

### 5.2.7 Interpolation in Several Variables and Bicubic Splines

The problem of interpolating in a two-way table is fairly common in many engineering and scientific applications. Specifically if  $f(x, y)$  is a function of two variables that has been tabulated at the points  $(x_i, y_j)$ ,  $0 \leq i \leq n$ ,  $0 \leq j \leq m$ , then we can construct a two-way table that represents this information:

	$x_0$	$x_1$	$\cdots$	$x_n$
$y_0$	$f(x_0, y_0)$	$f(x_1, y_0)$	$\cdots$	$f(x_n, y_0)$
$y_1$	$f(x_0, y_1)$	$f(x_1, y_1)$	$\cdots$	$f(x_n, y_1)$
$\vdots$	$\vdots$			$\vdots$
$y_m$	$f(x_0, y_m)$	$f(x_1, y_m)$	$\cdots$	$f(x_n, y_m)$

If we want to estimate  $f(x, y)$  at a point  $(\hat{x}, \hat{y})$ , which is not in the table, we are faced with an interpolation problem. We will see momentarily that we can construct a polynomial in  $x$  and  $y$  of the form

$$p(x, y) = \sum_{r=0}^n \sum_{s=0}^m a_{rs} x^r y^s,$$

which satisfies  $p(x_i, y_j) = f(x_i, y_j)$ ,  $0 \leq i \leq n$ ,  $0 \leq j \leq m$ . Given this result, we can estimate  $f(x, y)$  at  $(\hat{x}, \hat{y})$  by  $f(\hat{x}, \hat{y}) \approx p(\hat{x}, \hat{y})$ .

A variation of the problem of interpolating in a two-way table is the problem of approximating a function  $f(x, y)$  of two variables; this problem is sometimes called "surface fitting." The graph of the function  $z = f(x, y)$  is a surface in three-space, and we can ask for a simple function  $p(x, y)$  such that  $p(x, y) \approx f(x, y)$ . As an application, if we want the normal to the surface  $z = f(x, y)$  at the point  $(\hat{x}, \hat{y}, f(\hat{x}, \hat{y}))$ , we can estimate the normal by finding the normal to  $z = p(x, y)$ .

If we restrict our approximation problem and consider how we might approximate only functions  $f(x, y)$  that are defined over a *rectangular* region in the  $xy$ -plane, then we can very easily extend our previous results for polynomial and spline approximation. So for simplicity we consider how we can approximate  $f(x, y)$  only for  $(x, y)$  in  $R$  where

$$R = \{(x, y) : a \leq x \leq b, c \leq y \leq d\}.$$

We first treat the problem of interpolating  $f(x, y)$  by polynomials in two variables and then consider the case of bicubic spline approximation. To begin, suppose  $a \leq x_0 < x_1 < \dots < x_n \leq b$  and  $c \leq y_0 < y_1 < \dots < y_m \leq d$ ; and let  $\pi_{nm} = \{(x_i, y_j) : 0 \leq i \leq n, 0 \leq j \leq m\}$ ; that is,  $\pi_{nm}$  is a rectangular grid of  $(n+1)(m+1)$  points in  $R$ . Given a function  $f(x, y)$ , we would like to find a polynomial in two variables,  $p(x, y)$ , such that  $p(x_i, y_j) = f(x_i, y_j)$  for all  $(x_i, y_j)$  in the grid  $\pi_{nm}$ . Now if we define  $\mathcal{P}_{nm}$  by

$$\mathcal{P}_{nm} = \{p(x, y) \mid p(x, y) = \sum_{i=0}^n \sum_{j=0}^m a_{ij} x^i y^j, \text{ for all real } a_{ij}\},$$

then each  $p(x, y)$  has  $(n+1)(m+1)$  coefficients. We then hope that we could choose these coefficients to satisfy the  $(n+1)(m+1)$  interpolation constraints  $p(x_i, y_j) = f(x_i, y_j)$ ,  $0 \leq i \leq n$ ,  $0 \leq j \leq m$ .

Following the lines of one-variable polynomial interpolation in Section 5.2, we define the Lagrange polynomials of two variables by

$$\ell_{ij}(x, y) = \ell_i(x) \tilde{\ell}_j(y), \quad 0 \leq i \leq n, 0 \leq j \leq m$$

where  $\ell_i(x)$  is given in (5.2) and  $\tilde{\ell}_j(y)$  is defined similarly for the points  $c = y_0 < y_1 < \dots < y_m = d$ . Thus we have  $\ell_i(x_k) = \delta_{ik}$  and  $\tilde{\ell}_j(y_k) = \delta_{jk}$ ; and so  $\ell_{ij}(x_r, y_s) = 1$  if  $i = r$  and  $j = s$ ; and  $\ell_{ij}(x_r, y_s) = 0$  otherwise. From this we see that

$$p(x, y) = \sum_{i=0}^n \sum_{j=0}^m f(x_i, y_j) \ell_{ij}(x, y)$$

is a polynomial in  $\mathcal{P}_{nm}$  that interpolates  $f(x, y)$  on the set of points  $\pi_{nm}$ , and we will call this form of  $p(x, y)$  the Lagrange form of the interpolating polynomial.

To see that  $p(x, y)$  is unique in  $\mathcal{P}_{nm}$ , suppose that  $q(x, y)$  in  $\mathcal{P}_{nm}$  satisfies  $q(x_i, y_j) = f(x_i, y_j)$ ,  $0 \leq i \leq n$ ,  $0 \leq j \leq m$ , where  $q(x, y) = \sum_{i=0}^n \sum_{j=0}^m b_{ij} x^i y^j$ . Let us rewrite  $q(x, y)$  as

$$q(x, y) = \sum_{i=0}^n x^i \sum_{j=0}^m b_{ij} y^j = \sum_{i=0}^n x^i q_i(y)$$

where

$$q_i(y) = \sum_{j=0}^m b_{ij} y^j.$$

If we set  $y = y_s$  for a fixed value of  $s$ ,  $0 \leq s \leq m$ , then  $q(x, y_s)$  interpolates  $f(x, y_s)$  at  $x = x_0, x_1, \dots, x_n$ . Thus  $q(x, y_s)$  is a polynomial only in  $x$ , and its coefficients  $q_i(y_s)$  are uniquely determined (see Theorem 5.3) by the data  $f(x_0, y_s), f(x_1, y_s), \dots, f(x_n, y_s)$ .

Knowing that each  $q_i(y_s)$  is determined uniquely by the data  $f(x_i, y_s)$ ,  $0 \leq i \leq n$ ,  $0 \leq s \leq m$ , we can show that the  $b_{ij}$  are also uniquely determined. If we fix  $i$  and consider the system of  $(m + 1)$  equations

$$q_i(y_s) = b_{i0} + b_{i1}y_s + \dots + b_{im}y_s^m, \quad 0 \leq s \leq m,$$

then it is clear that  $b_{i0}, b_{i1}, \dots, b_{im}$  are uniquely determined by  $q_i(y_0), q_i(y_1), \dots, q_i(y_m)$  since the coefficient matrix for the system is a Vandermonde matrix. Thus we have shown that interpolation by  $p(x, y)$  is unique.

For computational purposes, it is frequently useful to express the interpolating polynomial  $p(x, y)$  in the form

$$p(x, y) = \sum_{i=0}^n \ell_i(x) \sum_{j=0}^m f(x_i, y_j) \tilde{\ell}_j(y) = \sum_{i=0}^n \ell_i(x) p_i(y)$$

where

$$p_i(y) = \sum_{j=0}^m f(x_i, y_j) \tilde{\ell}_j(y).$$

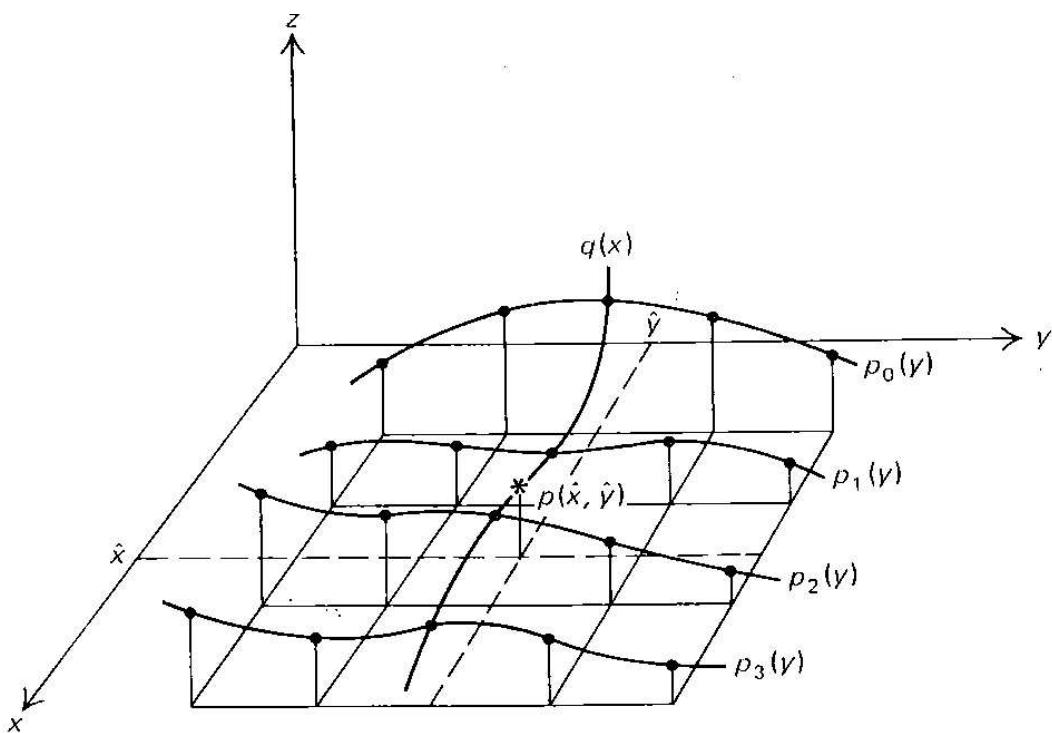
As we noted before,  $p_i(y)$  is a one-variable polynomial (in  $y$ ) interpolating  $f(x, y)$  along the line  $x = x_i$ , at the points  $(x_i, y_0), (x_i, y_1), \dots, (x_i, y_m)$ . This means that  $p(x, y)$  can be built up using one-variable polynomial interpolation. For example,  $p(\hat{x}, \hat{y})$  is given by

$$p(\hat{x}, \hat{y}) = \sum_{i=0}^n \ell_i(\hat{x}) p_i(\hat{y}).$$

From the expression above,  $p(\hat{x}, \hat{y}) = q(\hat{x})$  where  $q(x)$  is the polynomial interpolating the data  $p_0(\hat{y}), p_1(\hat{y}), \dots, p_n(\hat{y})$ , at the points  $x = x_i$ ,  $0 \leq i \leq n$ . To get the values  $p_i(\hat{y})$ , we interpolate the data  $f(x_i, y_j)$ ,  $0 \leq j \leq m$ , by  $p_i(y)$  and evaluate  $p_i(y)$  at  $y = \hat{y}$ .

This two-dimensional interpolation scheme can be rephrased thus.

1. For each fixed grid line  $x = x_i$ , interpolate the data  $f(x_i, y_j)$  in the  $y$ -direction at the knots  $y_0, y_1, \dots, y_m$  and evaluate the (one-variable) interpolating polynomial  $p_i(y)$  at  $y = \hat{y}$ .
2. Interpolate the values  $p_0(\hat{y}), p_1(\hat{y}), \dots, p_n(\hat{y})$  in the  $x$ -direction at the knots  $x_0, x_1, \dots, x_n$  and evaluate the (one-variable) interpolating polynomial  $q(x)$  at  $x = \hat{x}$ . The result,  $q(\hat{x})$ , is equal to  $p(\hat{x}, \hat{y})$ .



**Figure 5.7** A wire-frame model illustrating the calculation  $p(\hat{x}, \hat{y})$ .

This algorithm can be illustrated as a “wire-frame” model as in Fig. 5.7. We also emphasize what is probably obvious: an algorithm to implement two-dimensional interpolation can be constructed from any routine for one-dimensional interpolation. For step (1) above, we call a polynomial interpolator  $(n + 1)$  times to obtain the values  $p_0(\hat{y}), p_1(\hat{y}), \dots, p_n(\hat{y})$  and then call it once more to get  $q(\hat{x}) = p(\hat{x}, \hat{y})$ .

Writing  $p(x, y)$  in the form

$$p(x, y) = \sum_{i=0}^n \sum_{j=0}^m \ell_i(x) \tilde{\ell}_j(y) f(x_i, y_j)$$

makes it clear (see Fig. 5.7) that we could equally well construct one-dimensional interpolators in the  $x$ -direction (at the grid lines  $y = y_j$ ), evaluate these at  $x = \hat{x}$ , and then pass a one-dimensional interpolator through these data along the line  $x = \hat{x}$  in the  $y$ -direction. In addition, from the form of  $p(x, y)$  above, we see how to calculate quantities such as  $p_{xy}(\hat{x}, \hat{y})$ . In particular, the mixed partial  $p_{xy}$  is clearly given by

$$p_{xy}(x, y) = \sum_{i=0}^n \ell'_i(x) \sum_{j=0}^m \tilde{\ell}'_j(y) f(x_i, y_j).$$

Therefore,

$$p_{xy}(\hat{x}, \hat{y}) = \sum_{i=0}^n \ell'_i(\hat{x}) p'_i(\hat{y})$$

where

$$p'_i(\hat{y}) = \sum_{j=0}^m \tilde{\ell}'_j(\hat{y}) f(x_i, y_j).$$

By way of interpretation, we pass  $p_i(y)$  through the data  $f(x_i, y_j)$ ,  $0 \leq j \leq m$  (along the grid line  $x = x_i$ ), differentiate  $p_i(y)$ , and evaluate  $p'_i(\hat{y})$ . Given the values  $p'_i(\hat{y})$ ,  $0 \leq i \leq n$ , we pass  $q(x)$  through these values at the knots  $x_0, x_1, \dots, x_n$ , differentiate  $q(x)$ , and evaluate  $q'(\hat{x})$ . The net result is  $q'(\hat{x}) = p_{xy}(\hat{x}, \hat{y}) = f_{xy}(\hat{x}, \hat{y})$ .

An exactly analogous development can be given for spline interpolation. In particular, a function  $S(x, y)$  is called a *bicubic spline* if  $S(x, y)$  is a cubic spline in  $x$  for fixed  $y$  and a cubic spline in  $y$  for fixed  $x$ . The natural bicubic spline is easiest to treat, and so we restrict our attention to it. As with two-variable polynomial interpolation, we will see that a bicubic spline interpolator can be built up from one-dimensional cubic spline interpolators. The wire-frame representation in Fig. 5.7 is also valid for bicubic splines. We can run cubic spline interpolators along the grid lines  $x = x_i$ , through the data  $f(x_i, y_j)$  for  $0 \leq j \leq m$ , and then evaluate them at  $y = \hat{y}$ . We then run a cubic spline interpolator along  $y = \hat{y}$ , through the values obtained above, and evaluate this cubic spline at  $x = \hat{x}$ .

To see that the construction described above is valid, we can use the idea of a *cardinal* natural spline. Given the knots  $y_0, y_1, \dots, y_m$ , we say that  $S_j(y)$  is a cardinal natural spline if  $S_j(y)$  is a natural cubic spline on knots  $y_0, y_1, \dots, y_m$  and if  $S_j(y_k) = \delta_{jk}$ ,  $0 \leq k \leq m$ . Clearly, these cardinal natural splines always exist; and if  $g(y)$  is a function defined on  $[y_0, y_m]$ , then the natural cubic spline interpolator for  $g(y)$  can be represented as

$$S(y) = \sum_{j=0}^m S_j(y) g(y_j).$$

[Note that the sum of natural cubic splines is a natural cubic spline; so  $S(y)$  is obviously the cubic spline interpolator for  $g(y)$ ; the form of the representation is similar to the Lagrange form of the interpolating polynomial.]

Given the concept of a cardinal natural spline, it is easy to see how to construct a natural bicubic spline  $S(x, y)$ . In particular, if  $S_i(x)$ ,  $0 \leq i \leq n$ , denote the cardinal natural splines for  $x_0, x_1, \dots, x_n$ , and  $\tilde{S}_j(y)$  denote the cardinal natural splines for  $y_0, y_1, \dots, y_m$ , then

$$S(x, y) = \sum_{i=0}^n \sum_{j=0}^m S_i(x) \tilde{S}_j(y) f(x_i, y_j)$$

is a bicubic spline that interpolates  $f(x, y)$  on the grid  $\pi_{nm}$ . As before, we can write  $S(x, y)$  as

$$S(x, y) = \sum_{i=0}^n S_i(x) \phi_i(y)$$

where

$$\phi_i(y) = \sum_{j=0}^m \tilde{S}_j(y) f(x_i, y_j).$$

From the representation of  $S(x, y)$  above, we see that we can evaluate  $S(x, y)$  and various partial derivatives of  $S(x, y)$  by repeatedly calling a one-dimensional spline routine,  $(n + 1)$  times in the  $y$ -direction and then once in the  $x$ -direction. [Note, we do not need to calculate  $S_i(x)$  or  $\tilde{S}_j(y)$ .]

Bicubic spline approximation to functions  $f(x, y)$  are quite effective. A number of applications should be evident; we can use these approximations to estimate mixed partials, estimate normals to a surface, calculate approximations to surface area, etc. Finally, cardinal splines can be used to obtain higher-dimensional analogs of cubic splines. Given  $f(x, y, z)$  to approximate over a solid rectangular region that is gridded by  $(x_i, y_j, z_k)$ ,  $0 \leq i \leq n$ ,  $0 \leq j \leq m$ ,  $0 \leq k \leq \ell$ , we can form

$$S(x, y, z) = \sum_{i=0}^n \sum_{j=0}^m \sum_{k=0}^{\ell} S_i(x) \tilde{S}_j(y) S_k^*(z) f(x_i, y_j, z_k).$$

As before,

$$S(x, y, z) = \sum_{k=0}^{\ell} S_k^*(z) B_k(x, y)$$

where  $B_k(x, y)$  is the bicubic spline interpolator to  $f(x, y, z)$  at the level  $z = z_k$ . Again, the three-dimensional spline approximation to  $f(x, y, z)$  can be built up by calling a one-dimensional spline routine repeatedly. To evaluate  $f(\hat{x}, \hat{y}, \hat{z})$ , we call the routine  $(n + 2)$  times to get  $B_k(\hat{x}, \hat{y})$ , a total of  $(\ell + 1)(n + 2)$  calls. A final call in the  $z$ -direction interpolating  $B_k(\hat{x}, \hat{y})$ ,  $0 \leq k \leq \ell$  will produce  $S(\hat{x}, \hat{y}, \hat{z})$ .