



Demonstrating UPC++/Kokkos
Interoperability in a Heat Conduction
Simulation

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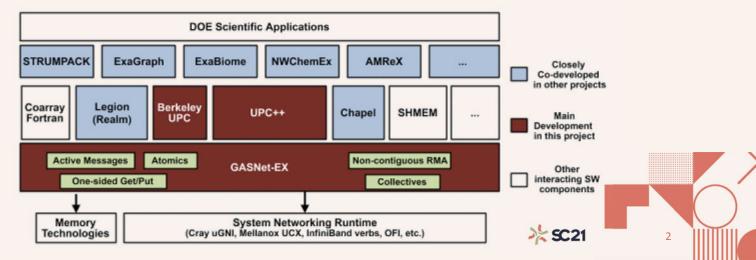
Video: doi:10.25344/S4DK5F



https://go.lbl.gov/pagoda

The Pagoda project

- Support for lightweight communication for exascale applications, frameworks and runtimes
 - GASNet-EX low-level layer that provides a network-independent interface suitable for Partitioned Global Address Space (PGAS) runtime developers
 - **UPC++** C++ PGAS library for application, framework and library developers, a productivity layer over GASNet-EX







Communication operations include:

- Remote Memory Access (RMA):
 - Get/put/atomics on a remote location in another address space
 - One-sided communication leverages low-overhead, zero-copy RDMA
- Remote Procedure Call (RPC):
 - Moves computation to the data

Design principles for performance and scalability

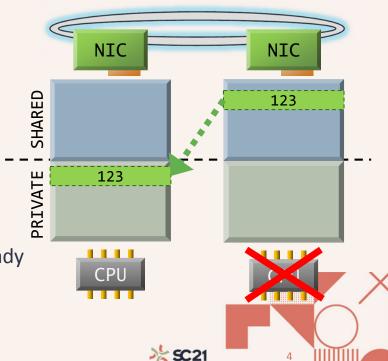
- All communication is syntactically explicit
- All communication is asynchronous: futures and promises
- Scalable data structures that avoid unnecessary replication



Asynchronous RMA in UPC++

- By default, all communication operations are split-phased
 - **Initiate** operation
 - Wait for completion

A UPC++ future holds a value and a state: ready/not-ready Wait returns the result when the rget completes



Remote procedure call (RPC)

- Execute a function on another process, sending arguments and returning an optional result
 - 1. Initiator injects the RPC to the target process, returning a future
 - 2. Target process executes fn(arg1, arg2) at some later time determined at the target
 - 3. Result becomes available to the initiator via the future
- Let's imagine that process 0 performs this RPC

 int area(int a, int b) { return a * b; }

 ... = upcxx::rpc(p, area, a, b).wait();

 upcxx::rpc(p, area, a, b)

 ["area", a, b]

 ["area", a, b]

Aggressive Asynchrony via Futures and callbacks

- RMA and RPC both return a future object, which represents an operation that may or may not be complete
- Callbacks can be chained through calls to <u>then</u>()
- Multiple futures can be *conjoined* with <u>when all()</u> into a single future that encompasses all their results.

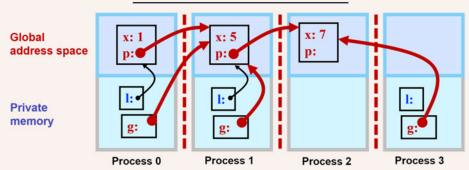
This code gets two remote values (an int and a double) and puts their product to another location:



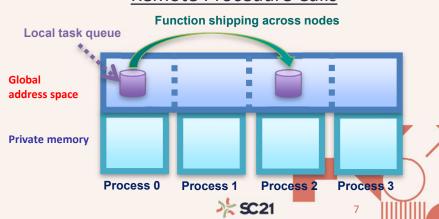
Overview of UPC++ Features

RMA and Global Pointers

- One-sided RMA and Global Pointers
- Remote Procedure Calls (RPC)
- Future-based asynchrony, continuations
- Remote atomics, non-contiguous RMA
- Serialization, distributed objects
- Teams and non-blocking collectives
- Hierarchical shared mem, node-level bypass
- Memory Kinds for GPU support
- Personas (multi-threading)
- Interoperability with other models
- For more details: https://upcxx.lbl.gov



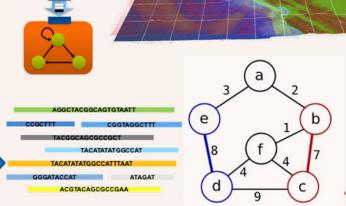
Remote Procedure Calls





UPC++ Application Examples

- Several applications have been written using UPC++, resulting in improved programmer productivity and runtime performance. Examples include:
- MetaHipMer, a genome assembler
- symPACK, a sparse symmetric matrix solver
- Pond, an actor-based shallow water simulation
- SIMCoV, an agent-based simulation of lungs with COVID-19
- Mel-UPX, a half-approximate graph matching solver

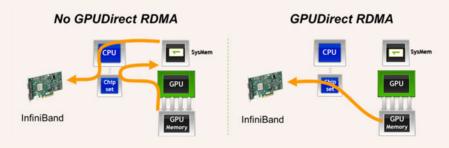






GPU Direct

- UPC++ RMA operations utilize Remote Direct Memory Access (RDMA) support when present in hardware
 - RDMA is an important performance feature in all modern HPC networks
 - GPUDirect RDMA (GDR) technology extends RDMA to GPU memory
- UPC++ "memory kinds" extend the PGAS model to encompass GPU memory
 - Distinction between host and GPU memory is part of the global pointer type
 - Permits static choice of appropriate communication code paths
- CUDA-aware MPI implementations lack equivalent static information







- Demonstrate interoperability of UPC++ and Kokkos
 - Kokkos: A C++ parallelism abstraction framework designed for portable leveraging of computational resources within a heterogeneous node, such as a GPGPU
 - UPC++: A C++ template library implementing an asynchronous one-sided PGAS programming model using gets, puts, and remote procedure calls complete with serialization of nontrivial C++ objects
 - Complementary C++ libraries emphasizing performance and productivity. Kokkos for on-node parallelism, UPC++ for between nodes.
- Demonstrate that UPC++ can be used to perform inter-node communication with comparable performance to MPI
 - MPI interoperability example provided by Kokkos project
 - MPI message passing is a two-sided model
 - Not an apples-to-apples comparison
 - Goal is merely to show performance isn't sacrificed

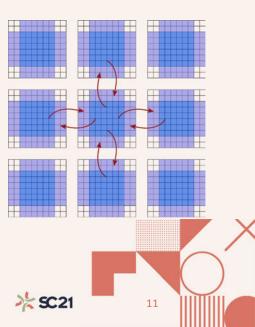


Heat Conduction Example

- 3D cube computation with halo exchange
- Ported to UPC++ without prior experience interoperating with Kokkos
- Both written naïvely, without excessive effort put into optimization

```
pack_T_halo();
compute_inner_dT();
exchange_T_halo();
compute_surface_dT();
Kokkos::fence();
double T_ave = compute_T();
```

Routines performed each timestep



Porting Communication from MPI to UPC++ : Pseudocode

```
for(n in neighbors) {
   n.outbuf.fence();
   MPI_Irecv(n.inbuf, n.inbuf.size(), MPI_DOUBLE, n.rank, ...);
   MPI_Isend(n.outbuf, n.outbuf.size(), MPI_DOUBLE, n.rank, ...);
}
MPI_Waitall(...);
```

GPU Memory Allocation

- Original two-sided MPI version interoperates with Kokkos Views
- UPC++ needs to interact with GPU memory with its own *memory kinds* interface
 - upcxx::device allocator allocates large GPU memory segment and registers it with the network adapter
 - Partitions of segment assigned to each halo boundary region
 - Use non-owning construction of Kokkos Views from pointer allow UPC++ control of allocation
- Exchange of global pointers:
 - upcxx::dist object's were used to communicate pointers for puts

```
using namespace upcxx;

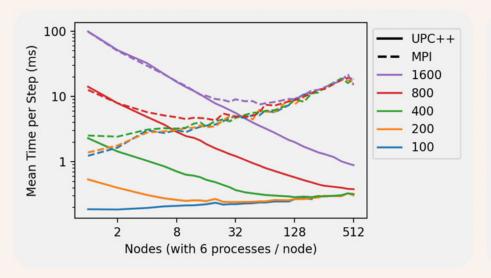
device allocator<cuda device> alloc(size, ...);
dist object<global ptr<double, memory kind::cuda device>> dptr(alloc->allocate<double>(size));
global ptr<double, memory kind::cuda device> nbr = dptr.fetch(nbr_rank).wait();
```

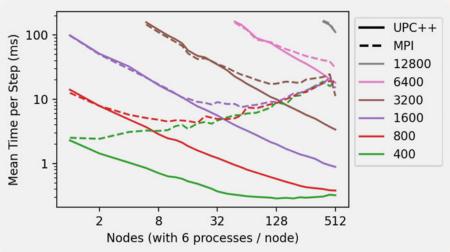
Benchmarking



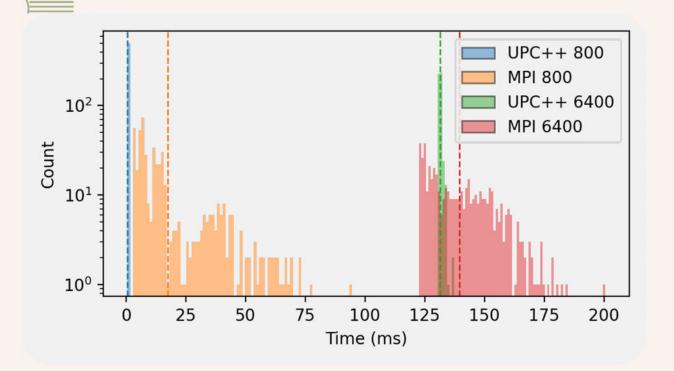
- Benchmarks performed on OLCF Summit
 - Each node has 6 NVIDIA V100 GPUs
 - One process per GPU
- CUDA-aware IBM Spectrum MPI 10.3.1.2-20200121 (for MPI version)
- GPUDirect RDMA enabled UPC++ 2021.3.0
- All benchmarks for a node count were executed within one job to minimize variability of network performance
- 500 timesteps, not including warmup iterations
- Profiled using sliding window of two timesteps used to ensure a communication happens within window, as load imbalance may result in early arrival of incoming data

Benchmark: Mean Time Per Step





Benchmark: Execution Time Variation



Vertical dotted lines indicate medians

128 Nodes

Conclusions

- UPC++ can interoperate with Kokkos without difficulty
- Porting from MPI to UPC++ was likewise straightforward
 - The intention is not to prove superiority to MPI performance
 - Well-tuned usage of both should get close to the hardware performance
- The new UPC++ memory kinds feature is performant
- UPC++ has many sophisticated high-level features, but don't need to sacrifice performance



Acknowledgements

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Thank you! Questions?