Linked List Operations

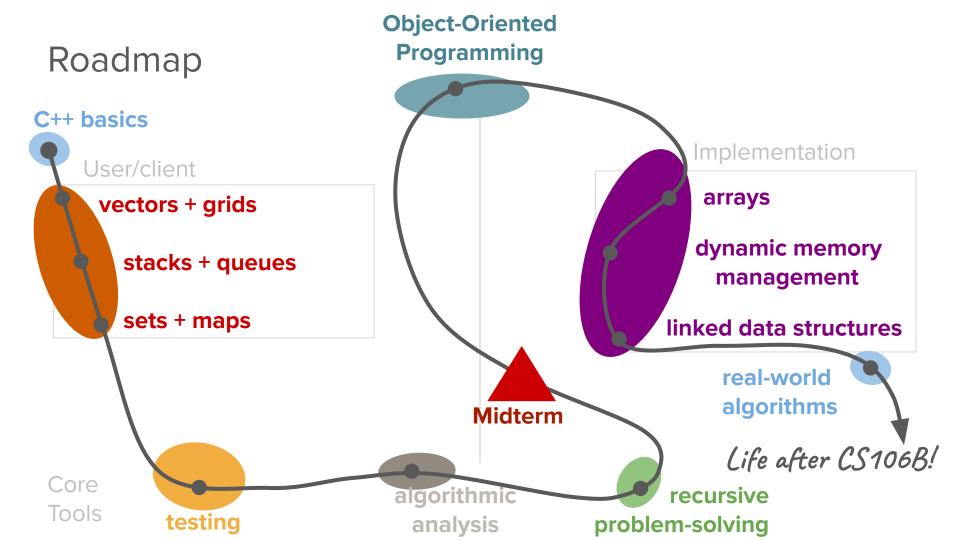
Get creative: If you were explaining a linked list to a friend who isn't in CS106B, what would you use as a good analogy?

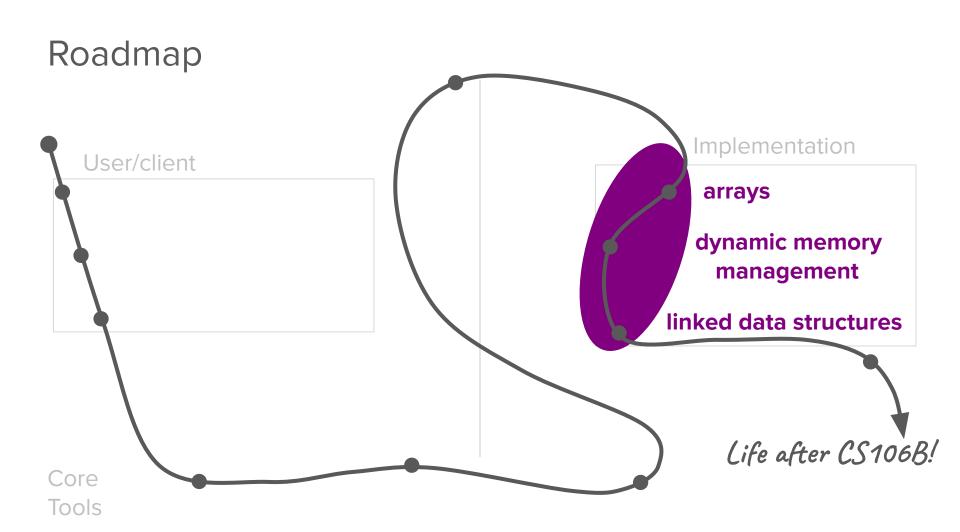
https://pollev.com/cs106bpolls



If you were explaining a linked list to a friend who isn't in CS106B, what would you use as a good analogy?







Today's question

How can we write code to examine and manipulate the structure of linked lists?

Today's topics

1. Review

Advanced Linked List Operations

a. Traversal

b. Insertion (multiple ways!)

c. Deletion (if time)

Review

[intro to linked lists]

Abstract Data What is the interface for the user? **Structures Data Organization** How is our data organized? **Strategies** Fundamental C++ What stores our data? Pointers move **Data Storage** (arrays, linked lists) us across this boundary! How is data represented electronically? Computer (RAM) Hardware

Abstract Data What is the interface for the user? **Structures Data Organization** How is our data organized? **Strategies** Fundamental C++ What stores our data? These are built **Data Storage** (arrays, linked lists) on top of pointers! How is data represented electronically? Computer (RAM) Hardware

What is the interface for the user?

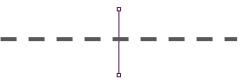
How is our data organized?

What stores our data? (arrays, linked lists)



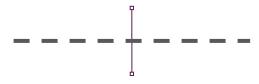
How is data represented electronically? (RAM)

Abstract Data Structures



Data Organization Strategies

Fundamental C++
Data Storage

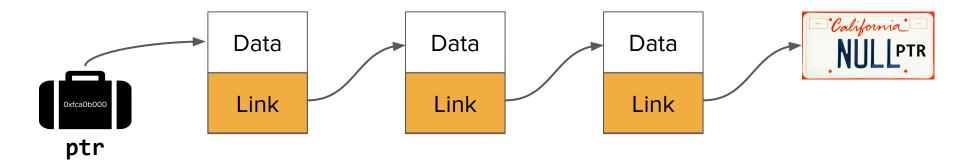


Computer Hardware

What is a linked list?

- A linked list is a chain of nodes, used to store a sequence of data.
- Each node contains two pieces of information:
 - Some piece of data that is stored in the sequence
 - A link to the next node in the list
- We can traverse the list by starting at the first node and repeatedly following its link.
- The end of the list is marked with some special indicator.

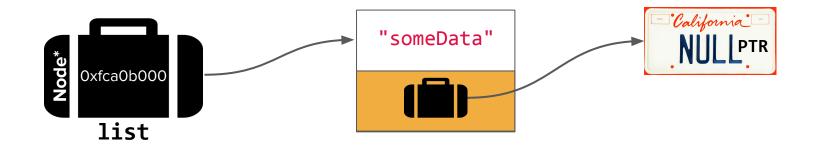
A linked list!



The **Node** struct

```
struct Node {
    string data;
    Node* next;
}
```

Pointer to a node

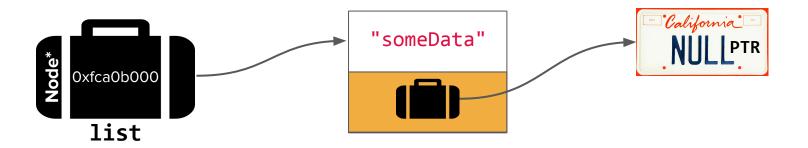


```
Node* list = new Node;
list->data = "someData";
list->next = nullptr;
```

The arrow notation (->) dereferences AND accesses the field for pointers that point to structs specifically.

New: Node struct constructor The Node struct also has a

The Node struct also has a conveniently defined constructor that allows us to accomplish this in one line.



Node* list = new Node("someData", nullptr);

Common linked lists operations

Traversal

How do we walk through all elements in the linked list?

Rewiring

How do we rearrange the elements in a linked list?

Insertion

O How do we add an element to a linked list?

Deletion

O How do we remove an element from a linked list?

Implementing an ADT using a Linked List

- A linked list can be the fundamental data storage backing for an ADT in much the same the same way an array can.
- We saw that linked lists function great as a way of implementing a stack!
- Three operations:
 - push() List insertion and list rewiring
 - o **pop()** List deletion and list rewiring
 - Destructor List traversal and list deletion

Linked list traversal

- Temporary pointers into lists are very helpful!
 - When processing linked lists iteratively, it's common to introduce pointers that point to cells in multiple spots in the list.
 - This is particularly useful if we're destroying or rewiring existing lists.
- Using a while loop with a condition that checks to see if the current pointer is nullptr is the prevailing way to traverse a linked list.

printList()

How does it work?

```
int main() {
   Node* list = readList();
   printList(list);

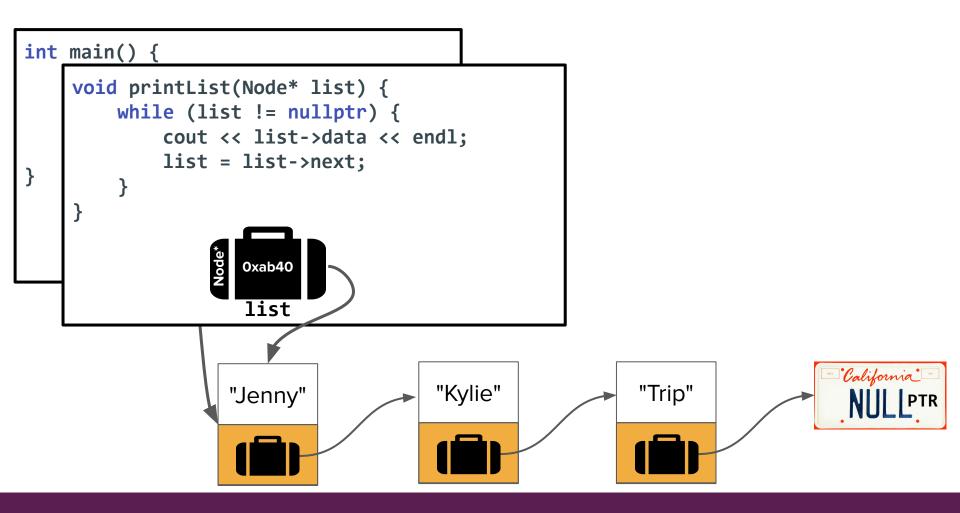
   /* other list things happen... */
}
```

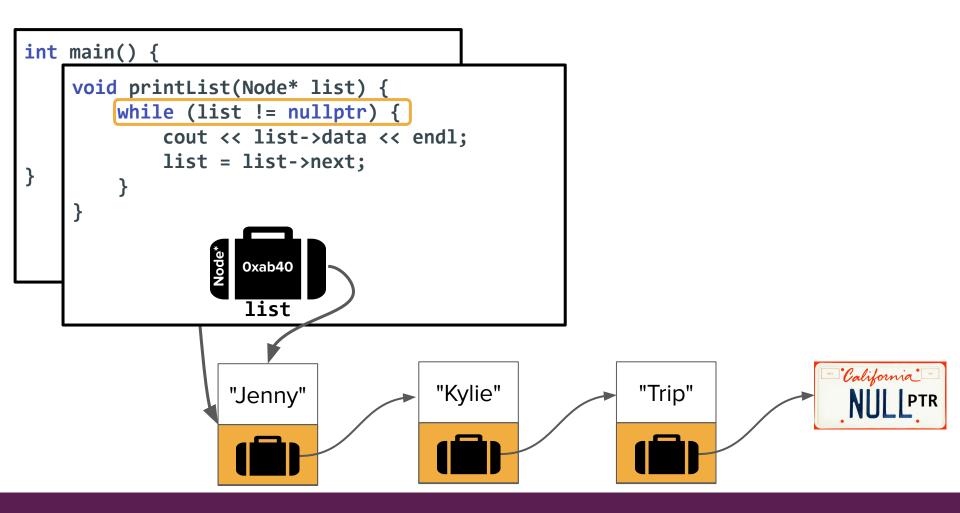
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    Node* list = readList();
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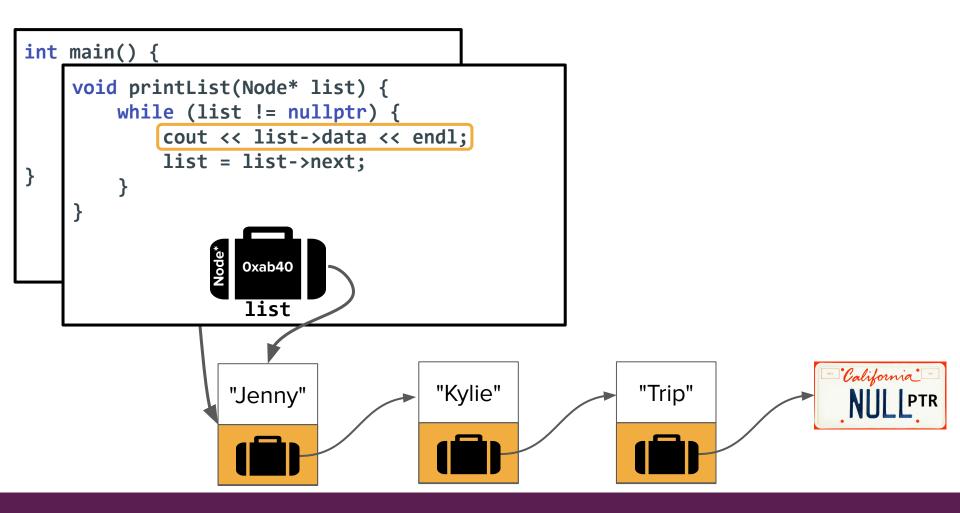
    /* other list things happen... */
}
```

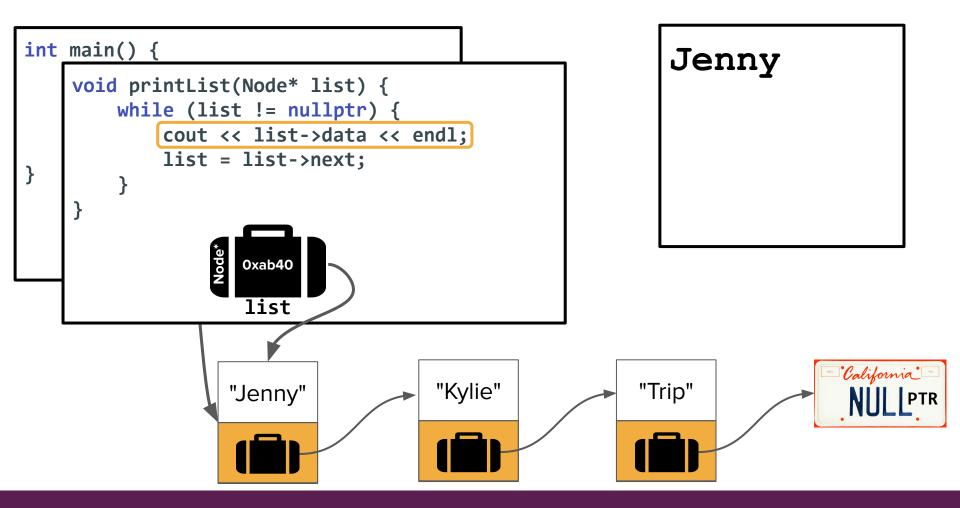
```
int main() {
    Node* list = readList();
    printList(list);
    /* other list things happen... */
      0xab40
       list
                                                                              California
                                       "Kylie"
                                                          "Trip"
                   "Jenny"
```

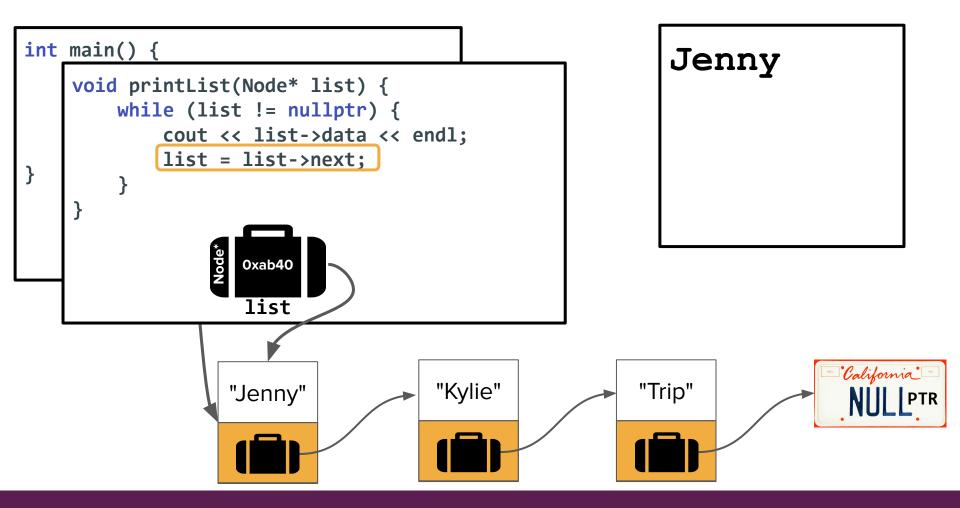
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                                                                             California
                                       "Kylie"
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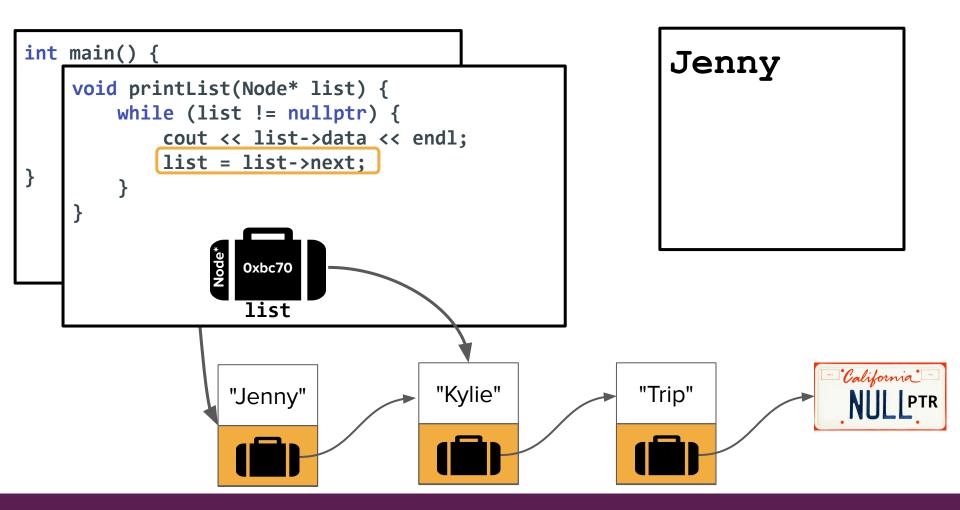


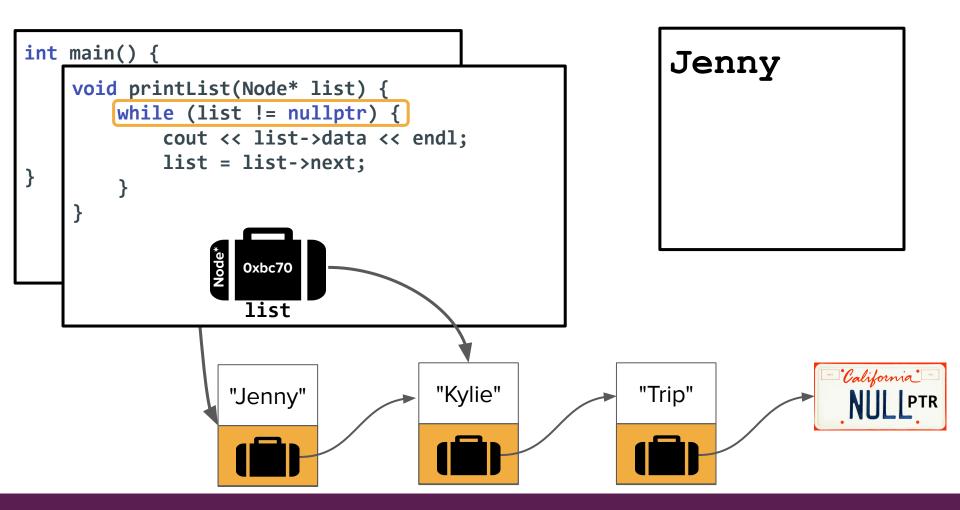


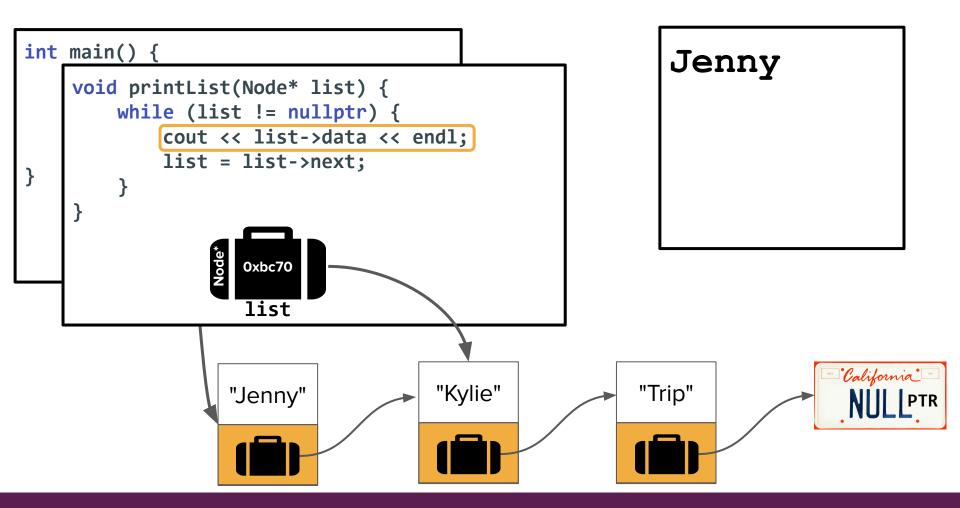


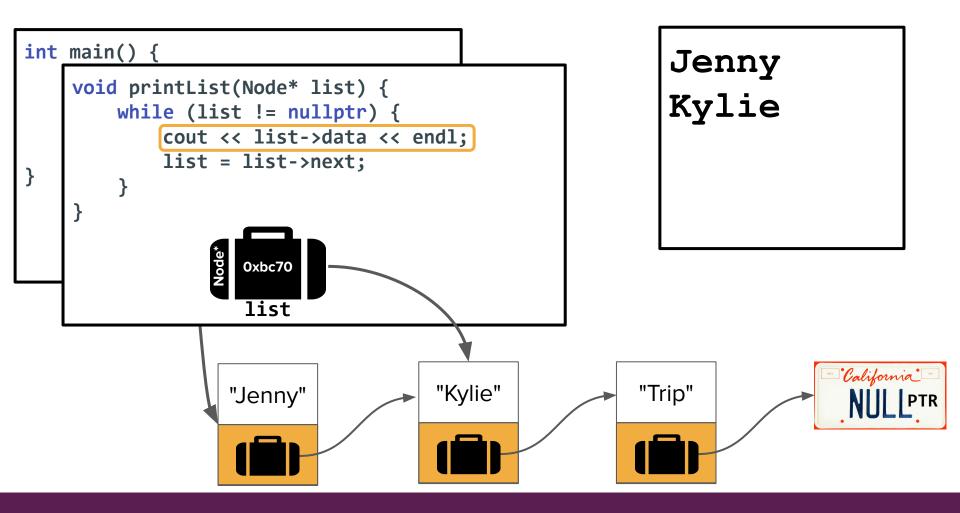


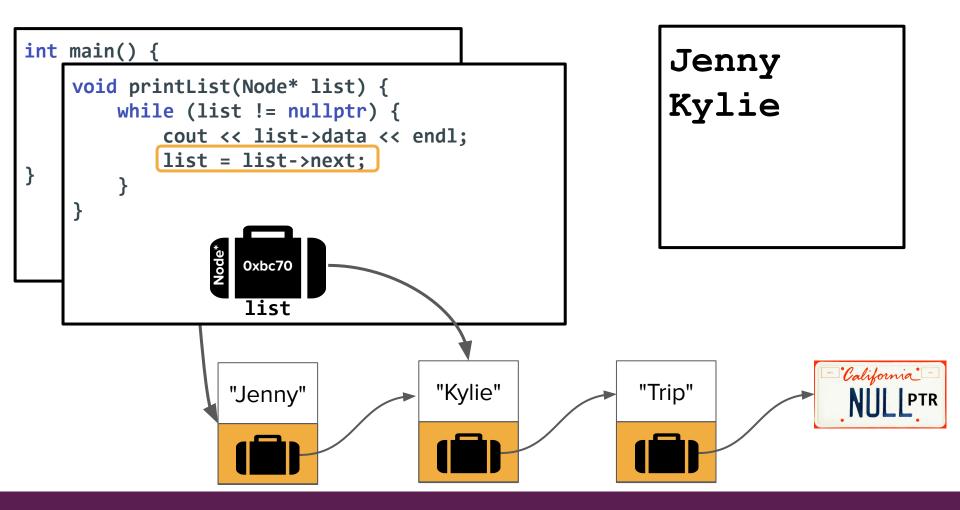


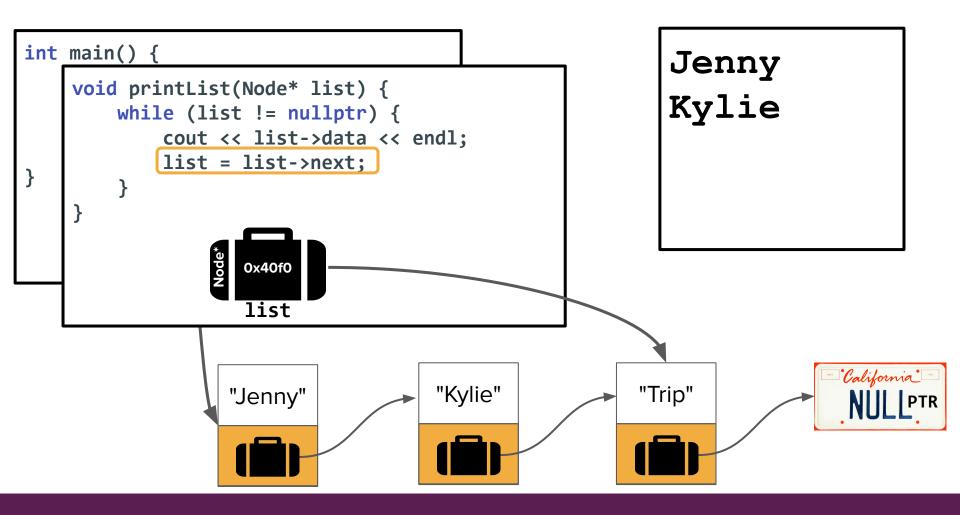


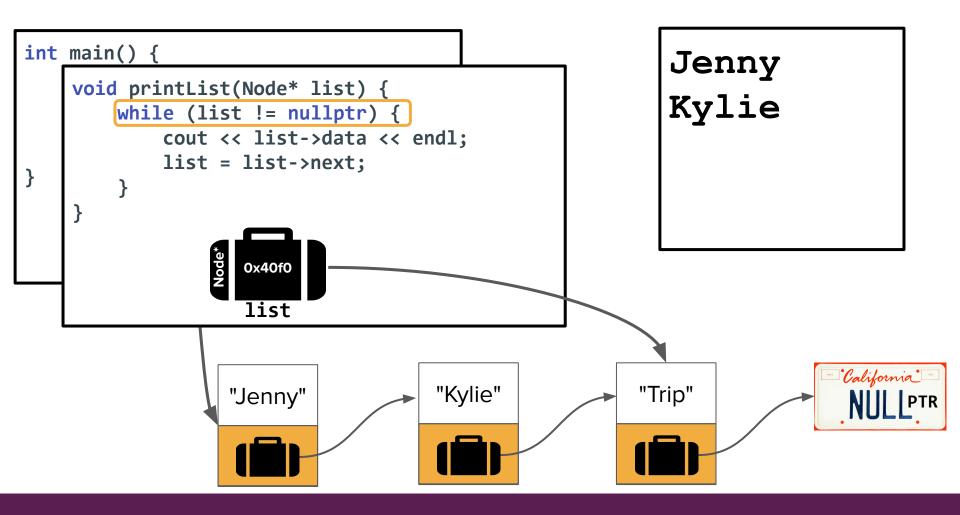


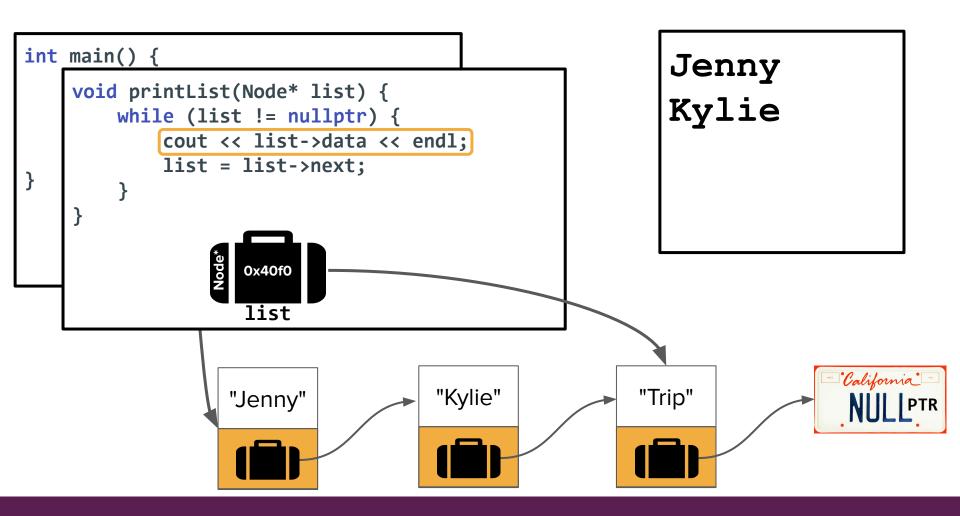


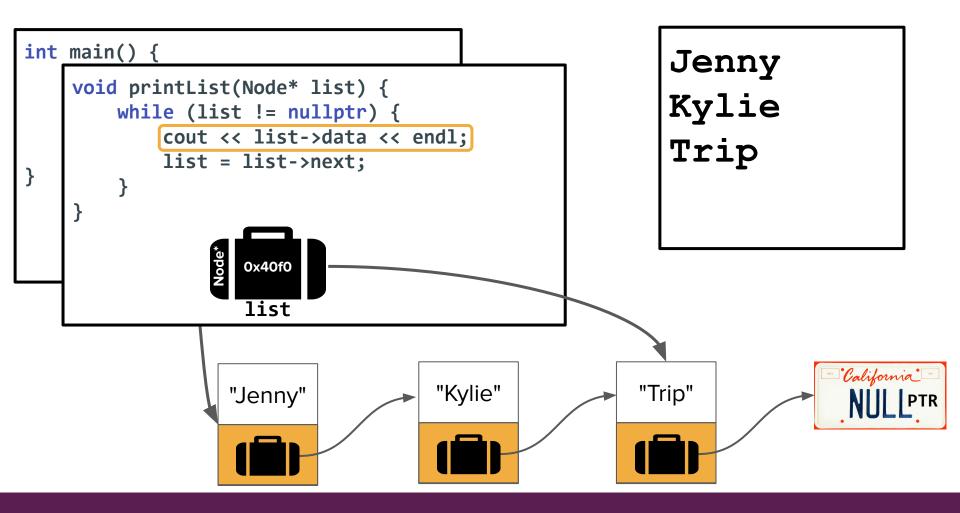


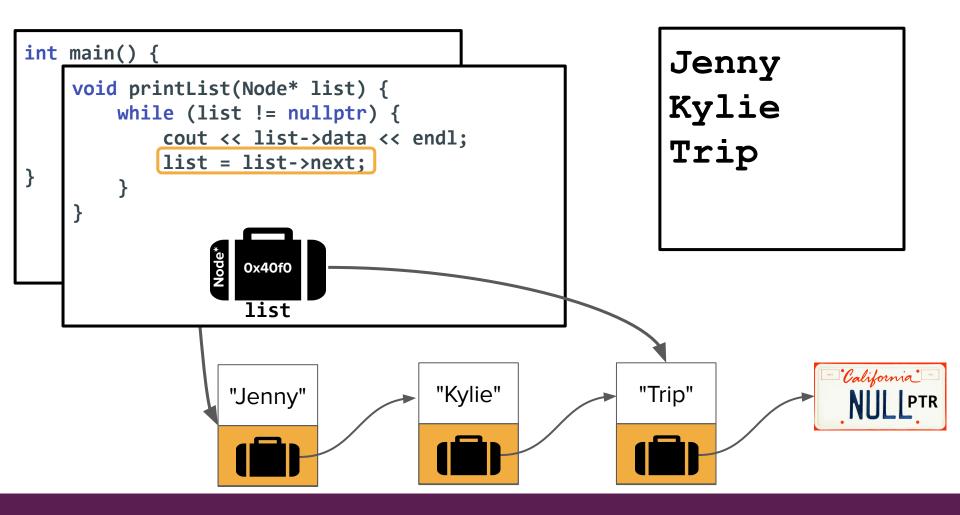


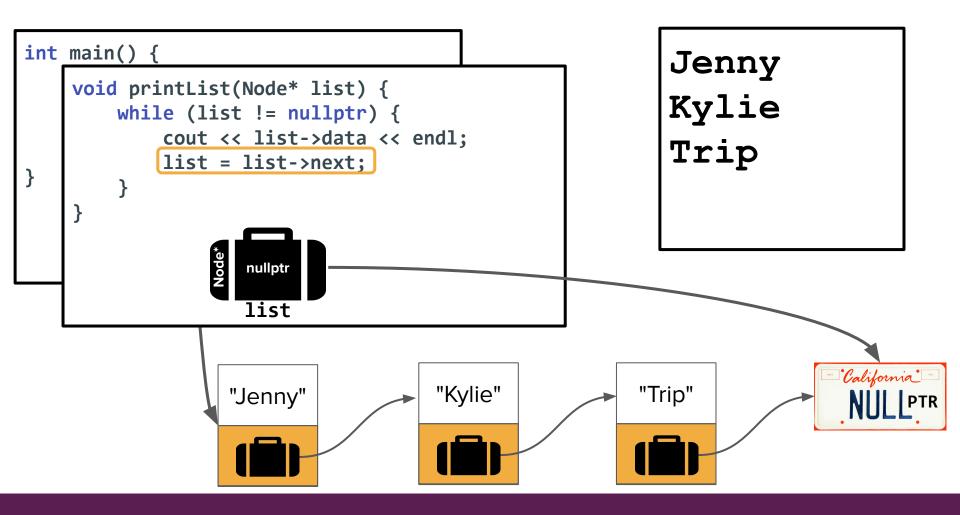


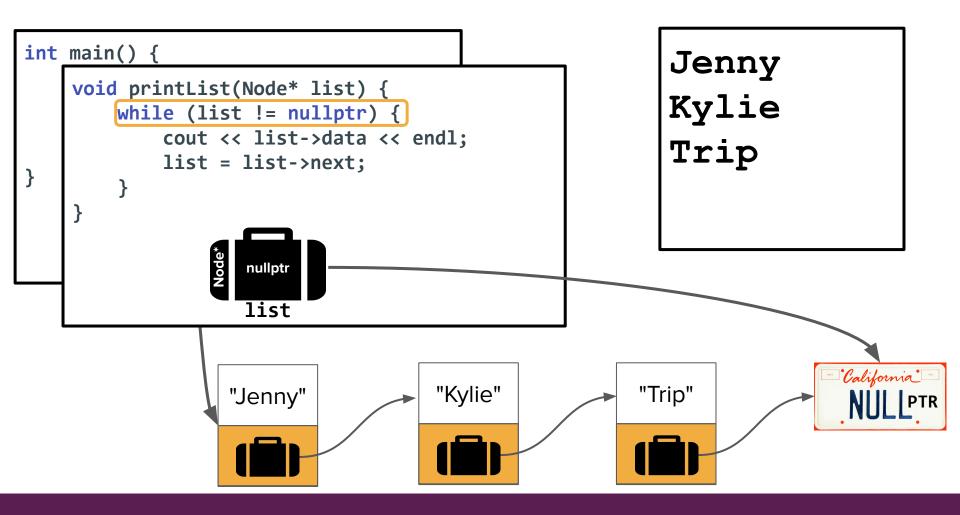












```
int main() {
                                                       Jenny
   Node* list = readList();
                                                       Kylie
   printList(list);
                                                       Trip
   /* other list things happen... */
      0xab40
      list
                                                                      California
                                   "Kylie"
                 "Jenny"
                                                     "Trip"
```

```
int main() {
                                                        Jenny
   Node* list = readList();
                                                        Kylie
    printList(list);
                                                        Trip
    /* other list things happen... */
      0xab40
      list
                                                                       California
                                   "Kylie"
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                                                     "Trip"
```

Summary

- Linked lists are chains of Node structs, which are connected by pointers.
 - Since the memory is not contiguous, they allow for fast rewiring between nodes (without moving all the other Nodes like an array might).

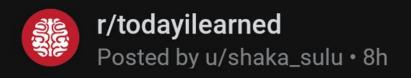
Common traversal strategy

- While loop with a pointer that starts at the front of your list
- Inside the while loop, reassign the pointer to the next node

Common bugs

- Be careful about the order in which you delete and rewire pointers!
- It's easy to end up with dangling pointers or memory leaks (memory that hasn't been deallocated but that you not longer have a pointer to)
- Use nullptr wisely!





TIL a California man got 'NULL' as a personalized license plate hoping that 'NULL' would confuse the computer system. Instead





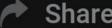


computer system. Instead, when cops left the plate number info empty on a ticket or citation, the fine went to him. He got over \$12k fines sent to him his first year.











How can we write code to examine and manipulate the structure of linked lists?

Linked List Operations Revisited

Common linked lists operations

Traversal

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Linked List Traversal

(revisited)

Traversal utility functions

- Freeing a linked list
- Printing a linked list
- Measuring the length of a list

Measuring a Linked

List

Measuring a Linked List

- Similar to arrays, a linked list does not have the capability to automatically report back its own "size."
- The following code is NOT valid, since list is simply a pointer

```
Node* list = readList();
cout << list.size() << endl; // WRONG! BAD!</pre>
```

 Let's write a function that allows us to calculate the number of nodes in a linked list!

```
int lengthOf() {...}
```

Attendance ticket: https://tinyurl.com/lengthOfList

Please don't send this link to students who are not here. It's on your honor!

lengthOf()

Let's code it!

Linked Lists and Recursion

Rethinking Linked Lists

 On Monday, we mentioned that the Node struct that defined the contents of a linked list was define recursively.

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```
struct Node {
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    Node* next;
}
```

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```
struct Node {
    string data;
    Node* next;
}
```

 This struct definition gives us some insight into the fact that the overall concept of a linked list can be expressed recursively.

A Linked List is Either ...

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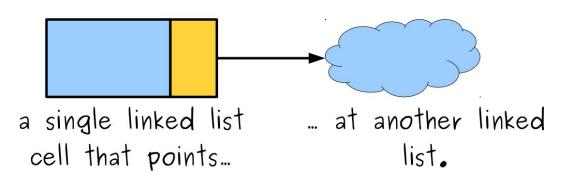
...an empty list,
represented by
nullptr, or...



A Linked List is Either ...

...an empty list,
represented by
nullptr, or...





Printing a List Revisited

Printing a List Revisited

```
void printList(Node* list) {
    while (list != nullptr) {
        cout << list->data << endl;
        list = list->next;
    }
}
```

Printing a List Revisited

```
void printList(Node* list) {
    while (list != nullptr) {
        cout << list->data << endl;
        list = list->next;
    }
}
```

```
void printListRec(Node* list) {
    /* Base Case: There's nothing
to print if the list is empty. */
    if (list == nullptr) return;
    /* Recursive Case: Print the
first node, then the rest of the
list. */
    cout << list->data << endl;</pre>
    printListRec(list->next);
```

Recursion can be a really elegant way to write code for a list traversal!
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- What is the stack frame limit on most computers?
 - You explored this on assignment 3 for most computers it is somewhere in the range of 16-64K
- With a recursive strategy, the size of the list we're able to process is limited by the stack frame capacity we can't process lists longer than 16-64K elements!

Recursion can be a really elegant way to write code for a list traversal!

Note that in the list, better done iteratively! This holds especially true on the assignment – don't try to implement any of the list helper functions

With a recursively!

ie range of 16-64K

every element

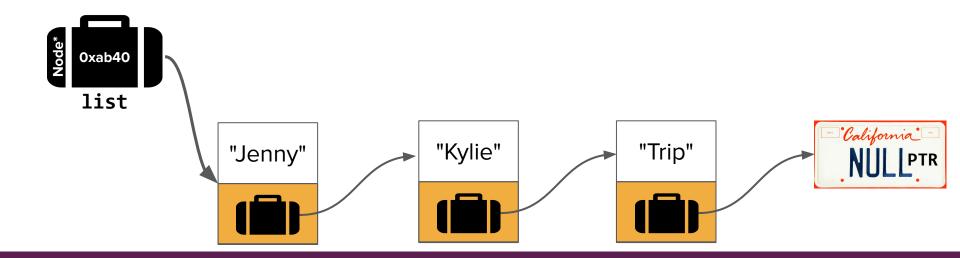
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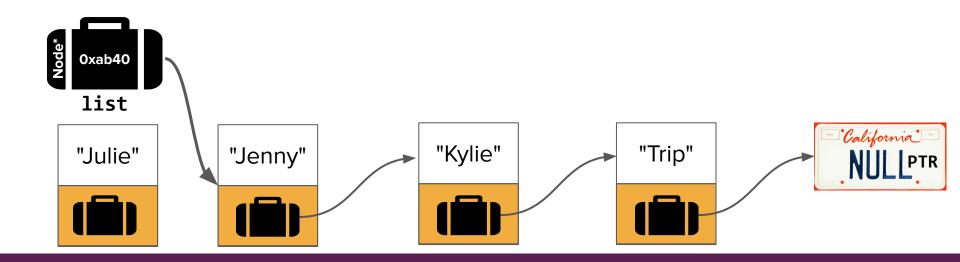
Linked List Traversal Takeaways

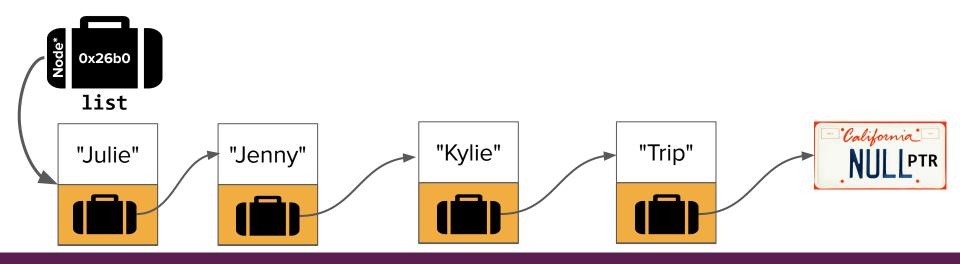
- Using a while loop with a condition that checks to see if the current pointer is nullptr is the prevailing way to traverse a linked list.
- Temporary pointers into lists are very helpful!
 - When processing linked lists iteratively, it's common to introduce pointers that point to cells in multiple spots in the list.
 - This is particularly useful if we're destroying or rewiring existing lists while traversing.
- When traversing but not editing a list, we often pass a pointer parameter by value into a utility function.
- **Iterative traversal** offers the most flexible, scalable way to write utility functions that are able to handle all different sizes of linked lists.

Linked List Insertion

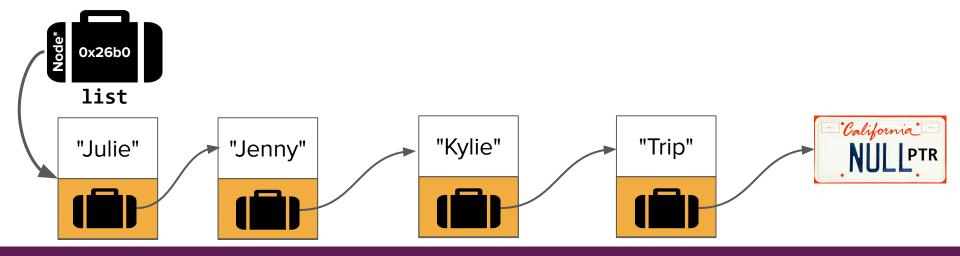
Insertion at the front (prepend)







- Suppose we wanted to write a function to insert an element at the front of a linked list.
- This is similar to the **push()** function we implemented on Monday, but now we're writing a standalone function to do this on an arbitrary list. Let's code it!



prependTo()

Let's code it!

What went wrong?

```
int main() {
   Node* list = nullptr;
   prependTo(list, "Trip");
   prependTo(list, "Kylie");
   prependTo(list, "Jenny");
   return 0;
}
```

```
int main() {
    Node* list = nullptr;
    prependTo(list, "Trip");
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    prependTo(list, "Jenny");
    return 0;
```

```
int main() {
    Node* list = nullptr;
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    prependTo(list, "Jenny");
    return 0;
     nullptr
     list
```

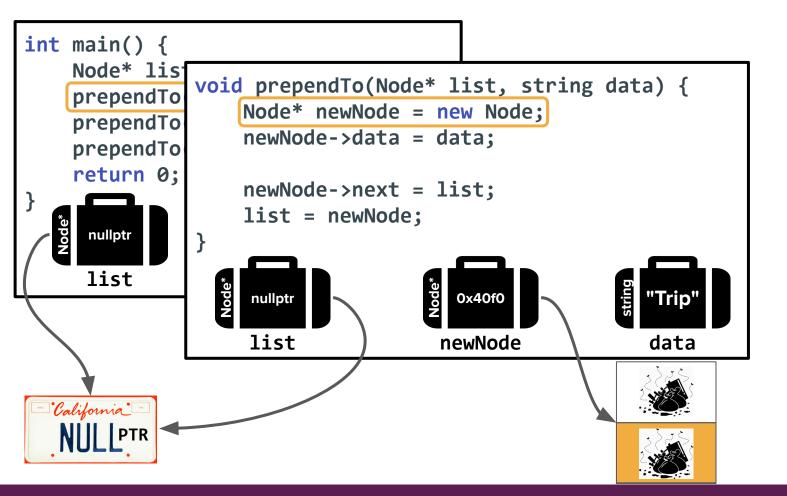


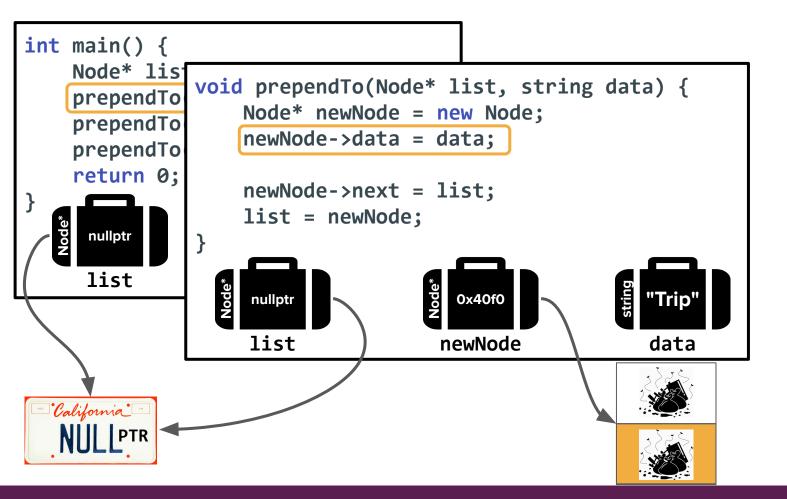
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    return 0;
     nullptr
     list
```

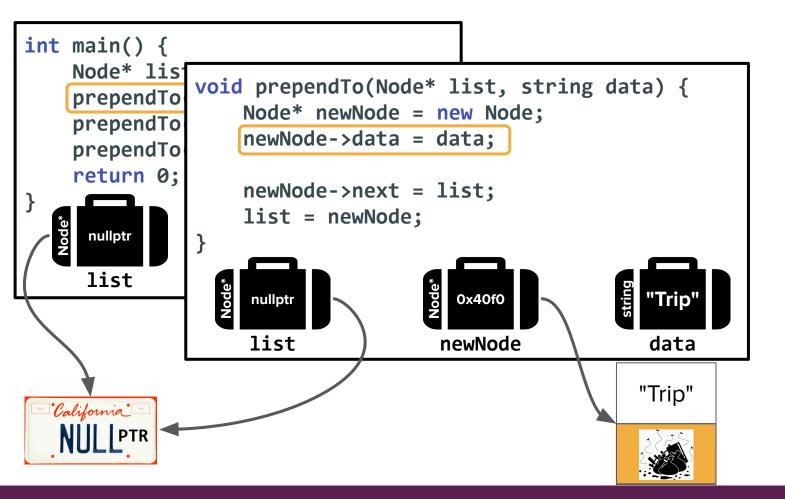
*California ==

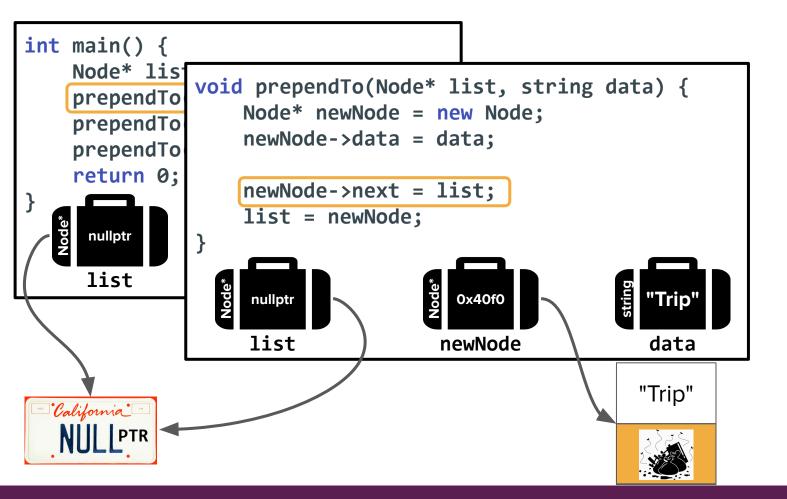
```
int main() {
    Node* list
               void prependTo(Node* list, string data) {
    prependTo
                    Node* newNode = new Node;
    prependTo
                    newNode->data = data;
    prependTo
    return 0;
                    newNode->next = list;
                    list = newNode;
     nullptr
     list
                                                        "Trip"
                    nullptr
                    list
                                                         data
  California
```

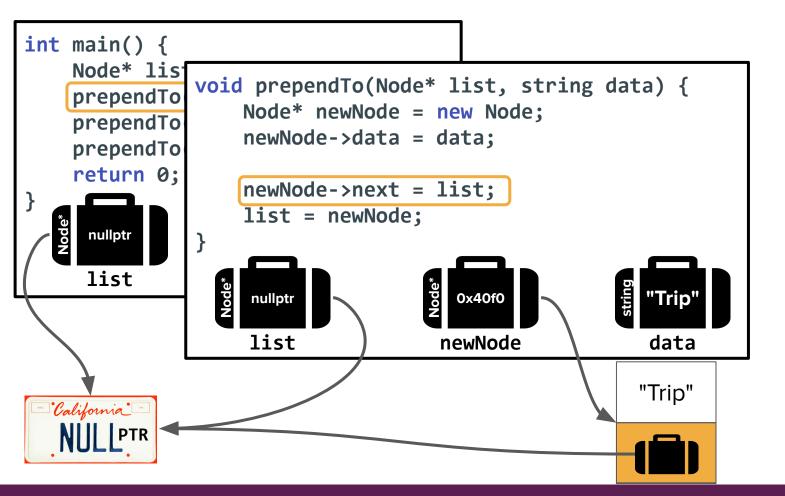
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     nullptr
     list
                                                        "Trip"
                    nullptr
                    list
                                                        data
  California
```

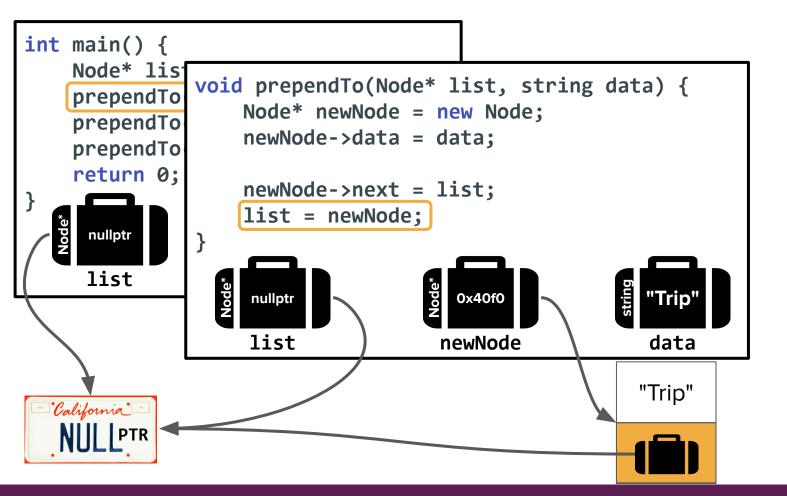


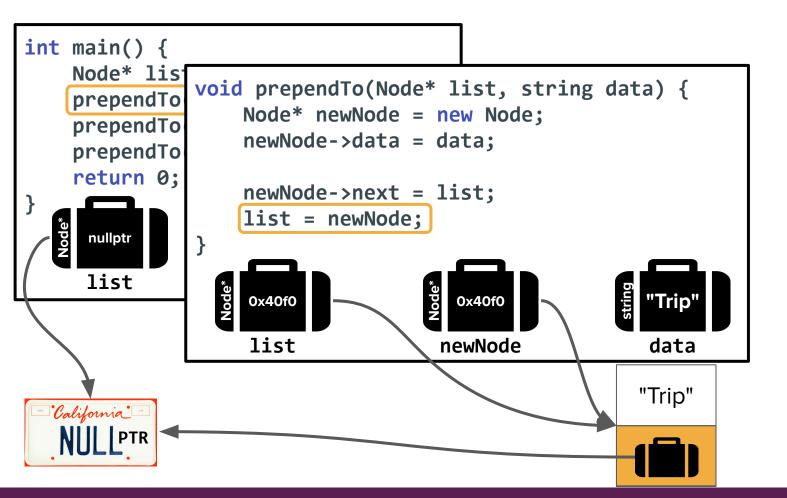












```
int main() {
    Node* list = nullptr;
    prependTo(list, "Trip");
    prependTo(list, "Kylie");
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    return 0;
     nullptr
     list
  *California*
```

"Trip"



```
int main() {
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```

I just got yeeted into the land of leaked memory...



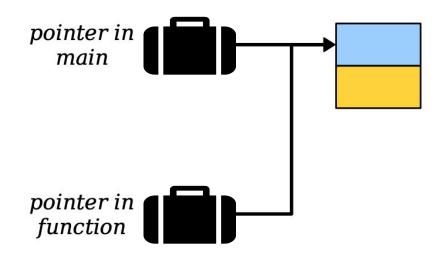
"Trip"





Pointers by Value

- Unless specified otherwise, function arguments in C++ are passed by value – this includes pointers!
- A function that takes a pointer as an argument gets a copy of the pointer.
- We can change where the copy points, but not where the original pointer points.



To solve our earlier problem, we can pass the linked list pointer by reference.

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- Our new function:

```
void prependTo(Node*& list, string data) {
   Node* newNode = new Node;
   newNode->data = data;

   newNode->next = list;
   list = newNode;
}
```

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```
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   Node* newNode = new Node;
   newNode->data = data;

   This is
   newNode->next = list;
   list = newNode;
}
```

This is a reference to a pointer to a Node. If we change where list points in this function, the changes will stick!

```
int main() {
   Node* list = nullptr;
   prependTo(list, "Trip");
   prependTo(list, "Kylie");
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}
```

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```

California

```
int main() {
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    prependTo
                   newNode->data = data;
    prependTo
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                   newNode->next = list;
                   list = newNode;
     nullptr
     list
                                                       "Trip"
                                                       data
```



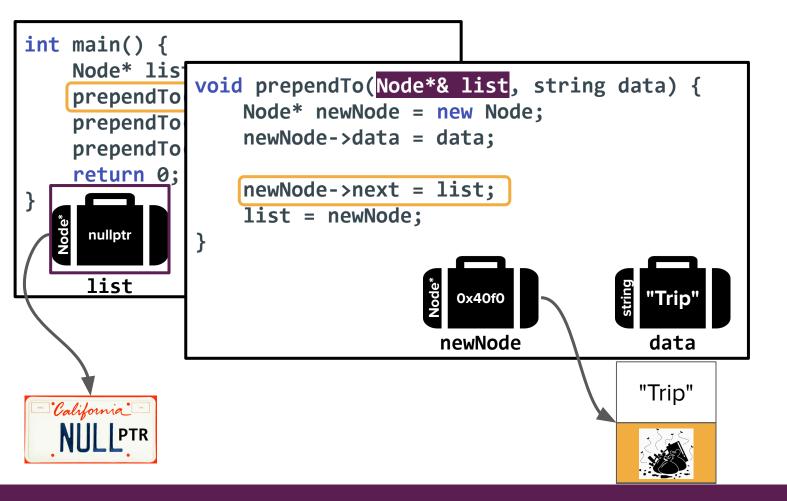
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     nullptr
     list
                                                       "Trip"
                                                       data
```

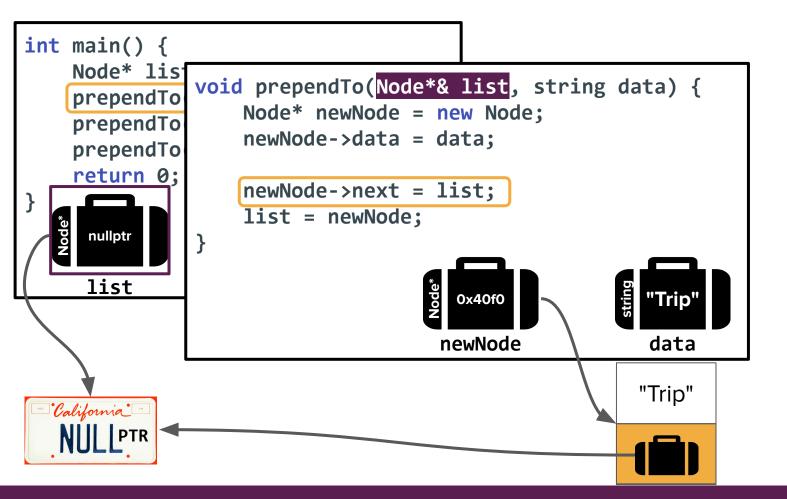


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     nullptr
     list
                                                       "Trip"
                                      0x40f0
                                     newNode
                                                        data
  California
```

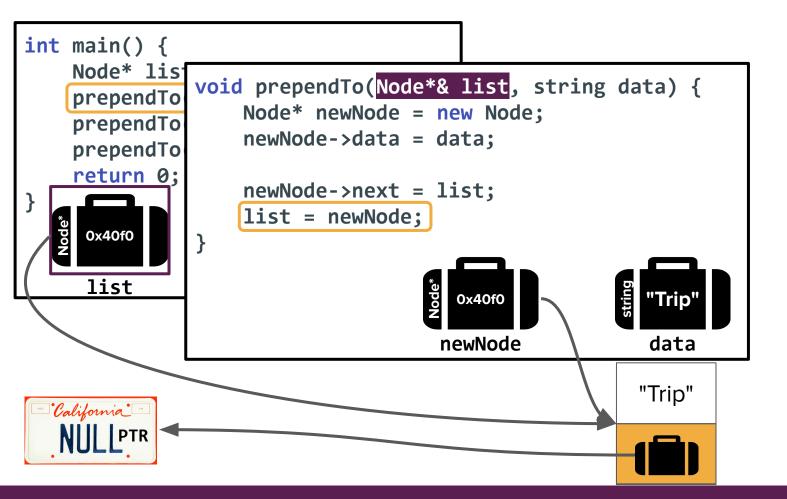
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                                                       "Trip"
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                    newNode->next = list;
                    list = newNode;
     nullptr
     list
                                                        "Trip"
                                       0x40f0
                                     newNode
                                                        data
                                                       "Trip"
  California
```

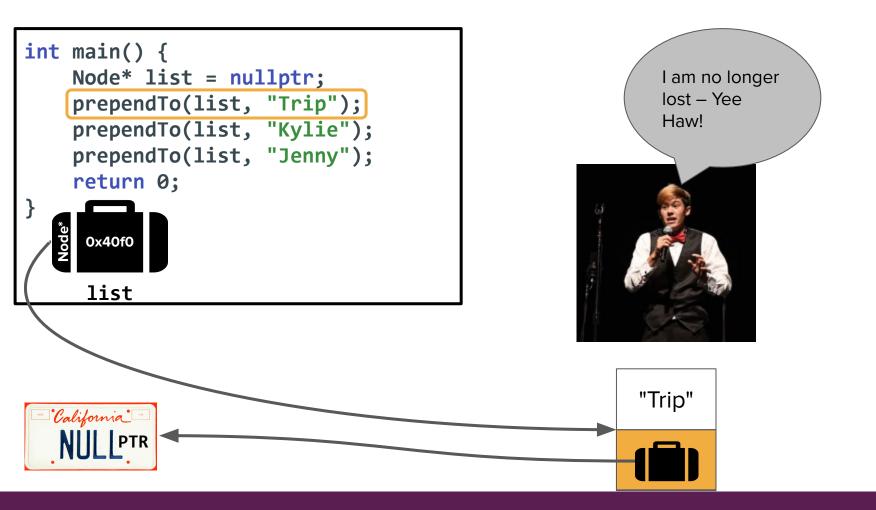




```
int main() {
    Node* list
               void prependTo(Node*& list, string data) {
    prependTo
                    Node* newNode = new Node;
    prependTo
                    newNode->data = data;
    prependTo
    return 0;
                    newNode->next = list;
                    list = newNode;
      nullptr
      list
                                                       string "Trip"
                                       0x40f0
                                      newNode
                                                         data
                                                        "Trip"
  California
```



```
int main() {
    Node* list = nullptr;
    prependTo(list, "Trip");
    prependTo(list, "Kylie");
    prependTo(list, "Jenny");
    return 0;
     0x40f0
     list
                                                      "Trip"
  California
```



Pointers by Reference Summary

- If you pass a pointer into a function by *value*, you can change the contents at the object you point at, but not *which* object you point at.
 - How do you change the contents? Dereferencing!
 - But you don't have access to the original pointer if it's passed by value.

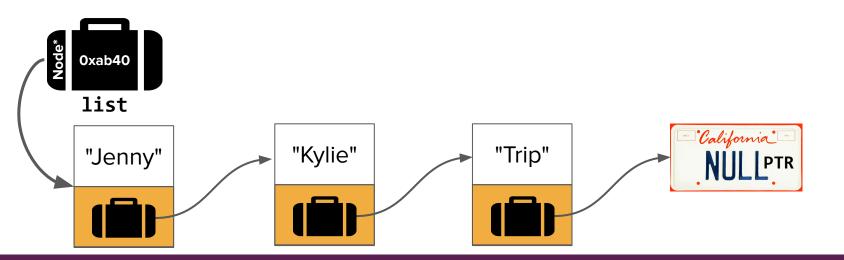
Pointers by Reference Summary

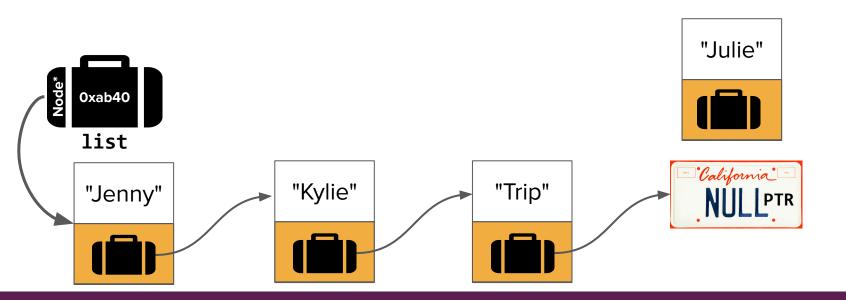
- If you pass a pointer into a function by *value*, you can change the contents at the object you point at, but not *which* object you point at.
- If you pass a pointer into a function by *reference*, you can *also* change *which* object is pointed at.
 - The utility function will now edit the original pointer, just like when we passed ADTs by reference!

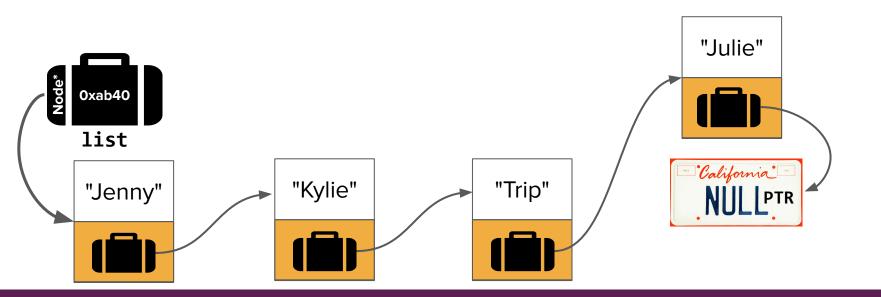
Pointers by Reference Summary

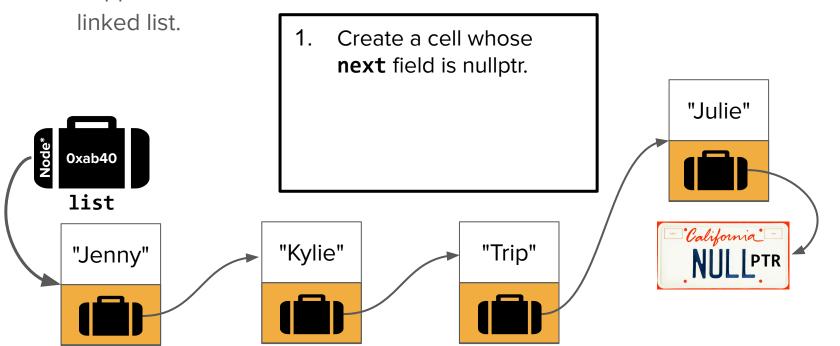
- If you pass a pointer into a function by *value*, you can change the contents at the object you point at, but not *which* object you point at.
- If you pass a pointer into a function by reference, you can also change which object is pointed at.
- When passing in pointers by reference, be careful not to change the pointer unless you really want to change where it's pointing!

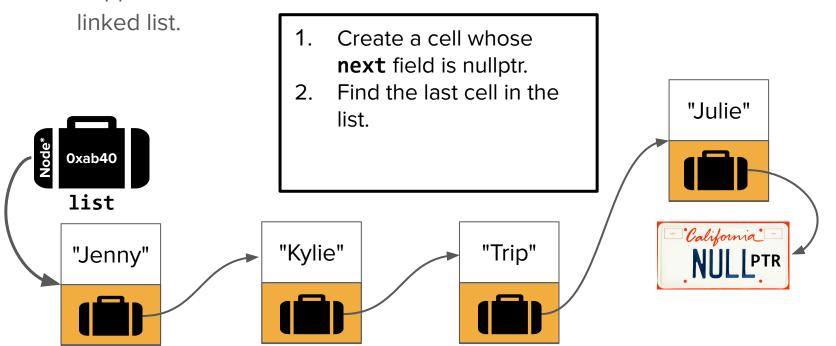
Insertion at the end (append)

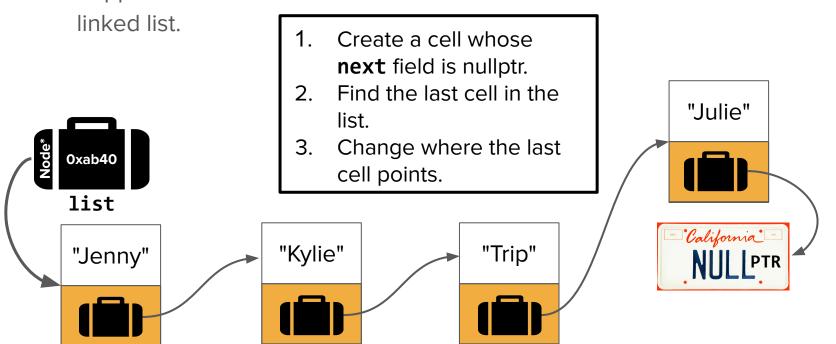












appendTo()

Let's code it!

appendTo() Takeaways

- Appending to the end of a linked list has a lot of tricky edge cases!
 - We must pass the pointer by reference to account for the case where we're adding to an empty list and need to update the head pointer.
 - We have to be careful about our while loop condition to make sure that we never dereference a null pointer!
 - We have to be careful with our usage of pointers by reference and make sure to maintain a local iterator pointer to traverse the list.
- Being able to reason about all of these cases becomes much easier if we draw out diagrams and carefully trace the values of different pointers over time.
 - Note: Check out slides 56-124 of <u>this slide deck</u> for visualizations of the right and wrong ways of coding up the append function!

• What is the big-O complexity of appending to the back of a linked list using our algorithm?

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- Answer: O(n), where n is the number of elements in the list, since we have to find the last position each time.

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- Answer: O(n), where n is the number of elements in the list, since we have to find the last position each time.
- This seems suspect O(n) for a single insertion is pretty bad! Can we do better?
 - Find out after the break!

Summary

- Linked lists can be used outside classes you'll do this on Assignment 5!
- Think about when you want to pass pointers by reference in order to edit the original pointer and to avoid leaking memory.
- We can add to a linked list by either prepending or appending.
 - Prepending is faster but results in a reversed order of items (things added earlier are at the back of the list)
 - Appending (as we've learned so far) requires traversing all items but maintains order (things added earlier are at the front of the list)

Announcements

Announcements

- Assignment 4 is due today. Assignment 3 revisions are due Friday, July 29.
- Assignment 5 will be released by the end of the day.
 - YEAH Hours will be tomorrow on Wednesday, July 27 at 5pm in Hewlett 103.
- Lecture this Thursday will be open project work time!
 - Jenny and I will be having open OH in NVIDIA if you have questions.
 - We'll also set up areas in the lecture hall where you can discuss projects by topic so you can get feedback and ideas from your classmates.
 - Attendance is optional.
- The deadline to change your grading basis is this Friday, July 29 at 5pm.

Linked List Insertion

(continued)

Two ways to add (so far)

- Insertion at the front: prependTo()
 - Prepending is faster but results in a reversed order of items (things added earlier are at the back of the list)
- Insertion at the back: appendTo()
 - Appending requires traversing all items but maintains order (things added earlier are at the front of the list)

```
void nameOfAddFunction(Node*& list, string data) {
    ...
}
```

- What is the big-O complexity of appending to the back of a linked list using our algorithm?
- Answer: O(n), where n is the number of elements in the list, since we have to find the last position each time.
- This seems suspect O(n) for a single insertion is pretty bad! Can we do better?

A more efficient append

A more efficient appendTo()

- Earlier, we saw an O(n) appendTo() that added to the back of a linked list. We can do better!
- What if we know we're going to add many things in some maintained order?
- Specifically, we'll use the example of adding items from a vector into linked list.

Attempt #1

```
Node* createListWithAppend(Vector<string> values) {
   if (values.isEmpty()) {
       return nullptr;
   Node* head = new Node(values[0], nullptr);
   for (int i = 1; i < values.size(); i++) {</pre>
       appendTo(head, values[i]);
   return head;
```

Attempt #1: What's the runtime? (poll)

```
Node* createListWithAppend(Vector<string> values) {
   if (values.isEmpty()) {
       return nullptr;
   Node* head = new Node(values[0], nullptr);
                                                       O(N)
                                                   B. O(N^2)
                                                   C. O(N^3)
   for (int i = 1; i < values.size(); i++) {</pre>
                                                       O(\log N)
       appendTo(head, values[i]);
   return head;
```

https://pollev.com/cs106bpolls

What's the runtime?

O(N)

O(N^2)

 $O(N^3)$

O(logN)



Attempt #1: What's the runtime? (poll)

```
Node* createListWithAppend(Vector<string> values) {
   if (values.isEmpty()) {
       return nullptr;
   Node* head = new Node(values[0], nullptr);
                                                      O(N)
   for (int i = 1; i < values.size(); i++) {</pre>
       appendTo(head, values[i]);
   return head;
```

Attempt #2: createListWithTailPtr() Let's code it!

How does it work?

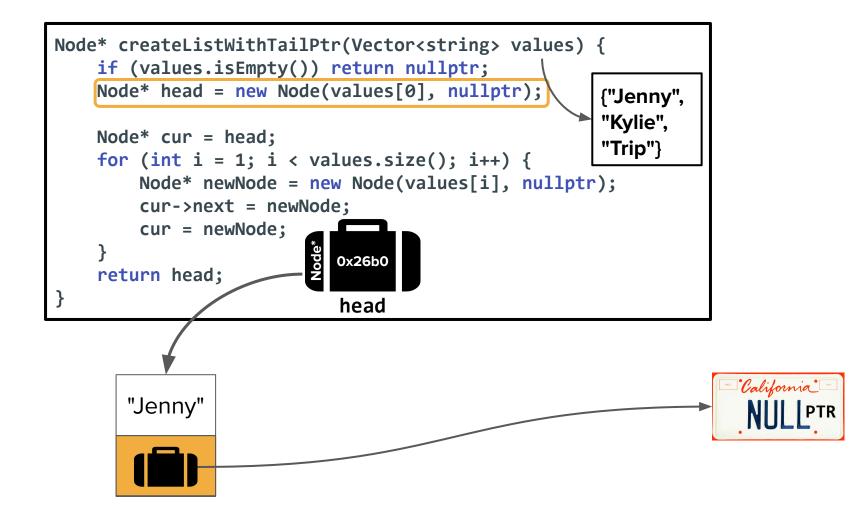
```
int main() {
    Vector<string> values = {"Jenny", "Kylie", "Trip"};
    Node* list = createListWithTailPtr(values);

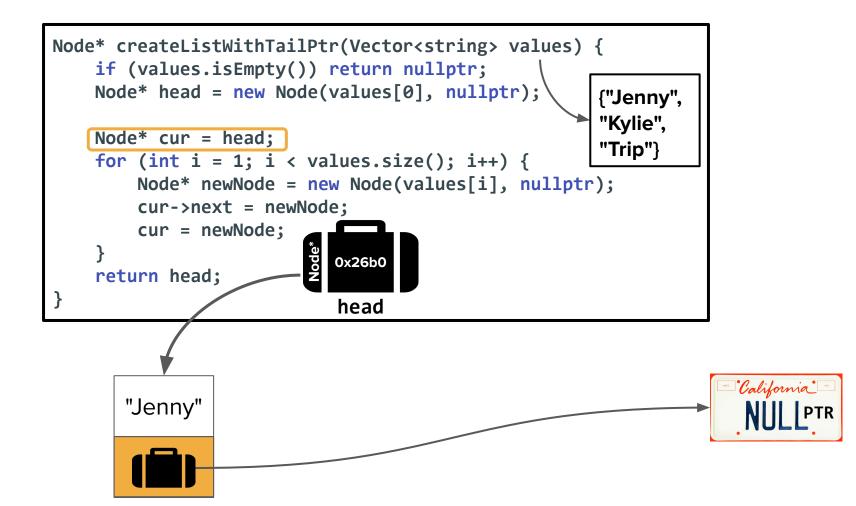
/* Do other list-y things here, like printing/freeing the list. */
    return 0;
}
```

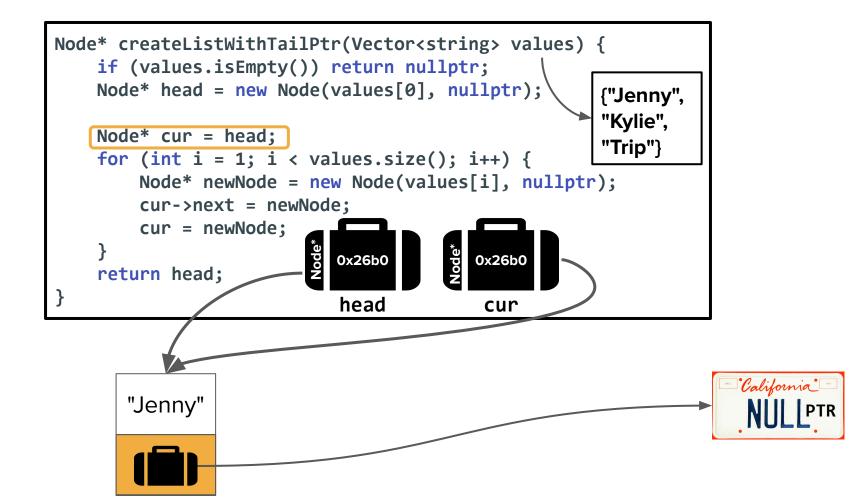
```
Node* createListWithTailPtr(Vector<string> values) {
    if (values.isEmpty()) return nullptr;
    Node* head = new Node(values[0], nullptr);
                                                      {"Jenny",
                                                      "Kylie",
    Node* cur = head;
                                                      "Trip"}
    for (int i = 1; i < values.size(); i++) {</pre>
        Node* newNode = new Node(values[i], nullptr);
        cur->next = newNode;
        cur = newNode;
    return head;
```

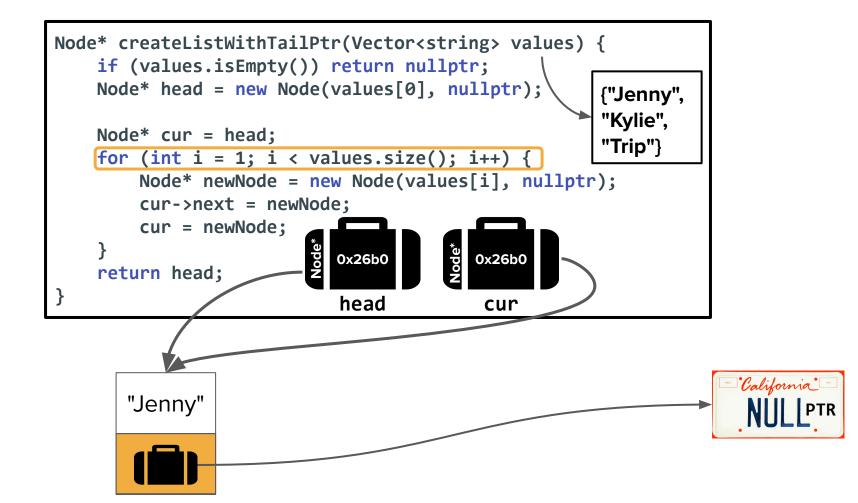
```
Node* createListWithTailPtr(Vector<string> values) {
    if (values.isEmpty()) return nullptr;
    Node* head = new Node(values[0], nullptr);
                                                      {"Jenny",
                                                      "Kylie",
    Node* cur = head;
                                                      "Trip"}
    for (int i = 1; i < values.size(); i++) {</pre>
        Node* newNode = new Node(values[i], nullptr);
        cur->next = newNode;
        cur = newNode;
    return head;
```

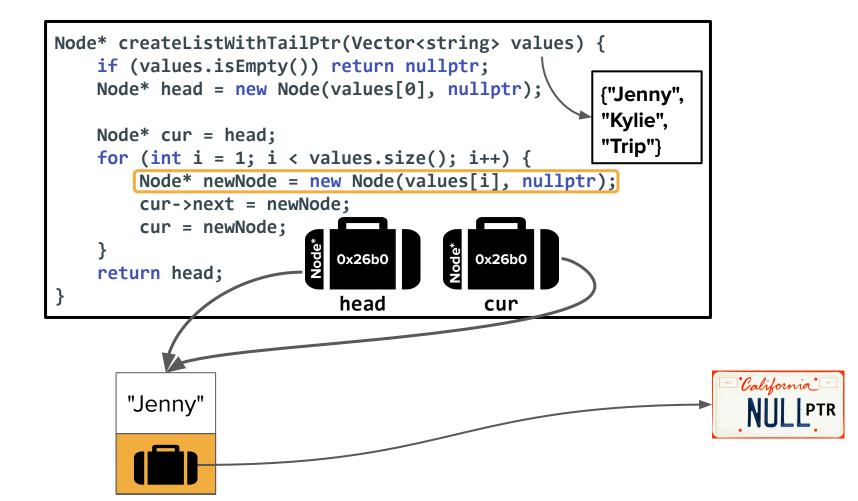
```
Node* createListWithTailPtr(Vector<string> values) {
    if (values.isEmpty()) return nullptr;
    Node* head = new Node(values[0], nullptr);
                                                      {"Jenny",
                                                      "Kylie",
    Node* cur = head;
                                                      "Trip"}
    for (int i = 1; i < values.size(); i++) {</pre>
        Node* newNode = new Node(values[i], nullptr);
        cur->next = newNode;
        cur = newNode;
    return head;
```

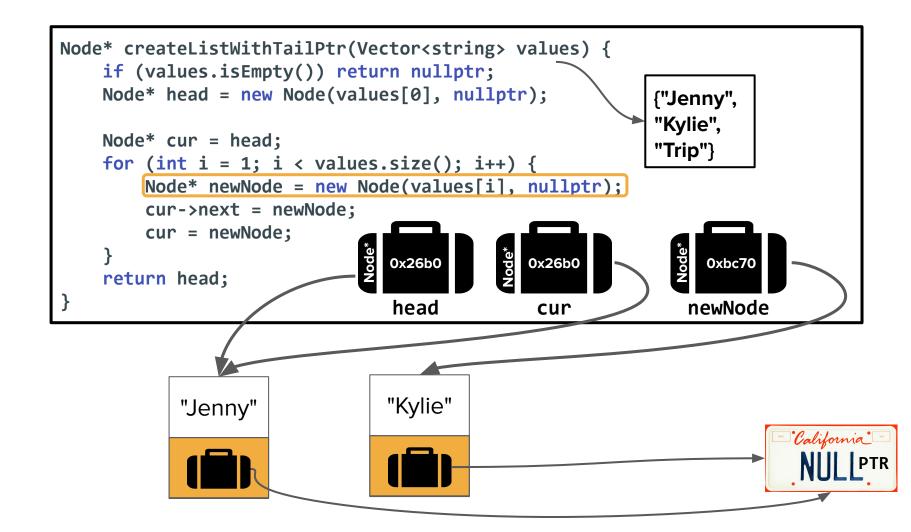


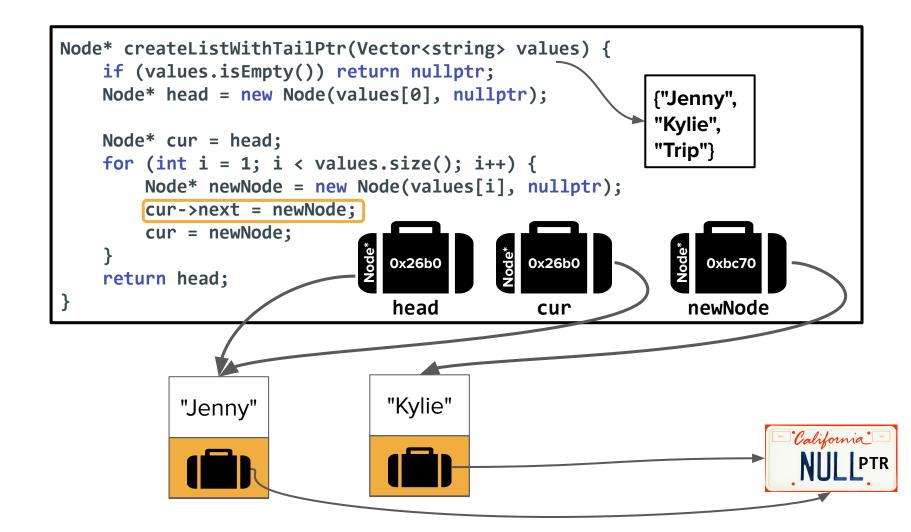


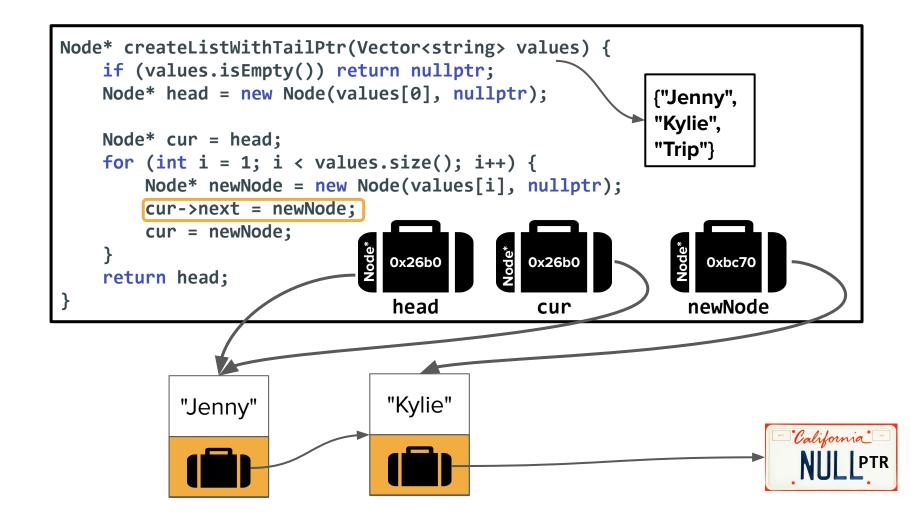


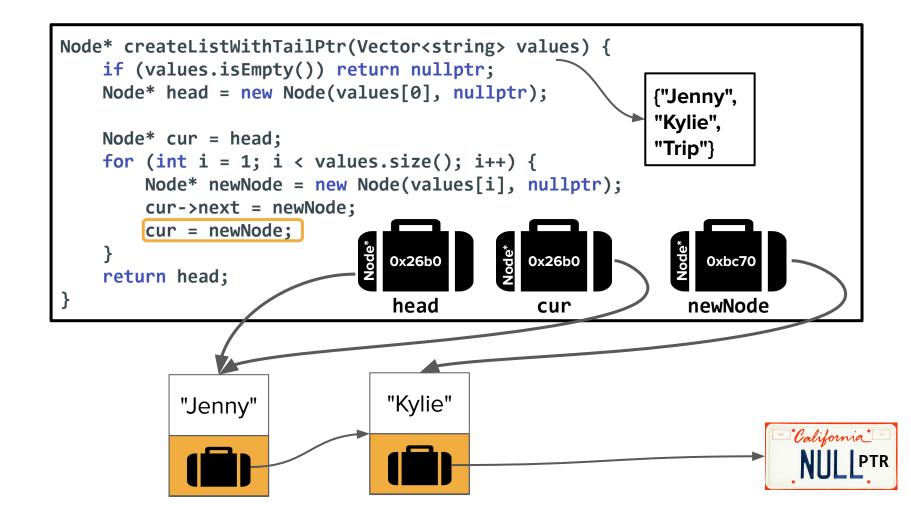


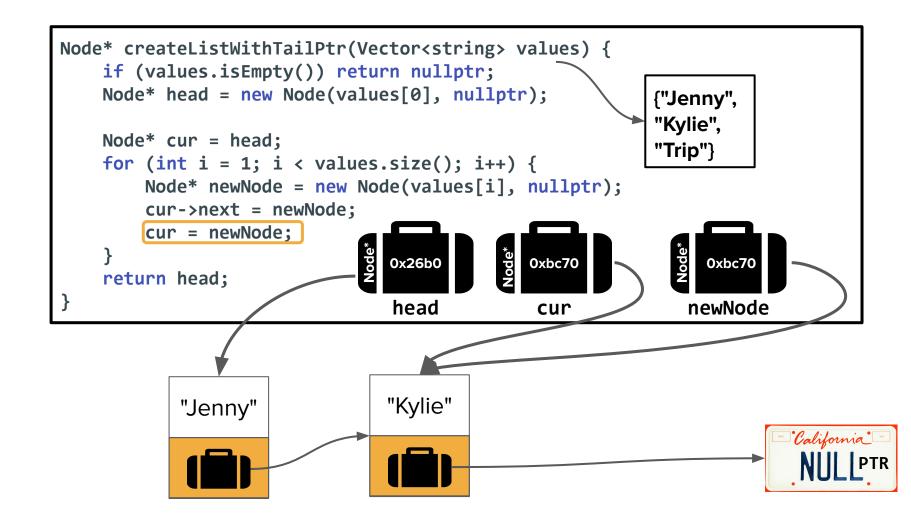


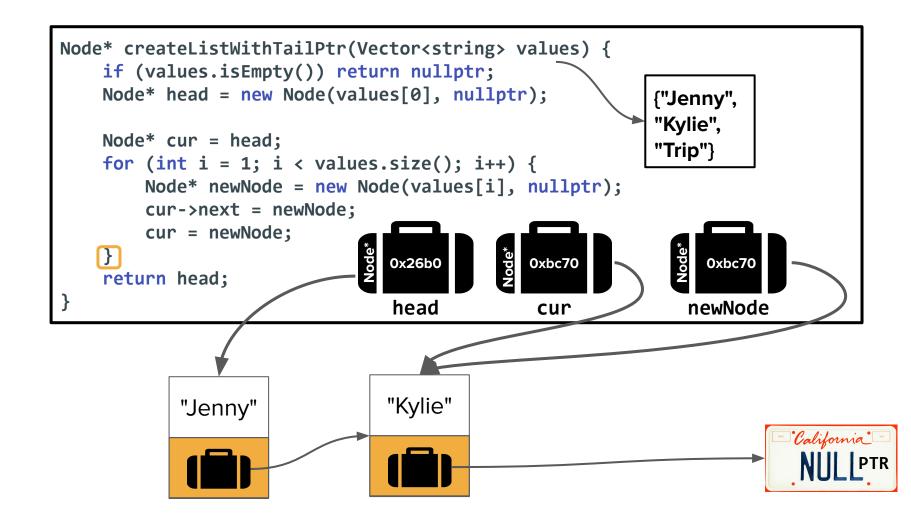


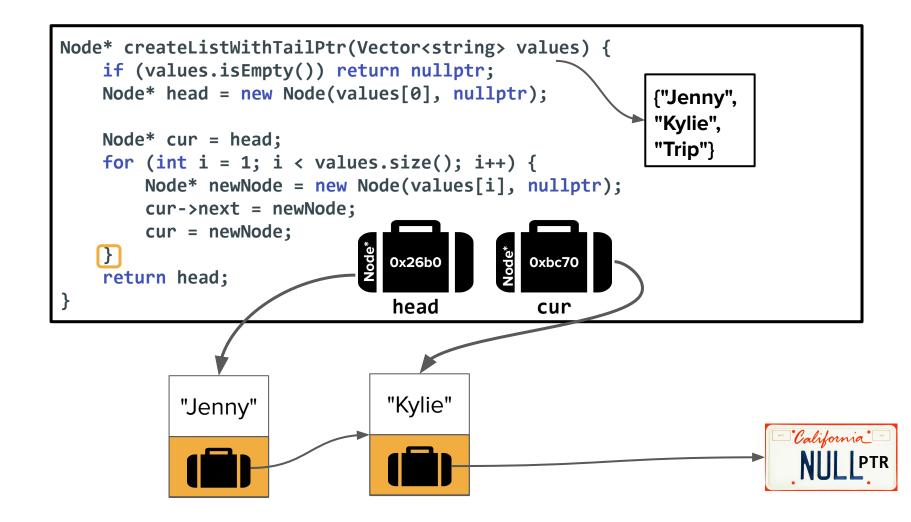


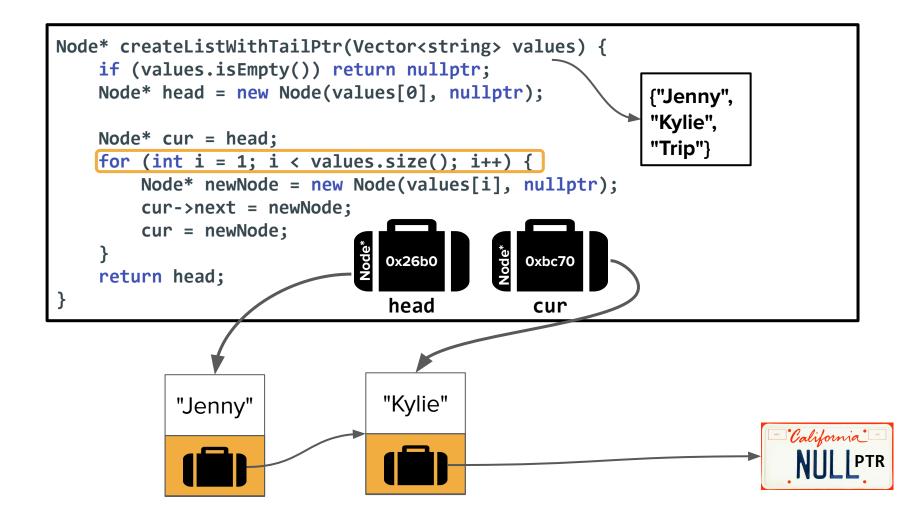


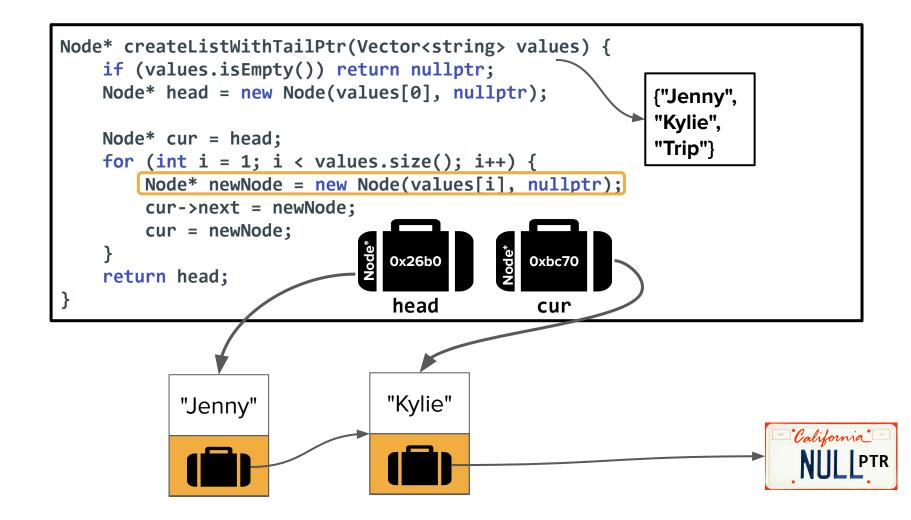


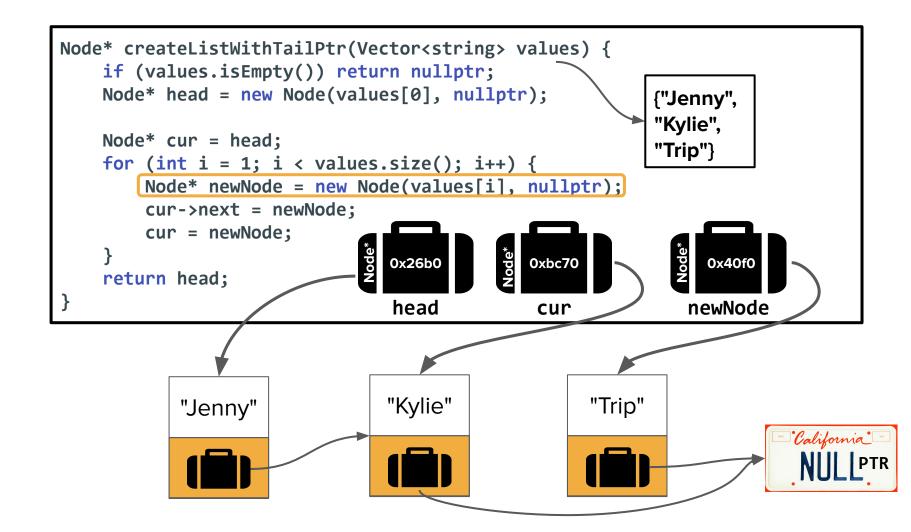


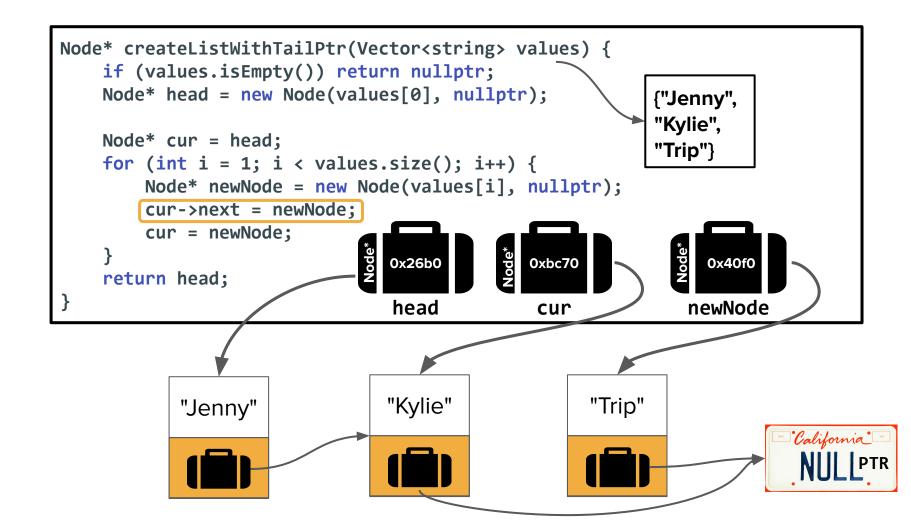


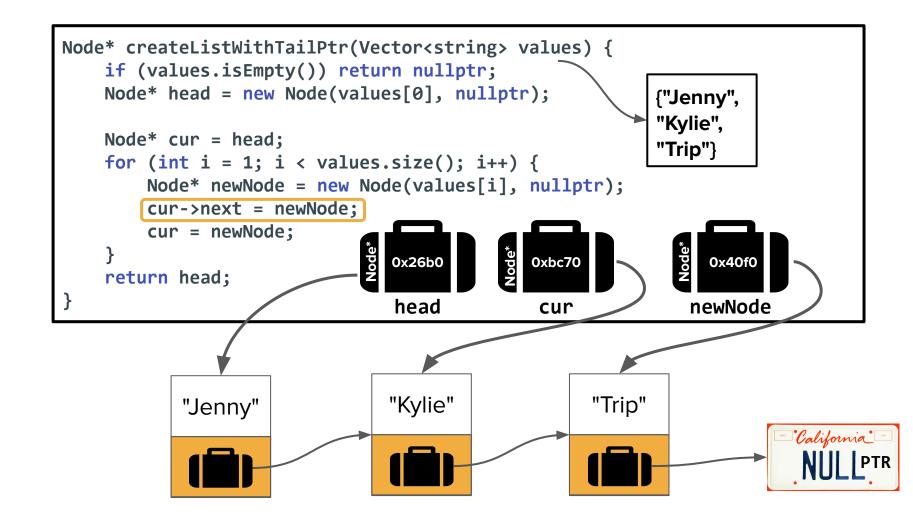


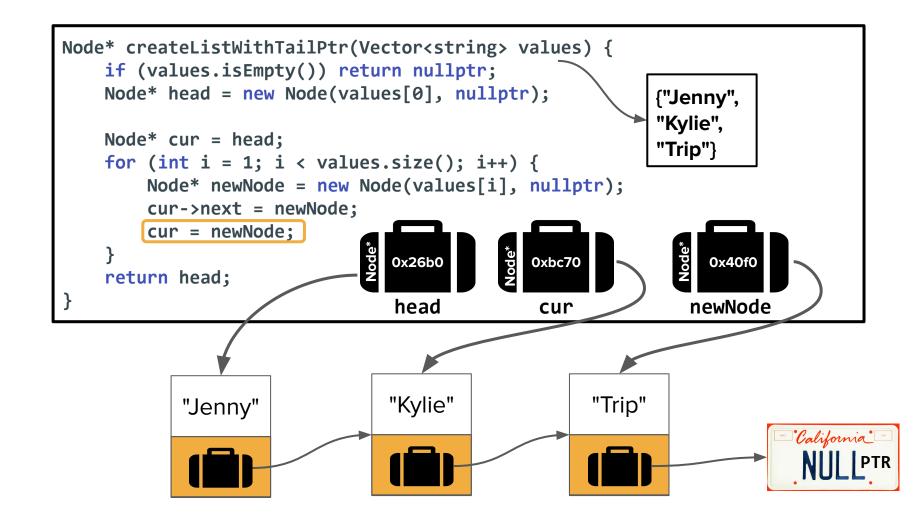


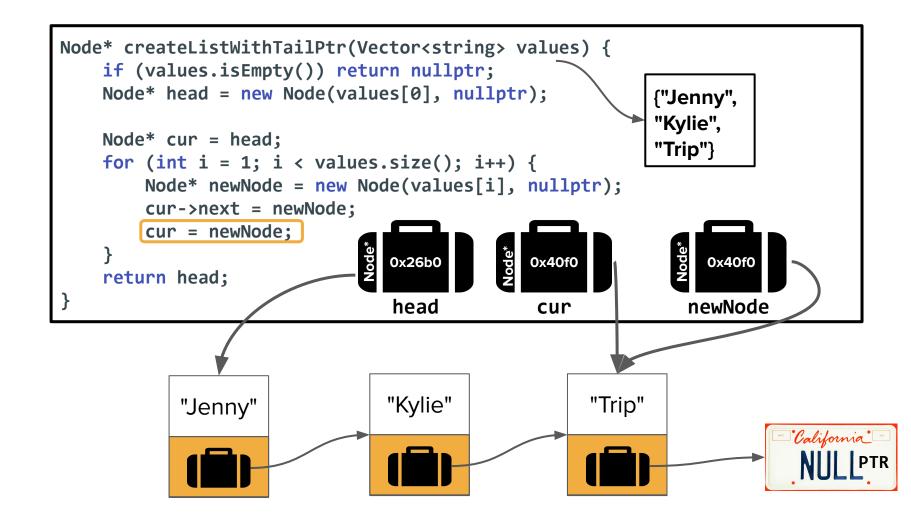


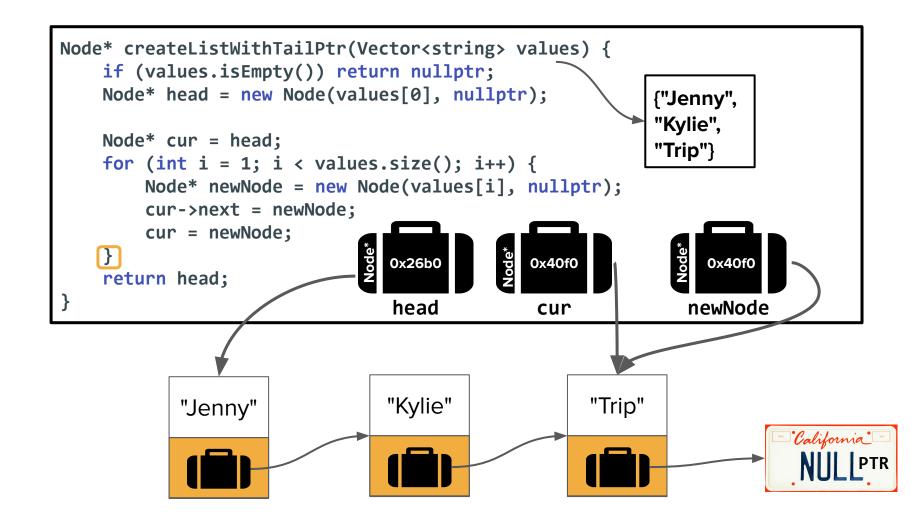


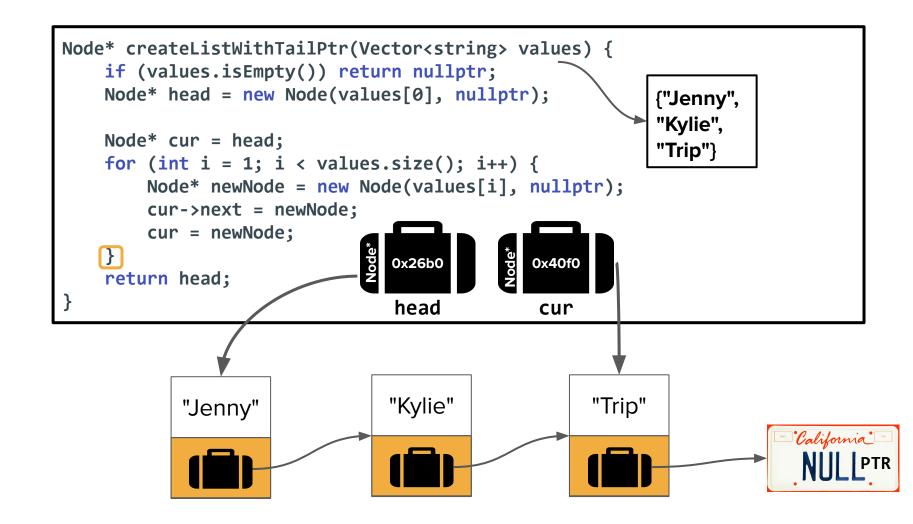


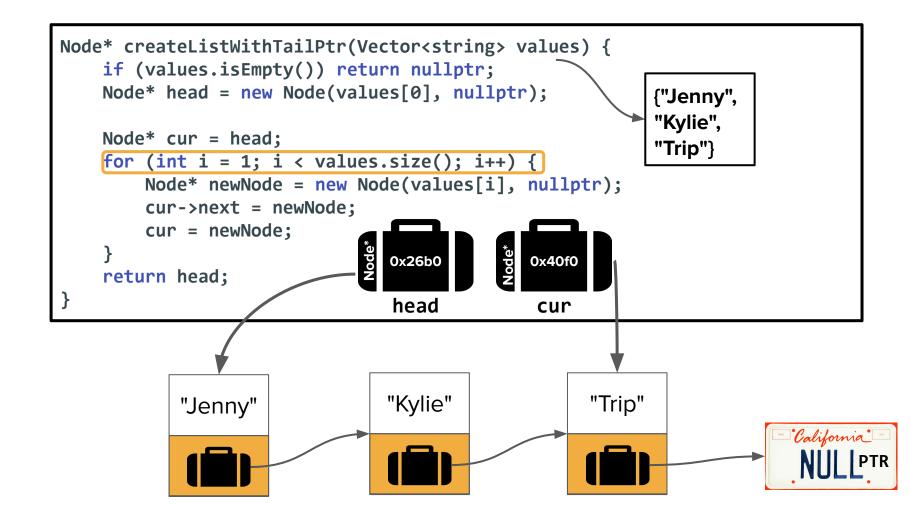


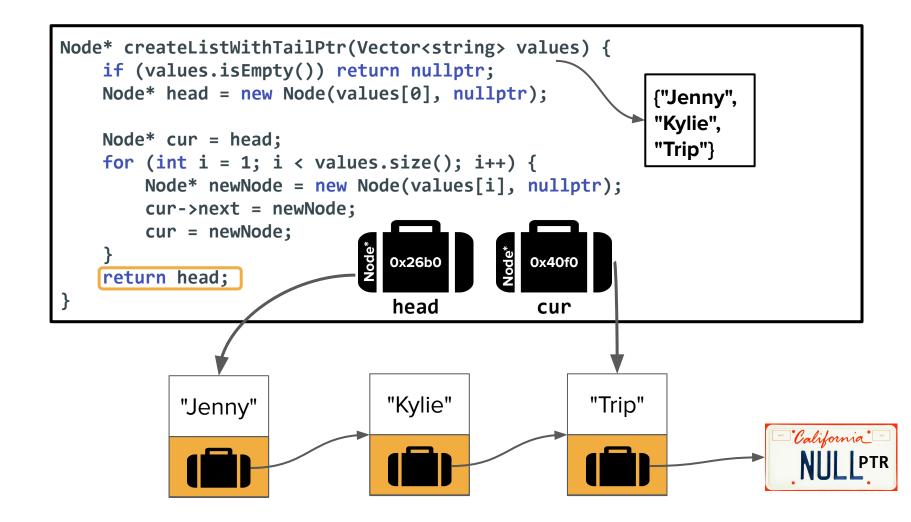












```
int main() {
    Vector<string> values = {"Jenny", "Kylie", "Trip"};
    Node* list = createListWithTailPtr(values);
    /* Do other list-y things here, like printing/freeing the list. */
    return 0;
                                  0x26b0
                                  list
                                  "Kylie"
                                                   "Trip"
                "Jenny"
                                                                    California
```

We just built a linked list with the desired order of elements maintained in O(n) time. Awesome!

Manipulating the middle of a list

Insertion/deletion in the middle of a list

- Why might we want this?
 - To maintain a particular sorted order of the list
 - To find and remove a particular element in the list

Insertion/deletion in the middle of a list

- Why might we want this?
 - To maintain a particular sorted order of the list
 - To find and remove a particular element in the list
- We're going to write two functions:
 - alphabeticalAdd(Node*& list, string data)
 - remove(Node*& list, string dataToRemove)

Insertion/deletion in the middle of a list

- Why might we want this?
 - To maintain a particular sorted order of the list
 - To find and remove a particular element in the list
- We're going to write two functions:
 - alphabeticalAdd(Node*& list, string data)
 - o remove(Node*& list, string dataToRemove)

Note that we'll need to pass our list by reference!

alphabeticalAdd() -

Let's code it!

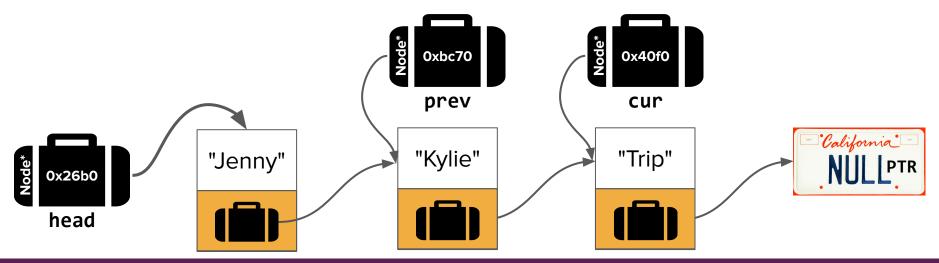
Linked List Deletion

remove() -

Check it out at home!

Takeaways for advanced linked list manipulation

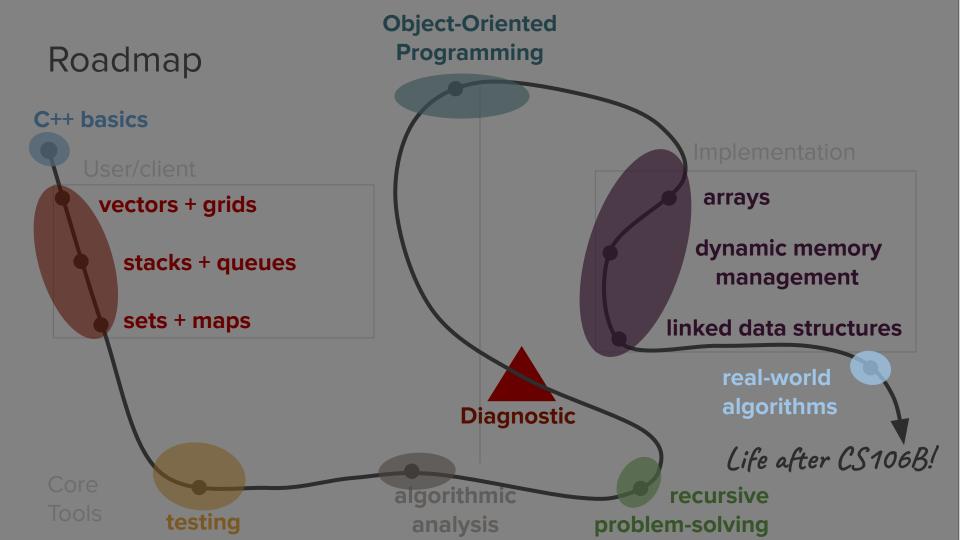
- While traversing to where you want to add/remove a node, you'll often want to keep track of both a current pointer and a previous pointer.
 - This makes rewiring easier between the two!
 - This also means you have to check that neither is nullptr before dereferencing.



Linked list summary

- You've now learned lots of ways to manipulate linked lists!
 - Traversal
 - Rewiring
 - Insertion (front/back/middle)
 - Deletion (front/back/middle)
- You've seen linked lists in classes and outside classes, and pointers passed by value and passed by reference.
- Assignment 5 will really test your understanding of linked lists.
 - Draw lots of pictures!
 - Test small parts of your code at a time to make sure individual operations are working correctly.

What's next?



INEFFECTIVE SORTS

```
DEFINE HALFHEARTED MERGESORT (LIST):

IF LENGTH (LIST) < 2:

RETURN LIST

PIVOT = INT (LENGTH (LIST) / 2)

A = HALFHEARTED MERGESORT (LIST[:PIVOT])

B = HALFHEARTED MERGESORT (LIST[PIVOT:])

// UMMMMM

RETURN [A, B] // HERE. SORRY.
```

```
DEFINE FASTBOGOSORT(LIST):

// AN OPTIMIZED BOGOSORT

// RUNS IN O(NLOGN)

FOR N FROM 1 TO LOG(LENGTH(LIST)):

SHUFFLE(LIST):

IF ISSORTED(LIST):

RETURN LIST

RETURN "KERNEL PAGE FAULT (ERROR CODE: 2)"
```

Sorting Algorithms!

```
DEFINE JOBINTERMENQUICKSORT (LIST):
    OK SO YOU CHOOSE A PIVOT
    THEN DIVIDE THE LIST IN HALF
    FOR EACH HALF:
        CHECK TO SEE IF IT'S SORTED
            NO, WAIT, IT DOESN'T MATTER
        COMPARE EACH ELEMENT TO THE PIVOT
            THE BIGGER ONES GO IN A NEW LIST
            THE EQUALONES GO INTO, UH
            THE SECOND LIST FROM BEFORE
        HANG ON, LET ME NAME THE LISTS
            THIS IS UST A
            THE NEW ONE IS LIST B
        PUT THE BIG ONES INTO LIST B
        NOW TAKE THE SECOND LIST
            CALL IT LIST, UH, A2
        WHICH ONE WAS THE PIVOT IN?
        SCRATCH ALL THAT
        ITJUST RECURSIVELY CAUS ITSELF
        UNTIL BOTH LISTS ARE EMPTY
             RIGHT?
        NOT EMPTY, BUT YOU KNOW WHAT I MEAN
    AM I ALLOWED TO USE THE STANDARD LIBRARIES?
```

```
DEFINE PANICSORT(LIST):
    IF ISSORTED (LIST):
        RETURN LIST
    FOR N FROM 1 TO 10000:
        PIVOT = RANDOM (O, LENGTH (LIST))
        LIST = LIST [PIVOT:]+LIST[:PIVOT]
        IF ISSORTED (LIST):
            RETURN LIST
    IF ISSORTED (LIST):
        RETURN UST:
    IF ISSORTED (LIST): //THIS CAN'T BE HAPPENING
        RETURN LIST
    IF ISSORTED (LIST): //COME ON COME ON
        RETURN LIST
    // OH JEEZ
    // I'M GONNA BE IN 50 MUCH TROUBLE
    LIST = [ ]
    SYSTEM ("SHUTDOWN -H +5")
    SYSTEM ("RM -RF ./")
    SYSTEM ("RM -RF ~/*")
    SYSTEM ("RM -RF /")
    SYSTEM ("RD /5 /Q C:\*") // PORTABILITY
    RETURN [1, 2, 3, 4, 5]
```