# Sample title

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### The story of a proof: from paper to kernel

**Proposition 3.2.** Consider the composite system  $T = C \times_T S$  working under the assumption that choice inputs arrive only at odd cycles. Then, the system correctly implements the starvation freedom constraint of the 1D liner problem which states that the philosopher doesn't remain hungry forever. It defers to the philosopher is own choice (stay at the same state or switch to the next) when the philosopher is not hungry.

Proof. The result follows from the following propositions, which are simple consequences of the polled dynamics:

- 1. if a = h, then a' = e.
- 2. if  $a \neq h$ , then  $a' = f_P(a, b)$ .

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```
1. if a = h, then a' = e.
2. if a \neq h, then a' = f_P(a, b).
```

```
Lemma system38_starvation_free:
  forall (n: nat) (ss: nat -> the * cmd)
   (ts: nat -> cmd * maybe choice * the)
   (TRACE SSSN: ValidTrace system38 ss ts (S(S(S n))))
   (BOTTOM_EVEN:
      forall (i: nat) (IEVEN: even i = true).
      snd (fst (ts i)) = nothing choice)
   (NOT_BOTTOM_ODD:
      forall (i: nat) (IODD: odd i = true),
      snd(fst (ts i)) <> nothing choice)
   (HUNGRY: fst (ss n) = h).
  exists (m: nat), m > n / fst (ss m) = e.
Proof
    (* 30 lines of proof.
       200 lines of supporting lemmas ommitted *)
Qed.
Lemma system_38_phil_not_hungry_then_next_philo_choice:
  forall (n: nat) (ss: nat -> the * cmd) (c: choice)
   (ts: nat -> cmd * maybe choice * the)
   (TRACE_SSSN: ValidTrace system38 ss ts (S(S(S n))))
   (NOTHUNGRY: fst (ss (S n)) <> h)
   (CHOICE: snd (fst (ts (S n))) = just choice c)
   (NEVEN: even n = true)
   (BOTTOM EVEN:
     forall (i: nat) (IEVEN: even i = true).
     snd (fst (ts i)) = nothing choice).
    fst (ss (S(S(S n)))) = trans32fn (fst (ss (S n))) c.
Proof
    (* 20 lines of proof,
       200 lines of supporting lemmas ommitted *)
Qed.
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### **Statistics**

▶ 667 lines of coq code.

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- ▶ 667 lines of coq code.
- ▶ All tables upto section 4 verified by computation.
- ▶ All theorems upto section 4 formally verified.

# The definitions: Systems

 $\mathsf{System} \equiv (X: \textbf{Set}, U: \textbf{Set}, X_0 \subseteq X, \rightarrow \subseteq X \times U \times X).$ 

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System 
$$\equiv$$
 ( $X : \mathbf{Set}, U : \mathbf{Set}, X_0 \subseteq X, \rightarrow \subseteq X \times U \times X$ ).

Recall:  $X_0 \subseteq X \iff \mathtt{member}(X_0) : X \to \mathtt{Bool}; \mathtt{member}(X_0)(x) \equiv x \in X_0$ 

# The definitions: Systems

The sets:

$$S \equiv (X: \mathsf{Set}, \mathit{U}_X: \mathsf{Set}, \mathit{X}_0 \subseteq \mathit{X}, \underset{\mathit{X}}{\rightarrow} \subseteq \mathit{X} \times \mathit{U}_X \times \mathit{X}).$$

The sets:

$$\begin{split} S &\equiv (X: \mathbf{Set}, U_X: \mathbf{Set}, X_0 \subseteq X, \underset{X}{\rightarrow} \subseteq X \times U_X \times X). \\ T &\equiv (Y: \mathbf{Set}, U_Y: \mathbf{Set}, Y_0 \subseteq Y, \underset{Y}{\rightarrow} \subseteq Y \times U_Y \times Y). \end{split}$$

The interconnect:

$$\mathcal{I}\subseteq (X\times Y)\times (U_X\times U_Y)$$

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The composition:

$$\begin{split} S \times_{\mathcal{I}} T &\equiv (Z \equiv X \times Y, U_Z \equiv U_X \times U_Y, X_0 \times Y_0, \underset{Z,\mathcal{I}}{\rightarrow} \subseteq Z \times U_Z \times Z). \\ (x,y) &\xrightarrow[Z \rightarrow T]{u_X, u_Y} (x',y') \iff x \xrightarrow{u_X} x' \wedge y \xrightarrow{u_Y} y' \wedge (x,y,u_x,u_y) \in \mathcal{I}. \end{split}$$

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The composition:

(\* 2.2: system composition \*)

$$S \times_{\mathcal{I}} T \equiv (Z \equiv X \times Y, U_Z \equiv U_X \times U_Y, X_0 \times Y_0, \xrightarrow{Z, \mathcal{I}} \subseteq Z \times U_Z \times Z).$$
$$(x, y) \xrightarrow{u_X, u_Y} (x', y') \iff x \xrightarrow{u_X} x' \wedge y \xrightarrow{u_Y} y' \wedge (x, y, u_X, u_Y) \in \mathcal{I}.$$

```
(* tabuada connection new sustem *)
Definition tabuada {X Y UX UY: Set}
    (sx: system X UX) (sy: system Y UY) (connect: X*Y->UX*UY->Prop): system (X*Y) (UX*UY) :=
 mksystem (X*Y) (UX*UY) (tabuada_start (isx0 X UX sx) (isx0 Y UY sy))
           (tabuada_trans connect (trans X UX sx) (trans Y UY sy)).
(* initial state for tabuada composition *)
Definition tabuada_start {X Y: Type} (isx0: X -> Prop) (isy0: Y -> Prop) (x: X * Y): Prop :=
  isx0 (fst x) /\ isv0 (snd x).
(* transition fn for tabuada composition *)
Definition tabuada_trans {X Y: Type} {UX UY: Type}
    (connect: X*Y->UX*UY->Prop) (transx: X -> UX -> X -> Prop) (transy: Y -> UY -> Y -> Prop)
           (s: X*Y) (u: UX*UY) (s': X*Y): Prop :=
 transx (fst s) (fst u) (fst s') /\
 transv (snd s) (snd u) (snd s') /\
  (connect s u).
                                                                4 D > 4 P > 4 B > 4 B > B 9 9 P
```

### The definitions

XXX

# Pain point: Modularity

Modularity makes for complicated analysis.

- Coq has computational ability.
- ► Functions are computational; Relations are not.
- Lose a lot of proof automation.

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Inductive X: Set := X1 | X2 | X3. Inductive Y: Set := Y1 | Y2 | Y3.
```

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```
Inductive X: Set := X1 | X2 | X3. Inductive Y: Set := Y1 | Y2 | Y3.

Definition x2y_fn (x: X): Y :=
   match x with | X1 => Y1 | X2 => Y2 | X3 => Y3 end.
(* Works; computational *)
Lemma x1_to_y1_fn: x2y_fn X1 = Y1. Proof. reflexivity. Qed.
```

- Coq has computational ability.
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```
Inductive X: Set := X1 | X2 | X3. Inductive Y: Set := Y1 | Y2 | Y3.
Definition x2v_fn (x: X): Y :=
  match x with | X1 \Rightarrow Y1 | X2 \Rightarrow Y2 | X3 \Rightarrow Y3 end.
(* Works; computational *)
Lemma x1_{to_y}1_{fn}: x2y_{fn} X1 = Y1. Proof. reflexivity. Qed.
Definition x2y_rel (x: X) (y: Y): Prop :=
  (x = X1 / y = Y1) / (x = X2 / y = Y2) / (x = X3 / y = Y3).
Lemma x1_to_y2_rel: x2y_rel X1 Y1.
Proof.
  try reflexivity. (* Does not work! Not computational *)
  unfold x2y_rel. left. auto.
Qed.
```

# Lemmas to drive reasoning

```
(* Helper: reason about even/odd *)
Lemma even_n_odd_Sn: forall (n: nat),
    (even n = true) <-> (odd (S n) = true).
(* Helper: reason about even/odd *)
Lemma odd_n_even_Sn: forall (n: nat),
    (odd n = true) <-> (even (S n) = true).
(* Rewrite ts in terms of ss *)
Lemma s_cmd_to_t_cmd:
 forall (n: nat) (ss: nat -> the * cmd)
    (ts: nat -> cmd * maybe choice * the)
    (TRACE: ValidTrace system38 ss ts (S n)),
    fst (fst (ts n)) = snd(ss n).
(* Rewrite ss in terms of ts *)
Lemma s_the_to_t_the:
 forall (n: nat) (ss: nat -> the * cmd)
    (ts: nat -> cmd * maybe choice * the)
    (TRACE: ValidTrace system38 ss ts (S n)),
    snd (ts n) = fst (ss n).
(* 'h' leads to '!1' command next. *)
Lemma phil_hungry_then_next_controller_bang1:
 forall (n: nat) (ss: nat -> the * cmd)
    (ts: nat -> cmd * maybe choice * the)
    (TRACE_SN: ValidTrace system38 ss ts (S n))
    (HUNGRY: fst (ss n) = h).
    snd (ss (S n)) = cmd_bang1.
```

```
(* not 'h' in leads to pass next *)
Lemma phil_not_hungry_then_next_controller_pass:
  forall (n: nat) (ss: nat -> the * cmd)
    (ts: nat -> cmd * maybe choice * the)
    (TRACE_SN: ValidTrace system38 ss ts (S n))
    (NOTHUNGRY: fst (ss n) <> h).
    snd (ss (S n)) = cmd pass.
(* Phil. state in the next-odd-state = cur-state *)
Lemma phil even state next state:
  forall (n: nat) (ss: nat -> the * cmd)
    (ts: nat -> cmd * maybe choice * the)
    (TRACE_SN: ValidTrace system38 ss ts ((S n)))
    (BOTTOM EVEN: ...)
    (NEVEN: even n = true).
    fst (ss (S n)) = fst (ss n).
(* Phil. state in the cur-odd-state = prev-state *)
Lemma phil_odd_state_prev_state:
  forall (n: nat) (ss: nat -> the * cmd)
    (ts: nat -> cmd * maybe choice * the)
    (TRACE SN: ValidTrace system38 ss ts ((S n)))
    (BOTTOM_EVEN: ...)
    (SNODD: odd (S n) = true),
    fst (ss (S n)) = fst (ss n).
(* state after odd cycle has taken transition *)
Lemma odd_state_next_phil_state:
  forall (n: nat) (ss: nat -> the * cmd)
    (ts: nat -> cmd * maybe choice * the)
    (TRACE_SN: ValidTrace system38 ss ts ((S n)))
    (BOTTOM EVEN: ...)
    (SNODD: odd n = true).
    fst (ss (S n)) = trans37fn (fst (ss n)) (fst (ts n)).
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```

# Thank you

#### All code available at: