Principles of Al Planning

5. Planning as search: progression and regression

Bernhard Nebel and Robert Mattmüller

Albert-Ludwigs-Universität Freiburg

November 4th, 2011

Al Planning

B. Nebel, R. Mattmüller

Searcn

Progression

egression

Planning as (classical) search

Al Planning

B. Nebel, R. Mattmüller

Search

Introduction Classification

Progression

Regression

What do we mean by search?

- Search is a very generic term.
- Every algorithm that tries out various alternatives can be said to "search" in some way.
 - Here, we mean classical search algorithms.
 - Search nodes are expanded to generate successor nodes.
 - Examples: breadth-first search, A*, hill-climbing, ...
 - To be brief, we just say search in the following (not "classical search").

Al Planning

B. Nebel, R. Mattmüller

Search

Introduction Classification

Progression

Regression

Do you know this stuff already?

- We assume prior knowledge of basic search algorithms:
 - uninformed vs. informed
 - systematic vs. local
- There will be a small refresher in the next chapter.
- Background: Russell & Norvig, Artificial Intelligence –
 A Modern Approach, Ch. 3 (all of it), Ch. 4 (local search)

Al Planning

B. Nebel, R. Mattmüller

Introduction Classification

Progression

Regression

Search in planning

- search: one of the big success stories of AI
- many planning algorithms based on classical Al search (we'll see some other algorithms later, though)
- will be the focus of this and the following chapters (the majority of the course)

Al Planning

B. Nebel, R. Mattmüller

Search

Introduction Classification

Progression

Regression

Satisficing or optimal planning?

Must carefully distinguish two different problems:

- satisficing planning: any solution is OK (although shorter solutions typically preferred)
- optimal planning: plans must have shortest possible length

Both are often solved by search, but:

- details are very different
- almost no overlap between good techniques for satisficing planning and good techniques for optimal planning
- many problems that are trivial for satisficing planners are impossibly hard for optimal planners

Al Planning

B. Nebel, R. Mattmüller

Search

Introduction Classification

Progression

legression

How to apply search to planning? → many choices to make!

Choice 1: Search direction

- progression: forward from initial state to goal
- regression: backward from goal states to initial state
- bidirectional search

Al Planning

B. Nebel, R. Mattmüller

Search

Classification

.

How to apply search to planning? → many choices to make!

Choice 2: Search space representation

- search nodes are associated with states
 (→ state-space search)
- search nodes are associated with sets of states

Al Planning

B. Nebel, R. Mattmüller

Search

Classification

Regression

How to apply search to planning? → many choices to make!

Choice 3: Search algorithm

- uninformed search: depth-first, breadth-first, iterative depth-first, . . .
- heuristic search (systematic):
 greedy best-first, A*, Weighted A*, IDA*, ...
- heuristic search (local):
 hill-climbing, simulated annealing, beam search, . . .

Al Planning

B. Nebel, R. Mattmüller

Search

Classification

Ŭ

6

How to apply search to planning? → many choices to make!

Choice 4: Search control

- heuristics for informed search algorithms
- pruning techniques: invariants, symmetry elimination, partial-order reduction, helpful actions pruning, . . .

Al Planning

B. Nebel, R. Mattmüller

Search

Classification

. .

C

Search-based satisficing planners

FF (Hoffmann & Nebel, 2001)

- search direction: forward search
- search space representation: single states
- search algorithm: enforced hill-climbing (informed local)
- heuristic: FF heuristic (inadmissible)
- pruning technique: helpful actions (incomplete)

→ one of the best satisficing planners

Al Planning

B. Nebel, R. Mattmüller

Jean Cil

Classification

ogression

Regression

Search-based optimal planners

Fast Downward Stone Soup (Helmert et al., 2011)

- search direction: forward search
- search space representation: single states
- search algorithm: A* (informed systematic)
- heuristic: multiple admissible heuristics combined into a heuristic portfolio (LM-cut, M&S, blind, ...)
- pruning technique: none

→ one of the best optimal planners

Al Planning

B. Nebel, R. Mattmüller

Search

Classification

ogression

Regression

Our plan for the next lectures

Choices to make:

- Search algorithm: uninformed/heuristic; systematic/local → next chapter
- search control: heuristics, pruning techniques→ following chapters

Al Planning

B. Nebel, R. Mattmüller

Search

Introduction Classification

Progression

Regression

Progression

Al Planning

B. Nebel, R. Mattmülle

Searc

Progression

Example

Regression

Planning by forward search: progression

Progression: Computing the successor state $app_o(s)$ of a state s with respect to an operator o.

Progression planners find solutions by forward search:

- start from initial state
- iteratively pick a previously generated state and progress it through an operator, generating a new state
- solution found when a goal state generated

pro: very easy and efficient to implement

AI Planning

B. Nebel, R. Mattmüller

Search

Overview Example

Regression

Search space representation in progression planners

Two alternative search spaces for progression planners:

- search nodes correspond to states
 - when the same state is generated along different paths, it is not considered again (duplicate detection)
 - pro: save time to consider same state again
 - con: memory intensive (must maintain closed list)
- search nodes correspond to operator sequences
 - different operator sequences may lead to identical states (transpositions); search does not notice this
 - pro: can be very memory-efficient
 - con: much wasted work (often exponentially slower)

 \leadsto first alternative usually preferable in planning (unlike many classical search benchmarks like 15-puzzle)

Al Planning

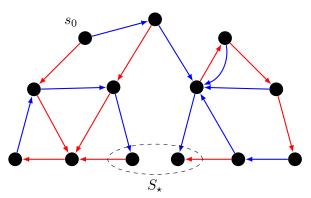
B. Nebel, R. Mattmüller

Search

Overview

Regression

Example where search nodes correspond to operator sequences (no duplicate detection)



Al Planning

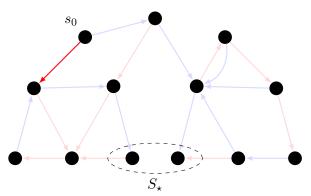
B. Nebel, R. Mattmüller

Search

Progressio Overview Example

Regression

Example where search nodes correspond to operator sequences (no duplicate detection)



Al Planning

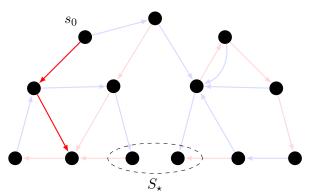
B. Nebel, R. Mattmüller

Search

Progressioi Overview Example

Regression

Example where search nodes correspond to operator sequences (no duplicate detection)



Al Planning

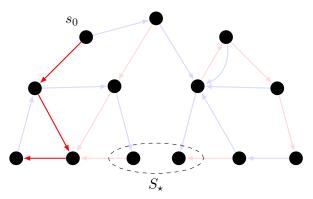
B. Nebel, R. Mattmüller

Search

Progressio Overview Example

Regression

Example where search nodes correspond to operator sequences (no duplicate detection)



Al Planning

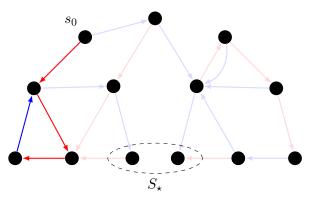
B. Nebel, R. Mattmüller

Search

Progressio Overview Example

Regression

Example where search nodes correspond to operator sequences (no duplicate detection)



Al Planning

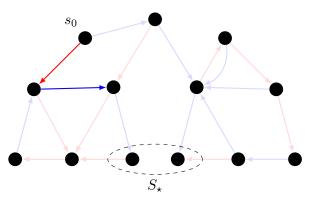
B. Nebel, R. Mattmüller

Search

Progressio Overview Example

Regression

Example where search nodes correspond to operator sequences (no duplicate detection)



Al Planning

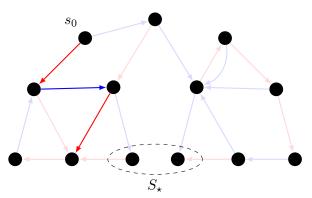
B. Nebel, R. Mattmüller

Search

Progressio Overview Example

Regression

Example where search nodes correspond to operator sequences (no duplicate detection)



Al Planning

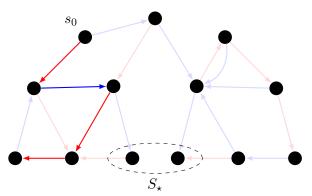
B. Nebel, R. Mattmüller

Search

Progressio Overview Example

Regression

Example where search nodes correspond to operator sequences (no duplicate detection)



Al Planning

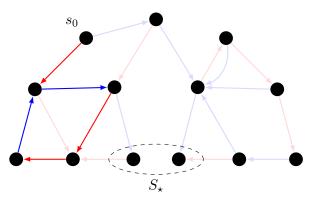
B. Nebel, R. Mattmüller

Search

Progressio Overview Example

Regression

Example where search nodes correspond to operator sequences (no duplicate detection)



Al Planning

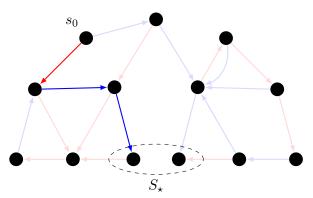
B. Nebel, R. Mattmüller

Search

Progressio Overview Example

Regression

Example where search nodes correspond to operator sequences (no duplicate detection)



Al Planning

B. Nebel, R. Mattmüller

Search

Progressio Overview Example

Regression

Regression

Al Planning

B. Nebel, R. Mattmülle

Search

Progression

Regression

Example STRIPS General case

Forward search vs. backward search

Going through a transition graph in forward and backward directions is **not symmetric**:

- forward search starts from a single initial state;
 backward search starts from a set of goal states
- when applying an operator o in a state s in forward direction, there is a unique successor state s'; if we applied operator o to end up in state s', there can be several possible predecessor states s

→ most natural representation for backward search in planning associates sets of states with search nodes

AI Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression
Overview
Example
STRIPS
General case
Practical issues

Planning by backward search: regression

Regression: Computing the possible predecessor states $regr_o(G)$ of a set of states G with respect to the last operator o that was applied.

Regression planners find solutions by backward search:

- start from set of goal states
- iteratively pick a previously generated state set and regress it through an operator, generating a new state set
- solution found when a generated state set includes the initial state

Pro: can handle many states simultaneously
Con: basic operations complicated and expensive

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Overview
Example
STRIPS
General case
Practical issues

Search space representation in regression planners

identify state sets with logical formulae (again):

- search nodes correspond to state sets
- each state set is represented by a logical formula: φ represents $\{s \in S \mid s \models \varphi\}$
- many basic search operations like detecting duplicates are NP-hard or coNP-hard

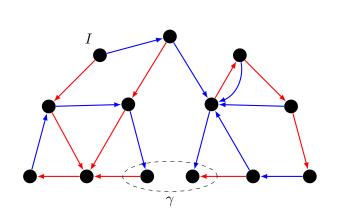
Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression
Overview
Example
STRIPS
General case



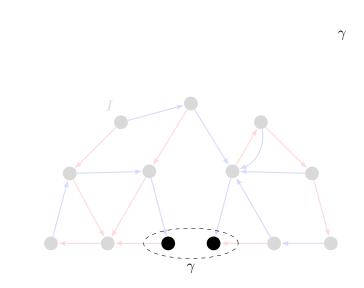
Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Overview
Example
STRIPS
General case



Al Planning

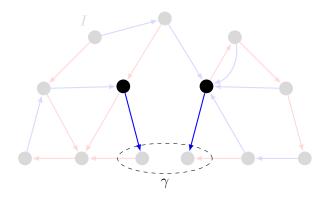
B. Nebel, R. Mattmüller

Search

Progressior

Overview
Example
STRIPS
General case
Practical issues

$$\varphi_1 = \mathit{regr}_{\longrightarrow}(\gamma)$$
 $\qquad \qquad \varphi_1 \longrightarrow \gamma$



Al Planning

B. Nebel, R. Mattmüller

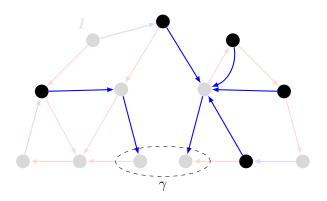
Search

Progression

Regression
Overview
Example
STRIPS
General case
Practical issues

$$\varphi_1 = \operatorname{regr}_{\longrightarrow}(\gamma)$$
$$\varphi_2 = \operatorname{regr}_{\longrightarrow}(\varphi_1)$$

$$\varphi_2 \longrightarrow \varphi_1 \longrightarrow \gamma$$



Al Planning

B. Nebel, R. Mattmüller

Search

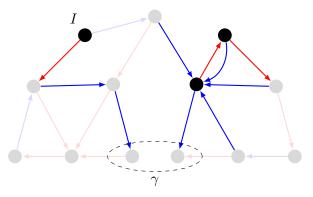
Progression

Overview
Example
STRIPS
General case
Practical issues

$$\varphi_{1} = \operatorname{regr}_{\longrightarrow}(\gamma) \qquad \varphi_{3} \longrightarrow \varphi_{2} \longrightarrow \varphi_{1} \longrightarrow \gamma$$

$$\varphi_{2} = \operatorname{regr}_{\longrightarrow}(\varphi_{1})$$

$$\varphi_{3} = \operatorname{regr}_{\longrightarrow}(\varphi_{2}), I \models \varphi_{3}$$



Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Overview

Example
STRIPS
General case
Practical issues

Regression for STRIPS planning tasks

Definition (STRIPS planning task)

A planning task is a STRIPS planning task if all operators are STRIPS operators and the goal is a conjunction of atoms.

Regression for STRIPS planning tasks is very simple:

- Goals are conjunctions of atoms $a_1 \wedge \cdots \wedge a_n$.
- First step: Choose an operator that makes none of a_1, \ldots, a_n false.
- Second step: Remove goal atoms achieved by the operator (if any) and add its preconditions.
- \rightsquigarrow Outcome of regression is again conjunction of atoms.

Optimization: only consider operators making some a_i true

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Overview
Example
STRIPS
General case

STRIPS regression

Definition (STRIPS regression)

Let $\varphi = \varphi_1 \wedge \cdots \wedge \varphi_n$ be a conjunction of atoms, and let $o = \langle \chi, e \rangle$ be a STRIPS operator which adds the atoms a_1, \ldots, a_k and deletes the atoms d_1, \ldots, d_l .

The STRIPS regression of φ with respect to o is

$$\mathit{sregr}_o(\varphi) := \begin{cases} \bot & \text{if } a_i = d_j \text{ for some } i, j \\ \bot & \text{if } \varphi_i = d_j \text{ for some } i, j \\ \chi \land \bigwedge(\{\varphi_1, \dots, \varphi_n\} \setminus \{a_1, \dots, a_k\}) & \text{otherwise} \end{cases}$$

Note: $sregr_o(\varphi)$ is again a conjunction of atoms, or \bot .

AI Planning

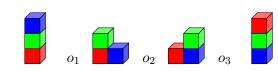
B. Nebel, R. Mattmüller

Jearch

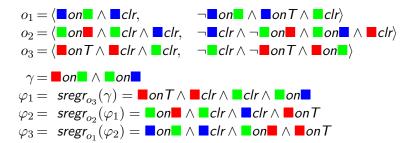
Progression

Overview
Example
STRIPS
General case
Practical issues

STRIPS regression example



Note: Predecessor states are in general not unique. This picture is just for illustration purposes.



Al Planning

B. Nebel, R. Mattmüller

bearch

Progression

Overview
Example
STRIPS
General case

Regression for general planning tasks

- With disjunctions and conditional effects, things become more tricky. How to regress $a \lor (b \land c)$ with respect to $\langle q, d \rhd b \rangle$?
- The story about goals and subgoals and fulfilling subgoals, as in the STRIPS case, is no longer useful.
- We present a general method for doing regression for any formula and any operator.
- Now we extensively use the idea of representing sets of states as formulae.

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Overview
Example
STRIPS
General case

Practical issu

Effect preconditions

Definition (effect precondition)

The effect precondition $EPC_l(e)$ for literal l and effect e is defined as follows:

$$\begin{array}{rcl} \textit{EPC}_l(l) & = & \top \\ \textit{EPC}_l(l') & = & \bot \text{ if } l \neq l' \quad \text{(for literals } l') \\ \textit{EPC}_l(e_1 \wedge \cdots \wedge e_n) & = & \textit{EPC}_l(e_1) \vee \cdots \vee \textit{EPC}_l(e_n) \\ \textit{EPC}_l(\chi \rhd e) & = & \textit{EPC}_l(e) \wedge \chi \end{array}$$

Intuition: $EPC_l(e)$ describes the situations in which effect e causes literal l to become true.

AI Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression

STRIPS
General case

Effect precondition examples

Example

$$\begin{array}{rcl} \textit{EPC}_a(b \wedge c) & = & \bot \vee \bot \equiv \bot \\ \textit{EPC}_a(a \wedge (b \rhd a)) & = & \top \vee (\top \wedge b) \equiv \top \\ \textit{EPC}_a((c \rhd a) \wedge (b \rhd a)) & = & (\top \wedge c) \vee (\top \wedge b) \equiv c \vee b \end{array}$$

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Overview Example STRIPS

General case Practical issues

Lemma (A)

Let s be a state, l a literal and e an effect. Then $l \in [e]_s$ if and only if $s \models \textit{EPC}_l(e)$.

Proof.

Induction on the structure of the effect e.

Base case 1, e = l: $l \in [l]_s = \{l\}$ by definition, and $s \models EPC_l(l) = \top$ by definition. Both sides of the equivalence are true.

Base case 2, e=l' for some literal $l' \neq l$: $l \notin [l']_s = \{l'\}$ by definition, and $s \not\models EPC_l(l') = \bot$ by definition. Both sides are false.

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression
Overview
Example

General case Practical issues

Lemma (A)

Let s be a state, l a literal and e an effect. Then $l \in [e]_s$ if and only if $s \models EPC_l(e)$.

Proof.

Induction on the structure of the effect e.

Base case 1, e=l: $l\in [l]_s=\{l\}$ by definition, and $s\models \mathit{EPC}_l(l)=\top$ by definition. Both sides of the equivalence are true.

Base case 2, e=l' for some literal $l' \neq l$: $l \notin [l']_s = \{l'\}$ by definition, and $s \not\models EPC_l(l') = \bot$ by definition. Both sides are false.

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression
Overview
Example
STRIPS

General case Practical issues

Lemma (A)

Let s be a state, l a literal and e an effect.

Then $l \in [e]_s$ if and only if $s \models EPC_l(e)$.

Proof.

Induction on the structure of the effect e.

Base case 1, e=l: $l \in [l]_s = \{l\}$ by definition, and $s \models \mathit{EPC}_l(l) = \top$ by definition. Both sides of the equivalence are true.

Base case 2, e=l' for some literal $l' \neq l$: $l \notin [l']_s = \{l'\}$ by definition, and $s \not\models EPC_l(l') = \bot$ by definition. Both sides are false.

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression
Overview
Example
STRIPS

General case Practical issues

```
Proof (ctd.)
Inductive case 1, e = e_1 \wedge \cdots \wedge e_n:
 l \in [e]_s iff l \in [e_1]_s \cup \cdots \cup [e_n]_s (Def [e_1 \wedge \cdots \wedge e_n]_s)
            iff l \in [e']_s for some e' \in \{e_1, \dots, e_n\}
            iff s \models EPC_l(e') for some e' \in \{e_1, \dots, e_n\}
                                                                                    (IH)
            iff s \models EPC_l(e_1) \lor \cdots \lor EPC_l(e_n)
                                                                           (Def EPC)
            iff s \models EPC_l(e_1 \land \cdots \land e_n).
```

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Overview
Example

General case Practical issues

```
Proof (ctd.)
```

```
Inductive case 1, e = e_1 \wedge \cdots \wedge e_n: l \in [e]_s \text{ iff } l \in [e_1]_s \cup \cdots \cup [e_n]_s \qquad \qquad \left( \begin{array}{c} \mathsf{Def} \ [e_1 \wedge \cdots \wedge e_n]_s \right) \\ \mathsf{iff} \ l \in [e']_s \ \mathsf{for some} \ e' \in \{e_1, \ldots, e_n\} \\ \mathsf{iff} \ s \models \mathit{EPC}_l(e') \ \mathsf{for some} \ e' \in \{e_1, \ldots, e_n\} \\ \mathsf{iff} \ s \models \mathit{EPC}_l(e_1) \vee \cdots \vee \mathit{EPC}_l(e_n) \\ \mathsf{iff} \ s \models \mathit{EPC}_l(e_1 \wedge \cdots \wedge e_n). \qquad \qquad \left( \begin{array}{c} \mathsf{Def} \ \mathit{EPC} \right) \\ \mathsf{Def} \ \mathit{EPC} \right) \end{array}
```

Inductive case 2, $e = \chi \triangleright e'$:

```
\begin{split} l \in [\chi \rhd e']_s & \text{ iff } l \in [e']_s \text{ and } s \models \chi \\ & \text{ iff } s \models \textit{EPC}_l(e') \text{ and } s \models \chi \\ & \text{ iff } s \models \textit{EPC}_l(e') \land \chi \\ & \text{ iff } s \models \textit{EPC}_l(\chi \rhd e'). \end{split} \tag{Def } \underbrace{[\chi \rhd e']_s}_{}
```

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression
Overview
Example

General case Practical issues

Effect preconditions: connection to normal form

Remark: EPC vs. effect normal form

Notice that in terms of $EPC_a(e)$, any operator $\langle \chi, e \rangle$ can be expressed in effect normal form as

$$\left\langle \chi, \bigwedge_{a \in A} \left((\mathit{EPC}_a(e) \rhd a) \land (\mathit{EPC}_{\neg a}(e) \rhd \neg a) \right) \right\rangle$$

where A is the set of all state variables.

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression

Example STRIPS General case

Practical issue

Regressing state variables

The formula $EPC_a(e) \lor (a \land \neg EPC_{\neg a}(e))$ expresses the value of state variable $a \in A$ after applying o in terms of values of state variables before applying o.

Fither:

- ullet a became true, or
- a was true before and it did not become false.

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression

Example STRIPS General case

Practical issue

Regressing state variables: examples

Example

Let $e = (b \rhd a) \land (c \rhd \neg a) \land b \land \neg d$.

$variable\ x$	$EPC_x(e) \lor (x \land \neg EPC_{\neg x}(e))$
\overline{a}	$b \lor (a \land \neg c)$
b	$\top \lor (b \land \neg \bot) \equiv \top$
c	$\perp \lor (c \land \neg \bot) \equiv c$
d	$\bot \lor (d \land \neg \top) \equiv \bot$

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression

Overview Example STRIPS

General case Practical issues

Lemma (B)

Let a be a state variable, $o = \langle \chi, e \rangle$ an operator, s a state, and $s' = \mathsf{app}_o(s)$. Then $s \models \mathsf{EPC}_a(e) \lor (a \land \neg \mathsf{EPC}_{\neg a}(e))$ if and only if $s' \models a$.

Proof.

(⇒): Assume $s \models EPC_a(e) \lor (a \land \neg EPC_{\neg a}(e))$. Do a case analysis on the two disjuncts.

- ① Assume that $s \models EPC_a(e)$. By Lemma A, we have $a \in [e]_s$ and hence $s' \models a$.
- ② Assume that $s \models a \land \neg EPC_{\neg a}(e)$. By Lemma A, we have $\neg a \notin [e]_s$. Hence a remains true in s'.

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression

Example STRIPS General case

ractical issu

Lemma (B)

Let a be a state variable, $o=\langle \chi,e\rangle$ an operator, s a state, and $s'=\operatorname{app}_o(s)$.

Then $s \models \mathsf{EPC}_a(e) \lor (a \land \neg \mathsf{EPC}_{\neg a}(e))$ if and only if $s' \models a$.

Proof.

 (\Rightarrow) : Assume $s \models EPC_a(e) \lor (a \land \neg EPC_{\neg a}(e))$.

Do a case analysis on the two disjuncts.

- ① Assume that $s \models EPC_a(e)$. By Lemma A, we have $a \in [e]_s$ and hence $s' \models a$.
- ② Assume that $s \models a \land \neg EPC_{\neg a}(e)$. By Lemma A, we have $\neg a \notin [e]_s$. Hence a remains true in s'.

AI Planning

B. Nebel, R. Mattmüller

Search

Progression

Overview Example

General case Practical issues

Lemma (B)

Let a be a state variable, $o = \langle \chi, e \rangle$ an operator, s a state, and $s' = \mathsf{app}_o(s)$. Then $s \models \mathsf{EPC}_a(e) \lor (a \land \neg \mathsf{EPC}_{\neg a}(e))$ if and only if $s' \models a$.

Proof.

 (\Rightarrow) : Assume $s \models EPC_a(e) \lor (a \land \neg EPC_{\neg a}(e))$.

Do a case analysis on the two disjuncts.

- Assume that $s \models EPC_a(e)$. By Lemma A, we have $a \in [e]_s$ and hence $s' \models a$.
- ② Assume that $s \models a \land \neg EPC_{\neg a}(e)$. By Lemma A, we have $\neg a \notin [e]_s$. Hence a remains true in s'.

AI Planning

B. Nebel, R. Mattmüller

Search

Progression

Overview
Example
STRIPS
General case

Lemma (B)

Let a be a state variable, $o = \langle \chi, e \rangle$ an operator, s a state, and $s' = \mathsf{app}_o(s)$. Then $s \models \mathsf{EPC}_a(e) \lor (a \land \neg \mathsf{EPC}_{\neg a}(e))$ if and only if $s' \models a$.

Proof.

 (\Rightarrow) : Assume $s \models EPC_a(e) \lor (a \land \neg EPC_{\neg a}(e))$.

Do a case analysis on the two disjuncts.

- **1** Assume that $s \models EPC_a(e)$. By Lemma A, we have $a \in [e]_s$ and hence $s' \models a$.
- ② Assume that $s \models a \land \neg EPC_{\neg a}(e)$. By Lemma A, we have $\neg a \notin [e]_s$. Hence a remains true in s'.

AI Planning

B. Nebel, R. Mattmüller

Search

Progression

Overview
Example
STRIPS
General case

Proof (ctd.)

 (\Leftarrow) : We showed that if the formula is true in s, then a is true in s'. For the second part, we show that if the formula is false in s, then a is false in s'.

- So assume $s \not\models EPC_a(e) \lor (a \land \neg EPC_{\neg a}(e))$.
- Then $s \models \neg EPC_a(e) \land (\neg a \lor EPC_{\neg a}(e))$ (de Morgan).
- Case distinction: a is true or a is false in s.
 - Assume that $s \models a$. Now $s \models EPC_{\neg a}(e)$ because $s \models \neg a \lor EPC_{\neg a}(e)$. Hence by Lemma A $\neg a \in [e]_s$ and we get $s' \not\models a$.
 - ② Assume that $s \not\models a$. Because $s \models \neg EPC_a(e)$, by Lemma A we get $a \notin [e]_s$ and hence $s' \not\models a$.

Therefore in both cases $s' \not\models a$.

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression

Example STRIPS General case

General case Practical issues

Proof (ctd.)

 (\Leftarrow) : We showed that if the formula is true in s, then a is true in s'. For the second part, we show that if the formula is false in s, then a is false in s'.

- So assume $s \not\models EPC_a(e) \lor (a \land \neg EPC_{\neg a}(e))$.
- Then $s \models \neg EPC_a(e) \land (\neg a \lor EPC_{\neg a}(e))$ (de Morgan).
- Case distinction: a is true or a is false in s.
 - ① Assume that $s \models a$. Now $s \models EPC_{\neg a}(e)$ because $s \models \neg a \lor EPC_{\neg a}(e)$. Hence by Lemma A $\neg a \in [e]_s$ and we get $s' \not\models a$.
 - ② Assume that $s \not\models a$. Because $s \models \neg EPC_a(e)$, by Lemma A we get $a \notin [e]_s$ and hence $s' \not\models a$.

Therefore in both cases $s' \not\models a$.

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Overview
Example

General case Practical issues

Proof (ctd.)

 (\Leftarrow) : We showed that if the formula is true in s, then a is true in s'. For the second part, we show that if the formula is false in s, then a is false in s'.

- So assume $s \not\models EPC_a(e) \lor (a \land \neg EPC_{\neg a}(e))$.
- Then $s \models \neg EPC_a(e) \land (\neg a \lor EPC_{\neg a}(e))$ (de Morgan).

Al Planning

B. Nebel. R. Mattmüller

General case

Proof (ctd.)

 (\Leftarrow) : We showed that if the formula is true in s, then a is true in s'. For the second part, we show that if the formula is false in s, then a is false in s'.

- So assume $s \not\models EPC_a(e) \lor (a \land \neg EPC_{\neg a}(e))$.
- Then $s \models \neg EPC_a(e) \land (\neg a \lor EPC_{\neg a}(e))$ (de Morgan).
- Case distinction: a is true or a is false in s.
 - ① Assume that $s \models a$. Now $s \models EPC_{\neg a}(e)$ because $s \models \neg a \lor EPC_{\neg a}(e)$. Hence by Lemma A $\neg a \in [e]_s$ and we get $s' \not\models a$.
 - ② Assume that $s \not\models a$. Because $s \models \neg EPC_a(e)$, by Lemma A we get $a \notin [e]_s$ and hence $s' \not\models a$.

Therefore in both cases $s' \not\models a$.

AI Planning

B. Nebel, R. Mattmüller

Search

Progression

Overview
Example
STRIPS

General case Practical issues

Proof (ctd.)

 (\Leftarrow) : We showed that if the formula is true in s, then a is true in s'. For the second part, we show that if the formula is false in s, then a is false in s'.

- So assume $s \not\models EPC_a(e) \lor (a \land \neg EPC_{\neg a}(e))$.
- Then $s \models \neg EPC_a(e) \land (\neg a \lor EPC_{\neg a}(e))$ (de Morgan).
- Case distinction: a is true or a is false in s.
 - **1** Assume that $s \models a$. Now $s \models EPC_{\neg a}(e)$ because $s \models \neg a \lor EPC_{\neg a}(e)$. Hence by Lemma A $\neg a \in [e]_s$ and we get $s' \not\models a$.
 - ② Assume that $s \not\models a$. Because $s \models \neg EPC_a(e)$, by Lemma A we get $a \notin [e]_s$ and hence $s' \not\models a$.

Therefore in both cases $s' \not\models a$.

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Overview
Example
STRIPS
General case

Proof (ctd.)

 (\Leftarrow) : We showed that if the formula is true in s, then a is true in s'. For the second part, we show that if the formula is false in s, then a is false in s'.

- So assume $s \not\models EPC_a(e) \lor (a \land \neg EPC_{\neg a}(e))$.
- Then $s \models \neg EPC_a(e) \land (\neg a \lor EPC_{\neg a}(e))$ (de Morgan).
- Case distinction: a is true or a is false in s.
 - **1** Assume that $s \models a$. Now $s \models EPC_{\neg a}(e)$ because $s \models \neg a \lor EPC_{\neg a}(e)$. Hence by Lemma A $\neg a \in [e]_s$ and we get $s' \not\models a$.
 - ② Assume that $s \not\models a$. Because $s \models \neg EPC_a(e)$, by Lemma A we get $a \notin [e]_s$ and hence $s' \not\models a$.

Therefore in both cases $s' \not\models a$

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Overview
Example
STRIPS
General case

Proof (ctd.)

 (\Leftarrow) : We showed that if the formula is true in s, then a is true in s'. For the second part, we show that if the formula is false in s, then a is false in s'.

- So assume $s \not\models EPC_a(e) \lor (a \land \neg EPC_{\neg a}(e))$.
- Then $s \models \neg EPC_a(e) \land (\neg a \lor EPC_{\neg a}(e))$ (de Morgan).
- Case distinction: a is true or a is false in s.
 - **1** Assume that $s \models a$. Now $s \models EPC_{\neg a}(e)$ because $s \models \neg a \lor EPC_{\neg a}(e)$. Hence by Lemma A $\neg a \in [e]_s$ and we get $s' \not\models a$.
 - ② Assume that $s \not\models a$. Because $s \models \neg EPC_a(e)$, by Lemma A we get $a \notin [e]_s$ and hence $s' \not\models a$.

Therefore in both cases $s' \not\models a$.

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Overview
Example
STRIPS
General case

Regression: general definition

We base the definition of regression on formulae $EPC_l(e)$.

Definition (general regression)

Let φ be a propositional formula and $o=\langle \chi,e\rangle$ an operator. The regression of φ with respect to o is

$$regr_o(\varphi) = \chi \wedge \varphi_{\mathsf{r}} \wedge \kappa$$

where

① φ_{r} is obtained from φ by replacing each $a \in A$ by $\mathit{EPC}_a(e) \lor (a \land \neg \mathit{EPC}_{\neg a}(e))$, and

The formula κ expresses that operators are only applicable in states where their change sets are consistent.

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression
Overview
Example
STRIPS

General case Practical issues

Regression examples

$$\bullet \ \operatorname{regr}_{\langle a,b\rangle}(b) \equiv a \wedge (\top \vee (b \wedge \neg \bot)) \wedge \top \equiv a$$

$$\begin{array}{l} \bullet \;\; \mathit{regr}_{\langle a,b\rangle}(b \wedge c \wedge d) \\ \equiv a \wedge (\top \vee (b \wedge \neg \bot)) \wedge (\bot \vee (c \wedge \neg \bot)) \wedge (\bot \vee (d \wedge \neg \bot)) \wedge \top \\ \equiv a \wedge c \wedge d \end{array}$$

- $\bullet \ \operatorname{regr}_{\langle a,c\rhd b\rangle}(b)\equiv a\wedge (c\vee (b\wedge\neg\bot))\wedge\top\equiv a\wedge (c\vee b)$
- $regr_{\langle a,(c\rhd b)\wedge(b\rhd \neg b)\rangle}(b)\equiv a\wedge(c\vee(b\wedge\neg b))\wedge\neg(c\wedge b)$ $\equiv a\wedge c\wedge\neg b$
- $regr_{\langle a,(c\rhd b)\land (d\rhd \neg b)\rangle}(b) \equiv a\land (c\lor (b\land \neg d))\land \neg (c\land d)$ $\equiv a\land (c\lor b)\land (c\lor \neg d)\land (\neg c\lor \neg d)$ $\equiv a\land (c\lor b)\land \neg d$

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression

Example STRIPS General case

Practical issues

oummary

Regression example: binary counter

$$(\neg b_0 \rhd b_0) \land \\ ((\neg b_1 \land b_0) \rhd (b_1 \land \neg b_0)) \land \\ ((\neg b_2 \land b_1 \land b_0) \rhd (b_2 \land \neg b_1 \land \neg b_0))$$

$$EPC_{b_2}(e) = \neg b_2 \land b_1 \land b_0 \\ EPC_{b_1}(e) = \neg b_1 \land b_0 \\ EPC_{b_0}(e) = \neg b_0 \\ EPC_{\neg b_2}(e) = \bot \\ EPC_{\neg b_1}(e) = \neg b_2 \land b_1 \land b_0 \\ EPC_{\neg b_0}(e) = (\neg b_1 \land b_0) \lor (\neg b_2 \land b_1 \land b_0) \equiv (\neg b_1 \lor \neg b_2) \land b_0$$

$$EPC_{\neg b_0}(e) = (\neg b_1 \land b_0) \lor (\neg b_2 \land b_1 \land b_0) \equiv (\neg b_1 \lor \neg b_2) \land b_0$$

Regression replaces state variables as follows:

$$\begin{array}{lll} b_2 & \text{by} & (\neg b_2 \wedge b_1 \wedge b_0) \vee (b_2 \wedge \neg \bot) \equiv (b_1 \wedge b_0) \vee b_2 \\ b_1 & \text{by} & (\neg b_1 \wedge b_0) \vee (b_1 \wedge \neg (\neg b_2 \wedge b_1 \wedge b_0)) \\ & & \equiv (\neg b_1 \wedge b_0) \vee (b_1 \wedge (b_2 \vee \neg b_0)) \\ b_0 & \text{by} & \neg b_0 \vee (b_0 \wedge \neg ((\neg b_1 \vee \neg b_2) \wedge b_0)) \equiv \neg b_0 \vee (b_1 \wedge b_2) \end{array}$$

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression Overview Example

STRIPS General case Practical issues

Theorem (correctness of $regr_o(\varphi)$)

Let φ be a formula, o an operator and s a state. Then $s \models \mathsf{regr}_o(\varphi)$ iff o is applicable in s and $\mathsf{app}_o(s) \models \varphi$.

Proof.

Let $o=\langle \chi,e\rangle$. Recall that $\operatorname{regr}_o(\varphi)=\chi\wedge\varphi_{\mathsf{r}}\wedge\kappa$, where φ_{r} and κ are as defined previously.

If o is inapplicable in s, then $s \not\models \chi \land \kappa$, both sides of the "iff" condition are false, and we are done. Hence, we only further consider states s where o is applicable. Let $s' := app_o(s)$.

We know that $s \models \chi \land \kappa$ (because o is applicable), so the "iff" condition we need to prove simplifies to:

$$s \models \varphi_{\mathsf{r}} \text{ iff } s' \models \varphi.$$

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression
Overview
Example

General case Practical issues

Theorem (correctness of $regr_o(\varphi)$)

Let φ be a formula, o an operator and s a state. Then $s \models regr_o(\varphi)$ iff o is applicable in s and $app_o(s) \models \varphi$.

Proof.

Let $o=\langle\chi,e\rangle$. Recall that $\mathit{regr}_o(\varphi)=\chi\wedge\varphi_\mathsf{r}\wedge\kappa$, where φ_r and κ are as defined previously.

If o is inapplicable in s, then $s \not\models \chi \land \kappa$, both sides of the "iff" condition are false, and we are done. Hence, we only further consider states s where o is applicable. Let $s' := app_o(s)$.

We know that $s \models \chi \land \kappa$ (because o is applicable), so the "iff" condition we need to prove simplifies to:

$$s \models \varphi_{\mathsf{r}} \text{ iff } s' \models \varphi.$$

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression
Overview
Example
STRIPS

General case Practical issues

Theorem (correctness of $regr_o(\varphi)$)

Let φ be a formula, o an operator and s a state. Then $s \models \mathit{regr}_o(\varphi)$ iff o is applicable in s and $\mathit{app}_o(s) \models \varphi$.

Proof.

Let $o=\langle \chi,e \rangle$. Recall that $\mathit{regr}_o(\varphi)=\chi \wedge \varphi_{\mathsf{r}} \wedge \kappa$, where φ_{r} and κ are as defined previously.

If o is inapplicable in s, then $s \not\models \chi \land \kappa$, both sides of the "iff" condition are false, and we are done. Hence, we only further consider states s where o is applicable. Let $s' := \mathsf{app}_o(s)$.

We know that $s \models \chi \land \kappa$ (because o is applicable), so the "iff" condition we need to prove simplifies to:

$$s \models \varphi_{\mathsf{r}} \text{ iff } s' \models \varphi.$$

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression
Overview
Example
STRIPS

General case Practical issues

Theorem (correctness of $regr_o(\varphi)$)

Let φ be a formula, o an operator and s a state. Then $s \models \mathit{regr}_o(\varphi)$ iff o is applicable in s and $\mathit{app}_o(s) \models \varphi$.

Proof.

Let $o=\langle \chi,e \rangle$. Recall that $\mathit{regr}_o(\varphi)=\chi \wedge \varphi_{\mathsf{r}} \wedge \kappa$, where φ_{r} and κ are as defined previously.

If o is inapplicable in s, then $s \not\models \chi \land \kappa$, both sides of the "iff" condition are false, and we are done. Hence, we only further consider states s where o is applicable. Let $s' := \textit{app}_o(s)$.

We know that $s \models \chi \land \kappa$ (because o is applicable), so the "iff" condition we need to prove simplifies to:

$$s \models \varphi_{\mathsf{r}} \text{ iff } s' \models \varphi.$$

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression
Overview
Example
STRIPS
General case

Practical issue

Proof (ctd.)

To show: $s \models \varphi_r$ iff $s' \models \varphi$.

We show that for all formulae ψ , $s \models \psi_r$ iff $s' \models \psi$, where ψ_r is ψ with every $a \in A$ replaced by $EPC_a(e) \lor (a \land \neg EPC_{\neg a}(e))$.

The proof is by structural induction on ψ .

Induction hypothesis $s \models \psi_{\mathsf{r}}$ if and only if $s' \models \psi$

Base cases 1 & 2 $\psi = \top$ or $\psi = \bot$: trivial, as $\psi_{\rm r} = \psi_{\rm r}$

Base case 3 $\psi = a$ for some $a \in A$:

Then $\psi_{\mathsf{r}} = EPC_a(e) \vee (a \wedge \neg EPC_{\neg a}(e)).$

By Lemma B, $s \models \psi_r$ iff $s' \models \psi$.

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression

Overview Example STRIPS

General case Practical issues

Proof (ctd.)

To show: $s \models \varphi_r$ iff $s' \models \varphi$.

We show that for all formulae ψ , $s \models \psi_r$ iff $s' \models \psi$, where ψ_r is ψ with every $a \in A$ replaced by $EPC_a(e) \lor (a \land \neg EPC_{\neg a}(e))$.

The proof is by structural induction on ψ .

Induction hypothesis $s \models \psi_{\mathsf{r}}$ if and only if $s' \models \psi$.

Base cases 1 & 2 $\psi = \top$ or $\psi = \bot$: trivial, as $\psi_r = \psi$.

Base case 3 $\psi = a$ for some $a \in A$:

Then $\psi_{r} = EPC_{a}(e) \lor (a \land \neg EPC_{\neg a}(e))$ By Lemma B, $e \vdash \psi_{r}$ iff $e' \vdash \psi_{r}$

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression

Overview Example STRIPS

General case Practical issues

Proof (ctd.)

To show: $s \models \varphi_r$ iff $s' \models \varphi$.

We show that for all formulae ψ , $s \models \psi_r$ iff $s' \models \psi$, where ψ_r is ψ with every $a \in A$ replaced by $EPC_a(e) \lor (a \land \neg EPC_{\neg a}(e))$.

The proof is by structural induction on ψ .

Induction hypothesis $s \models \psi_r$ if and only if $s' \models \psi$.

Base cases $1\ \&\ 2\ \psi = \top$ or $\psi = \bot$: trivial, as $\psi_{\mathsf{r}} = \psi$.

Base case 3 $\psi = a$ for some $a \in A$:

Then $\psi_r = EPC_a(e) \lor (a \land \neg EPC_{\neg a}(e))$ By Lemma B, $s \models \psi_r$ iff $s' \models \psi$. Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression

Example STRIPS General case

General case Practical issues

Proof (ctd.)

To show: $s \models \varphi_r$ iff $s' \models \varphi$.

We show that for all formulae ψ , $s \models \psi_r$ iff $s' \models \psi$, where ψ_r is ψ with every $a \in A$ replaced by $EPC_a(e) \lor (a \land \neg EPC_{\neg a}(e))$.

The proof is by structural induction on ψ .

Induction hypothesis $s \models \psi_r$ if and only if $s' \models \psi$.

Base cases 1 & 2 $\psi = \top$ or $\psi = \bot$: trivial, as $\psi_r = \psi$.

Base case 3 $\psi = a$ for some $a \in A$:

Then $\psi_r = EPC_a(e) \lor (a \land \neg EPC_{\neg a}(e))$. By Lemma B, $s \models \psi_r$ iff $s' \models \psi$. Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression

Example STRIPS General case

Practical issues

Proof (ctd.)

To show: $s \models \varphi_r$ iff $s' \models \varphi$.

We show that for all formulae ψ , $s \models \psi_r$ iff $s' \models \psi$, where ψ_r is ψ with every $a \in A$ replaced by $EPC_a(e) \lor (a \land \neg EPC_{\neg a}(e))$.

The proof is by structural induction on ψ .

Induction hypothesis $s \models \psi_r$ if and only if $s' \models \psi$.

Base cases 1 & 2 $\psi = \top$ or $\psi = \bot$: trivial, as $\psi_{\mathsf{r}} = \psi$.

Base case 3 $\psi = a$ for some $a \in A$:

Then $\psi_{\mathbf{r}} = EPC_a(e) \lor (a \land \neg EPC_{\neg a}(e))$. By Lemma B, $s \models \psi_{\mathbf{r}}$ iff $s' \models \psi$. AI Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression Overview

STRIPS
General case

Practical issue

Proof (ctd.)

Inductive case 1 $\psi = \neg \psi'$:

$$s \models \psi_{\mathsf{r}} \text{ iff } s \models (\neg \psi')_{\mathsf{r}} \text{ iff } s \models \neg (\psi'_{\mathsf{r}}) \text{ iff } s \not\models \psi'_{\mathsf{r}}$$

$$\text{iff (IH) } s' \not\models \psi' \text{ iff } s' \models \neg \psi' \text{ iff } s' \models \psi$$

Inductive case 2 $\psi = \psi' \lor \psi''$:

$$s \models \psi_{\mathsf{r}} \text{ iff } s \models (\psi' \lor \psi'')_{\mathsf{r}} \text{ iff } s \models \psi'_{\mathsf{r}} \lor \psi''_{\mathsf{r}}$$

$$\text{iff } s \models \psi'_{\mathsf{r}} \text{ or } s \models \psi''_{\mathsf{r}}$$

$$\text{iff (IH, twice) } s' \models \psi' \text{ or } s' \models \psi''$$

$$\text{iff } s' \models \psi' \lor \psi'' \text{ iff } s' \models \psi$$

Inductive case 3 $\psi=\psi'\wedge\psi''$: Very similar to inductive case 2, just with \wedge instead of \vee and "and" instead of "or".

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression

STRIPS
General case

Proof (ctd.)

Inductive case 1 $\psi = \neg \psi'$:

$$s \models \psi_{\mathsf{r}} \text{ iff } s \models (\neg \psi')_{\mathsf{r}} \text{ iff } s \models \neg (\psi'_{\mathsf{r}}) \text{ iff } s \not\models \psi'_{\mathsf{r}}$$

$$\text{iff (IH) } s' \not\models \psi' \text{ iff } s' \models \neg \psi' \text{ iff } s' \models \psi$$

Inductive case 2 $\psi = \psi' \lor \psi''$:

$$s \models \psi_{\mathsf{r}} \text{ iff } s \models (\psi' \lor \psi'')_{\mathsf{r}} \text{ iff } s \models \psi'_{\mathsf{r}} \lor \psi''_{\mathsf{r}}$$

$$\text{iff } s \models \psi'_{\mathsf{r}} \text{ or } s \models \psi''_{\mathsf{r}}$$

$$\text{iff (IH, twice) } s' \models \psi' \text{ or } s' \models \psi''$$

$$\text{iff } s' \models \psi' \lor \psi'' \text{ iff } s' \models \psi$$

Inductive case 3 $\psi=\psi'\wedge\psi''$: Very similar to inductive case 2, just with \wedge instead of \vee and "and" instead of "or".

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression Overview

STRIPS
General case
Practical issues

Proof (ctd.)

Inductive case 1 $\psi = \neg \psi'$:

$$s \models \psi_{\mathsf{r}} \text{ iff } s \models (\neg \psi')_{\mathsf{r}} \text{ iff } s \models \neg (\psi'_{\mathsf{r}}) \text{ iff } s \not\models \psi'_{\mathsf{r}}$$

$$\text{iff (IH) } s' \not\models \psi' \text{ iff } s' \models \neg \psi' \text{ iff } s' \models \psi$$

Inductive case 2 $\psi = \psi' \lor \psi''$:

$$s \models \psi_{\mathsf{r}} \text{ iff } s \models (\psi' \lor \psi'')_{\mathsf{r}} \text{ iff } s \models \psi'_{\mathsf{r}} \lor \psi''_{\mathsf{r}}$$

$$\text{iff } s \models \psi'_{\mathsf{r}} \text{ or } s \models \psi''_{\mathsf{r}}$$

$$\text{iff (IH, twice)} \ s' \models \psi' \text{ or } s' \models \psi''$$

$$\text{iff } s' \models \psi' \lor \psi'' \text{ iff } s' \models \psi$$

Inductive case 3 $\,\psi=\psi'\wedge\psi''$: Very similar to inductive case 2, just with \wedge instead of \vee and "and" instead of "or".

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Overview
Example
STRIPS
General case

Emptiness and subsumption testing

The following two tests are useful when performing regression searches to avoid exploring unpromising branches:

- Test that $regr_o(\varphi)$ does not represent the empty set (which would mean that search is in a dead end). For example, $regr_{\langle a, \neg v \rangle}(p) \equiv a \land \bot \equiv \bot$.
- Test that $regr_o(\varphi)$ does not represent a subset of φ (which would make the problem harder than before). For example, $regr_{\langle b,c\rangle}(a)\equiv a\wedge b$.

Both of these problems are NP-hard.

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Overview
Example
STRIPS
General case
Practical issues

Formula growth

The formula $\mathit{regr}_{o_1}(\mathit{regr}_{o_2}(\ldots \mathit{regr}_{o_{n-1}}(\mathit{regr}_{o_n}(\varphi))))$ may have size $O(|\varphi||o_1||o_2|\ldots |o_{n-1}||o_n|)$, i. e., the product of the sizes of φ and the operators.

 \rightsquigarrow worst-case exponential size $O(m^n)$

Logical simplifications

- $\bullet \perp \land \varphi \equiv \bot, \ \top \land \varphi \equiv \varphi, \ \bot \lor \varphi \equiv \varphi, \ \top \lor \varphi \equiv \top$
- $a \lor \varphi \equiv a \lor \varphi[\bot/a]$, $\neg a \lor \varphi \equiv \neg a \lor \varphi[\top/a]$, $a \land \varphi \equiv a \land \varphi[\top/a]$, $\neg a \land \varphi \equiv \neg a \land \varphi[\bot/a]$
- idempotency, absorption, commutativity, associativity, ...

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression

Example
STRIPS
General case
Practical issues

Restricting formula growth in search trees

Problem very big formulae obtained by regression

Cause disjunctivity in the (NNF) formulae

(formulae without disjunctions easily convertible to

is at most the number of state variables.)

small formulae $l_1 \wedge \cdots \wedge l_n$ where l_i are literals and n

Idea handle disjunctivity when generating search trees

AI Planning

B. Nebel, R. Mattmüller

Search

Progression

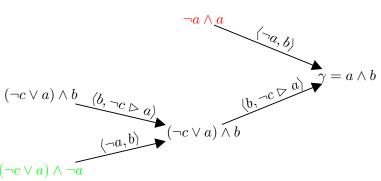
Regression

Example STRIPS General case Practical issues

Unrestricted regression: search tree example

Unrestricted regression: do not treat disjunctions specially

 $\text{Goal } \gamma = a \wedge b \text{, initial state } I = \{a \mapsto 0, b \mapsto 0, c \mapsto 0\}.$



Al Planning

B. Nebel, R. Mattmüller

Search

Progression

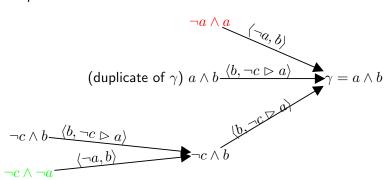
Regression Overview

STRIPS
General case
Practical issues

Full splitting: search tree example

Full splitting: always remove all disjunctivity

Goal
$$\gamma = a \wedge b$$
, initial state $I = \{a \mapsto 0, b \mapsto 0, c \mapsto 0\}$. $(\neg c \vee a) \wedge b$ in DNF: $(\neg c \wedge b) \vee (a \wedge b)$ \rightsquigarrow split into $\neg c \wedge b$ and $a \wedge b$



Al Planning

B. Nebel, R. Mattmüller

Search

Progression

D.

Overview Example STRIPS General case

Practical issues

General splitting strategies

Alternatives:

- 1 Do nothing (unrestricted regression).
- Always eliminate all disjunctivity (full splitting).
- Reduce disjunctivity if formula becomes too big.

Discussion:

- With unrestricted regression the formulae may have size that is exponential in the number of state variables.
- With full splitting search tree can be exponentially bigger than without splitting.
- The third option lies between these two extremes.

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression
Overview
Example
STRIPS
General case
Practical issues

Summary

- (Classical) search is a very important planning approach.
- Search-based planning algorithms differ along many dimensions, including
 - search direction (forward, backward)
 - what each search node represents
 (a state, a set of states, an operator sequence)
- Progression search proceeds forwards from the initial state.
 - If we use duplicate detection, each search node corresponds to a unique state.
 - If we do not use duplicate detection, each search node corresponds to a unique operator sequence.

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

egression

Summary (ctd.)

• Regression search proceeds backwards from the goal.

- Each search node corresponds to a set of states represented by a formula.
- Regression is simple for STRIPS operators.
- The theory for general regression is more complex.
- When applying regression in practice, additional considerations such as when and how to perform splitting come into play.

Al Planning

B. Nebel, R. Mattmüller

Search

Progression

Regression