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Team Note of National Potato

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ALL BELOW HERE ARE USELESS IF YOU READ THE STATEMENT WRONG

1 Network Flow

1.1 Dinic's Algorithm

```
1. build level graph with bfs ( O(E) )
 2. flow blocking flow in level graph ( O(VE) )
 In each step, level grows at least by 1, and eventually grows upto O(V)
 So total time complexity is O(V^2E)
struct dinic { // O(V^2 E)
 struct edg { int v, c, r; };
 int n;
 vi dis. itr:
 vector<vector<edg>> g;
 dinic(int n) : n(n), g(n, vector<edg>()), dis(n), itr(n) { }
 void addedge(int u, int v, int c) {
   g[u].pb({ v, c, sz(g[v]) });
   g[v].pb({u, 0, sz(g[u])-1});
 bool bfs(int s, int t) { // build level graph
   dis.assign(n, 0), itr.assign(n, 0);
   queue<int> q;
   q.push(s);
   dis[s] = 1;
   while(q.size()) {
     int u = q.front(); q.pop();
     for(auto& [v, c, r] : g[u]) {
       if(c > 0 && !dis[v]) {
         dis[v] = dis[u] + 1;
         q.push(v);
       }
     }
   }
   return dis[t] > 0:
 int dfs(int u, int t, int f) { // get blocking flow
   if(u == t) return f:
   for( ; itr[u] < g[u].size(); itr[u]++) {</pre>
     auto& [v, c, r] = g[u][itr[u]];
     if(c > 0 \&\& dis[v] == dis[u] + 1) {
       int w = dfs(v, t, min(f, c));
       if(w) {
         g[u][itr[u]].c -= w;
         g[v][r].c += w;
         return w:
       }
     }
   }
   return 0;
 i64 nflow(int s, int t) { // network flow
   i64 \text{ ret} = 0:
   while(bfs(s, t)) {
     int r;
     while(r = dfs(s, t, 2e9)) ret += r, debug("----");
```

```
return ret;
 }
1.2 Hofcroft-Karp Bipartite Matching
    - alternating path: path consists of (x, y, x, y, x) for X = { x | x is matched edge}, Y
= { y | y is not mathed edge }
    - augmenting path: path consists of (y, x, y, x, y)
    - if augmenting path exists, we can match one more edge with fliping matched state (x,
y, x, y, x
   For maximaum matching A, B
   1. lv[0] = { v | v in A and v is not matched }
    2. starting from lv[0] vertices, get alternating path with bfs
   3. starting from lv[0] vertices, get augmenting path with dfs
   * min cover: selecting minimum vertices to cover all edges
   * max independent set: selecting maximum vertices not connected with edge
    * V - min cover = max independent set
struct hofcroft {
 int n, m;
 vi dis. l. r. vis. chk:
 hofcroft(int n, int m) : n(n), m(m), g(n, vi()) { }
  void addedge(int u, int v) { g[u].pb(v); }
 bool bfs() { // build alternating path starts from lv[0] nodes
   queue<int> q:
   bool ok = 0;
    dis.assign(n, 0);
   FOR(u, 0, n) {
     if(l[u] == -1 && !dis[u]) {
       q.push(u);
        dis[u] = 1;
   }
    while(q.size()) {
     int u = q.front(); q.pop();
     for(int v : g[u]) {
        if(r[v] == -1) ok = 1; // v is not matched
        else if(!dis[r[v]]) { // if v is matched, u>v>r[v] can be path
          dis[r[v]] = dis[u] + 1;
         q.push(r[v]);
     }
   }
   return ok;
  bool dfs(int u) { // find augmenting path and flip it!
   if(vis[u]) return 0; // augmenting path start/end with non-matched vertices
   vis[u] = 1;
   for(int v : g[u]) {
     if(r[v] == -1 \mid | (dis[r[v]] == dis[u] + 1 && dfs(r[v]))) {
       l[u] = v; r[v] = u;
       return 1:
```

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```
}
   }
   return 0;
 int match() { // bipartite match
   1.assign(n, -1);
   r.assign(m, -1);
   int ret = 0;
   while(bfs()) {
     vis.assign(n, 0);
     FOR(u, 0, n) if(1[u] == -1 \&\& dfs(u)) ++ret;
   }
   return ret;
  void rdfs(int u) { // dfs matched
   if(chk[u]) return;
   chk[u] = 1;
   for(int v : g[u]) {
     chk[v + n] = 1:
     rdfs(r[v]);
   }
 }
 vi getcover() { // get min cover vertices
   match();
   chk.assign(n+m, 0);
   FOR(u, 0, n) if(1[u] == -1) rdfs(u);
   vi ret:
   FOR(u, 0, n) if(!chk[u]) ret.pb(u);
   FOR(u, n, n+m) if(chk[u]) ret.pb(u);
   return ret:
 }
};
    Minimum Cost Maximum Flow
const int WINF = 0x3ffffffff, FINF = 0x3ffffffff; // weight/flow inf
struct mcmf {
 struct edg { int v, c, r, w; };
 int n;
 vi dis, par, peg;
 vector<bool> ing;
 vector<vector<edg> > g;
 mcmf(int n) : n(n), g(n, vector < edg > ()), par(n), peg(n) { }
 void addedge(int u, int v, int c, int w) {
   g[u].pb({ v, c, sz(g[v]), w });
   g[v].pb({u, 0, sz(g[u])-1, -w});
 bool spfa(int s, int t) {
   dis.assign(n, WINF);
   inq.assign(n, 0);
   queue<int> q;
   dis[s] = 0;
   inq[s] = 1;
   q.push(s);
   bool ok = 0:
   while(q.size()) {
     int u = q.front(); q.pop();
     if(u == t) ok = 1;
```

```
ina[u] = 0:
      FOR(eidx, 0, g[u].size()) {
        auto [v, c, r, w] = g[u][eidx];
        if(c > 0 && dis[v] > dis[u] + w) {
          dis[v] = dis[u] + w;
          par[v] = u;
          peg[v] = eidx;
          if(!inq[v]) {
           inq[v] = 1;
            q.push(v);
     }
   }
   return ok;
 ii flow(int s, int t) { // return (max_flow, min_cost)
   int cost = 0, flow = 0;
    while(spfa(s, t)) {
     int cur = FINF;
     for(int u = t; u != s; u = par[u]) cur = min(cur, g[par[u]][peg[u]].c);
     for(int u = t; u != s; u = par[u]) {
        int r = g[par[u]][peg[u]].r;
        g[par[u]][peg[u]].c -= cur;
        g[u][r].c += cur;
     flow += cur:
      cost += dis[t] * cur:
   return { flow, cost };
 }
};
1.4 Ford-Fulkerson Algorithm
// Caution: All vertices' idx > 0 (par[S] = 0)
const int MAXND = 500, S = 1, T = 2, INF = 0x3ffffffff:
struct nflow {
 struct edge {
    edge* rev:
   int v, c, f; // need initialized
   edge(int v, int c) : v(v), c(c), f(0) { }
   int res() { return c-f; }
   int flow(int x) { f += x, rev -> f -= x; }
 };
  vector<edge*> g[MAXND];
  int par[MAXND];
  edge* pedg[MAXND]:
 nflow() {
   FOR(i, 0, MAXN) g[i] = vector<edge*>();
  void addedge(int u, int v, int c) {
   edge *uv = new edge(v, c), *vu = new edge(u, 0);
   uv->rev = vu, vu->rev = uv;
    g[u].pb(uv), g[v].pb(vu);
  i64 maxflow() {
   i64 \text{ ret} = 0;
```

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```
while(true) {
      memset(par, 0, sizeof(par));
      queue<int> q;
      q.push(S);
      par[S] = S;
      while(q.size()) {
        int u = q.front(); q.pop();
        for(auto e : g[u]) {
          if(e->res() && !par[e->v]) {
            q.push(e->v);
            par[e->v] = u;
            pedg[e->v] = e;
            if(e->v == T) break;
        if(par[T]) break;
      if(!par[T]) break;
      int flow = INF:
      for(int u = T; u != S; u = par[u]) flow = min(flow, pedg[u]->res());
      for(int u = T; u != S; u = par[u]) pedg[u]->flow(flow);
      ret += flow:
    return ret;
};
     Naive Bipartite Match
const int MAXN = 5e2+10:
int vis[MAXN], ato[MAXN], bto[MAXN];
bool dfs(int u) {
 if(vis[u]) return 0;
  vis[u] = 1;
  for(int v : g[u]) {
   if(bto[v] == -1 || dfs(bto[v])) {
      ato[u] = v;
      bto[v] = u:
      return 1;
 }
 return 0;
int bimatch() {
  memset(ato, -1, sizeof(ato)), memset(bto, -1, sizeof(bto));
 int ret = 0:
  FOR(u, 0, n) {
   memset(vis, 0, sizeof(vis));
   if(dfs(u)) ++ret;
  return ret;
2 Graph
2.1 \quad 2\text{-SAT} + SCC
  2-SAT: (A | | B) && (C | | D) && (E | | F) ...
 1. X \mid \mid Y = !X \rightarrow Y, !Y \rightarrow X (Proposition)
```

```
False: T -> F. True: Others
 2. !X, X in same SCC: no solution
 3. For every SCC, each node in same SCC must have same flag (if both T, F exists in same
SCC. T->F exists)
 4. Assign False to (Don't have in edge & Unassigned node) and erase node
   - sort nodes topologically, iterate nodes with assigning False to var if var is
unassigned
      - !X node: X = True, X node: X = False
struct sat2 {
 struct tarjan {
   int n, ncnt, scnt;
   vi scc, dis;
   vvi g;
   stack<int> sta;
   tarjan(int n): n(n), g(n, vi()) { } // n: number of variables (NOT NODES!)
   void addedge(int u, int v) { g[u].pb(v); } // directed graph
   int f(int u) {
     int ret = dis[u] = ncnt++:
     sta.push(u);
     for(int v : g[u]) {
       if(dis[v] == -1) ret = min(ret, f(v));
        else if(scc[v] == -1) ret = min(ret, dis[v]);
     if(ret == dis[u]) {
       while(1) {
         int t = sta.top(); sta.pop();
         scc[t] = scnt:
         if(t == u) break;
        ++scnt;
     return ret;
   }
   vi& get_scc() {
     ncnt = scnt = 0:
     scc = dis = vi(n, -1);
     sta = stack<int>():
     FOR(i, 0, n) if(dis[i] == -1) f(i);
     dis.clear();
     return scc:
   }
 };
 int n;
 vi res:
 tarjan tj;
 sat2(int n) : n(n), tj(2*n) { }
 int nd(int u, int neg) { return u + neg*n; } // var u's node
 int neg(int u) { return (u+n)%(2*n); } // ~u
 void addedge(int u, int nu, int v, int nv) { // add (X || Y) clauses
   u = nd(u, nu), v = nd(v, nv);
   tj.addedge(neg(u), v);
   tj.addedge(neg(v), u);
 vi& solve() { // return solved vars, if no solution return vi()
   vi& scc = tj.get_scc();
```

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```
FOR(u, 0, n) if(scc[u] == scc[u+n]) return res;
   res.assign(n, -1);
   vi ord(2*n);
   FOR(i, 0, 2*n) \text{ ord}[i] = i;
   sort(ALL(ord), [&](int u, int v) { return scc[u] > scc[v]; });
   FOR(i, 0, 2*n) {
     int u = ord[i];
     if(res[u/n] == -1) res[u/n] = !(u < n);
   return res;
 }
};
     Cut-Vertex
// Find Bridge
const int MAXN = 3e5. INF = 0x7ffffffff:
int ncnt, vid[MAXN], par[MAXN];
vii g[MAXN]; // (v, bridge flag)
int dfs(int u) {
 int ret = vid[u] = ncnt++:
 for(auto& e : g[u]) {
   int v = e.se, c = INF;
   if(par[u] == v) continue; // Tree edge
   if(vid[v] == -1) { // Tree Edge}
     par[v] = u;
     c = min(c, dfs(v)):
   } else c = min({ c, vid[v], vid[u] }); // Forward/Backward Edge
   if(c > vid[u]) e.fi = 1; // Bridge
   ret = min(ret, c):
 }
 return ret;
   String
3.1 KMP Algorithm
string s, t;
vi getpi(const string& str) {
 int n = str.size(), len = 0:
 vi pi(n, 0);
 FOR(i, 1, n) {
   while(len && str[len] != str[i]) len = pi[len-1];
   if(str[len] == str[i]) pi[i] = ++len;
 return pi:
vi kmp() {
 int n = s.size(), m = t.size(), len = 0;
 vi ret, pi = getpi(t);
 FOR(i, 0, n) {
   while(len && s[i] != t[len]) len = pi[len-1]:
   if(s[i] == t[len] \&\& ++len == m) ret.pb(i-len+1), len = pi[len-1];
 }
 return ret;
```

```
3.2 Rabin-Karp Algorithm
// f(p) = s[0] + s[1] * p + s[2] * p^2 + ... + s[n-1] * p^{n-1}
// h[i+1] = p * (h[i] - s[i] * s^(m-1)) + s[i+m]
// --> sub first character from hash > hash degree up > add last character to hash
const i64 MUL = 232153, MD = 1012924417; // be careful for MD not MOD
void mod(i64& x) { x %= MD; if(x < 0) x += MD; }
vi rabin(string& s, string& t) { // return start indexs
 vi ret:
 i64 \text{ ht} = 0, hs = 0, mul = 1;
 RFOR(i, sz(t)-1, 0) { // get t's hash
   mod(ht += mul * t[i] % MD):
   mod(mul = mul * MUL);
 mul = 1;
 RFOR(i, sz(t)-1, 0) { // get s's hash for first sz(t) string
   mod(hs += mul * s[i] % MD):
   if(i != 0) mod(mul = mul * MUL); // mul must be p^{m-1)
 if(hs == ht) ret.pb(0):
 FOR(i, sz(t), sz(s)) {
   mod(hs -= mul*s[i-sz(t)]%MD):
   mod(hs *= MUL);
   mod(hs += s[i]);
   if(hs == ht) ret.pb(i-sz(t)+1):
 }
 return ret:
3.3 Trie(Array)
const int MAX NODE = 1e6+10:
int cld[MAX_NODE][30];
i64 cnt[MAX_NODE];
int ncnt = 1:
void push(const string& x) {
 int u = 0:
 FOR(i, 0, x.size()) {
   if(!cld[u][x[i]-'a']) cld[u][x[i]-'a'] = ncnt++;
   u = cld[u][x[i]-'a'];
 ++cnt[u]:
void calc_back(int u) {
 FOR(i, 0, 30) {
   if(cld[u][i]) {
      calc_back(cld[u][i]);
      cnt[u] += cnt[cld[u][i]]:
 }
 cnt[u] %= MOD:
3.4 Aho-Corasick Algorithm
const int MAXNODE = 1e5+10, MAXC = 26, INITCHAR = 'a';
struct ahocorasick {
 int ncnt. t[MAXNODE][MAXC]. f[MAXNODE]. chk[MAXNODE];
 ahocorasick(): ncnt(0) { memset(t, 0, sizeof(t)), memset(f, 0, sizeof(f)), memset(chk, 0,
 sizeof(chk)): }
```

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```
void insert(const string& s) {
   int u = 0;
   for(auto i : s) {
     i -= INITCHAR:
     if(!t[u][i]) t[u][i] = ++ncnt;
     u = t[u][i]:
   }
   chk[u] = 1; // end of string
 void precalc() {
   queue<int> q;
   FOR(i, 0, MAXC) if(t[0][i]) q.push(t[0][i]);
   // calculdate fail, chk with bfs
   while(q.size()) {
     int x = q.front(); q.pop();
     FOR(i, 0, MAXC) {
       if(t[x][i]) {
         int u = x, p = f[u];
         while(p && !t[p][i]) p = f[p]; // find fail link
         u = t[u][i], p = t[p][i]; // goto original target node
         f[u] = p;
         if(chk[p]) chk[u] = 1;
         q.push(u);
       }
     }
   }
 bool query(const string& s) {
   int u = 0;
   for(auto i : s) {
     i -= INITCHAR:
     while(u && !t[u][i]) u = f[u];
     if(chk[u = t[u][i]]) return true;
   }
   return false:
 }
};
3.5
     Suffix Array + LCP
 sa[i] = ordered suffix (suffix's start position)
 ord[i] = [i:]'s index in sa (ord[sa[i]] = i)
 lcp[i] = longest common prefix length of two suffix [i-1:], [i:]
 LCP's Lemma
 1. Two adjacent in SA suffixes' LCP is always bigger than which of non-adjacents
 2. lcp(sa[i-1], sa[i]) = h, h >= 1 then
   lcp(sa[i-1]+1, sa[i]+1) = h-1
 So that lcp[sa[i]+1] >= h-1 because it is always bigger than lcp(sa[i-1]+1, sa[i]+1) by
Lemma 1
 and by Lemma 2, lcp(sa[i-1]+1, sa[i]+1) = h-1
*/
struct sfxarray {
 int n;
 string& str;
 vi sa, lcp, ord;
```

```
sfxarrav(string& str) : str(str). n(str.size()) { }
 void getsa() {
   sa = ord = vi(n+1);
   FOR(i, 0, n) sa[i] = i, ord[i] = str[i]; ord[n] = 0;
   for(int t = 1; t \le n; t *= 2) {
     int sz = max(257, n+1):
     vi cnt, tmp;
     cnt = tmp = vi(sz, 0);
     FOR(i, 0, n) ++cnt[ord[min(n, i+t)]];
     FOR(i, 1, sz) cnt[i] += cnt[i-1];
     FOR(i, 0, n) tmp[--cnt[ord[min(n, i+t)]]] = i;
     cnt = vi(sz. 0):
     FOR(i, 0, n) ++cnt[ord[i]];
     FOR(i, 1, sz) cnt[i] += cnt[i-1]:
     RFOR(i, n-1, 0) sa[--cnt[ord[tmp[i]]]] = tmp[i];
     tmp[sa[0]] = 1;
     FOR(i, 1, n) {
       int u = sa[i-1], v = sa[i];
       tmp[v] = tmp[u] + (ord[u] < ord[v] || ord[u+t] < ord[v+t]);
     ord = tmp;
     if(ord[sa[n-1]] == n) break;
   FOR(i, 0, n) --ord[i];
 void getlcp() {
   lcp = vi(n, 0);
   for(int i = 0, len = 0; i < n; ++i, len = max(0, len-1)) {
     if(ord[i]) {
       for(int j = sa[ord[i]-1]; str[i+len] == str[j+len]; ++len);
       lcp[ord[i]] = len:
   }
 }
 tuple<vi, vi, vi> build() { getsa(), getlcp(); return { sa, lcp, ord }; }
3.6 Manacher's Algorithm
const int MAXN = 1000005:
int aux[2 * MAXN - 1];
void solve(int n, int *str, int *ret){
 // *ret : number of nonobvious palindromic character pair
 for(int i=0: i<n: i++){
   aux[2*i] = str[i];
   if(i != n-1) aux[2*i+1] = -1;
 int p = 0, c = 0:
 for(int i=0; i<2*n-1; i++){
   int cur = 0;
   if(i \le p) cur = min(ret[2 * c - i], p - i);
   while(i - cur - 1 >= 0 && i + cur + 1 < 2*n-1 && aux[i-cur-1] == aux[i+cur+1]){
   }
   ret[i] = cur:
   if(i + ret[i] > p){
     p = i + ret[i];
      c = i:
```

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```
}
4 Math
4.1 Fraction
// dependency: GCD(i64 a, i64 b)
struct frac {
 i64 a, b;
 frac(i64 _a=0, i64 _b=1) : a(_a), b(_b) { if(a == 0 && b == 0) b = 1; assert(b != 0);}
 // Essential: Basic Operations
 void relax() { i64 g = GCD(abs(a), abs(b)); a /= g, b /= g; }
 frac operator + (const frac &ot) const { return { a * ot.b + ot.a * b, b * ot.b }; }
 frac operator - (const frac &ot) const { return { a * ot.b - ot.a * b, b * ot.b }; }
 frac operator * (const frac &ot) const { return { a * ot.a, b * ot.b }; }
 frac operator / (const frac &ot) const { return { a * ot.b, b * ot.a }; }
 frac operator - () { return { -a, b }; }
 // Essential: Basic Comparation
 bool operator == (const frac& ot) const { return a * ot.b == ot.a * b; }
 bool operator < (const frac& ot) const { return a * ot.b < ot.a * b: }
 bool operator <= (const frac& ot) const { return a * ot.b <= ot.a * b; }</pre>
 bool operator > (const frac& ot) const { return ot <= *this; }</pre>
 bool operator >= (const frac& ot) const { return ot < *this; }</pre>
 // Optional: Advanced Operations
 const frac& operator += (const frac &ot) { return *this = *this + ot; }
 const frac& operator -= (const frac &ot) { return *this = *this - ot; }
 const frac& operator *= (const frac &ot) { return *this = *this * ot; }
 const frac& operator /= (const frac &ot) { return *this = *this / ot: }
// fraction IO
ostream& operator<< (ostream& os, const frac& frac_x) { return os << frac_x.a << "/" <<
istream& operator>> (istream& os, frac& frac_x) {
 os >> frac_x.a >> frac_x.b;
 frac x.relax():
 return os:
4.2 Matrix
// Do not use this class as const
typedef i64 ELEM;
struct mat {
 int n, m;
 vector<vector<ELEM> > ar:
 // ---- constructor, assignment ---- //
 mat(int n, int m, ELEM x = 0) : n(n), m(m), ar(n, vector < ELEM > (m, x)) { }
 mat(int n = 0) : mat(n, n) \{ \}
 mat(const mat& o) { n = o.n, m = o.m, ar = o.ar; }
 mat(const vector<vector<ELEM>>& ar) : n(ar.size()), m(ar.size() ? ar[0].size() : 0),
 // ---- get field ---- //
 operator const vector<vector<ELEM> >& () const { return ar; }
 vector<ELEM>& operator[](int i) { return ar[i]; }
  const vector<ELEM>& operator[](int i) const { return ar[i]; }
 // ---- calculate ---- //
```

```
mat pow(i64 x) const {
  assert(n == m && 0 <= x);
  mat a(*this), ret = eye(n);
  while(x) {
    if(x\%2) ret = ret * a;
    a = a * a:
    x /= 2;
  return ret:
mat operator * (const mat& o) const {
  assert(m == o.n):
  mat ret(n, o.m);
  FOR(i, 0, n) {
    FOR(i, 0, o.m) {
      FOR(k, 0, m) {
       ret[i][j] += ar[i][k] * o[k][j] % MOD;
       ret[i][j] %= MOD;
    }
  }
  return ret:
mat operator + (const mat& o) const {
  assert(n == o.n \&\& m == o.m):
  mat ret(n, m);
  FOR(i, 0, n) FOR(j, 0, n) ret[i][j] = (ar[i][j] + o[i][j]) % MOD;
  return ret:
}
mat operator - (const mat& o) const {
  assert(n == o.n && m == o.m);
  mat ret(n, m);
  FOR(i, 0, n) FOR(j, 0, n) ret[i][j] = (ar[i][j] - o[i][j]) % MOD;
mat operator * (const ELEM x) const {
  mat ret = ar;
  FOR(i, 0, n) FOR(j, 0, m) ret[i][j] = ret[i][j] * x % MOD;
  return ret:
mat operator / (const ELEM x) const {
  mat ret = ar;
  FOR(i, 0, n) FOR(j, 0, m) ret[i][j] = ret[i][j] / x % MOD;
  return ret:
const mat& operator - () {
 FOR(i, 0, n) FOR(j, 0, m) ar[i][j] = -ar[i][j];
  return *this:
}
  // If use dp matrix as: state = state * dpmat
  // use rotated dpmat and horizontal state mat
mat rotate() const {
  mat ret(m, n);
  FOR(i, 0, n) FOR(j, 0, m) ret[j][i] = ar[i][j];
  return ret;
```

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```
static mat eve(const int size) {
   mat ret(size);
   FOR(i, 0, size) ret[i][i] = 1;
   return ret:
 }
 // return matrix dp
   // dp[i] = ar[0] * dp[i-n] + ar[1] * dp[i-n+1] + ... + ar[n-1] * dp[i-1]
 // data matrix br: br[0] = dp[i], br[1] = dp[i-1] ...
 // If DP equation contains constant, fix one element for constant
  static mat dpmat(const vector<ELEM>& ar) {
   int n = ar.size();
   mat ret(n, n):
   FOR(i, 0, n-1) ret[i][i+1] = 1; // transition prev dp values
   FOR(i, 0, n) ret[n-1][i] = ar[i]; // DP equation
   return ret;
 }
};
ostream& operator<<(ostream& os, const mat& v) { for(auto vv : (vector<vector<ELEM> >)v) os
<< vv << ENDL: return os: }
4.3 Miller-Rabin Test
namespace miller_rabin{ // O(logP)
 i64 mul(i64 x, i64 y, i64 mod){ return (__int128) x * y % mod; }
 i64 ipow(i64 x, i64 y, i64 p){
   i64 ret = 1, piv = x \% p;
   while(v){
     if(y&1) ret = mul(ret, piv, p);
     piv = mul(piv, piv, p);
     y >>= 1;
   return ret:
  bool miller_rabin(i64 x, i64 a){
   if(x \% a == 0) return 0:
   i64 d = x - 1;
   while(1){
     i64 \text{ tmp} = ipow(a, d, x);
     if(d&1) return (tmp != 1 && tmp != x-1);
     else if(tmp == x-1) return 0;
     d >>= 1;
   }
 bool isprime(i64 x){
   for(auto &i : {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37}){
     if(x == i) return 1:
     if (x > 40 \&\& miller_rabin(x, i)) return 0;
   }
   if(x \le 40) return 0;
   return 1;
 }
4.4 Euler's Sieve
const int RANGE = 2e7:
int pn, spf[RANGE], pr[RANGE]; // spf[x] = min prime factor of x
void eulerSieve() {
 FOR(x, 2, RANGE) {
```

```
if(!spf[x]) spf[x] = pr[pn++] = x; // if x is prime, spf[x] = x
   for(int j = 0; x*pr[j] < RANGE; ++j) {</pre>
     spf[x*pr[j]] = pr[j];
     if(x % pr[j] == 0) break;
   }
 }
}
4.5 Binomial Coefficient
// 1. c[n][r] = c[n-1][r-1] + c[n-1][r]
// 2. \text{ nCr} = n! / ((n-r)! * r!)
const int RANGE = 1e6;
// precalc: n!, n!^(-1) --> O(NlogP)
void precalc_1() {
 i64 ftr[RANGE], iftr[RANGE];
 ftr[0] = iftr[0] = 1;
 FOR(i, 1, RANGE) {
   ft[i] = (ft[i-1] * (i64)i) % MOD;
   ift[i] = (ift[i-1] * POW((i64)i, MOD-2)) % MOD:
}
// (n-1)!^{(-1)} = n*n!^{(-1)} --> O(n+logP)
void precalc_2() {
 i64 ftr[RANGE], iftr[RANGE];
 ftr[0] = iftr[0] = 1;
 FOR(i, 1, RANGE) ftr[i] = ftr[i-1] * i;
 iftr[RANGE-1] = POW(ftr[RANGE-1], MOD-2);
 RFOR(i, RANGE-2, 0) iftr[i] = (i * iftr[i+1]) % MOD;
// inv(1) = 1, inv(1) = -floor(p/i) + inv(p%i) --> O(n)
void precalc_3() {
 i64 inv[RANGE+1], ftr[RANGE], iftr[RANGE]:
 inv[1] = 1;
 FOR(i, 2, RANGE+1) {
   inv[i] = inv[MOD % i] - (MOD / i);
   if(inv[i] < 0) inv[i] += MOD;</pre>
   inv[i] %= MOD;
 }
4.6 Inclusion-Exclusion Principle
FOR(i, 1, (1 << n)) { // get Union(A, B, C, D ...)
   int bits = __builtin_popcount(i);
   if(bits % 2): // add to ans
   else; // sub to ans
4.7 Josephus Problem
/* O(n)
 f(n, k) = last survived person for n-people, k-cycle
 < basic idea >
 except 1 element from f(n, k), then answer is f(n-1, k)
 but f(n-1, k) need to be repositioned to starting from kth's next person
 < 1-indexed >
 f(1, k) = 1
 f(n, k) = ((f(n-1, k) + k-1) \% n) + 1
```

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if(k < 0) return 0:

```
< 0-indexed >
 f(1, k) = 0;
 f(n, k) = ((f(n-1, k) + k) \% n)
// O(KlogN) algorithm
long long joseph (long long n,long long k) {
    if (n==1LL) return OLL;
   if (k==1LL) return n-1LL;
   if (k>n) return (joseph(n-1LL,k)+k)%n;
   long long cnt=n/k;
   long long res=joseph(n-cnt,k);
   res-=n%k:
   if (res<OLL) res+=n;</pre>
    else res+=res/(k-1LL):
    return res;
    Segment Tree
5.1 Fenwick Tree Tricks
struct fenwick {
  int n:
  vector<i64> t:
  void init(int _n) { n = _n, t = vector<i64>(n+1, 0); }
  void add(int u, i64 x) { for(++u: u < t.size(): u += (u&-u)) t[u] += x: }
  i64 sum(int u) {
        i64 \text{ ret} = 0;
   for(++u; u > 0; u -= (u\&-u)) ret += t[u];
   return ret;
  i64 operator[](int u) {
   int ret = t[u], p = u - (u&-u);
    while(u != p) {
     ret -= t[u]:
     u = (u\&-u);
   }
   return ret:
 // Can use all elements are positive.
  // find x which ( sum[a[0]..a[x-1]] < k <= sum[a[0]..a[x]] )
  int lower(i64 k) {
   if(k < 0) return 0;
   int l = (1 << (8*sizeof(int) - __builtin_clz(n)) - 1);</pre>
    while(1 > 0 && u <= n) {
     int tu = u + 1;
     if(k > t[tu]) u = tu, k -= t[tu];
     do 1 >>= 1;
      while(1 > 0 & u + 1 > n):
   }
   return u;
  // find x which ( sum[a[0]..a[x-1]] \le k \le sum[a[0]..a[x]] )
 int upper(i64 k) {
```

```
int l = (1 << (8*sizeof(int) - __builtin_clz(n)) - 1);</pre>
   int u = 0;
   while(1 > 0 && u <= n) {
     int tu = u + 1;
     if(k >= t[tu]) u = tu, k -= t[tu]:
     do 1 >>= 1;
     while(1 > 0 \&\& u + 1 > n);
   return u;
 }
};
5.2 Fenwick Tree 2D (Sparse)
struct segtree {
 vi vs[RANGE], t[RANGE];
 // Notify segtree update access on (x, y)
 void initpos(int x, int y) {
   for(++x; x < RANGE; x += (x\&-x)) {
     ys[x].pb(y);
   }
 7
 // Execute after notifying (x, y)
 void init() {
   FOR(i, 0, RANGE) sort(ALL(ys[i])), UNIQUE(ys[i]), t[i].assign(ys[i].size()+1, 0);
 // add (x, y) to c
 void add(int x, int y, int c) {
   for(++x; x < RANGE; x += (x\&-x)) {
     for(int j = getidx(ys[x], y)+1; j < sz(t[x]); j += (j&-j)) {
       t[x][j] += c;
   }
 // partial sum of ([..x], [..y])
 int sum(int x, int v) {
   int ret = 0;
   for(++x; x > 0; x -= (x\&-x)) {
     int j = getidx(ys[x], y);
     if(j == ys[x].size() || ys[x][j] > y) --j;
     for(++j; j > 0; j = (j\&-j)) {
       ret += t[x][i];
   }
   return ret;
5.3 Fenwick Tree Range Update/Query
struct rfenwick { // using 2 basic fenwick tree
 fenwick ax, b;
 void init(int n) { ax.init(n), b.init(n); }
 void add(int u, i64 x) { b.add(u, x); }
 void add(int s, int e, i64 x) {
   ax.add(s, x);
   ax.add(e+1, -x);
   b.add(s, -x * (s-1));
```

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```
b.add(e+1, x * e):
 }
 i64 sum(int u) {
   return u * ax.sum(u) + b.sum(u);
};
5.4 Segment Tree (Loop)
template < class T, class C>
struct segtree {
 int n:
 vector<T> t:
 void build(const vector<T>& ar) {
   n = ar.size():
   t.assign(n*2, 0);
   FOR(i, 0, n) t[n+i] = ar[i];
   RFOR(i, n-1, 1) t[i] = C()(t[i*2], t[i*2+1]);
 void mod(int u. T k) {
   for(t[u += n] = k; u > 1; u /= 2) t[u/2] = C()(t[u], t[u^1]);
 T query(int s, int e) {
   T ret = 0;
   for(s += n, e += n; s < e; s /= 2, e /= 2) {
     if(s & 1) ret = C()(ret, t[s++]):
     if (e \& 1) ret = C()(ret, t[--e]);
   }
   return ret:
 }
};
5.5 Segment Tree 2D (Dense)
struct segtree {
 int n, m;
 void init(const vvi& ar) {
   n = ar.size(), m = ar[0].size();
   t.assign(2*n, vi(2*m, 0));
   FOR(y, 0, n) \{ // push in leaf \}
     FOR(x, 0, m) {
       t[n+y][m+x] = ar[y][x];
   RFOR(y, 2*n-1, 1) { // construct
     RFOR(x, 2*m-1, 1) {
       if(y < n) t[y][x] = t[y*2][x] + t[y*2+1][x];
       if(x < m) t[v][x] = t[v][x*2] + t[v][x*2+1]:
   }
 }
 void modify(int y, int x, int c) {
   t[v + n][x + m] = c: // leaf update
   for(y += n; y > 0; y /= 2) {
     for(int x2 = x + m; x2 > 0; x2 /= 2) {
       if(y < n) t[y][x2] = t[y*2][x2] + t[y*2+1][x2];
       if(x2 < m) t[y][x2] = t[y][x2*2] + t[y][x2*2+1];
     }
```

```
}
 int query(int sy, int sx, int ey, int ex) {
   int ret = 0:
   for(sy += n, ey += n; sy < ey; sy /= 2, ey /= 2) {
     for(int x1 = sx + m, x2 = ex + m; x1 < x2; x1 /= 2, x2 /= 2) {
        if(sy&1) {
          if(x1\&1) ret += t[sy][x1];
         if(x2\&1) ret += t[sy][x2-1];
        if(ey&1) {
         if(x1\&1) ret += t[ey-1][x1];
         if(x2\&1) ret += t[ey-1][x2-1];
       if(x1&1) ++x1;
        if(x2&1) --x2;
     if(sv&1) ++sv;
     if(ey&1) --ey;
   return ret;
5.6 Lazy Propagation
const int ST_MAX = 1<<21, lf = ST_MAX/2;</pre>
struct segtree{
 i64 t[ST_MAX], d[ST_MAX];
 segtree(){ memset(t, 0, sizeof(t)), memset(d, 0, sizeof(d)); }
 void build(){ RFOR(i, lf-1, 1) t[i] = t[i*2]+ t[i*2+1]; } // !! BUILD !!
 void propagate(int u, int ns, int ne){
   if(!d[u]) return;
   if(u < lf){ // propagate to childs
     d[u*2] += d[u];
     d[u*2+1] += d[u];
   t[u] += d[u] * (ne-ns); // update node
   d[u] = 0:
 void add(int s, int e, int x){ add(s, e, x, 1, 0, lf); } // [s, e)
 void add(int s, int e, int x, int u, int ns, int ne){
   propagate(u, ns, ne);
   if(e <= ns || ne <= s) return:
   if(s <= ns && ne <= e){
     d[u] += x;
     propagate(u, ns, ne);
     return:
   }
   int mid = (ns+ne)/2;
   add(s, e, x, u*2, ns, mid), add(s, e, x, u*2+1, mid, ne);
   t[u] = t[u*2] + t[u*2+1];
 i64 sum(int s, int e){ return sum(s, e, 1, 0, lf); } // [s, e)
 i64 sum(int s, int e, int u, int ns, int ne){
   propagate(u, ns, ne);
   if(e <= ns || ne <= s) return 0;
   if(s <= ns && ne <= e) return t[u];</pre>
```

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```
int mid = (ns+ne)/2:
    return sum(s, e, u*2, ns, mid) + sum(s, e, u*2+1, mid, ne);
 }
};
5.7 Persistent Segment Tree
struct node {
  int x;
  node *1, *r;
  node(int _x = 0, node* 1 = 0, node* r = 0) : x(_x), 1(1), r(r) { }
  node* addtree(int u, int c, int ns = 0, int ne = MAXN-1) {
    if(ns <= u && u <= ne) {
     if (ns == ne) return new node(x + c, 0, 0):
     int mid = (ns+ne)/2:
     return new node(x + c, 1->addtree(u, c, ns, mid), r->addtree(u, c, mid+1, ne));
   return this;
  int querv(int s, int e, int ns = 0, int ne = MAXN-1) {
    if (s <= ns && ne <= e) return x;
    if (ne < s || e < ns) return 0;
   int mid = (ns+ne)/2:
    return 1->query(s, e, ns, mid) + r->query(s, e, mid+1, ne);
} *root[MAXN+1];
5.8 Persistent Segment Tree (Array)
struct pst {
  i64 x[MAXN*LOGN]:
  int 1[MAXN*LOGN], r[MAXN*LOGN], tcnt;
  int base(int ns = 0, int ne = MAXN-1) { // make 0th tree
   int u = tcnt++:
   1[u] = u, r[u] = u;
  int make(int idx, int c, int u, int ns = 0, int ne = MAXN-1) { // update from u-rooted
   if(idx < ns || ne < idx) return u;
    int v = tcnt++:
    if (ns == ne) x[v] = (x[u] + c) \% MOD;
    else {
      int m = (ns+ne)/2:
     1[v] = make(idx, c, l[u], ns, m);
     r[v] = make(idx, c, r[u], m+1, ne);
     x[v] = (x[l[v]] + x[r[v]]) % MOD:
   }
   return v;
  i64 query(int s, int e, int u, int ns = 0, int ne = MAXN-1) { // query from u-rooted
    if(s <= ns && ne <= e) return x[u]:
   if (ne < s || e < ns) return 0;
   int m = (ns+ne)/2:
    return (query(s, e, l[u], ns, m) + query(s, e, r[u], m+1, ne)) % MOD;
 }
};
     HLD (Vertex)
  HLD with costed vertex.
```

```
usually (dfs init, lca, decomposite, eidx, query) don't need to be changed.
 just modify (segtree, init_segs), and if segtree function changed modify (update,
query_to)
const int MAXN = 1e5+10, LOGN = 18, INF = 0x7ffffffff;
struct hld vtx {
 struct segtree {
   int n;
   vi t:
   void init(int _n) { n = _n; t.assign(2*n, INF); }
   void update(int u. int x) {
     for(t[u += n] = x; u > 1; u /= 2) t[u/2] = min(t[u], t[u^1]);
   }
   int querv(int s. int e) {
     int ret = INF:
     for(s += n, e += n; s < e; s /= 2, e /= 2) {
       if(s&1) ret = min(ret, t[s++]):
       if(e\&1) ret = min(ret, t[--e]);
     return ret;
   }
 };
 int n, rt;
 vi ssz, dep, hidx;
 vvi g, par, hvy;
 vector<segtree> segs;
 hld_vtx(vvi\& g, int rt) : g(g), rt(rt), n(g.size()), ssz(n, 0), dep(n, 0), hidx(n, -1),
 par(n, vi(LOGN, 0)) {
   par[rt][0] = rt;
   dfs init(rt):
   decomposite(rt);
   init_segs();
 void dfs_init(int u) { // initialize dfs info
   ssz[u] = 1:
   FOR(j, 1, LOGN) par[u][j] = par[par[u][j-1]][j-1];
   for(int v : g[u]) {
     if(par[u][0] == v) continue:
     par[v][0] = u;
     dep[v] = dep[u] + 1;
     dfs init(v):
     ssz[u] += ssz[v];
 7
 int lca(int u, int v) { // consider par[root] = root
   if(dep[u] < dep[v]) swap(u, v);</pre>
   int dif = dep[u] - dep[v];
   FOR(j, 0, LOGN) if (dif & (1<<j)) u = par[u][j];
   if(u != v) {
     RFOR(j, LOGN-1, 0) if (par[u][j] != par[v][j]) u = par[u][j], v = par[v][j];
     u = par[u][0];
   }
   return u;
 void decomposite(int rt) { // decomposite tree
   queue<int> q;
```

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```
q.push(rt);
   while(q.size()) {
     int u = q.front(); q.pop();
     for(int v : g[u]) if(par[v][0] == u) q.push(v);
     int p = par[u][0];
     if (u != rt \&\& ssz[u]*2 >= ssz[p]) { // extend h-path}
       hidx[u] = hidx[p];
       hvy[hidx[u]].pb(u);
     } else { // create h-path
       hidx[u] = hvy.size();
       hvy.pb(vi());
       hvy[hidx[u]].pb(u);
     }
   }
  void init_segs() { // initialize segtrees
   segs.assign(hvy.size(), segtree());
   FOR(i, 0, hvy.size()) segs[i].init(hvy[i].size()); // m nodes
 int vidx(int u) { // get v's index in h-path
   return dep[u] - dep[hvy[hidx[u]][0]];
 void update(int u, int x) { // update v's cost
   if(x == 0) segs[hidx[u]].update(vidx(u), INF);
   else segs[hidx[u]].update(vidx(u), vidx(u));
 int query(int v) { // root->v query
   return query_to(0, v);
 int query_to(int u, int v) { // return u->v path's query
   if(hidx[u] == hidx[v]) {
     int res = segs[hidx[u]].query(vidx(u), vidx(v)+1);
     if(res == INF) return INF;
     return hvv[hidx[u]][res];
   int res = query_to(u, par[hvy[hidx[v]][0]][0]);
   if(res != INF) return res;
   res = segs[hidx[v]].query(0, vidx(v)+1);
   if(res == INF) return INF;
   return hvv[hidx[v]][res];
 }
};
5.10 HLD (Edge)
 HLD with costed edge.
 Unlike the normal HLD, top edge of each chains are also belongs to chain.
 usually (dfs_init, lca, decomposite, eidx, query) don't need to be changed.
 just modify (segtree, init_segs), and if segtree function changed modify (update,
query_to)
*/
const int LOGN = 18;
struct hld edge {
 struct segtree { // just edit segtree ( currently half-open interval [s, e) )
   int n;
   vi t;
```

```
void init(int n) { n = n: t.assign(2*n, 0): }
  void update(int u, int x) {
    for(t[u += n] = x; u > 1; u /= 2) t[u/2] = \max(t[u], t[u^1]);
  int query(int s, int e) {
    int ret = 0:
    for(s += n, e += n; s < e; s /= 2, e /= 2) {
      if(s\&1) ret = max(ret, t[s++]);
      if(e&1) ret = max(ret, t[--e]):
    return ret;
  }
};
int n. rt:
vi ssz, dep, hidx;
vvi g, par, hvv;
vector<segtree> segs;
hld_edge(vvi\& g, int rt) : g(g), rt(rt), n(g.size()), ssz(n, 0), dep(n, 0), hidx(n, -1),
par(n, vi(LOGN, 0)) {
  par[rt][0] = rt;
  dfs_init(rt);
  decomposite(rt);
  init_segs();
}
void dfs init(int u) { // initialize dfs info
  ssz[u] = 1;
  FOR(j, 1, LOGN) par[u][j] = par[par[u][j-1]][j-1];
  for(int v : g[u]) {
    if(par[u][0] == v) continue;
    par[v][0] = u;
    dep[v] = dep[u] + 1;
    dfs_init(v);
    ssz[u] += ssz[v];
  }
int lca(int u, int v) { // consider par[root] = root
  if(dep[u] < dep[v]) swap(u, v);</pre>
  int dif = dep[u] - dep[v];
  FOR(j, 0, LOGN) if (dif & (1<<j)) u = par[u][j];
  if(u != v) {
    RFOR(j, LOGN-1, 0) if(par[u][j] != par[v][j]) u = par[u][j], v = par[v][j];
    u = par[u][0];
  }
  return u;
void decomposite(int rt) { // decomposite tree
  hidx[rt] = -1;
  queue<int> q;
  q.push(rt);
  while(q.size()) {
    int u = q.front(); q.pop();
    for(int v : g[u]) if(par[v][0] == u) q.push(v);
    if(u != rt) {
      int p = par[u][0]:
      if(p != rt && ssz[u]*2 >= ssz[p]) { // extend h-path (only if h-path)
        hidx[u] = hidx[p]:
```

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```
}
   }
  void init_segs() { // initialize segtrees
   segs.assign(hvv.size(), segtree());
   FOR(i, 0, hvy.size()) segs[i].init(hvy[i].size()-1); // m vertices: m-1 edges
 int eidx(int v) { // get u->v edge index in h-path
   return dep[par[v][0]] - dep[hvv[hidx[v]][0]];
 void update(int u, int v, int x) { // u->v edge update
   if(par[u][0] == v) swap(u, v):
   assert(par[v][0] == u);
   segs[hidx[v]].update(eidx(v), x);
  int query_to(int u, int v) { // return u->v path's query
   if(u == v) return 0;
   // modify range if segtree use closed interval [s. e]
   if(hidx[u] == hidx[v]) return segs[hidx[u]].query(eidx(u)+1, eidx(v)+1); // e(u)+1
   because target is edge
   return max(query_to(u, hvy[hidx[v]][0]), segs[hidx[v]].query(0, eidx(v)+1)); // query
   tail path + recur
 int query(int u, int v) {
   int t = lca(u, v):
   return max(query_to(t, u), query_to(t, v));
 }
};
   Miscellaneous
6.1 Preset
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template < class key, class value, class cmp = less < key >>
using treemap = tree<key, value, less<int>, rb_tree_tag, tree_order_statistics_node_update>;
// key, val, comp, implements, 노드 불변 규칙
template < class key, class cmp = less < key >>
using treeset = tree<key, null_type, cmp, rb_tree_tag, tree_order_statistics_node_update>;
#ifdef LOCAL BOOKNU
#define debug(...) cerr << "[" << #__VA_ARGS__ << "]:", debug_out(__VA_ARGS__)
#else
#define debug(...) 42
#endif
#define FOR(i, f, n) for(int (i) = (f); (i) < (int)(n); ++(i))
#define RFOR(i, f, n) for(int (i) = (f); (i) >= (int)(n); --(i))
#define pb push_back
```

hvv[hidx[u]].pb(u):

hvy[hidx[u]].pb(p);

hvy[hidx[u]].pb(u);

hvy.pb(vi());

hidx[u] = hvy.size();

} else { // create h-path (l-path or root-h-path)

```
#define emb emplace_back
#define fi first
#define se second
#define ENDL '\n'
#define sz(A) (int)(A).size()
#define ALL(A) A.begin(), A.end()
#define UNIQUE(c) (c).resize(unique(ALL(c)) - (c).begin())
typedef pair<int, int> ii;
typedef pair<int, ii> iii;
typedef vector<int> vi;
typedef vector<vi> vvi;
typedef vector<ii> vii:
typedef vector<vii> vvii;
typedef long long i64:
typedef unsigned long long ui64;
inline int getidx(const vi& ar, int x) { return lower_bound(ALL(ar), x) - ar.begin(); }
inline i64 GCD(i64 a, i64 b) { i64 n; if(a < b) swap(a, b); while(b != 0) { n = a % b; a =
b: b = n:  return a:  }
inline i64 LCM(i64 a, i64 b) { if(a == 0 || b == 0) return GCD(a, b); return a / GCD(a, b) *
inline i64 CEIL(i64 n, i64 d) { return n / d + (i64)(n % d != 0); } // for positive numbers
inline i64 ROUND(i64 n, i64 d) { return n / d + (i64)((n % d) * 2 >= d); }
const i64 \text{ MOD} = 1e9+7;
inline i64 POW(i64 a, i64 n) {
 i64 ret:
 for(ret = 1; n; a = a*a%MOD, n /= 2) { if(n%2) ret = ret*a%MOD; }
 return ret:
template <class T> ostream& operator<<(ostream& os, vector<T> v) {
 os << "[":
 int cnt = 0:
 for(auto vv : v) { os << vv; if(++cnt < v.size()) os << ","; }</pre>
 return os << "]":
template <class T> ostream& operator<<(ostream& os, set<T> v) {
 os << "[":
 int cnt = 0;
 for(auto vv : v) { os << vv: if(++cnt < v.size()) os << ".": }
 return os << "]":
template <class L. class R> ostream& operator<<(ostream& os. pair<L. R> p) { return os <<
"(" << p.fi << "," << p.se << ")"; }
void debug_out() { cerr << endl; }</pre>
template <typename Head, typename... Tail> void debug_out(Head H, Tail... T) { cerr << " "
<< H. debug out(T...): }
// ..... MAIN ..... //
void input() { }
int solve() { return 0; }
void execute() { input(), solve(); }
int main(void) {
#ifdef LOCAL BOOKNU
 freopen("__IO/input.txt", "r", stdin);
 // freopen("__IO/out.txt", "w", stdout);
#endif
  cin.tie(0), ios_base::sync_with_stdio(false);
  execute():
```

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```
return 0:
}
6.2 3D-Partial Sum
const int RANGE = 256;
int n, k, ps[RANGE] [RANGE] [RANGE], ar[RANGE] [RANGE];
int f(int x, int y, int z) {
 return (x < 0 \mid | y < 0 \mid | z < 0 \mid | x >= RANGE | | y >= RANGE | | z >= RANGE ? 0 :
 ps[x][y][z]);
int sum(int x1, int y1, int z1, int 1) {
 int x^2 = \min(RANGE - 1, x^1 + 1), y^2 = \min(RANGE - 1, y^1 + 1), z^2 = \min(RANGE - 1, z^1 + 1);
 --x1, --y1, --z1;
 return
   f(x2, y2, z2)
   - f(x1, y2, z2) - f(x2, y1, z2) - f(x2, y2, z1)
   + f(x1, y1, z2) + f(x1, y2, z1) + f(x2, y1, z1)
   - f(x1, y1, z1);
void init() {
 FOR(x, 0, RANGE) {
   FOR(y, 0, RANGE) {
     FOR(z, 0, RANGE) {
       ps[x][y][z] +=
         f(x-1, y, z) + f(x, y-1, z) + f(x, y, z-1)
         - f(x - 1, y - 1, z) - f(x - 1, y, z - 1) - f(x, y - 1, z - 1)
         + f(x - 1, y - 1, z - 1);
     }
   }
 }
    Knuth's Optimization
```

If three conditions satisfies in DP equation, the time complexity can be reduced from $O(n^3)$ to $O(n^2)$

- 1. DP Equation Form
 - $D[i][j] = \min_{i < k < j} (D[i][k] + D[k][j]) + C[i][j]$
- 2. Quadrangle Inequality
 - $C[a][c] + C[b][d] \le C[a][d] + C[b][c], a \le b \le c \le d$
- 3. Monotonicity
 - $C[b][c] \le C[a][d], a \le b \le c \le d$

Define A[i][j] = k for D[i][j] becomes minimum If condition 2, 3 been satisfied, bellowing inequality holds. $A[i][j-1] \leq A[i][j] \leq A[i+1][j]$

6.4 Bit Tricks

```
__builtin_clz(int x); // count leading-zero
__builtin_ctz(int x); // count tailing-zero
__builtin_clzll(i64 x);
__builtin_ctzll(i64 x);
__builtin_popcount(int x); // number of 1-bits
__builtin__ffs(int x); // lsb index (1-based, x = 0 -> 0)
```

```
floor(log2(n)): 31 - __builtin_clz(n|1);
// 00111, 01011, 01101, 01110, 10011, 10101...
i64 next_perm(i64 x) {
    i64 t = x | (x-1);
    return (t + 1) \mid (((^t \& -^t) - 1) >> (\_builtin_ctz(x) + 1))
6.5 Fast IO
class FastIO {
  int fd, bufsz;
  char *buf, *cur, *end;
 public:
  FastIO(int _fd = 0, int _bufsz = 1 << 20) : fd(_fd), bufsz(_bufsz) {
    buf = cur = end = new char[bufsz];
  }
  "FastIO() { delete[] buf; }
  bool readbuf() {
    cur = buf;
    end = buf + bufsz;
    while(true) {
      size_t res = fread(cur, sizeof(char), end - cur, stdin);
      if(res == 0) break:
      cur += res:
    }
    end = cur:
    cur = buf;
    return buf != end;
  bool hasNext() {
    while(true) {
      if(cur == end && !readbuf()) return false;
      if(isdigit(*cur) || *cur == '-') break;
      ++cur:
    return true;
  }
  int r() {
    while(true) {
      if(cur == end) readbuf();
      if(isdigit(*cur) || *cur == '-') break;
      ++cur:
    bool sign = true;
    if(*cur == '-') {
       sign = false;
       ++cur:
    int ret = 0;
    while(true) {
      if(cur == end && !readbuf()) break;
      if(!isdigit(*cur)) break;
      ret = ret * 10 + (*cur - '0'):
      ++cur;
    return sign ? ret : -ret;
} sc;
```

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```
6.6 Input Format
while(scanf("%d", &n) > 0) { // until input ends
   scanf("%d: (%d)", &x, &y); // formatted input
   for(int i = 0; i < x; ++i) scanf("%d", &z);</pre>
}
getline(cin, line);
6.7 String Parsing
vector<string> split(string& target, string regex) { // using regex
 vector<string> ret;
 std::regex rgx(regex);
  std::sregex_token_iterator iter(target.begin(),
   target.end(),
   rgx,
   -1):
  std::sregex_token_iterator end;
 for( ; iter != end; ++iter) ret.pb(*iter);
 return ret;
vector<string> space_split(string& s) { // split by whitespace
 istringstream iss(s);
 vector<string> ret(istream_iterator<string>{iss}, istream_iterator<string>());
 return ret;
}
std::to_string(42);
std::atoi("42");
6.8 Ordered Statistics Tree in g++
s.find_by_order(x); // 0-indexed
s.order_of_key(x) // 0-indexed, find first element x <= ar[idx]</pre>
6.9 Prime Numbers
  < 10<sup>k</sup>
                 number
                           divisors 2 3 5 71113171923293137
                    6
                                  4 1 1
                    60
                                 12 2 1 1
                   840
                                 32 3 1 1 1
                  7560
                                 64 3 3 1 1
                 83160
                                128 3 3 1 1 1
                 720720
                                240 4 2 1 1 1 1
               8648640
                                448 6 3 1 1 1 1
 8
               73513440
                                768 5 3 1 1 1 1 1
              735134400
                               1344 6 3 2 1 1 1 1
             6983776800
 10
                               2304 5 3 2 1 1 1 1 1
  11
            97772875200
                               4032 6 3 2 2 1 1 1 1
  12
           963761198400
                               6720 6 4 2 1 1 1 1 1 1
 13
          9316358251200
                              10752 6 3 2 1 1 1 1 1 1 1
 14
         97821761637600
                              17280 5 4 2 2 1 1 1 1 1 1
  15
        866421317361600
                              26880 6 4 2 1 1 1 1 1 1 1 1
  16
       8086598962041600
                              41472 8 3 2 2 1 1 1 1 1 1 1
  17
      74801040398884800
                              64512 6 3 2 2 1 1 1 1 1 1 1 1
  18 897612484786617600
                             103680 8 4 2 2 1 1 1 1 1 1 1 1
  < 10^k prime # of prime
                                      < 10^k
                                      10
                                                   999999967
 2
             97
                          25
                                      11
                                                  9999999977
```

997

168

12

99999999989

```
1229
                                                 999999999971
 5
           99991
                         9592
                                       14
                                                9999999999973
 6
          999983
                        78498
                                       15
                                               99999999999989
 7
         9999991
                       664579
                                              99999999999937
 8
        99999989
                      5761455
                                       17
                                             999999999999997
 9
       99999937
                     50847534
                                           999999999999999
6.10 Random
mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
mt19937 rng(0x14004); // 0x94949
int randint(int s, int e) { return uniform_int_distribution<int>(s, e)(rng); }
6.11 Hashing
struct chash {
 const int RANDOM = (long long)(make_unique<char>().get()) ^
  chrono::high_resolution_clock::now().time_since_epoch().count();
 static unsigned long long hash_f(unsigned long long x) {
   x += 0x9e3779b97f4a7c15;
   x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
   x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
   return x ^{\circ} (x >> 31):
 }
 static unsigned hash_combine(unsigned a, unsigned b) { return a * 31 + b; }
 int operator()(int x) const { return hash_f(x)^RANDOM; }
};
gp_hash_table<key, int, chash> mp;
int main() {
 mp[1] = 1:
```