

Foundations of Computer Science HW 8

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22.14(a)-(c) Answer True or False

- (a) *A bijection must be an injection.*
- (b) *There is a bijection from \mathbb{Q} to \mathbb{R}*
- (c) *There is a bijection from \mathbb{Q} to \mathbb{R}*

22.18(a)-(c) Prove or Disprove

- (a) *\mathbb{Z}^2 is the set of pairs $\{(z_1, z_2) \mid z_1, z_2 \in \mathbb{Z}\}$. \mathbb{Z}^2 is countable.*
- (b) *\mathbb{Q} is the set of rational numbers, $\mathbb{Q} = \{r \mid r = \frac{a}{b}, a \in \mathbb{Z}, \text{ and } b \in \mathbb{N}\}$. \mathbb{Q} is countable.*
- (c) *F is the set of all functions from \mathbb{N} to \mathbb{N} , $F = \{f \mid f : \mathbb{N} \rightarrow \mathbb{N}\}$. F is countable.*

24.5(d)-(f) Give DFA's for the following

- (d) *The strings which begin with 10 and end with 01.*
- (e) *$\mathcal{L} = \{01^n \mid n \geq 0\}$*
- (f) *The language with all strings whose length is divisible by 3.*

25.4(a)-(c) Give a DFA and a CFG for the following

- (a) *$\mathcal{L} = \{01^n \mid n \geq 0\}$*
- (b) *$\mathcal{L} = \{0^n 1^n \mid 0 \leq n \leq 5\}$*
- (c) *$\mathcal{L} = \{\text{strings which end in a 1}\}$*

25.15(a)-(c) Consider the language $\mathcal{L} = \{\epsilon, 1, 11, 111, \dots\} = \{1\}^*$

- (a) *Show that the CFG $S \rightarrow \epsilon \mid 1 \mid 1S$ generates \mathcal{L} . Give a derivation of 111.*
- (b) *Show that the CFG $S \rightarrow \epsilon \mid 1 \mid SS$ generates \mathcal{L} . Give two different derivations of 111.*
- (c) *A leftmost (rightmost) derivation replaces the leftmost (rightmost) variable at every step. For the grammar in (b), give left and rightmost derivations of 111.*

25.16(a)-(c) Given the CFG $S \rightarrow A1B; A \rightarrow \epsilon \mid 0A; B \rightarrow \epsilon \mid 0B \mid 1B$, Give leftmost and rightmost derivations, and parse trees, of

(a) 00101

(b) 1001

(c) 00011

26.4(a)-(b) Give high-level pseudocode for Turing Machines that solve these problems. In some cases you are asked for a decider. In others, you are asked for a transducer.

(a) *Regular languages:* $\mathcal{L}_1 = \{ *01* \}$ and $\mathcal{L}_2 = \{ *01 \}$

(b) *Not CFL:* $\mathcal{L} = \{ 0^{\bullet n} \# 1^{\bullet n} \# 0^{\bullet n} \}$ where $\#$ is a punctuation symbol.