

Table: Language for Tables

- Hybrid paradigm: imperative + relational
- Data types: atomic (`integer`, `string`, `boolean`) and tables (`table`)
- Program = `program` [*statements*] `end`
 - Definition of variables
 - Assignments (`id = expr`)
 - Control statements: `if-then`, `if-then-else`, `while-do`
 - Input of variables and output of expressions: `read`, `write`
- Expressions:
 - Logical (`and`, `or`, `not`): in short circuit
 - Comparison (`==`, `!=`, `>`, `>=`, `<`, `<=`)
 - Arithmetic (`+`, `-`, `*`, `/`)
 - Relational (`project`, `select`, `exists`, `all`, `join`, `update`, `extend`, `rename`)

Relational Expressions

<i>Operator</i>	<i>Example</i>
project	project [a,b] R
select	select [a==b and c != 25] R
exists	exists [a!=b or c] R
all	all [a and c] R
join	R join [a==b and c==d] S
update	update [a = b+c] R
extend	extend [integer d = b-c] R
rename	rename [x,y,z] R

Precedence, Associativity, Order of Evaluation

<i>Operator</i>	<i>Type</i>	<i>Associativity</i>
and, or	binary	left
==, !=, >, >=, <, <=,	binary	nonassoc
+, -	binary	left
*, /, join	binary	left
-, not, project, select, exists, all, update, extend, rename	unary	right

increasing



- Evaluation order of operands: from left to right

Conditional Statement

- **if** *expr* **then** *stat-list* [**else** *stat-list*] **end**

```
integer a, b, c;  
table(integer x, string y) t, r;  
...  
if a==b then  
    t = select [x>0] r  
else  
    a = b+c  
end;  
...  
if t != r then  
    write "Different tables!"  
end;  
...
```

Loop

- **while** *expr* **do** *stat-list* **end**

```
integer n, fact;  
read n;  
fact = 1;  
while n > 1 do  
    fact = fact * n;  
    n = n - 1  
end;  
write fact;
```

Input

- **read id**
- **read [*filename*] id**

```
integer i, j;  
table(integer a, string b) t, r, z;  
...  
read t;  
r = select [a == i] t;  
read ["t.dat"] t;  
z = select [a > j] t;
```

Output

- **write** *expr*
- **write** [*filename*] *expr*

```
integer i, j;  
string name;  
table(integer a, string b) t;  
...  
write select [a==i] t;  
read name;  
write [name] select [a>j] t;
```

BNF

```
program → program stat-list end
stat-list → stat ; stat-list | stat
stat → def-stat | assign-stat | if-stat | while-stat | read-stat | write-stat
def-stat → type id-list
id-list → id , id-list | id
type → atomic-type | table-type
atomic-type → integer | string | boolean
table-type → table ( attr-list )
attr-list → attr-decl , attr-list | attr-decl
attr-decl → atomic-type id
assign-stat → id = expr
expr → expr bool-op bool-term | bool-term
bool-op → and | or
bool-term → comp-term comp-op comp-term | comp-term
comp-op → == | != | > | >= | < | <=
comp-term → comp-term low-bin-op low-term | low-term
low-bin-op → + | -
low-term → low-term high-bin-op factor | factor
high-bin-op → * | / | join-op
factor → unary-op factor | ( expr ) | id | constant
unary-op → - | not | project-op | select-op | exists-op | all-op |
            update-op | extend-op | rename-op
```

```
join-op → join [ expr ]
project-op → project [ id-list ]
select-op → select [ expr ]
exists-op → exists [ expr ]
all-op → all [ expr ]
extend-op → extend [ atomic-type id = expr ]
update-op → update [ id = expr ]
rename-op → rename [ id-list ]
constant → atomic-const | table-const
atomic-const → intconst | strconst | boolconst
table-const → { table-instance }
table-instance → tuple-list | atomic-type-list
tuple-list → tuple-const , tuple-list | tuple-const
tuple-const → ( atomic-const-list )
atomic-const-list → atomic-const , atomic-const-list |
                    atomic-const
atomic-type-list → atomic-type , atomic-type-list |
                    atomic-type
if-stat → if expr then stat-list else-part end
else-part → else stat-list | ε
while-stat → while expr do stat-list end
read-stat → read specifier id
specifier → [ expr ] | ε
write-stat → write specifier expr
```




EBNF

```
program → program stat-list end
stat-list → stat { ; stat }
stat → def-stat | assign-stat | if-stat | while-stat | read-stat | write-stat
def-stat → type id-list
id-list → id { , id }
type → atomic-type | table-type
atomic-type → integer | string | boolean
table-type → table ( attr-list )
attr-list → attr-decl { , attr-decl }
attr-decl → atomic-type id
assign-stat → id = expr
expr → bool-term { (and | or) bool-term }
bool-term → comp-term [ (== | != | > | >= | < | <=) comp-term ]
comp-term → low-term { (+ | - ) low-term }
low-term → factor { (* | / | join-op) factor }
factor → unary-op factor | ( expr ) | id | constant
unary-op → - | not | project-op | select-op | exists-op | all-op |
           update-op | extend-op | rename-op
```

```
join-op → join [ expr ]
project-op → project [ id-list ]
select-op → select [ expr ]
exists-op → exists [ expr ]
all-op → all [ expr ]
extend-op → extend [ atomic-type id = expr ]
update-op → update [ id = expr ]
rename-op → rename [ id-list ]
constant → atomic-const | table-const
atomic-const → intconst | strconst | boolconst
table-const → { tuple-list | atomic-type-list }
tuple-list → tuple-const { , tuple-const }
atomic-type-list → atomic-type { , atomic-type }
tuple-const → ( atomic-const { , atomic-const } )
if-stat → if expr then stat-list [ else stat-list ] end
while-stat → while expr do stat-list end
read-stat → read specifier id
specifier → [ expr ] | ε
write-stat → write specifier expr
```

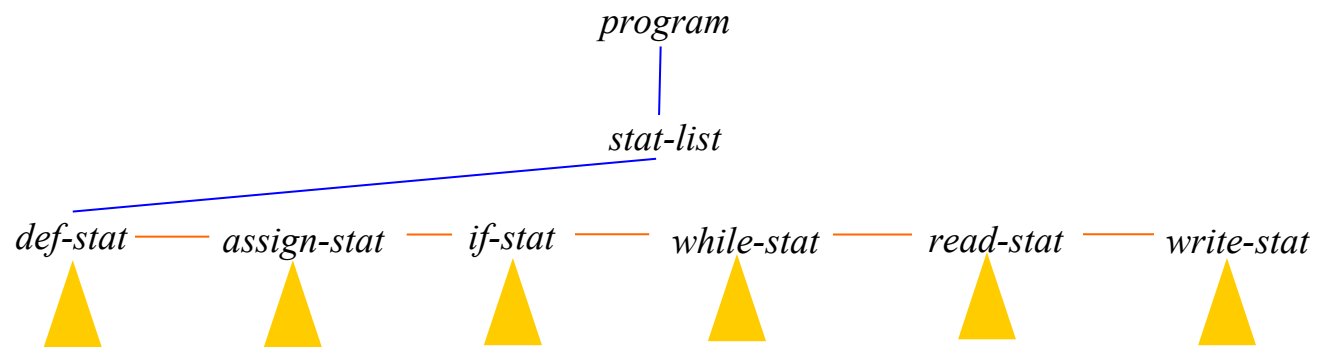
Abstract Tree

- Node structure:

type	value	line	child	brother
				

 (value: ival, sval)
- $\text{type} \in \{ \text{N_ASSIGN_STAT}, \text{N_ATOMIC_TYPE}, \text{N_ATTR_DECL}, \text{N_BOOLCONST}, \text{N_COMP_EXPR}, \text{N_DEF_STAT}, \text{N_EXTEND_EXPR}, \text{N_ID}, \text{N_IF_STAT}, \text{N_INTCONST}, \text{N_JOIN_EXPR}, \text{N_LOGIC_EXPR}, \text{N_MATH_EXPR}, \text{N_NEG_EXPR}, \text{N_PROGRAM}, \text{N_PROJECT_EXPR}, \text{N_READ_STAT}, \text{N_RENAME_EXPR}, \text{N_SELECT_EXPR}, \text{N_SPECIFIER}, \text{N_STAT_LIST}, \text{N_STRCONST}, \text{N_TABLE_CONST}, \text{N_TABLE_TYPE}, \text{N_TUPLE_CONST}, \text{N_TYPE}, \text{N_UPDATE_EXPR}, \text{N_WHILE_STAT}, \text{N_WRITE_STAT} \}$
- $\text{type} = \text{N_ATOMIC_TYPE} \rightarrow \text{value.ival} \in \{ \text{INTEGER}, \text{STRING}, \text{BOOLEAN} \}$
- $\text{type} = \text{N_COMP_EXPR} \rightarrow \text{value.ival} \in \{ \text{EQ}, \text{NE}, '>', \text{GE}, '<', \text{LE} \}$
- $\text{type} = \text{N_LOGIC_EXPR} \rightarrow \text{value.ival} \in \{ \text{AND}, \text{OR} \}$
- $\text{type} = \text{N_MATH_EXPR} \rightarrow \text{value.ival} \in \{ '+', '-', '*', '/' \}$
- $\text{type} = \text{N_NEG_EXPR} \rightarrow \text{value.ival} \in \{ '-', \text{NOT} \}$
- $\text{type} = \text{N_SELECT_EXPR} \rightarrow \text{value.ival} \in \{ \text{SELECT}, \text{EXISTS}, \text{ALL} \}$

Abstract Tree (ii)



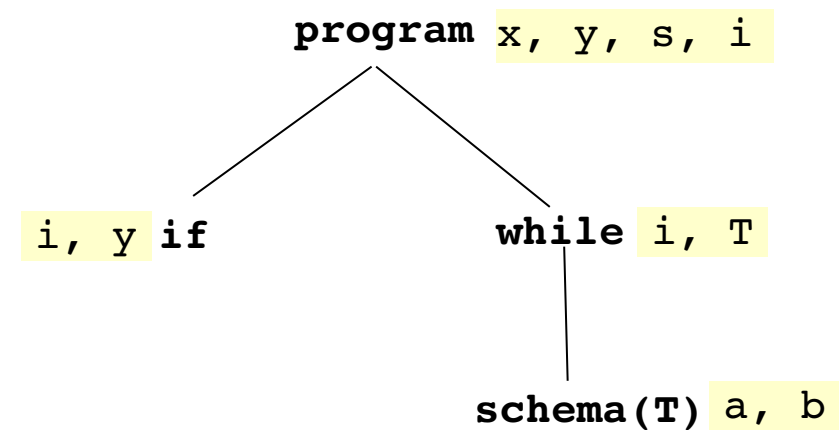
Type Checking

- Traversing of the tree starting from the root
- **Checks:**
 1. Uniqueness of names within environments (attributes in table, variables in blocks)
 2. Visibility within the **environment hierarchy** of a referenced identifier:
$$\text{environment hierarchy} = \text{context hierarchy} + \text{block hierarchy}$$
 3. Compatibility of an operator with its operands
 4. Compatibility of the identifier with the assignment expression
 5. Compatibility of expressions with the relevant statements
- **Type inference:** computation of result schema in each operation

Environment Hierarchy

- Problem: binding of a name with the relevant definition
- Search within the environment hierarchy (static scope)

```
program
  real x, y;
  string s;
  integer i;
  ...
  if(x == y) then
    integer i;
    real y;
    ...
  end;
  ...
  while i > j do
    integer i;
    table(integer a, string b) T;
    ...
    write select [ a > i and b == s ] T
  end
end
```

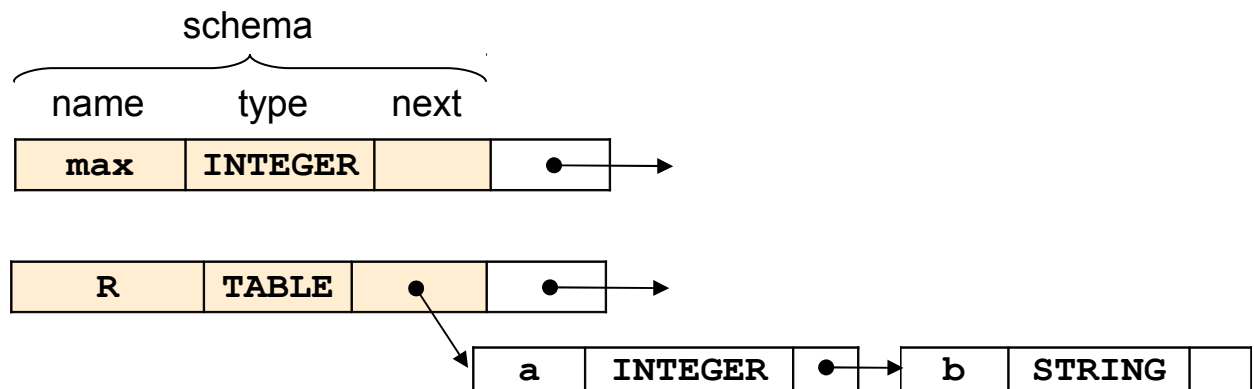
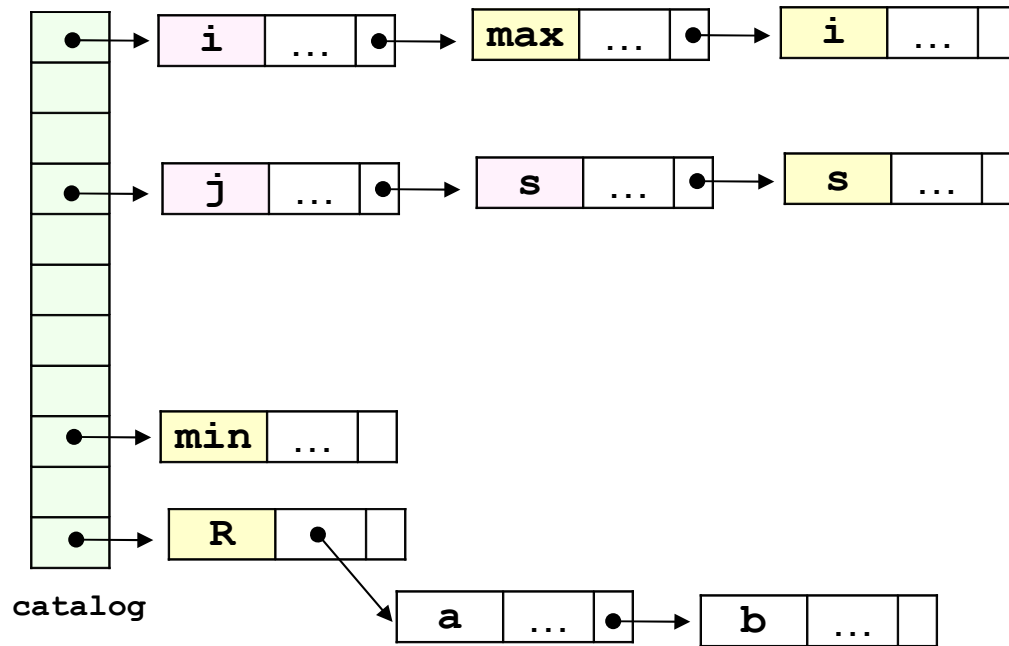


Symbol Table

```

program
  integer min, max, i;
  string s;
  table(integer a, string b) R;
  ...
  if min == max then
    table(string c, boolean d) T;
    read T;
    write project [a, d]
      R join [ b == c ] T;
  else
    while(min != max) do
      integer i, j, s;
      ...
      min = i + j - max;
      ...
      boolean ok;
      ...
    end
  end;
  ...
end

```



Context Hierarchy

- Contextualizing operators: `select`, `exists`, `all`, `join`, `update`, `extend`

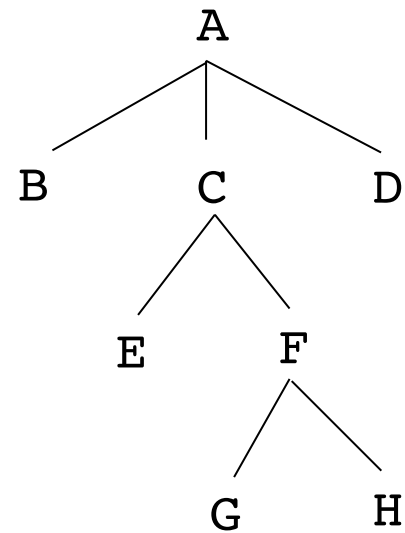
```
table(int a, string x) A;  
table(int b, bool y) B, C;  
...  
C = select [ exists [ a > b ] A ] B
```

A		B		C	
a	x	b	y	b	y
3	alpha	13	alpha	4	beta
8	beta	4	beta	6	delta
5	gamma	15	gamma	3	zeta
2	delta	6	delta		
6	epsilon	24	epsilon		
		3	zeta		

Context Hierarchy (ii)

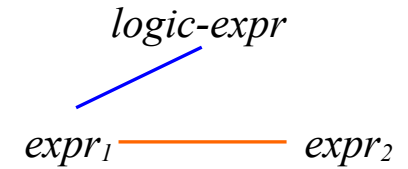
- Search for the name by ascending the context hierarchy

```
[  
  [  
    [  
      [  
        [  
          [  
            [  
              [  
                [  
                  [  
                    [  
                      [  
                        [  
                          [  
                        ] G  
                      ] H  
                    ] F  
                  ] C  
                ] D  
              ] A  
            ]  
          ]  
        ]  
      ]  
    ]  
  ]  
]
```



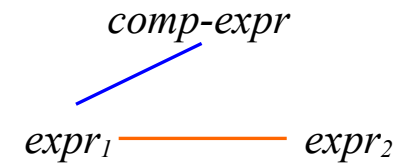
- Type checking: keeps a **context stack** for searching identifiers

Logical Expressions (*logic-expr*)



- Qualifier: **AND, OR**
- Constraint: `type(expr1) = boolean, type(expr2) = boolean`
- Result schema: `type(logic-expr) = boolean`

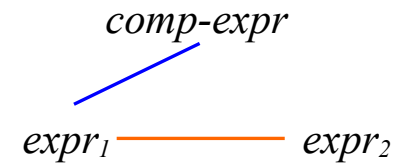
Comparison Expressions (*comp-expr*)



- Qualifier: **EQ, NE**
- Constraint: $\text{type}(\text{expr}_1) = \text{type}(\text{expr}_2)$
- Result schema: $\text{type}(\text{comp-expr}) = \mathbf{boolean}$

Comparison Expressions (*comp-expr*) (ii)

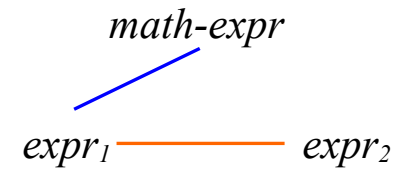
- Qualifier: ' $>$ ', **GE**, ' $<$ ', **LE**



- Constraint: $(\text{type}(\text{expr}_1) = \text{int}, \text{type}(\text{expr}_2) = \text{int})$ or $(\text{type}(\text{expr}_1) = \text{string}, \text{type}(\text{expr}_2) = \text{string})$
- Result schema: $\text{type}(\text{comp-expr}) = \text{boolean}$

Arithmetic Expressions (*math-expr*)

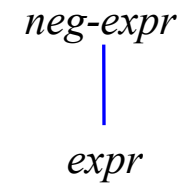
- Qualifier: `'+', '-', '*', '/'`



- Constraint: `type(expr1) = integer, type(expr2) = integer`
- Result schema: `type(math-expr) = integer`

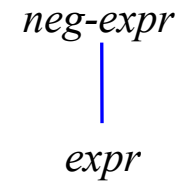
Negation (*neg-expr*)

- Qualifier: ' - '
- Constraint: `type(expr) = integer`
- Result schema: `type(neg-expr) = integer`

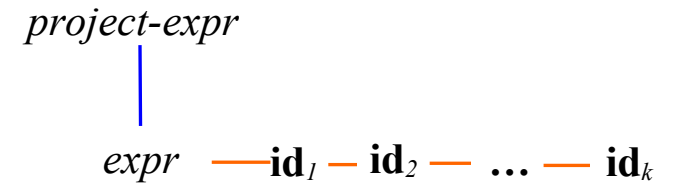


Negation (*neg-expr*) (ii)

- Qualifier: **NOT**
- Constraint: $\text{type}(\text{expr}) = \text{boolean}$
- Result schema: $\text{type}(\text{neg-expr}) = \text{boolean}$



Projection (*project-expr*)



- Constraint: $\text{type}(\text{expr}) = \mathbf{table}$,
 $\text{names}(\text{id-list}) \subset \text{names}(\text{attributes}(\text{type}(\text{expr})))$, $\text{nodup}(\text{names}(\text{id-list}))$
- Result schema: $\text{type}(\text{project-expr}) =$
 $\mathbf{table}(a \mid a \in \text{attributes}(\text{type}(\text{expr})), \text{name}(a) \in \text{names}(\text{id-list}))$

Rename (*rename-expr*)

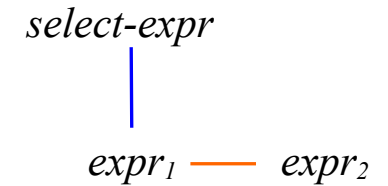
rename-expr

expr — **id**₁ — **id**₂ — ... — **id**_n

- Constraint: $\text{type}(\text{expr}) = \mathbf{table}$,
 $|\text{names}(\text{id-list})| = |\text{attributes}(\text{type}(\text{expr}))| = n$,
 $\text{nodup}(\text{id-list})$
- Result schema: $\text{type}(\text{rename-expr}) = \mathbf{table} (\text{id-list}[i]: \text{type}(\text{expr})[i] \mid i \in [1..n])$

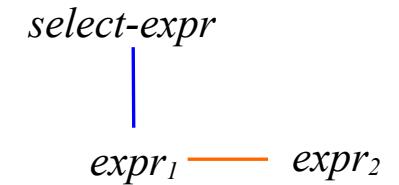
Selection (*select-expr*)

- Qualifier: **SELECT**
- Constraint: $\text{type}(\text{expr}_2) = \mathbf{table}$, $\text{type}(\text{expr}_1) = \mathbf{boolean}$
- Result schema: $\text{type}(\text{select-expr}) = \mathbf{type}(\text{expr}_2)$

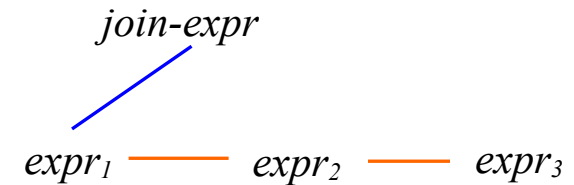


Selection (*select-expr*) (ii)

- Qualifier: **EXISTS**, **ALL**
- Constraint: $\text{type}(\text{expr}_2) = \mathbf{table}$, $\text{type}(\text{expr}_1) = \mathbf{boolean}$
- Result schema: $\text{type}(\text{select-expr}) = \mathbf{boolean}$



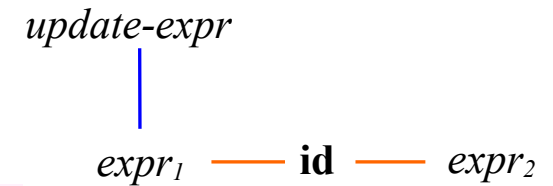
Join (*join-expr*)



- Constraint: $\text{type}(\text{expr}_1) = \mathbf{table}$, $\text{type}(\text{expr}_3) = \mathbf{table}$, $\text{type}(\text{expr}_2) = \mathbf{boolean}$,
 $\text{names}(\text{attributes}(\text{type}(\text{expr}_1))) \cap \text{names}(\text{attributes}(\text{type}(\text{expr}_3))) = \emptyset$
- Result schema: $\text{type}(\text{join-expr}) = \text{type}(\text{expr}_1) \cup \text{type}(\text{expr}_3)$

Update (*update-expr*)

- Constraint: $\text{type}(\text{expr}_1) = \mathbf{table}$,
 $(\text{name}(\mathbf{id}): \text{domain}) \in \text{attributes}(\text{type}(\text{expr}_1))$,
 $\text{type}(\text{expr}_2) = \text{domain}$



- Result schema: $\text{type}(\text{update-expr}) = \text{type}(\text{expr}_1)$

Extension (*extend-expr*)

extend-expr

$expr_1 \text{ --- } atomic\text{-}type \text{ --- } id \text{ --- } expr_2$

- Constraint: $type(expr_1) = \mathbf{table}$,
 $(name(id)) \notin names(attributes(type(expr_1)))$,
 $type(expr_2) = type(atomic\text{-}type)$
- Result schema: $type(extend\text{-}expr) = type(expr_1) \cup (name(id): type(atomic\text{-}type))$

Definition of Variables (*def-stat*)

def-stat
|
type — **id**₁ — **id**₂ — ... — **id**_k

- Constraint: $\forall i \in [1 .. k], \forall j \in [1 .. k], i \neq j, (\text{name}(\mathbf{id}_i) \neq \text{name}(\mathbf{id}_j)),$
 $\forall i \in [1 .. k] (\text{name}(\mathbf{id}_i) \notin \text{Environment})$

Assignment (*assign-stat*)

- Constraint: `visible(name(id)),`
`type(id) = type(expr)`



Conditional Statement (*if-stat*)

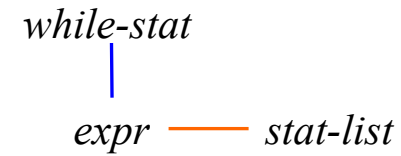
- Constraint: `type(expr) = boolean`

if-stat
|
expr — *stat-list*

if-stat
|
expr — *stat-list* — *stat-list*

While Statement (*while-stat*)

- Constraint: `type(expr) = boolean`



Reading a Variable (*read-stat*)

read-stat
|
specifier — **id**

- Constraint: `visible(name(id)),`
`(type(specifier) = nil or type(specifier) = string)`

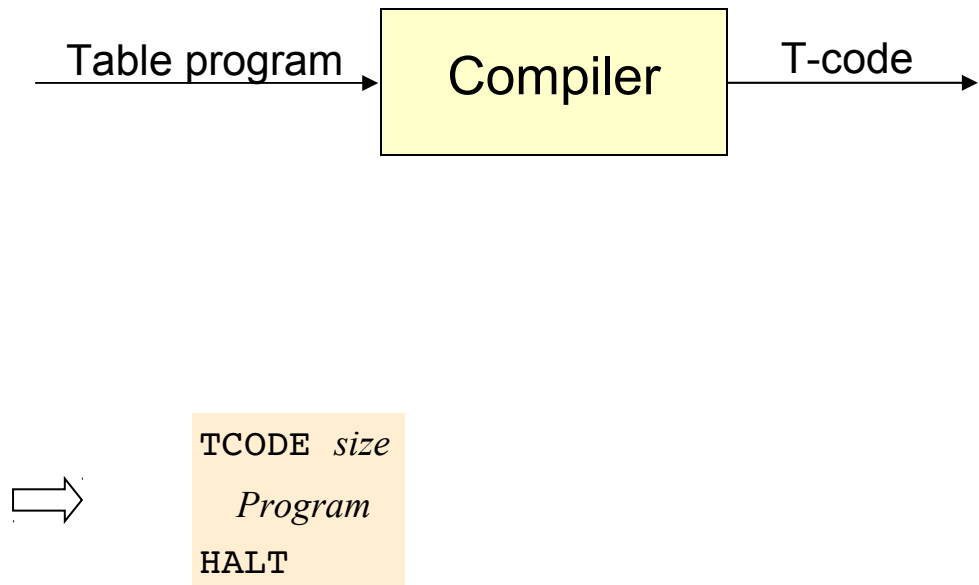
Writing a Value (*write-stat*)

write-stat
|
specifier — *expr*

- Constraint: $(\text{type}(\textit{specifier}) = \text{nil} \text{ or } \text{type}(\textit{specifier}) = \text{string})$

Intermediate Code Generation

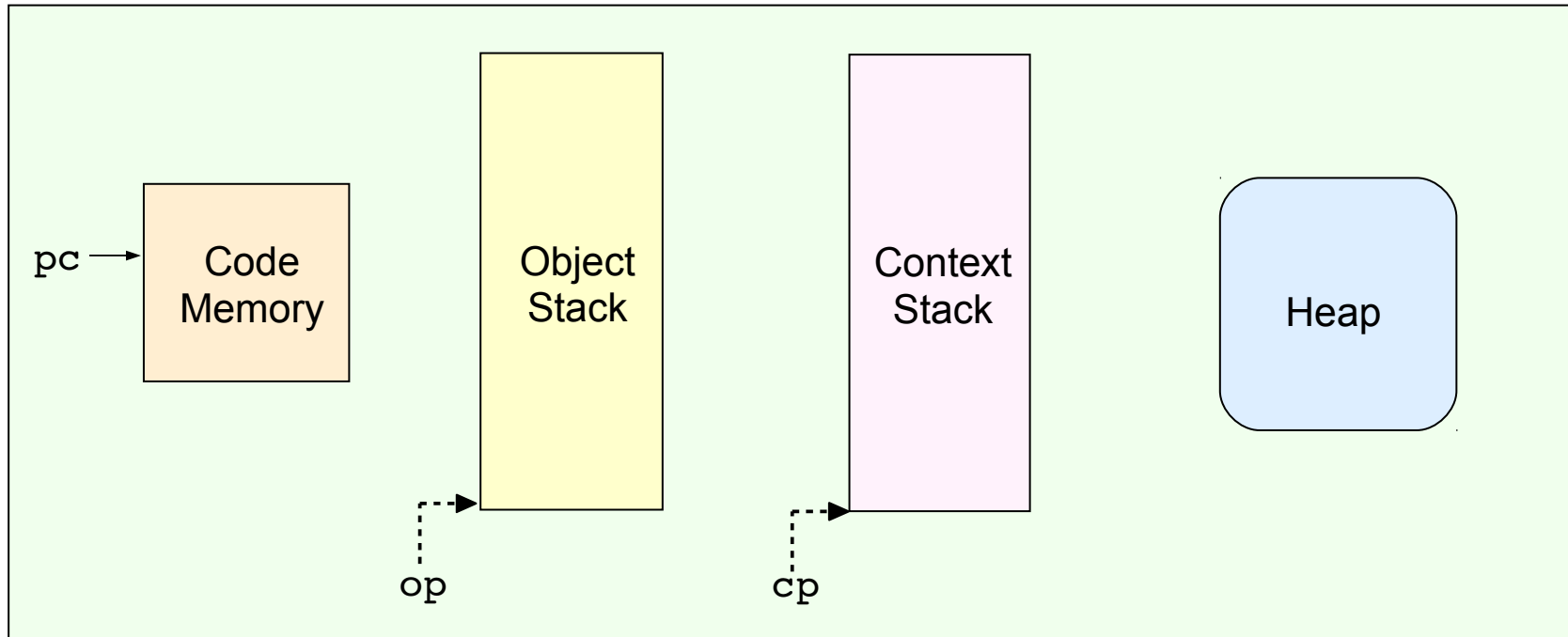
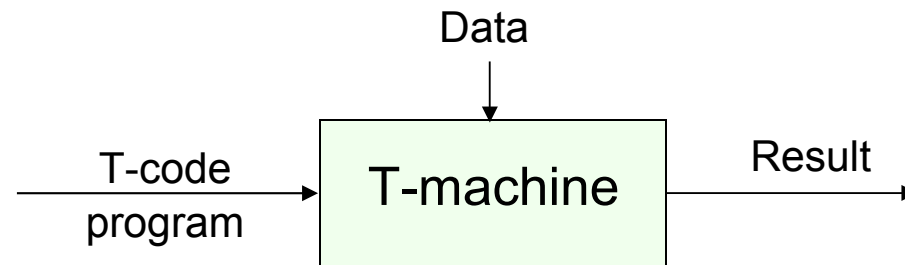
```
program
  integer min, max, i;
  string s;
  table(integer a, string b) R;
  ...
  if min == max then
    table(string c, boolean d) T;
    read T;
    write project [a, d]
      R join [ b == c ] T;
  else
    while(min != max) do
      integer i, j, s;
      ...
      min = i + j - max;
      ...
      boolean ok;
      ...
    end
  end;
  ...
end
```



Design Choices

- Code directly addressable (without explicit labels)
- Address of a T-code stat = position of the stat within the generated code
- Variable descriptors allocated in the order in which variables are defined
- Variable identification: *oid* = position of the variable within the stack (0 .. *n*)
- Attribute identification: *context-offset*, *position-within-schema*, *attribute-size*
- Management of blocks by stack: exit from block → pop variables in block

Abstract Machine



Definition of Variables

- Two types of objects $\left\{ \begin{array}{l} \text{Atoms} \\ \text{Tables} \end{array} \right.$

```
integer min, max, i;  
string s;  
boolean b;  
table(integer a, string b) R;
```



```
NEWATOM |integer|  
NEWATOM |integer|  
NEWATOM |integer|  
NEWATOM |string|  
NEWATOM |integer|  
NEWTAB <|integer|+|string|>
```

- Notes:**

- atom (embedded instance): `NEWATOM` *object-size*
- table (instance in heap): `NEWTAB` *tuple-size*

Reference to Atomic Constants

- Integer constant `y = 25;` \Rightarrow `LDINT 25`
- String constant `s = "alpha";` \Rightarrow `LDSTR "alpha"`
- Boolean constant `b = true;` \Rightarrow `LDINT 1`

- **Note:**

- Boolean values: **true**, **false** \rightarrow surrogated by integers: 1, 0

Reference to Table Constants

```
table (integer i, string s, boolean b) t;  
t = {(1, "alpha", true),  
      (2, "beta", false),  
      (3, "gamma", true)};
```



```
LDTAB <|integer|+|string|+|integer|> 3  
  IATTR 1  
  SATTR "alpha"  
  IATTR 1  
  IATTR 2  
  SATTR "beta"  
  IATTR 0  
  IATTR 3  
  SATTR "gamma"  
  IATTR 1  
ENDTAB
```

```
table (integer i, string s, boolean b) t;  
t = {integer, string, boolean};
```



```
LDTAB <|integer|+|string|+|integer|> 0  
ENDTAB
```

• Notes:

- *tuple-size* = size of tuples
- *num-tuples* = number of tuples
- *<attribute loadings>* = sequence of attribute-constant loadings

```
LDTAB tuple-size num-tuples  
      <attribute loadings>  
ENDTAB
```

Reference to Identifiers

- Variable

```
x = y + 1;
```



```
LOB ^y
```

```
t1 = t2;
```



```
LOB ^t2
```

```
integer x, y;
table(integer a, boolean b) t1, t2;
table(integer c, string d) t3;
```

- Attribute within current context

```
write
select [d == "alpha"] t3;
```



```
LAT 0 &d |string|
```

- Attribute within external context

```
write
select [ exists [a > c] t1 ] t3;
```



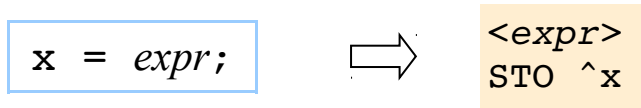
```
LAT 1 &c |integer|
```

- Notes:

- Argument of LOB = *oid* (object identifier)
- Arguments of LAT = *context-offset*, *attribute-offset*, *attribute-size*
- For reference to attributes within predicate of join, *attribute-offset* in LAT refers to result schema (concatenation of operand schemas)

Assignment

- Computation of assignment expression + assignment statement (STO)



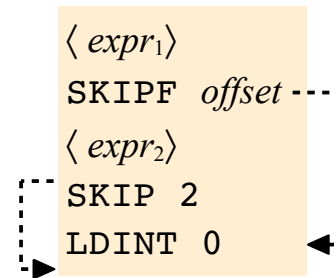
- **Note:**
 - Arguments of `STO` = *oid* (of assigned object)

Logical Operations (and, or)

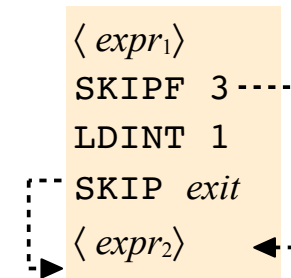
$logic\text{-}expr \rightarrow expr_1 \ expr_2$



and

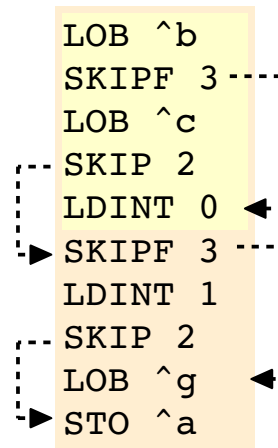


or



```

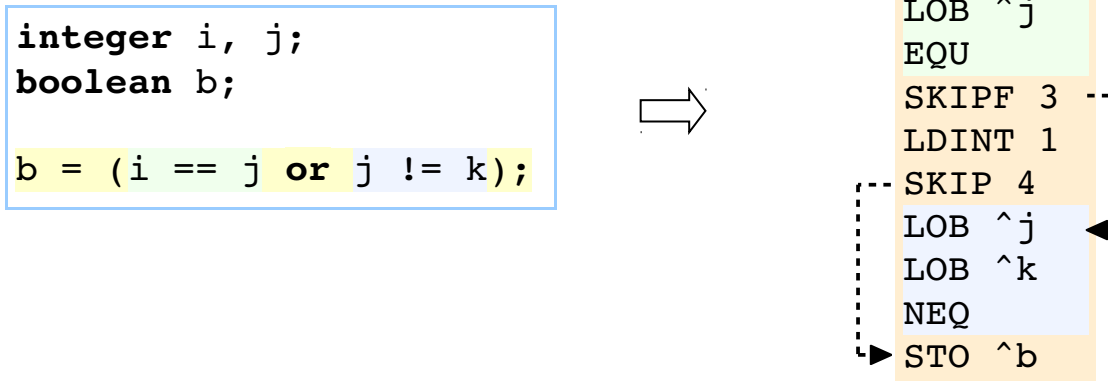
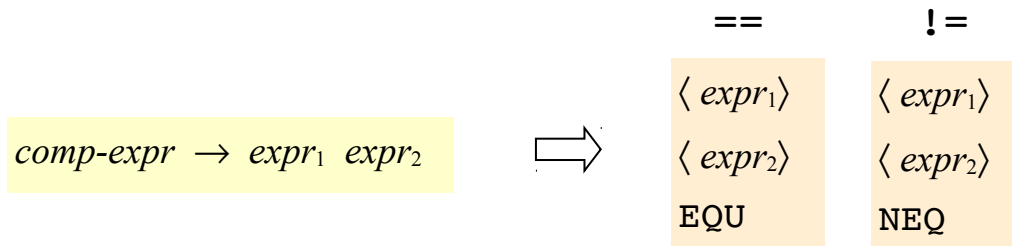
boolean a, b, c;
...
a = (b and c) or g;
    
```



• Notes:

- Short circuit evaluation
- SKIP = unconditional jump
- SKIPF = conditional jump
- Argument of SKIP, SKIPF = length of the jump (offset) $\begin{cases} exit = |\langle expr_2 \rangle| + 1 \\ offset = |\langle expr_2 \rangle| + 2 \end{cases}$

Comparison Operations: ==, !=



• Note:

- EQU, NEQ: polymorphic for all types

Comparison Operations: >, >=, <, <=

$comp\text{-}expr \rightarrow expr_1\ expr_2$



integer

>	>=	<	<=
$\langle expr_1 \rangle$	$\langle expr_1 \rangle$	$\langle expr_1 \rangle$	$\langle expr_1 \rangle$
$\langle expr_2 \rangle$	$\langle expr_2 \rangle$	$\langle expr_2 \rangle$	$\langle expr_2 \rangle$
IGT	IGE	ILT	ILE

string

>	>=	<	<=
$\langle expr_1 \rangle$	$\langle expr_1 \rangle$	$\langle expr_1 \rangle$	$\langle expr_1 \rangle$
$\langle expr_2 \rangle$	$\langle expr_2 \rangle$	$\langle expr_2 \rangle$	$\langle expr_2 \rangle$
SGT	SGE	SLT	SLE

```
integer i, j;
string x, y;
boolean b;

b = (i > j or x < y);
```



```
LOB ^i
LOB ^j
IGT
SKIPF 3
LDINT 1
SKIP 4
LOB ^x
LOB ^y
SLT
STO ^b
```

Arithmetic Operations: +, −, *, /,

math-set-expr → *expr*₁ *expr*₂



+	−	*	/
⟨ <i>expr</i> ₁ ⟩	⟨ <i>expr</i> ₁ ⟩	⟨ <i>expr</i> ₁ ⟩	⟨ <i>expr</i> ₁ ⟩
⟨ <i>expr</i> ₂ ⟩	⟨ <i>expr</i> ₂ ⟩	⟨ <i>expr</i> ₂ ⟩	⟨ <i>expr</i> ₂ ⟩
PLUS	MINUS	TIMES	DIV

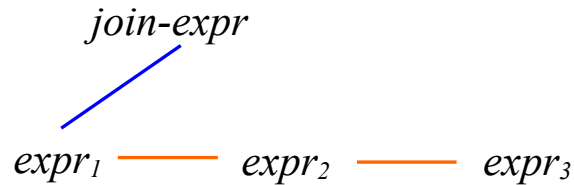
```
integer i, j, k;
```

```
i = (i + 5) * (j − k);
```



```
LOB ^i
LDINT 5
PLUS
LOB ^j
LOB ^k
MINUS
TIMES
STO ^i
```

Join



```

< expr1 >
< expr3 >
JOIN gap
  < expr2 >
ENDJOIN gap
  
```

```

table(integer a, boolean b, integer c) r;
table(integer d, integer e) s;

w = r join [ c == a+d-5 and b ] s;
  
```



```

LOB ^r
LOB ^s
JOIN 11
  LAT 0 &c |integer|
  LAT 0 &a |integer|
  LAT 0 &d |integer|
  IPLUS
  LDINT 5
  IMINUS
  EQU
  SKIPF 3
  LAT 0 &b |integer|
  SKIP 2
  LDINT 0
ENDJOIN 11
STO ^w
  
```

• Notes:

- $expr_2$ = predicate of join
- $gap = |\langle expr_2 \rangle|$
- Argument *attribute-offset* of LAT: computed on result (also for external contexts)
- Argument of ENDJOIN: = implicit jump (backward) to predicate of JOIN

Change of Sign

neg-expr
|
expr



< expr >
UMI

```
integer i, j;  
j = -i + 10;
```



```
LOB ^i  
UMI  
LDINT 10  
PLUS  
STO ^j
```

Logical Negation

neg-expr

|

expr



< expr >

NEG

```
integer i, j, k;  
boolean a, b;  
b = i > j * k and not (a or j);
```



```
LOB ^i  
LOB ^j  
LOB ^k  
TIMES  
IGT
```

```
SKIPF 8
```

```
LOB ^a
```

```
SKIPF 3
```

```
LDINT 1
```

```
SKIP 2
```

```
LOB ^j
```

```
NEG
```

```
SKIP 2
```

```
LDINT 0
```

```
STO ^b
```

Projection

project-expr

expr — **id₁** — **id₂** — ... — **id_k**



$\langle expr \rangle$

PROJ *k*

ATTR *attr-offset₁* *size₁*

ATTR *attr-offset₂* *size₂*

...

ATTR *attr-offset_k* *size_n*

ENDPROJ

REMDUP

```
table(integer a, boolean b, string c) t;
```

```
w := project [ a, c ] t;
```



LOB ^t

PROJ 2

ATTR &a |integer|

ATTR &c |string|

ENDPROJ

REMDUP

STO ^w

• Notes:

- Argument of PROJ (*k*) = | *id-list* |
- Each projection attribute is specified by an ATTR
- Need for duplicate removal by REMDUP

Rename



```
table(integer a, boolean b, string c) t;
w := rename [ x, y, z ] t;
```

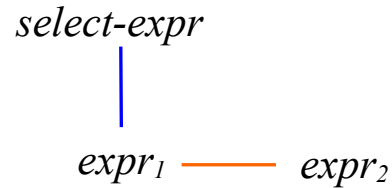


LOB	^t
STO	^w

• Note:

- Rename is relevant to semantic analysis only
- Code of rename coincides with code of its operand

Selection Operations: SELECT, EXISTS, ALL

select-expr

expr₁ — *expr₂*



SELECT

```
< expr2 >
SEL gap
  < expr1 >
ENDSEL gap
```

EXISTS

```
< expr2 >
EXS gap
  < expr1 >
ENDEXS gap
```

ALL

```
< expr2 >
ALL gap
  < expr1 >
ENDALL gap
```

```
table(integer a, boolean b) r;
table(integer c, string d) s;

w = select [ b and exists [ c == a ] s ] r;
```



```
LOB ^r
SEL 10
  LAT 0 &b |integer|
  SKIPF 8
  LOB ^s
  EXS 3
    LAT 0 &c |integer|
    LAT 1 &a |integer|
    EQU
  ENDEXS 4
  SKIP 2
  LDINT 0
  ENDSSEL 10
  STO ^w
```

• Notes:

- *expr₁* = selection predicate
- *gap* = |< *expr₁*>|

Update

update-expr

*expr*₁ — **id** — *expr*₂



```
< expr1 >  
UPD attr-offset size gap  
  < expr2 >  
ENDUPD gap  
REMDUP
```

```
integer i;  
table(integer a, string b) t;  
  
w = update [ a = i + 8 ] t;
```



```
LOB ^t  
UPD &a |integer| 3  
  LOB ^i  
  LDINT 8  
  PLUS  
ENDUPD 3  
REMDUP  
STO ^w
```

• Notes:

- *expr*₂ = update expression
- *gap* = |< *expr*₂ >|
- need for duplicate removal (REMDUP)

Extend

extend-expr

*expr*₁ — *atomic-type* — **id** — *expr*₂



```

< expr1 >
EXT size gap
    < expr2 >
ENDEXT gap
    
```

```

integer i;
table(integer a, string b) t;

w = extend [ integer c = a - i ] t;
    
```



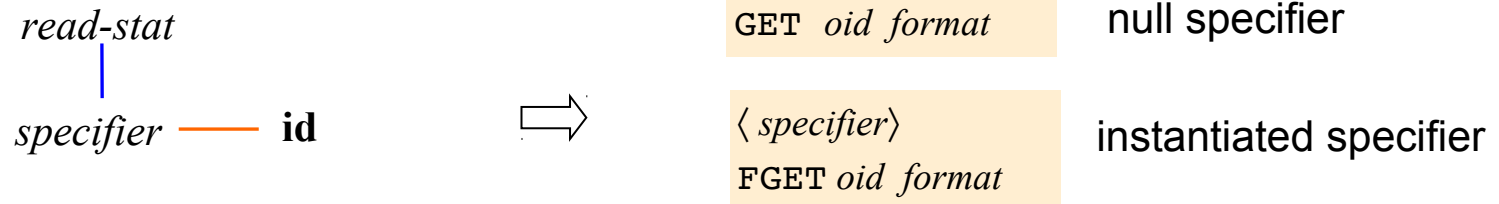
```

LOB ^t
EXT |integer| 3
    LAT 0 &a |integer|
    LOB ^i
    MINUS
ENDEXT 3
STO ^w
    
```

• Notes:

- *expr*₂ = extension expression
- *size* = size of the new attribute
- *gap* = |< *expr*₂>|

Read



```
boolean k;  
string s;  
table(integer a, string b) t;  
  
read k;  
read t;  
read [ s ] t;
```



```
GET ^k "b"  
GET ^t "(a:i,b:s)"  
LOB ^s  
FGET ^t "(a:i,b:s)"
```

• Notes:

- *oid* = object to be instantiated
- *format* = string specifying the schema of the variable to instantiate

Write

write-stat

specifier — *expr*



<expr>
PRINT *format*

null specifier

<expr>
<specifier>
FPRINT *format*

instantiated specifier

```
integer i;
string s;
table(integer a, string b) t;
```

```
write i + 25;
write [ s ] select [ a >= i ] t;
write {(1, "alpha"), (2, "beta")};
```



```
LOB ^i
LDINT 25
IPLUS
PRINT "i"
```

```
LOB ^t
SEL 3
  LAT 0 &a |integer|
  LOB ^i
  IGE
ENDSEL 3
LOB ^s
FPRINT "(a:i,b:s)"
```

```
LDTAB <|integer|+|string|> 2
  IATTR 1
  SATTR "alpha"
  IATTR 2
  SATTR "beta"
ENDTAB
PRINT "(?:i,?:s)"
```

• Notes:

- *expr* = object to print
- *specifier* (if instantiated) = file name where instance is printed
- *format* = string specifying the schema of the object to instantiate
- If attribute without name → ?

Conditional Statement: if-then

if-stat
|
expr — *stat-list*



<expr>
SKIPF *offset*
<stat-list>

```
integer i, j, k;  
table(integer a, string b) t, r;  
  
if i == j then  
    t = select [ a > 0 ] r  
end;
```



```
LOB ^i  
LOB ^j  
EQU  
SKIPF 8  
LOB ^r  
SEL 3  
    LAT 0 &a |integer|  
    LDINT 0  
    IGT  
ENDSEL 3  
STO ^t
```

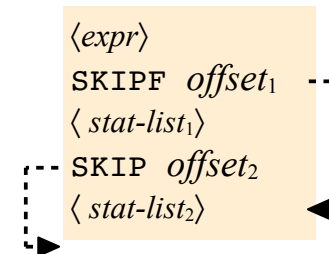
- **Note:**

- $offset = |\langle stat-list \rangle| + 1$

Conditional Statement: if-then-else

if-stat

expr — *stat-list*₁ — *stat-list*₂

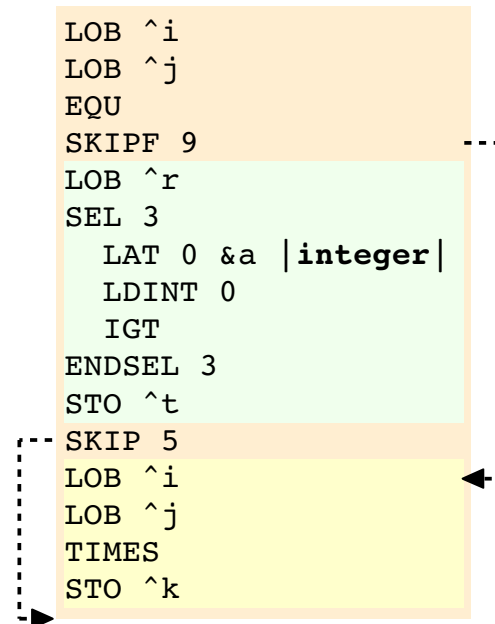


```

integer i, j, k;
table(integer a, string b) t, r;

if i == j then
    t = select [ a > 0 ] r
else
    k = i * j
end;

```



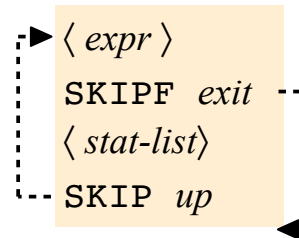
• Notes:

- $offset_1 = |\langle stat-list_1 \rangle| + 2$
- $offset_2 = |\langle stat-list_2 \rangle| + 1$

Loop

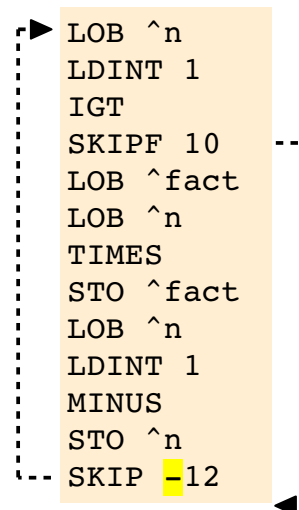
while

expr — *stat-list*



```

integer n, fact;
read n;
fact = 1;
while n > 1 do
    fact = fact * n;
    n = n - 1
end;
write fact;
  
```

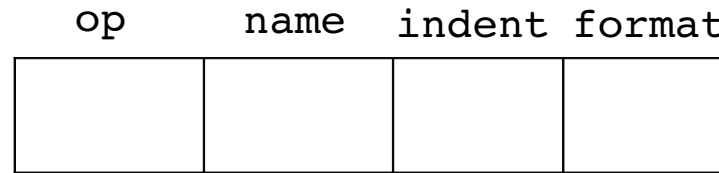


• Notes:

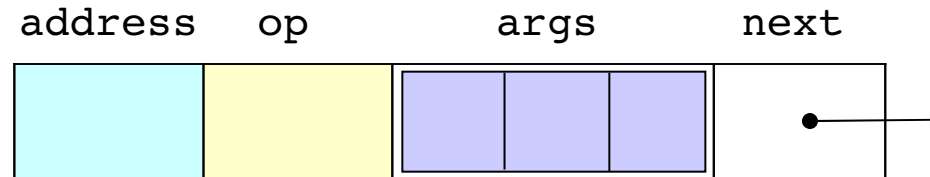
- $exit = |\langle stat-list \rangle| + 2$
- $up = -(|\langle expr \rangle| + |\langle stat-list \rangle| + 1)$

Data Structures for Code Generation

- Opdescr:



- Tstat:



- Code:



- Representation of a code segment (list of T-code statements):

