

# Ellipsoid Algorithm

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**Abstract**—We motivate the Ellipsoid algorithm, discuss its original theoretical importance, and remark on its practical efficiency.

**Index Terms**—component, formatting, style, styling, insert

## I. INTRODUCTION

The ellipsoid algorithm was discovered in Naum Z. Shor. Later, Leonid Genrikhovich Khachiyan proved that the algorithm runs in polynomial time. This was a breakthrough in the theory of linear programming, which proved that solving LP's is in PTIME.

## II. DESCRIPTION THE ALGORITHM: CHECKING FOR NON-EMPTYNESS

## III. USING NON-EMPTYNESS TO SOLVE LP'S

So far, all we can do using the ellipsoid algorithm is to check if some system of equations  $Ax \leq b$  is *non-empty*. In other words, we can check the non-emptiness of a given polyhedron. Here, we will describe how to use this to solve *optimisation problems*.

Consider a linear program and its dual:

$$\begin{aligned} x, c &\in \mathbb{R}^{n \times 1} \quad A \in \mathbb{R}^{m \times n} \quad b, y \in \mathbb{R}^{m \times 1} \\ P_{\text{primal}} &\equiv \max_x c^T x \text{ subject to } Ax = b \\ P_{\text{dual}} &\equiv \min_y b^T y \text{ subject to } A^T y \geq c \end{aligned}$$

Let  $x^*$  be the optimal value of  $x$  for  $P_{\text{primal}}$ , and  $y^*$  be the optimal value of  $y$  for  $P_{\text{dual}}$ . From strong duality, we know that the value of  $c^T x^* = b^T y^*$ .

So, we can create a *combined* linear program, whose feasibility will force us to provide a point such that  $c^T x = b^T y$ . That is, we create a new polyhedra  $Q$  defined by the equations:

$$Ax \leq b \quad A^T y \geq c \quad c^T x = b^T y$$

Now, if a feasible point  $(x_0, y_0) \in Q$ , then it must be the case that  $Ax_0 = b$ ,  $A^T y_0 \geq c$ , and  $c^T x_0 = b^T y_0$ . At this point, strong duality tells us that  $(x_0, y_0) = (x^*, y^*)$ .

Hence, we can find the optimal value of the linear program by evaluating  $c^T x_0$ .

Thus, the ellipsoid algorithm can be used to solve for the optimality of a linear program, by starting from a non-emptiness check! This is beautiful, and proves a deep result of LP's: A certificate of non-emptiness is as good as a certificate of optimality.