# CECS 474: Computer Network Interoperability

Lab 2

#### **CSMA/CD Performance Evaluation**

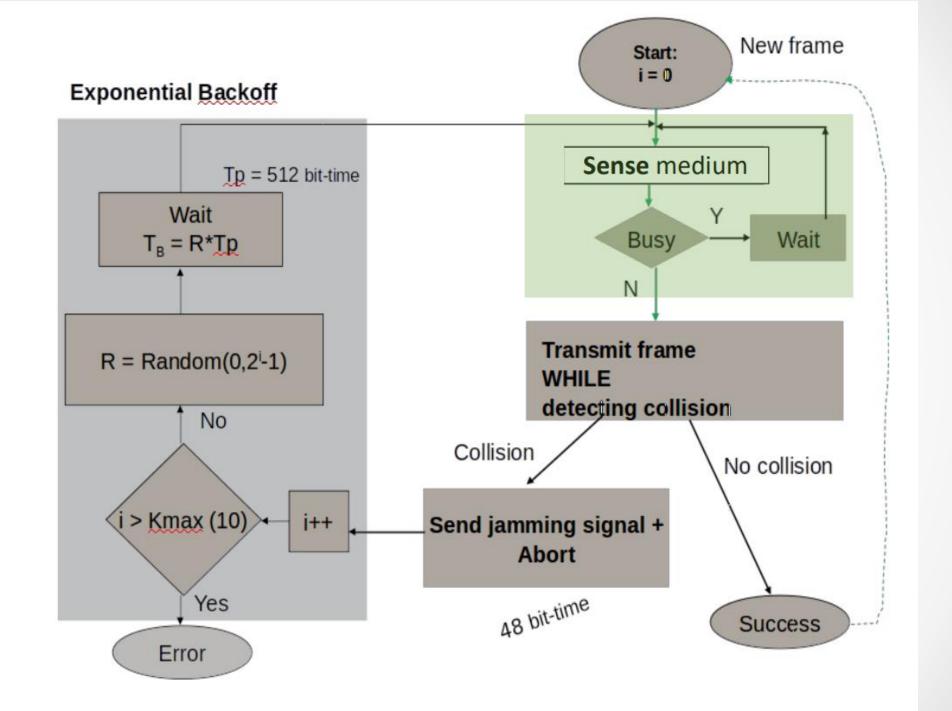
Haixia Peng
Department of CECS
California State University Long Beach

#### Goals

- In this lab, you will once again be implementing a Discrete Event Simulator (similar to what you did in Lab 1)
- You will be simulating the behavior of a CSMA/CD LAN
- You will be monitoring the performance of the LAN

#### How does CSMA/CD work?

- Listen to the network, and wait for it to be idle
- Transmit a frame, and while transmitting, check for collisions
- If no collisions, the frame was sent successfully
- If there were collisions, back off for a bit and then try again



# Exponential Backoff

- If everyone on the LAN tries to transmit as soon as the medium is available, you'll get a lot of collisions
- Instead, we have the nodes wait a random amount of time before retransmitting
- The random amount of time increases exponentially, so when network isn't busy you don't wait long
- Specifically, wait Random $(0, 2^i 1)^*$  512 bit-times where i is the number of times you've retried
- Retry up to  $K_{max}$  times, and discard the frame if you go past that  $(K_{max})$  == 10 in this lab)

# Sensing the Medium

- There are two alternatives when it comes to sensing whether the medium is busy
- Persistent
  - o If the channel is busy, continue sensing until it becomes idle
- Non-persistent
  - o If the channel is busy, wait a while before checking again
  - o By "a while" we mean a random period of time
  - o For this lab, use the same algorithm as for exponential backoff (previous slide)
- Persistent mode is "greedy", non-persistent is friendlier

# Simplifying Assumptions

- Real-world implementations can be quite complex
- To make it easier, you can assume the following:
  - o There is no need to simulate the "jamming signal" shown in the figure
  - All nodes see a collision at the same time (aside from propagation delay)
  - o The distance between consecutive nodes on the network is the same
  - All the packets are the same length

#### What you need to do

- You need to create an application in C, C++, Python or Java that simulates the a CSMA/CD-based network
- We provide the speed of the network in bits/second, the length of the packet, the distance between nodes and the propagation speed (these are all constant throughout your simulations)
- You will simulate both a 1-persistent and a non-persistent CSMA-CD protocol
- For each, you will vary the number of nodes N and the average packet arrival rate (using the same Poisson process as for Lab 1), compute the efficiency and throughput, and plot the results

#### How to do it

- You can implement this lab however you like
- What follows is a description of one possible approach
- If you choose a different approach, that's fine, but be sure to document the approach you're using, how it works and why you chose it.

# Setup and transmission

#### Setup

- Create a queue of packets for each node
- Each packet has an arrival timestamp
- o Populate the queues, generating the timestamps the same way you did for Lab 1
- Select a simulation time T (the same way you did for Lab 1)

#### Sending a packet

- Determine which node transmits next (the one whose queue head has the lowest timestamp)
- Set the current simulation time to the timestamp of that packet
- In other words, the time of the next discrete event in the simulator is the time of the next packet (no matter which queue it's in)
- The packet is transmitted at that time (i.e. that is the time that the first bit of the packet is sent)

# Propagation Delay

- The first bit of the packet arrives at the adjacent nodes immediately after some fixed propagation time T<sub>prop</sub>, and later than that at more distant nodes (after some multiple of T<sub>prop</sub>)
- You can also compute T<sub>transmission</sub> which is the length of time the bits of the packet are being transmitted, based on the (fixed) length of a packet L and the (fixed) transmission rate R

# Detecting Bus Busy

- When node B sends a packet at time  $T_b$ , the bus is busy (at node C) between  $T_b + T_{prop}$  and  $T_b + T_{prop} + T_{transmission}$
- If the arrival time of the first packet in the queue at C is greater than  $T_b$  +  $T_{prop}$  and less than  $T_b$  +  $T_{prop}$  +  $T_{transmission}$ , C will see the bus as busy and will wait
- "Waiting" means updating the arrival times of C's queued packets so that any packets that were due to be sent before the end B's packet transmission are rescheduled to the end of the packet (i.e. to  $T_b + T_{prop}$ )

# Detecting Collisions

- If the arrival time of the next packet in the queue at C is less than  $T_b + T_{prop}$ , then C does not know that there is a packet on the bus already and will therefore start transmitting
- If the arrival time of the first packet in the queue is between  $T_b + T_{prop}$  and  $T_b + T_{prop} + T_{transmission}$ , a collision has occurred and you need to use exponential backoff
- Reschedule the packets in the queue at C as discussed on the previous slide, setting their time to the end of the waiting period

# Exponential Backoff

- When you transmit a packet for the first time, you reset a counter to zero
- Every time you try to transmit and encounter a collision, you increment the counter
- If the counter reaches ten, the packet is discarded
- Each time you increment the counter (call it iteration) you generate a random waiting period between zero and 2<sup>iteration</sup>-1 times 512 bit-times (the time it takes to transmit a bit, based on R)

#### Non-Persistent Mode

The non-persistent case is exactly the same as the persistent case we
just described, except that instead of setting the waiting time to be
the end of the packet from B it's set to a random (same way you do
the backoff, so you need a separate iteration counter)

# Efficiency and Throughput

- The efficiency is the number of successfully transmitted packets divided by the total number of packets pulled from the queues
- The throughput (in Mbps) can be computed using the number of successfully transmitted packets, the packet length L, and the simulation time T

#### What You Need to Find

- You are given fixed values for the packet length L, the transmission rate R, the distance between adjacent nodes D and the propagation speed S
- You will run the simulation multiple times, using three different values for the packet arrival rate A (7, 10 and 20) and varying the number of nodes N from 20 to 100 inclusive in steps of 20
- You are expected to compute the efficiency and throughput for both the persistent and non-persistent modes
- You are expected to graph the efficiency vs number of nodes for each of the three arrival rates

17

# Marking Criteria

- You will be marked based on the correctness of your results (e.g. shape of the curves)
- You will also be marked based on your written report and the description you provide of how your simulation works
- The written report is especially important if you choose an approach to simulation that differs from the one described here

#### Notes

- Programming languages: you can use C, C++ or Python or any other programming language
- You must submit your code and your report on BeachBoard
- You should provide a single command that builds and runs your application and describe (in your report) how to use it

#### Questions?