4. Convex optimization problems

- optimization problem in standard form
- convex optimization problems
- quasiconvex optimization
- linear optimization
- quadratic optimization
- geometric programming
- generalized inequality constraints
- semidefinite programming
- vector optimization

Optimization problem in standard form

minimize
$$f_0(x)$$

subject to $f_i(x) \leq 0, \quad i = 1, \dots, m$
 $h_i(x) = 0, \quad i = 1, \dots, p$

- $x \in \mathbb{R}^n$ is the optimization variable
- $f_0: \mathbf{R}^n \to \mathbf{R}$ is the objective or cost function
- $f_i: \mathbf{R}^n \to \mathbf{R}, \ i=1,\ldots,m$, are the inequality constraint functions
- $h_i: \mathbf{R}^n \to \mathbf{R}$ are the equality constraint functions

optimal value:

$$p^* = \inf\{f_0(x) \mid f_i(x) \le 0, \ i = 1, \dots, m, \ h_i(x) = 0, \ i = 1, \dots, p\}$$

- $p^* = \infty$ if problem is infeasible (no x satisfies the constraints)
- $p^* = -\infty$ if problem is unbounded below

Optimal and locally optimal points

x is **feasible** if $x \in \operatorname{dom} f_0$ and it satisfies the constraints a feasible x is **optimal** if $f_0(x) = p^\star$; X_{opt} is the set of optimal points x is **locally optimal** if there is an R > 0 such that x is optimal for

minimize (over
$$z$$
) $f_0(z)$ subject to
$$f_i(z) \leq 0, \quad i=1,\ldots,m, \quad h_i(z)=0, \quad i=1,\ldots,p$$
 $\|z-x\|_2 \leq R$

examples (with n = 1, m = p = 0)

- $f_0(x) = 1/x$, $\operatorname{dom} f_0 = \mathbf{R}_{++}$: $p^* = 0$, no optimal point
- $f_0(x) = -\log x$, $\operatorname{dom} f_0 = \mathbf{R}_{++}$: $p^* = -\infty$
- $f_0(x) = x \log x$, $\operatorname{dom} f_0 = \mathbf{R}_{++}$: $p^* = -1/e$, x = 1/e is optimal
- $f_0(x) = x^3 3x$, $p^* = -\infty$, local optimum at x = 1

Local and global optima

any locally optimal point of a convex problem is (globally) optimal **proof**: suppose x is locally optimal and y is optimal with $f_0(y) < f_0(x)$ x locally optimal means there is an R>0 such that

z feasible,
$$||z-x||_2 \leq R \implies f_0(z) \geq f_0(x)$$

consider $z = \theta y + (1 - \theta)x$ with $\theta = R/(2||y - x||_2)$

- $||y x||_2 > R$, so $0 < \theta < 1/2$
- z is a convex combination of two feasible points, hence also feasible
- $||z x||_2 = R/2$ and

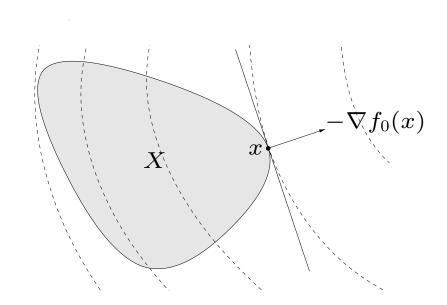
$$f_0(z) \le \theta f_0(x) + (1 - \theta) f_0(y) < f_0(x)$$

which contradicts our assumption that x is locally optimal

Optimality criterion for differentiable f_0

x is optimal if and only if it is feasible and

$$\nabla f_0(x)^T(y-x) \ge 0$$
 for all feasible y



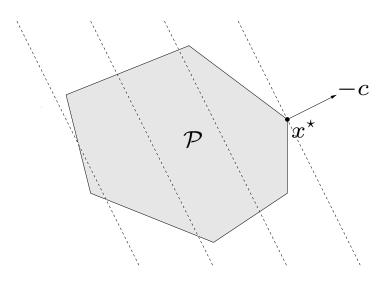
if nonzero, $\nabla f_0(x)$ defines a supporting hyperplane to feasible set X at x

Linear program (LP)

minimize
$$c^T x + d$$

subject to $Gx \leq h$
 $Ax = b$

- convex problem with affine objective and constraint functions
- feasible set is a polyhedron



Examples

diet problem: choose quantities x_1, \ldots, x_n of n foods

- ullet one unit of food j costs c_j , contains amount a_{ij} of nutrient i
- ullet healthy diet requires nutrient i in quantity at least b_i

to find cheapest healthy diet,

minimize
$$c^T x$$

subject to $Ax \succeq b$, $x \succeq 0$

piecewise-linear minimization

minimize
$$\max_{i=1,...,m} (a_i^T x + b_i)$$

equivalent to an LP

minimize
$$t$$
 subject to $a_i^T x + b_i \leq t, \quad i = 1, \dots, m$

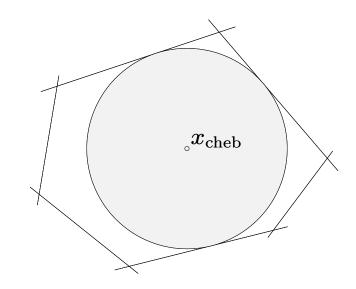
Chebyshev center of a polyhedron

Chebyshev center of

$$\mathcal{P} = \{ x \mid a_i^T x \le b_i, \ i = 1, \dots, m \}$$

is center of largest inscribed ball

$$\mathcal{B} = \{x_c + u \mid ||u||_2 \le r\}$$



• $a_i^T x \leq b_i$ for all $x \in \mathcal{B}$ if and only if

$$\sup\{a_i^T(x_c+u) \mid ||u||_2 \le r\} = a_i^T x_c + r||a_i||_2 \le b_i$$

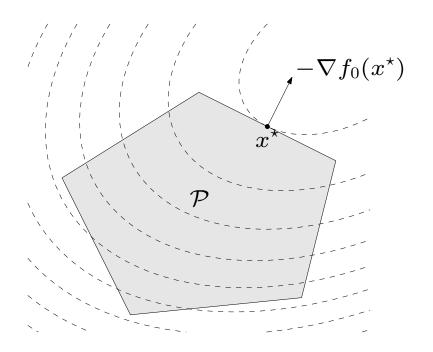
ullet hence, x_c , r can be determined by solving the LP

maximize
$$r$$
 subject to $a_i^T x_c + r ||a_i||_2 \leq b_i, \quad i = 1, \dots, m$

Quadratic program (QP)

minimize
$$(1/2)x^TPx + q^Tx + r$$
 subject to $Gx \leq h$ $Ax = b$

- $P \in \mathbf{S}_{+}^{n}$, so objective is convex quadratic
- minimize a convex quadratic function over a polyhedron



Examples

least-squares

minimize
$$||Ax - b||_2^2$$

- analytical solution $x^* = A^{\dagger}b$ (A^{\dagger} is pseudo-inverse)
- can add linear constraints, e.g., $l \leq x \leq u$

linear program with random cost

minimize
$$\bar{c}^T x + \gamma x^T \Sigma x = \mathbf{E} \, c^T x + \gamma \, \mathbf{var}(c^T x)$$
 subject to $Gx \leq h$, $Ax = b$

- ullet c is random vector with mean \bar{c} and covariance Σ
- ullet hence, c^Tx is random variable with mean \bar{c}^Tx and variance $x^T\Sigma x$
- \bullet $\gamma>0$ is risk aversion parameter; controls the trade-off between expected cost and variance (risk)

Norm approximation

minimize
$$||Ax - b||$$

 $(A \in \mathbf{R}^{m \times n} \text{ with } m \geq n, \| \cdot \| \text{ is a norm on } \mathbf{R}^m)$ interpretations of solution $x^* = \operatorname{argmin}_x \|Ax - b\|$:

- **geometric**: Ax^* is point in $\mathcal{R}(A)$ closest to b
- estimation: linear measurement model

$$y = Ax + v$$

y are measurements, x is unknown, v is measurement error given y=b, best guess of x is x^{\star}

• **optimal design**: x are design variables (input), Ax is result (output) x^* is design that best approximates desired result b

examples

• least-squares approximation $(\|\cdot\|_2)$: solution satisfies normal equations

$$A^T A x = A^T b$$

$$(x^* = (A^T A)^{-1} A^T b \text{ if } \mathbf{rank} A = n)$$

• Chebyshev approximation $(\|\cdot\|_{\infty})$: can be solved as an LP

minimize
$$t$$
 subject to $-t\mathbf{1} \preceq Ax - b \preceq t\mathbf{1}$

• sum of absolute residuals approximation $(\|\cdot\|_1)$: can be solved as an LP

minimize
$$\mathbf{1}^T y$$
 subject to $-y \leq Ax - b \leq y$

Penalty function approximation

minimize
$$\phi(r_1) + \cdots + \phi(r_m)$$

subject to $r = Ax - b$

 $(A \in \mathbf{R}^{m \times n}, \phi : \mathbf{R} \to \mathbf{R} \text{ is a convex penalty function})$

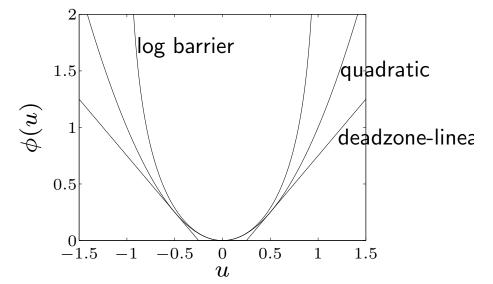
examples

- quadratic: $\phi(u) = u^2$
- deadzone-linear with width *a*:

$$\phi(u) = \max\{0, |u| - a\}$$

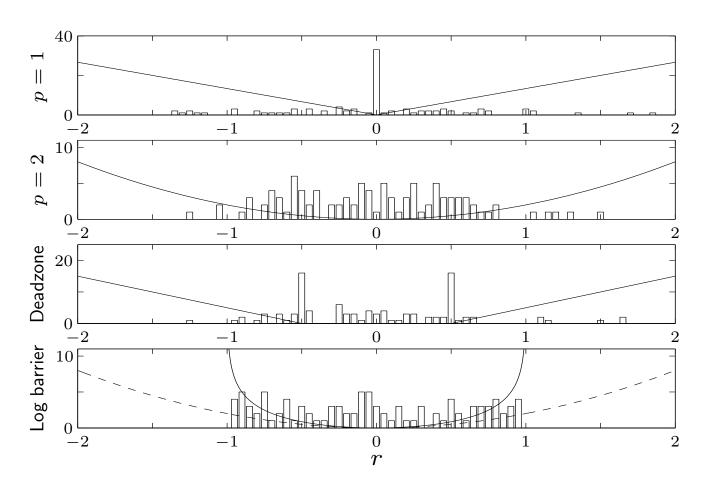
• log-barrier with limit *a*:

$$\phi(u) = \begin{cases} -a^2 \log(1 - (u/a)^2) & |u| < a \\ \infty & \text{otherwise} \end{cases}$$



example (m = 100, n = 30): histogram of residuals for penalties

$$\phi(u) = |u|, \quad \phi(u) = u^2, \quad \phi(u) = \max\{0, |u| - a\}, \quad \phi(u) = -\log(1 - u^2)$$

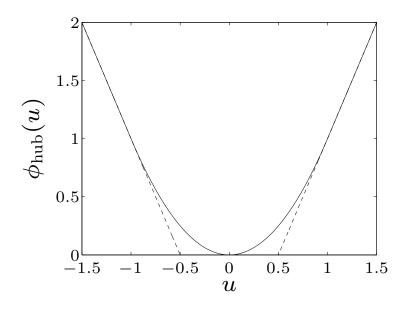


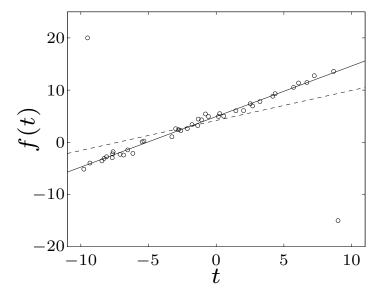
shape of penalty function has large effect on distribution of residuals

Huber penalty function (with parameter M)

$$\phi_{\text{hub}}(u) = \begin{cases} u^2 & |u| \le M \\ M(2|u| - M) & |u| > M \end{cases}$$

linear growth for large u makes approximation less sensitive to outliers





- left: Huber penalty for M=1
- right: affine function $f(t) = \alpha + \beta t$ fitted to 42 points t_i , y_i (circles) using quadratic (dashed) and Huber (solid) penalty

Least-norm problems

minimize
$$||x||$$
 subject to $Ax = b$

 $(A \in \mathbf{R}^{m \times n} \text{ with } m \leq n, \| \cdot \| \text{ is a norm on } \mathbf{R}^n)$

interpretations of solution $x^* = \operatorname{argmin}_{Ax=b} ||x||$:

- **geometric:** x^* is point in affine set $\{x \mid Ax = b\}$ with minimum distance to 0
- estimation: b = Ax are (perfect) measurements of x; x^* is smallest ('most plausible') estimate consistent with measurements
- **design:** x are design variables (inputs); b are required results (outputs) x^* is smallest ('most efficient') design that satisfies requirements

examples

• least-squares solution of linear equations ($\|\cdot\|_2$): can be solved via optimality conditions

$$2x + A^T \nu = 0, \qquad Ax = b$$

• minimum sum of absolute values $(\|\cdot\|_1)$: can be solved as an LP

tends to produce sparse solution x^{\star}

extension: least-penalty problem

minimize
$$\phi(x_1) + \cdots + \phi(x_n)$$

subject to $Ax = b$

 $\phi: \mathbf{R} \to \mathbf{R}$ is convex penalty function

Regularized approximation

minimize (w.r.t.
$$\mathbf{R}_{+}^{2}$$
) $(\|Ax - b\|, \|x\|)$

 $A \in \mathbf{R}^{m \times n}$, norms on \mathbf{R}^m and \mathbf{R}^n can be different

interpretation: find good approximation $Ax \approx b$ with small x

- estimation: linear measurement model y = Ax + v, with prior knowledge that ||x|| is small
- **optimal design**: small x is cheaper or more efficient, or the linear model y = Ax is only valid for small x
- robust approximation: good approximation $Ax \approx b$ with small x is less sensitive to errors in A than good approximation with large x

Scalarized problem

minimize
$$||Ax - b|| + \gamma ||x||$$

- ullet solution for $\gamma>0$ traces out optimal trade-off curve
- other common method: minimize $||Ax b||^2 + \delta ||x||^2$ with $\delta > 0$

Tikhonov regularization

minimize
$$||Ax - b||_2^2 + \delta ||x||_2^2$$

can be solved as a least-squares problem

minimize
$$\left\| \begin{bmatrix} A \\ \sqrt{\delta}I \end{bmatrix} x - \begin{bmatrix} b \\ 0 \end{bmatrix} \right\|_2^2$$

solution
$$x^* = (A^T A + \delta I)^{-1} A^T b$$

Signal reconstruction

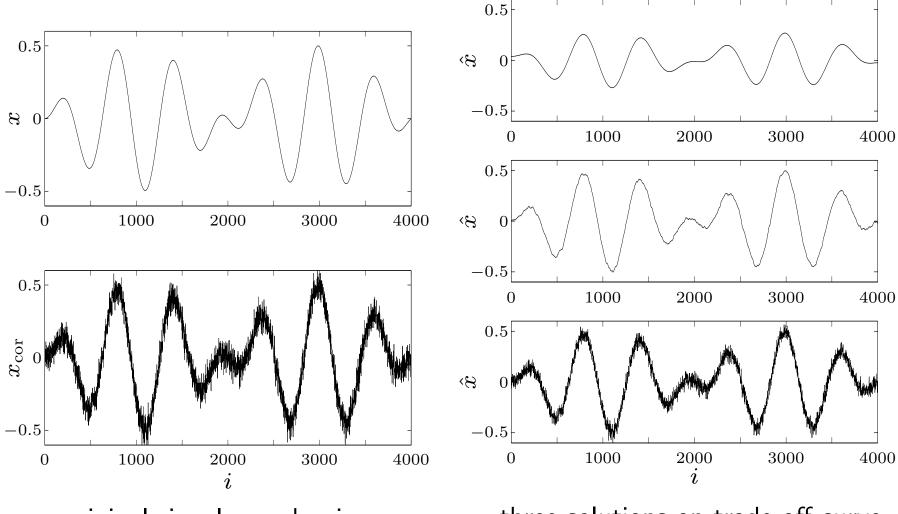
minimize (w.r.t.
$$\mathbf{R}_{+}^{2}$$
) $(\|\hat{x} - x_{\text{cor}}\|_{2}, \phi(\hat{x}))$

- $x \in \mathbf{R}^n$ is unknown signal
- $x_{cor} = x + v$ is (known) corrupted version of x, with additive noise v
- variable \hat{x} (reconstructed signal) is estimate of x
- $\phi: \mathbf{R}^n \to \mathbf{R}$ is regularization function or smoothing objective

examples: quadratic smoothing, total variation smoothing:

$$\phi_{\text{quad}}(\hat{x}) = \sum_{i=1}^{n-1} (\hat{x}_{i+1} - \hat{x}_i)^2, \qquad \phi_{\text{tv}}(\hat{x}) = \sum_{i=1}^{n-1} |\hat{x}_{i+1} - \hat{x}_i|$$

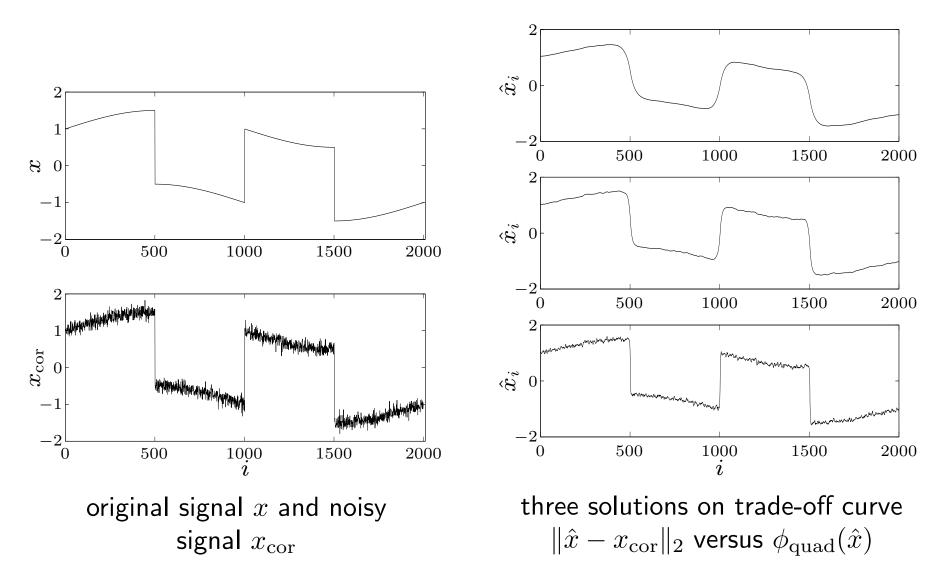
quadratic smoothing example



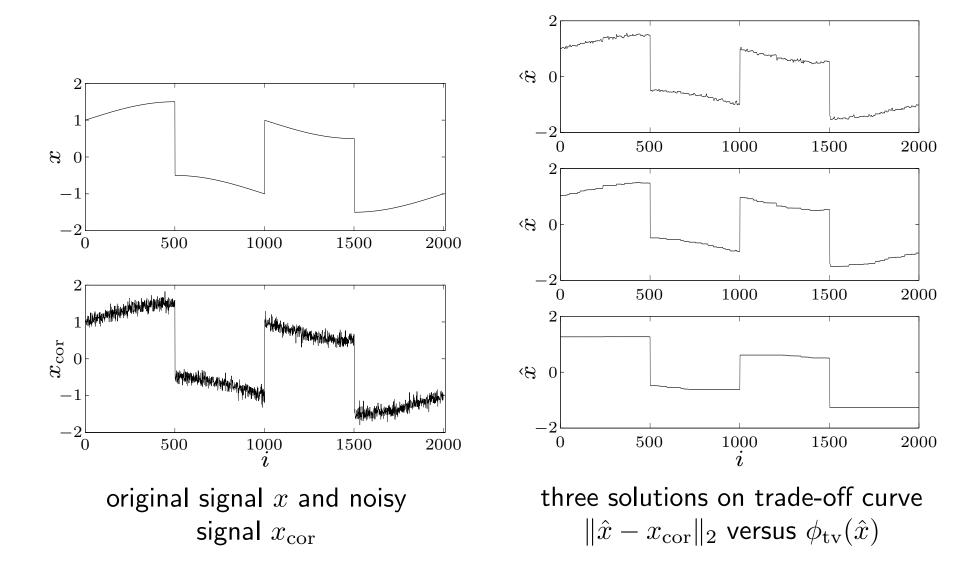
original signal x and noisy signal $x_{\rm cor}$

three solutions on trade-off curve $\|\hat{x} - x_{\text{cor}}\|_2$ versus $\phi_{\text{quad}}(\hat{x})$

total variation reconstruction example



quadratic smoothing smooths out noise and sharp transitions in signal



total variation smoothing preserves sharp transitions in signal

Linear measurements with IID noise

linear measurement model

$$y_i = a_i^T x + v_i, \quad i = 1, \dots, m$$

- $x \in \mathbf{R}^n$ is vector of unknown parameters
- v_i is IID measurement noise, with density p(z)
- y_i is measurement: $y \in \mathbf{R}^m$ has density $p_x(y) = \prod_{i=1}^m p(y_i a_i^T x)$

maximum likelihood estimate: any solution x of

maximize
$$l(x) = \sum_{i=1}^{m} \log p(y_i - a_i^T x)$$

(y is observed value)

examples

ullet Gaussian noise $\mathcal{N}(0,\sigma^2)$: $p(z)=(2\pi\sigma^2)^{-1/2}e^{-z^2/(2\sigma^2)}$,

$$l(x) = -\frac{m}{2}\log(2\pi\sigma^2) - \frac{1}{2\sigma^2}\sum_{i=1}^{m}(a_i^T x - y_i)^2$$

ML estimate is LS solution

• Laplacian noise: $p(z) = (1/(2a))e^{-|z|/a}$

$$l(x) = -m \log(2a) - \frac{1}{a} \sum_{i=1}^{m} |a_i^T x - y_i|$$

ML estimate is ℓ_1 -norm solution

• uniform noise on [-a, a]:

$$l(x) = \begin{cases} -m \log(2a) & |a_i^T x - y_i| \le a, \quad i = 1, \dots, m \\ -\infty & \text{otherwise} \end{cases}$$

ML estimate is any x with $|a_i^T x - y_i| \le a$

Logistic regression

random variable $y \in \{0,1\}$ with distribution

$$p = \mathbf{prob}(y = 1) = \frac{\exp(a^T u + b)}{1 + \exp(a^T u + b)}$$

- a, b are parameters; $u \in \mathbf{R}^n$ are (observable) explanatory variables
- ullet estimation problem: estimate a, b from m observations (u_i,y_i)

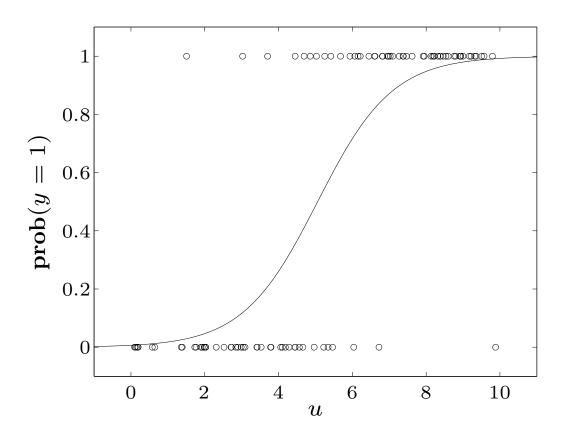
log-likelihood function (for $y_1 = \cdots = y_k = 1$, $y_{k+1} = \cdots = y_m = 0$):

$$l(a,b) = \log \left(\prod_{i=1}^{k} \frac{\exp(a^{T}u_{i} + b)}{1 + \exp(a^{T}u_{i} + b)} \prod_{i=k+1}^{m} \frac{1}{1 + \exp(a^{T}u_{i} + b)} \right)$$

$$= \sum_{i=1}^{k} (a^{T}u_{i} + b) - \sum_{i=1}^{m} \log(1 + \exp(a^{T}u_{i} + b))$$

concave in a, b

example (n = 1, m = 50 measurements)



- circles show 50 points (u_i, y_i)
- solid curve is ML estimate of $p = \exp(au + b)/(1 + \exp(au + b))$

Minimum volume ellipsoid around a set

Löwner-John ellipsoid of a set C: minimum volume ellipsoid \mathcal{E} s.t. $C \subseteq \mathcal{E}$

- parametrize \mathcal{E} as $\mathcal{E} = \{v \mid ||Av + b||_2 \leq 1\}$; w.l.o.g. assume $A \in \mathbf{S}_{++}^n$
- $\operatorname{vol} \mathcal{E}$ is proportional to $\det A^{-1}$; to compute minimum volume ellipsoid,

minimize (over
$$A$$
, b) $\log \det A^{-1}$ subject to $\sup_{v \in C} \|Av + b\|_2 \le 1$

convex, but evaluating the constraint can be hard (for general C)

finite set
$$C = \{x_1, ..., x_m\}$$
:

minimize (over
$$A$$
, b) $\log \det A^{-1}$ subject to $||Ax_i + b||_2 \le 1, \quad i = 1, \dots, m$

also gives Löwner-John ellipsoid for polyhedron $\mathbf{conv}\{x_1,\ldots,x_m\}$

Maximum volume inscribed ellipsoid

maximum volume ellipsoid \mathcal{E} inside a convex set $C \subseteq \mathbf{R}^n$

- parametrize \mathcal{E} as $\mathcal{E} = \{Bu + d \mid ||u||_2 \le 1\}$; w.l.o.g. assume $B \in \mathbf{S}_{++}^n$
- $\operatorname{vol} \mathcal{E}$ is proportional to $\det B$; can compute \mathcal{E} by solving

maximize
$$\log \det B$$

subject to $\sup_{\|u\|_2 \le 1} I_C(Bu+d) \le 0$

(where $I_C(x) = 0$ for $x \in C$ and $I_C(x) = \infty$ for $x \notin C$) convex, but evaluating the constraint can be hard (for general C)

polyhedron
$$\{x \mid a_i^T x \leq b_i, i = 1, \dots, m\}$$
:

maximize
$$\log \det B$$

subject to $\|Ba_i\|_2 + a_i^T d \leq b_i, \quad i = 1, \dots, m$

(constraint follows from $\sup_{\|u\|_{2} \le 1} a_{i}^{T}(Bu + d) = \|Ba_{i}\|_{2} + a_{i}^{T}d$)

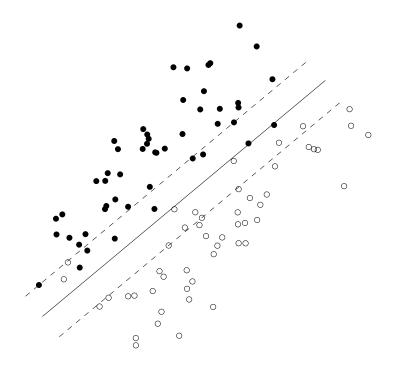
Support vector classifier

minimize
$$\|a\|_2 + \gamma (\mathbf{1}^T u + \mathbf{1}^T v)$$

subject to $a^T x_i + b \ge 1 - u_i, \quad i = 1, \dots, N$
 $a^T y_i + b \le -1 + v_i, \quad i = 1, \dots, M$
 $u \succeq 0, \quad v \succeq 0$

produces point on trade-off curve between inverse of margin $2/\|a\|_2$ and classification error, measured by total slack $\mathbf{1}^T u + \mathbf{1}^T v$

same example as previous page, with $\gamma=0.1$:



quadratic discrimination: $f(z) = z^T P z + q^T z + r$

$$x_i^T P x_i + q^T x_i + r \ge 1,$$
 $y_i^T P y_i + q^T y_i + r \le -1$

can add additional constraints (e.g., $P \leq -I$ to separate by an ellipsoid) **polynomial discrimination**: F(z) are all monomials up to a given degree



separation by ellipsoid

separation by 4th degree polynomial

Placement and facility location

- N points with coordinates $x_i \in \mathbf{R}^2$ (or \mathbf{R}^3)
- ullet some positions x_i are given; the other x_i 's are variables
- ullet for each pair of points, a cost function $f_{ij}(x_i,x_j)$

placement problem

minimize
$$\sum_{i\neq j} f_{ij}(x_i, x_j)$$

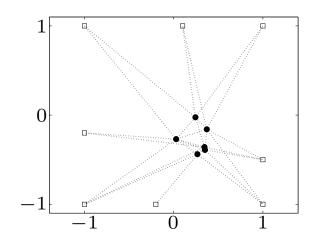
variables are positions of free points

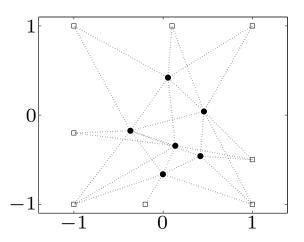
interpretations

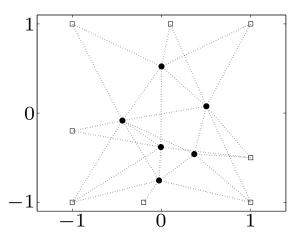
- ullet points represent plants or warehouses; f_{ij} is transportation cost between facilities i and j
- ullet points represent cells on an IC; f_{ij} represents wirelength

example: minimize $\sum_{(i,j)\in\mathcal{A}} h(\|x_i - x_j\|_2)$, with 6 free points, 27 links

optimal placement for h(z)=z, $h(z)=z^2$, $h(z)=z^4$







histograms of connection lengths $||x_i - x_j||_2$

