

# 2020 FALL MAS583 HW7

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## 1 Problem 1

### 1.1 Question

Let  $G = (V, E)$  be the graph whose vertices are all  $7^n$  vectors of length  $n$  over  $\mathbb{Z}_7$ , in which two vertices are adjacent iff they differ in precisely one coordinate. Let  $U \subset V$  be a set of  $7^{n-1}$  vertices of  $G$ , and let  $W$  be the set of all vertices of  $G$  whose distance from  $U$  exceeds  $(c+2)\sqrt{n}$ , where  $c > 0$  is a constant. Prove that  $|W| \leq 7^n \cdot e^{-c^2/2}$ .

### 1.2 Answer

To show the inequality, suffice it to show that if we pick a vertex  $x \in V$  uniformly at random, the probability that  $x$  has a distance from the given  $U$  exceeds  $(c+2)\sqrt{n}$  is at most  $e^{-c^2/2}$ . Consider  $A = \mathbb{Z}_7$  and  $B = [n]$ , then  $\Omega = A^B$ , the set of functions  $g : B \rightarrow A$ , is equivalent to  $V = V(G)$ . Let

$$p_{ab} = \Pr[g(b) = a] = 1/7,$$

for each  $a \in A$  and  $b \in B$ . And fix a gradation

$$\emptyset = B_0 \subset B_1 \subset \cdots \subset B_n = B,$$

where  $B_i = [i]$ , for each  $i \in [n]$ . For given  $U$ , let  $L : A^B \rightarrow \mathbb{R}$  be the function denoting the distance from  $U$  of the corresponding vertex. We define a martingale  $X_0, X_1, \dots, X_n$  by setting

$$X_i(h) = \mathbb{E}[L(g)|g(b) = h(b), \forall b \in B_i = [i]],$$

in specific  $X_0 = \mathbb{E}[L(g)|g \in \Omega]$  is a constant and  $X_n = L$ . We can easily check that  $L$  satisfies the Lipschitz condition relative to the gradation above as changing in one coordinate can make the distance change by at most 1. Thus, our corresponding martingale satisfies

$$|X_{i+1}(h) - X_i(h)| \leq 1$$

for all  $0 \leq i < n$  and  $h \in A^B$ . By setting  $\mu = X_0 = \mathbb{E}[L(g)|g \in \Omega]$ , we have for all  $\lambda > 0$ ,

$$\Pr[L(g) \geq \mu + \lambda\sqrt{n}] < e^{-\lambda^2/2} \tag{1}$$

and

$$\Pr[L(g) \leq \mu - \lambda\sqrt{n}] < e^{-\lambda^2/2}. \tag{2}$$

Now, with (1), we only need to show that  $\mu \leq 2\sqrt{n}$  to complete the proof. Actually, by setting  $\lambda = \sqrt{2\ln 7}$  in (2), we have

$$\Pr[L(g) \leq \mu - \sqrt{2\ln 7}\sqrt{n}] < 1/7,$$

which implies that  $\mu \leq \sqrt{2\ln 7}\sqrt{n} < 2\sqrt{n}$  as  $\Pr[L(g) \leq 0] = |U|/|V| = 1/7$ , completing the proof.

## 2 Problem 2

### 2.1 Question

(\*) Let  $G = (V, E)$  be a graph with chromatic number  $\chi(G) = 1000$ . Let  $U \subset V$  be a random subset of  $V$  chosen uniformly among all  $2^{|V|}$  subsets of  $V$ . Let  $H = G[U]$  be the induced subgraph of  $G$  on  $U$ . Prove that

$$\Pr[\chi(H) \leq 400] < 1/100.$$

### 2.2 Answer

First, we show the following lemma.

**Lemma 2.1** *Given  $G = (V, E)$  with chromatic number  $\chi(G)$ , let  $A \cup B = V$  be a partition of  $V$ , and let  $H_A$  and  $H_B$  be the subgraphs of  $G$  induced by  $A$  and  $B$ , respectively. Then  $\chi(A) + \chi(B) \geq \chi(G)$ .*

**Proof.** Consider a coloring where  $G[A]$  is properly colored by  $\chi(A)$  colors and  $G[B]$  is properly colored with another  $\chi(B)$  different colors, which is a proper coloring of  $G$  with  $\chi(A) + \chi(B)$  colors, completing the proof. ■

As  $\chi(G) = 1000$ , we can find a partition of  $V$  into 1000 independent sets  $\{V_1, \dots, V_{1000}\}$ . For any  $U_0 \subset V$ , we build a martingale  $\mathbb{E}[\chi(G[U])] = X_0, X_1, \dots, X_{1000} = \chi(G[U_0])$  by

$$X_i(H) = \mathbb{E} \left[ \chi(G[H]) : H \cap \bigcup_{j \in [i]} V_j = U \cap \bigcup_{j \in [i]} V_j \right].$$

It is easy to check that the above martingale satisfies the Lipschitz condition  $|X_{i+1} - X_i| \leq 1$ , because at worst you can only change the coloring of the newly exposed part  $H \cap V_{i+1}$ , where at most one color should be used in a minimum proper coloring. By the uniformity, we have

$$\mathbb{E}[\chi(G[U])] = \mathbb{E}[\chi(G[V \setminus U])],$$

with the lemma, we further have

$$\mathbb{E}[\chi(G[U])] = \frac{1}{2}(\mathbb{E}[\chi(G[U])] + \mathbb{E}[\chi(G[V \setminus U])]) \geq \chi(G)/2 = 500.$$

Finally, by Azuma's Inequality, we have

$$\Pr[\chi(H) \leq 400] = \Pr[X_{1000} \leq 400] \leq \Pr[X_{1000} - \mathbb{E}[\chi(G[U])] \leq -100] \leq e^{-100^2/2000} < 1/100,$$

completing the proof.

## 2.3 References

<https://math.stackexchange.com/questions/2118314>

## 3 Problem 3

### 3.1 Question

Prove that there is an absolute constant  $c$  such that for every  $n > 1$  there is an interval  $I_n$  of at most  $c\sqrt{n}/\log n$  consecutive integers such that the probability that the chromatic number of  $G(n, 0.5)$  lies in  $I_n$  is at least 0.99.

### 3.2 Answer

Let  $\epsilon > 0$  be arbitrarily small and let  $u = u(n, p, \epsilon)$  be the least integer so that

$$\Pr[\chi(G) \leq u] > \epsilon.$$

Now define  $Y(G)$  to be the minimal size of a set of vertices  $S$  for which  $G - S$  may be  $u$ -colored. This  $Y$  satisfies the vertex Lipschitz condition since at worst one could add a vertex to  $S$ . Apply the vertex exposure martingale on  $G(n, p)$  to  $Y$ . Letting  $\mu = E[Y]$ , for all  $\lambda > 0$ ,

$$\Pr[Y \leq \mu - \lambda\sqrt{n-1}] < e^{-\lambda^2/2}, \quad (3)$$

$$\Pr[Y \geq \mu + \lambda\sqrt{n-1}] < e^{-\lambda^2/2}. \quad (4)$$

Let  $\lambda$  satisfy  $e^{-\lambda^2/2} = \epsilon$ , with (3), we have

$$\Pr[Y \leq \mu - \lambda\sqrt{n-1}] < \epsilon \leq \Pr[Y \leq 0],$$

which implies that  $\mu \leq \lambda\sqrt{n-1}$ . Apply this to (4), we have

$$\Pr[Y \geq 2\lambda\sqrt{n-1}] \leq \Pr[Y \geq \mu + \lambda\sqrt{n-1}] < \epsilon,$$

i.e., with probability at least  $1 - \epsilon$  there is a  $u$ -coloring of all but at most  $2\lambda\sqrt{n-1}$  vertices, where  $c_1$  is a positive constant.

To complete the proof, it remains to show the following lemma.

**Lemma 3.1** *Given  $\epsilon > 0$ , there is an absolute constant  $c = c(\epsilon)$  such that with probability at least  $1 - \epsilon$ , every  $2\lambda\sqrt{n-1}$  vertices of  $G = G(n, 0.5)$  may be  $(c\sqrt{n}/\log n)$ -colored, where  $\lambda = \lambda(\epsilon)$  satisfies  $e^{-\lambda^2/2} = \epsilon$ .*

**Proof.** As a known result, by setting  $f(k) = \binom{m}{k} 2^{-\binom{k}{2}}$ ,  $k_0$  so that  $f(k_0 - 1) > 1 > f(k_0)$ ,  $k = k_0 - 4$  so that  $k \sim 2 \log_2 m$  and  $f(k) > m^{3+o(1)}$ , we have

$$\Pr[\omega(G(m, 0.5)) < k] = e^{-m^{2+o(1)}}.$$

By the property of  $G(m, 0.5)$ , we have

$$\Pr[\alpha(G(m, 0.5)) \leq \log m < k] \leq e^{-m^{2+o(1)}}.$$

Let  $m = n^{1/3}$ , it follow that the probability that there is a set of vertices  $S$  of size at least  $n^{1/3}$  that does not contain an independent set of size at least  $\log m = \frac{1}{3} \log n$  is at most

$$\binom{n}{n^{1/3}} e^{-n^{2/3+o(1)}} \leq n^{n^{1/3}} e^{-n^{2/3+o(1)}} \leq \exp(n^{1/3} \log n - n^{2/3+o(1)}) = o(1),$$

i.e., in  $G(n, 0.5)$ , almost always, every set of at least  $n^{1/3}$  contains an independent set of size at least  $\frac{1}{3} \log n$ . In particular, with probability at least  $1 - \epsilon$ , this holds for large  $n$ . Now, let  $S_0 = S$  be a set of size at most  $2\lambda\sqrt{n-1}$ , when  $n$  is sufficiently large, with probability at least  $1 - \epsilon$ , we can do the following for  $i \geq 1$  until the algorithm stops:

- If  $|S_i| \leq \sqrt{n}/\log n$ , stop.
- Let  $U_i$  be an independent set of size at least  $\frac{1}{3} \log n$  in  $S_i$ .
- Set  $S_{i+1} = S_i \setminus U_i$

Clearly, the above algorithm ends in  $l \leq 6\lambda\sqrt{n}/\log n$  steps. Let  $W = S \setminus \bigcup_{i \in [l]} U_i$ , we obtain a proper coloring of  $G[S]$  with at most  $(6\lambda + 1)\sqrt{n}/\log n$  by coloring each  $U_i$  with color  $i$  and coloring each vertex in  $W$  with a distinct new color. It follows that, by setting  $c \geq (6\lambda + 1)$

$$\chi(G[S]) \leq (6\lambda + 1)\sqrt{n}/\log n \leq c\sqrt{n}/\log n,$$

completing the proof. ■

Now we have

- With probability at least  $1 - \epsilon$ , there is a  $u$ -coloring of all but at most  $2\lambda\sqrt{n-1}$  vertices.
- With probability at least  $1 - \epsilon$ , these at most  $2\lambda\sqrt{n-1}$  vertices can be properly colored with  $(c\sqrt{n}/\log n)$  more colors, where  $c$  is an absolute constant only depending on  $\epsilon$ .
- With probability at least  $1 - \epsilon$ , by the definition of  $u$ ,  $\chi(G) \geq u$ .

Altogether we have

$$\Pr[u \leq \chi(G) \leq u + c\sqrt{n}/\log n] \geq 1 - 3\epsilon.$$

By setting  $\epsilon \leq 1/300$ , we complete the proof.

### 3.3 References

<http://gxyau.github.io/pdf-documents/ETHZ/Autumn%202018/Probabilistic%20Methods%20in%20Combinatorics/Assignment%2012/solutions12.pdf>