ALGORITHMS AND LAB (CSE130) INTRODUCTION TO ALGORITHMS

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Note: These notes are prepared from the following resources.

- (main text)Foundations of Algorithms, by Richard Neapolitan and Kumarss Naimipour
- OPython Algorithm (파이썬 알고리즘) by Y.K. Choi (2021) (Korean)
- Introduction to the Design and Analysis of Algorithms (3rd Edition) by Anany Levitin
- Introduction to Algorithms, (3rd Edition) by By Thomas H. Cormen, Charles E. Leiserson, Ronald L. Rivest and Clifford Stein
- https://www.geeksforgeeks.org

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- Course Overview
- WHAT IS AN ALGORITHM?
- Problems
 ⇔ Algorithms
- PROBLEM SOLVING PROCESS
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Course Overview

Course Objectives

- Understand algorithms for the most common problems(searching, sorting, shortest path, string manipulations, scheduling).
- Analyze the performance of algorithms (asymptotically).
- Apply methods for algorithm analysis.
- Learn important algorithmic design paradigms.
- Apply important algorithmic design paradigms in problem solving.
- Develop efficient algorithms for common science and engineering problems.

COURSE OVERVIEW (CONT...)

Course Prerequisites: It is assumed that you already have studied the following courses

Data Structures

• Programming C/C++ or JAVA or Python

Discrete mathematics

Grading Scheme

• Midterm Exam: 30%

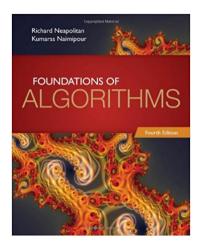
• Final Exam: 30%

• Exercises/Homework's/Labs : 30%

Participation, Q/A sessions: 10%

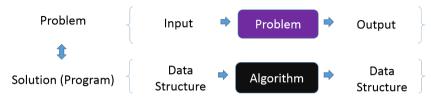
Course Overview (cont...)

- Main Textbook: Foundations of Algorithms, Fourth Edition, Richard Neapolitan, Kumarss Naimipour
- Additional References:
- Python Algorithm (파이썬 알고리즘) by Y.K. Choi (2021) (Korean)
- Introduction to Algorithm by Cormen, Thomas H., Charles E. Leiserson, Ronald L. Rivest, and Clifford Stein, 3rd ed. Cambridge, MA: MIT Press.
- Introduction to the design and analysis of algorithms by Anany Levitin, Pearson Addison Wesley.
- https://www.geeksforgeeks.org



What is an Algorithm?

What is an Algorithm?

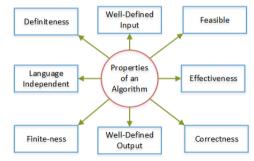


- A computer algorithm is a detailed step-by-step method for solving a problem using a computer.
- An algorithm is a sequence of unambiguous instructions for solving a problem, i.e., for obtaining a
 required output for any legitimate input in a finite amount of time.
- An algorithm is a well-defined computational procedure that takes some value or a set of values as input
 and produces some value or a set of values as output.
- Applying a technique to solve a problem results in a step by step procedure. This step by step procedure
 is called an algorithm for the problem.

Properties of an algorithm:

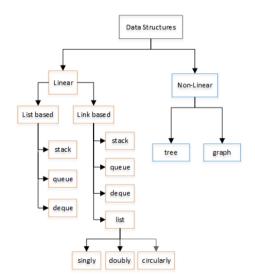
- Well-Defined Input An algorithm takes zero or more number of input values.
- Well-Defined Output An algorithm produces an output as result.
- Definiteness Every statement/instruction in an algorithm must be clear and unambiguous
- Finite-ness For all different cases, the algorithm must produce result within a finite number of steps.
- Effectiveness Every instruction must be basic enough to be carried out
- Feasible: The algorithm must be simple, generic and practical, such that it can be executed upon will the available resources.

- Correctness: Correct set of output values must be produced from the each set of inputs.
- Language Independent: The Algorithm designed must be language-independent, i.e. it must be just plain instructions that can be implemented in any language



Relationship between Data structures and Algorithms

- Data structures are entities designed to hold the data that is used by the algorithms.
- Each data structure has a way to store the elements in memory and functions that will help you to manipulate the data stored efficiently.
- Some data structures are better than others to be used to solve determinate problems efficiently.
- Note that, pragram = DataStructure + Algorithm



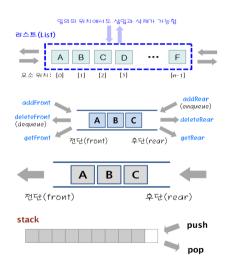
Linear Data Structures

 List is a more general data structure. Items can be inserted and deleted at any position. Deque, Queue and Stack are special cases of List data structure

 Deque is a specific data structure. Items can only be inserted and deleted at rear and front ends

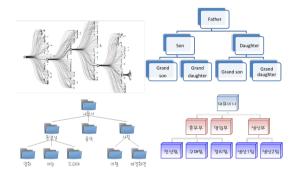
 Queue is a more specific data structure. Items can only be inserted at rear end and can be deleted at front end

 Stack is also a more specific data structure. Items can be inserted and deleted one end



Tree Data Structures

- There are a number of applications where linear data structures are not appropriate.
- For example, applications that require searching, linear data structures are not suitable.
- A linear linked list will not be able to capture the tree-like relationship with ease.
- Tree data structure is useful to organize complex data and it has numerous applications
- It has a hierarchical structure
 - ▶ Tree is finite set of one or more nodes such that
 - there is a special node called root
 - ② remaining nodes are partitioned into $n \ge 0$ disjoint trees T_1, T_2, \dots, T_n where each of these is a tree; we call each T_i subtree of the root
 - ► A tree is **Acyclic graph**(a graph that contains no cycle)



Graph Data Structures

- A graph data structure is a collection of nodes that have data and are connected to other nodes.
- Example: On facebook, everything is a node including User, Photo, Album, Event, Group, Page, Comment, Story, Video, Link, Note...anything that has data is a node.
- Every relationship is an edge from one node to another. Whether you post a photo, join a group, like a page, etc., a new edge is created for that relationship.



 All of facebook is then a collection of these nodes and edges. This is because facebook uses a graph data structure to store its data.

PROBLEMS

What is a Problem?

- A problem is a question to which we seek an answer.
- A problem may contain variables that are not assigned specific values in the statement of the problem.
- These variables are called parameters to the problem.
- A solution to an instance of a problem is the answer (algorithm) to the question asked by the problem.

Examples

- Is the key x in the array S of n keys? (Searching Problem)
- Add all the numbers in the array S of n numbers (Arithmetic Problem)
- Sort n keys in nondecreasing order (Sorting Problem)
- $\bullet \ \, \text{Matrix Multiplication} \Rightarrow \text{Determine the product of two n} \times \text{n matrices (Matrix Multiplication Problem)}$
- Determine whether x is in the sorted array S of n keys. (Sorting Problem)
- Determine the nth term in the Fibonacci sequence (Arithmetic Problem)

PROBLEMS (CONT...)

- Determine whether x is in the sorted array S of size n. (Searching Problem)
- Sort n keys in nondecreasing order (Sorting Problem)
- Merge two sorted arrays into one sorted array. (Merging Problem)
- Partition ⇒ Partition the array S for Quicksort (Partition Problem)
- Determine the product of two n x n matrices where n is a power of 2. (Matrix Multiplication Problem)
- Large Integer Multiplication ⇒ Multiply two large integers, u and v. (Integer Manipulation Problem)
- Compute the binomial coefficient (Computing Binomial Coefficients)
- Compute the shortest paths from each vertex in a weighted graph to each of the other vertex. The weights are nonnegative numbers (Shortest Path Problem)
- Print the intermediate vertices on a shortest path from one vertex to another vertex in a weighted graph.
 (Shortest Path Problem)

PROBLEMS (CONT...)

- Determining the minimum number of elementary multiplications needed to multiply n matrices (Matrix Multiplication Problem)
- Print the optimal order for multiplying n matrices. (Matrix Multiplication Problem)
- Determine the node containing a key in a binary search tree. It is assumed that the key is in the tree (Searching Problem)
- Determine an optimal binary search tree for a set of keys, each with a given probability of being the search key. (Searching Problem)
- Build an optimal binary search tree (Searching Problem)
- Determine an optimal tour in a weighted, directed graph. The weights are nonnegative numbers. (Optimization Problem)
- Determine a minimum spanning tree (Minimum spanning Tree Problem)
- Determine the shortest paths from v1 to all other vertices in a weighted, directed graph. (Shortest Path Problem)

PROBLEMS (CONT...)

- Determine the schedule with maximum total profit given that each job has a profit that will be obtained only if the job is scheduled by its deadline. (Scheduling Problem)
- Determine a binary character code by constructing a binary tree corresponding to an optimal code (Encoding Problem)
- Position n queens on a chessboard so that no two are in the same row, column, or diagonal. (Searching Problem)
- Given n positive integers (weights) and a positive integer W, determine all combinations of the integers that sum to W. (Searching Problem)
- Determine all ways in which the vertices in an undirected graph can be colored, using only m colors, so that adjacent vertices are not the same color. (coloring Problem)
- Determine all Hamiltonian Circuits in a connected, undirected graph. (Searching Problem)
- Let n items be given, where each item has a weight and a profit. The weights and profits are positive integers. Furthermore, let a positive integer W be given. Determine a set of items with maximum total profit, under the constraint that the sum of their weights cannot exceed W. (Optimization Problem)

Problems ⇔ Algorithms

Problems ⇔ **Algorithms**

- Searching (Sequential Search, Binary Search)
- Sorting (exchange Sort, bubble sort, heap sort, quick sort, merge sort)
- Fibonacci Terms (Recursive and iterative algorithms)
- Binomial Coefficients (Divide-and-Conquer algorithm, Dynamic Programming algorithm)
- Shortest Paths (Floyd's Algorithm, Dijkstra's Algorithm)
- Sum-of-Subsets(Backtracking Algorithm, Branch-and-Bound)
- Minimum spanning Tree (Prim's Algorithm, Kruskal's Algorithm)
- Encoding Problems (Huffman's Algorithm)
- Traveling Salesperson Problem (Greedy Algorithm, Backtracking Algorithm, Branch-and-Bound)
- Knapsack Problem(Dynamic Programming Approach, Grady Approach, Backtracking Algorithm, Branch-and-Bound)

PROBLEM SOLVING PROCESS

Problem Solving Process

- Requirements analysis and specifications: It involves Maintenance: It involves in modifying software to in determining 1) input and output required to solve the problem. 2) other information, such as software usage, feasibility of the solution
- Design: It involves in 1) selecting appropriate data structures to organize the input/output data and designing algorithms, needed to process the data efficiently
- Implementation: It involves in selecting appropriate programming language and translating algorithms into well-structured, well-documented, readable code
- Testing: It involves in executing the software and correcting errors and proving correctness of algorithms and modifying algorithms

improve performance and adding new features, etc Problem Requirement Analysis **Specifications** Design Solving Implementation Testing **Proces**

Maintenance

DESIGN AND IMPLEMENTATION PHASES



Algorithm design techniques

- There are general approaches to construct efficient solutions to problems. They provide templates suited to solving a broad range of diverse problems.
 - Brute Force
 - Divide and conquer
 - Dynamic programming
 - Greedy approach
 - State space search techniques



Expressing Algorithms (Pseudocode)

- Pseudocode specifies the steps required to accomplish the algorithm.
- Pseudocode cannot be compiled nor executed, and there are no syntax rules.
- Can be written any human language
- Example: getting and printing records from a file

```
1: procedure Example()
       open subscriber file
       Get a record
 3:
       while more records do
          if count > 3 then
 5:
              Output the record
6.
 7:
          else
8:
              Get another record
          end if
g.
       end while
10:
11: end procedure
```

DESIGN AND IMPLEMENTATION PHASES



Algorithm Analysis

- The analysis of algorithms is to determine the computational complexity in terms of time and storage.
- Two approaches: Empirical/Experimental analysis and Theoretical/formal analysis
- Empirical/Experimental analysis:
 - It compares the time of execution for various algorithms.
 - lt requires implementation of algorithms before analysis for a specific problem
 - It depends on
 - Actual number of CPU cycles (Computer)
 - Number of instructions
 - Programming language
 - Programmer



Algorithm Analysis

- The analysis of algorithms is to determine the computational complexity in terms of time and storage.
- Two approaches: Empirical/Experimental analysis and theoretical/formal analysis
- Theoretical/Formal analysis:
 - It involves determining a function (the time complexity function) that relates the algorithm's **input size** to the number of steps it takes or the number of storage locations it uses (its space complexity function).
 - It is common to estimate algorithm complexity in the asymptotic sense i.e., using order notations, Big-O notation, Big-omega(Ω) notation and Big-theta(Θ) notation
 - lt dose not depend on
 - Actual number of CPU cycles (Computer)
 - Number of instructions
 - Programming language
 - Programmer



Algorithm Correctness Analysis

- This means to verify if the algorithm leads to the correct solution of the problem after a finite number of processing steps.
- Two common strategies: Experimental analysis and Formal analysis
- Experimental analysis (Program Correctness):
 - ▶ The main advantage of this approach is its simplicity.
 - ▶ The main disadvantage is the fact that testing cannot cover always all possible instances of input data (it is difficult to know how much testing is enough).



Algorithm Correctness Analysis

- This means to verify if the algorithm leads to the correct output of the problem instance (specific input) after a finite number of processing steps.
- Two common strategies: Experimental Analysis and Formal Analysis
- Formal Analysis (Algorithm Correctness):
 - ▶ The main advantage of this approach is that if it is rigorously applied it guarantees the correctness of the algorithm.
 - ▶ The main disadvantage is the difficulty of finding a proof, mainly for complex algorithms.
 - ▶ A common technique for proving correctness is to use **mathematical induction** because an algorithm's iterations provide a natural sequence of steps needed for such proofs

EXAMPLES

Example-1: Add Array Members

- Problem: Add all the numbers in the array S of n numbers.
- Inputs: positive integer n, array of numbers S indexed from 1 to n.
- Outputs: sum, the sum of the numbers in S.
- Pseudo-code

Algorithm 1 Add Array Members (Pseudo-code)

- 1: **procedure** SUMARRAY(integer n, number S[])
- 2: integer i
 - 3: number *sum*
- 4: sum = 0
- 5: **for** $(i = 1; i \le n; i + +)$ **do**
- 6: sum = sum + S[i]
- 7: end for
- 8: return sum
- 9: end procedure

Example-2: Exchange Sort

- Problem: Sort *n* keys in nondecreasing order.
- Inputs: positive integer n, array of keys S indexed from 1 to n.
- Outputs: the array *S* containing the keys in nondecreasing order.
- Pseudo-code

Algorithm 2 Exchange Sort

```
1: procedure EXCHANGESORT(integer n, number S[])
2: integer i,j
3: for (i=1;i<=n-1;i++) do
4: for (j=i+1;j<=n;j++) do
5: if (S[j] < S[i]) then
6: exchange S[i] and S[j]
7: end if
8: end for
9: end procedure
```

Example 3: A Recursive Algorithm for computing the n-th power of 2

- Problem: Computing the n-th power of 2.
- Inputs: n, a natural number.
- Outputs: n-th power of 2.
- Pseudo-code

Algorithm 3 n-th power of 2

```
1: procedure Power2(n)
```

2: **if** (n == 0) then

3: return 1;

4: else

5: return 2 * Power2(n-1)

6: end if

7: end procedure

Example 4: Matrix Multiplication

- Problem: Determine the product of two nxn matrices.
- Inputs: a positive integer n, two-dimensional arrays of numbers A and B, each of which has both its rows and columns indexed from I to n.
- Outputs: a two-dimensional array of numbers C, which has both its rows and columns indexed from 1 to n, containing the product of A and B.
- Pseudo-code

Algorithm 4 Matrix Multiplication

```
1: procedure MATRIXMULTUPLY(integer n, number
   A[][], number B[][])
       integer i, j, k
 2:
 3:
       number C[][]
       for (i = 1; i <= n; i + +) do
           for (i = 1; i \le n; i + +) do
 5:
              C[i][i] = 0
 6:
              for (k = 1; k \le n; k + +) do
 7:
                  C[i][i] = C[i][i] + A[i][k] * B[k][i]
              end for
 g.
           end for
10:
       end for
11:
       return C
12.
13: end procedure
```

Example 5: nth Fibonacci Term (Iterative)

- Problem: Determine the nth term in the Fibonacci sequence.
- **Inputs:** a nonnegative integer *n*.
- Outputs: fib2, the *nth* term in the Fibonacci sequence.
- Pseudo-code

Algorithm 5 Algorithm 1.7: nth Fibonacci Term (Dynamic Programming)

```
1: procedure FIB2(integer n)
2: integer i
3: integer f[0..n]
4: f[0] = 0
5: f[1] = 1
6: for (i = 2; i <= n; i + +) do
7: f[i] = f[i - 2] + f[i - 1]
8: end for
9: return f[n]
10: end procedure
```

Example 6: Recursive Algorithm for nth Fibonacci term

- Problem: Determine the n_{th} term in the Fibonacci sequence.
- Inputs: a nonnegative integer n.
- Outputs: fib, the n_{th} term of the Fibonacci sequence.

Algorithm 6 Recursive Algorithm for computing Fibonacci Sequence

```
1: procedure FIB(integer n)
2: if (n <= 1) then
3: return n
4: else
5: return fib(n-1) + fib(n-2)
6: end if
7: end procedure
```

Example 7: Sequential Search

- **Problem:** Is the key x in the array S of n keys?
- Inputs (parameters): positive integer n, array of keys S indexed from 1 to n, and a key x.
- Outputs: location, the location of x in S (0 if x is not in S.)
- Pseudo-code

Algorithm 7 Sequential Search (Iterative)

```
1: procedure SEQSEARCH(integer
                                            number
   S[], number x)
       integer location
2:
       location = 1
 3:
       while location \leq n \&\& S[location]! = x
4:
   do
          location = location + 1
5.
6:
          if (location > n) then
              location = 0
7.
          end if
      end while
g.
       return location
10:
11: end procedure
```

Example 8: Binary Search (Recursive)

- **Problem:** Determine whether *x* is in the sorted array *S* of size *n*.
- Inputs: positive integer n, sorted (nondecreasing order) array of keys S indexed from 1 to n, a key x.
- Outputs: location, the location of x in S (0 if x is not in S).
- Pseudo-code

```
1: procedure LOCATION(low, high)
 2:
       integer mid
       if (low > high) then
           return 0
 4.
       else
 5:
           mid = \left| \frac{(low + high)}{2} \right|
6.
           if (x == S[mid]) then
 7:
               return mid
 8:
           else
 g.
               if (x < S[mid]) then
10:
                   return location(low, mid - 1)
11.
               else
12.
                   return location(mid + 1, high)
13:
               end if
14:
           end if
15:
       end if
16.
17: end procedure
```