

1. Given a following grammar:

1 $E \rightarrow FT, \checkmark$

2 $T \rightarrow ADFT, \checkmark$

3 $T \rightarrow e, \checkmark$

4 $AD \rightarrow +, - \checkmark$

5 $AD \rightarrow \cdot, \checkmark$

6 $F \rightarrow LK, \checkmark$

7 $K \rightarrow MUL LK, \checkmark$

8 $K \rightarrow e, \checkmark$ empty string

9 $MUL \rightarrow *, \checkmark$

10 $MUL \rightarrow /, \checkmark$

11 $L \rightarrow (E), \checkmark$

12 $L \rightarrow id, \checkmark$

a. Find the first and follow sets of the grammar.

b. The parsing table of the grammar

	First	Follow
E	id, (\$,)
F	id, (+, -, \$,)
T	e, +, -	\$,)
AD	+, -	id, (
K	e, *, /	\$, +, -,)
MUL	*, /	id, (
L	id, (*, /, +, -,), \$

	(id	+	-	\$	*	/)
E	1	1						
T			2	2	3			3
AD			4	5				
F	6	6						
K			8	8	8	7	7	8
MUL						9	10	
L	11	12						

2. From the parsing table in (1), use stack to simulate leftmost derivation as the LL(1)

parsing for stream of tokens id + id * (id + id).

No.	Stack	Tokens	Action
1.	\$	id + id * (id + id)\$	
2.			
...

No.	Stack	Tokens	Action	No.	Stack	Tokens	Action
1	\$	id + id * (id + id)\$		20	\$TWTKL	id + id)\$	L \rightarrow id
2	\$E		E \rightarrow FT	21	\$TWTKid		pop id
3	\$TF		F \rightarrow LK	22	\$TWTK + id)\$		K \rightarrow e
4	\$TKL		L \rightarrow id	23	\$TW		T \rightarrow ADFT
5	\$TKid		pop id	24	\$TK)TFAD		AD \rightarrow +
6	\$TK + id * (id + id)\$		K \rightarrow e	25	\$TK)TF +		pop +
7	\$T		T \rightarrow ADFT	26	\$TK)TF id)\$		F \rightarrow LK
8	\$TFAD		AD \rightarrow +	27	\$TK)TKL		L \rightarrow id
9	\$TF +		pop +	28	\$TK)TKid		pop id
10	\$TF id * (id + id)\$		F \rightarrow LK	29	\$TK)TK)\$		K \rightarrow e
11	\$TKL		L \rightarrow id	30	\$TK)T		T \rightarrow e
12	\$TKid		pop id	31	\$TK)		pop)
13	\$TK * (id + id)\$		K \rightarrow MUL LK	32	\$TK \$		K \rightarrow e
14	\$TKLMUL		MUL \rightarrow *	33	\$T		T \rightarrow e
15	\$TKL *		pop *	34	\$		accept
16	\$TKL (id + id)\$		L \rightarrow (E)				
17	\$TK)E(pop (
18	\$TK)E id + id)\$		E \rightarrow FT				
19	\$TK)TF		F \rightarrow LK				

3. Given a following grammar:

- 1 $E \rightarrow E \text{ AD } F,$
- 2 $E \rightarrow F,$
- 3 $\text{AD} \rightarrow +,$
- 4 $\text{AD} \rightarrow -,$
- 5 $F \rightarrow F \text{ MUL } L,$
- 6 $F \rightarrow L,$
- 7 $\text{MUL} \rightarrow *,$
- 8 $\text{MUL} \rightarrow /,$
- 9 $L \rightarrow (E),$
- 10 $L \rightarrow \text{id}$

- a. Is the grammar LL(1)? Justify your answer.
- b. If it's not LL(1), how to change the grammar to LL(1)?

a) Not LL(1), production 1, 5 has left-recursion

b) $E \rightarrow FT$

$T \rightarrow \text{AD } FT$

$T \rightarrow e$

$\text{AD} \rightarrow +$

$\text{AD} \rightarrow -$

$F \rightarrow LK$

$K \rightarrow \text{MUL } L K$

$K \rightarrow e$

$\text{MUL} \rightarrow *$

$\text{MUL} \rightarrow /$

$L \rightarrow (E)$

$L \rightarrow \text{id}$

4. The following is a grammar for regular expressions over symbols a and b only, using + in place of | for union, to avoid conflict with the use of vertical bar as a metasymbol in grammars:

```

rexpr -> rexpr + rterm | rterm
rterm -> rterm rfactor | rfactor
rfactor -> rfactor * | rprimary
rprimary -> a | b
    
```

- a. Left factor this grammar.
- b. Does left factoring make the grammar suitable for top-down parsing?
- c. In addition to left factoring, eliminate left recursion from the original grammar.
- d. Is the resulting grammar suitable for top-down parsing? Justify your answer.

a. the grammar cannot be left factored further.

b. not suitable

c. $\text{rexpr} \rightarrow \text{rterm } A$

$A \rightarrow + \text{rterm } A \mid e$

$\text{rterm} \rightarrow \text{rfactor } B$

$B \rightarrow \text{rfactor } B \mid e$

$\text{rfactor} \rightarrow \text{rprimary } C$

$C \rightarrow * C \mid e$

$\text{rprimary} \rightarrow a \mid b$

d. Suitable; there's no more left-recursion or left factor.

Exercise 3

1. Given a following grammar:

$$E \rightarrow F T,$$

$$T \rightarrow AD F T,$$

$$T \rightarrow e,$$

$$AD \rightarrow +,$$

$$AD \rightarrow -,$$

$$F \rightarrow L K,$$

$$K \rightarrow MUL L K,$$

$$K \rightarrow e,$$

$$MUL \rightarrow *,$$

$$MUL \rightarrow /,$$

$$L \rightarrow (E),$$

$$L \rightarrow id$$

- a. Find the first and follow sets of the grammar.
 - b. The parsing table of the grammar
2. From the parsing table in (1), use stack to simulate leftmost derivation as the LL(1) parsing for stream of tokens $id + id * (id + id)$.

No.	Stack	Tokens	Action
1.	\$	$id + id * (id + id)$ \$	
2.			
...

3. Given a following grammar:

$E \rightarrow E \text{ AD } F,$

$E \rightarrow F,$

$\text{AD} \rightarrow +,$

$\text{AD} \rightarrow -,$

$F \rightarrow F \text{ MUL } L,$

$F \rightarrow L,$

$\text{MUL} \rightarrow *,$

$\text{MUL} \rightarrow /,$

$L \rightarrow (E),$

$L \rightarrow \text{id}$

- Is the grammar LL(1)? Justify your answer.
 - If it's not LL(1), how to change the grammar to LL(1)?
4. The following is a grammar for regular expressions over symbols a and b only, using + in place of | for union, to avoid conflict with the use of vertical bar as a metasymbol in grammars:

```
rexpr -> rexpr + rterm | rterm
rterm -> rterm rfactor | rfactor
rfactor -> rfactor * | rprimary
rprimary -> a | b
```

- Left factor this grammar.
- Does left factoring make the grammar suitable for top-down parsing?
- In addition to left factoring, eliminate left recursion from the original grammar.
- Is the resulting grammar suitable for top-down parsing? Justify your answer.