(1, sp, 2048)) {
        for (n = 0; n <= 4; n++)
            kp[n] = r[k[n] ^ s[84+n]];

        l[127] = kp[1] | (1 << 8) | (kp[0] << 9);
        m[127] = kp[4] | (kp[3] << 8) | ((kp[2] & 0x1f) << 16) |
            (1 << 21) | ((kp[2] & 0xe0) << 17);
        x[127] = 0;

        for (n = 128; n <= 2047; n++) {
            lp[n-1] = l[n-1] ^ (l[n-1] >> 14);
            l[n] = (l[n-1] >> 8) ^ (lp[n-1] << 9) ^
                (lp[n-1] << 12) ^ (lp[n-1] << 15);
            l[n] &= 0x1ffff;

            mp[n-1] = m[n-1] ^ (m[n-1] >> 3) ^ (m[n-1] >> 4) ^
                (m[n-1] >> 12);
            m[n] = (m[n-1] >> 8) ^ (mp[n-1] << 17);
            m[n] &= 0xffffffff;

            x[n] = (0xff & ~(l[n] >> 9)) + (m[n] >> 17) +
                (x[n-1] >> 8);
        }

        for (n = 0; n <= 127; n++)
            sp[n] = s[n];
        for (n = 128; n <= 2047; n++)
            sp[n] = p[s[n]] ^ x[n];
    }
}
```

}

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