Final Project CS 2420

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Introduction

Our final project for CS 2420 was implemented in Python. In addition to completing all of the requirements, we also included some bonus features. The search engine has a ranking algorithm, and returns results sorted by relevancy. It can accommodate extended combinations of boolean operators, including all combinations of AND, OR, and NOT. It eliminates stop words for efficiency. All of these features come packaged in a fancy interface that was implemented using JavaScript, HTML/CSS, and a dash of PHP.

Algorithms

The three main factors that influence the speed of our engine are the number of words in documents (W), the number of documents (D), and the number of words in the query (Q). The size and quantity of documents affects the running time due to the process of reading each file for indexing. The primary data structure that we use is a dictionary (basically a hash table). We use a dictionary for each document, to track the number of occurrences of each word in the given document. We also use a dictionary of arrays, which contain all words in all the documents (excluding stop words) and an array of documents which contain said word as key value pairs. Filling these arrays is O(s), based on the quantity and size of the documents. This is unavoidable for any given data structure. However, once these hash tables are filled, lookup time is approximately O(1). This is the primary reason for utilizing a hash table.

The algorithm that we use to address the boolean operators in the search query (parseQuery) is recursive. Because it must iterate over each word in each document for each part of the search query, in the worst case, it is $O(Q \cdot W \cdot D^2)$. Given that the majority of queries will not be very complicated, this is approximately $O(W \cdot D^2)$, which is not the fastest, but is not noticeably slow until the number of documents or the diversity of words within the various documents increases. The code itself is clean, readable and functions well. In larger scales, a more robust method involving a database engine, perhaps such as MySQL, would clearly be desirable to reduce the time complexity in relation to the number of documents indexed. We chose to implement this algorithm recursively to allow for more flexibility in search query structure.

We also implement document ranking. The algorithm considers two factors: Word Frequency (how often a word from the query appears in the document), and Document Frequency (how often a word appears in all documents). In principle, if a word from the search query appears in the document more times, then it is more likely to be the document the user is searching for. Conversely if the same

word exists in many documents, then it is less specific to a given document and, therefore, it is less likely that particular word from the search query has an influence on what the user is searching for. The time to look up both of these factors from the index is O(1) because of our use of a hash table, making this portion of the algorithm very nice, indeed.

Instead of writing our own sorting method to sort the documents by ranking, we used Python's built in Timsort, because its time complexity is, in best case, better than both Quicksort and Mergesort: O(n). The average and worst cases are equivalent for the Mergesort: $O(n \cdot log(n))$. Timsort makes use of already sorted subsets of data, storing them in temporary arrays before merging them. The only potential hazard with Timsort is space complexity, which is the worst case is O(n). With such a small scale project, however, space was not of much concern.

Design Issues

At first glance, handling the boolean operators of the search queries seemed trivial. Once we began to write the function, however, it quickly became apparent that we would need to either limit the number of search terms so that we could focus on one boolean operator at a time, or we would need to write a recursive function. We chose to go with recursion so that we could have flexibility in the number of search items. This required a restructuring of the existing parsing algorithm so that it could accommodate recursion.

Our parseQuery function initially takes the dictionary of all the words from the indexing function, and, treating spaces and punctuation as an "AND" operator, uses the "AND", "OR" and "NOT" operators in the query to cut and merge each set obtained from the hash table, returning a list of documents that match the search query. It's a pretty rudimentary implementation of the Set ADT functions.

Another design challenge was the delivery of our search engine. We knew the end result would be a web page, but as we looked at AWS it seemed to be much more sophisticated a service than we needed, making it difficult to configure it as we needed. Instead we setup a server with the basic necessities: apache, PHP, python 2.7.

We know that python is used by many web services and there are frameworks to help python integrate with apache, but similarly these seemed to be much more than we bargained for.

The result was we added a small PHP script to call the python script (since all we ever needed from apache was the POST array sent to the server asynchronously with JavaScript). This does, however, leave the minor issue that shell arguments are executed with user input. This is not the best design practice for a website and would not be advised for a production environment.

Optimization and Scalability

With regards to optimization, not much can be done to improve the indexing. In any situation, all of the words in a document and all of the documents themselves have to be indexed. However, as we've discussed in class, portions of the document could be indexed in parallel across multiple computers to reduce the time required.

Within the searching part of the algorithm, however, some optimization could be achieved with dividing and conquering the search query. The parseQuery function has a worst case of $\approx O(Q \cdot W \cdot D^2)$, if the query is comprised of common words that are present in many documents. If the algorithm could be broken up in such a way that common words are treated differently than specific words, then the overall time complexity could be reduced.

There are some things that need to be addressed due to the nature of web browsers and their interaction with the server. Our website asynchronously loads content based on the results of a given query. As the results get larger, the resources and time required by the browser to dynamically generate and insert that HTML content is significant. The first optimization to be performed, therefore, would be to implement a pagination feature. Similar to how Google only shows a limited number of search results on the first page, and results after that are unlikely to be what you're looking for regardless.

The second optimization that needs to occur based on the nature of a web server is that the indexing loads with each query, after which memory is released

GREETINGS, STRANGER.

WHATEVER QUEST DRIVES YOU,
ABANDON IT.

YOU SHALL FIND NO ANSWERS
IN THESE DESOLATE WASTES.

I KNEW I WOULDN'T.
I GUESS I ... JUST
HAD TO SEE.

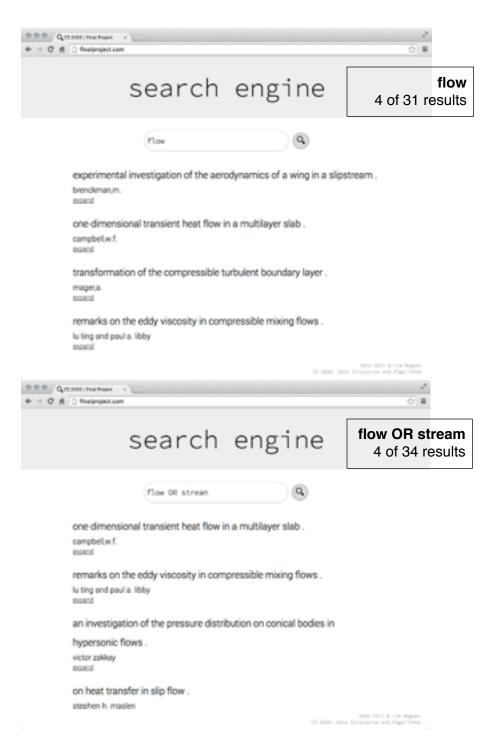
I HATE FEELING DESPERATE ENOUGH TO VISIT THE SECOND PAGE OF GOOGLE RESULTS.

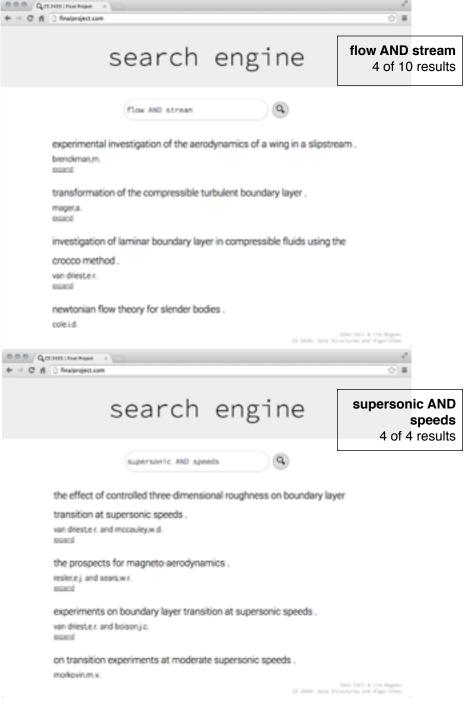
Source: http://xkcd.com/1334/

back to the operating system once this script is finished executing. In an ideal scenario our indexer would remain active as long as the server is on, or better yet store the index to the file system for long term usage, and wait for requests from the web server. This requires the use of sockets, or other methods, that we do not yet have much practice in implementing.

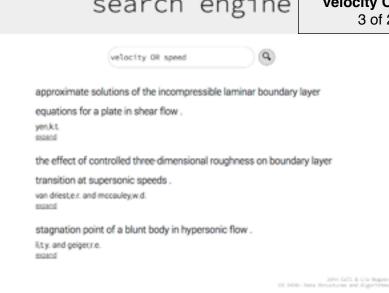
Graphic Design Notes

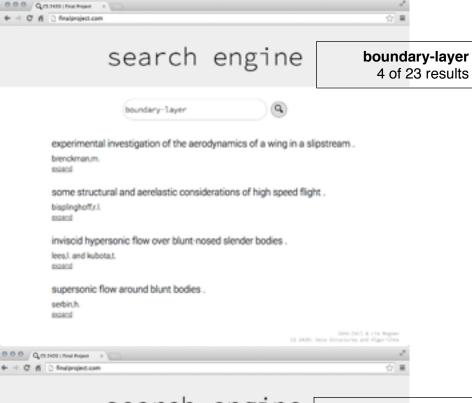
The face of any programming project inherently influences how people react to it. Aside from its incredibly complex and groundbreaking algorithms, Google's simple, minimalist design draws people in. Inspired by Google's originality, we implemented a clean HTML interface to showcase our search engine. We felt that showing the search query in context with highlighting would enhance the usefulness of our engine, and that the design as a whole elevates the experience of searching.











search engine

reynolds AND number OR reynolds AND numbers 3 of 7 results

reynolds AND number OR reynolds

measurements of the effect of two-dimensional and three-dimensional roughness elements on boundary layer transition .

Webanoff,p.s.

a new technique for investigating heat transfer and surface phenomena under hypersonic flow conditions .

femi,a. and libby,p.a.

expand

inviscid hypersonic flow over blunt-nosed slender bodies.

lees.I. and kubota.t.

expand

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2420: Unio Structures and Algorithm

Five additional queries + Help Screen

```
jcall@callServ:/sites/finalProject$ python searchEngine.py -i
                  Enter a query or '-q' to exit:
                  analysis
   analysis
                  01 - Calculated Rank: 0.07692 - File: cranfield0010
                  02 - Calculated Rank: 0.07692 - File: cranfield0015
                  03 - Calculated Rank: 0.07692 - File: cranfield0016
                  04 - Calculated Rank: 0.07692 - File: cranfield0021
                  05 - Calculated Rank: 0.07692 - File: cranfield0025
                  06 - Calculated Rank: 0.07692 - File: cranfield0029
                  07 - Calculated Rank: 0.07692 - File: cranfield0035
                  08 - Calculated Rank: 0.07692 - File: cranfield0036
                  09 - Calculated Rank: 0.07692 - File: cranfield0046
                  10 - Calculated Rank: 0.07692 - File: cranfield0049
                  11 - Calculated Rank: 0.15385 - File: cranfield0030
                  12 - Calculated Rank: 0.15385 - File: cranfield0039
                  13 - Calculated Rank: 0.23077 - File: cranfield0047
                  Enter a query or '-q' to exit:
                  analytic
   analytic
                  01 - Calculated Rank: 0.25000 - File: cranfield0005
                  02 - Calculated Rank: 0.25000 - File: cranfield0006
                  03 - Calculated Rank: 0.25000 - File: cranfield0034
                  04 - Calculated Rank: 0.50000 - File: cranfield0049
                  Enter a query or '-q' to exit:
                  analytic and not analysis
analytic AND
                  01 - Calculated Rank: 0.25000 - File: cranfield0005
                  02 - Calculated Rank: 0.25000 - File: cranfield0006
NOT analysis
                  03 - Calculated Rank: 0.33333 - File: cranfield0034
                  Enter a query or '-q' to exit:
                  analytic or analysis
 analytic OR
                  01 - Calculated Rank: 0.07692 - File: cranfield0010
                  02 - Calculated Rank: 0.07692 - File: cranfield0016
   analysis
                  03 - Calculated Rank: 0.07692 - File: cranfield0021
                  04 - Calculated Rank: 0.07692 - File: cranfield0035
                  05 - Calculated Rank: 0.07692 - File: cranfield0046
                  06 - Calculated Rank: 0.15385 - File: cranfield0030
                  07 - Calculated Rank: 0.15385 - File: cranfield0039
                  08 - Calculated Rank: 0.17692 - File: cranfield0015
                  09 - Calculated Rank: 0.17692 - File: cranfield0025
                  10 - Calculated Rank: 0.17692 - File: cranfield0029
                  11 - Calculated Rank: 0.23077 - File: cranfield0047
                  12 - Calculated Rank: 0.25000 - File: cranfield0005
                  13 - Calculated Rank: 0.25000 - File: cranfield0006
                  14 - Calculated Rank: 0.25000 - File: cranfield0034
                  15 - Calculated Rank: 0.27692 - File: cranfield0036
                  16 - Calculated Rank: 0.77692 - File: cranfield0049
                  Enter a query or '-q' to exit:
```

```
jcall@callServ:/sites/finalProject$ python searchEngine.py -q "not analysis"
                 01 - Calculated Rank: 0.00000 - File: cranfield0001
NOT analysis 02 - Calculated Rank: 0.00000 - File: cranfield0002
                 03 - Calculated Rank: 0.00000 - File: cranfield0003
                 04 - Calculated Rank: 0.00000 - File: cranfield0004
                 05 - Calculated Rank: 0.00000 - File: cranfield0005
                 06 - Calculated Rank: 0.00000 - File: cranfield0006
                 07 - Calculated Rank: 0.00000 - File: cranfield0007
                 08 - Calculated Rank: 0.00000 - File: cranfield0008
                 09 - Calculated Rank: 0.00000 - File: cranfield0009
                 10 - Calculated Rank: 0.00000 - File: cranfield0011
                 11 - Calculated Rank: 0.00000 - File: cranfield0012
                 12 - Calculated Rank: 0.00000 - File: cranfield0013
                 13 - Calculated Rank: 0.00000 - File: cranfield0014
                 14 - Calculated Rank: 0.00000 - File: cranfield0017
                 15 - Calculated Rank: 0.00000 - File: cranfield0018
                 16 - Calculated Rank: 0.00000 - File: cranfield0019
                 17 - Calculated Rank: 0.00000 - File: cranfield0020
                 18 - Calculated Rank: 0.00000 - File: cranfield0022
                 19 - Calculated Rank: 0.00000 - File: cranfield0023
                 20 - Calculated Rank: 0.00000 - File: cranfield0024
                 21 - Calculated Rank: 0.00000 - File: cranfield0026
                 22 - Calculated Rank: 0.00000 - File: cranfield0027
                 23 - Calculated Rank: 0.00000 - File: cranfield0028
                 24 - Calculated Rank: 0.00000 - File: cranfield0031
                 25 - Calculated Rank: 0.00000 - File: cranfield0032
                 26 - Calculated Rank: 0.00000 - File: cranfield0033
                 27 - Calculated Rank: 0.00000 - File: cranfield0034
                 28 - Calculated Rank: 0.00000 - File: cranfield0037
                 29 - Calculated Rank: 0.00000 - File: cranfield0038
                  30 - Calculated Rank: 0.00000 - File: cranfield0040
                 31 - Calculated Rank: 0.00000 - File: cranfield0041
                 32 - Calculated Rank: 0.00000 - File: cranfield0042
                  33 - Calculated Rank: 0.00000 - File: cranfield0043
                 34 - Calculated Rank: 0.00000 - File: cranfield0044
                 35 - Calculated Rank: 0.00000 - File: cranfield0045
                  36 - Calculated Rank: 0.00000 - File: cranfield0048
                 37 - Calculated Rank: 0.00000 - File: cranfield0050
                 jcall@callServ:/sites/finalProject$
```

help menu

Keyword highlighting:

