

# Code Generation

Prabhas Chongstitvatana

Department of Computer Engineering

Chulalongkorn University

Thailand

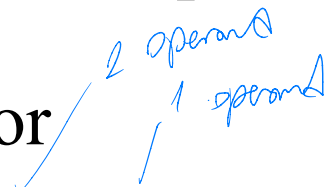
# Approach

- Beginning with Parse Tree
- Abstract Interpretation
  - Interpreter (take a parse tree and "run" it)
- Machine Abstraction
  - Machine Instructions (what kind of machine to "run" the abstract program)
- Code Generation

# Parse Tree

data structure: list

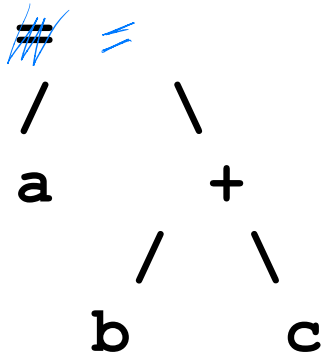
consists of operator and operand\*

Operator   
binary, unary, n-ary, function call

Operand  
constant, local variable, global variable (include vector)

# Illustrative example

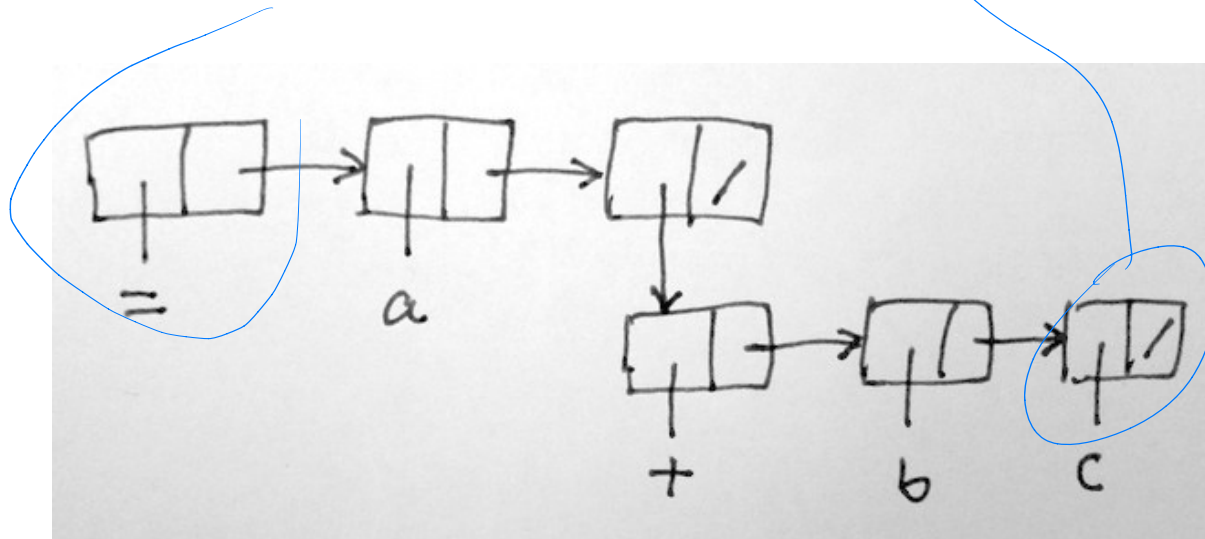
Source language:  $a = b + c$



Printed parsed tree:  $(= a (+ b c ))$

# Data structure

- This data structure is built from a number of cells.
- A cell contains two fields, head and tail.
- There are two types of cell: atom and dot-pair.
- An atom is the end of a branch (a kind of leaf node).
- A dot-pair is a pointer to another list.



# a program to copy a parse tree

```
copy(a)
```

```
  if(isatom(a)) return newatom(copyof(a))
```

```
  else return cons(copy(head(a)), copy(tail(a)))
```

# Interpreter

traverse a parse tree and execute the operator.

```
eval(e)
  if (e == nil) stop
  if isatom(e) do_atom(e)
  // then e is a list (operator operand*)
  a = head(e)
  b = tail(e)
  switch typeof(a)
    PLUS: ret eval(arg1(b)) + eval(arg2(b))
    MINUS:
    ...
    IF:   if( eval(arg1(b)) )
          ret eval(arg2(b))
          ret nil
    ...
```

# Interpreter (2)

```
do_atom(a)
  switch typeof(a)
    NUM:
    GLOBAL:
    LOCAL:
    . . .
```



# Machine Instructions (S-code)

- A stack-based instruction set.
- S-code is zero-address instruction. The main working storage is a stack and is comparable to the registers in a processor.
- A stack has two operations: push and pop.
- It is not necessary to "address" the stack (hence the name "zero-address").

# S-code

stack-based, fix 32-bit width

instruction format:    argument:24-bit    opcode:8-bit

arith:            add, sub, mul, div, mod

logic:            band, bor, bxor, shl, shr, not,  
eq, ne, lt, le, gt, ge

data:            ld, st, ldx, stx, lit, get, put

control:          call, ret, case, fun, jt, jf,  
jmp

other:            inc, dec, callt, sys

# Example: local var a,b,c

source :  $a = b + c$

S-code :

```
get.2 b  
get.3 c  
add  
put.1
```

# global var a,b,c

source:  $a = b + c$

S-code:

bad  
ld.b  
ld.c  
add  
st.a

# control: local var a,b,c

source: if( a == b ) c = 1 else c = 2

S-code:

get.1

get.2

eq

*jump iff (not)*

jf xx

*if a == b*

lit.1

put.3

jmp yy

:xx

<lit.2

# function call

source:   mysum(a,b){ return a + b; }

S-code:

```
fun.x  
get.1  
get.2  
add  
ret.y
```

# function call (2)

source: `main() { mysum(3,4); }`

S-code:

```
fun.x  
lit.3  
lit.4  
call.mysum  
ret.y
```

literal

# Code Generator

- The code generator works very much like an interpreter.
- Instead of "run" the parse tree, it "generates" the machine code that will give the same result as "run".



```
eval(e)
  if (e == nil) stop
  if isatom(e) do_atom(e)
  // then e is a list
  (operator operand*)
  a = head(e)
  b = tail(e)
  switch (typeof(a))
    PLUS:
      eval(arg1(b))
      eval(arg2(b))
      out(ADD)
```

# Example: generate if statement

source: if (a == b) return b;

parse tree

```
(if (== #1 #2 ) (return #2 ))
```

machine code

```
get.1
```

```
get.2
```

```
eq
```

```
jf.L18
```

```
get.2
```

```
ret.3
```

```
.L18
```

# Example of code generation

## Source

```
sum(n, m){  
  if( n == m ) return m;  
  else return n + sum(  
n+1, m);  
}
```

```
main(){  
  print(sum(1,10));  
}
```

## Parse Tree

```
(fun main  
  (print (call sum 1 10 )))  
(fun sum  
  (if-else  
    (== #1 #2 )  
    (return #2 )  
    (return (+ #1 (call sum (+ #1 1 )#2 )))))
```

# Parse Tree

```
(fun main (print (call sum 1 10 )))  
(fun sum  
  (if-else (== #1 #2 )  
    (return #2 )  
    (return (+ #1 (call sum (+ #1 1 )#2 )))))
```

## Machine code

:main	:sum	get.2
fun.1	fun.1	get.2
lit.1	get.2	lit.1
lit.10	get.1	add
call.sum	eq	get.1
sys.1	jf.L18	call.sum
ret.1	get.1	add
	ret.3	ret.3
	jmp.L26	:L26
	:L18	ret.3

# Summary

- Begin with abstract interpretation of a parse tree
- (interpreter)
- then modify it to output "sequence of machine code"