Décodage de codes polaires sur des architectures programmables

Mathieu Léonardon

13 Décembre 2018





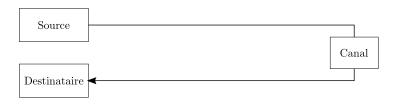




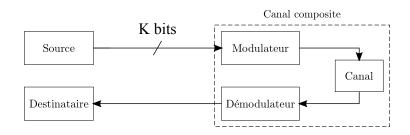
Introduction

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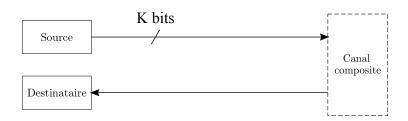
Chaîne de communication



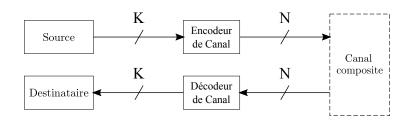
Message transmis au travers d'un canal de communication



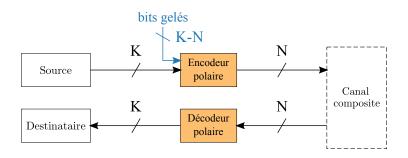
- Message transmis au travers d'un canal de communication
- Nécessité d'une modulation



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- Nécessité d'une modulation
- Modèle simplifié

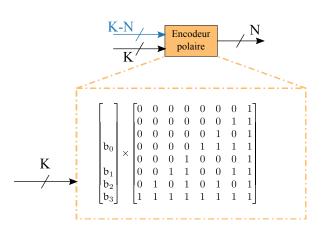


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- Nécessité d'une modulation
- Modèle simplifié
- Correction d'erreurs par ajout de redondance (N > K)

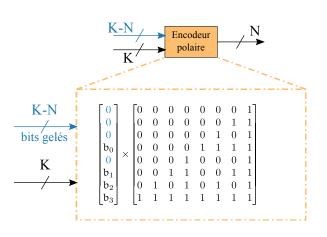


- Message transmis au travers d'un canal de communication
- Nécessité d'une modulation
- Modèle simplifié
- Correction d'erreurs par ajout de redondance (N > K)
- Le cas des codes polaires

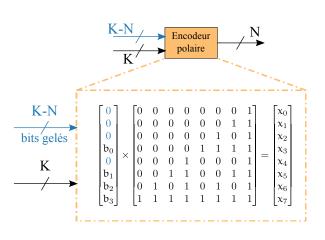
- 1948 : Théorie de l'information (Shannon)
- 1950 : Codes de Hamming
- 1955 : Codes convolutifs
- 1960 : Codes BCH
- 1960 : Codes Reed-Solomon
- 1960 : Codes LDPC
- 1966 : Codes concaténés (Forney)
- 1993 : Codes turbos
- 1996 : LDPC rediscovery
- 2008 : Codes polaires



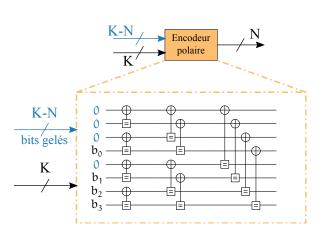
Matrice d'encodage



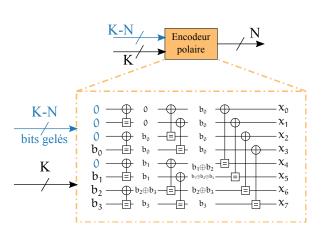
- Matrice d'encodage
- Bits gelés



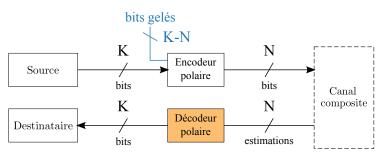
- Matrice d'encodage
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- Mot de code : x_i



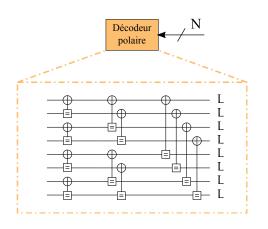
- Matrice d'encodage
- Bits gelés
- Mot de code : x_i
- Graphe de factorisation



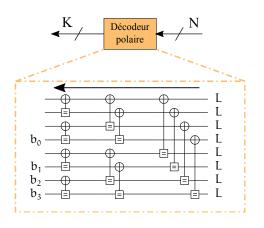
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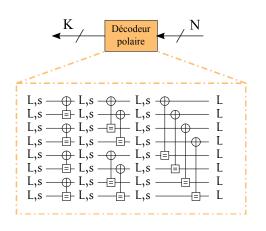
- Estimations : LLR
- Signe: valeur binaire la plus probable
- Valeur absolue : fiabilité de l'information



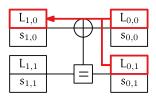
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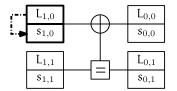


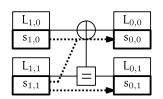
- L : Log Likelihood Ratios (LLR)
- s : Sommes Partielles



$$f(L_a,L_b) pprox \mathrm{sign}(L_a.L_b).\min(|L_a|,|L_b|)$$

$$g(L_a, L_b, \hat{s}_a) = (1 - 2\hat{s}_a)L_a + L_b$$

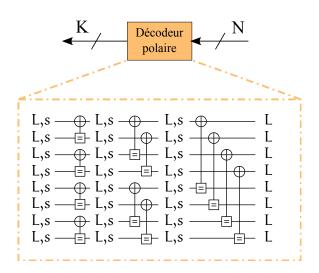


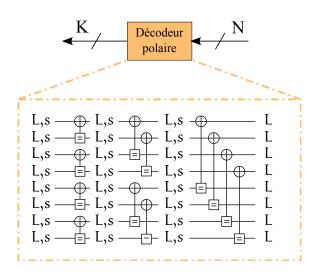


$$exttt{R1}(L_a) = \left\{egin{array}{l} 0 ext{ si } L_a \geqslant 0 \ 1 ext{ si } L_a < 0 \end{array}
ight.$$

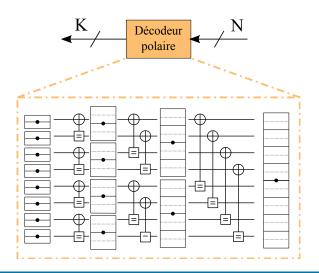
$$h(\hat{\mathbf{s}}_a,\hat{\mathbf{s}}_b)=(\hat{\mathbf{s}}_a\oplus\hat{\mathbf{s}}_b,\hat{\mathbf{s}}_b)$$

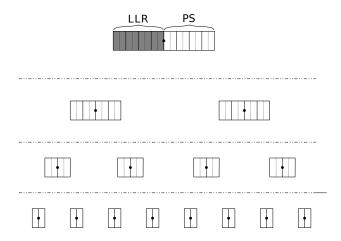
Le décodage SC

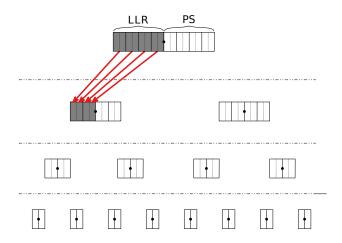


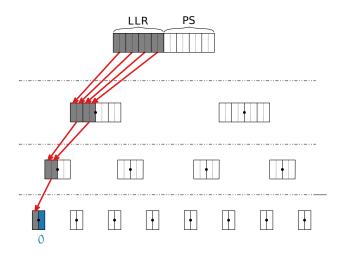


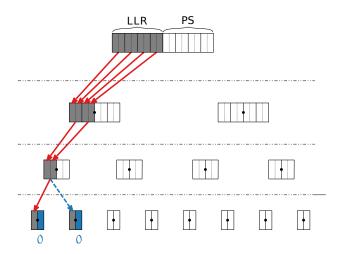
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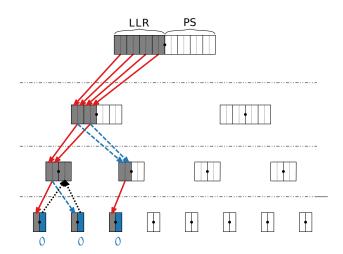


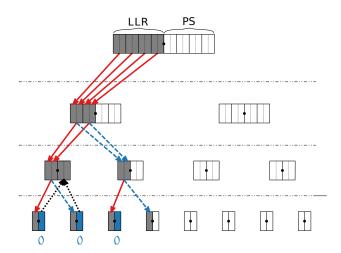


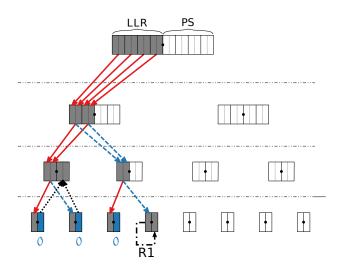


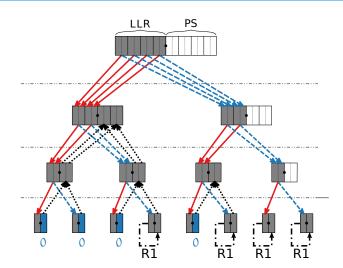


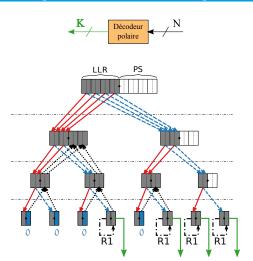






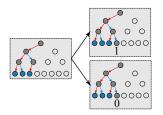




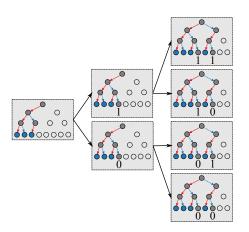




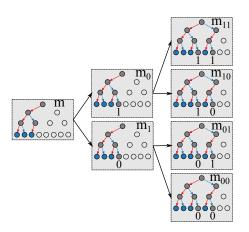




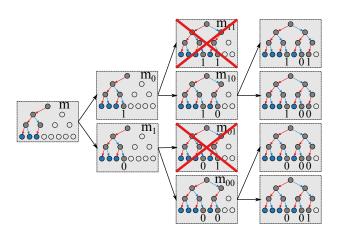
 Pas de décision dure



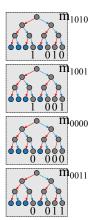
- Pas de décision dure
- **Duplication des** chemins



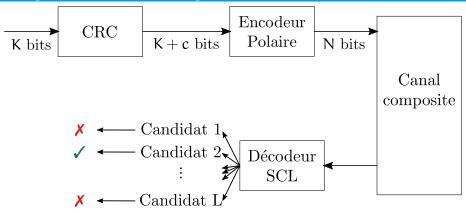
- Pas de décision dure
- Duplication des chemins
- Métrique de chemins



- Pas de décision dure
 - Duplication des chemins
- Métrique de chemins
- Tri des métrique
 & élimination



- Pas de décision dure
- Duplication des chemins
- Métrique de chemins
- Tri des métrique
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- Mot de code décodé

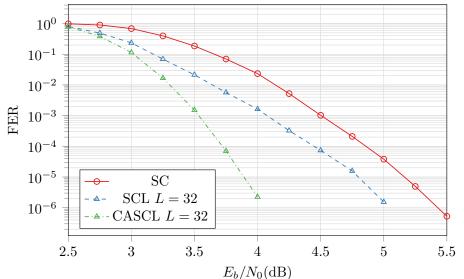


Vérification

CRC

Distance faible des codes polaires

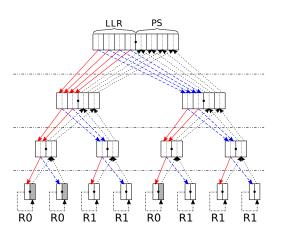
Performances de décodage

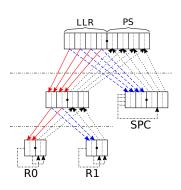


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Élagage de l'arbre





Algorithme	Performances	Débit	Latence
de décodage	BER & FER		Maximum
SC	faibles	haut	faible
SCL	moyennes	bas	moyenne
CASCL	élevées	bas	moyenne
Adaptive-SCL	élevées	haut	forte

Résumé

- Chaîne de communications
- Role des codes correcteurs d'erreurs
- Codes polaires
- Algorithme de décodage SC
- Algorithme de décodage SCL (et CASCL)
- Compromis entre performances de décodage, débit et latence