Optimizing smallpt
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### Optimizing smallpt

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Haskell Exchange

November 4th, 2020



- ▶ 99 LoC C++ raytracer.
- Perfect for an optimization case study.
- Ported to many languages, including Haskell! (Thanks to Vo Minh Thu(noteed)).
- Start from noteed's original source; SHA the output image for baseline, keep optimizing.
- ▶ Plan: Quick walk through Haskell code, end up at C++ (clang++) performance.

### Optimizing smallpt

20-11-04

-What is smallpt anyway?



What is smallpt anyway?

Note for first slide: what is smallpt anyway?

```
struct Vec {
 double x, y, z; // position, also color (r,q,b)
  ... methods...
struct Ray { Vec o, d; Ray(Vec o_, Vec d_) : o(o_), d(d_) {} };
enum Refl_t { DIFF, SPEC, REFR }; // material types, used in radiance()
struct Sphere {
 double rad; // radius
 Vec p, e, c; // position, emission, color
 Refl_t refl; // reflection type (DIFFuse, SPECular, REFRactive)
  ... methods ...
 double intersect(const Ray &r) const // returns distance, 0 if nohit
Sphere spheres[] = {//Scene: radius, position, emission, color, material
 Sphere(1e5, Vec( 1e5+1,40.8,81.6), Vec(), Vec(.75,.25,.25), DIFF), //Left
  ... initialization ...
};
```

### Optimizing smallpt

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What is smallpt anyway?

### What is smallpt anyway?

struct Vec {
 double x, y, z; // position, also color (r,g,b)

); struct Ray { Vec o, d; Ray(Vec o\_, Vec d\_) : o(o\_), d(d\_) {} }; enum Refl\_t ( DIFF, SPEC, MEFR ); // material types, used in radiance( double rad; // radio

Redl\_t well; // reflection type (SIFFace, SPECular, SEFEaction) double interment(court May dr) count // returns distance, d if mobile 2phere spherea[] = {//2ces: radius, position, emission, color, materia 2phere(1s5, Vec( 1s5=1,60.0,81.6), Vec(),Vec(.75,.25,.25),DIFF),//ief

Vec radiance(const Ray &r, int depth, unsigned short \*Xi){

```
Optimizing smallpt
```

What is smallpt anyway?

What is smallpt anyway?

Say that the core function is radiance

```
What is smallpt anyway?
```

radiance

radiance

Vec radiance(const Ray &r, int depth, unsigned short \*Xi){

```
radiance
radiance
```

radiance

radiance

Optimizing smallpt

What is smallpt anyway?

What is smallpt anyway?

radiance

```
Vec radiance(const Ray &r, int depth, unsigned short *Xi){
```

```
if (
           ) if (
                                         else
if (
                  ){
                      radiance
} else if (
                      radiance
if (
                      radiance
 radiance
                             radiance
```

radiance

### Optimizing smallpt

What is smallpt anyway? What is smallpt anyway?

...with control flow

radiance

```
Vec radiance(const Ray &r, int depth, unsigned short *Xi){
```

```
Vec x=r.o+r.d*t, n=(x-obj.p).norm(), nl=n.dot(r.d)<0?n:n*-1, f=obj.c;
            ) if (
if (
                                           else
if (
                  ){
                       radiance
} else if (
                       radiance
if ((cos2t=1-nnt*nnt*(1-ddn*ddn))<0)
                       radiance
  radiance
                               radiance
```

radiance

Optimizing smallpt

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What is smallpt anyway?

What is smallpt anyway?

and lots of arithmetic

```
Vec radiance(const Ray &r. int depth, unsigned short *Xi){
 double t;
                                         // distance to intersection
 int id=0:
                                         // id of intersected object
 if (!intersect(r, t, id)) return Vec(); // if miss, return black
 const Sphere &obj = spheres[id];
                                         // the hit object
 Vec x=r.o+r.d*t, n=(x-obj.p).norm(), nl=n.dot(r.d)<0?n:n*-1, f=obj.c;
 double p = f.x > f.y && f.x > f.z ? f.x : f.y > f.z ? f.y : f.z; // max refl
 if (++depth>5) if (erand48(Xi)<p) f=f*(1/p); else return obj.e; //R.R.
 if (obj.refl == DIFF){
                                         // Ideal DIFFUSE reflection
   double r1=2*M PI*erand48(Xi), r2=erand48(Xi), r2s=sgrt(r2);
   Vec w=n1, u=((fabs(w.x)>.1?Vec(0,1):Vec(1))\%w).norm(), v=w\%u;
   Vec d = (u*cos(r1)*r2s + v*sin(r1)*r2s + w*sart(1-r2)).norm():
   return obj.e + f.mult(radiance(Ray(x,d),depth,Xi));
 } else if (obi.refl == SPEC)
                                         // Ideal SPECULAR reflection
   return obj.e + f.mult(radiance(Ray(x,r.d-n*2*n.dot(r.d)),depth,Xi));
  Ray reflRay(x, r.d-n*2*n.dot(r.d)); // Ideal dielectric REFRACTION
  bool into = n.dot(n1)>0;
                                         // Ray from outside going in?
 double nc=1, nt=1.5, nnt=into?nc/nt:nt/nc, ddn=r.d.dot(nl), cos2t;
 if ((cos2t=1-nnt*nnt*(1-ddn*ddn))<0) // Total internal reflection
   return obj.e + f.mult(radiance(reflRay,depth,Xi));
 Vec tdir = (r.d*nnt - n*((into?1:-1)*(ddn*nnt+sqrt(cos2t)))).norm();
 double a=nt-nc, b=nt+nc, R0=a*a/(b*b), c = 1-(into?-ddn:tdir.dot(n));
 double Re=RO+(1-RO)*c*c*c*c*c.Tr=1-Re.P=.25+.5*Re.RP=Re/P.TP=Tr/(1-P);
  return obj.e + f.mult(depth>2 ? (erand48(Xi) < P ? // Russian roulette
   radiance(reflRav.depth.Xi)*RP:radiance(Rav(x.tdir).depth.Xi)*TP) :
   radiance(reflRay,depth,Xi)*Re+radiance(Ray(x,tdir),depth,Xi)*Tr);
```



### Optimizing smallpt

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—What is smallpt anyway?

### What is smallpt anyway?

of Chinamanda, A. alil years have f(x) = f(x) + f

display the whole thing

# Initial Haskell Code: radiance $(1 \times)$

```
radiance :: Ray -> CInt -> Ptr CUShort -> IO Vec
radiance ray@(Ray o d) depth xi = case intersects ray of
(Nothing,_) -> return zerov
(Just t,Sphere _r p e c refi) -> do
```

```
continue f = case refl of
DIFF -> do
```

### radiance

```
SPEC -> do
  rad <- radiance
REFR -> do
```

```
if
  then do
  rad <- radiance reflRay depth' xi</pre>
```





show the same thing, this time in haskell

### Initial Haskell Code: radiance $(1\times)$

```
radiance :: Ray -> CInt -> Ptr CUShort -> IO Vec
radiance ray@(Ray o d) depth xi = case intersects ray of
 (Nothing,_) -> return zerov
 (Just t,Sphere _r p e c refl) -> do
   let x = o 'addv' (d 'mulvs' t)
       n = norm $ x `subv` p
       nl = if n 'dot' d < 0 then n else n 'mulvs' (-1)
        pr = maxv c
       depth' = depth + 1
       continue f = case refl of
         DIFF -> do
           r1 <- ((2*pi)*) 'fmap' erand48 xi
           r2 <- erand48 xi
           let r2s = sqrt r2
               w@(Vec wx _ _) = nl
               u = norm $ (if abs wx > 0.1 then (Vec 0 1 0) else (Vec 1 0 0)) `cross` w
               v = w `cross` u
               d' = norm $ (u`mulvs`(cos r1*r2s)) `addv` (v`mulvs`(sin r1*r2s)) `addv` (w`mulvs`sqrt (1-r2))
           rad <- radiance (Rav x d') depth' xi
           return $ e 'addy' (f 'mulv' rad)
         SPEC -> do
           let d' = d `subv` (n `mulvs` (2 * (n'dot'd)))
           rad <- radiance (Ray x d') depth' xi
           return $ e 'addy' (f 'mulv' rad)
         REFR -> do
           let reflRay = Ray x (d `subv` (n `mulvs` (2* n`dot`d)))
               into = n'dot'n1 > 0
               nc = 1
               nnt = if into then nc/nt else nt/nc
               ddn= d'dot'nl
               cos2t = 1-nnt*nnt*(1-ddn*ddn)
            if cos2t<0
             then do
               rad <- radiance reflRay depth' xi
```



# Optimizing smallpt Initial Haskell Code: radiance (1x) Initial Haskell Code: radiance (1x)

full code

```
Initial Haskell Code: Entry point (1\times)
```

```
smallpt :: Int -> Int -> Int -> IO ()
smallpt w h nsamps = do
 c <- VM.replicate (w * h) 0
 allocaArray 3 \xi -> -- Create mutable memory
   flip mapM_ [0..h-1] $ \y -> do -- Loop
     writeXi xi v
     for_ [0..w-1] \ x -> do -- Loop
       let i = (h-y-1) * w + x
       for_ [0..1] \sy -> do -- Loop
         for_ [0..1] \sx -> do -- Loop
           r <- newIORef 0 -- Create mutable memory
           for_ [0..samps-1] \_s -> do -- Loops, Loops
             r1 <- (2*) <$> erand48 xi
             rad <- radiance (Ray (org+d.*140) (norm d)) 0 xi -- Crunch
              . . .
             modifyIORef r (+ rad .* recip (fromIntegral samps)) -- Write
            ci <- VM.unsafeRead c i
            Vec rr rg rb <- readIORef r
           VM.unsafeWrite c i $
               ci + Vec (clamp rr) (clamp rg) (clamp rb) .* 0.25 -- Write
```

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Initial Haskell Code: Entry point  $(1 \times)$ 

 $\begin{aligned} & \min_{k \in \mathbb{N}} \left\{ \begin{array}{ll} - \log_{k} \log_{k}$ 

W. unandeWrite c 1 0 ci + Yes (class sw) (class sw) (class sh) . \* 0.25 - Frite

Initial Haskell Code: Entry point (1x)

### Initial Haskell Code: File I/O $(1\times)$

```
withFile "image.ppm" WriteMode $ \hdl -> do
    hPrintf hdl "P3\n%d %d\n%d\n" w h (255::Int)
    flip mapM_ [0..w*h-1] \i -> do
        Vec r g b <- VM.unsafeRead c i
        hPrintf hdl "%d %d %d " (toInt r) (toInt g) (toInt b)</pre>
```

```
foreign import ccall unsafe "erand48"
  erand48 :: Ptr CUShort -> IO Double
```



Optimizing smallpt

-module Main there -module Main (main) there

Restrict export list to 'main' (1.13×)

```
-module Main where
+module Main (main) where
```

```
Restrict export list to 'main' (1.13\times)
```

```
-module Main where
```

- +module Main (main) where
- Exported functions could be used by something unknown.
- original versions must be available.



### Restrict export list to 'main' $(1.13\times)$

```
-module Main where
+module Main (main) where
```

Exported functions could be used by something unknown.

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- original versions must be available.
- ▶ Make many optimizations unsound.



The very first thing to do is to let the compiler actually optimize. Exported functions could be used by something unknown to the compiler and thus their original versions must be available. This makes many optimizations look bad or be unreasonable. Explicit export lists tell the compiler what you care about. They aren't just about encapsulation.

### Mark entries of Ray and Sphere as UNPACK and Strict $(1.07 \times)$



### Optimizing smallpt

 $^-$ Mark entries of Ray and Sphere as UNPACK and Strict (1.07imes)

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Mark entries of Ray and Sphere as UNPACK and Strict (1.07×)

Strict means that they're evaluated ahead of time, not lazily of course. This can save a bit in their evaluation cost but means we always pay it. It is a trade off but here we know it will be a good one because we have almost perfect demand. Unpacking is different. It removes indirection of doing a memory lookup for components but means we have to copy everything into the data structure that it is unpacked into. We don't unpack ray because we'll be doing a lot of calculations on it and we want the Vec math to fuse and not get copied, when we use the Vecs they should be hot in cache anyway. We do inline Sphere because we statically create them and while they'll also be hot in cache there is no benefit to taking any extra cost at all.

### Use a pattern synonym to unpack Refl in Sphere $(1.07\times)$



### Optimizing smallpt

Use a pattern synonym to unpack Refl in Sphere (1.07imes)

```
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**March Mark (2007) Sept 1987 - Sept 1988 - Sept 1
```

Use a pattern synonym to unpack Refl in Sphere (1.07×)

Until the recent addition of unboxed sums in GHC there was no way to unpack a type with multiple constructors. While unboxed sums exist now, their syntax is quite unpleasant. We're using an older trick to fake the unboxing here instead. In this case it isn't much of a win, but it illustrates the technique.

### Change from maximum on a list to max $(1.08\times)$

```
-maxv (Vec a b c) = maximum [a,b,c]
+\max v (Vec a b c) = \max a (\max b c)
    let x = o `addv` (d `mulvs` t)
        n = norm $ x `subv` p
         nl = if n `dot` d < 0 then n else n `mulvs` (-1)</pre>
         pr = maxv c
         depth' = depth + 1
         continue f = case refl of
           DIFF -> do
    if depth'>5
       then do
        er <- erand48 xi
        let !pr = maxv c
```



# Optimizing smallpt Change from maximum on a list to max (1.08×) Change from maximum on a list to max (1.08×)

finicky optimization. GHC does not evaluate at compile time, making optimizations like these necessary  ${\sf necessary}$ 

### Convert erand48 to pure Haskell $(1.09\times)$



# Optimizing smallpt Convert erand48 to pure Haskell (1.09×) Convert erand48 to pure Haskell (1.09×)

This function isn't used much so it doesn't really impact our performance but often building unnecessary data structures is detrimental.

### Remove mutability: Erand48 Monad

```
-erand48 :: IORef Word64 -> IO Double
-erand48 !t = do
- r <- readIORef t
+data ET a = ET !Word64 !a deriving Functor
+newtype Erand48 a = Erand48 { runErand48' :: Word64 -> ET a } deriving Functor
+instance Applicative Erand48 where
+instance Monad Erand48 where
+runWithErand48 :: Int -> Erand48 a -> a
+erand48 :: Erand48 Double
-radiance :: Ray -> Int -> IORef Word64 -> IO Vec
-radiance ray@(Ray o d) depth xi = case intersects ray of
+radiance :: Rav -> Int -> Erand48 Vec
+radiance rav@(Rav o d) depth = case intersects rav of
            r1 <- (2*pi*) <$> erand48 xi
            r2 <- erand48 xi
             r1 <- (2*pi*) <$> erand48
            r2 <- erand48
                              then (.* rp) <$> radiance reflRav depth' xi
                              else (.* tp) <$> radiance (Ray x tdir) depth' xi
                              then (.* rp) <$> radiance reflRav depth'
                               else (.* tp) <$> radiance (Ray x tdir) depth'
```



### Optimizing smallpt

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Remove mutability: Erand48 Monad

```
Remove matality Exade Manie 

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```

The entire premise of this talk is that Haskell can be as fast as C. Further any impedance mismatch, such as FFI almost universally has to have, carries some bookkeeping overhead. If our Haskell code was as fast as the C code moving the code into Haskell would be a win, if it was slightly slower it could still be a win. Often considering your Haskell code's performance is a better option and easier than reimplementing something in C. As is the way with optimizations, this is not universally true.

### Removing mutation: eliminate IORef

```
- c <- VM.replicate (w * h) 0
- xi <- newTORef O
- flip mapM_ [0..h-1] $ \y -> do
      writeXi xi y
      for_{0..w-1} \ x \rightarrow do
        let i = (h-v-1) * w + x
        for_{0..1} \ v \rightarrow do
          for_ [0..1] \sx -> do
            r <- newIORef 0
            for [0..samps-1] \setminus s \rightarrow do
               r1 <- (2*) <$> erand48 xi
      img = (`concatMap` [(h-1),(h-2)..0]) $\y -> runWithErand48 y do
        for [0..w-1] \x -> do
           (pf - foldl m pf 0 [(sy, sx) | sy < [0,1], sx < [0,1]]) (ci (sy, sx) -> do
             Vec rr rg rb <- (\f -> foldlM f 0 [0..samps-1]) \ !r _s \rightarrow do
               r1 <- (2*) <$> erand48
               modifyIORef r (+ rad .* recip (fromIntegral samps))
             ci <- VM.unsafeRead c i
             Vec rr rg rb <- readIORef r
             VM.unsafeWrite c i $ ci + Vec (clamp rr) (clamp rg) (clamp rb) .* 0.25
               pure (r + rad .* recip (fromIntegral samps))
             pure (ci + Vec (clamp rr) (clamp rg) (clamp rb) .* 0.25)
```



### Optimizing smallpt

Removing mutation: eliminate IORef



All these mutability locations throw in extra RTS code, extra sequencing that blocks the compiler's optimization, and dependency chains. Sometimes we need mutability for performance but in fact one primary optimization technique powering modern compilers is SSA. We almost start there as a functional language don't break it when you don't have a good reason.

### Set everything in smallpt to be strict $(1.17\times)$

```
smallpt :: Int -> Int -> Int -> IO ()
smallpt w h nsamps = do
- let samps = nsamps `div` 4
       org = Vec 50 52 295.6
      dir = norm \$ Vec 0 (-0.042612) (-1)
       cx = Vec (fromIntegral w * 0.5135 / fromIntegral h) 0 0
       cy = norm (cx `cross` dir) `mulvs` 0.5135
+ let !samps = nsamps `div` 4
       !org = Vec 50 52 295.6
       !dir = norm \$ Vec 0 (-0.042612) (-1)
       !cx = Vec (fromIntegral w * 0.5135 / fromIntegral h) 0 0
       !cv = norm (cx `cross` dir) `mulvs` 0.5135
- r1 <- (2*) `fmap` erand48 xi
- let dx = if r1<1 then sqrt r1-1 else 1-sqrt(2-r1)
- r2 <- (2*) `fmap` erand48 xi
- let dy = if r2<1 then sqrt r2-1 else 1-sqrt(2-r2)
      d = ...
- rad <- radiance (Ray (org`addv`(d`mulvs`140)) (norm d)) 0 xi
+ !r1 <- (2*) `fmap` erand48 xi
+ let !dx = if r1<1 then sqrt r1-1 else 1-sqrt(2-r1)
+ !r2 <- (2*) `fmap` erand48 xi
+ let !dy = if r2<1 then sqrt r2-1 else 1-sqrt(2-r2)
       !d = ...
              pure $! r + rad .* recip (fromIntegral samps)
             pure $! ci + Vec (clamp rr) (clamp rg) (clamp rb) .* 0.25
```



# Optimizing smallpt Set everything in smallpt to be strict (1.17×) Set everything in smallpt to be strict (1.17×)

This is not a recommendation, this is a warning. We get a speedup here but it can also regress performance. Some of these bangs are regressions that are hidden. Don't senselessly bang everything in sight.

### Reduce to only useful strictnesses in smallpt(1.17 $\times$ )

```
- let !samps = nsamps `div` 4
       !org = Vec 50 52 295.6
       !dir = norm \$ Vec 0 (-0.042612) (-1)
       !cx = Vec (fromIntegral w * 0.5135 / fromIntegral h) 0 0
       !cv = norm (cx `cross` dir) .* 0.5135
  let samps = nsamps `div` 4
       org = Vec 50 52 295.6
       dir = norm \$ Vec 0 (-0.042612) (-1)
       cx = Vec (fromIntegral w * 0.5135 / fromIntegral h) 0 0
       cv = norm (cx `cross` dir) .* 0.5135
               !r1 <- (2*) <$> erand48
               r1 <- (2*) <$> erand48
               !r2 <- (2*) < $> erand 48
               r2 <- (2*) <$> erand48
               !rad <- radiance (Ray (org+d.*140) (norm d)) 0
               rad <- radiance (Ray (org+d.*140) (norm d)) 0
               pure $! r + rad .* recip (fromIntegral samps)
             pure $! ci + Vec (clamp rr) (clamp rg) (clamp rb) .* 0.25
               pure (r + rad .* recip (fromIntegral samps))
             pure (ci + Vec (clamp rr) (clamp rg) (clamp rb) .* 0.25)
```





Bangs force evaluation. The computation might diverge. Thus the compiler can no longer move the computation around or simplify it. Also if the work doesn't have to get done it still might be. A little thinking about how the variables are used or looking at core allows us to select which ones we bang selectively.

### Use strictness strategically in entire project

```
...
- if det<0 then Nothing else f (b-sdet) (b+sdet)
- where op = p - o
- eps = 1e-4
- b = dot op d
- det = b*b - dot op op + r*r
- sdet = sqrt det
- f a s = if a>eps then Just a else if s>eps then Just s else Nothing
+ if det<0
+ then Nothing
+ else
+ let !eps = 1e-4
+ !sdet = sqrt det
+ !a = b-sdet
+ !s = b+sdet
+ in if a>eps then Just a else if s>eps then Just s else Nothing
```



### Optimizing smallpt

20-11-0

Use strictness strategically in entire project



Use strictness strategically in entire project

Sometimes (point out 'intersect') we have to rearrange the code though when we use bangs. Bangs tell the compiler to make more efficient code, but take away the compiler's options in how to do so. Only take away the compiler's liberties when it's using them poorly. After a little trial and error or reading core how to use bangs in your code will be intuitive.

### Remove Maybe from intersect(s) $(1.32\times)$

```
| Old: Use Maybe Double to represent (was-hit?:bool, hit-distance: Double)
 New: use (1/0) to represent not (was-hit?)
-intersect :: Ray -> Sphere -> Maybe Double
+intersect :: Rav -> Sphere -> Double
intersect (Ray o d) (Sphere r p _e _c _refl) =
- if det<0 then Nothing else f (b-sdet) (b+sdet)
+ if det<0 then (1/0.0) else f (b-sdet) (b+sdet)
   where op = p `subv` o
         f a s = if a>eps then Just a else if s>eps then Just s else Nothing
         f a s = if a>eps then a else if s>eps then s else (1/0.0)
-intersects :: Ray -> (Maybe Double, Sphere)
+intersects :: Ray -> (Double, Sphere)
intersects rav = (k, s)
- where (k,s) = foldl' f (Nothing, undefined) spheres
         f (k',sp) s' = case (k',intersect ray s') of
                   (Nothing, Just x) -> (Just x,s')
                   (Just y, Just x) \mid x < y \rightarrow (Just x,s')
                   _ -> (k',sp)
+ where (k,s) = foldl' f (1/0.0,undefined) spheres
         f(k', sp) s' = let !x = intersect ray s' in if <math>x < k' then (x, s') else (k', sp)
radiance :: Ray -> Int -> STRefU s Word64 -> ST s Vec
 radiance ray@(Ray o d) depth xi = case intersects ray of
- (Nothing, ) -> return zerov
- (Just t,Sphere _r p e c refl) -> do
+ (t,_) \mid t == (1/0.0) -> return zerov
+ (t,Sphere _r p e c refl) -> do
```





20-11-0

Remove Maybe from intersect(s)  $(1.32\times)$ 

Figure 2. The state of the stat

This is a far more performance critical version of what we saw with 'maximum' vs. 'max'. Our innermost functions are of critical importance. Here we remove a Maybe which significantly reduces the boxing (which could have been mitigated with a StrictMaybe) and the cases. Since a Ray that fails to intersect something can be said to intersect at infinity, Double already actually covers the structure at play.

This also reduces allocation.

### Hand unroll the fold in intersects $(1.35\times)$

```
intersects :: Ray -> (Double, Sphere)
-intersects ray = (k, s)
- where (k,s) = foldl' f (1/0.0, undefined) spheres
+intersects ray =
    f (... (f (f (intersect ray sphLeft, sphLeft) sphRight) ...)
+ where
    f (k', sp) s' = let !x = intersect ray s' in if x < k' then (x, s') else (k', sp)
-spheres :: [Sphere]
-spheres = let s = Sphere ; z = zerov ; (.*) = mulvs ; v = Vec in
- [ s 1e5 (v (1e5+1) 40.8 81.6) z (v 0.75 0.25 0.25) DIFF --Left
- , s 1e5 (v (-1e5+99) 40.8 81.6) z (v 0.25 0.25 0.75) DIFF --Rght
. . .
+sphLeft, sphRight, ... :: Sphere
+sphLeft = Sphere 1e5 (Vec (1e5+1) 40.8 81.6)
                                                  zerov (Vec 0.75 0.25 0.25) DIFF
+sphRight = Sphere 1e5 (Vec (-1e5+99) 40.8 81.6) zerov (Vec 0.25 0.25 0.75) DIFF
```





20-11-0

Hand unroll the fold in intersects  $(1.35\times)$ 



Hand unroll the fold in intersects (1.35×)

Sadly 'intersects' remains one of our most costly functions even after that improvement so we'll focus on it a bit. A common optimization is loop unrolling. Many compilers do this for us, and there are special versions of it like Duff's Device. Sadly GHC is the one compiler I've used that doesn't. We can though implement every variety of it I know of by hand in Haskell, and this is one of the simplest. I believe we could produce a rule that did it for a given function but I've not actually explored that.

### Custom datatype for intersects parameter passing

```
Old: Tuple with possibly-uenevaluated Double and Sphere
New: Reference to a guaranteed-to-be-evaluated Double and Sphere
-intersects :: Ray -> (Double, Sphere)
+data T = T !Double !Sphere
+intersects :: Rav -> T
intersects rav =
    f ( ... f (intersect ray sphLeft, sphLeft) sphRight) ... sphLite
   f ( ... f (T (intersect ray sphLeft) sphLeft) sphRight) ... sphLite
  where
    f(k', sp) s' =
        let !x = intersect ray s' in if x < k' then (x, s') else (k', sp)
    f !(T k' sp) !s' =
        let !x = intersect ray s' in if x < k' then T x s' else T k' sp
radiance :: Ray -> Int -> Erand48 Vec
radiance ray@(Ray o d) depth = case intersects ray of
- (!t,_) \mid t == 1/0.0 \rightarrow return 0
- (!t,!Sphere _r p e c refl) -> do
+ (T t_{-}) | t == 1/0.0 -> return 0
+ (T t (Sphere _r p e c refl)) -> do
    let !x = o + d .* t
         !n = norm \$ x - p
         !nl = if dot n d < 0 then n else negate n
```



### Optimizing smallpt

0-11-04

Custom datatype for intersects parameter passing

Custom distingue for intercrectin parameter passing

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ta = norm \$ x - p tal = if dot n d < 0 then n else negate n

Still trying to focus on the performance of 'intersects' we can optimize how it passes data to itself and its caller. We want the data to be strict so the compiler knows the tuple's members will always be evaluated, but we want to avoid copying. A normal tuple lacks strictness information. An unboxed tuple sadly forces copying because of its unpacked nature. Sphere is quite large, and will be expensive to copy. What we need is a strict tuple, and we create one for just this usage. This exists in libraries of course, but we wanted to illustrate it.

### Optimize file writing

```
build-depends:
      base >= 4.12 && < 4.15
    , bytestring ^>= 0.11
-toInt :: Double -> Int
-toInt x = floor $ clamp x ** recip 2.2 * 255 + 0.5 
+toInt :: Double -> BB.Builder -- O(1) concatenation
+toInt x = BB.intDec (floor (clamp x ** recip 2.2 * 255 + 0.5)) <> BB.char8 ' '
  withFile "image.ppm" WriteMode $ \hdl -> do
         hPrintf hdl "P3\n%d \d\n%d\n" w h (255::Int)
        for_ img \(Vec r g b) -> do
          hPrintf hdl "%d %d %d " (toInt r) (toInt g) (toInt b)
         BB.hPutBuilder hdl $
           BB.string8 "P3\n" <> -- efficient builders for ASCII
          BB.intDec w <> BB.char8 ' ' <> BB.intDec h <> BB.char8 '\n' <>
          BB.intDec 255 <> BB.char8 '\n' <>
           (mconcat $ fmap (\( (Vec r g b) -> toInt r <> toInt g <> toInt b) img)
```



### Optimizing smallpt build-depends: base >= 4.12 && < 4.15 -tolat :: Bushle -> lat -tolat x = floor 8 clasp x == recip 2.2 = 255 + 0.5 Optimize file writing vithFile "image.ppm" WriteHode \$ \hdl -> do SPRINGS BELL "PROMISE SENSONS" w. h. (266:: Dat)

Optimize file writing

for\_ing \(Vec r g b) -> do 

We're getting quite a bit faster now, but there's still some small pieces of low lying fruit. Strings are fairly inefficient and the conversion functions to them we use aren't particularly efficient. 'bytestring' has some efficient writing code, so we just convert to that for a modest gain.

Use LLVM backend  $(1.87\times)$ 

+package smallpt-opt

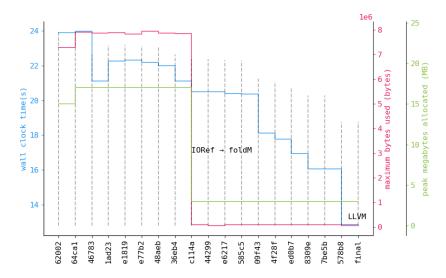
+ ghc-options: -fllvm

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Finally, this particular code is quite numeric heavy. There are optimizations for numeric heavy code we're missing in GHC. LLVM has an extensive library of laws to optimize low level numeric ops. LLVM is too low-level to understand haskell as haskell. It makes decisions with the tacit assumption that the assembly came from a C-like language, which is often to the detriment of a Haskell-like language. In this case, as the code is "fortran-like", LLVM wins.

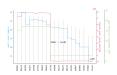
### The view from the mountaintop





### Optimizing smallpt

The view from the mountaintop



The view from the mountaintop

### **Takeaways**

- ► The unrolling in 'intersects' is ugly.
- ▶ (We feel) the maintainability of this code hasn't been significantly harmed.
- ▶ We're faster than clang++ and within 6% of g++
- ► Haven't exhausted the optimization opportunities.
- ▶ GHC could learn to do several of these optimizations for us.
- Others are just good Haskell style.
- Clean Haskell is often performant Haskell.





The unrolling in 'intersects' is ugly but other than that optimization it is our opinion that the maintainability of this code hasn't been significantly harmed by improving its performance. We're faster than clang++ and within 6We haven't exhausted the optimization opportunities. I believe GHC could learn to do several of these optimizations for us. Many of the others are just good Haskell style. Clean Haskell is often performant Haskell.