Optimizing smallpt
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November 4th, 2020

Optimizing smallpt

Davean Scies, Siddharth Bhat

Haskell Exchange

November 4th, 2020



- 1. S
- 2. Perfect for an optimization case study.
- 3. Plan: Quick walk through Haskell code, end up at C++ (clang++) performance.



- ▶ 99 LoC C++: **small p**ath **t**racer.
- Ported to many languages, including Haskell! (Thanks to Vo Minh Thu/noteed).
- Start from noteed's original source; SHA the output image from the Haskell source for baseline.

Optimizing smallpt

What is smallpt anyway?



What is smallpt anyway?

- 1. S
- 2. Perfect for an optimization case study.
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```
struct Vec {
 double x, y, z; // position, also color (r,q,b)
  ... methods...
struct Ray { Vec o, d; Ray(Vec o_, Vec d_) : o(o_), d(d_) {} };
enum Refl_t { DIFF, SPEC, REFR }; // material types, used in radiance()
struct Sphere {
 double rad; // radius
 Vec p, e, c; // position, emission, color
 Refl_t refl: // reflection type (DIFFuse, SPECular, REFRactive)
  ... methods ...
 double intersect(const Ray &r) const // returns distance, 0 if nohit
Sphere spheres[] = {//Scene: radius, position, emission, color, material
 Sphere(1e5, Vec( 1e5+1,40.8,81.6), Vec(), Vec(.75,.25,.25), DIFF), //Left
  ... initialization ...
};
```

Optimizing smallpt

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What is smallpt anyway?

where the ξ $= \xi_1$ $= \xi_2$ $= \xi_3$ $= \xi_3$

- 1. S
- 2. Has geometric primitives: vectors, spheres, materials
- 3. Entirely number-based, no real data structure

```
Vec radiance(const Ray &r, int depth, unsigned short *Xi){
```

```
4□ > 4률 > 4를 > 4를 > 를 되었다.
```

Optimizing smallpt

What is smallpt anyway?

- 1. **S**
- 2. Most of the compute cost is spent in the function that traces rays.
- 3. is called radiance

radiance

radiance

```
Vec radiance(const Ray &r, int depth, unsigned short *Xi){
```

```
radiance
radiance
```

radiance

radiance

```
Optimizing smallpt

What is smallpt anyway?

What is smallpt anyway?
```

- 1. **S**
- 2. Recursively calls itself a bunch of times

```
Vec radiance(const Ray &r, int depth, unsigned short *Xi){
```

```
if (
           ) if (
                                          else
if (
                  ){
                      radiance
} else if (
                      radiance
if (
                      radiance
 radiance
                              radiance
 radiance
                              radiance
```

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What is smallpt anyway?

- 1. S
- 2. Recursion is guarded by a lot of control flow

```
Vec radiance(const Ray &r, int depth, unsigned short *Xi){
```

```
Vec x=r.o+r.d*t, n=(x-obj.p).norm(), nl=n.dot(r.d)<0?n:n*-1, f=obj.c;
            ) if (
if (
                                           else
if (
                   ) {
                       radiance
} else if (
                       radiance
if ((cos2t=1-nnt*nnt*(1-ddn*ddn))<0)
                       radiance
  radiance
                               radiance
  radiance
                               radiance
```

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- 1. S
- 2. The control flow and computation is very numeric in nature

```
Vec radiance(const Ray &r, int depth, unsigned short *Xi){
```

```
Vec x=r.o+r.d*t, n=(x-obj.p).norm(), nl=n.dot(r.d)<0?n:n*-1, f=obj.c;
            ) if (erand48(Xi) )
if (
                                           else
if (
                   ){
                  erand48(Xi)
                                  erand48(Xi)
                       radiance
} else if (
                       radiance
if ((cos2t=1-nnt*nnt*(1-ddn*ddn))<0)
                       radiance
                                erand48(Xi)
  radiance
                               radiance
  radiance
                               radiance
```

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What is smallpt anyway?



- 1. **S**
- 2. We use erand48 for randomness

```
Vec radiance(const Ray &r. int depth, unsigned short *Xi){
 double t;
                                         // distance to intersection
 int id=0:
                                         // id of intersected object
 if (!intersect(r, t, id)) return Vec(); // if miss, return black
  const Sphere &obj = spheres[id];
                                         // the hit object
 Vec x=r.o+r.d*t, n=(x-obj.p).norm(), nl=n.dot(r.d)<0?n:n*-1, f=obj.c;
 double p = f.x > f.y && f.x > f.z ? f.x : f.y > f.z ? f.y : f.z; // max refl
 if (++depth>5) if (erand48(Xi)<p) f=f*(1/p); else return obj.e; //R.R.
 if (obj.refl == DIFF){
                                         // Ideal DIFFUSE reflection
   double r1=2*M PI*erand48(Xi), r2=erand48(Xi), r2s=sgrt(r2);
   Vec w=n1, u=((fabs(w.x)>.1?Vec(0,1):Vec(1))\%w).norm(), v=w\%u;
   Vec d = (u*cos(r1)*r2s + v*sin(r1)*r2s + w*sart(1-r2)).norm():
   return obj.e + f.mult(radiance(Ray(x,d),depth,Xi));
 } else if (obi.refl == SPEC)
                                         // Ideal SPECULAR reflection
   return obj.e + f.mult(radiance(Ray(x,r.d-n*2*n.dot(r.d)),depth,Xi));
  Ray reflRay(x, r.d-n*2*n.dot(r.d)); // Ideal dielectric REFRACTION
  bool into = n.dot(n1)>0;
                                         // Ray from outside going in?
 double nc=1, nt=1.5, nnt=into?nc/nt:nt/nc, ddn=r.d.dot(nl), cos2t;
  if ((cos2t=1-nnt*nnt*(1-ddn*ddn))<0) // Total internal reflection
   return obj.e + f.mult(radiance(reflRay,depth,Xi));
  Vec tdir = (r.d*nnt - n*((into?1:-1)*(ddn*nnt+sqrt(cos2t)))).norm();
 double a=nt-nc, b=nt+nc, R0=a*a/(b*b), c = 1-(into?-ddn:tdir.dot(n));
 double Re=RO+(1-RO)*c*c*c*c*c.Tr=1-Re.P=.25+.5*Re.RP=Re/P.TP=Tr/(1-P);
  return obj.e + f.mult(depth>2 ? (erand48(Xi) < P ? // Russian roulette
   radiance(reflRav.depth.Xi)*RP:radiance(Rav(x.tdir).depth.Xi)*TP) :
   radiance(reflRay,depth,Xi)*Re+radiance(Ray(x,tdir),depth,Xi)*Tr);
```



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—What is smallpt anyway?

What is smallpt anyway?

We consider the property of t

- 1. S
- 2. The full code continues to be more of the same

Initial Haskell Code: radiance $(1\times)$

```
radiance :: Ray -> CInt -> Ptr CUShort -> IO Vec
radiance ray@(Ray o d) depth xi = case intersects ray of
(Nothing,__) -> return zerov
(Just t,Sphere _r p e c refl) -> do
```

```
continue f = case refl of - BRANCHING
DIFF -> do
    r1 <- ((2*pi)*) `fmap` erand48 xi -- RMG</pre>
```

radiance

```
SPEC -> do
  rad <- radiance -- RECURSION
REFR -> do
```

```
if
then do
rad <- radiance
```

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Initial Haskell Code: radiance $(1\times)$



- 1. S
- 2. the original source has the same computation in haskell

Initial Haskell Code: Data structures

```
cross :: Vec -> Vec -> Vec
(.*) :: Vec -> Double -> Vec
infixl 7 .*
len :: Vec -> Double
norm :: Vec -> Vec
norm v = v .* recip (len v)
dot :: Vec -> Vec -> Double
maxv :: Vec -> Double

data Ray = Ray Vec Vec -- origin, direction
data Refl = DIFF | SPEC | REFR -- material types, used in radiance
-- / radius, position, emission, color, reflection
data Sphere = Sphere Double Vec Vec Refl
```

data Vec = Vec {-# UNPACK #-} !Double {-# UNPACK #-} !Double {-# UNPACK #-} !Double

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Initial Haskell Code: Data structures

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Initial Haskell Code: Data structures

- 1. S
- 2. We implement the same data structures in Haskell
- 3. Note that Ray, Sphere not having unpack

Initial Haskell Code: scene data

```
spheres :: [Sphere]
spheres =
 [ Sphere 1e5 (Vec (1e5+1) 40.8 81.6) 0 (Vec 0.75 0.25 0.25) DIFF --Left
 , Sphere 1e5 (Vec (99-1e5) 40.8 81.6) 0 (Vec 0.25 0.25 0.75) DIFF --Rght
 , Sphere 1e5 (Vec 50 40.8 1e5)
                                      0 0.75 DIFF --Back
 , Sphere 1e5 (Vec 50 40.8 (170-1e5)) 0 0
                                             DIFF --Frnt
 , Sphere 1e5 (Vec 50 1e5 81.6)
                                      0 0.75 DIFF --Botm
 , Sphere 1e5 (Vec 50 (81.6-1e5) 81.6) 0 0.75 DIFF -- Top
 , Sphere 16.5 (Vec 27 16.5 47)
                                0 0.999 SPEC --Mirr
 , Sphere 16.5 (Vec 73 16.5 78)
                                0 0.999 REFR --Glas
 , Sphere 600 (Vec 50 681.33 81.6)
                                             DIFF] -- Lite
                                     12 0
```





- 1. S
- 2. this list will be walked many times, as it contains our scene information.

Initial Haskell code: Sphere intersection

```
intersect :: Ray -> Sphere -> Maybe Double
intersect (Ray o d) (Sphere r p _e _c _refl) =
 if det<0 then Nothing else f (b-sdet) (b+sdet)
 where op = p - o -- Numeric
        eps = 1e-4
        b = dot op d
        det = b*b - dot op op + r*r -- Numeric
        sdet = sqrt det
       f a s = if a>eps then Just a else if s>eps then Just s else Nothing
intersects :: Ray -> (Maybe Double, Sphere)
intersects ray = (k, s)
 where (k,s) = foldl' f (Nothing, undefined) spheres -- Spheres iterated over
        f(k',sp) s' = case(k',intersect ray s') of
                  (Nothing, Just x) -> (Just x,s')
                  (Just y, Just x) \mid x < y \rightarrow (Just x,s')
                  _ -> (k',sp)
```



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Initial Haskell code: Sphere intersection

-> Sphere -> Maybe Double 4) (Sphere F p _e _c _pef1) -

interest. (May a d) (Sphere r p .e. c. _nex1) —
If detCO them Subling also r (G-adet) (B-adet)
Union op - p - a - - Fanoria
op - 1 = - Fanoria
op - 1 = - Fanoria
op - 1 = - Op - Op - PV - - Fanoria
obst - beta - du op op - PV - - Fanoria
obst - apt det

and - age one per the last a close if a sep then last a close little per the last a close if a sep then last a close little per - (height boddle, piners) interested any - (x_1, x_2) interested any - (x_1, x_2) interested per - (x_1, x_2)

Initial Haskell code: Sphere intersection

- 1. S
- 2. Responsible for figuring out what the ray hits.
- 3. We iterate over the list of spheres.
- 4. Once again, numeric heavy.
- 5. Use a Maybe to indicate whether we've found an answer or not.

Initial Haskell Code: radiance $(1\times)$

```
radiance :: Ray -> CInt -> Ptr CUShort -> IO Vec
radiance ray@(Ray o d) depth xi = case intersects ray of
 (Nothing,_) -> return zerov
 (Just t,Sphere _r p e c refl) -> do
   let x = o 'addy' (d 'mulvs' t)
       n = norm $ x 'subv' p
       nl = if n 'dot' d < 0 then n else n 'mulvs' (-1)
       pr = maxv c
       depth' = depth + 1
       continue f = case refl of
         DIFF -> do
           r1 <- ((2*pi)*) 'fmap' erand48 xi
           r2 <- erand48 xi
           let r2s = sqrt r2
               w@(Vec wx _ _) = nl
               u = norm $ (if abs wx > 0.1 then (Vec 0 1 0) else (Vec 1 0 0)) `cross` w
               v = w `cross` u
               d' = norm $ (u`mulvs`(cos r1*r2s)) `addv` (v`mulvs`(sin r1*r2s)) `addv` (w`mulvs`sqrt (1-r2))
           rad <- radiance (Ray x d') depth' xi
           return $ e 'addv' (f 'mulv' rad)
         SPEC -> do
           let d' = d `subv` (n `mulvs` (2 * (n'dot'd)))
           rad <- radiance (Ray x d') depth' xi
           return $ e 'addy' (f 'mulv' rad)
         REFR -> do
           let reflRay = Ray x (d `subv` (n `mulvs` (2* n`dot`d)))
               into = n'dot'n1 > 0
               nc = 1
               nnt = if into then nc/nt else nt/nc
               ddn= d'dot'nl
               cos2t = 1-nnt*nnt*(1-ddn*ddn)
           if cos2t<0
             then do
               rad <- radiance reflRay depth' xi
```



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Initial Haskell Code: radiance $(1\times)$



- 1. S
- 2. Branch heavy
- 3. Recursive
- 4. Uses an RNG

Initial Haskell Code: Entry point $(1\times)$

```
smallpt :: Int -> Int -> Int -> IO ()
smallpt w h nsamps = do
 c <- VM.replicate (w * h) 0
 allocaArray 3 \xi -> -- Create mutable memory
    flip mapM_ [0..h-1] $ \y -> do -- Loop
     writeXi xi v
     for_[0..w-1] \ x \rightarrow do -- Loop
        let i = (h-y-1) * w + x
        for_ [0..1] \sy -> do -- Loop
          for_ [0..1] \sx -> do -- Loop
            r <- newIORef 0 -- Create mutable memory
            for_ [0..samps-1] \_s -> do -- Loops, Loops
              r1 <- (2*) <$> erand48 xi
              rad <- radiance (Ray (org+d.*140) (norm d)) 0 xi -- Crunch
              modifyIORef r (+ rad .* recip (fromIntegral samps)) -- Write
            ci <- VM.unsafeRead c i
            Vec rr rg rb <- readIORef r</pre>
            VM.unsafeWrite c i $
                ci + Vec (clamp rr) (clamp rg) (clamp rb) .* 0.25 -- Write
            . . .
```

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Initial Haskell Code: Entry point $(1\times)$

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ci + Yec (class sw) (class sw) (class sh) .* 0.25 -- Svite

Initial Haskell Code: Entry point (1x)

- 1. S
- 2. Uses mutability
- 3. Performs number crunchy loops
- 4. Finally, writes results out

Initial Haskell Code: File I/O $(1\times)$

```
withFile "image.ppm" WriteMode $ \hdl -> do
    hPrintf hdl "P3\n%d %d\n%d\n" w h (255::Int)
    flip mapM_ [0..w*h-1] \i -> do
        Vec r g b <- VM.unsafeRead c i
        hPrintf hdl "%d %d %d " (toInt r) (toInt g) (toInt b)</pre>
```

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Initial Haskell Code: File I/O (1×)

```
Initial Haskell Code: RNG (1\times)
```

```
foreign import ccall unsafe "erand48"
erand48 :: Ptr CUShort -> IO Double
```





- 1. S
- 2. We use the RNG to decide randomly in which direction to send rays
- 3. Point out the use of foreign CCall.

Restrict export list to 'main' $(1.13 \times)$

-module Main where

+module Main (main) where





- 1. S
- 2. The very first thing to do is to let the compiler actually optimize.
- 3. Compiler can't know how exported functions are used.
- 4. Export lists not just about encapsulation.

Restrict export list to 'main' $(1.13\times)$

```
-module Main where
```

+module Main (main) where

- Exported functions could be used by something unknown.
- Original versions must be available.





Restrict export list to 'main' $(1.13\times)$

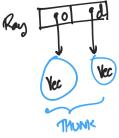


Restrict export list to 'main' (1.13×)

- 1. S
- 2. The very first thing to do is to let the compiler actually optimize.
- 3. Compiler can't know how exported functions are used.
- 4. Export lists not just about encapsulation.

Mark entries of Ray and Sphere as UNPACK and Strict $(1.07 \times)$

data Ray = Ray Vec Vec



- ▶ By default, all fields are *thunks* to rest of computation
- Pure, allow equational reasoning.

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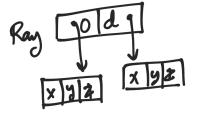
Mark entries of Ray and Sphere as UNPACK and Strict $(1.07 \times)$

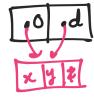
Mark entries of Ray and Sphere as UNPACK and Strict (1.07×)

► By default, all fields are thunks to rest of computation Pure, allow equational reasoning.

Mark entries of Ray and Sphere as UNPACK and Strict $(1.07 \times)$

data Ray = Ray !Vec !Vec





- ▶ When strict, elements are *pointers* to known structures
- pointers enable sharing!

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Mark entries of Ray and Sphere as UNPACK and Strict $(1.07\times)$

Mark entries of Ray and Sphere as UNPACK and Strict (1.07 \times)



data Ray - Ray (Sec (Sec



► When strict, elements are pointers to known structures
► pointers enable sharing!



- ▶ When unpacked, elements are *members* of the parent.
- Larger, but eliminate pointer chasing.



Mark entries of Ray and Sphere as UNPACK and Strict $(1.07 \times)$

```
data Vec = Vec {-# UNPACK #-} !Double
              {-# UNPACK #-} !Double
              {-# UNPACK #-} !Double
-data Ray = Ray Vec Vec -- origin, direction
+data Ray = Ray {-# UNPACK #-} !Vec {-# UNPACK #-} !Vec -- origin, direction
data Refl = DIFF | SPEC | REFR -- material types, used in radiance
 -- radius, position, emission, color, reflection
-data Sphere = Sphere Double Vec Vec !Refl
+data Sphere = Sphere {-# UNPACK #-} !Double
                     {-# UNPACK #-} !Vec
                     {-# UNPACK #-} !Vec
                     {-# UNPACK #-} !Vec !Refl
struct Vec { double x, y, z; }
struct Ray { std::function<Vec()> v; std::function<Vec()> w; };
struct RayUnpack { double xv, yv, int zv;
                  double xw, yw, zw; };
```



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Mark entries of Ray and Sphere as UNPACK and Strict (1.07×)

- 1. D
- Strictness in the arguments means that they're evaluated when instantiated, not when demanded.
- 3. Unpacking removes indirection from doing a memory lookup for components.
- 4. Means we have to copy everything into the data structure that it is unpacked into.
- 5. We don't unpack ray (Lots of calculations on its components, want those to fuse)
- 6. Unpack Sphere its static

Use a pattern synonym to unpack Refl in Sphere $(1.07\times)$



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Use a pattern synonym to unpack Refl in Sphere (1.07 imes)



Use a pattern synonym to unpack Refl in Sphere (1.07×)

- D
- 2. UnboxedSums are recent
- 3. UnboxedSums are very unpleasant
- 4. We're using an older trick to fake the unboxing here instead.
- 5. In this case it isn't much of a win, but it illustrates the technique.

Change from maximum on a list to max $(1.08\times)$

```
-maxv (Vec a b c) = maximum [a,b,c]
+\max v (Vec a b c) = \max a (\max b c)
    let x = o `addv` (d `mulvs` t)
         n = norm $ x `subv` p
         nl = if n `dot` d < 0 then n else n `mulvs` (-1)</pre>
         pr = maxv c
         depth' = depth + 1
         continue f = case refl of
           DIFF -> do
    if depth'>5
       then do
         er <- erand48 xi
        let !pr = maxv c
```

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Change from maximum on a list to max $(1.08 \times)$

Change from maximum on a list to max (1.08×)

- 1. D
- 2. finicky optimization.
- 3. Minimal usage of this function.
- 4. GHC does not evaluate at compile time, only has RULES

Convert erand48 to pure Haskell $(1.09\times)$

```
-foreign import ccall unsafe "erand48"
- erand48 :: Ptr CUShort -> IO Double
+erand48 :: IORef Word64 -> IO Double
+erand48 !t = do -- | Some number crunchy thing.
+ r <- readIORef t
+ let x' = 0x5deece66d * r + 0xb
       d_word = 0x3ff000000000000 . | . ((x' .&. 0xfffffffffff) `unsafeShiftL` 4)
       d = castWord64ToDouble d word - 1.0
+ writeTORef t x'
+ pure d
-radiance :: Ray -> CInt -> Ptr CUShort -> IO Vec
+radiance :: Ray -> Int -> IORef Word64 -> IO Vec -- IORef with state
 radiance ray@(Ray o d) depth xi = case intersects ray of
   c <- VM.replicate (w * h) zerov
- allocaArrav 3 $ \xi -> -- Old RNG state
       flip mapM_ [0..h-1] $ \y -> do
+ xi <- newIORef O -- New RNG state
+ flip mapM_ [0..h-1] $ \y -> do
       writeXi xi y
```



Optimizing smallpt Convert erand48 to pure Haskell (1.09×) Convert erand48 to pure Haskell (1.09×)

- 1. S
- 2. The entire premise of this talk is that Haskell can be as fast as C.
- 3. We're opening the black box of what erand48 does to GHC
- Further any impedance mismatch, such as FFI almost universally has to have, carries some bookkeeping overhead.
- 5. If our Haskell code was as fast as the C code moving the code into Haskell would be a win, if it was slightly slower it could still be a win.
- 6. Often considering your Haskell code's performance is a better option and easier than reimplementing something in C.
- 7. As is the way with optimizations, this is not universally true.

Remove mutability: Erand48 Monad

```
-erand48 :: IORef Word64 -> IO Double
-erand48 !t = do
- r <- readIORef t
+data ET a = ET !Word64 !a deriving Functor
+newtype Erand48 a = Erand48 { runErand48' :: Word64 -> ET a } deriving Functor
+instance Applicative Erand48 where
+instance Monad Erand48 where
+runWithErand48 :: Int -> Erand48 a -> a
+erand48 :: Erand48 Double
-radiance :: Ray -> Int -> IORef Word64 -> IO Vec
-radiance ray@(Ray o d) depth xi = case intersects ray of
+radiance :: Rav -> Int -> Erand48 Vec
+radiance ray@(Ray o d) depth = case intersects ray of
            r1 <- (2*pi*) <$> erand48 xi
            r2 <- erand48 xi
             r1 <- (2*pi*) <$> erand48
            r2 <- erand48
                              then (.* rp) <$> radiance reflRay depth' xi
                              else (.* tp) <$> radiance (Ray x tdir) depth' xi
                              then (.* rp) <$> radiance reflRav depth'
                               else (.* tp) <$> radiance (Ray x tdir) depth'
```



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Remove mutability: Erand48 Monad



- 1. D
- 2. All these mutability locations throw in extra RTS code, extra sequencing that blocks the compiler's optimization, and dependency chains.
- 3. Sometimes we need mutability for performance.
- 4. SSA is normal to compilers though.
- 5. We almost start at SSA as a functional language.
- 6. don't break it when you don't have a good reason.

Removing mutation: eliminate IORef and Data. Vector. Mutable

```
- c <- VM.replicate (w * h) 0
- xi <- newIORef 0
- flip mapM_ [0..h-1] $ \y -> do
       writeXi xi y
       for_{0..w-1} \ x \rightarrow do
         let i = (h-v-1) * w + x
         for_{0..1} \ v \rightarrow do
          for_ [0..1] \sx -> do
             r <- newIORef 0
             for [0..samps-1] \setminus s \rightarrow do
               r1 <- (2*) <$> erand48 xi
       img = (`concatMap` [(h-1), (h-2)..0]) $\y -> runWithErand48 y do
         for [0..w-1] \ x \rightarrow do
           (\protect\) foldlM pf 0 [(sy, sx) | sy <- [0,1], sx <- [0,1]]) \ci (sy, sx) -> do
             Vec rr rg rb <- (\f -> foldlM f 0 [0..samps-1]) \ !r _s -> do
               r1 <- (2*) <$> erand48
               modifyIORef r (+ rad .* recip (fromIntegral samps))
             ci <- VM.unsafeRead c i
             Vec rr rg rb <- readIORef r
             VM.unsafeWrite c i $ ci + Vec (clamp rr) (clamp rg) (clamp rb) .* 0.25
               pure (r + rad .* recip (fromIntegral samps))
             pure (ci + Vec (clamp rr) (clamp rg) (clamp rb) .* 0.25)
```

```
Removing mutation: eliminate IORef and Data . Vector . Mutable
Optimizing smallpt
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                        \begin{aligned} & \text{for}_{i} = \{ x_{i}, x_{i} \in \mathcal{Y}_{i} \mid x_{i} = x \\ & \text{for}_{i} = \{ x_{i}, x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i}, x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i}, x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i}, x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i}, x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i}, x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i}, x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i}, x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i}, x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i}, x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i}, x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i}, x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i}, x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i}, x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i}, x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i}, x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i}, x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i}, x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} = \{ x_{i} \neq x_{i} \} \mid x_{i} > x \\ & \text{for}_{i} 
                                                                                                                                                                                                                                                                                     Removing mutation: eliminate IORef and
                                                                                                                                                                                                                                                                                 Data Vector Mutable
```

1. D



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Set everything in smallpt to be strict $(1.17\times)$

```
smallpt :: Int -> Int -> Int -> IO ()
smallpt w h nsamps = do
- let samps = nsamps `div` 4
       org = Vec 50 52 295.6
      dir = norm \$ Vec 0 (-0.042612) (-1)
       cx = Vec (fromIntegral w * 0.5135 / fromIntegral h) 0 0
       cy = norm (cx `cross` dir) `mulvs` 0.5135
+ let !samps = nsamps `div` 4
       !org = Vec 50 52 295.6
       !dir = norm \$ Vec 0 (-0.042612) (-1)
       !cx = Vec (fromIntegral w * 0.5135 / fromIntegral h) 0 0
       !cv = norm (cx `cross` dir) `mulvs` 0.5135
- r1 <- (2*) `fmap` erand48 xi
- let dx = if r1<1 then sqrt r1-1 else 1-sqrt(2-r1)
- r2 <- (2*) `fmap` erand48 xi
- let dy = if r2<1 then sqrt r2-1 else 1-sqrt(2-r2)
      d = ...
- rad <- radiance (Ray (org`addv`(d`mulvs`140)) (norm d)) 0 xi
+ !r1 <- (2*) `fmap` erand48 xi
+ let !dx = if r1<1 then sqrt r1-1 else 1-sqrt(2-r1)
+ !r2 <- (2*) `fmap` erand48 xi
+ let !dy = if r2<1 then sqrt r2-1 else 1-sqrt(2-r2)
       !d = ...
              pure $! r + rad .* recip (fromIntegral samps)
             pure $! ci + Vec (clamp rr) (clamp rg) (clamp rb) .* 0.25
```

Optimizing smallpt

-11-05

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Set everything in smallpt to be strict $(1.17 \times)$



Don't senselessly bang everything in sight.

- 1.
- 2. This is not a recommendation, this is a warning.
- 3. We get a speedup here but it can also regress performance. Some of these bangs are regressions that are hidden.

Why strictness may be bad

```
let foo = let x = error "ERR" in \y -> y
let foo' = let !x = error "ERR" in \y -> y
```



- 1. S
- 2. bangs remove freedom from the compiler
- 3. bangs force large objects to be evaluated
- 4. Haskell is awesome because equational reasoning.
- 5. Bangs are not great with equational reasoning.
- 6. judiciously use bangs. Too many remove choice from the compiler.

Why strictness may be bad

```
let foo = let x = error "ERR" in \y -> y
let foo' = let !x = error "ERR" in \y -> y
let fooOpt = \y -> y
let foo'Opt = \y -> y -- ERROR! forcing foo' should give "ERR"
```





Why strictness may be bad



Why strictness may be bad

- 1. S
- 2. bangs remove freedom from the compiler
- 3. bangs force large objects to be evaluated
- 4. Haskell is awesome because equational reasoning.
- 5. Bangs are not great with equational reasoning.
- 6. judiciously use bangs. Too many remove choice from the compiler.

Reduce to only useful strictnesses in smallpt($1.17 \times$)

```
- let !samps = nsamps `div` 4
       !org = Vec 50 52 295.6
       !dir = norm \$ Vec 0 (-0.042612) (-1)
       !cx = Vec (fromIntegral w * 0.5135 / fromIntegral h) 0 0
       !cv = norm (cx `cross` dir) .* 0.5135
  let samps = nsamps `div` 4
       org = Vec 50 52 295.6
       dir = norm \$ Vec 0 (-0.042612) (-1)
       cx = Vec (fromIntegral w * 0.5135 / fromIntegral h) 0 0
       cv = norm (cx `cross` dir) .* 0.5135
               !r1 <- (2*) <$> erand48
               r1 <- (2*) <$> erand48
               !r2 <- (2*) < $> erand 48
               r2 <- (2*) <$> erand48
               !rad <- radiance (Ray (org+d.*140) (norm d)) 0
               rad <- radiance (Ray (org+d.*140) (norm d)) 0
               pure $! r + rad .* recip (fromIntegral samps)
             pure $! ci + Vec (clamp rr) (clamp rg) (clamp rb) .* 0.25
               pure (r + rad .* recip (fromIntegral samps))
             pure (ci + Vec (clamp rr) (clamp rg) (clamp rb) .* 0.25)
```





trad <- radiance (Ray (orged.*560) (norm d)) (

pure \$1 x + rad .* recip (from[ategral samps) pure \$1 ci + Vec (clamp rV) (clamp rg) (clamp rb) .* 0.25

- 1. D
- 2. Bangs force evaluation.
- 3. The computation might diverge.
- 4. Thus the compiler can no longer move the computation around or simplify it.
- Useless work.
- 6. A little thinking about how the variables are used or looking at core allows us to select which ones we bang selectively.

Use strictness strategically in entire project

```
if det<0 then Nothing else f (b-sdet) (b+sdet)

where op = p - o

eps = 1e-4

b = dot op d

det = b*b - dot op op + r*r

sdet = sqrt det

f a s = if a>eps then Just a else if s>eps then Just s else Nothing

if det<0

then Nothing

else

let !eps = 1e-4

!sdet = sqrt det

!a = b-sdet

!s = b+sdet

in if a>eps then Just a else if s>eps then Just s else Nothing
```



Optimizing smallpt

Use strictness strategically in entire project



Use strictness strategically in entire project

- 1. D
- Sometimes (point out 'intersect') we have to rearrange the code though when we use bangs.
- 3. Bangs tell the compiler to make more efficient code, but take away the compiler's options in how to do so.
- 4. Only take away the compiler's liberties when it's using them poorly.
- 5. Becomes intuitive.

Remove Maybe from intersect(s) $(1.32\times)$

```
Old: Use Maybe Double to represent (was-hit?:bool, hit-distance: Double)
 New: use (1/0) to represent not (was-hit?)
-intersect :: Ray -> Sphere -> Maybe Double
+intersect :: Rav -> Sphere -> Double
intersect (Ray o d) (Sphere r p _e _c _refl) =
- if det<0 then Nothing else f (b-sdet) (b+sdet)
+ if det<0 then (1/0.0) else f (b-sdet) (b+sdet)
   where op = p `subv` o
         f a s = if a>eps then Just a else if s>eps then Just s else Nothing
         f a s = if a>eps then a else if s>eps then s else (1/0.0)
-intersects :: Ray -> (Maybe Double, Sphere)
+intersects :: Ray -> (Double, Sphere)
intersects rav = (k, s)
- where (k.s) = foldl' f (Nothing, undefined) spheres
         f (k',sp) s' = case (k',intersect ray s') of
                   (Nothing, Just x) -> (Just x,s')
                   (Just y, Just x) \mid x < y \rightarrow (Just x,s')
                   _ -> (k',sp)
+ where (k,s) = foldl' f (1/0.0,undefined) spheres
         f(k', sp) s' = let !x = intersect ray s' in if <math>x < k' then (x, s') else (k', sp)
radiance :: Ray -> Int -> STRefU s Word64 -> ST s Vec
radiance ray@(Ray o d) depth xi = case intersects ray of
- (Nothing, ) -> return zerov
- (Just t,Sphere _r p e c refl) -> do
+ (t,_) \mid t == (1/0.0) -> return zerov
+ (t,Sphere _r p e c refl) -> do
```

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Optimizing smallpt

Remove Maybe from intersect(s) $(1.32 \times)$

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Remove Maybe from intersect(s) (1.32×)

- 1. S
- 2. This is a far more performance critical version of what we saw with 'maximum' vs. 'max'.
- 3. innermost functions are of critical importance. remove Maybe which significantly reduces the boxing
- 4. Since a Ray that fails to intersect something can be said to intersect at infinity, Double already actually covers the structure at play
- 5. This also reduces allocation.

Hand unroll the fold in intersects $(1.35\times)$

```
intersects :: Ray -> (Double, Sphere)
-intersects ray = (k, s)
- where (k,s) = foldl' f (1/0.0, undefined) spheres
+intersects ray =
+ f (... (f (f (intersect ray sphLeft, sphLeft) sphRight) ...)
+ where
    f(k', sp) s' = let !x = intersect ray s' in if x < k' then (x, s') else (k', sp)
-spheres :: [Sphere]
-spheres = let s = Sphere ; z = zerov ; (.*) = mulvs ; v = Vec in
- [ s 1e5 (v (1e5+1) 40.8 81.6) z (v 0.75 0.25 0.25) DIFF --Left
- , s 1e5 (v (-1e5+99) 40.8 81.6) z (v 0.25 0.25 0.75) DIFF --Rght
. . .
+sphLeft, sphRight, ... :: Sphere
+sphLeft = Sphere 1e5 (Vec (1e5+1) 40.8 81.6)
                                                 zerov (Vec 0.75 0.25 0.25) DIFF
+sphRight = Sphere 1e5 (Vec (-1e5+99) 40.8 81.6) zerov (Vec 0.25 0.25 0.75) DIFF
```



Optimizing smallpt

Hand unroll the fold in intersects $(1.35\times)$

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Hand unroll the fold in intersects (1.35×)

- 1. D
- 2. 'intersects' is very hot
- 3. Loop unrolling
- 4. Many compilers do this for us, and there are special versions of it like Duff's Device.
- 5. Sadly GHC doesn't
- 6. Can do variants by hand.
- 7. RULE could handle each one specificly (only exactly that one)?

Custom datatype for intersects parameter passing

```
Old: Tuple with possibly-uenevaluated Double and Sphere
New: Reference to a guaranteed-to-be-evaluated Double and Sphere
-intersects :: Ray -> (Double, Sphere)
+data T = T !Double !Sphere
+intersects :: Rav -> T
intersects rav =
    f ( ... f (intersect ray sphLeft, sphLeft) sphRight) ... sphLite
+ f ( ... f (T (intersect ray sphLeft) sphRight) ... sphLite
  where
    f(k', sp) s' =
        let !x = intersect ray s' in if x < k' then (x, s') else (k', sp)
    f !(T k' sp) !s' =
        let !x = intersect ray s' in if x < k' then T x s' else T k' sp
radiance :: Ray -> Int -> Erand48 Vec
radiance ray@(Ray o d) depth = case intersects ray of
- (!t,_) \mid t == 1/0.0 \rightarrow return 0
- (!t,!Sphere _r p e c refl) -> do
+ (T t_{-}) | t == 1/0.0 -> return 0
+ (T t (Sphere _r p e c refl)) -> do
    let !x = o + d .* t
         !n = norm \$ x - p
         !nl = if dot n d < 0 then n else negate n
```



Optimizing smallpt

20-11-05

Custom datatype for intersects parameter passing

Custom datatype for intersectic parameter parameter passing
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for indersection is appointed to be reliable that the originate finals and figure
cases to Tambali tiplace.

***Comment of the Tambali tiplace cases the case of the Case

ta = norm \$ x - p tal = if dot n d < 0 then n else negate n

- 1. D
- 2. We can optimize data passing.
- 3. Want: Data strict, but not unpacked.
- 4. Compiler knows its evaluated but no copying
- 5. A normal tuple lacks strictness information.
- 6. An unboxed tuple forces copying
- 7. Strict Tuple.
- 8. This exists in libraries of course, but we wanted to illustrate it.

Optimize file writing

```
build-depends:
      base >= 4.12 && < 4.15
    , bytestring >= 0.11
-toInt :: Double -> Int
-toInt x = floor $ clamp x ** recip 2.2 * 255 + 0.5 
+toInt :: Double -> BB.Builder -- O(1) concatenation
+toInt x = BB.intDec (floor (clamp x ** recip 2.2 * 255 + 0.5)) <> BB.char8 ' '
. . .
  withFile "image.ppm" WriteMode $ \hdl -> do
         hPrintf hdl "P3\nd" \d\n\d\n" w h (255::Int)
        for_ img \(Vec r g b) -> do
          hPrintf hdl "%d %d %d " (toInt r) (toInt g) (toInt b)
         BB.hPutBuilder hdl $
          BB.string8 "P3\n" <> -- efficient builders for ASCII
          BB.intDec w <> BB.char8 ' ' <> BB.intDec h <> BB.char8 '\n' <>
          BB.intDec 255 <> BB.char8 '\n' <>
           (mconcat $ fmap (\( (Vec r g b) -> toInt r <> toInt g <> toInt b) img)
```



Optimizing smallpt

Optimize file writing



Optimize file writing

build-depends: base >= 4.12 && < 4.15

- 1. D
- 2. Strings are inefficient
- 3. 'bytestring' has some efficient writing code, so we just convert to that for a modest gain.

Use LLVM backend $(1.87\times)$

+package smallpt-opt

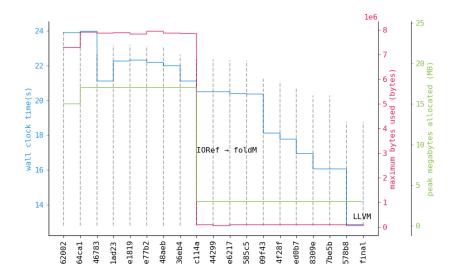
+ ghc-options: -fllvm





- 1. S
- 2. Finally, this particular code is quite numeric heavy.
- 3. There are optimizations for numeric heavy code we're missing in GHC.
- 4. LLVM has an extensive library of laws to optimize low level numeric ops.
- 5. LLVM is too low-level to understand haskell as haskell.
- 6. LLVM makes decisions with the tacit assumption that the assembly came from a C-like language, which is often to the detriment of a Haskell-like language.
- 7. In this case, as the code is "fortran-like", LLVM wins.

The view from the mountaintop





1. D



Avoid CPU ieee754 slow paths

```
intersect :: Ray -> Sphere -> Double
 intersect (Ray o d) (Sphere r p _e _c _refl) =
  if det<0
- then 1/0.0
+ then 1e20
   else
    . . .
    in if a>eps then a else if s>eps then s else 1/0.0
    in if a>eps then a else if s>eps then s else 1e20
radiance :: Ray -> Int -> Erand48 Vec
radiance ray@(Ray o d) depth = case intersects ray of
- (T t_) | t == 1/0.0 \rightarrow return 0
+ (T 1e20 _) -> return 0
```



Optimizing smallpt

-11-05

-Avoid CPU ieee754 slow paths

Avoid CPU ieee754 slow paths

intersect :: Ray -> Sphere -> Dushle intersect (May o d) (Sphere v p _e _c _vef1) = :f det0 - then 1/6.0 - then 1/6.0 - else

radiance rayO(ky o 4) depth = case intersects may of

- (T t _) | t = 1/0.0 -> return 0

- (T 1600 _) -> return 0

- (T 1600 _) -> return 0

- 1. D
- 2. We used +Inf to match the Maybeness
- 3. C++ code set 1e20 s the horizon
- 4. Mechanical sympathy is important.
- 5. Know how the CPU (abstractly) executes slow path / fast path.

Fix differences with C++ version

```
if depth>2
+ if depth'>2 -- depth' = depth + 1
```





- 1. S
- 2. Since the sha1 of the output didn't match the C++ version we started investigating.
- 3. clang++, g++ actually produce different sha1s
- 4. unincremented depth was being used in one branch, causing us to do more work
- 5. now confident to say we're doing the same computation as C++

Takeaways

- ► The unrolling in 'intersects' is ugly.
- ▶ (We feel) the maintainability of this code hasn't been significantly harmed.
- ▶ We're faster than clang++ and within 6% of g++
- ► Haven't exhausted the optimization opportunities.
- ▶ GHC could learn to do several of these optimizations for us.
- ▶ Others are just good Haskell style.
- Clean Haskell is often performant Haskell.



➤ The servolling in 'intersectin' in egly:
➤ (Wis first) the maintriabality of this code have't been significantly harmed.
➤ Wire future that charge+ and whith 6% of g++
➤ Havest enhanted the optimization apparamelies.
— GHC could favor not ownerful of these optimization for us.
➤ Codes now juig good Habidal egls.
➤ Class Habidal is other performant braid.

1. D



Experimental data

▶ All test were on an otherwise idle Equinix Metal c3.small.x86 (Intel Xeon E-2278G with 32GiB RAM, Ubuntu 20.04).

