

1. Consider the context-free grammar:

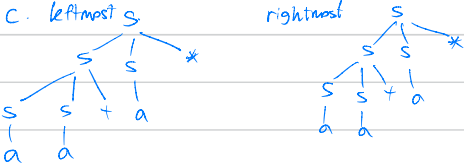
$S \rightarrow SS+ \mid SS^* \mid a$

and the string $aa+a^*$.

- a. Give a leftmost derivation for the string.
- b. Give a rightmost derivation for the string.
- c. Give a parse tree for the string.
- d. Is the grammar ambiguous or unambiguous? Justify your answer.

a. $S \Rightarrow SS^*$
 $\Rightarrow SS+S^*$
 $\Rightarrow aS+S^*$
 $\Rightarrow aa+S^*$
 $\Rightarrow aa+a^*$

b. $S \Rightarrow SS^*$
 $\Rightarrow Sa^*$
 $\Rightarrow SS+a^*$
 $\Rightarrow Sa+a^*$
 $\Rightarrow aa+a^*$



d. unambiguous, leftmost and rightmost derivatives give the same parse tree.

2. Consider the context-free grammar:

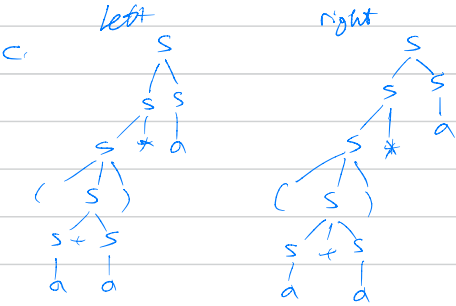
$S \rightarrow S+S \mid S(S) \mid S^* \mid a$

and the string $(a+a)^*a$.

- a. Give a leftmost derivation for the string.
- b. Give a rightmost derivation for the string.
- c. Give a parse tree for the string.
- d. Is the grammar ambiguous or unambiguous? Justify your answer.

a. $S \Rightarrow SS$
 $\Rightarrow S^*S$
 $\Rightarrow (S)^*S$
 $\Rightarrow (S+S)^*S$
 $\Rightarrow (b+S)^*S$
 $\Rightarrow (a+a)^*S$
 $\Rightarrow (a+a)^*a$

b. $S \Rightarrow SS$
 $\Rightarrow Sa$
 $\Rightarrow S^*a$
 $\Rightarrow (S)^*a$
 $\Rightarrow (S+S)^*a$
 $\Rightarrow (S+a)^*a$
 $\Rightarrow (a+a)^*a$



d. unambiguous, leftmost and rightmost derivatives give the same parse tree.

3. Design grammars for the following languages:

- a. The set of all strings of 0s and 1s such that every 0 is immediately followed by at least one 1.
- b. The set of all strings of 0s and 1s that are palindromes; that is, the string reads the same backward as forward.
- c. The set of all strings of 0s and 1s with an equal number of 0s and 1s.
- d. The set of all strings of 0s and 1s in which 011 does not appear as a substring.

a. $S \Rightarrow (0?1)^*$

or

$S \Rightarrow AB$
 $A \Rightarrow 1A \mid \epsilon$
 $B \Rightarrow 0B \mid \epsilon$
 $C \Rightarrow 0A$

b. $S \Rightarrow 0s0 \mid 1s1 \mid 0 \mid 1 \mid \epsilon$

c. $S \Rightarrow 0s1s \mid 1s0s \mid \epsilon$

d. $S \Rightarrow 1^*(0+1?)*$

4. The following is a grammar for regular expressions over symbols a and b only, using + in place of | for union, to avoid conflict with the use of vertical bar as a metasymbol in grammars:

rexpr \rightarrow rexpr + rterm | rterm
rterm \rightarrow rterm rfactor | rfactor
rfactor \rightarrow rfactor * | rprimary
rprimary \rightarrow a | b

- a. Left factor this grammar.
- b. Does left factoring make the grammar suitable for top-down parsing?
- c. In addition to left factoring, eliminate left recursion from the original grammar.
- d. Is the resulting grammar suitable for top-down parsing? Justify your answer.

a. Cannot do left factor

b. not suitable

c. $rexpr \rightarrow rterm A$
 $A \Rightarrow +rterm A \mid \epsilon$
 $rterm \Rightarrow rfactor B \mid \epsilon$
 $B \Rightarrow rfactor B \mid \epsilon$
 $rfactor \Rightarrow rprimary C$
 $C \Rightarrow *C \mid \epsilon$
 $rprimary \Rightarrow a \mid b$

d. suitable

Exercise 2

1. Consider the context-free grammar:

$$S \rightarrow S S + \mid S S * \mid a$$

and the string $aa + a^*$.

- Give a leftmost derivation for the string.
- Give a rightmost derivation for the string.
- Give a parse tree for the string.
- Is the grammar ambiguous or unambiguous? Justify your answer.

2. Consider the context-free grammar:

$$S \rightarrow S + S \mid S S \mid (S) \mid S * \mid a$$

and the string $(a+a)^*a$.

- Give a leftmost derivation for the string.
- Give a rightmost derivation for the string.
- Give a parse tree for the string.
- Is the grammar ambiguous or unambiguous? Justify your answer.

3. Design grammars for the following languages:

- The set of all strings of 0s and 1s such that every 0 is immediately followed by at least one 1.
- The set of all strings of 0s and 1s that are palindromes; that is, the string reads the same backward as forward.
- The set of all strings of 0s and 1s with an equal number of 0s and 1s.
- The set of all strings of 0s and as in which 011 does not appear as a substring.

4. The following is a grammar for regular expressions over symbols a and b only, using $+$ in place of $|$ for union, to avoid conflict with the use of vertical bar as a metasymbol in grammars:

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repr -> repr + rterm | rterm
rterm -> rterm rfactor | rfactor
rfactor -> rfactor * | rprimary
rprimary -> a | b

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- Left factor this grammar.
- Does left factoring make the grammar suitable for top-down parsing?
- In addition to left factoring, eliminate left recursion from the original grammar.
- Is the resulting grammar suitable for top-down parsing? Justify your answer.